

# 10.555 Bioinformatics

## Spring 2003

### Lecture 2 (part B)

# My Not-So-Fair Casino

- we play one game only:
  - ▲ heads / tails
- we have both fair and biased coins
- at the beginning of the day my croupier picks a coin and heads to his or her table
- can you win?
- can you loose?
- can you tell which coin the croupier has selected?

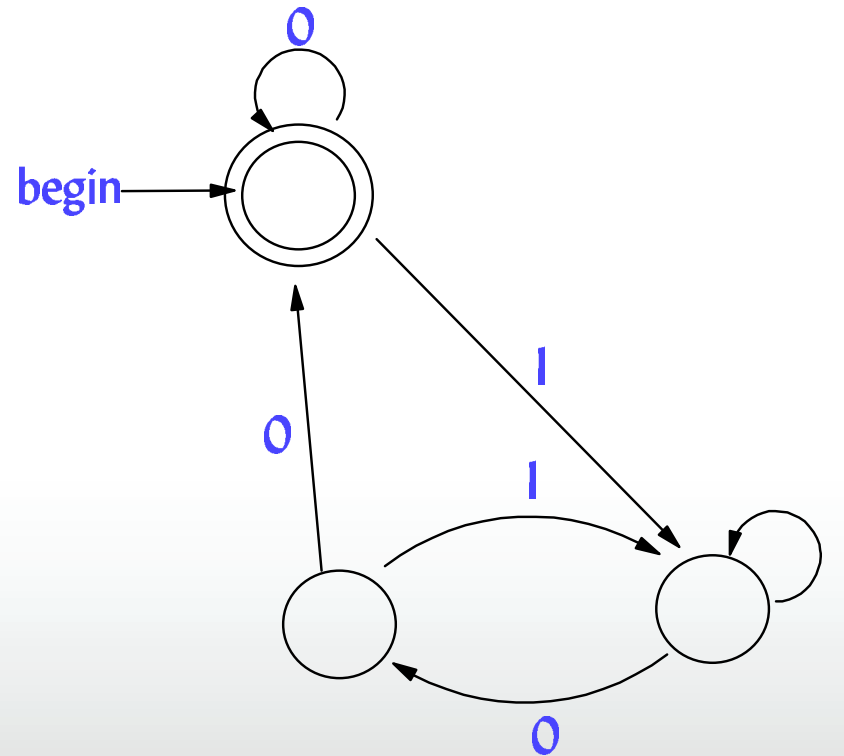
# My Not-So-Fair Casino (cont.)

- what do you need to do in order to be able to tell which coin is in use?
- how can you tell which coin is in use?
- Is the following information useful?
  - ▲  $\Pr(\text{head/coin is fair}) = \Pr(\text{tail/coin is fair}) = 1/2$
  - ▲  $\Pr(\text{head/coin is biased}) = 3/4$
  - ▲  $\Pr(\text{tail/coin is biased}) = 1/4$

# Finite Automata

- a finite automaton comprises
  - ▲ *states*
  - ▲ *transitions* occurring on input symbols that come from an alphabet  $\Sigma$
- a special "begin" state / one or more special "end" states
- a (directed) *transition diagram* is associated with an fsa
- a f.a. *accepts* a string  $x$  iff the string  $x$  leads to a sequence of transitions from "begin" to "end" (i.e. output of an f.a. is "accept" or "don't accept")

# Finite Automata (cont.)

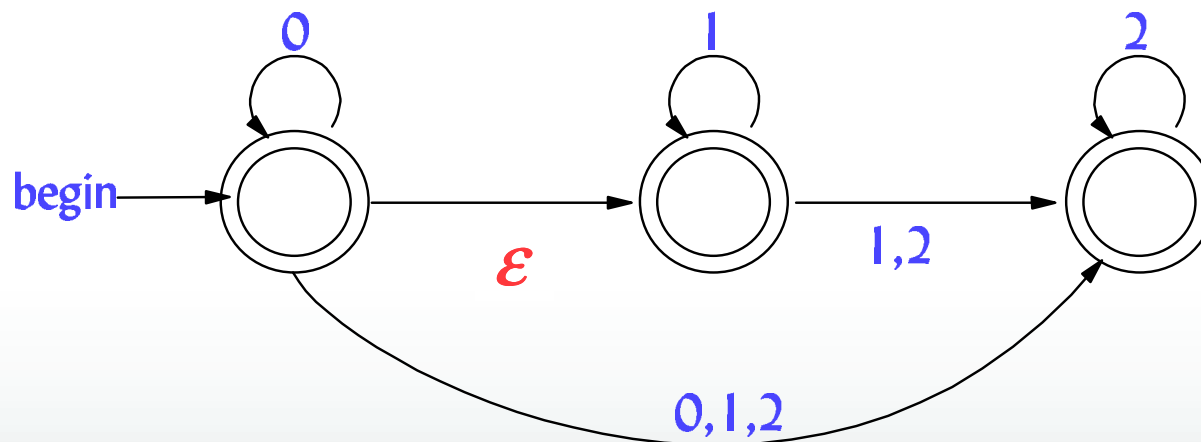


# Finite Automata (cont.)

- the set of all strings accepted by a f.a. form the *language* accepted by the f.a.
- a language is *regular* iff it is accepted by an f.a.
  
- an f.a. can be
  - ▲ deterministic
  - ▲ non-deterministic (allows *multiple* transitions out of a state on the *same* symbol)
  
- d-f.a. and non-d-f.a. are equivalent!

# Finite Automata (cont.)

- example of a non-deterministic finite automaton



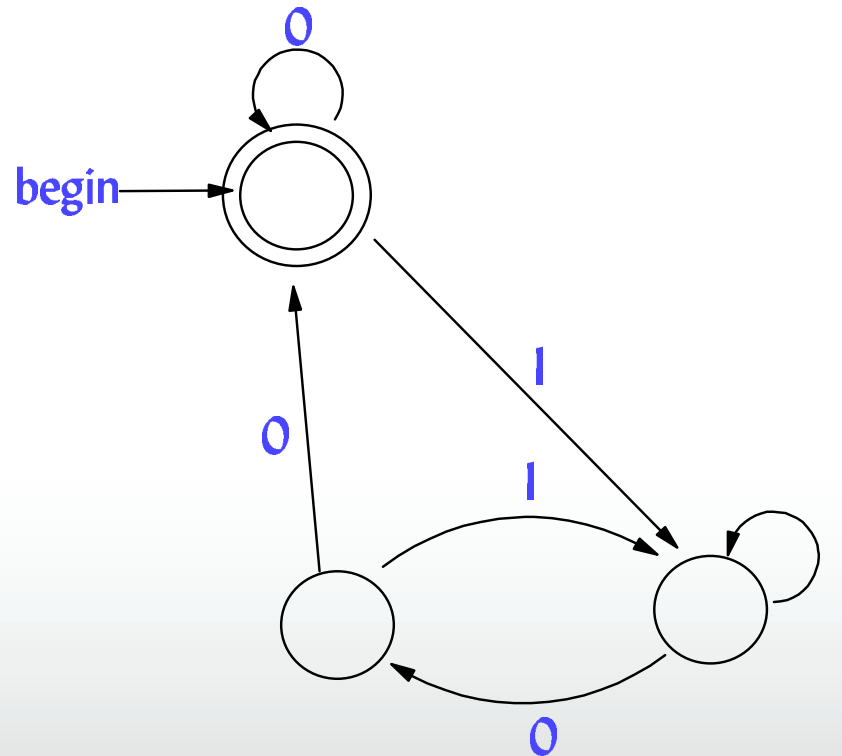
# Regular Expressions

- simple expressions defined in conjunction with an alphabet  $\Sigma$  and describing regular languages
- definition:
  - ▲  $\emptyset$  is a r.e. and denotes  $\{ \}$
  - ▲  $\varepsilon$  is a r.e. that denotes  $\{ \varepsilon \}$
  - ▲ for each  $a$  in  $\Sigma$ ,  $a$  is a r.e. that denotes  $\{ a \}$
  - ▲ if  $e_1$  and  $e_2$  are r.e. that denote the languages  $E_1$  and  $E_2$  respectively, then  $e_1 + e_2$ ,  $e_1 e_2$  and  $e_1^*$  denote languages  $E_1 + E_2$ ,  $E_1 E_2$  and  $E_1^*$  respectively



# Regular Expressions (cont.)

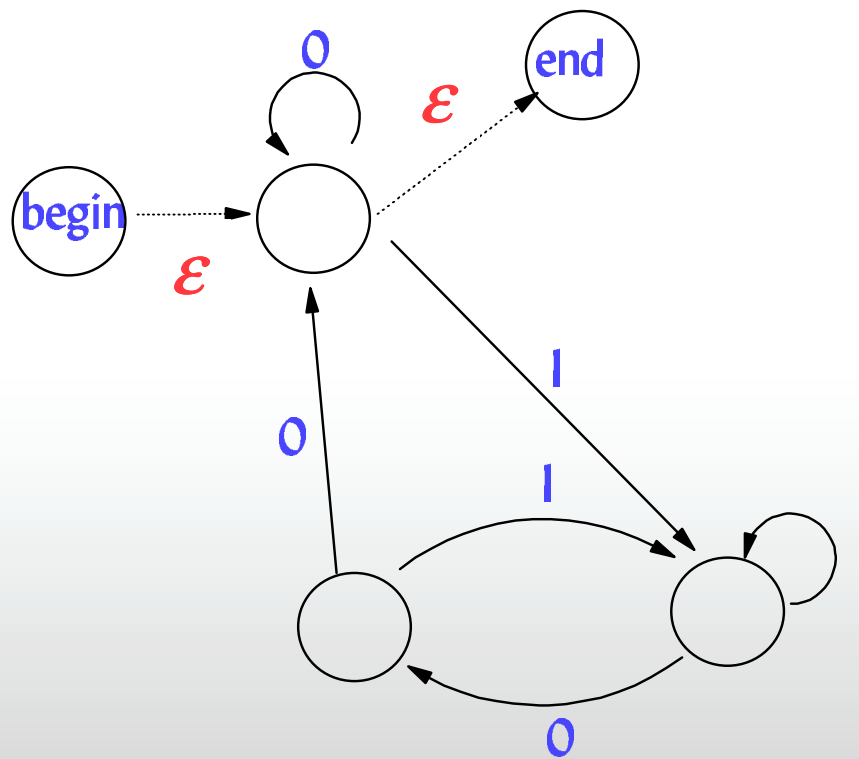
- let's derive the regular expression for this f.a. ...



# Finite Automata with Output

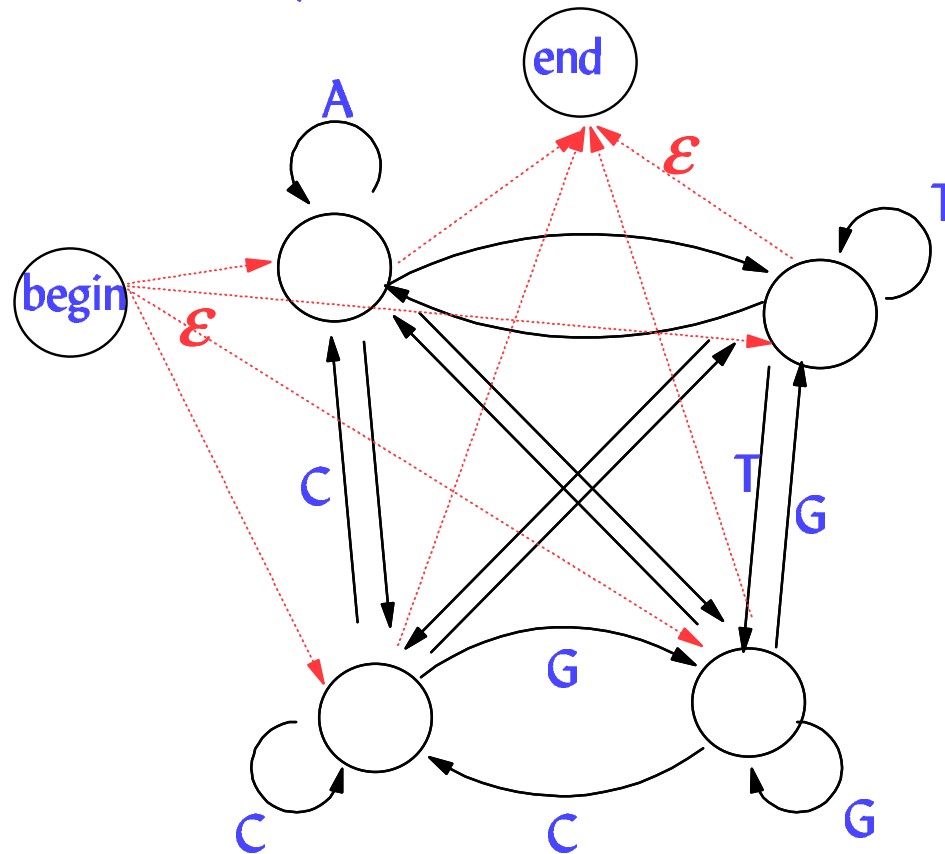
- two choices:
  - ▲ attach output to the state that is entered (Moore machines)
  - ▲ attach output to the transition that is traversed (Mealy machines)
  
- Moore and Mealy machines are equivalent!

# Trivial Extensions



# More Extensions

- a f.a. that "accepts" all DNA inputs  
(not all transitions are labelled)



- what is the limitation of this state diagram?

# My Not-So-Fair Casino / Part 2

- what if my croupier could switch coins on-the-fly while at his/her table?
- can you win?
- can you lose?
- can you tell?
- how is this different from the previous situation?
- Markov Chain vs. Hidden Markov Model