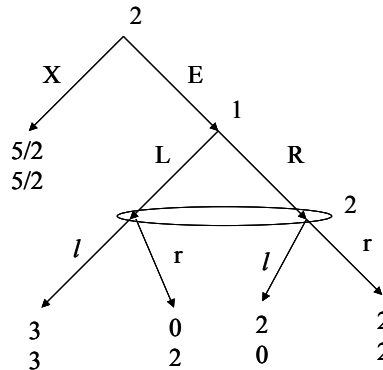


14.12 Game Theory – Midterm II
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Instructions. This is an open book exam; you can use any written material. You have one hour and 20 minutes. Each question is 25 points. Good luck!

1. Consider the following game:



Compute all the pure-strategy subgame-perfect equilibria. Use a forward induction argument to eliminate one of these equilibria.

2. Below, there are pairs of stage games and strategy profiles. For each pair, check whether the strategy profile is a subgame-perfect equilibrium of the game in which the stage game is repeated infinitely many times. Each agent tries to maximize the discounted sum of his expected payoffs in the stage game, and the discount rate is $\delta = 0.99$.

(a) **Stage Game:**

1\2	L	M	R
T	2,-1	0,0	-1,2
M	0,0	0,0	0,0
B	-1,2	0,0	2,-1

Strategy profile: Until some player deviates, player 1 plays T and player 2 alternates between L and R. If anyone deviates, then each play M thereafter.

(b) **Stage Game:**

1\2	A	B
A	2,2	1,3
B	3,1	0,0

Strategy profile: The play depends on three states. In state S_0 , each player plays A; in states S_1 and S_2 , each player plays B. The game starts at state S_0 . In state S_0 , if each player plays A or if each player plays B, we stay at S_0 , but if a player i plays B while the other is playing A, then we switch to state S_i . At any S_i , if player i plays B, we switch to state S_0 ; otherwise we stay at state S_i .

3. Consider the following first-price, sealed-bid auction where an indivisible good is sold. There are $n \geq 2$ buyers indexed by $i = 1, 2, \dots, n$. Simultaneously, each buyer i submits a bid $b_i \geq 0$. The agent who submits the highest bid wins. If there are $k > 1$ players submitting the highest bid, then the winner is determined randomly among these players — each has probability $1/k$ of winning. The winner i gets the object and pays his bid b_i , obtaining payoff $v_i - b_i$, while the other buyers get 0, where v_1, \dots, v_n are independently and identically distributed with probability density function f where

$$f(x) = \begin{cases} 3x^2 & x \in [0, 1] \\ 0 & \text{otherwise.} \end{cases}$$

- (a) Compute the symmetric, linear Bayesian Nash equilibrium.
- (b) What happens as $n \rightarrow \infty$?

[Hint: Since v_1, v_2, \dots, v_n is independently distributed, for any w_1, w_2, \dots, w_k , we have

$$\Pr(v_1 \leq w_1, v_2 \leq w_2, \dots, v_k \leq w_k) = \Pr(v_1 \leq w_1) \Pr(v_2 \leq w_2) \dots \Pr(v_k \leq w_k).]$$

4. This question is about a thief and a policeman. The thief has stolen an object. He can either hide the object INSIDE his car or in the TRUNK. The policeman stops the thief. He can either check INSIDE the car or the TRUNK, but not both. (He cannot let the thief go without checking, either.) If the policeman checks the place where the thief hides the object, he catches the thief, in which case the thief gets -1 and the police gets 1; otherwise, he cannot catch the thief, and the thief gets 1, the police gets -1.

- (a) Compute all the Nash equilibria.
- (b) Now imagine that we have 100 thieves and 100 policemen, indexed by $i = 1, \dots, 100$, and $j = 1, \dots, 100$. In addition to their payoffs above, each thief i gets extra payoff b_i from hiding the object in the TRUNK, and each policeman j gets extra payoff d_j from checking the TRUNK. We have

$$\begin{aligned} b_1 &< b_2 < \dots < b_{50} < 0 < b_{51} < \dots < b_{100}, \\ d_1 &< d_2 < \dots < d_{50} < 0 < d_{51} < \dots < d_{100}. \end{aligned}$$

Policemen cannot distinguish the thieves from each other, nor can the thieves distinguish the policemen from each other. Each thief has stolen an object, hiding it either in the TRUNK or INSIDE the car. Then, each of them is randomly matched to a policeman. Each matching is equally likely. Again, a policeman can either check INSIDE the car or the TRUNK, but not both. Compute a pure-strategy Bayesian Nash equilibrium of this game.

The game for problem 1.

