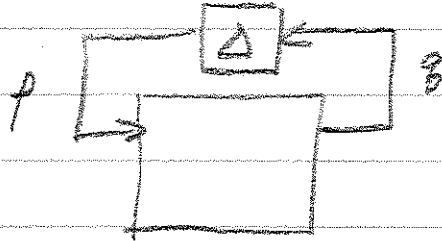


Structured Uncertainties

So far, have look at unstructured uncertainties:

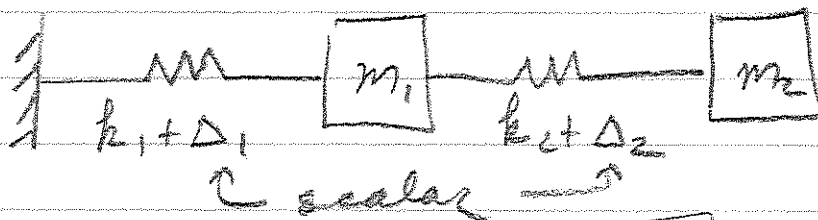


p, q vectors Δ matrix

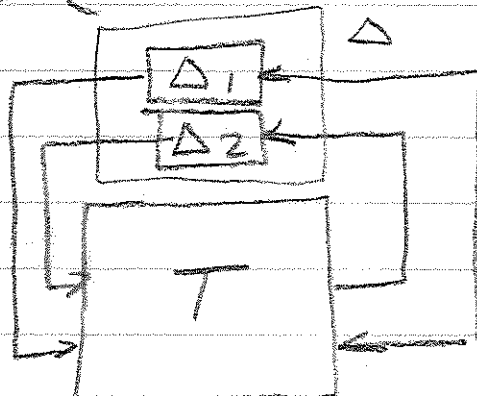
$$\|\Delta\|_{\infty} < 1 \quad (\text{or } \bar{\sigma}(\Delta(j\omega)) < 1)$$

In practice, the Δ matrix may be structured

Example



model:



so $\Delta = \begin{bmatrix} \Delta_1 & 0 \\ 0 & \Delta_2 \end{bmatrix}$

Suppose we know $\overline{\sigma}(\Delta_1) \leq 1$,
 $\overline{\sigma}(\Delta_2) \leq 1$. Define

$$\mathcal{D} = \left\{ \Delta : \Delta = \begin{bmatrix} \Delta_1 & 0 \\ 0 & \Delta_2 \end{bmatrix}, \overline{\sigma}(\Delta_1) \leq 1, \overline{\sigma}(\Delta_2) \leq 1 \right\}$$

I.e., allowable Δ 's are

$$\Delta \in \mathcal{D}$$

Can we use earlier methods to determine stability of the loop?

Sort of.

Note that

$$\Delta \in \mathcal{D} \implies \overline{\sigma}(\Delta) \leq 1$$

~~\leftarrow~~

Example $\Delta = \begin{pmatrix} 0.1 & 0.1 \\ 0.1 & 0.1 \end{pmatrix}$

$$\Rightarrow \bar{\sigma}(\Delta) = 0.2$$

But $\Delta \notin \mathcal{D}$!

So system is stable if

$$\bar{\sigma}(T) < 1$$

But this is sufficient, not necessary

Example

$$T = \begin{bmatrix} 1/2 & 100 \\ 0 & 1/2 \end{bmatrix}$$

$$\bar{\sigma}(T) = 100.025$$

But, stable for all $\Delta \in \mathcal{D}$

$$(I - \Delta T)^{-1}$$

The structured singular value

What we need is a way to describe how big T is, that takes into account the structure of the uncertainty.

Define the function

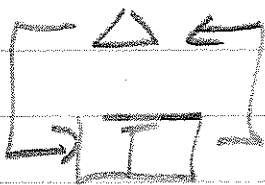
$$\mu(T) = \frac{1}{k},$$

where k is the largest k such that

$$\det(I + k\Delta T) \neq 0 \text{ for all } \Delta \in \mathcal{D}$$

$\mu(T)$ is the structured singular value of T .

Theorem. The loop



is stable for all $\Delta \in \mathcal{D}$ iff $\mu(T) < 1$. $\forall \omega$.

Hard to calculate $\mu(T)$, but can bound, and bounds allow for control design.

Uncertainty blocks:

Repeated real scalars:

$$\Delta_i = \delta_i^r \mathbf{I}$$

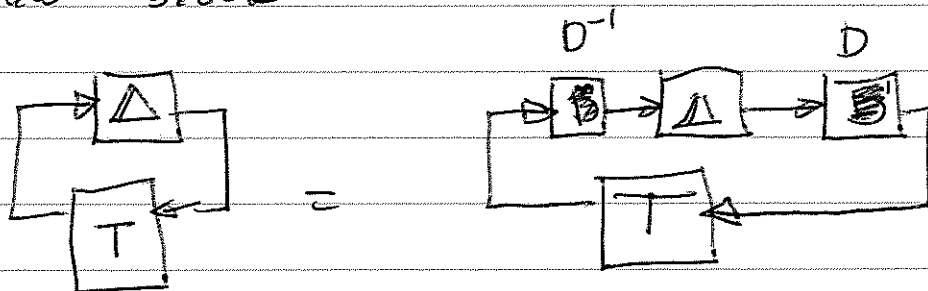
Repeated complex scalars:

$$\Delta_i = \delta_i^c \mathbf{I}$$

Complex matrix

$$\Delta_i = \text{complex matrix}$$

Basic idea for computing μ : ~~if~~ if D_i commutes with Δ_i , then can redraw block



where $D = \text{diag}(D_1, D_2, D_3, \dots)$

Note that since $\mu(T) \leq \bar{\nu}(T)$,
it is also true that

$$\mu(T) \leq \bar{\nu}(D^* T D^{-1})$$

So can optimize over D to find
best (smallest) upper bound.

For $\Delta_i = \delta_i^c I$, $D_i = \text{any matrix}$

For $\Delta_i = \text{complex matrix}$, $D_i = d_i I$

Will come back to real case.

Note that

$$\bar{\nu}(D T D^{-1}) = \alpha$$

$$\Rightarrow p^* D^{-*} T^* D^* D T D^{-1} p \leq \alpha p^* p$$

$$\Rightarrow \cancel{p^* D^{-*} T^* D^* D T D^{-1} p}$$

$$\bar{p}^* \cancel{D^{-*}} T^* D^* D T \bar{p} \leq \alpha \bar{p}^* \underbrace{D^* D}_{\bar{D}} \bar{p}$$

$$\Rightarrow \bar{p}^* T^* \bar{D} T \bar{p} \leq \alpha \bar{p}^* \bar{D} \bar{p}$$

$$\Rightarrow T^* \bar{D} T \leq \alpha \bar{D}$$

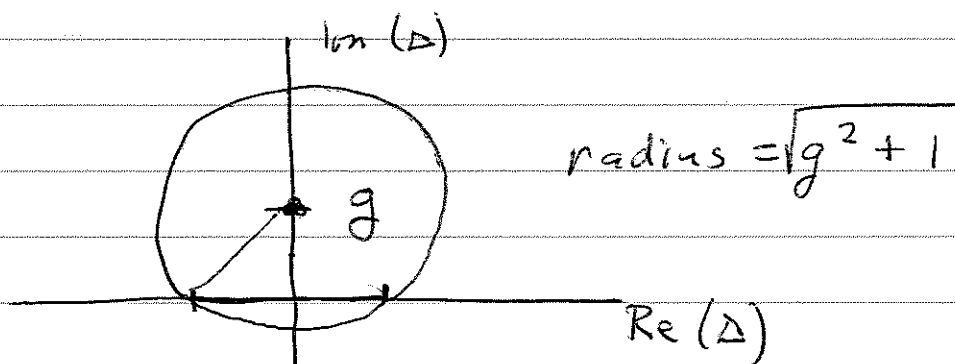
This is a linear matrix inequality (LMI), which is relatively easily solved. So can do a " γ -Iteration" to find (bound on) μ , which also generates \bar{D} .

With real uncertainties, can add term:

$$T^* \bar{D} T + j(GT - T^*G) \leq \alpha \bar{D}$$

Note that $\bar{D} = \bar{D}^* > 0$, $G = G^*$. G must also commute appropriately with the real Δ blocks. ($G = 0$ for complex blocks.)

G term corresponds to covering a real Δ with off-axis circle:



D-K iteration:

1. For D (and G) given, find controller $K(s)$ that minimizes ∞ -norm of $\begin{bmatrix} D & Tzw \\ 0 & D^{-1} \end{bmatrix}$. This requires H_2 controller and γ iteration.
2. For given $K(s)$, find D (and G) that minimizes $\bar{\sigma}(DTzwD^{-1})$. This can be done with efficient LMI techniques. Also requires γ -iteration.
3. Repeat until desired convergence.