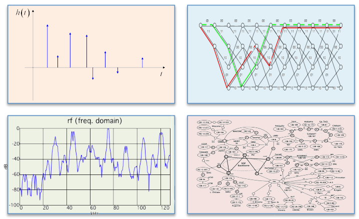


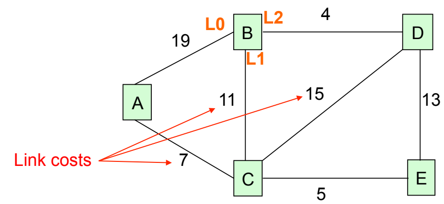
INTRODUCTION TO EECS II

DIGITAL COMMUNICATION SYSTEMS

6.02 Fall 2011
Lecture #20

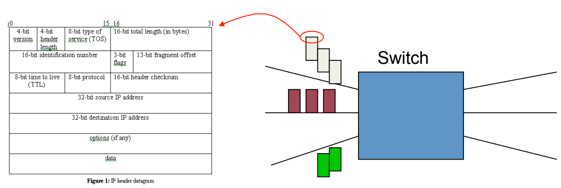
- addressing, forwarding, routing
- liveness, advertisements, integration
- distance-vector routing
- routing loops, counting to infinity

The Problem: Distributed Methods for Finding Paths in Networks



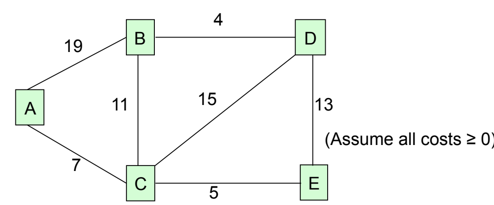
- Addressing (how to name nodes?)
  - Unique identifier for global addressing
  - Link name for neighbors
- Forwarding (how does a switch process a packet?)
- Routing (building and updating data structures to ensure that forwarding works)
- Functions of the **network layer**

Forwarding



- Core function is conceptually simple
  - `lookup(dst_addr)` in routing table returns *route* (i.e., *outgoing link*) for packet
  - `enqueue(packet, link_queue)`
  - `send(packet)` along outgoing link
- And do some bookkeeping before enqueue
  - Decrement hop limit (TTL); if 0, discard packet
  - Recalculate checksum (in IP, header checksum)

Shortest Path Routing



- Each node wants to find the path with *minimum total cost* to other nodes
  - We use the term "shortest path" even though we're interested in min cost (and not min #hops)
- Several possible distributed approaches
  - Vector protocols, esp. *distance vector* (DV)
  - *Link-state* protocols (LS)

Routing Table Structure

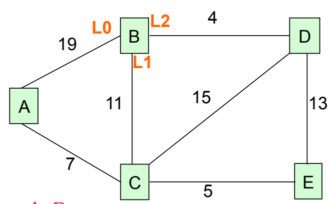


Table @ node B

Destination	**Link (next-hop)**	***Cost***
A	ROUTE L1	18
B	'Self'	0
C	L1	11
D	L2	4
E	L1	16

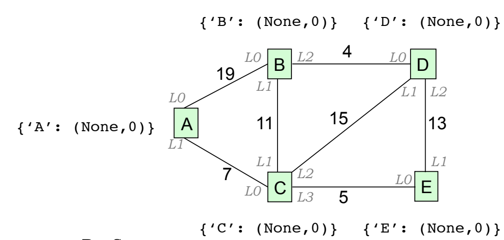
Distributed Routing: A Common Plan

- Determining live neighbors
  - Common to both DV and LS protocols
  - HELLO protocol (periodic)
    - Send HELLO packet to each neighbor to let them know who's at the end of their outgoing links
    - Use received HELLO packets to build a list of neighbors containing an information tuple for each link: (timestamp, neighbor addr, link)
    - Repeat periodically. Don't hear anything for a while → link is down, so remove from neighbor list.
- Advertisement step (periodic)
  - Send some information to all neighbors
  - Used to determine connectivity & costs to reachable nodes
- Integration step
  - Compute routing table using info from advertisements
  - Dealing with stale data

Distance-Vector Routing

- DV advertisement
  - Send info from routing table entries: (dest, cost)
  - Initially just (self,0)
- DV integration step [Bellman-Ford]
  - For each (dest,cost) entry in neighbor's advertisement
    - Account for cost to reach neighbor: (dest,my_cost)
    - my_cost = cost_in_advertisement + link_cost
  - Are we currently sending packets for dest to this neighbor?
    - See if link matches what we have in routing table
    - If so, update cost in routing table to be my_cost
  - Otherwise, is my_cost smaller than existing route?
    - If so, neighbor is offering a better deal! Use it…
    - update routing table so that packets for dest are sent to this neighbor

DV Example: round 1



Node A: update routes to B_B, C_C
Node B: update routes to A_A, C_C, D_D
Node C: update routes to A_A, B_B, D_D, E_E
Node D: update routes to B_B, C_C, E_E
Node E: update routes to C_C, D_D

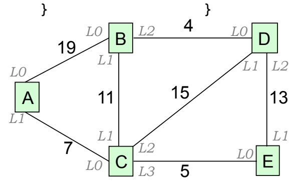
Subscript indicates node that gave better route

DV Example: round 2

```
{'A': (L0,19),
 'B': (None,0),
 'C': (L1,11),
 'D': (L2,4)
}
```

```
{'B': (L0,4),
 'C': (L1,15),
 'D': (None,0),
 'E': (L2,13)
}
```

```
{'A': (None,0),
 'B': (L0,19),
 'C': (L1,7)
}
```




```
{'A': (L0,7),
 'B': (L1,11),
 'C': (None,0),
 'D': (L2,15),
 'E': (L3,5)
}
```

```
{'C': (L0,5),
 'D': (L1,13),
 'E': (None,0)
}
```

Node A: update routes to B_C, D_C, E_C
Node B: update routes to A_C, E_C
Node C: no updates
Node D: update routes to A_C
Node E: update routes to A_C, B_C

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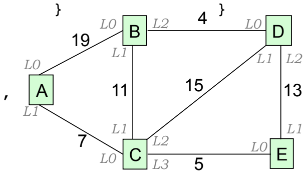
DV Example: round 3

```
{'A': (L1,18),
 'B': (None,0),
 'C': (L1,11),
 'D': (L2,4),
 'E': (L1,16)
}
```

```
{'A': (L1,22),
 'B': (L0,4),
 'C': (L1,15),
 'D': (None,0),
 'E': (L2,13)
}
```

```
{'A': (None,0),
 'B': (L1,18),
 'C': (L1,7),
 'D': (L1,22),
 'E': (L1,12)
}
```




```
{'A': (L0,7),
 'B': (L1,11),
 'C': (None,0),
 'D': (L2,15),
 'E': (L3,5)
}
```

```
{'A': (L0,12),
 'B': (L0,16),
 'C': (L0,5),
 'D': (L1,13),
 'E': (None,0)
}
```

Node A: no updates
Node B: no updates
Node C: no updates
Node D: no updates
Node E: no updates

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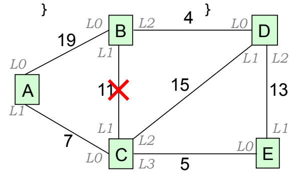
DV Example: Break a Link

~~{'A': (L1,18),~~
'B': (None,0),
~~'C': (L1,11),~~
'D': (L2,4),
~~'E': (L1,16)~~
}

```
{'A': (L1,22),
 'B': (L0,4),
 'C': (L1,15),
 'D': (None,0),
 'E': (L2,13)
}
```

```
{'A': (None,0),
 'B': (L1,18),
 'C': (L1,7),
 'D': (L1,22),
 'E': (L1,12)
}
```



{'A': (L0,7),
~~'B': (L1,11),~~
'C': (None,0),
'D': (L2,15),
'E': (L3,5)
}

```
{'A': (L0,12),
 'B': (L0,16),
 'C': (L0,5),
 'D': (L1,13),
 'E': (None,0)
}
```

When link breaks: eliminate routes that use that link.

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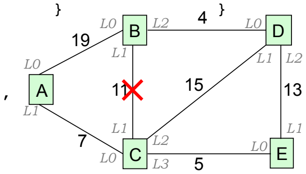
DV Example: round 4

```
{'A': (None,∞),
 'B': (None,0),
 'C': (None,∞),
 'D': (L2,4),
 'E': (None,∞)
}
```

```
{'A': (L1,22),
 'B': (L0,4),
 'C': (L1,15),
 'D': (None,0),
 'E': (L2,13)
}
```

```
{'A': (None,0),
 'B': (L1,18),
 'C': (L1,7),
 'D': (L1,22),
 'E': (L1,12)
}
```




```
{'A': (L0,7),
 'B': (None,∞),
 'C': (None,0),
 'D': (L2,15),
 'E': (L3,5)
}
```

```
{'A': (L0,12),
 'B': (L0,16),
 'C': (L0,5),
 'D': (L1,13),
 'E': (None,0)
}
```

Node A: update cost to B_C
Node B: update routes to A_A, C_D, E_D
Node C: update routes to B_D
Node D: no updates
Node E: update routes to B_D

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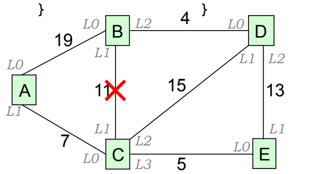
DV Example: round 5

```
{'A': (L0,19),     {'A': (L1,22),
 'B': (None,0),     'B': (L0,4),
 'C': (L2,19),      'C': (L1,15),
 'D': (L2,4),       'D': (None,0),
 'E': (L2,17)       'E': (L2,13)
}                  }
```

Update cost

```
{'A': (None,0),
 'B': (L1, ∞),
 'C': (L1,7),
 'D': (L1,22),
 'E': (L1,12)
}
```



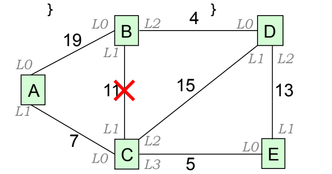

```
{'A': (L0,7),      {'A': (L0,12),
 'B': (L2,19),      'B': (L1,17),
 'C': (None,0),     'C': (L0,5),
 'D': (L2,15),      'D': (L1,13),
 'E': (L3,5)        'E': (None,0)
}                  }
```

Node A: update route to B_B
Node B: no updates
Node C: no updates
Node D: no updates
Node E: no updates

DV Example: final state

```
{'A': (L0,19),     {'A': (L1,22),
 'B': (None,0),     'B': (L0,4),
 'C': (L2,19),      'C': (L1,15),
 'D': (L2,4),       'D': (None,0),
 'E': (L2,17)       'E': (L2,13)
}                  }
```

```
{'A': (None,0),
 'B': (L0,19),
 'C': (L1,7),
 'D': (L1,22),
 'E': (L1,12)
}
```




```
{'A': (L0,7),      {'A': (L0,12),
 'B': (L2,19),      'B': (L1,17),
 'C': (None,0),     'C': (L0,5),
 'D': (L2,15),      'D': (L1,13),
 'E': (L3,5)        'E': (None,0)
}                  }
```

Node A: no updates
Node B: no updates
Node C: no updates
Node D: no updates
Node E: no updates

Correctness & Performance

- Optimal substructure property fundamental to correctness of both Bellman-Ford and Dijkstra's shortest path algorithms
  - ***Suppose shortest path from X to Y goes through Z. Then, the sub-path from X to Z must be a shortest path.***
- Proof of Bellman-Ford via induction on number of walks on shortest (min-cost) paths
  - Easy when all costs > 0 and *synchronous model* (see notes)
  - Harder with distributed async model (not in 6.02)
- How long does it take for distance-vector routing protocol to *converge*?
  - Time proportional to largest number of hops considering all the min-cost paths

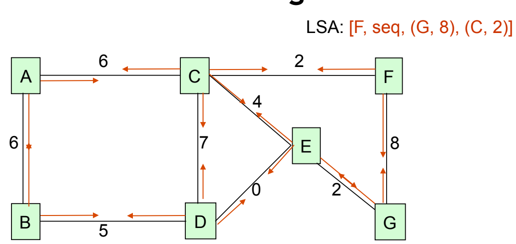
Link-State Routing

- Advertisement step
  - Send information about its links to its neighbors (aka **link state advertisement** or LSA):

    [seq#, [(nbhr1, linkcost1), (nbhr2, linkcost2), …]

  - Do it periodically (liveness, recover from lost LSAs)
- Integration
  - If seq# in incoming LSA > seq# in saved LSA for source node:
    update LSA for node with new seq#, neighbor list
    rebroadcast LSA to neighbors (→ **flooding**)
  - Remove saved LSAs if seq# is too far out-of-date
  - Result: Each node discovers current map of the network
- Build routing table
  - Periodically each node runs the same *shortest path algorithm* over its map (e.g., Dijkstra's alg)
  - If each node implements computation correctly and each node has the same map, then routing tables will be correct

LSA Flooding



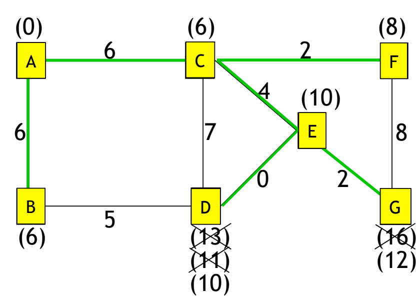
- Periodically originate LSA
- LSA travels each link in each direction
  - Don't bother with figuring out which link LSA came from
- Termination: each node rebroadcasts LSA exactly once
  - Use sequence number to determine if new, save latest seq
- Multiple opportunities for each node to hear any given LSA
  - Time required: number of links to cross network

Dijkstra's Shortest Path Algorithm

- Initially
  - nodeset = [all nodes] = set of nodes we haven't processed
  - spcost = {me:0, all other nodes: ∞} # shortest path cost
  - routes = {me:--, all other nodes: ?} # routing table
- while nodeset isn't empty:
  - find u, the node in nodeset with smallest spcost
  - remove u from nodeset
  - for v in [u's neighbors]:
    - d = spcost(u) + cost(u,v) # distance to v via u
    - if d < spcost(v): # we found a shorter path!
      - spcost[v] = d
      - routes[v] = routes[u] (or if u == me, enter link from me to v)
- Complexity: N = number of nodes, L = number of links
  - Finding u (N times): linear search=O(N), using heapq=O(log N)
  - Updating spcost: O(L) since each link appears twice in neighbors

Integration Step: Dijkstra's Algorithm (Example)

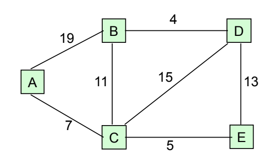
Suppose we want to find paths from A to other nodes



Another Example

Finding shortest paths from A:

LSAs:
```
A: [(B,19), (C, 7)]
B: [(A,19), (C,11), (D, 4)]
C: [(A, 7), (B,11), (D,15), (E, 5)]
D: [(B, 4), (C,15), (E,13)]
E: [(C, 5), (D,13)]
```



Step	u	Nodeset	spcost					route				
			A	B	C	D	E	A	B	C	D	E
0		[A,B,C,D,E]	0	∞	∞	∞	∞	--	?	?	?	?
1	A	[B,C,D,E]	0	19	7	∞	∞	--	L0	L1	?	?
2	C	[B,D,E]	0	18	7	22	12	--	L1	L1	L1	L1
3	E	[B,D]	0	18	7	22	12	--	L1	L1	L1	L1
4	B	[D]	0	18	7	22	12	--	L1	L1	L1	L1
5	D	[]	0	18	7	22	12	--	L1	L1	L1	L1