

Recitation 18

Today: ① MAC Protocols

② Filters (if time permits)

Examples of Shared Media: ① satellite network ② wireless (radio) network ③ Wired Bus (e.g. Ethernet) ④ FDM radio & TV broadcast

MAC Protocol: Defines rules governing how a common (shared) communication channel (medium) is shared among different nodes.

Performance Metrics:

① Throughput (bps) ~ Num packets delivered successfully per unit time (s).

* (A) Utilization ~ (independent of channel rate): ratio of total throughput to max data rate of channel.
($0 \leq U \leq 1$)

(eg1) Given: channel rate ~ 10 Mbps.

~ 4 nodes share channel with throughputs 2Mbps, 2Mbps, 2Mbps, 1Mbps

Req'd: Utilization?

$$\text{Soln: } U = \frac{\text{Total Throughput}}{\text{Max Channel Rate}} = \frac{(2+2+2+1)}{10} = \frac{7}{10} = 0.7$$

* (B) Fairness ~ notions of fairness?

"Allocation with smaller g.d. fairer than one with larger standard deviation"

$$\text{Fairness Index } F_i = \frac{\left(\sum_N x_i\right)^2}{N \sum_N x_i^2}, \quad x_i \text{ is throughput of node } i \text{ and } N \text{ backlogged nodes.}$$

(eg2) For eg1 above, what is F-Index? $F = \frac{(2+2+2+1)^2}{4 \cdot (2^2+2^2+2^2+1^2)} = \frac{49}{52} = 0.94$ Fair? $\frac{1}{N} \leq F \leq 1$
(Yes ~ 1-perfect fairness)

2. TDMA

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A simple TDMA protocol:

If the current time is t , and N nodes share the channel, then the node with ID i transmits if and only if it is backlogged and $(t \bmod N) = i$ $i \in \{0, N-1\}$

Pros:

- It is fair
- Simple to implement if there is a central coordinator.

Cons:

- under-utilize the medium in skewed traffic pattern
- become complex if there is not a central coordinator (i.e. distributed)

3. Aloha

Protocol:

If a node is backlogged, it sends a packet with probability p .

Utilization:

count slots without collisions! (Exactly one node transmits)

$$U(p) = P(\text{Exactly one node sends})$$

$$= N p (1-p)^{N-1}$$

where N is the number of backlogged nodes
 p each node transmits with probability p .

The maximum value of U occurs

$$\text{when } p = \frac{1}{N}$$

and the maximum value of U is

$$U_{\max} = \left(1 - \frac{1}{N}\right)^{N-1}$$

As $N \rightarrow \infty$, $U_{\max} \rightarrow \frac{1}{e} \approx 37\%$

Example: Given: A and B share a channel.

The probability that A sends a packet is: P_a

The probability that B sends a packet is P_b .

A is $\frac{2}{3}$ as likely to send a packet as B,

$$\text{i.e. } P_a = \frac{2}{3} P_b$$

Assume both A and B are backlogged and each packet is one slot long.

Question: ① what is the utilization?

② what is P_a and P_b such that the utilization is maximized?

③ what is maximized utilization?

Solution:

① utilization is the probability that exactly one user sends.

$$\text{So: } U = P(A \text{ xmits}) * P(B \text{ doesn't xmits}) \\ + P(B \text{ xmits}) * P(A \text{ doesn't xmits})$$

$$= P_a * (1 - P_b) + P_b * (1 - P_a) \quad \left. \begin{array}{l} \\ \end{array} \right\} \text{ plug } P_a = \frac{2}{3} P_b$$

$$= \frac{2}{3} P_b (1 - P_b) + P_b (1 - \frac{2}{3} P_b)$$

$$= \frac{5}{3} P_b - \frac{4}{3} P_b^2$$

② the utilization U is maximized when $\frac{dU}{dP_b} = 0$.

$$\text{So: } 0 = \frac{dU}{dP_b} = \frac{d\left(\frac{5}{3}P_b - \frac{4}{3}P_b^2\right)}{dP_b}$$

$$= \frac{5}{3} - \frac{8}{3}P_b$$

$$\Rightarrow P_b = \frac{5}{8}, \quad P_a = \frac{2}{3}P_b = \frac{2}{3} * \frac{5}{8} = \frac{5}{12}$$

③ Plug in $P_b = \frac{5}{8}$, we got maximized U

$$U_{\text{max}} = \frac{5}{3}P_b - \frac{4}{3}P_b^2$$

$$= \frac{5}{3} * \frac{5}{8} - \frac{4}{3} * \left(\frac{5}{8}\right)^2$$

$$\approx 0.52$$

Example 2: Given: A and B shares a channel.

The probability that A sends packet is: $P_a = 0.3$.

We want B get three time the throughput as A whenever both nodes are backlogged.

Question? what is the transmission probability of B, P_b in order to achieve our throughput goal?

Solution:

$$\text{A's throughput} = P(\text{A xmits}) * P(\text{B doesn't})$$

$$= 0.3 * (1 - P_b)$$

$$\text{B's throughput} = P(\text{B xmits}) * P(\text{A doesn't})$$

$$= P_b * (1 - P_a)$$

$$= P_b * (1 - 0.3)$$

$$= 0.7 P_b$$

$$\text{B's throughput} = 3 * (\text{A's throughput})$$

$$\Rightarrow 0.7 P_b = 3 * [0.3 * (1 - P_b)] \Rightarrow 1.6 P_b = 0.9$$

$$\Rightarrow P_b \approx 0.56$$

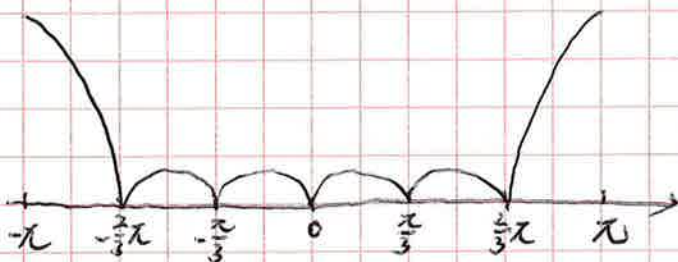
Filter:

Given: Consider three LTI filters, denoted by A, B, C.

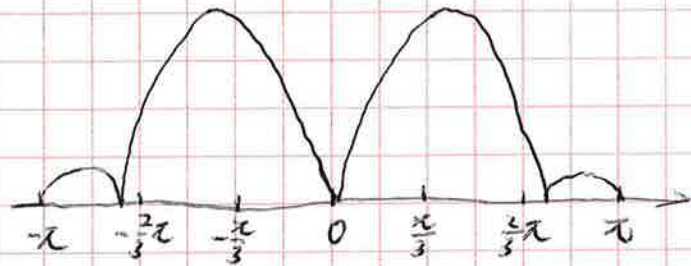
The magnitudes of their frequency response $|H_A(e^{j\omega})|$,

$|H_B(e^{j\omega})|$, $|H_C(e^{j\omega})|$ are given below.

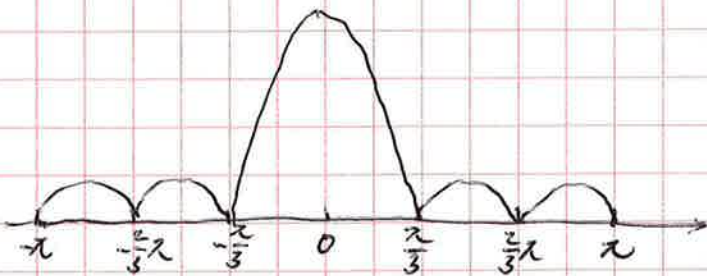
① $|H_A(e^{j\omega})|$



② $|H_B(e^{j\omega})|$

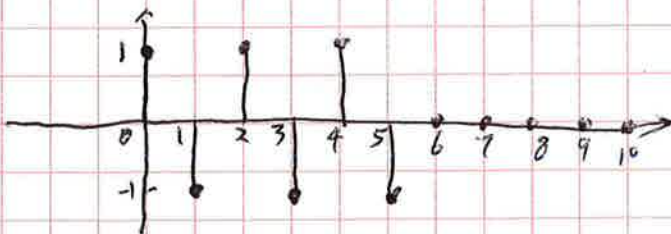


③ $|H_C(e^{j\omega})|$



Question:

which frequency response of filters (A, B, or C) correspondingly to the following unit sample response?



Solution:

(Evaluate $|H(e^{j\omega})|$ at several frequencies and see what we found)

Noticed: $h[0]=1, h[1]=-1, h[2]=1$

$h[3]=-1, h[4]=1, h[5]=-1, h[m]=0$ otherwise.

Then, $H(z) = \sum h[m] e^{-j\omega m}$

At $\omega=0, |H(e^{j\omega})| = |H(e^{j0})| = \sum h[m] e^{-j0 \cdot m}$

$\rightarrow e^{-j0 \cdot m} = e^0 = 1$

$= \sum h[m]$

$= h[0] + h[1] + h[2] + h[3] + h[4] + h[5]$

$= 0$

At $\omega=\pi, |H(e^{j\omega})| = |H(e^{j\pi})| = \sum h[m] e^{-j\pi \cdot m}$

$\rightarrow e^{-j\pi m} = (-1)^m$

$= \sum h[m] \cdot (-1)^m$

$= h[0] - h[1] + h[2] - h[3] + h[4] - h[5]$

$= 6$

The only plot with a frequency response of 0 at frequency 0,

and non-zero value at frequency π is the one for filter A.

So: The filter corresponding is A.