

INTRODUCTION TO EECS II DIGITAL COMMUNICATION SYSTEMS

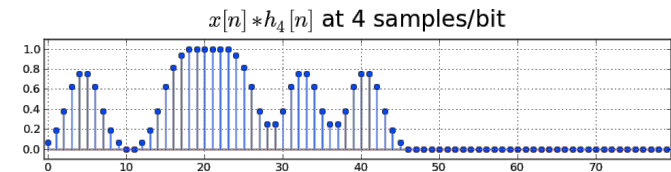
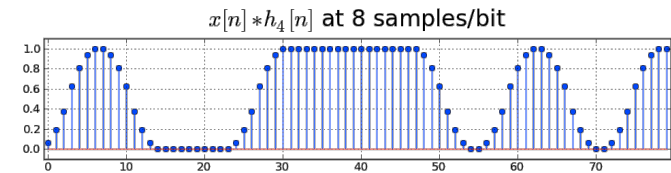
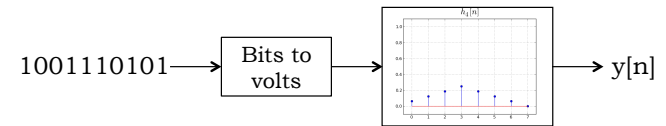
6.02 Spring 2011 Lecture #5

- Intersymbol interference
- Deconvolution
- Stability & noise, approximate deconvolvers

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Lecture 5, Slide #1

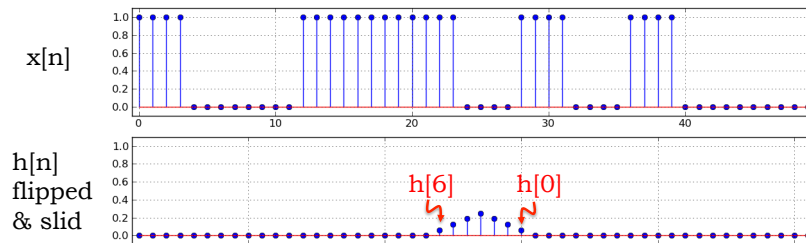
Transmission Over a Channel



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Lecture 5, Slide #2

Convolution sum: “flip and slide”



$$y[28] = x[28]h[0] + x[27]h[1] + \dots + x[22]h[6]$$

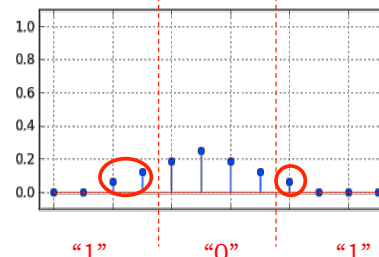
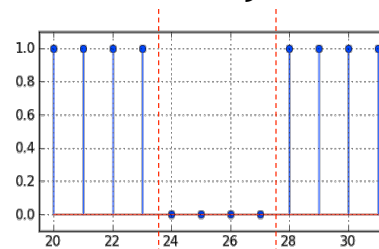
Visual representation of convolution sum: do a horizontal flip of the graph of $h[n]$, then slide along under $x[n]$.

To compute $y[m]$, slide flipped $h[n]$ until $h[0]$ is under $x[m]$, then compute sum of element-by-element product of the two sequences.

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Lecture 5, Slide #3

Intersymbol Interference (ISI)



“1” “0” “1”

Issue:

If we send a small number of samples/bit, the active portion of $h[n]$ may cover more than one bit cell when doing convolution sum.

Result:

$y[n]$ values for a particular bit cell include contributions from neighboring cells.

Example: $y[28]$ is the lowest voltage received for the “0” bit, but includes contributions from the neighboring “1” bits.

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Lecture 5, Slide #4

Given $h[n]$, how bad is ISI?

Recipe:

1. Compute B , the number bits “covered” by $h[n]$. Let $N =$ samples/bit

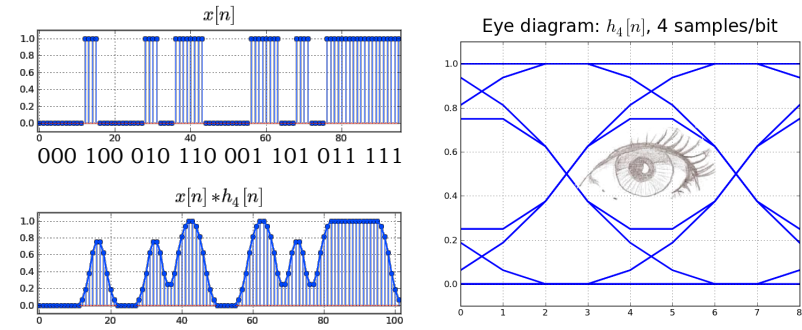
$$B = \left\lfloor \frac{\text{length of active portion of } h[n]}{N} \right\rfloor + 2$$
2. Generate a test pattern that contains all possible combinations of B bits – want all possible combinations of neighboring cells.
 If B is big, randomly choose a large number of combinations.
3. Transmit the test pattern over the channel ($2^N \cdot B$ samples)
4. Instead of one long plot of $y[n]$, plot the response as an *eye diagram*:
 - a. break the plot up into short segments each containing $2N+1$ samples, starting at sample 0, N , $2N$, $3N$, ...
 - b. plot all the short segments on top of each other

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Lecture 5, Slide #5

Eye Diagram Example

Using $h_4[n]$ and samples_per_bit=4: $B = 3$

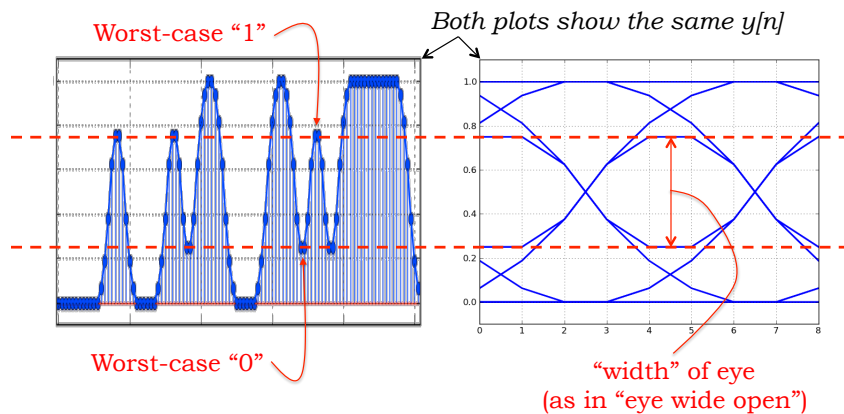


Eye diagrams make it easy to find the worst-case signaling conditions at the receiving end.

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Lecture 5, Slide #6

“Width” of Eye



To maximize noise margins:

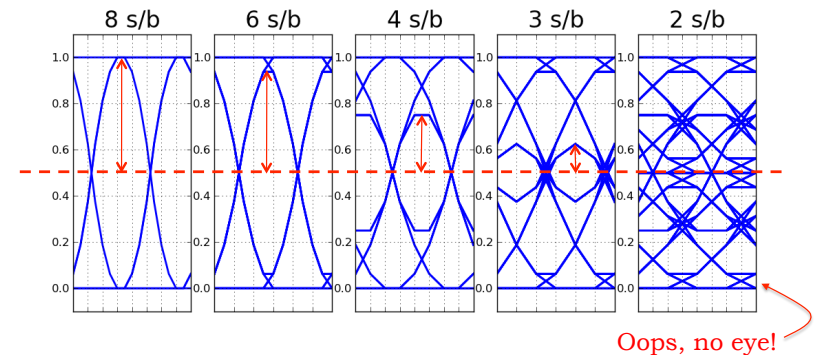
Pick the best sample point → widest point in the eye

Pick the best digitization threshold → half-way across width

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Lecture 5, Slide #7

Choosing Samples/Bit



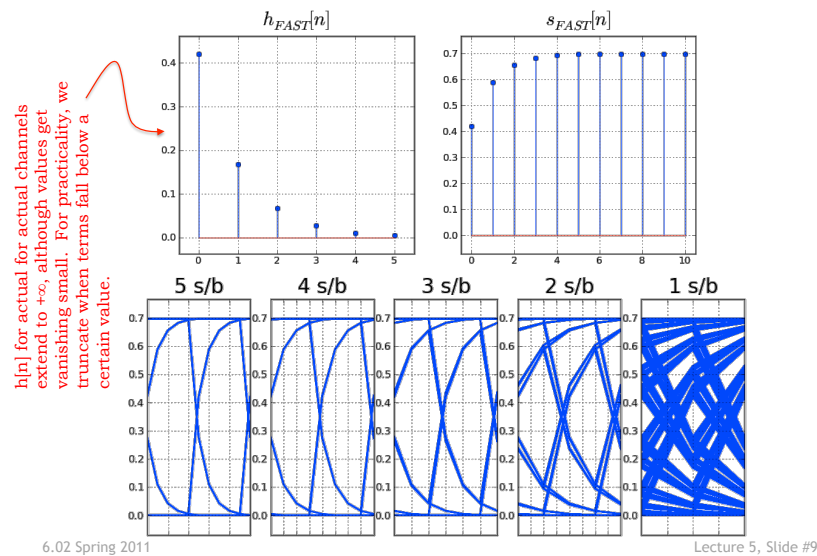
Given $h[n]$, you can use the eye diagram to pick the number of samples transmitted for each bit (N):

Reduce N until you reach the noise margin you feel is the minimum acceptable value.

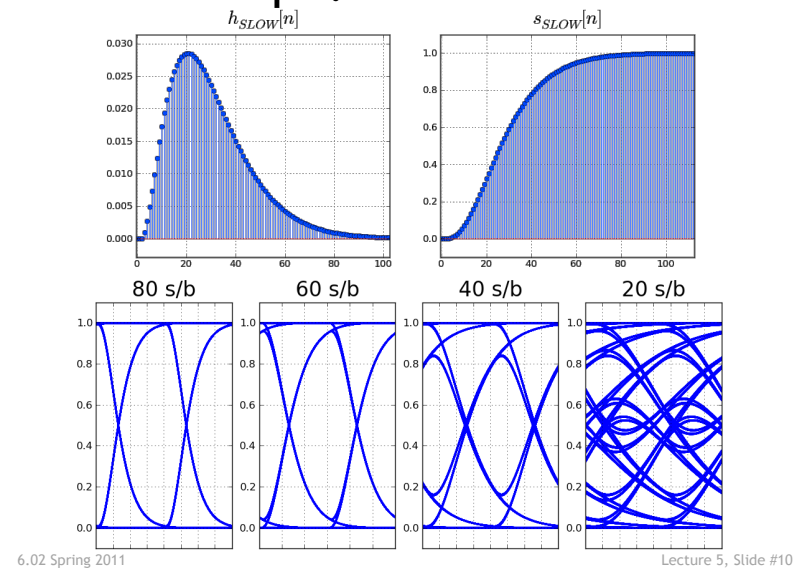
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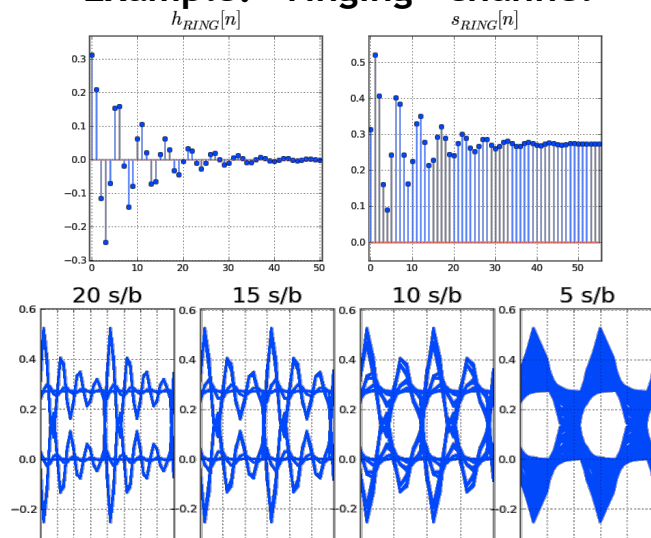
Example: “fast” channel



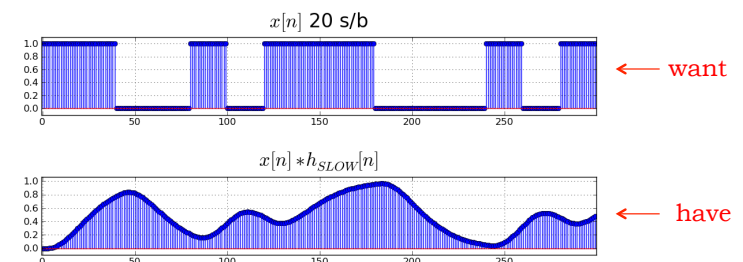
Example: “slow” channel



Example: “ringing” channel



Can We Recover From ISI?



After all, in a perfect world (no noise), no information has been lost, only spread out over many samples.

Given $y[n]$ and $h[n]$, can we develop an estimate $w[n]$ for the actual input waveform $x[n]$? We could, of course, easily receive $x[n]$!

Difference Equation for $w[n]$

If $w[n]$ was a perfect estimate of $x[n]$, it would satisfy:

$$y[n] = w[n]h[0] + w[n-1]h[1] + w[n-2]h[2] + \dots + w[n-K]h[K]$$

Simplifying assumption: $h[K]$ is last non-zero element ↴

Let's solve this for $w[n]$:

$$w[n] = \frac{1}{h[0]} \left(y[n] - (w[n-1]h[1] + w[n-2]h[2] + \dots + w[n-K]h[K]) \right)$$

Given $y[n]$ and $h[n]$, we can incrementally compute sequence $w[n]$ using a straightforward “plug and chug” approach:

$$w[0] = \frac{1}{h[0]} (y[0])$$

$$\begin{matrix} h[i] = 0 & i < 0 \text{ or } i > K \\ w[j] = 0 & j < 0 \end{matrix}$$

$$w[1] = \frac{1}{h[0]} (y[1] - w[0]h[1])$$

$$w[2] = \frac{1}{h[0]} (y[2] - w[1]h[1] - w[0]h[2])$$

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Lecture 5, Slide #13

What if $h[0]=0$?

$$w[n] = \frac{1}{h[0]} \left(y[n] - (w[n-1]h[1] + w[n-2]h[2] + \dots + w[n-K]h[K]) \right)$$

Oops! Division by 0 isn't a good idea...

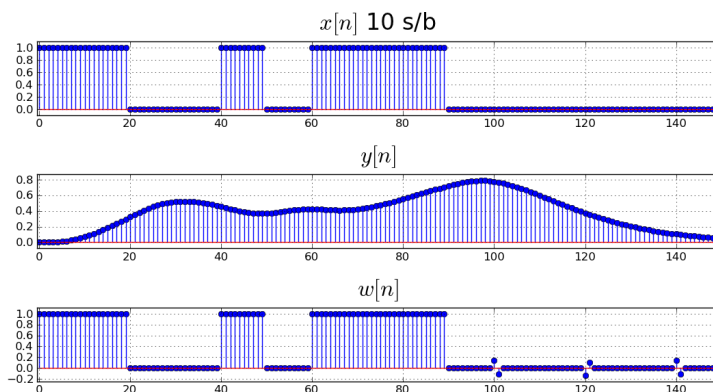
Zeros at the beginning $h[n]$ represent a channel with a delay: m zeros would mean a m -sample delay. We can eliminate the delay without affecting our estimate for $x[n]$. So

1. Count the number of zeros at the front of $h[n] = m$
2. Eliminate the first m elements of $h[n]$, and eliminate the first m elements of $y[n]$
3. Now use the equation above on the shortened $h[n]$ and $y[n]$

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Lecture 5, Slide #14

Deconvolution Example



???
(hint: see slide #10)

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Lecture 5, Slide #15

Sensitivity to Noise

Let's consider what happens if some small amount of noise (ϵ) is added to the first sample of the response ($y[0]$):

Estimate	Error
$w[0] = \frac{1}{h[0]} (y[0] + \epsilon)$	$\frac{\epsilon}{h[0]}$
$w[1] = \frac{1}{h[0]} (y[1] - w[0]h[1])$	$-\frac{\epsilon}{h[0]} \frac{h[1]}{h[0]}$
$w[2] = \frac{1}{h[0]} (y[2] - w[1]h[1] - w[0]h[2])$	$-\frac{\epsilon}{h[0]} \left(\frac{h[1]}{h[0]} \right)^2 - \frac{\epsilon}{h[0]} \frac{h[2]}{h[0]}$

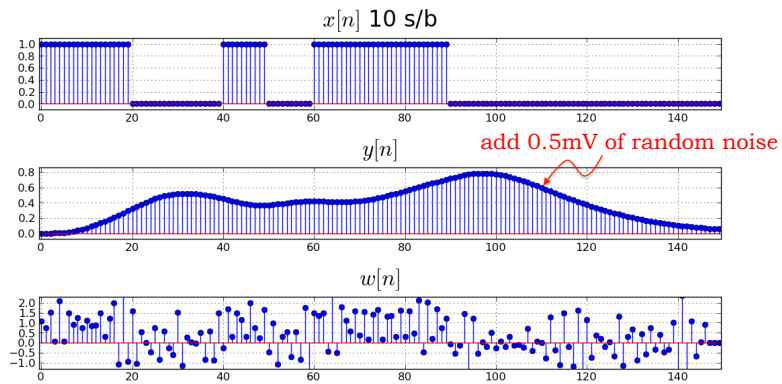
Question: is the error growing as we compute more w 's?

Answer: depends on $h[0]$ and the ratios $h[m]/h[0]$. Small values of $h[0]$ and $(h[m]/h[0]) > 1$ are troublesome...

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Lecture 5, Slide #16

Noisy Deconvolution Example



Urkl!

Stability Criterion

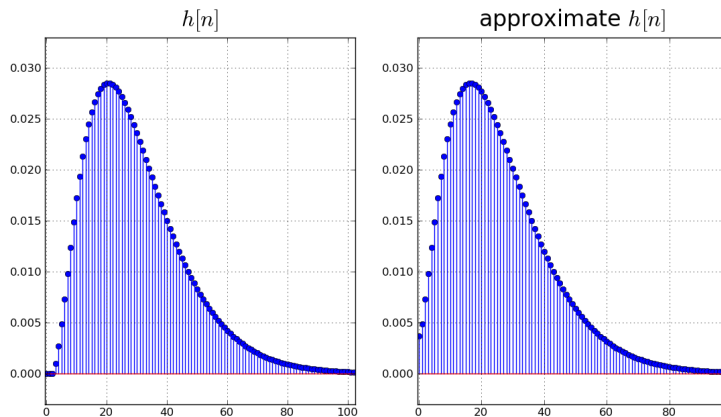
The notes have a derivation of the following sufficient (very conservative) condition that will ensure the stability of the deconvolver operating on a noisy $y[n]$:

$$\sum_{m=1}^K \left| \frac{h[m]}{h[0]} \right| < 1 \quad \text{or, perhaps more usefully} \quad \sum_{m=1}^K |h[m]| < |h[0]|$$

What if my $h[n]$ doesn't meet this criterion?

Make a new "approximate" $h[n]$ that does! Combine samples at the beginning of $h[n]$ to make a bigger $h[0]$.

Example Approximate $h_{\text{SLOW}}[n]$



Approximation: combine first 5 samples of $h_{\text{SLOW}}[n]$

(Less) Noisy Deconvolution Example

