DNS

What is the relationship between a domain name (e.g., youtube.com) and an IP address?

DNS is the system that determines this mapping.

Basic idea:

You contact a DNS server, give it a query

It responds with the answer, or forwards your answer on to some other server

Example:

Whenever a server answers a query, it adds it to a cache

When you get a response, you add it to a cache

Each cached value has a duration after which you are supposed to look up the name again

At the top of the DNS hierarchy are root servers

Then servers for .com, .edu, etc ...

Then servers for MIT, youtube, etc....

Show cached

dig +norec csail.mit.edu
A record vs NS Record

TTL

Observe that A records expire much sooner than NS Records

Example:

csail.mit.edu
dig
dig @E.ROOT-SERVERS.NET edu
dig @G.GTLD-SERVERS.NET mit.edu
dig @bitsy.mit.edu csail.mit.edu

How do you get your DNS server initially?

- Your admin tells it to you
- DHCP (Domain Host Configuration Protocol) -- you send out a broadcast on the local link -- DHCP server responds (via another broadcast message) with:
  - An IP address for you to use
  - The IP address of a "gateway" router for you to use
  - The IP address of a DNS server to use
  - ....

DNS has some fancy features:

- One physical machine can have multiple names (web.mit.edu, www.mit.edu)
- One name can correspond to multiple IP addresses
  DNS load balancing
  DNS server picks which name to return
  most DNS servers cycle through these in "round robin" fashion

Example (multiple times)
> dig yahoo.com
yahoo.com. 13810 IN A 72.30.2.43
yahoo.com. 13810 IN A 98.137.149.56
yahoo.com. 13810 IN A 209.131.36.159
yahoo.com. 13810 IN A 209.191.93.53
yahoo.com. 13810 IN A 209.191.122.70
yahoo.com. 13810 IN A 67.195.160.76

MX records -- when I send email to mit.edu, I ask DNS for an MX record, which gives me the name of a mail server for MIT. DNS is a fairly general system, in the sense that it can have many different record types.
dig mit.edu mx

dig nytimes.com all

**Content delivery networks (CDNs -- e.g., Akamai)**
Use DNS to provide scalability and adapt to load

Suppose I have some content that is accessed a lot -- e.g., a video on youtube

I can balance load amongst my local servers using DNS load balancing, but I still have to own the servers.

Puts a huge load on my servers to deliver it; can't adapt to load spikes (e.g., the "slashdot" effect.)

Also, for users that are far away (e.g., in Asia or Australia), they have to download that content over long distance and thin pipes, and ISPs in Australia have to pay a lot for than bandwidth. (How expensive -- apparently around .10/GB for ISP to deliver data to you-- according to David Clark -- [http://web.mit.edu/newsoffice/2009/fcc-neutrality.html](http://web.mit.edu/newsoffice/2009/fcc-neutrality.html))

For content -- **like this video** -- that is accessed repeatedly, would be be better to not have to go all the way to San Bruno CA everytime. (55M views) Youtube pays $1M / day for bandwidth!

Idea: create a local cache

Solution 1: Proxy cache. For every URL, look up in a local cache (perhaps run by your ISP) to see if the content is there. If it is, fetch it. Otherwise, get original data. Just like DNS requests, web pages can have cache lifetimes associated with them which proxy caches respect.

Problems:

- Helps ISP but not necessarily content provider
- requires clients to configure their browsers to use proxy caches
- Doesn't address slashdot effect

Exist so called "transparent proxies" that can do this filtering automatically, but this may be distasteful to users, especially if their ISP is doing it (companies do this all the time.)
Solution 2: Content Distribution Networks (e.g., Akamai)

Show diagram:

```
  DNS  akamai.net
     |     |
     v     v
  akamai.net
     |     |
     v     v
  akamai.net
     |
     v
nytimes.com
```

Give example:
- nytimes.com
- dig graphics8.nytimes.com

From both MIT network and cellular network

traceroute 18.7.20.70 (cellular)

Observe that TTL for answer is very short -- Why? -- Handle slashdot effect -- can dynamically start using more and more akamai servers for a particular request

Akamai -- company

Gives a way for content provider to offload load from server

Also helps ISP if server is inside ISP (creates an incentive for ISPs to participate!)

Akamai has thousands of caches all over the world

When a request for content -- like from images8.nytimes.com arrives -- it uses dns to forward user to the nearest server.

Determining the "nearest" server and how many servers to allocate is their secret sauce.

For dynamic content, Akamai also works.
Example -- show attempt to select best route back to Akamai, maintenance of reachability info, etc.

Claim Boston to Beijing has 237 ms latency on Akamai vs 318 on public internet

Akamai is a more general example of something called an **overlay network**.

An overlay is a way to create a network with new features or a different structure by building on top of an existing network.

Akamai creates an overlay on top of the IP network that chooses the best route from amongst a collection of IP servers.

Overlays are used widely to extend the network with new features, for example:

- Provide a different topology (e.g., direct connection of clients)
  - E.g., for administrative reasons (VPNs)
  - Or for performance reasons (Akamai)
- Provide a different addressing mechanism (e.g., content addressability, P2P)

VPN example:

Companies want to allow remote users or sites to have access to their corporate Internet sites

Good old days might have done this with a modem, but that's slow, and a pain to run.

VPN provides a way for remote users to appear to be on internal network while actually being external.
Idea is to "tunnel" traffic over public internet using an overlay (client and server inside of network):

Diagram:

This is a simple example, but in principle can create complex multi-hop VPNs.

Recitation will talk about an idea like this for dealing with mobile clients.