Fault isolation with VM (2/15/2004)

- x86 look more thin
- need 5 minutes, set first page
- or make x86 a running example

1. Effecting modularity
   - module per computer
   - one per virtual computer
   - virtual - viewable program

2. Address space
   - possible for C to link Ps returns
   - (and vice versa)

3. Use VM (6-c00)
   - each module brings device
   - register defines which process
   - p + c can only go what in page map
   - simple thread

4. Case study: x86/PC
   - fault isolation
   - h/s/3 interface
   - user addressable
   - protection function
   - cost
   - evolution

5. Kernel address space
   - supervisor call
   - internal tables of techniques
   - (share privilege, include
   - new pas, p store
   - also for interrupt
   - illegal instruction)

6. Entry level
   - 1. U -> K
   - 2. Change privilege level
   - 3. Save PC loaded in RC

Who manages page maps?
- in management
- in management