0. Introduction
- Currently: building reliable systems out of unreliable components. We're working on implementing transactions which provide
  - atomicity
  - isolation
- So far: have a poorly-performing version of atomicity via shadow copies.
- Today: Logging, which will give us reasonable performance for atomicity. Logging also works when we have multiple concurrent transactions, even though for today we're not thinking about concurrency.

1. Motivating example
- begin      // T1
  A = 100
  B = 50
  commit     // At commit: A=100; B=50

begin      // T2
  A = A - 20
  B = B + 20
  commit     // At commit: A=80; B=70

begin      // T3
  A = A + 30
  --CRASH--
- Problem: A = 110, but T3 didn't commit. We need to revert.

2. Basic idea
- Keep a log of all changes and whether a transaction commits or aborts
  - every transaction gets a unique ID
  - UPDATE records include old and new values of a variable
  - COMMIT records specify that transaction committed
  - ABORT records specify that transaction aborted
  - Not always needed
  - (See slides for the log for this example)
- Nice: updates are small appends

3. How to use a log for transactions
- On begin: allocate new transaction ID (TID)
- On write: append entry to log
- On read: scan log to find last committed value
- On commit: write commit record
- This is the commit point
- Atomic because we can assume it's a single-sector write
- Another way to do it would be to put checksums on each record and ignore partially-written records
- On abort: nothing (could write an ABORT record but not strictly needed)
- On recover: nothing
- (see slide for code)

4. Performance of log
   - Writes: good. sequential = fast.
   - Reads: terrible. Must scan entire log.
   - Recovery: instantaneous

5. Cell Storage
   - Improve read performance with cell storage.
   - (For us) stored on disk, i.e., non-volatile storage
   - Updates go to log and cell storage
   - Read from cell storage
   - "Log" = write to log. "Install" = write to cell storage
   - How to recover
     - Scan the log backwards, determine what actions aborted, and undo them
     - (see slide for code)
     - What if we crash during recovery? No worries; recover() is idempotent. Can do it repeatedly.
   - How to write
     - Log before install, not the other way; otherwise, can't recover from a crash in between the two writes.
     - This is write-ahead logging

6. Performance of log + cell storage
   - Writes: Okay, but now we write to disk twice instead of once
   - Reads: fast
   - Recovery: Bad. Have to scan the entire log.

7. Improving performance
   - Improve writes: use a (volatile) cache
     - Reads go to cache first, writes go to cache and are eventually flushed to cell storage
     - Problem: After crash, there may be updates that didn't make it to cell storage (were in cache but not flushed)
     - Also could be updates in cell storage that need to be undone, but we had that problem before
     - Solution: We need a redo phase in addition to an undo phase in our recovery (see slide for code)
   - Improving recovery
     - Problem: recovery takes longer and longer as the log grows
     - Solution: truncate the log
     - How?
- Assuming no pending actions
  - Flush all cached updates to cell storage
  - Write a CHECKPOINT record
  - Truncate the log prior to the CHECKPOINT record
    - Usually amounts to deleting a file
  - With pending actions, delete before the checkpoint and earliest undecided record.
- ABORT records
  - Can be used to help recovery and skip undo-ing aborted transaction. Not necessary for correctness -- can always just pretend we crashed -- but can help.

8. What about un-undo-able actions?
- What if our transaction fires a missile and then aborts?
  - Typically: wait for software that controls the action to commit and then take the action, but have a special way to detect whether the action has/will happened

9. Summary
- Logging is a general technique for achieving atomicity
  - Writes are fast, reads can be fast with cell storage
  - Need to log before installing (write-ahead), and need a recovery process
- Tomorrow is recitation: logging for file systems
- Now: we're good with atomicity
  - In fact, logging will work fine with concurrent transactions; the problem will be figuring out which steps we can actually run in parallel
- Wednesday: isolation
- Next week: distributed transactions