Finite State Machines

- Design methodology for sequential logic
  -- identify distinct states
  -- create state transition diagram
  -- choose state encoding
  -- write combinational Verilog for next-state logic
  -- write combinational Verilog for output signals
- Lots of examples

Module Instantiation

- Given submodule mux32two:

2-to-1 MUX

```
module mux32two
     (input [31:0] i0,i1,
      input sel,
      output [31:0] out);

   assign out = sel ? i1 : i0;
endmodule
```

- Instantiation of mux32two

```
module zyz (input xin,... output yout,...)

  mux32two adder_mux(.i0(b), .i1(32'd1), .sel(f[0]), .out(addmux_out));
  mux32two sub_mux(.i0(b), .i1(32'd1), .sel(f[0]), .out(submux_out));
```

module names

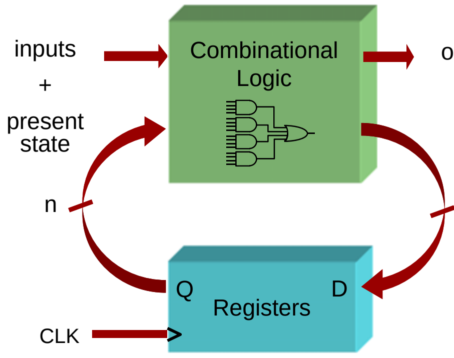
(unique) instance names

corresponding inputs/outputs for mux32two explicitly declared for each instance

---

Finite State Machines

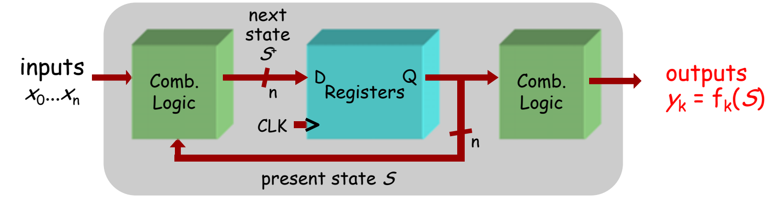
- Finite State Machines (FSMs) are a useful abstraction for sequential circuits with centralized "states" of operation
- At each clock edge, combinational logic computes *outputs* and *next state* as a function of *inputs* and *present state*



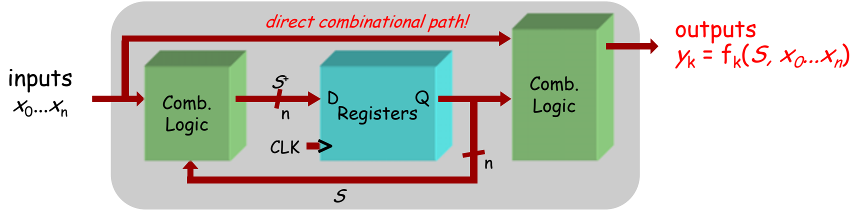
Two Types of FSMs

Moore and Mealy FSMs : different output generation

- Moore FSM:

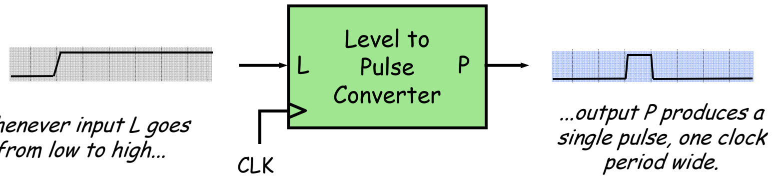


- Mealy FSM:



Design Example: Level-to-Pulse

- A level-to-pulse converter produces a single-cycle pulse each time its input goes high.
- It's a synchronous rising-edge detector.
- Sample uses:
  - Buttons and switches pressed by humans for arbitrary periods of time
  - Single-cycle enable signals for counters

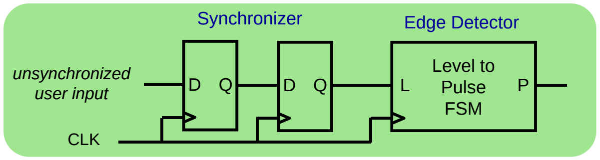


Whenever input L goes from low to high...

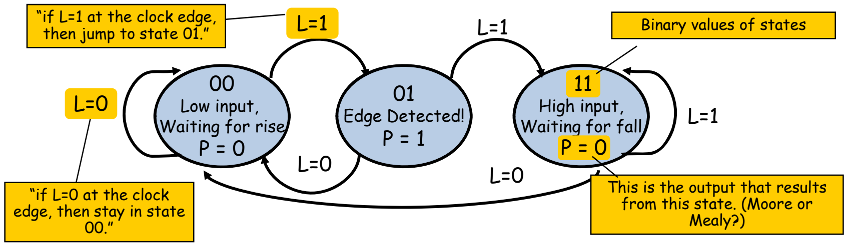
...output P produces a single pulse, one clock period wide.

Step 1: State Transition Diagram

- Block diagram of desired system:



- State transition diagram is a useful FSM representation and design aid:



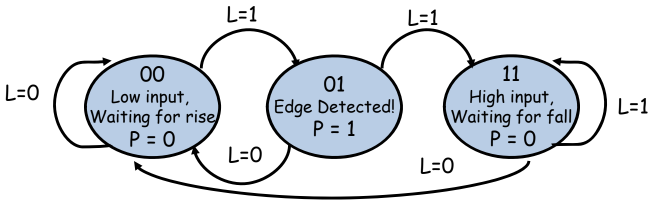
"if L=1 at the clock edge, then jump to state 01."

"if L=0 at the clock edge, then stay in state 00."

Binary values of states

This is the output that results from this state. (Moore or Mealy?)

Valid State Transition Diagrams



- Arcs leaving a state are mutually exclusive, i.e., for any combination input values there's at most one applicable arc
- Arcs leaving a state are collectively exhaustive, i.e., for any combination of input values there's at least one applicable arc
- So for each state: for any combination of input values there's <u>exactly one</u> applicable arc
- Often a starting state is specified
- Each state specifies values for all outputs (Moore)

Choosing State Representation

Choice #1: binary encoding

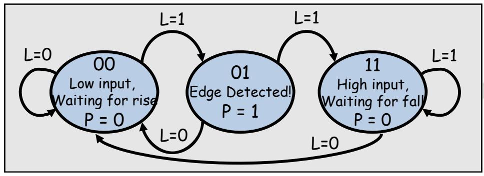
For N states, use ceil($\log_2 N$) bits to encode the state with each state represented by a unique combination of the bits. Tradeoffs: most efficient use of state registers, but requires more complicated combinational logic to detect when in a particular state.

Choice #2: "one-hot" encoding

For N states, use N bits to encode the state where the bit corresponding to the current state is 1, all the others 0. Tradeoffs: more state registers, but often much less combinational logic since state decoding is trivial.

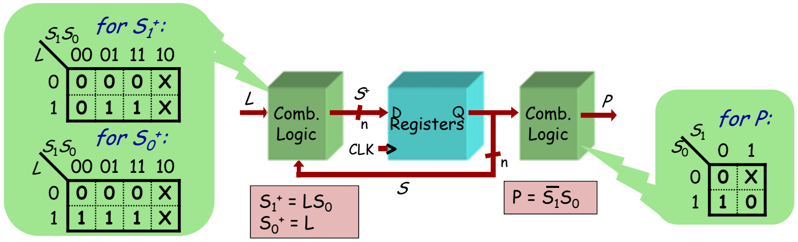
Step 2: Logic Derivation

Transition diagram is readily converted to a state transition table (just a truth table)



Current State		In	Next State		Out
S_1	S_0	L	S_1^+	S_0^+	P
0	0	0	0	0	0
0	0	1	0	1	0
0	1	0	0	0	1
0	1	1	1	1	1
1	1	0	0	0	0
1	1	1	1	1	0

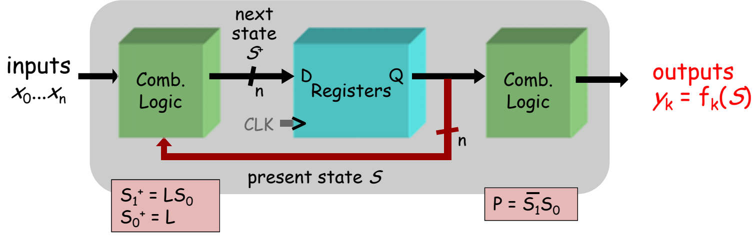
- Combinational logic may be derived using Karnaugh maps



$S_1^+ = LS_0$
$S_0^+ = L$

$P = \overline{S_1}S_0$

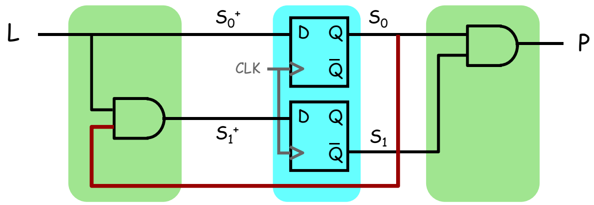
Moore Level-to-Pulse Converter



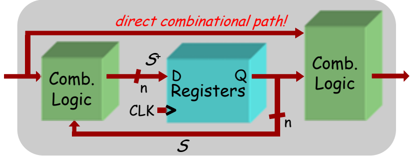
$S_1^+ = LS_0$
$S_0^+ = L$

$P = \overline{S_1}S_0$

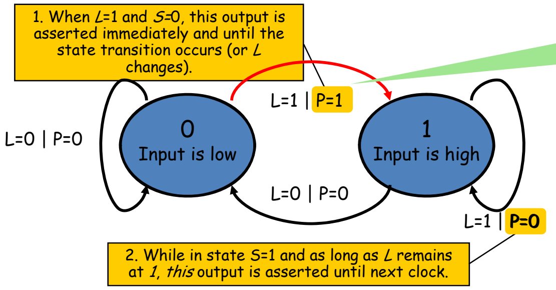
Moore FSM circuit implementation of level-to-pulse converter:



Design of a Mealy Level-to-Pulse

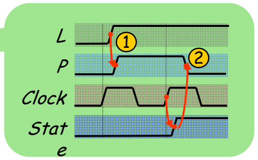


- Since outputs are determined by state *and* inputs, Mealy FSMs may need fewer states than Moore FSM implementations



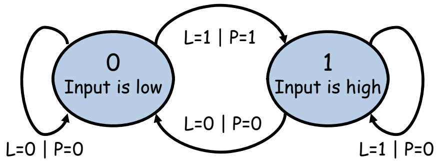
1. When *L*=1 and *S*=0, this output is asserted immediately and until the state transition occurs (or *L* changes).

2. While in state S=1 and as long as *L* remains at 1, *this* output is asserted until next clock.



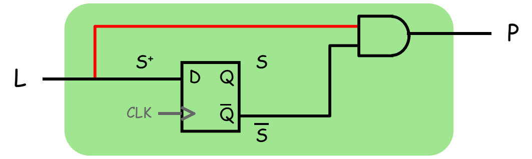
Output transitions immediately. State transitions at the clock edge.

Mealy Level-to-Pulse Converter



Pres. State	In	Next State	Out
S	L	S^+	P
0	0	0	0
0	1	1	1
1	1	1	0
1	0	0	0

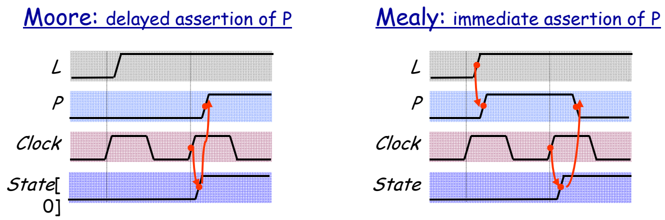
Mealy FSM circuit implementation of level-to-pulse converter:



- FSM's state simply remembers the previous value of L
- Circuit benefits from the Mealy FSM's implicit single-cycle assertion of outputs during state transitions

Moore/Mealy Trade-Offs

- How are they different?
  - Moore: outputs = f(state) only
  - Mealy outputs = f(state *and input*)
  - Mealy outputs generally occur <u>one cycle earlier</u> than a Moore:


Moore: delayed assertion of P
Mealy: immediate assertion of P

- Compared to a Moore FSM, a Mealy FSM might...
  - Be more difficult to conceptualize and design
  - Have fewer states

Example: Intersection Traffic Lights

- Design a controller for the traffic lights at the intersection of two streets – two sets of traffic lights, one for each of the streets.
- Step 1: Draw starting state transition diagram. Just handle the usual green-yellow-red cycle for both streets. How many states? Well, how many different combinations of the two sets of lights are needed?
- Step 2: add support for a walk button and walk lights to your state transition diagram.
- Step 3: add support for a traffic sensor for each of the streets – when the sensor detects traffic the green cycle for that street is extended.

Example to be worked collaboratively on the board...

FSM Example

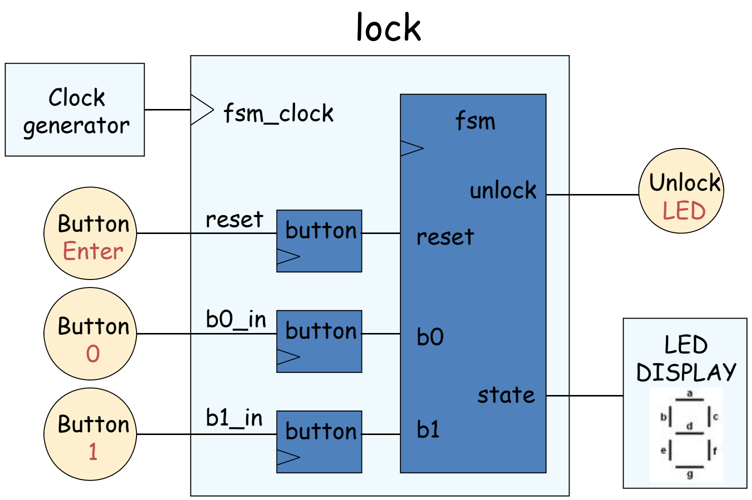
GOAL:

Build an electronic combination lock with a reset button, two number buttons (0 and 1), and an unlock output. The combination should be 01011.

STEPS:

1. Design lock FSM (block diagram, state transitions)
2. Write Verilog module(s) for FSM

Step 1A: Block Diagram


lock

Clock generator

fsm_clock

fsm

unlock

Unlock LED

Button Enter

reset

button

reset

Button 0

b0_in

button

b0

Button 1

b1_in

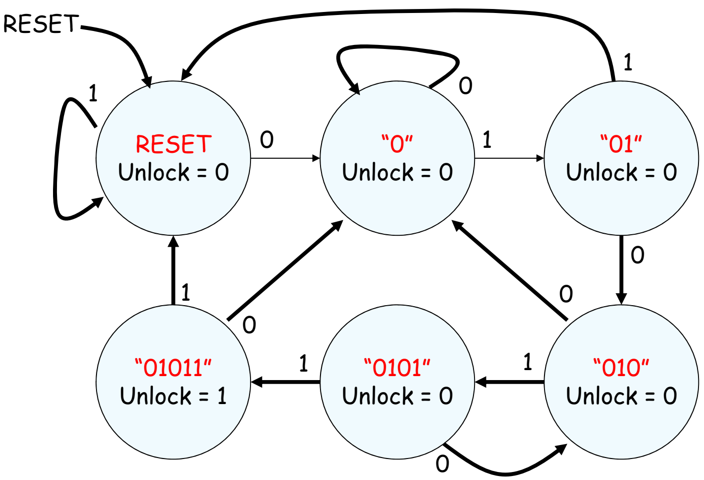
button

b1

state

LED DISPLAY

Step 1B: State transition diagram

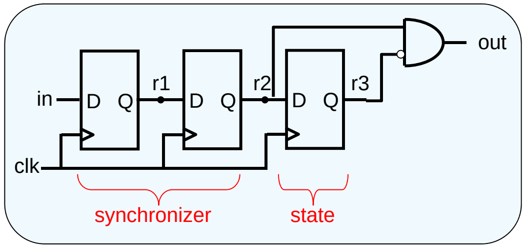


6 states → 3 bits

Step 2: Write Verilog

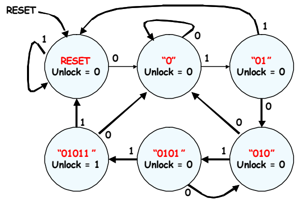
```
module lock(input clk,reset_in,b0_in,b1_in,
            output out);

   // synchronize push buttons, convert to pulses

   // implement state transition diagram
   reg [2:0] state,next_state;
   always @(*) begin
     // combinational logic!
     next_state = ???;
   end
   always @(posedge clk) state <= next_state;

   // generate output
   assign out = ???;

   // debugging?
endmodule
```

Step 2A: Synchronize buttons



```
// button
// push button synchronizer and level-to-pulse converter
// OUT goes high for one cycle of CLK whenever IN makes a
// low-to-high transition.

module button(
  input clk,in,
  output out
);
  reg r1,r2,r3;
  always @(posedge clk)
  begin
    r1 <= in;   // first reg in synchronizer
    r2 <= r1;   // second reg in synchronizer, output is in sync!
    r3 <= r2;   // remembers previous state of button
  end

  // rising edge = old value is 0, new value is 1
  assign out = ~r3 & r2;
endmodule
```

Step 2B: state transition diagram



```
parameter S_RESET = 0; // state assignments
parameter S_0 = 1;
parameter S_01 = 2;
parameter S_010 = 3;
parameter S_0101 = 4;
parameter S_01011 = 5;

reg [2:0] state, next_state;
always @(*) begin
  // implement state transition diagram
  if (reset) next_state = S_RESET;
  else case (state)
    S_RESET: next_state = b0 ? S_0   : b1 ? S_RESET : state;
    S_0:     next_state = b0 ? S_0   : b1 ? S_01    : state;
    S_01:    next_state = b0 ? S_010 : b1 ? S_RESET : state;
    S_010:   next_state = b0 ? S_0   : b1 ? S_0101  : state;
    S_0101:  next_state = b0 ? S_010 : b1 ? S_01011 : state;
    S_01011: next_state = b0 ? S_0   : b1 ? S_RESET : state;
    default: next_state = S_RESET;  // handle unused states
  endcase
end

always @(posedge clk) state <= next_state;
```

Step 2C: generate output

```
// it’s a Moore machine!  Output only depends on current state

assign out = (state == S_01011);
```

Step 2D: debugging?

```
// hmmm.  What would be useful to know?  Current state?
// hex_display on labkit shows 16 four bit values

assign hex_display = {60’b0, 1'b0, state[2:0]};
```

Step 2: final Verilog implementation

```
module lock(input clk,reset_in,b0_in,b1_in,
            output out, output [3:0] hex_display);

  wire reset, b0, b1; // synchronize push buttons, convert to pulses
  button b_reset(clk,reset_in,reset);
  button b_0(clk,b0_in,b0);
  button b_1(clk,b1_in,b1);

  parameter S_RESET = 0; parameter S_0 = 1; // state assignments
  parameter S_01 = 2;  parameter S_010 = 3;
  parameter S_0101 = 4;  parameter S_01011 = 5;

  reg [2:0] state,next_state;
  always @(*) begin                      // implement state transition diagram
    if (reset) next_state = S_RESET;
    else case (state)
      S_RESET: next_state = b0 ? S_0   : b1 ? S_RESET : state;
      S_0:     next_state = b0 ? S_0   : b1 ? S_01    : state;
      S_01:    next_state = b0 ? S_010 : b1 ? S_RESET : state;
      S_010:   next_state = b0 ? S_0   : b1 ? S_0101  : state;
      S_0101:  next_state = b0 ? S_010 : b1 ? S_01011 : state;
      S_01011: next_state = b0 ? S_0   : b1 ? S_RESET : state;
      default: next_state = S_RESET;     // handle unused states
    endcase
  end
  always @ (posedge clk) state <= next_state;

  assign out = (state == S_01011);   // assign output: Moore machine
  assign hex_display = {1'b0,state};      // debugging
endmodule
```

Real FSM Security System

The 6.111 Vending Machine

- Lab assistants demand a new soda machine for the 6.111 lab. You design the FSM controller.
- All selections are $0.30.
- The machine makes change. (Dimes and nickels only.)
- Inputs: limit 1 per clock
  - Q - quarter inserted
  - D - dime inserted
  - N - nickel inserted
- Outputs: limit 1 per clock
  - DC - dispense can
  - DD - dispense dime
  - DN - dispense nickel

What States are in the System?

- A starting (idle) state:

  idle

- A state for each possible amount of money captured:

  got5c **got10c** **got15c** ...

- What's the maximum amount of money captured before purchase?

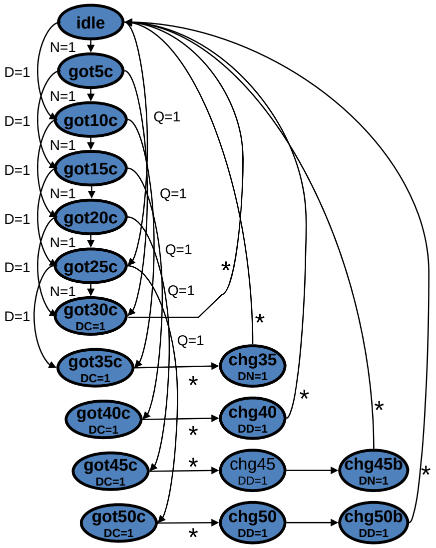
  25 cents (just shy of a purchase) + one quarter (largest coin)

  ... **got35c** **got40c** **got45c** **got50c**

- States to dispense change (one per coin dispensed):

A Moore Vender

Here's a first cut at the state transition diagram.

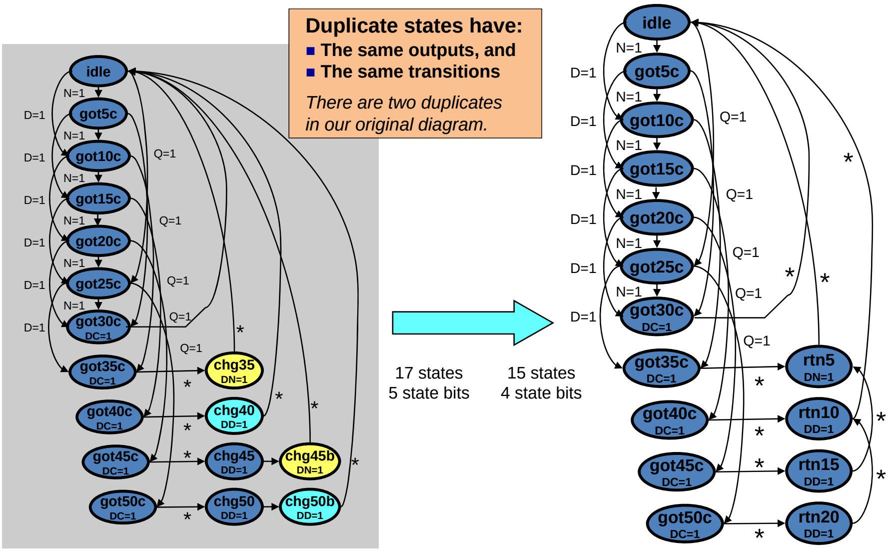


State Reduction

Duplicate states have:

- **The same outputs, and**
- **The same transitions**

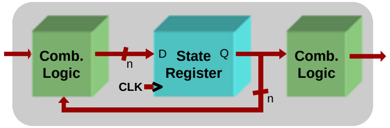
There are two duplicates in our original diagram.



17 states
5 state bits

15 states
4 state bits

Verilog for the Moore Vender



FSMs are easy in Verilog. Simply write one of each:

- **State register (sequential always block)**
- **Next-state combinational logic (comb. always block with case)**
- **Output combinational logic block (comb. always block *or* assign statements)**

```
module mooreVender (
  input N, D, Q, clk, reset,
  output DC, DN, DD,
  output reg [3:0] state);

  reg  next;
```

States defined with parameter keyword

```
  parameter IDLE = 0;
  parameter GOT_5c = 1;
  parameter GOT_10c = 2;
  parameter GOT_15c = 3;
  parameter GOT_20c = 4;
  parameter GOT_25c = 5;
  parameter GOT_30c = 6;
  parameter GOT_35c = 7;
  parameter GOT_40c = 8;
  parameter GOT_45c = 9;
  parameter GOT_50c = 10;
  parameter RETURN_20c = 11;
  parameter RETURN_15c = 12;
  parameter RETURN_10c = 13;
  parameter RETURN_5c = 14;
```

State register defined with sequential always block

```
  always @ (posedge clk or negedge reset)
    if (!reset)    state <= IDLE;
    else           state <= next;
```

Verilog for the Moore Vender

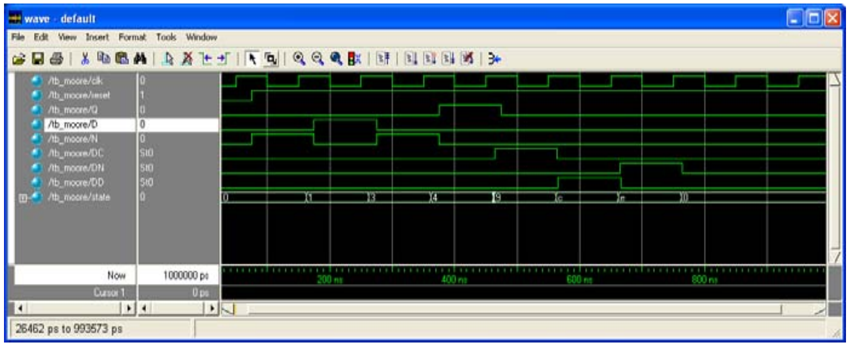
Next-state logic within a combinational always block

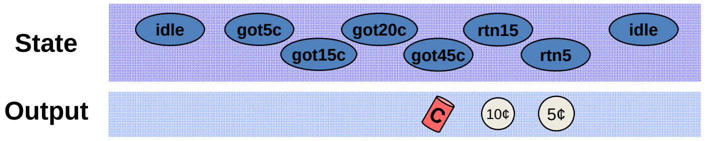
```
always @ (state or N or D or Q) begin

  case (state)
    IDLE:       if (Q) next = GOT_25c;
          else if (D) next = GOT_10c;
            else if (N) next = GOT_5c;
            else next = IDLE;

    GOT_5c:     if (Q) next = GOT_30c;
          else if (D) next = GOT_15c;
            else if (N) next = GOT_10c;
            else next = GOT_5c;

    GOT_10c:    if (Q) next = GOT_35c;
          else if (D) next = GOT_20c;
            else if (N) next = GOT_15c;
            else next = GOT_10c;

    GOT_15c:    if (Q) next = GOT_40c;
          else if (D) next = GOT_25c;
            else if (N) next = GOT_20c;
            else next = GOT_15c;

    GOT_20c:    if (Q) next = GOT_45c;
          else if (D) next = GOT_30c;
            else if (N) next = GOT_25c;
            else next = GOT_20c;

    GOT_25c:    if (Q) next = GOT_50c;
          else if (D) next = GOT_35c;
            else if (N) next = GOT_30c;
            else next = GOT_25c;

    GOT_30c:  next = IDLE;
    GOT_35c:  next = RETURN_5c;
    GOT_40c:  next = RETURN_10c;
    GOT_45c:  next = RETURN_15c;
    GOT_50c:  next = RETURN_20c;

    RETURN_20c:  next = RETURN_10c;
    RETURN_15c:  next = RETURN_5c;
    RETURN_10c:  next = IDLE;
    RETURN_5c:   next = IDLE;

    default: next = IDLE;
  endcase
end
```

Combinational output assignment

```
 assign DC = (state == GOT_30c || state == GOT_35c ||
              state == GOT_40c || state == GOT_45c ||
              state == GOT_50c);
 assign DN = (state == RETURN_5c);
 assign DD = (state == RETURN_20c || state == RETURN_15c ||
              state == RETURN_10c);
endmodule
```

Simulation of Moore Vender





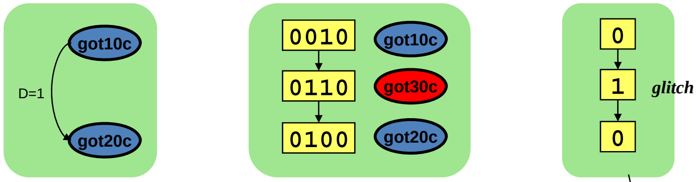
FSM Output Glitching

- FSM state bits may not transition at precisely the same time
- Combinational logic for outputs may contain hazards
- Result: your FSM outputs may glitch!

during this state transition...

...the state registers may transtion like this...

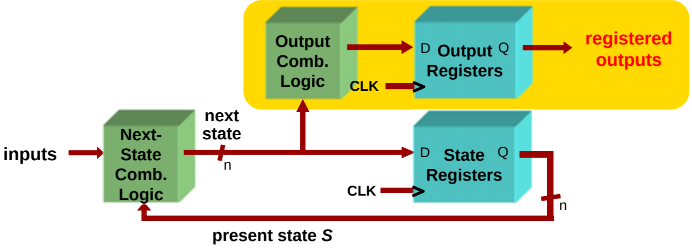
...causing the DC output to ***glitch*** *like this!*



```
assign DC = (state == GOT_30c || state == GOT_35c ||
             state == GOT_40c || state == GOT_45c ||
             state == GOT_50c);
```

If the soda dispenser is glitch-sensitive, your customers can get a 20-cent soda!

Registered FSM Outputs are Glitch-Free

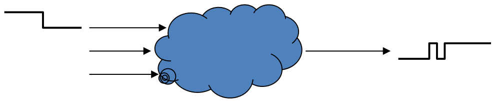


- Move output generation into the sequential always block
- Calculate outputs based on <u>next</u> state
- Delays outputs by one clock cycle. Problematic in some application.

```
reg DC,DN,DD;

 // Sequential always block for state assignment
 always @ (posedge clk or negedge reset) begin
   if (!reset)   state <= IDLE;
   else if (clk) state <= next;

   DC <= (next == GOT_30c || next == GOT_35c ||
          next == GOT_40c || next == GOT_45c ||
          next == GOT_50c);
   DN <= (next == RETURN_5c);
   DD <= (next == RETURN_20c || next == RETURN_15c ||
          next == RETURN_10c);
 end
```

Where should CLK come from?

- Option 1: external crystal
  - Stable, known frequency, typically 50% duty cycle
- Option 2: internal signals
  - Option 2A: output of combinational logic



    - No! If inputs to logic change, output may make several transitions before settling to final value → several rising edges, not just one! Hard to design away output glitches…
  - Option 2B: output of a register
    - Okay, but timing of CLK2 won't line up with CLK1