



L3: Introduction to Verilog (Combinational Logic)



Acknowledgements : Rex Min

Verilog References:

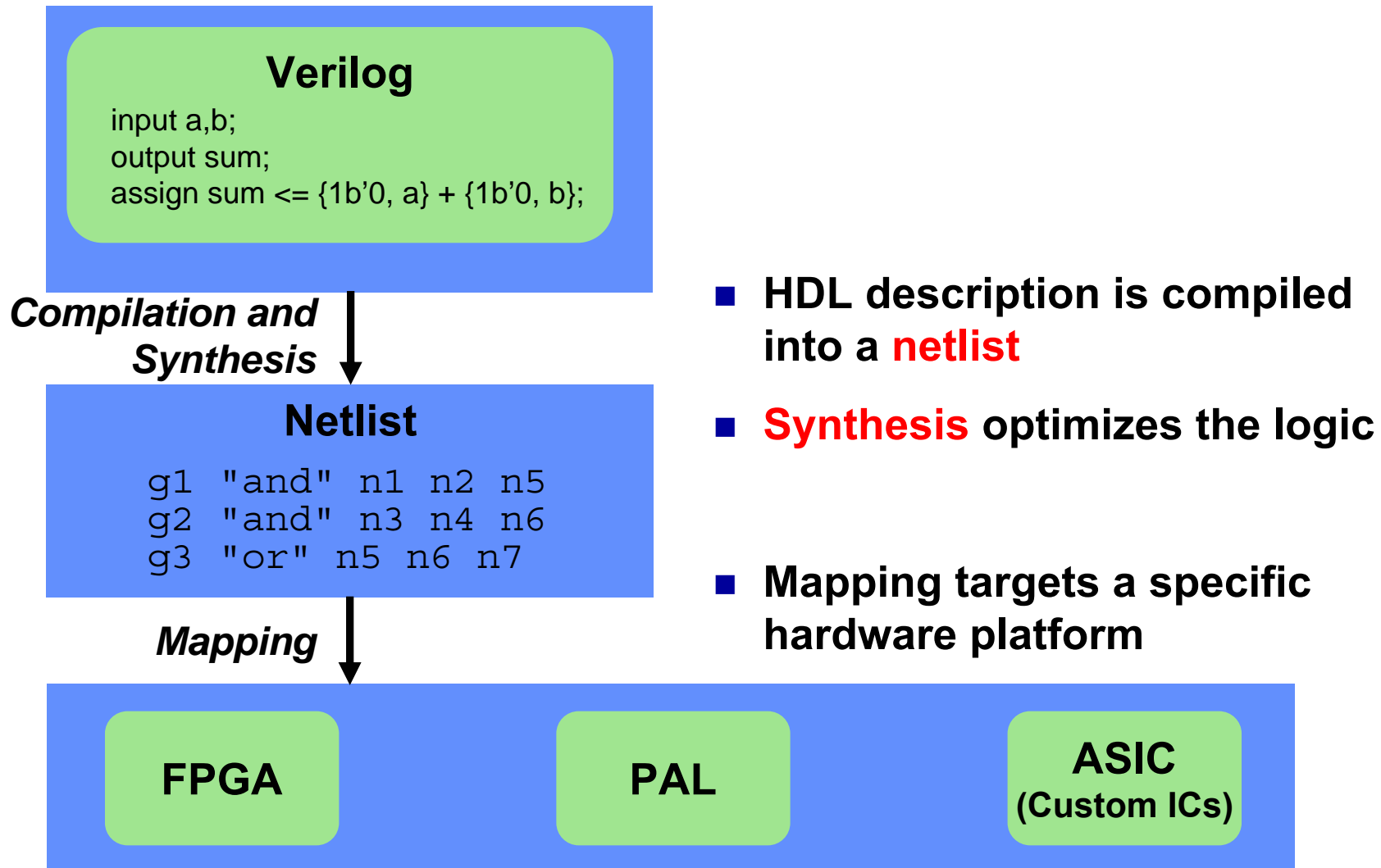
- **Samir Palnitkar, Verilog HDL, Pearson Education (2nd edition).**
- **Donald Thomas, Philip Moorby, The Verilog Hardware Description Language, Fifth Edition, Kluwer Academic Publishers.**
- **J. Bhasker, Verilog HDL Synthesis (A Practical Primer), Star Galaxy Publishing**



Synthesis and HDLs



- Hardware description language (HDL) is a convenient, device-independent representation of digital logic

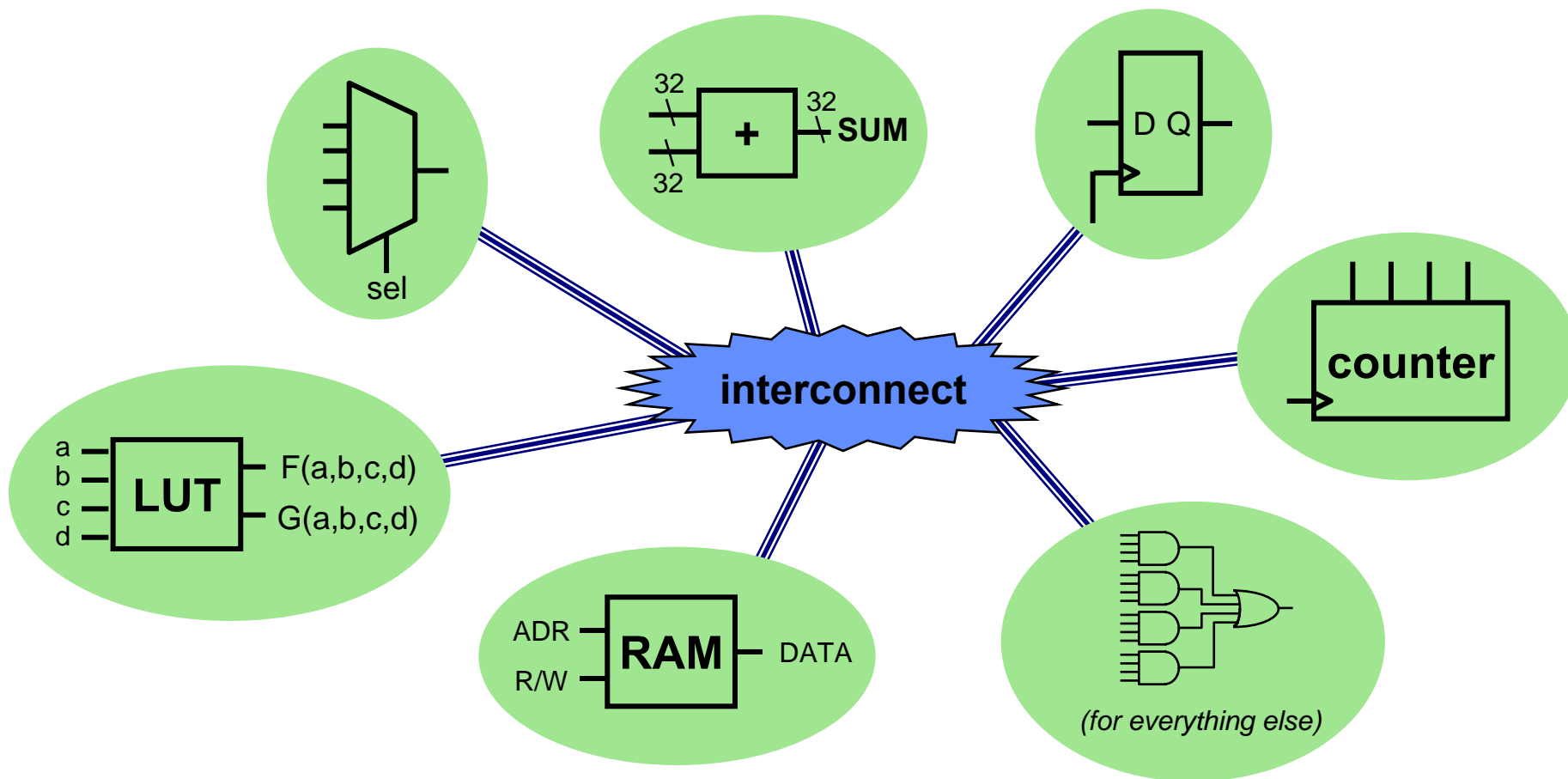




The FPGA: A Conceptual View



- An FPGA is like an electronic breadboard that is wired together by an automated **synthesis tool**
- Built-in components are called **macros**





Synthesis and Mapping for FPGAs



- **Infer macros:** choose the FPGA macros that efficiently implement various parts of the HDL code

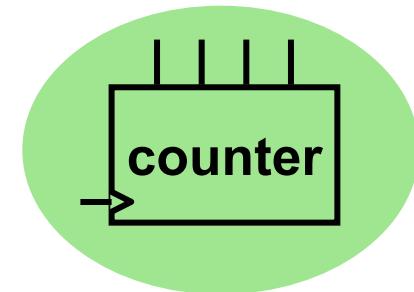
```

...
always @ (posedge clk)
begin
    count <= count + 1;
end
...

```

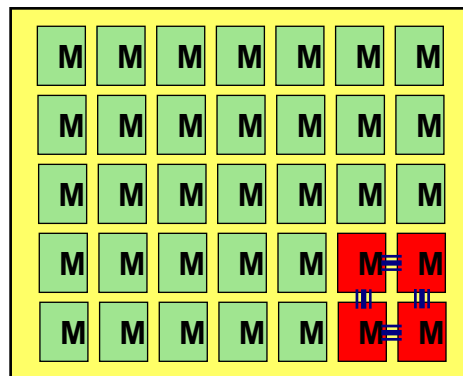
HDL Code

“This section of code looks like a counter. My FPGA has some of those...”



Inferred Macro

- **Place-and-route:** with area and/or speed in mind, choose the needed macros by location and route the interconnect



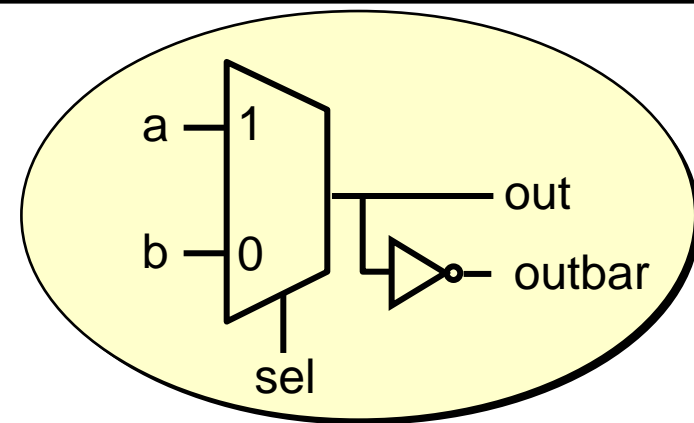
“This design only uses 10% of the FPGA. Let’s use the macros in one corner to minimize the distance between blocks.”



Verilog: The Module



- Verilog designs consist of interconnected **modules**.
- A module can be an element or collection of lower level design blocks.
- A simple module with combinational logic might look like this:



$$\text{Out} = \text{sel} \bullet a + \overline{\text{sel}} \bullet b$$

2-to-1 multiplexer with inverted output

```
module mux_2_to_1(a, b, out,
                 outbar, sel);
```

```
// This is 2:1 multiplexor
```

```
input a, b, sel;
output out, outbar;
```

```
assign out = sel ? a : b;
assign outbar = ~out;
```

```
endmodule
```

Declare and name a module; list its ports. Don't forget that semicolon.

Comment starts with //
Verilog skips from // to end of the line

Specify each port as input, output, or inout

Express the module's behavior. Each statement executes in parallel; order does not matter.

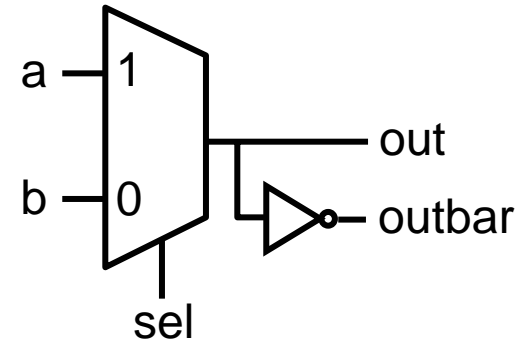
Conclude the module code.



Continuous (Dataflow) Assignment



```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
  
    input a, b, sel;  
    output out, outbar;  
  
    assign out = sel ? a : b;  
    assign outbar = ~out;  
  
endmodule
```



- Continuous assignments use the `assign` keyword
- A simple and natural way to represent combinational logic
- Conceptually, the right-hand expression is continuously evaluated as a function of arbitrarily-changing inputs...just like dataflow
- The target of a continuous assignment is a net driven by combinational logic
- Left side of the assignment must be a scalar or vector net or a concatenation of scalar and vector nets. It can't be a scalar or vector register (*discussed later*). Right side can be register or nets
- Dataflow operators are fairly low-level:
 - Conditional assignment: `(conditional_expression) ? (value-if-true) : (value-if-false);`
 - Boolean logic: `~, &, |`
 - Arithmetic: `+, -, *`
- Nested conditional operator (4:1 mux)
 - `assign out = s1 ? (s0 ? i3 : i2) : (s0 ? i1 : i0);`

MAX+plusII: Simulator, Synthesis, Mapping

- Must be synthesizable Verilog files
- Step by step instructions on the course WEB site

Create *.v file (module name same as file name)

The screenshot displays the MAX+plus II software interface. The main window is titled "mux_2_to_1.v - Text Editor" and contains the following Verilog code:

```
// this is a 2:1 verilog description of a multiplexor  
  
module mux_2_to_1(a, b, out, outbar, sel);  
  
    input a, b, sel;  
    output out, outbar;  
    assign out = sel ? a : b;  
    assign outbar = ~out;  
  
endmodule
```

To the right of the code editor is a "Compiler" window with buttons for "Compiler Netlist Extractor", "Database Builder", "Logic Synthesizer", "Partitioner", and "Fitter". Below these buttons is a progress bar and a "Start" button.

Below the code editor is a "mux_2_to_1.scf - Waveform Editor" window. It shows a timing diagram with signals: sel, b, a, out, and outbar. The time axis ranges from 0 to 800 ns. A vertical line is drawn at 600 ns. Two areas are circled: one on the sel signal at approximately 100 ns and another on the outbar signal at 600 ns, which shows a sharp spike labeled as a "Glitch".

Select area and set inputs through overwrite or insert menu (under edit)

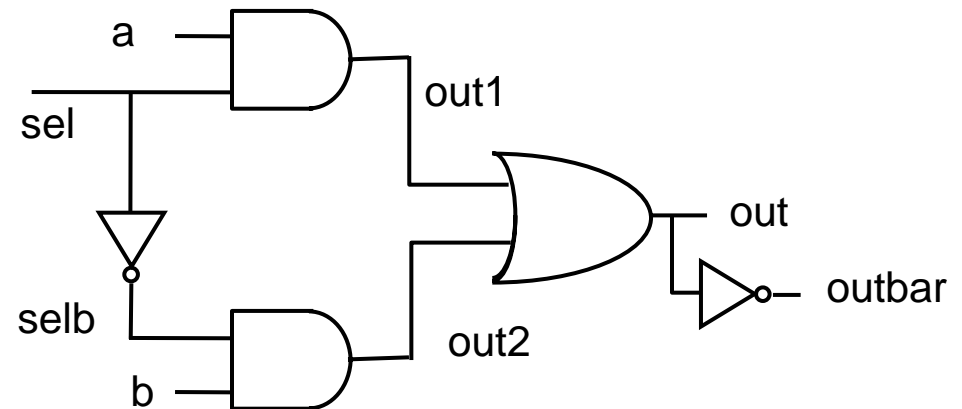
Glitch



Gate Level Description



```
module muxgate (a, b, out,  
outbar, sel);  
input a, b, sel;  
output out, outbar;  
wire out1, out2, selb;  
and a1 (out1, a, sel);  
not i1 (selb, sel);  
and a2 (out2, b, selb);  
or o1 (out, out1, out2);  
assign outbar = ~out;  
endmodule
```



- Verilog supports basic logic gates as primitives

- and, nand, or, nor, xor, xnor, not, buf

- can be extended to multiple inputs: e.g., nand nand3in (out, in1, in2,in3);

- buif1 and buif0 are tri-state buffers

- Net represents connections between hardware elements. Nets are declared with the keyword `wire`.



Procedural Assignment with `always`



- Procedural assignment allows an alternative, often higher-level, behavioral description of combinational logic
- Two structured procedure statements: `initial` and `always`
- Supports richer, C-like control structures such as `if`, `for`, `while`, `case`

```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
  input a, b, sel;  
  output out, outbar;
```

Exactly the same as before.

```
  reg out, outbar;
```

Anything assigned in an `always` block must *also* be declared as type `reg` (next slide)

```
  always @ (a or b or sel)
```

Conceptually, the `always` block runs *once* whenever a signal in the **sensitivity list** changes value

```
  begin
```

```
    if (sel) out = a;  
    else out = b;  
  
    outbar = ~out;
```

Statements within the `always` block are executed sequentially. Order matters!

```
  end
```

Surround multiple statements in a single `always` block with `begin/end`.

```
endmodule
```



Verilog Registers



- In digital design, registers represent memory elements (we will study these in the next few lectures)
- Digital registers need a clock to operate and update their state on certain phase or edge
- Registers in Verilog should not be confused with hardware registers
- **In Verilog, the term register (`reg`) simply means a variable that can hold a value**
- Verilog registers don't need a clock and don't need to be driven like a net. Values of registers can be changed anytime in a simulation by assuming a new value to the register



Mix-and-Match Assignments



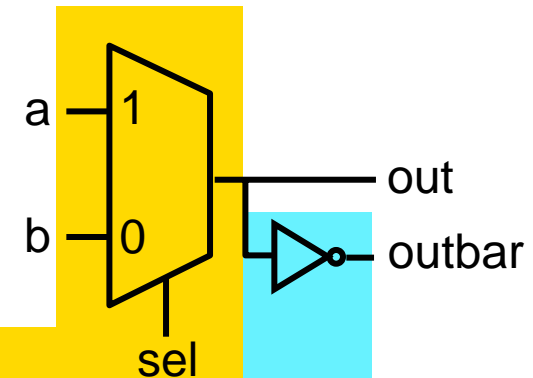
- Procedural and continuous assignments can (and often do) co-exist within a module
- Procedural assignments update the value of `reg`. The value will remain unchanged till another procedural assignment updates the variable. This is the main difference with continuous assignments in which the right hand expression is constantly placed on the left-side

```
module mux_2_to_1(a, b, out,  
                  outbar, sel);  
  input a, b, sel;  
  output out, outbar;  
  reg out;
```

```
  always @ (a or b or sel)  
  begin  
    if (sel) out = a;  
    else out = b;  
  end
```

```
  assign outbar = ~out;
```

```
endmodule
```



*procedural
description*

*continuous
description*



The case Statement



- case and if may be used interchangeably to implement conditional execution within always blocks
- case is easier to read than a long string of if...else statements

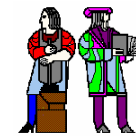
```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
    input a, b, sel;  
    output out, outbar;  
    reg out;  
  
    always @ (a or b or sel)  
    begin  
        if (sel) out = a;  
        else out = b;  
    end  
  
    assign outbar = ~out;  
  
endmodule
```

```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
    input a, b, sel;  
    output out, outbar;  
    reg out;  
  
    always @ (a or b or sel)  
    begin  
        case (sel)  
            1'b1: out = a;  
            1'b0: out = b;  
        endcase  
    end  
  
    assign outbar = ~out;  
  
endmodule
```

Note: Number specification notation: `<size>'<base><number>`
(`4'b1010` is a 4-bit binary value, `16'h6cda` is a 16 bit hex number, and `8'd40` is an 8-bit decimal value)

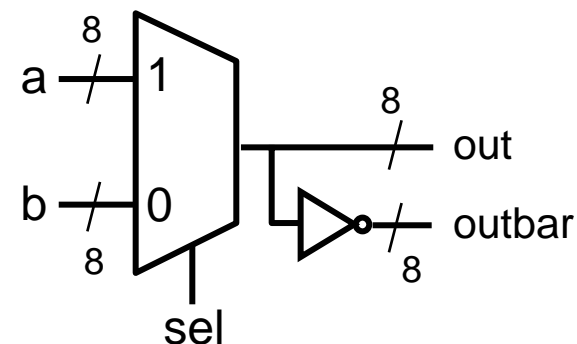


The Power of Verilog: *n*-bit Signals



- Multi-bit signals and buses are *easy* in Verilog.
- 2-to-1 multiplexer with *8-bit operands*:

```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
    input [7:0] a, b;  
    input sel;  
    output [7:0] out, outbar;  
    reg [7:0] out;  
    always @ (a or b or sel)  
    begin  
        if (sel) out = a;  
        else out = b;  
    end  
    assign outbar = ~out;  
endmodule
```



Concatenate signals using the **{ }** operator

```
assign {b[7:0], b[15:8]} = {a[15:8], a[7:0]};  
effects a byte swap
```



The Power of Verilog: Integer Arithmetic



- Verilog's built-in arithmetic makes a 32-bit adder easy:

```
module add32(a, b, sum);  
    input [31:0] a,b;  
    output [31:0] sum;  
    assign sum = a + b;  
endmodule
```

- A 32-bit adder with carry-in and carry-out:

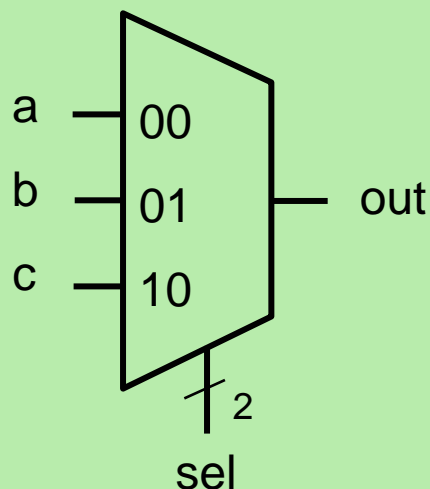
```
module add32_carry(a, b, cin, sum, cout);  
    input [31:0] a,b;  
    input cin;  
    output [31:0] sum;  
    output cout;  
    assign {cout, sum} = a + b + cin;  
endmodule
```



Dangers of Verilog: Incomplete Specification



Goal:



3-to-1 MUX
(‘11’ input is a don’t-care)

Proposed Verilog Code:

```
module maybe_mux_3to1(a, b, c,
                      sel, out);

    input [1:0] sel;
    input a,b,c;
    output out;
    reg out;

    always @(a or b or c or sel)
    begin
        case (sel)
            2'b00: out = a;
            2'b01: out = b;
            2'b10: out = c;
        endcase
    end
endmodule
```

Is this a 3-to-1 multiplexer?



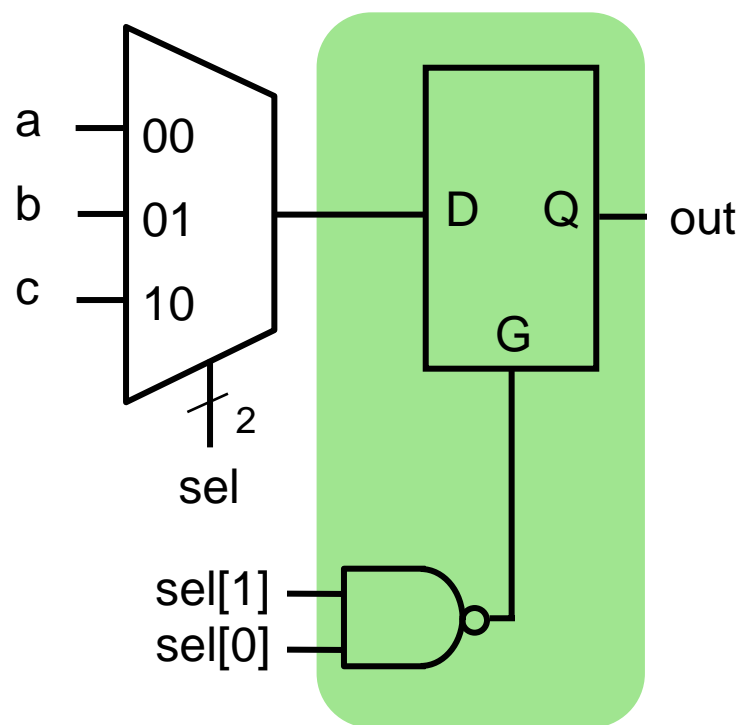
Incomplete Specification Infers Latches



```
module maybe_mux_3to1(a, b, c,  
                      sel, out);  
  
  input [1:0] sel;  
  input a,b,c;  
  output out;  
  reg out;  
  
  always @(a or b or c or sel)  
  begin  
    case (sel)  
      2'b00: out = a;  
      2'b01: out = b;  
      2'b10: out = c;  
    endcase  
  end  
endmodule
```

if out is not assigned during any pass through the always block, then **the previous value must be retained!**

Synthesized Result:



- Latch memory “latches” old data when $G=0$ (we will discuss latches later)
- In practice, we almost *never* intend this



Avoiding Incomplete Specification



- Precede all conditionals with a default assignment for all signals assigned within them...

```
always @(a or b or c or sel)
begin
    out = 1'bx;
    case (sel)
        2'b00: out = a;
        2'b01: out = b;
        2'b10: out = c;
    endcase
end
endmodule
```

```
always @(a or b or c or sel)
begin
    case (sel)
        2'b00: out = a;
        2'b01: out = b;
        2'b10: out = c;
        default: out = 1'bx;
    endcase
end
endmodule
```

- ...or, fully specify all branches of conditionals and assign all signals from all branches
 - For each `if`, include `else`
 - For each `case`, include `default`

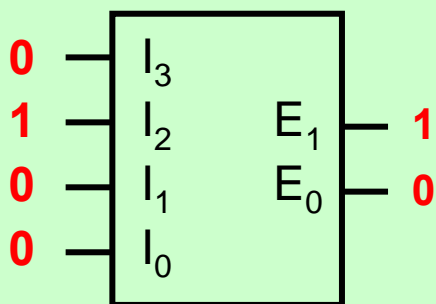


Dangers of Verilog: Priority Logic



Goal:

4-to-2 Binary Encoder



I_3	I_2	I_1	I_0	E_1	E_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1
all others				X	X

Proposed Verilog Code:

```

module binary_encoder(i, e);
  input [3:0] i;
  output [1:0] e;
  reg e;

  always @(i)
  begin
    if (i[0]) e = 2'b00;
    else if (i[1]) e = 2'b01;
    else if (i[2]) e = 2'b10;
    else if (i[3]) e = 2'b11;
    else e = 2'bxx;
  end
endmodule

```

What is the resulting circuit?



Priority Logic



Intent: if more than one input is 1, the result is a don't-care.

I_3	I_2	I_1	I_0	E_1	E_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1
all others				X	X

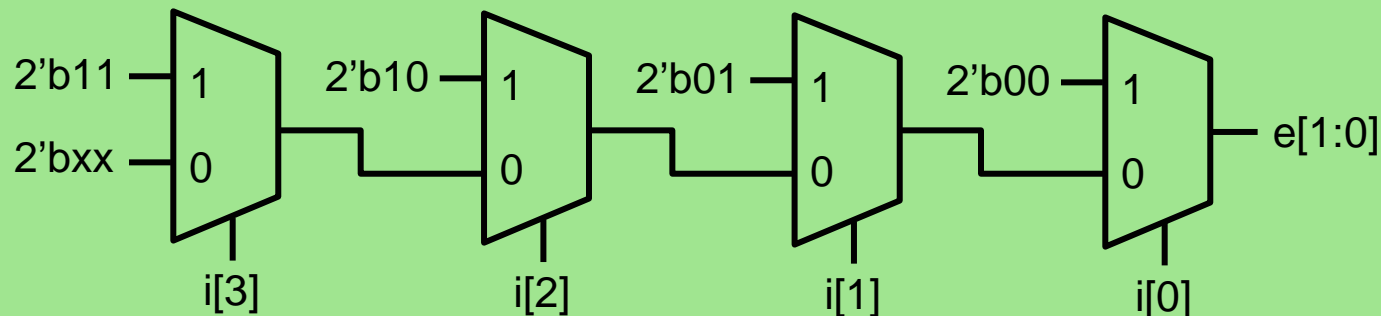
Code: if $i[0]$ is 1, the result is 00 regardless of the other inputs. $i[0]$ takes the highest priority.

```

if (i[0]) e = 2'b00;
else if (i[1]) e = 2'b01;
else if (i[2]) e = 2'b10;
else if (i[3]) e = 2'b11;
else e = 2'bxx;
end

```

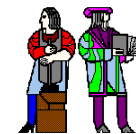
Inferred Result:



- **if-else and case statements are interpreted very literally!**
Beware of unintended **priority logic**.



Avoiding (Unintended) Priority Logic



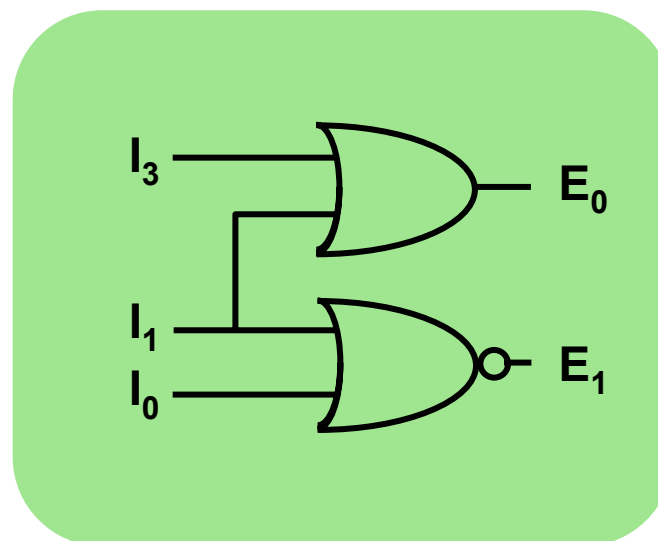
- Make sure that `if-else` and `case` statements are *parallel*
 - If **mutually exclusive conditions** are chosen for each branch...
 - ...then synthesis tool can generate a simpler circuit that evaluates the branches in parallel

Parallel Code:

```
module binary_encoder(i, e);
  input [3:0] i;
  output [1:0] e;
  reg e;

  always @(i)
  begin
    if (i == 4'b0001) e = 2'b00;
    else if (i == 4'b0010) e = 2'b01;
    else if (i == 4'b0100) e = 2'b10;
    else if (i == 4'b1000) e = 2'b11;
    else e = 2'bxx;
  end
endmodule
```

Minimized Result:



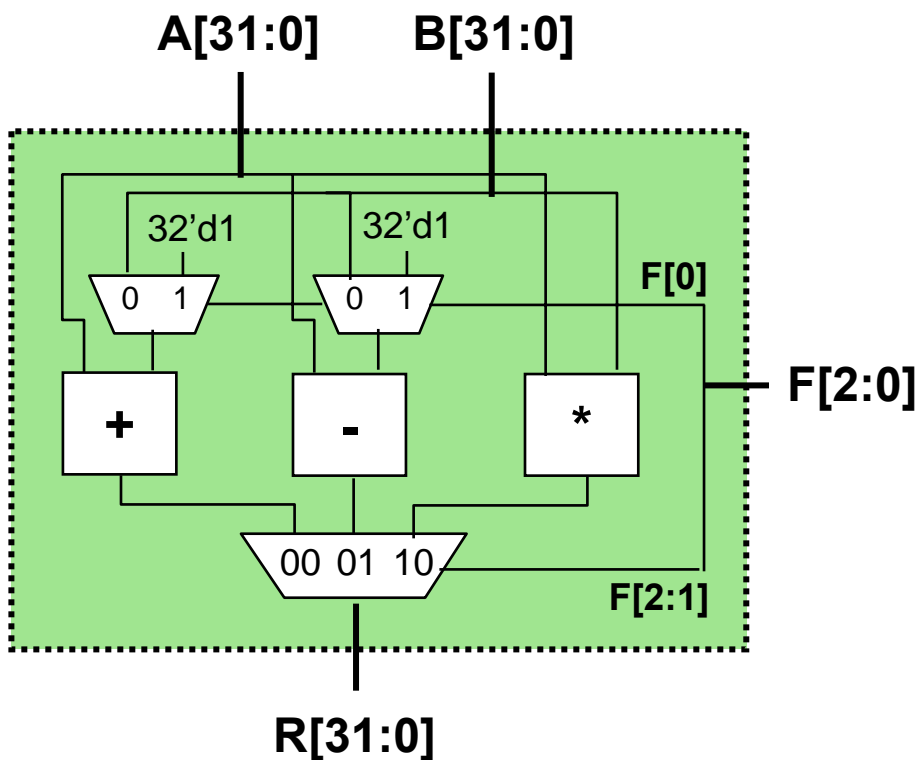


Interconnecting Modules



- Modularity is essential to the success of large designs
- A Verilog `module` may contain submodules that are “wired together”
- High-level primitives enable direct synthesis of behavioral descriptions (functions such as additions, subtractions, shifts (<< and >>), etc.

Example: A 32-bit ALU



Function Table

F2	F1	F0	Function
0	0	0	A + B
0	0	1	A + 1
0	1	0	A - B
0	1	1	A - 1
1	0	X	A * B



Module Definitions



2-to-1 MUX

```
module mux32two(i0,i1,sel,out);
input [31:0] i0,i1;
input sel;
output [31:0] out;

assign out = sel ? i1 : i0;

endmodule
```

3-to-1 MUX

```
module mux32three(i0,i1,i2,sel,out);
input [31:0] i0,i1,i2;
input [1:0] sel;
output [31:0] out;
reg [31:0] out;

always @ (i0 or i1 or i2 or sel)
begin
    case (sel)
        2'b00: out = i0;
        2'b01: out = i1;
        2'b10: out = i2;
        default: out = 32'bx;
    endcase
end
endmodule
```

32-bit Adder

```
module add32(i0,i1,sum);
input [31:0] i0,i1;
output [31:0] sum;

assign sum = i0 + i1;

endmodule
```

32-bit Subtractor

```
module sub32(i0,i1,diff);
input [31:0] i0,i1;
output [31:0] diff;

assign diff = i0 - i1;

endmodule
```

16-bit Multiplier

```
module mul16(i0,i1,prod);
input [15:0] i0,i1;
output [31:0] prod;

// this is a magnitude multiplier
// signed arithmetic later
assign prod = i0 * i1;

endmodule
```



Top-Level ALU Declaration



Given submodules:

```

module mux32two(i0,i1,sel,out);
module mux32three(i0,i1,i2,sel,out);
module add32(i0,i1,sum);
module sub32(i0,i1,diff);
module mul16(i0,i1,prod);

```

Declaration of the ALU Module:

```

module alu(a, b, f, r);
  input [31:0] a, b;
  input [2:0] f;
  output [31:0] r;

```

```

wire [31:0] addmux_out, submux_out;
wire [31:0] add_out, sub_out, mul_out;

```

```

mux32two    adder_mux(b, 32'd1, f[0], addmux_out);
mux32two    sub_mux(b, 32'd1, f[0], submux_out);
add32       our_adder(a, addmux_out, add_out);
sub32       our_subtractor(a, submux_out, sub_out);
mul16       our_multiplier(a[15:0], b[15:0], mul_out);
mux32three  output_mux(add_out, sub_out, mul_out, f[2:1], r);

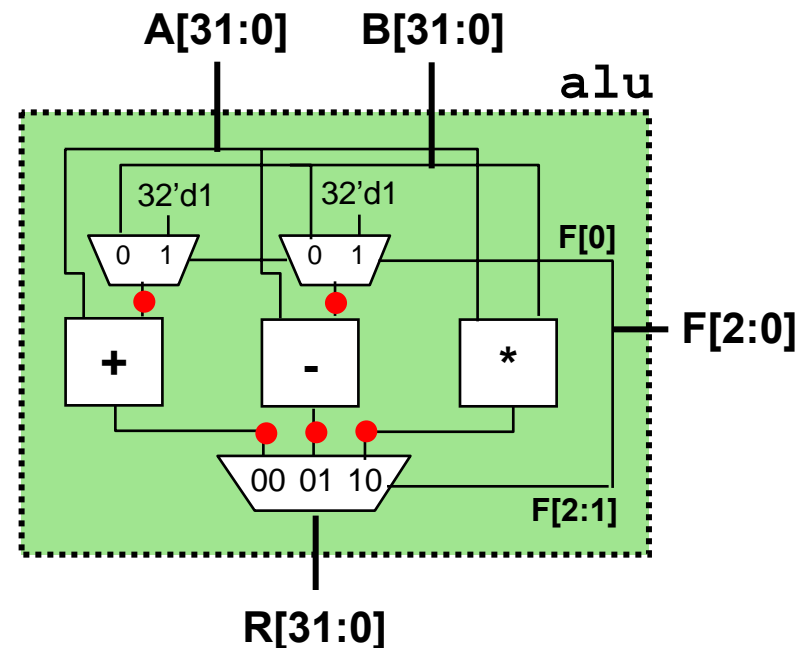
```

endmodule

module names

(unique) instance names

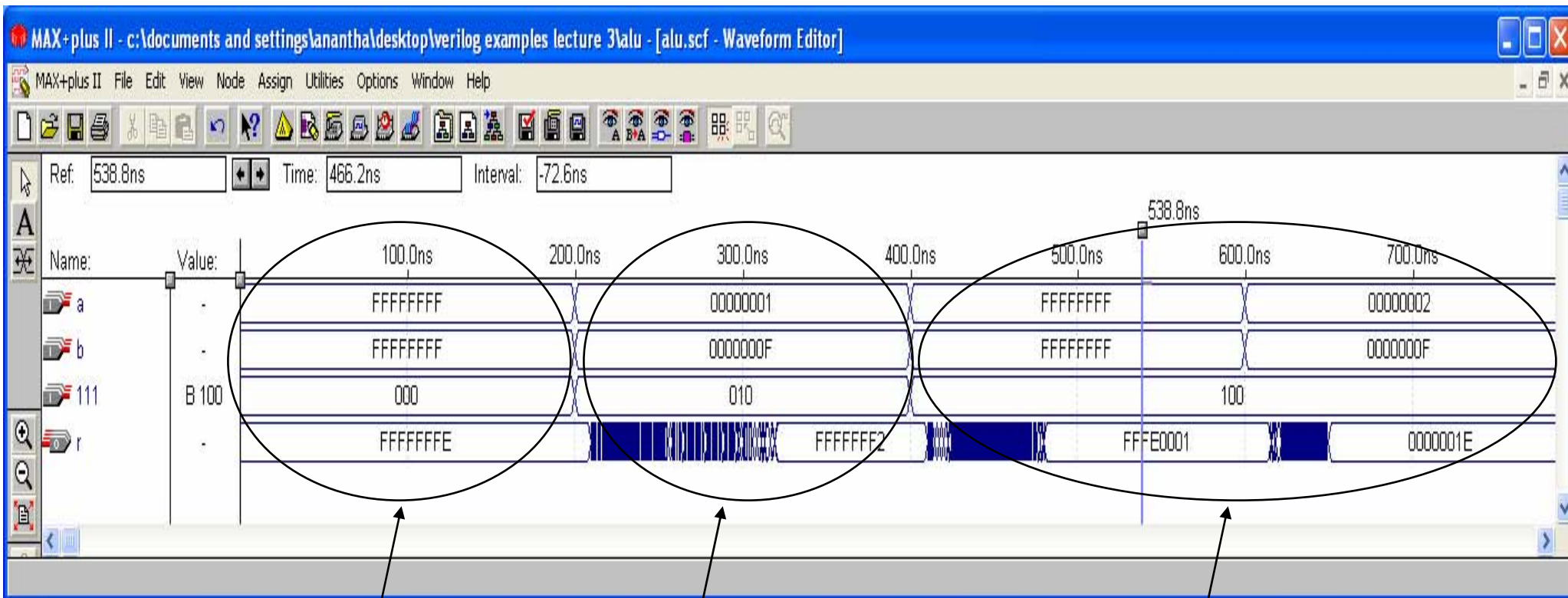
corresponding wires/regs in module alu



intermediate output nodes ●



Simulation



addition

subtraction

multiplier



More on Module Interconnection



- **Explicit port naming allows port mappings in arbitrary order: better scaling for large, evolving designs**

Given Submodule Declaration:

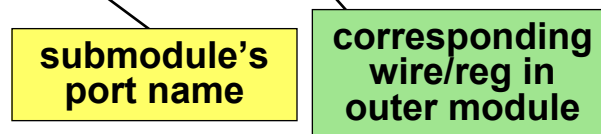
```
module mux32three(i0,i1,i2,sel,out);
```

Module Instantiation with Ordered Ports:

```
mux32three output_mux(add_out, sub_out, mul_out, f[2:1], r);
```

Module Instantiation with Named Ports:

```
mux32three output_mux(.sel(f[2:1]), .out(r), .i0(add_out),  
                      .i1(sub_out), .i2(mul_out));
```



- **Built-in Verilog gate primitives may be instantiated as well**
 - **Instantiations may omit instance name and must be ordered:**

```
buf(out1,out2,...,outN, in); and(in1,in2,...inN,out);
```



Useful Boolean Operators



- **Bitwise operators** perform bit-sliced operations on vectors
 - $\sim(4'b0101) = \{\sim 0, \sim 1, \sim 0, \sim 1\} = 4'b1010$
 - $4'b0101 \& 4'b0011 = 4'b0001$
- **Logical operators** return one-bit (true/false) results
 - $!(4'b0101) = \sim 1 = 1'b0$
- **Reduction operators** act on each bit of a single input vector
 - $\&(4'b0101) = 0 \& 1 \& 0 \& 1 = 1'b0$
- **Comparison operators** perform a Boolean test on two arguments

Bitwise

$\sim a$	NOT
$a \& b$	AND
$a b$	OR
$a \wedge b$	XOR
$a \sim \wedge b$	XNOR

Logical

$!a$	NOT
$a \&\& b$	AND
$a b$	OR

Reduction

$\&a$	AND
$\sim \&$	NAND
$ $	OR
$\sim $	NOR
\wedge	XOR

Comparison

$a < b$ $a > b$ $a \leq b$ $a \geq b$	Relational
$a == b$ $a != b$	[in]equality returns x when x or z in bits. Else returns 0 or 1
$a === b$ $a !== b$	case [in]equality returns 0 or 1 based on bit by bit comparison

Note distinction between $\sim a$ and $!a$



Testbenches (ModelSim) – Demo this week in Lab by TAs



Full Adder (1-bit)

```

module full_adder (a, b, cin,
                  sum, cout);
    input  a, b, cin;
    output sum, cout;
    reg   sum, cout;

    always @(a or b or cin)
    begin
        sum = a ^ b ^ cin;
        cout = (a & b) | (a & cin) | (b & cin);
    end
endmodule

```

Full Adder (4-bit)

```

module full_adder_4bit (a, b, cin, sum,
                       cout);
    input[3:0] a, b;
    input      cin;
    output [3:0] sum;
    output      cout;
    wire       c1, c2, c3;

    // instantiate 1-bit adders
    full_adder FA0(a[0],b[0], cin, sum[0], c1);
    full_adder FA1(a[1],b[1], c1, sum[1], c2);
    full_adder FA2(a[2],b[2], c2, sum[2], c3);
    full_adder FA3(a[3],b[3], c3, sum[3], cout);
endmodule

```

Testbench

```

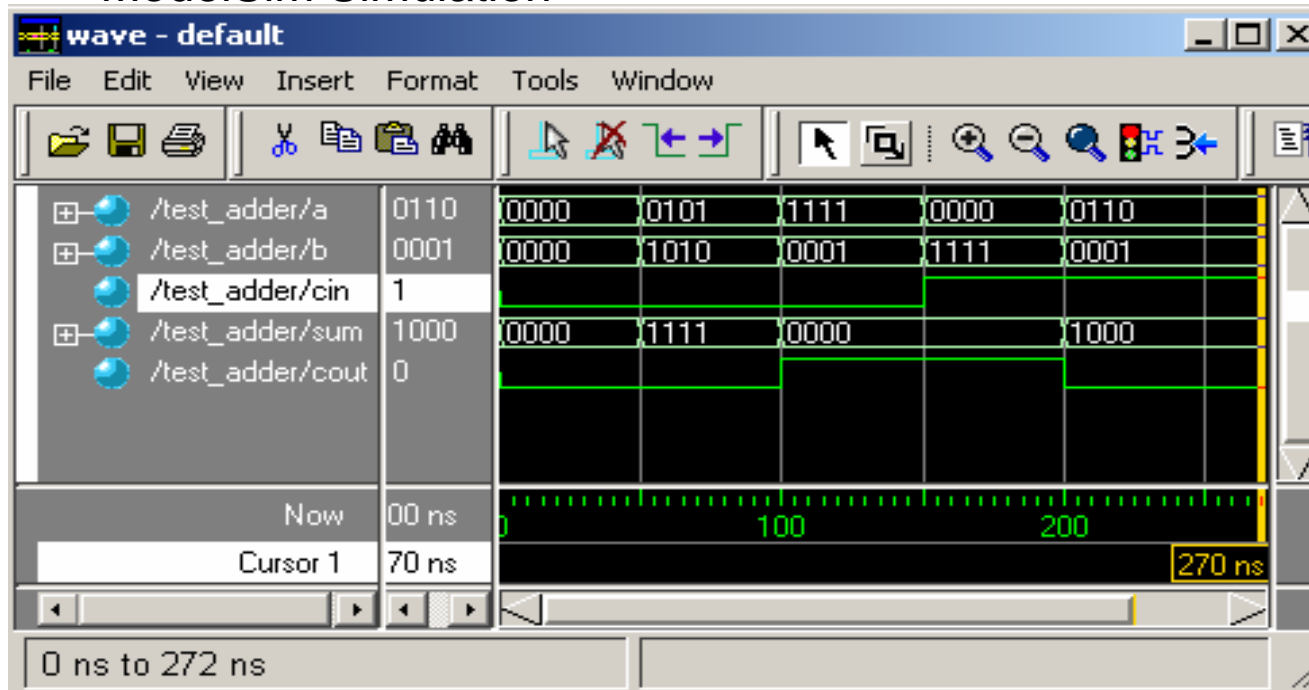
module test_adder;
    reg [3:0] a, b;
    reg       cin;
    wire [3:0] sum;
    wire       cout;

    full_adder_4bit dut(a, b, cin,
                       sum, cout);

    initial
    begin
        a = 4'b0000;
        b = 4'b0000;
        cin = 1'b0;
        #50;
        a = 4'b0101;
        b = 4'b1010;
        // sum = 1111, cout = 0
        #50;
        a = 4'b1111;
        b = 4'b0001;
        // sum = 0000, cout = 1
        #50;
        a = 4'b0000;
        b = 4'b1111;
        cin = 1'b1;
        // sum = 0000, cout = 1
        #50;
        a = 4'b0110;
        b = 4'b0001;
        // sum = 1000, cout = 0
    end // initial begin
endmodule // test_adder

```

ModelSim Simulation



Courtesy of F. Honore, D. Milliner



Summary



- Multiple levels of description: behavior, dataflow, logic and switch (not used in 6.111)
- Gate level is typically not used as it requires working out the interconnects
- Continuous assignment using `assign` allows specifying dataflow structures
- Procedural Assignment using `always` allows efficient behavioral description. Must carefully specify the sensitivity list
- Incomplete specification of `case` or `if` statements can result in non-combinational logic
- Verilog registers (`reg`) is not to be confused with a hardware memory element
- Modular design approach to manage complexity