

# L3: Introduction to Verilog (Combinational Logic)

## Acknowledgements :

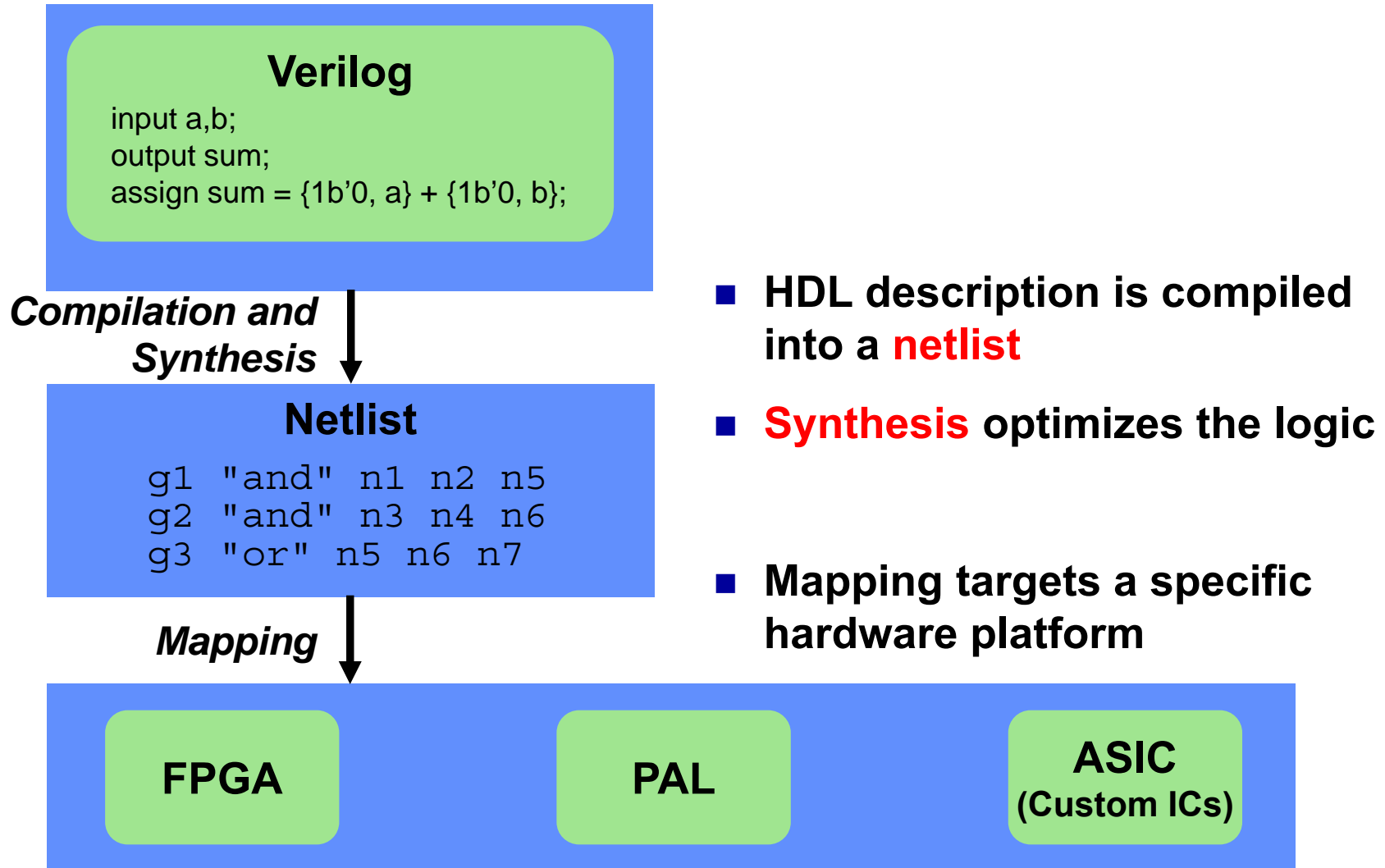
Rex Min

Lecture Notes prepared by Professor Anantha Chandrakasan

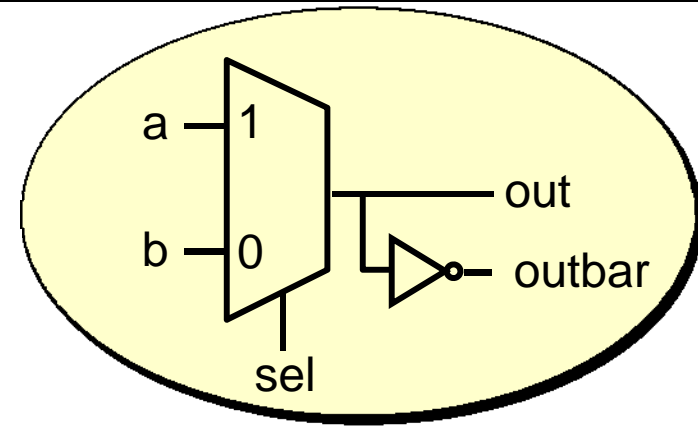
## Verilog References:

- Samir Palnitkar, Verilog HDL, Pearson Education (2nd edition).
- Donald Thomas, Philip Moorby, The Verilog Hardware Description Language, Fifth Edition, Kluwer Academic Publishers.
- J. Bhasker, Verilog HDL Synthesis (A Practical Primer), Star Galaxy Publishing

- Hardware description language (HDL) is a convenient, device-independent representation of digital logic



- Verilog designs consist of interconnected **modules**.
- A module can be an element or collection of lower level design blocks.
- A simple module with combinational logic might look like this:



$$\text{Out} = \text{sel} \bullet a + \overline{\text{sel}} \bullet b$$

*2-to-1 multiplexer with inverted output*

```
module mux_2_to_1(a, b, out,
                 outbar, sel);
```

```
// This is 2:1 multiplexor
```

```
input a, b, sel;
output out, outbar;
```

```
assign out = sel ? a : b;
assign outbar = ~out;
```

```
endmodule
```

**Declare and name a module; list its ports. Don't forget that semicolon.**

**Comment starts with //  
Verilog skips from // to end of the line**

**Specify each port as input, output, or inout**

**Express the module's behavior. Each statement executes in parallel; order does not matter.**

**Conclude the module code.**

```

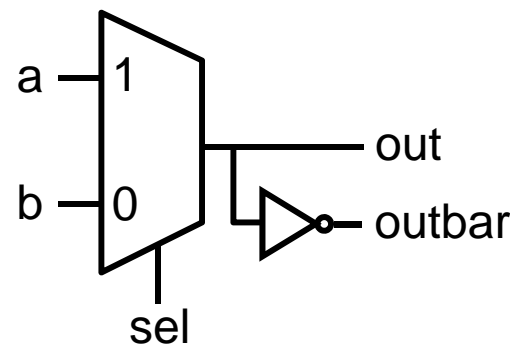
module mux_2_to_1(a, b, out,
                  outbar, sel);

    input a, b, sel;
    output out, outbar;

    assign out = sel ? a : b;
    assign outbar = ~out;

endmodule

```

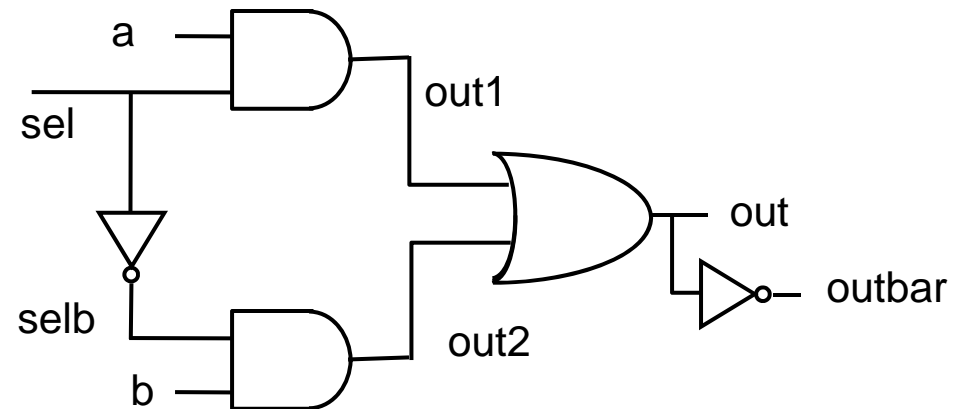


- Continuous assignments use the `assign` keyword
- A simple and natural way to represent combinational logic
- Conceptually, the right-hand expression is continuously evaluated as a function of arbitrarily-changing inputs...just like dataflow
- The target of a continuous assignment is a net driven by combinational logic
- Left side of the assignment must be a scalar or vector net or a concatenation of scalar and vector nets. It can't be a scalar or vector register (*discussed later*). Right side can be register or nets
- Dataflow operators are fairly low-level:
  - Conditional assignment: `(conditional_expression) ? (value-if-true) : (value-if-false);`
  - Boolean logic: `~, &, |`
  - Arithmetic: `+, -, *`
- Nested conditional operator (4:1 mux)
  - `assign out = s1 ? (s0 ? i3 : i2) : (s0 ? i1 : i0);`

```

module muxgate (a, b, out,
outbar, sel);
input a, b, sel;
output out, outbar;
wire out1, out2, selb;
and a1 (out1, a, sel);
not i1 (selb, sel);
and a2 (out2, b, selb);
or o1 (out, out1, out2);
assign outbar = ~out;
endmodule

```



- Verilog supports basic logic gates as primitives

- and, nand, or, nor, xor, xnor, not, buf

- can be extended to multiple inputs: e.g., nand nand3in (out, in1, in2,in3);

- bufif1 and bufif0 are tri-state buffers

- Net represents connections between hardware elements. Nets are declared with the keyword `wire`.

- Procedural assignment allows an alternative, often higher-level, behavioral description of combinational logic
- Two structured procedure statements: `initial` and `always`
- Supports richer, C-like control structures such as `if`, `for`, `while`, `case`

```
module mux_2_to_1(a, b, out,
                 outbar, sel);
  input a, b, sel;
  output out, outbar;
```

Exactly the same as before.

```
  reg out, outbar;
```

Anything assigned in an `always` block must *also* be declared as type `reg` (next slide)

```
  always @ (a or b or sel)
```

Conceptually, the `always` block runs *once* whenever a signal in the **sensitivity list** changes value

```
  begin
```

```
    if (sel) out = a;
    else out = b;
```

```
    outbar = ~out;
```

Statements within the `always` block are executed sequentially. Order matters!

```
  end
```

Surround multiple statements in a single `always` block with `begin/end`.

```
endmodule
```

- In digital design, registers represent memory elements (we will study these in the next few lectures)
- Digital registers need a clock to operate and update their state on certain phase or edge
- Registers in Verilog should not be confused with hardware registers
- **In Verilog, the term register (`reg`) simply means a variable that can hold a value**
- Verilog registers don't need a clock and don't need to be driven like a net. Values of registers can be changed anytime in a simulation by assuming a new value to the register

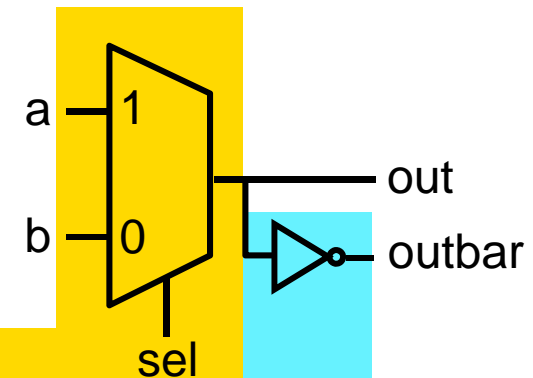
- Procedural and continuous assignments can (and often do) co-exist within a module
- Procedural assignments update the value of `reg`. The value will remain unchanged till another procedural assignment updates the variable. This is the main difference with continuous assignments in which the right hand expression is constantly placed on the left-side

```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
  input a, b, sel;  
  output out, outbar;  
  reg out;
```

```
  always @ (a or b or sel)  
  begin  
    if (sel) out = a;  
    else out = b;  
  end
```

```
  assign outbar = ~out;
```

```
endmodule
```



*procedural  
description*

*continuous  
description*



- case and if may be used interchangeably to implement conditional execution within always blocks
- case is easier to read than a long string of if...else statements

```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
    input a, b, sel;  
    output out, outbar;  
    reg out;  
  
    always @ (a or b or sel)  
    begin  
        if (sel) out = a;  
        else out = b;  
    end  
  
    assign outbar = ~out;  
  
endmodule
```

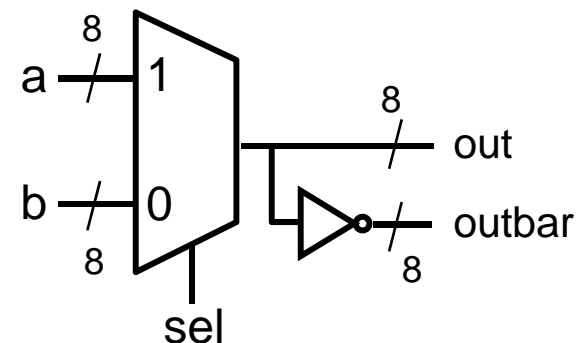
```
module mux_2_to_1(a, b, out,  
                 outbar, sel);  
    input a, b, sel;  
    output out, outbar;  
    reg out;  
  
    always @ (a or b or sel)  
    begin  
        case (sel)  
            1'b1: out = a;  
            1'b0: out = b;  
        endcase  
    end  
  
    assign outbar = ~out;  
  
endmodule
```

**Note: Number specification notation: <size>'<base><number>**

**(4'b1010 if a 4-bit binary value, 16'h6cda is a 16 bit hex number, and 8'd40 is an 8-bit decimal value)**

- Multi-bit signals and buses are easy in Verilog.
- 2-to-1 multiplexer with 8-bit operands:

```
module mux_2_to_1(a, b, out,  
                  outbar, sel);  
    input [7:0] a, b;  
    input sel;  
    output [7:0] out, outbar;  
    reg [7:0] out;  
    always @ (a or b or sel)  
    begin  
        if (sel) out = a;  
        else out = b;  
    end  
    assign outbar = ~out;  
endmodule
```



- {m,n} Concatenate m to n, creating larger vector

```
// if the MSB of a is high, this module
// concatenates 1111 to the vector.  With signed
// binary numbers, this is called sign extension.

module sign_extend(a, out);
    input [3:0] a;
    output [7:0] out;

    assign out = a[3] ? {4'b1111,a} : {4'b0000,a};
endmodule
```

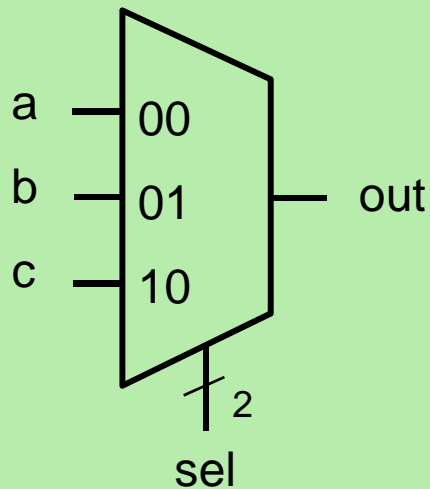
- Verilog's built-in arithmetic makes a 32-bit adder easy:

```
module add32(a, b, sum);  
    input [31:0] a,b;  
    output [31:0] sum;  
    assign sum = a + b;  
endmodule
```

- A 32-bit adder with carry-in and carry-out:

```
module add32_carry(a, b, cin, sum, cout);  
    input [31:0] a,b;  
    input cin;  
    output [31:0] sum;  
    output cout;  
    assign {cout, sum} = a + b + cin;  
endmodule
```

## Goal:



### **3-to-1 MUX**

('11' input is a don't-care)

## Proposed Verilog Code:

```

module maybe_mux_3to1(a, b, c,
                      sel, out);

    input [1:0] sel;
    input a,b,c;
    output out;
    reg out;

    always @(a or b or c or sel)
    begin
        case (sel)
            2'b00: out = a;
            2'b01: out = b;
            2'b10: out = c;
        endcase
    end
endmodule

```

*Is this a 3-to-1 multiplexer?*

```

module maybe_mux_3to1(a, b, c,
                      sel, out);

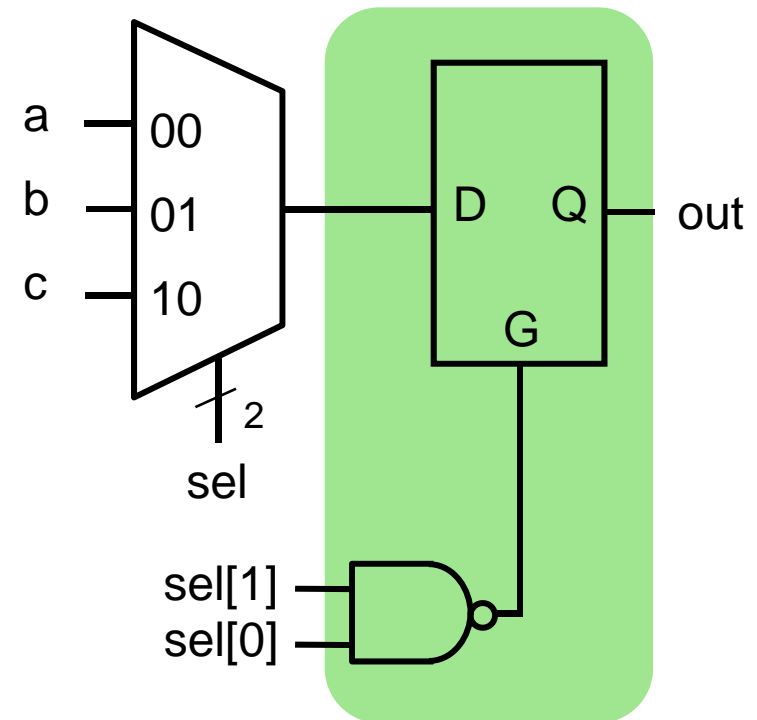
  input [1:0] sel;
  input a,b,c;
  output out;
  reg out;

  always @(a or b or c or sel)
  begin
    case (sel)
      2'b00: out = a;
      2'b01: out = b;
      2'b10: out = c;
    endcase
  end
endmodule

```

if out is not assigned during any pass through the always block, then **the previous value must be retained!**

## Synthesized Result:



- Latch memory “latches” old data when  $G=0$  (we will discuss latches later)
- In practice, we almost *never* intend this

- Precede all conditionals with a default assignment for all signals assigned within them...

```
always @(a or b or c or sel)
begin
    out = 1'bx;
    case (sel)
        2'b00: out = a;
        2'b01: out = b;
        2'b10: out = c;
    endcase
end
endmodule
```

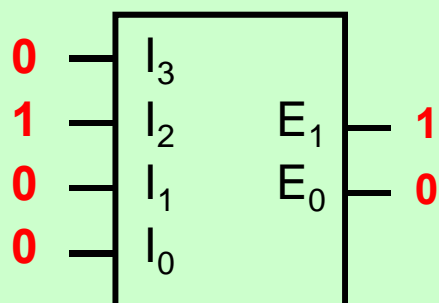
---

```
always @(a or b or c or sel)
begin
    case (sel)
        2'b00: out = a;
        2'b01: out = b;
        2'b10: out = c;
        default: out = 1'bx;
    endcase
end
endmodule
```

- ...or, fully specify all branches of conditionals and assign all signals from all branches
  - For each if, include else
  - For each case, include default

## Goal:

### 4-to-2 Binary Encoder



$I_3$	$I_2$	$I_1$	$I_0$	$E_1$	$E_0$
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1
all others				X	X

## Proposed Verilog Code:

```

module binary_encoder(i, e);
  input [3:0] i;
  output [1:0] e;
  reg [1:0] e;

  always @(i)
  begin
    if (i[0]) e = 2'b00;
    else if (i[1]) e = 2'b01;
    else if (i[2]) e = 2'b10;
    else if (i[3]) e = 2'b11;
    else e = 2'bxx;
  end
endmodule

```

***What is the resulting circuit?***



**Intent:** if more than one input is 1, the result is a don't-care.

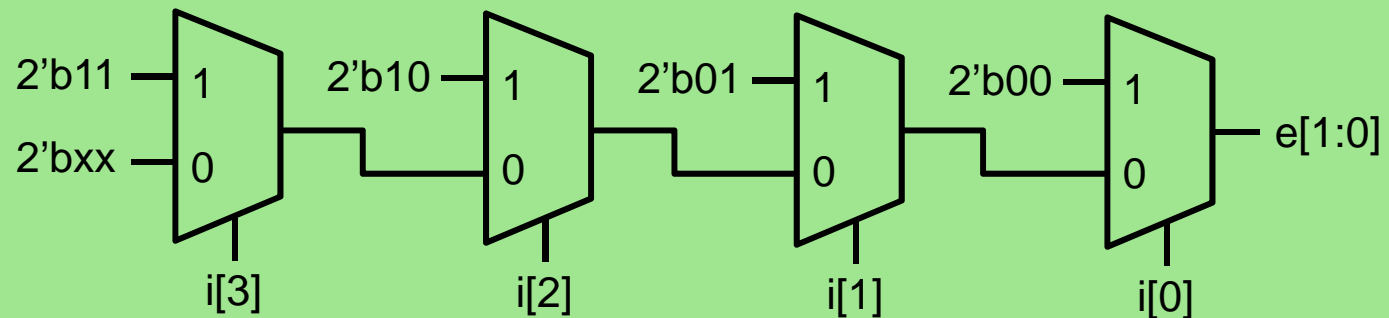
$I_3$	$I_2$	$I_1$	$I_0$	$E_1$	$E_0$
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1
all others				X	X

**Code:** if  $i[0]$  is 1, the result is 00 regardless of the other inputs.  $i[0]$  takes the highest priority.

```

if (i[0]) e = 2'b00;
else if (i[1]) e = 2'b01;
else if (i[2]) e = 2'b10;
else if (i[3]) e = 2'b11;
else e = 2'bxx;
end
    
```

**Inferred Result:**



- **if-else and case statements are interpreted very literally!**  
Beware of unintended **priority logic**.

- Make sure that `if-else` and `case` statements are *parallel*
  - If **mutually exclusive conditions** are chosen for each branch...
  - ...then synthesis tool can generate a simpler circuit that evaluates the branches in parallel

## Parallel Code:

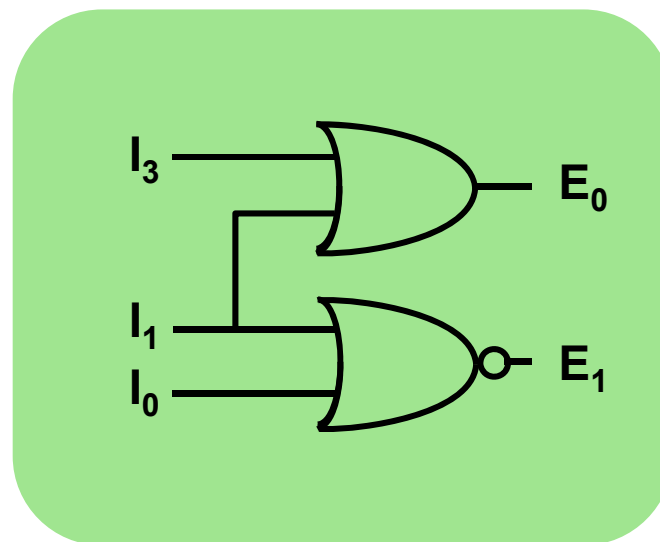
```

module binary_encoder(i, e);
  input [3:0] i;
  output [1:0] e;
  reg [1:0] e;

  always @(i)
  begin
    if (i == 4'b0001) e = 2'b00;
    else if (i == 4'b0010) e = 2'b01;
    else if (i == 4'b0100) e = 2'b10;
    else if (i == 4'b1000) e = 2'b11;
    else e = 2'bxx;
  end
endmodule

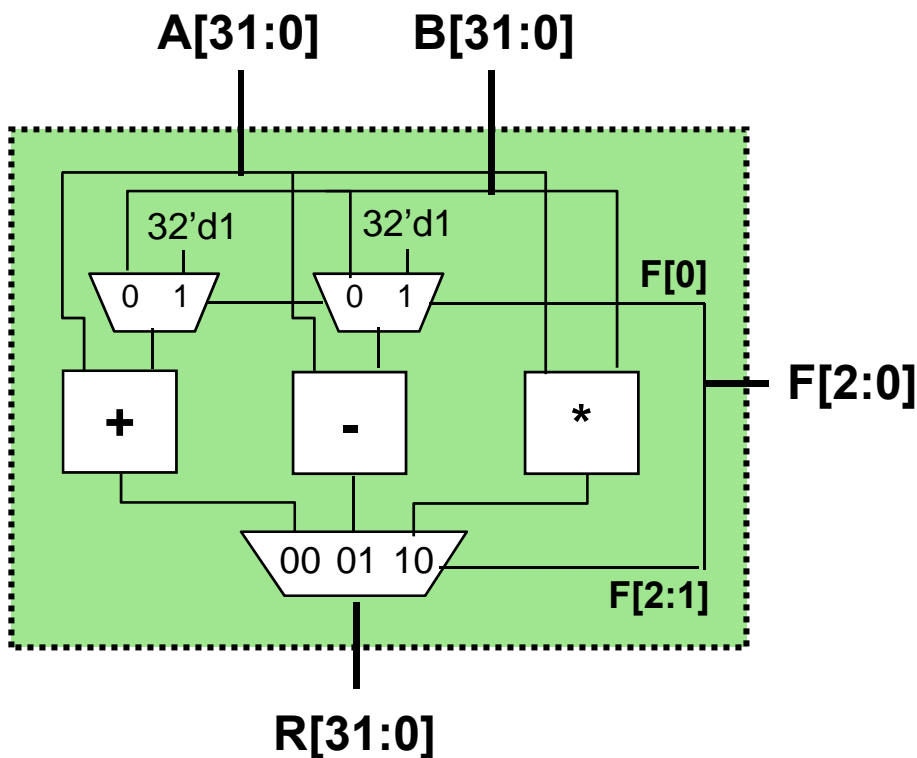
```

## Minimized Result:



- Modularity is essential to the success of large designs
- A Verilog `module` may contain submodules that are “wired together”
- High-level primitives enable direct synthesis of behavioral descriptions (functions such as additions, subtractions, shifts (<< and >>), etc.

## Example: A 32-bit ALU



## Function Table

F2	F1	F0	Function
0	0	0	A + B
0	0	1	A + 1
0	1	0	A - B
0	1	1	A - 1
1	0	X	A * B

**2-to-1 MUX**

```

module mux32two(i0,i1,sel,out);
input [31:0] i0,i1;
input sel;
output [31:0] out;

assign out = sel ? i1 : i0;

endmodule

```

**3-to-1 MUX**

```

module mux32three(i0,i1,i2,sel,out);
input [31:0] i0,i1,i2;
input [1:0] sel;
output [31:0] out;
reg [31:0] out;

always @ (i0 or i1 or i2 or sel)
begin
    case (sel)
        2'b00: out = i0;
        2'b01: out = i1;
        2'b10: out = i2;
        default: out = 32'bx;
    endcase
end
endmodule

```

**32-bit Adder**

```

module add32(i0,i1,sum);
input [31:0] i0,i1;
output [31:0] sum;

assign sum = i0 + i1;

endmodule

```

**32-bit Subtractor**

```

module sub32(i0,i1,diff);
input [31:0] i0,i1;
output [31:0] diff;

assign diff = i0 - i1;

endmodule

```

**16-bit Multiplier**

```

module mul16(i0,i1,prod);
input [15:0] i0,i1;
output [31:0] prod;

// this is a magnitude multiplier
// signed arithmetic later
assign prod = i0 * i1;

endmodule

```

## Given submodules:

```

module mux32two(i0,i1,sel,out);
module mux32three(i0,i1,i2,sel,out);
module add32(i0,i1,sum);
module sub32(i0,i1,diff);
module mul16(i0,i1,prod);

```

## Declaration of the ALU Module:

```

module alu(a, b, f, r);
  input [31:0] a, b;
  input [2:0] f;
  output [31:0] r;

```

```

wire [31:0] addmux_out, submux_out;
wire [31:0] add_out, sub_out, mul_out;

```

```

mux32two    adder_mux(b, 32'd1, f[0], addmux_out);
mux32two    sub_mux(b, 32'd1, f[0], submux_out);
add32      our_adder(a, addmux_out, add_out);
sub32      our_subtractor(a, submux_out, sub_out);
mul16      our_multiplier(a[15:0], b[15:0], mul_out);
mux32three output_mux(add_out, sub_out, mul_out, f[2:1], r);

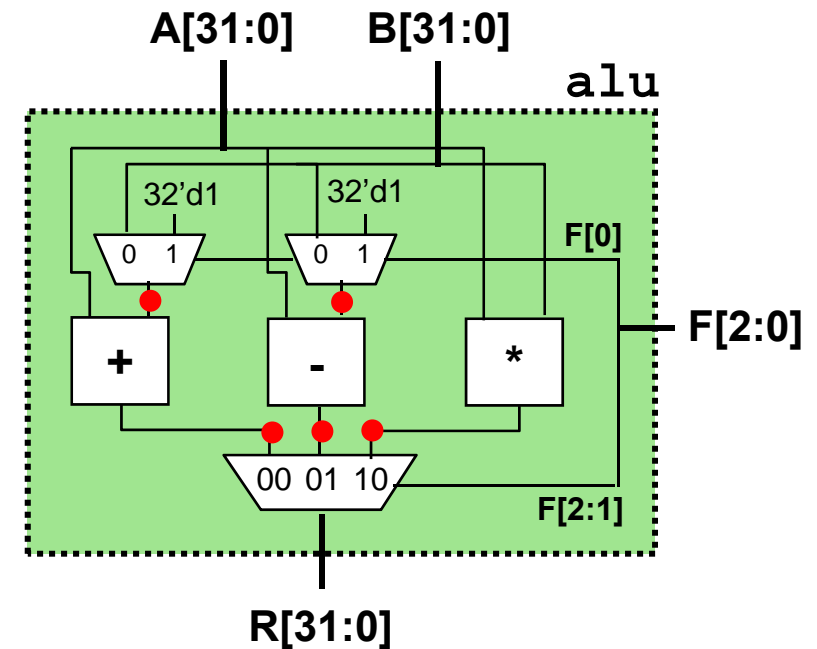
```

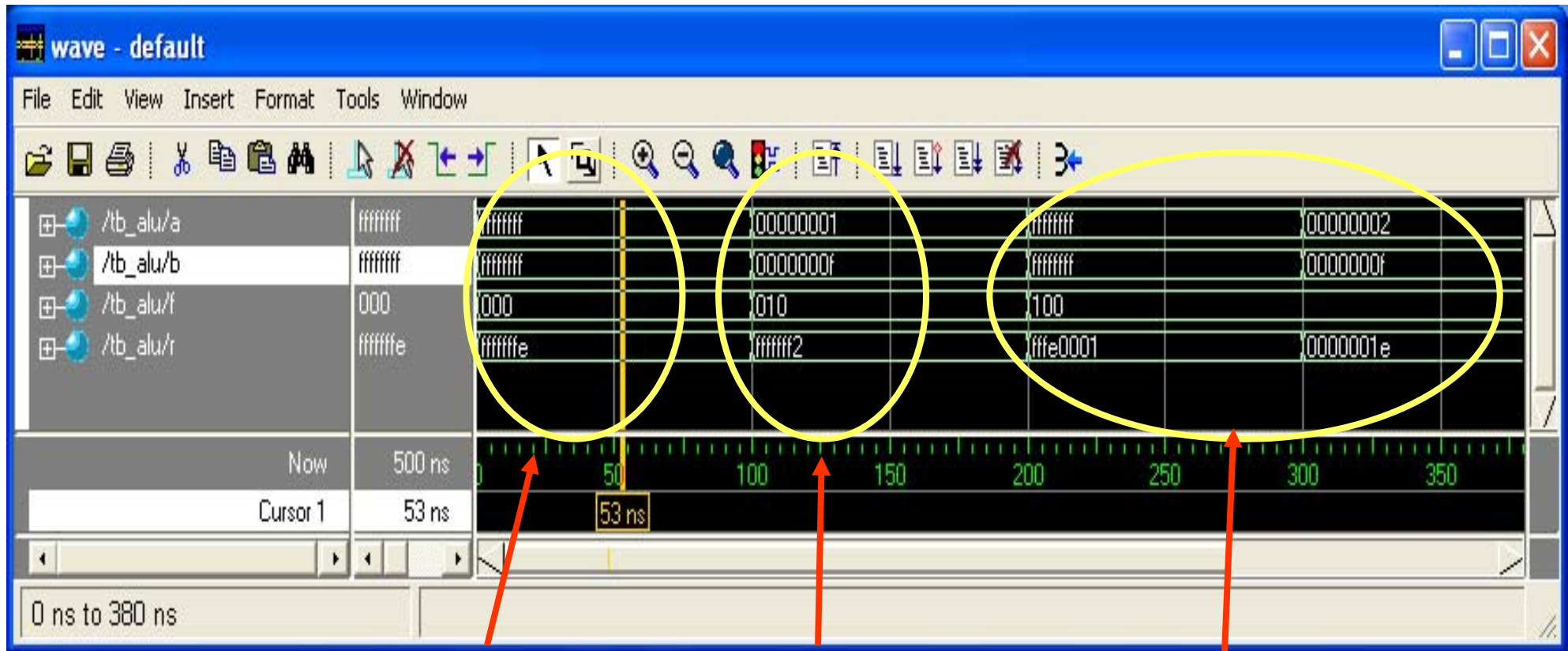
endmodule

module names

(unique) instance names

corresponding wires/regs in module alu





addition

subtraction

multiplication

- ModelSim used for behavior level simulation (pre-synthesis) – no timing information
- ModelSim can be run as a stand alone tool or from Xilinx ISE which allows simulation at different levels including **Behavioral and Post-Place-and-Route**

- **Explicit port naming allows port mappings in arbitrary order: better scaling for large, evolving designs**

Given Submodule Declaration:

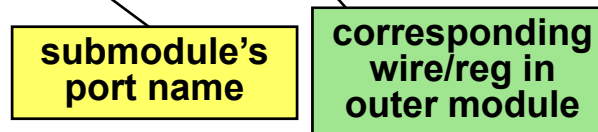
```
module mux32three(i0,i1,i2,sel,out);
```

Module Instantiation with Ordered Ports:

```
mux32three output_mux(add_out, sub_out, mul_out, f[2:1], r);
```

Module Instantiation with Named Ports:

```
mux32three output_mux(.sel(f[2:1]), .out(r), .i0(add_out),  
                      .i1(sub_out), .i2(mul_out));
```



- **Built-in Verilog gate primitives may be instantiated as well**
  - **Instantiations may omit instance name and must be ordered:**

```
and(out, in1, in2, ... inN);
```

- **Bitwise operators** perform bit-sliced operations on vectors
  - $\sim(4'b0101) = \{\sim 0, \sim 1, \sim 0, \sim 1\} = 4'b1010$
  - $4'b0101 \& 4'b0011 = 4'b0001$
- **Logical operators** return one-bit (true/false) results
  - $!(4'b0101) = \sim 1 = 1'b0$
- **Reduction operators** act on each bit of a single input vector
  - $\&(4'b0101) = 0 \& 1 \& 0 \& 1 = 1'b0$
- **Comparison operators** perform a Boolean test on two arguments

## Bitwise

$\sim a$	NOT
$a \& b$	AND
$a   b$	OR
$a \wedge b$	XOR
$a \sim \wedge b$	XNOR

## Logical

$!a$	NOT
$a \&\& b$	AND
$a    b$	OR

## Reduction

$\&a$	AND
$\sim \&$	NAND
$ $	OR
$\sim  $	NOR
$\wedge$	XOR

## Comparison

$a < b$ $a > b$ $a \leq b$ $a \geq b$	Relational
$a == b$ $a != b$	[in]equality returns x when x or z in bits. Else returns 0 or 1
$a === b$ $a !== b$	case [in]equality returns 0 or 1 based on bit by bit comparison

*Note distinction between  $\sim a$  and  $!a$*



# ModelSim/Testbench Introduction: Demo this week in Lab by TAs

## Full Adder (1-bit)

```

module full_adder (a, b, cin,
                  sum, cout);
    input  a, b, cin;
    output sum, cout;
    reg    sum, cout;

    always @(a or b or cin)
    begin
        sum = a ^ b ^ cin;
        cout = (a & b) | (a & cin) | (b & cin);
    end
endmodule

```

## Full Adder (4-bit)

```

module full_adder_4bit (a, b, cin, sum,
                       cout);
    input[3:0] a, b;
    input      cin;
    output [3:0] sum;
    output      cout;
    wire       c1, c2, c3;

    // instantiate 1-bit adders
    full_adder FA0(a[0],b[0], cin, sum[0], c1);
    full_adder FA1(a[1],b[1], c1, sum[1], c2);
    full_adder FA2(a[2],b[2], c2, sum[2], c3);
    full_adder FA3(a[3],b[3], c3, sum[3], cout);
endmodule

```

## Testbench

```

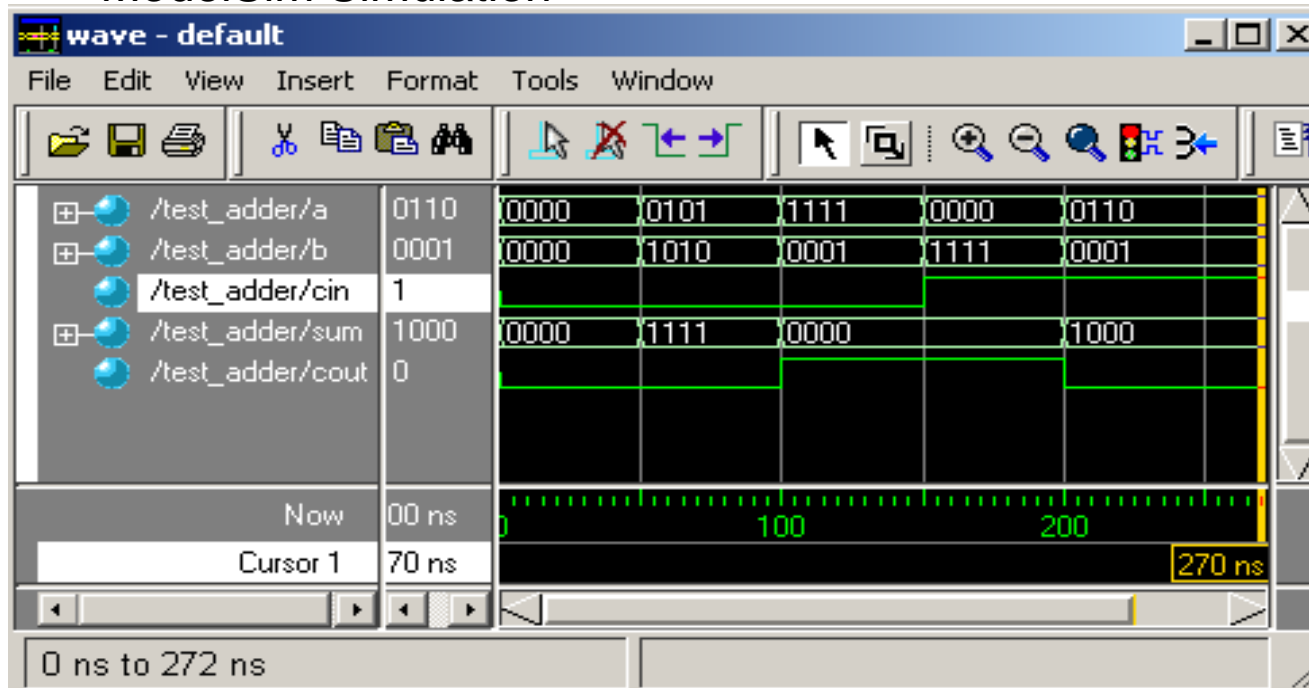
module test_adder;
    reg [3:0] a, b;
    reg      cin;
    wire [3:0] sum;
    wire      cout;

    full_adder_4bit dut(a, b, cin,
                       sum, cout);

    initial
    begin
        a = 4'b0000;
        b = 4'b0000;
        cin = 1'b0;
        #50;
        a = 4'b0101;
        b = 4'b1010;
        // sum = 1111, cout = 0
        #50;
        a = 4'b1111;
        b = 4'b0001;
        // sum = 0000, cout = 1
        #50;
        a = 4'b0000;
        b = 4'b1111;
        cin = 1'b1;
        // sum = 0000, cout = 1
        #50;
        a = 4'b0110;
        b = 4'b0001;
        // sum = 1000, cout = 0
    end // initial begin
endmodule // test_adder

```

## ModelSim Simulation



Courtesy of F. Honore, D. Milliner

- Multiple levels of description: behavior, dataflow, logic and switch (not used in 6.111)
- Gate level is typically not used as it requires working out the interconnects
- Continuous assignment using `assign` allows specifying dataflow structures
- Procedural Assignment using `always` allows efficient behavioral description. Must carefully specify the sensitivity list
- Incomplete specification of `case` or `if` statements can result in non-combinational logic
- Verilog registers (`reg`) is not to be confused with a hardware memory element
- Modular design approach to manage complexity