

Cooperative Diversity for Wireless Networks

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Introduction

- Fading is a major limitation in wireless systems
- Different forms of diversity used to overcome fading
 - Spatial Diversity
 - Temporal Diversity
 - Frequency Diversity
- Focus on specific form of spatial diversity named cooperative diversity
 - Other terminals in network act as relays between source and dest.
 - Relays experience independent fading and act as a virtual antenna array

Background

- Scalar Fading Channel
 - Slow/Fast Fading
 - Transmitter CSI
- MIMO Communication System
 - Diversity Gain
 - Multiplexing Gain
- Connection to Cooperative-Diversity

Scalar Fading Channel

Input signal experiences Rayleigh fading h and additive Gaussian noise w

$$y = hx + w$$

where $x \sim CN(0, \sigma_x^2)$, $w \sim CN(0, \sigma_w^2)$

SNR:

$$\rho \triangleq SNR = \sigma_x^2 / \sigma_w^2$$

Two major questions:

- How often does the realization of h change?
- Does the transmitter have knowledge of h ?

Transmission Cases

Fading:

- Slow Fading - Channel does not change within time period of interest
- Fast Fading - Channel changes on the order of symbols/codewords

Channel State Information (CSI):

- Tx CSI - Transmitter has knowledge of h
- No Tx CSI - Transmitter does not know h , only its distribution

	Slow Fading	Fast Fading
Tx CSI	Deterministic Capacity $C = \log(1 + h ^2\rho)$	Time Water-filling
No Tx CSI	Outage Capacity R, p^{out}	Ergodic Capacity $C = E[\log(1 + h ^2\rho)]$

Slow Fading/No Tx CSI

- Most difficult case
- Any non-zero rate chosen by the transmitter may be above the supported rate of the channel
- Channel is not changing with time (non-ergodic), the transmission may always fail with some non-zero probability
- Shannon capacity is zero, i.e., system cannot guarantee that any amount of information can be transmitted reliably.
- Outage capacity: given R and SNR , there is a non-zero probability that the channel does not support this rate; this probability is known as the outage probability (p^{out}).
- Can use relays to improve transmission, act like a virtual MIMO system

Multiple Antenna Systems

- The discrete-time MIMO channel model for system with m transmit and n receive antennas is

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{w}.$$

- \mathbf{H} is an $n \times m$ matrix of channel gains whose entries are i.i.d., $CN(0, 1)$
- \mathbf{H} drawn from such a distribution causes
 - Outages in the system
 - Increases degrees of freedom for communications

Diversity Gain

- High-SNR regime: the diversity gain d measures the rate at which the error probability decays with SNR ($1/\rho^d$)

$$d = - \lim_{\rho \rightarrow \infty} \frac{\log P_e(\rho)}{\log \rho}$$

- For a system with m transmit and n receive antennas, in the best case, the error can be made to decay with SNR as $1/\rho^{mn}$
- Diversity is a result of the mn independent paths between the transmit and receive antennas.

Multiplexing Gain

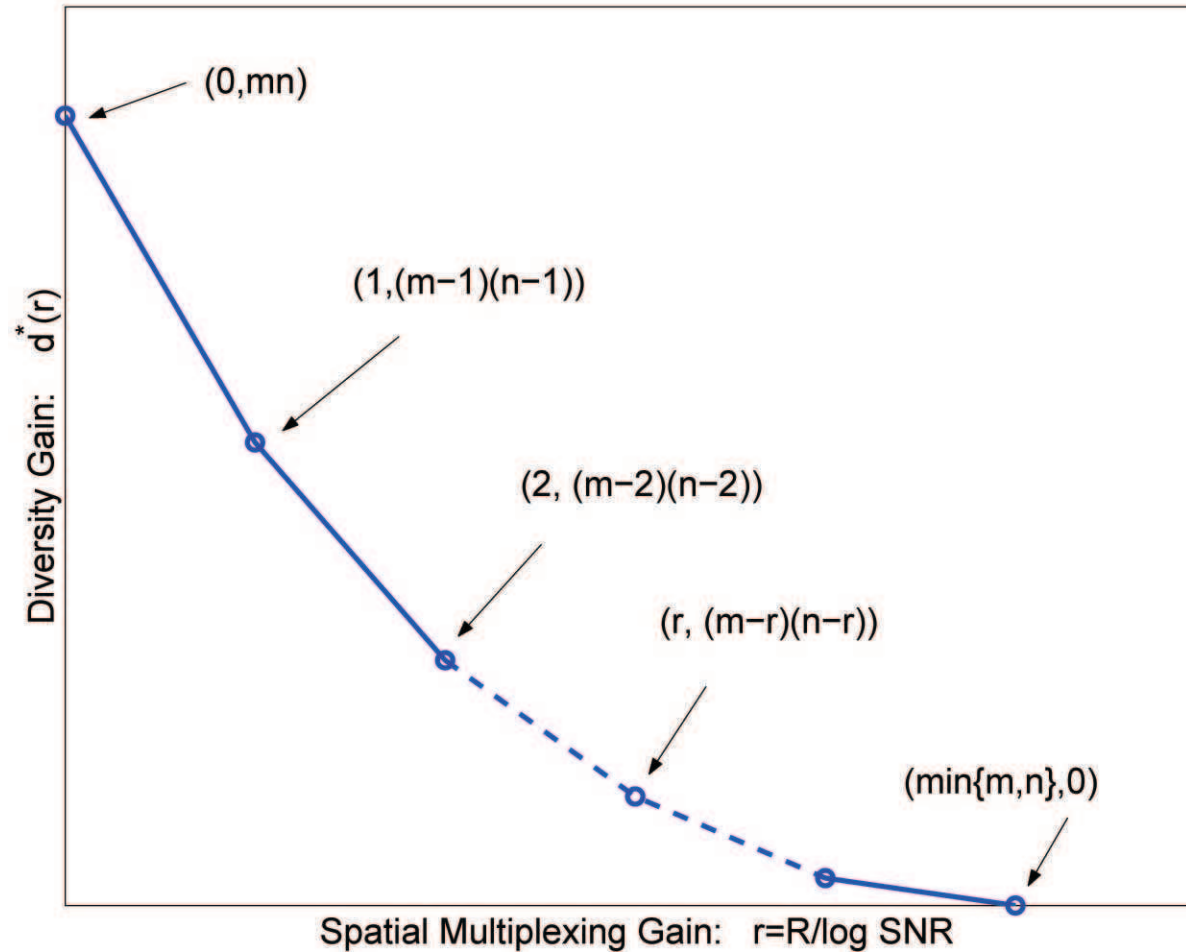
- Because the path gains are independent, the channel matrix is well-conditioned with high probability
- Can view the MIMO system as $\min(m, n)$ independent spatial channels between the transmitter and receiver.
- Fast fading ergodic capacity: $C(\rho) = \min(m, n) \log \rho + O(1)$
- For slow fading, spatial multiplexing gain r measures the rate at which R increases with respect $\log \rho$ and is formally defined as:

$$r = \lim_{\rho \rightarrow \infty} \frac{R(\rho)}{\log \rho}, \quad R \doteq r \log \rho$$

- Multiplexing gain can be thought of as the number of independent spatial channels being used optimally, which is upper-bounded by $\min(m, n)$.

Diversity-Multiplexing Tradeoff

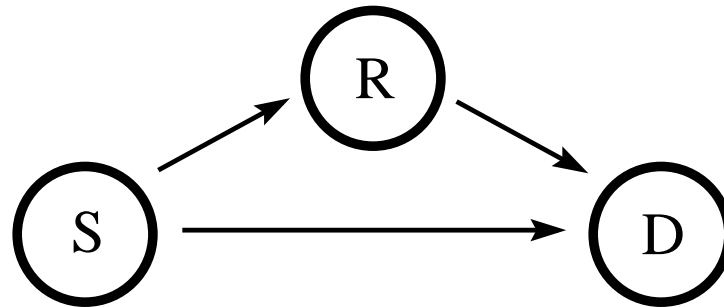
- Natural to talk about a tradeoff between the r and d



Connection to Cooperative-Diversity

- Scalar fading channel provides a lower bound for the performance of any valid cooperative-diversity scheme
 - Can always ignore use of the relay (non-cooperative protocol) and achieve scalar channel performance
- 2×1 MIMO system provides an upper bound
 - In MIMO system, the transmitter has full control over both antennas and can code across both of them, allowing more flexibility.
 - Often referred to as genie-aided protocol

System Model



- $x_s[n]$ and $x_r[n]$ are the transmitted signals from the source and relay
- $y_r[n]$ and $y_d[n]$ are the received signals of the relay and destination
- Direct source to destination transmission modeled as:

$$y_d[n] = a_{s,d}x_s[n] + z_d[n]$$

- Fading $a_{i,j} \sim CN(0, \sigma_{i,j}^2)$ and independent
- Noise $z_j[n] \sim CN(0, N_0)$, mutually independent

Model Assumptions

- Half duplex: antennas cannot transmit and receive simultaneously
- Channels experience independent flat Rayleigh-fading and white complex Gaussian noise.
- Slow fading environment
- Transmitters have no CSI. Receivers have full CSI.
- Each terminal has one antenna with power constraint, P .

Case Analysis

1. Single Relay Systems

- Amplify-and-Forward
- Decode-and-Forward

2. Multiple Terminal System

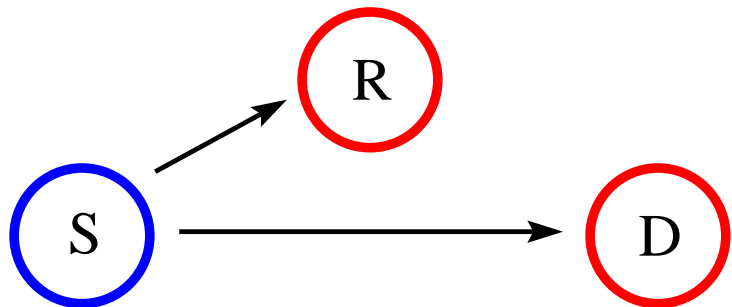
- AF Revisited
- DF Revisited
- Cooperative Broadcast Channel
- Cooperative Multiple-Access Channel

Single Relay Systems

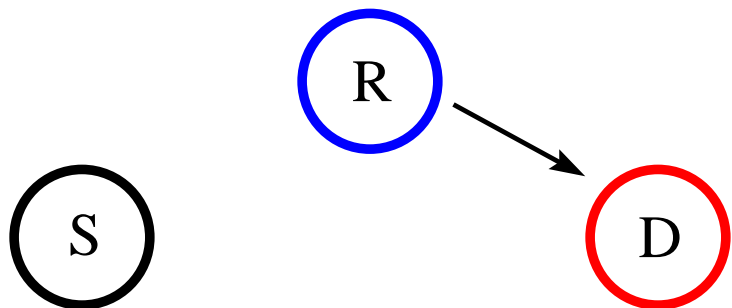
- Three terminals: source, relay, destination
- Relay can either amplify or decode, re-encode, and transmit
- Work by Laneman-Wornell-Tse
 - Additional constraint that only one terminal is transmitting at a time
- Work by Azarian-El Gamal-Schniter
 - Allow multiple terminals to transmit simultaneously
 - Protocols optimal in certain regimes.

Basic Amplify-and-Forward

- Relay can only transmit amplified version of received signal



	$n = 1, \dots, N/2$
R	$a_{s,r}x_s[n] + z_r[n]$
D	$a_{s,d}x_s[n] + z_d[n]$



	$n = N/2 + 1, \dots, N$
R	
D	$a_{r,d}(\beta y_r[n - N/2]) + z_d[n]$

- Relay power constraint: $\beta \leq \sqrt{\frac{P}{|a_{s,r}|^2 P + N_o}}$

Basic AF Performance

- Mutual information between the input and 2 outputs as

$$I_{AF} = \frac{1}{2} \log(1 + \rho |a_{s,d}|^2 + f(\rho |a_{s,r}|^2, \rho |a_{r,d}|^2))$$

where

$$f(x, y) = \frac{xy}{x + y + 1}.$$

- For an outage event (small fading coefficients), $f(x, y)$ can be approximated as $\min(x, y) \triangleq \gamma$.

Basic AF Performance cont.

- Thus, our outage event becomes ($I \leq R$):

$$\begin{aligned}\frac{1}{2} \log(\rho |a_{s,d}|^2 + \rho\gamma) &< R \\ \log(\rho) + \log(|a_{s,d}|^2 + \gamma) &< 2r \log \rho \\ \log(|a_{s,d}|^2 + \gamma) &< (2r - 1) \log \rho \\ |a_{s,d}|^2 + \gamma &< \rho^{(2r-1)}\end{aligned}$$

- Since γ and $|a_{s,d}|^2$ are both exponential, the probability of outage becomes $(\rho^{2r-1})^2$
- Diversity-multiplexing tradeoff of $d(r) = 2(1 - 2r)$

Non-Orthogonal Amplify-and-Forward

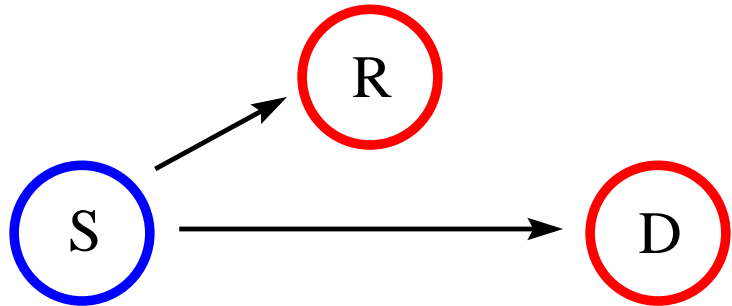
AF general form:

$$\mathbf{y} = \begin{bmatrix} a_{s,d}A_1 & 0 \\ a_{s,r}a_{r,d}\beta B & a_{s,d}A_2 \end{bmatrix} \mathbf{x} + \begin{bmatrix} 0 \\ a_{r,d}B \end{bmatrix} \mathbf{z}_r + \mathbf{z}_d$$

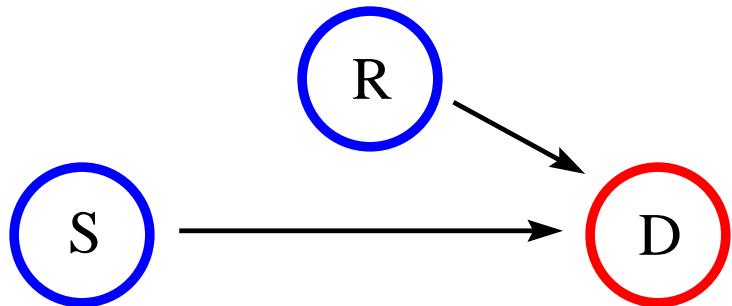
where $\mathbf{y}, \mathbf{x}, \mathbf{z}_r, \mathbf{z}_d \in \mathbb{C}^N$, $\mathbf{A}_1 \in \mathbb{C}^{N',N'}$, $\mathbf{B} \in \mathbb{C}^{N-N',N'}$, $\mathbf{A}_2 \in \mathbb{C}^{N-N',N-N'}$.

- Initial protocol: $N' = N/2$, $A_1 = B = I_{N/2}$, and $A_2 = 0$.
- NAF protocol: $N' = N/2$ and $A_1 = A_2 = B = I_{N/2}$
 - $A_1 = A_2 = I_{N/2}$ maximizes mutual information between x and y
 - B diagonal to prevent ISI, size $N/2$ for max symbol repetition

NAF cont.



	$n = 1, \dots, N/2$
R	$a_{s,r}x_s[n] + z_r[n]$
D	$a_{s,d}x_s[n] + z_d[n]$

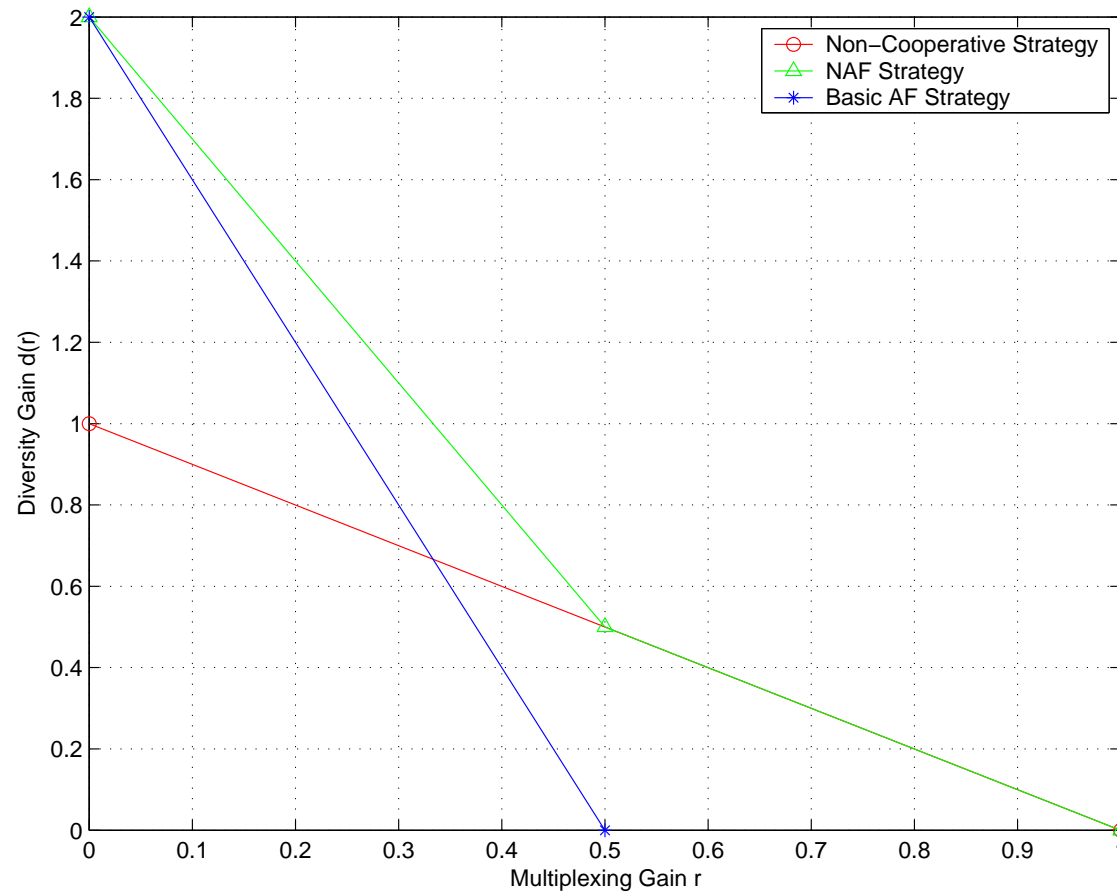


	$n = N/2 + 1, \dots, N$
R	
D	$(a_{r,d}(\beta y_r[n - N/2]) + a_{s,d}x_s[n] + z_d[n])$

- Achieves optimal diversity-multiplexing curve for single AF systems

$$d(r) = (1 - r) + (1 - 2r)^+$$

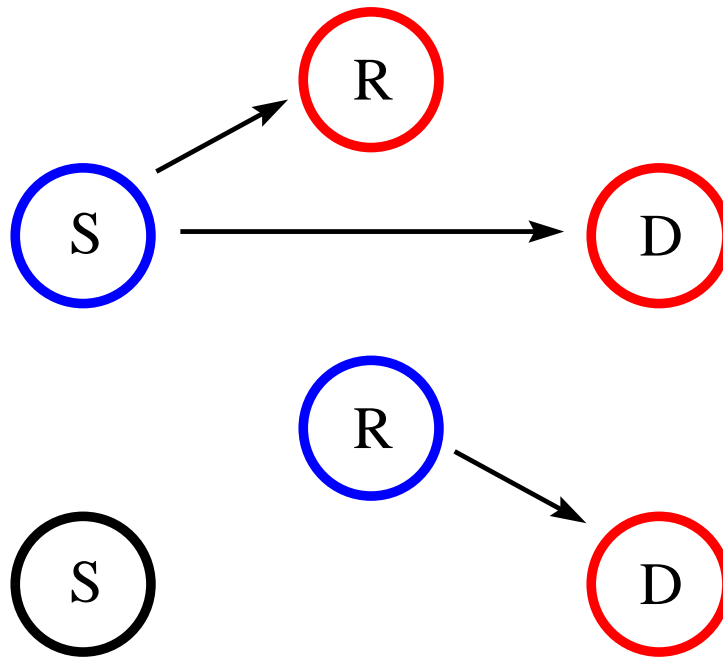
NAF cont.



- For $r > 0.5$, the NAF is equivalent to the non-cooperative method. Relay provides no benefit in this regime.

Decode-and-Forward

- Relay decodes message from source, re-encodes and transmits it



	$n = 1, \dots, N/2$
R	$a_{S,r}x_S[n] + z_r[n]$
D	$a_{S,d}x_S[n] + z_d[n]$

	$n = N/2 + 1, \dots, N$
R	
D	$a_{r,d}\hat{x}_S[n - N/2] + z_d[n]$

- Relay required to decode correctly

Decode-and-Forward cont.

Mutual information is given by

$$I_{DF} = \frac{1}{2} \min \left(\log(1 + \rho |a_{s,r}|^2), \log(1 + \rho |a_{s,d}|^2 + \rho |a_{r,d}|^2) \right)$$

An outage (either relay or destination cannot decode) occurs for a transmission rate R if

$$\min \left(|a_{s,r}|^2, |a_{s,d}|^2 + |a_{r,d}|^2 \right) < \frac{2^{2R} - 1}{\rho}.$$

This occurs with probability

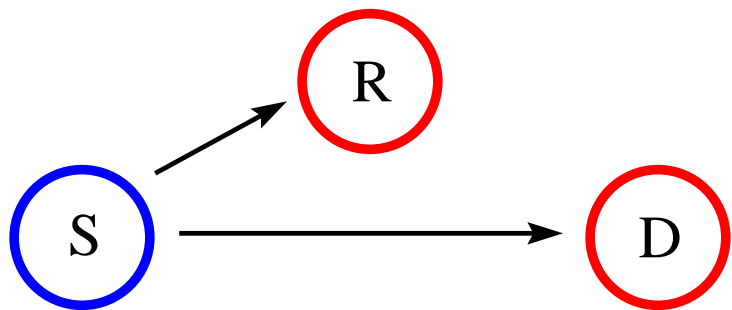
$$P_{DF}^{out}(\rho, R) \sim \frac{1}{\rho^{(1-2r)}}.$$

Thus, the diversity multiplexing tradeoff is

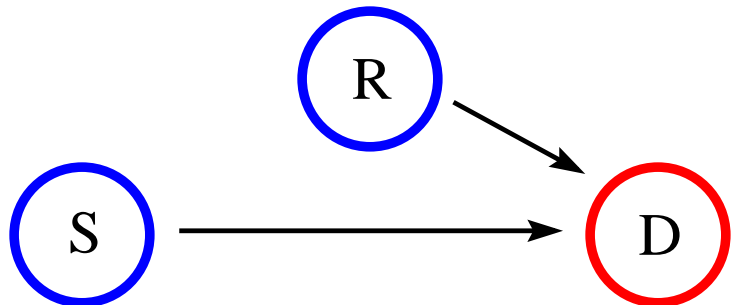
$$d(r) = 1 - 2r$$

Dynamic Decode and Forward

- Relay listens to the source until it collects sufficient energy to decode, then begins retransmitting to the destination using an independent codebook



	$n = 1, \dots, N'$
R	$a_{s,r}x_s[n] + z_r[n]$
D	$a_{s,d}x_s[n] + z_d[n]$



	$n = N' + 1, \dots, N$
R	
D	$a_{s,d}x_s[n] + a_{r,d}\hat{x}_r[n] + z_d[n]$

$$N' = \min \left[N, \left\lceil N \frac{R}{\log_2(1 + |a_{s,r}|^2 \rho)} \right\rceil \right]$$

DDF Analysis

It can be shown that the diversity-multiplexing tradeoff achieved via this DDF scheme is:

$$d(r) = \begin{cases} 2(1-r) & \text{for } \frac{1}{2} \geq r \geq 0, \\ (1-r)/r & \text{for } 1 \geq r \geq \frac{1}{2} \end{cases}$$

Outline of Proof:

- For a Gaussian random variable g , we define the exponential order of $1/|g|^2$ as ν , where

$$\nu = \lim_{\rho \rightarrow \infty} \frac{\log(|g|^2)}{\log \rho}.$$

This implies that in the high-SNR regime, $|g|^2 \doteq \rho^{-\nu}$. The resulting pdf for ν is equal to $\rho^{-\nu}$ for $\nu \geq 0$ and 0 for $\nu < 0$. Let us define ν_1 , ν_2 , and u as the exponential orders of $a_{s,d}$, $a_{r,d}$, and $a_{s,r}$ respectively.

DDF Analysis cont.

- An outage occurs if the mutual information is less than the chosen rate ($I < R$):

$$\begin{aligned}
 & Pr(N' \log(1 + |a_{s,d}|^2 \rho) + (N - N') \log(1 + (|a_{s,d}|^2 + |a_{r,d}|^2) \rho) < Nr \log(\rho)) \\
 = & Pr(N' \log(1 + \rho^{-v_1} \rho) + (N - N') \log(1 + (\rho^{-v_1} + \rho^{-v_2}) \rho) < Nr \log(\rho)) \\
 \approx & Pr(N'(1 - v_1) \log(\rho) + (N - N')(1 - \min(v_1, v_2)) \log(\rho) < Nr \log(\rho)) \\
 = & Pr((N'/N)(1 - v_1) + ((N - N')/N)(1 - \min(v_1, v_2)) < r)
 \end{aligned}$$

- Outage region O^+ :

$$O^+ = \{(v_1, v_2, u) \mid (N'/N)(1 - v_1) + ((N - N')/N)(1 - \min(v_1, v_2)) < r\}.$$

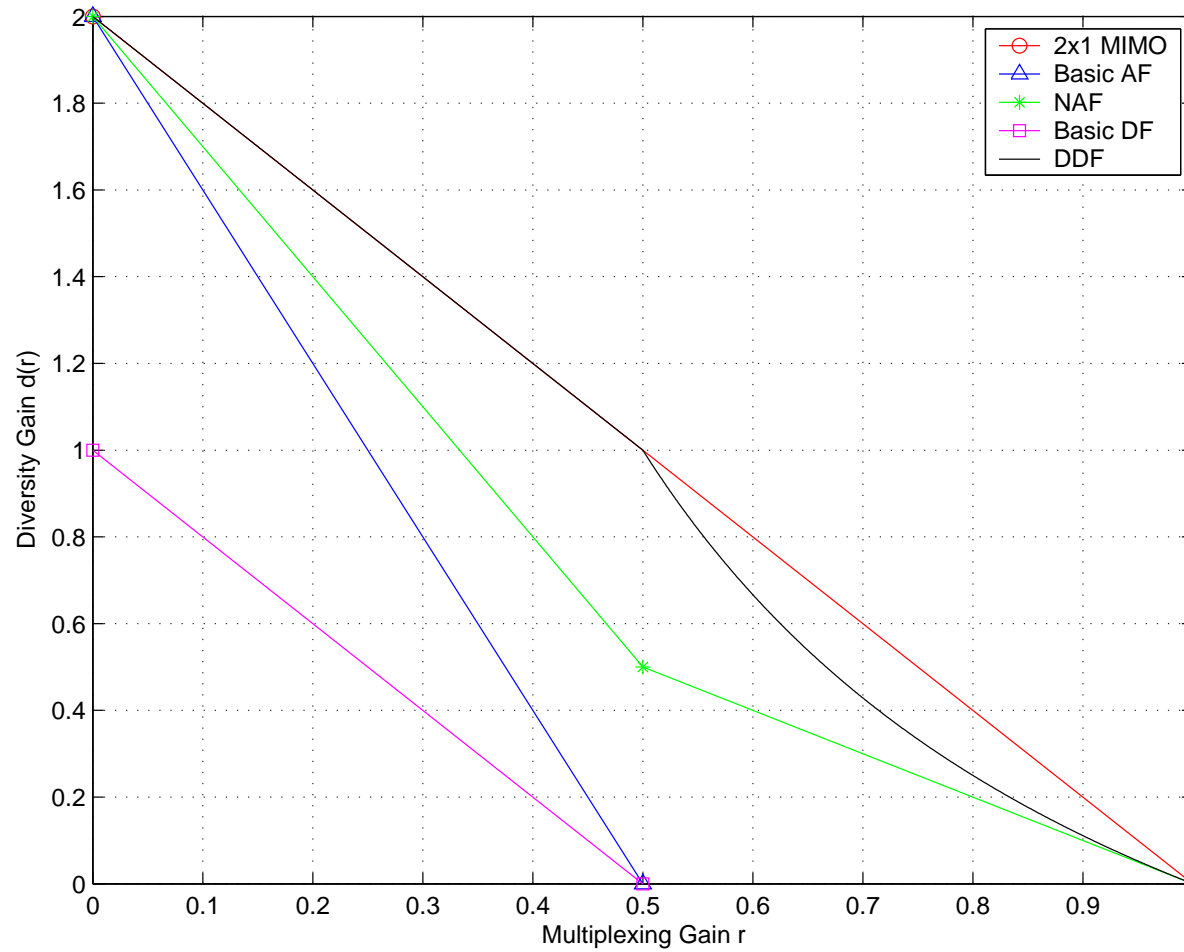
Probability of outage:

$$\begin{aligned}
 p_{out} &= \int_{(v_1, v_2) \in O^+} \rho^{-(v_1 + v_2)} dv_1 dv_2 \\
 &\approx \rho^{-d}
 \end{aligned}$$

where

$$d = \inf_{(v_1, v_2) \in O^+} (v_1 + v_2).$$

DDF Analysis cont.



DDF Remarks

Comments:

- Optimal diversity-multiplexing tradeoff for single DF relay is unknown for $r > 0.5$
- Extremely difficult to implement, requires a high level of synchronization and coordination among terminals
- Mixed AF/DF strategy cannot improve performance. Any event outage event for the DDF strategy will remain in outage through mixed strategies.

Multiple Terminals

- AF Revisited
- DF Revisited
- Cooperative Broadcast Channel
- Cooperative Multiple-Access Channel

AF Revisited

- NAF protocol generalized to $M - 1$ relays
 - Relay takes turns helping the source node (first relay is active during first time slots, second relay is active during second time slots, ...)
 - Due to power constraints on the relays, nothing can be gained by having multiple relays repeat the same symbol
- Diversity-multiplexing tradeoff

$$d(r) = (1 - r) + (M - 1)(1 - 2r)^+$$

DF Revisited

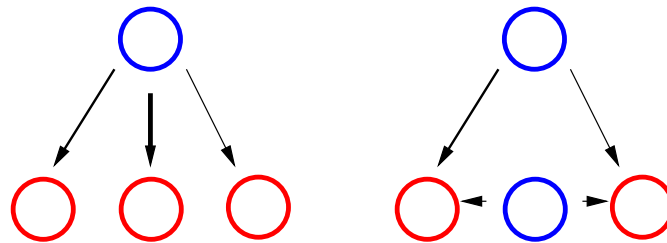
- DDF protocol generalized to $M - 1$ relay system
- When a relay can properly decode, it begins transmitting codeword using an independent codebook (DDF)
- Relays use beacon signal used to inform other relays of transmission starting time
- Diversity Multiplexing tradeoff:

$$d(r) = \begin{cases} M(1-r) & \text{for } \frac{1}{M} \geq r \geq 0, \\ 1 + \frac{(M-1)(1-2r)}{1-r} & \text{for } \frac{1}{2} \geq r \geq \frac{1}{M} \cdot \\ \frac{1-r}{r} & \text{for } 1 \geq r \geq \frac{1}{2} \end{cases}.$$

Note that this tradeoff is not known to be optimal for $r > 1/M$.

Cooperative Broadcast Channel

- Single source broadcasts information to M different destinations
- Worse case scenario: all users want to decode the full rate $R = R_c$



- Can use protocol similar to DDF
 - Source transmits codeword
 - After a destination decodes, it begins transmitting to the rest
 - D-M tradeoff same as relay channel with M relays

Multiple-Access Channel

- M sources each with independent message to the destination
- CMA-NAF protocol:
 - Sources alternate turns to transmit
 - In each interval, a single source transmits a linear combination of its own signal and the signal it received from the source before it
 - Reaches optimal diversity-multiplexing tradeoff for CMA systems:

$$d(r) = M(1 - r)$$

Summary of Results

Protocol	D-M Tradeoff $d(r)$
Non-Cooperative	$1 - r$
Basic AF	$2(1 - 2r)$
NAF	$(1 - r) + (1 - 2r)^+$
Basic DF	$1 - 2r$
DDF	$\begin{cases} 2(1 - r) & \text{for } \frac{1}{2} \geq r \geq 0, \\ (1 - r)/r & \text{for } 1 \geq r \geq \frac{1}{2} \end{cases}$
2×1 MIMO	$2(1 - r)$

- Diversity-Multiplexing Tradeoff
- Performance-Complexity Tradeoff