

## Lists and Algebraic Types

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#### Algebraic types • Algebraic types are tagged unions of products • Example data Shape = Line Pnt Pnt "union" Triangle Pnt Pnt Pnt keyword Quad Pnt Pnt Pnt new type "products" (fields) - new "constructors" (a.k.a. "tags", "disjuncts", "summands") - a k-ary constructor is applied to k type expressions October 2, 2002 http://www.csg.lcs.mit.edu/6.827

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#### Constructors are functions

 Constructors can be used as functions to create values of the type

```
let
     11 :: Shape
     11 = Line     e1     e2

     t1 :: Shape = Triangle     e3     e4     e5
     q1 :: Shape = Quad     e6     e7     e8     e9
in
...
```

where each "eJ" is an expression of type "Pnt"

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## Pattern-matching on algebraic types

• Pattern-matching is used to examine values of an algebraic type

- A pattern-match has two roles:
  - A test: "does the given value match this pattern?"
  - Binding ("if the given value matches the pattern, bind the variables in the pattern to the corresponding parts of the value")

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## Pattern-matching scope & don't cares

- Each clause starts a new scope: can reuse bound variables
- Can use "don't cares" for bound variables

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# Pattern-matching more syntax

 Functions can be defined directly using pattern-matching

```
anchorPnt :: Shape -> Pnt
anchorPnt (Line    p1 _) = p1
anchorPnt (Triangle p1 _ _) = p1
anchorPnt (Quad    p1 _ _ _) = p1
```

 Pattern-matching can be used in list comprehensions (later)

```
(Line p1 p2) <- shapes
```

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## Pattern-matching Type safety

- Given a "Line" object, it is impossible to read "the field corresponding to the third point in a Triangle object" because:
  - all unions are tagged unions
  - fields of an algebraic type can only be examined via pattern-matching

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# Special syntax

Function type constructor

Int -> Bool

Conceptually:

Function Int Bool

i.e., the arrow is an "infix" type constructor

• Tuple type constructor

(Int, Bool)

Conceptually:

Tuple2 Int Bool

Similarly for Tuple3, ...

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## Type Synonyms

```
data Point = Point Int Int
Versus

type Point = (Int,Int)
```

Type Synonyms do not create new types. It is just a convenience to improve readability.

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# **Abstract Types**

A rational number is a pair of integers but suppose we want to express it in the reduced form only. Such a restriction cannot be enforced using an algebraic type.

No pattern matching on abstract data types

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## **Examples of Algebraic types**

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#### Lists

data List t = Nil | Cons t (List t)

A list data type can be constructed in two different ways:

an empty list
or a non-empty list

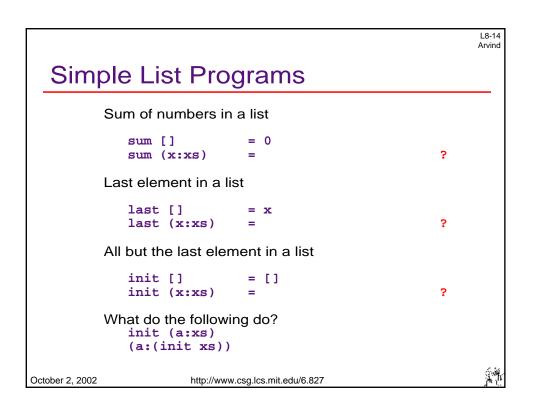
Cons x xs

the first element the rest of the elements

- All elements of a list have the same type
- The list type is recursive and polymorphic

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## Example: Split a list

```
data Token = Word String | Number Int
```

Split a list of tokens into two lists - a list words and a list of numbers.

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# Higher-order List abstractions

filter p [] = []
filter p (x:xs) =

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## Using maps and folds

- 1. Write sum in terms of fold
- 2. Write split using foldr
   split :: (List Token) -> ((List String),(List Int))
- 3. What does function fy do?

```
fy xys = map second xys
second (x,y) = y
fy ::
```

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## Flattening a List of Lists

```
append :: (List t) -> (List t) -> (List t)
append [] ys = ys
append (x:xs) ys = (x:(append xs ys))
```

```
flatten :: (List (List t)) -> (List t)
flatten [] = []
flatten (xs:xss) = append xs (concat xss)
```

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## Arithmetic Sequences: Special Lists

```
[1 .. 4] \equiv [1,2,3,4]
[1,3 .. 10] \equiv [1,3,5,7,9]
[5,4 .. 1] \equiv [5,4,3,2,1]
[5,5 .. 10] \equiv [5,5,5,...]
[5 .. ] \equiv [5,6,7,...]
```

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List Comprehensions a convenient syntax [ e | gen, gen, ...] **Examples** [fx|x<-xs] means map f xs  $[ x \mid x < -xs, (p x)]$ means filter p xs [ f x y | x <- xs, y <- ys ] means the list [(f x1 y1),...(f x1 yn), $(f \times 2 y1), \dots (f \times m yn)]$ which is defined by flatten (map ( $\ x \rightarrow (map (\ y \rightarrow e) \ ys) \ xs))$ October 2, 2002 http://www.csg.lcs.mit.edu/6.827

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#### **Three-Partitions**

Generate a list containing all three-partitions (nc1, nc2, nc3) of a number m, such that

```
• nc1 < nc2 < nc3
```

• nc1 + nc2 + nc3 = m

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# The Power of List Comprehensions [ (i,j) | i <- [1..n], j <- [1..m] ] using map point i j = (i,j) points i = map (point i) [1..m] all\_points = map points [1..n] ? October 2, 2002 http://www.csg.lcs.mit.edu/6.827

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#### Infinite Data Structures

2. ones = 1:ones
 nth 50 ones -->

3. xs = [fx | x <- a:xs]
nth 10 xs -->
?

These are well defined but *deadly* programs in pH. You will get an answer but the program may *not* terminate.

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#### Primes: The Sieve of Eratosthenes

```
primes = sieve [2..]
sieve (x:xs) = x:(sieve (filter (p x) xs))
p x y = (y mod x) ≠ 0
    nth 100 primes
```

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