

MASSACHUSETTS INSTITUTE OF TECHNOLOGY

Department of Economics

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Introduction to Moral Hazard

1. Moral Hazard in Teams

Setup:

- efforts are unobserved
- assume no uncertainty
- I players, output a product of many players

$$x = f(e_1, \dots, e_I)$$

- free rider problem (assume risk neutrality, z_i is money)

$$u_i(z_i, e_i) = z_i - e_i$$

- $s_i(x)$ sharing rule - how much money goes to agent i
- budget balancing (BB)

$$\sum s_i(x) = x$$

Question: Can we find $s_i(x)$ such that BB $\forall x$ and efficient output?

Answer: No, if s_i is differentiable and $\underbrace{e_i^*}_{\text{optimal effort}} > 0$

PROOF:

Nash Equilibrium: $e_i = \arg \max_{e_i} [s_i(f(e)) - e_i]$

First Order Condition (FOC)

$$\underbrace{s'_i}_{\text{share}} \underbrace{f_i(e)}_{i^{th} \text{ partial der of } f} = 1, \quad \forall i$$

Efficiency: $\max f(e) - \sum e_i, \quad f_i = 1$

Contradiction:

– for every agent $s'_i = 1$

– BB: $\sum s'_i(x) = 1$

Interpretation: personal benefits go one for one with costs, can't do this if benefits are shared.

- 2 Solutions:

- Boot Camp, group punishments (violates s_i differentiability assumption)
- Improve Measurement

2. Single Principal / Single Agent (1 of each)

- Tradeoffs: Insurance vs. Incentives

- efficient for principal to be insuring the workers

- 3 Formulations:

(a) State Space - output can take on a limited # of values (linear program)

- easy to understand

- analytically difficult
- (b) Parameterized Distribution - analytically elegant
- (c) General Distribution
 - * Question: $f(x | e)$ - why only a single parameter?
 - * Answer: actions can change distribution (in more than one-dimensional way); this expand richness of agent's space
 - * conceptually nice
 - * analytically difficult

3. Parameterized Distribution Formulation

- $e \rightarrow F(\underbrace{x}_{\text{output}} | e) \quad \text{density } [f(x | e)]$
 - $x \sim N(e, 1)$, $\text{variance} = 1$
 - same as $x = e + \varepsilon$, $\varepsilon \sim N(0, 1)$
- Principal: $v(w)$, $v' > 0$
- Agent: $u(w) - c(e)$, $u' > 0$, $c' > 0$, $c'' > 0$
Principal chooses $s(x)$, $v(x - s(x))$
- Behavioral Assumptions: for a given $s(x)$, e is chosen by agent iff

$$\int u(s(x)) f(x | e) dx - c(e) \geq \int u(s(x)) f(x | e') dx - c(e'), \quad \forall e' \quad (\text{IC})$$

$$\int u(s(x)) f(x | e) dx - c(e) \geq \underline{u} \quad (\text{IR})$$

Critique: agent can be given very low utility for some realizations of output: what about limited liability?

- Program:

$$\max_{s(x)} \int v(x - s(x)) f(x | e) dx$$

subject to IR, IC

- maximize Principal's utility, given agent's constraint
- Can incorporate pointwise constraint
- trying to push agent to reservation utility - Pareto optimality
- Analyze as constrained optimization
 - principal chooses sharing rule
 - remove agent entirely as a strategic player
- To solve this equation, use a 2 step procedure.
 - For any e , find $s(x)$ that satisfies IC, IR “implements” e

$$EV(e) = \int v(x - s_e(x)) f(x | e) dx$$
 - Choose e to maximize $EV(e)$