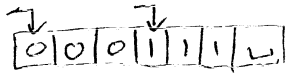


Today:
 Comp History Method
 Recursion Thm
 Review

Define a 2-headed DFA (2HDFA) to be a TM with two read-only heads, both of which can only move right or stay put (not necessarily in sync). Moreover, there is only one "L" symbol to mark the end of the input.



A 2HDFA can, for example, recognize $\{0^n 1^n \mid n \geq 0\}$
 But probably cannot recognize $\{ww^R \mid w \in \Sigma^*\}$

Q) Show $E_{2HDFA} = \{ \langle D \rangle \mid D \text{ is a 2HDFA and } L(D) = \emptyset \}$ is undecidable

Proof] Comp History method.

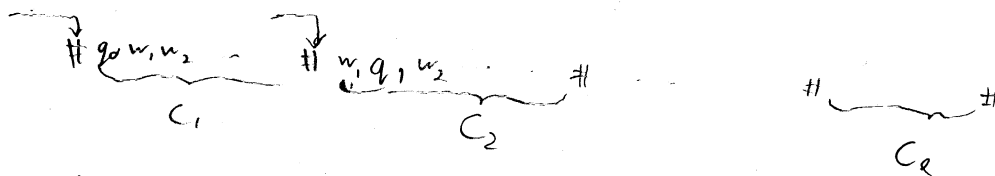
For particular M and w , show how to construct a 2HDFA $D_{M,w}$ which only accepts valid, accepting comp histories for M on w

Then, given a decider for E_{2HDFA} , we decide A_{TM} as follows

On input $\langle M, w \rangle$

1. Construct $\langle D_{M,w} \rangle$
2. Run decider for E_{2HDFA} on $\langle D_{M,w} \rangle$, accept if decider rejects, reject if decider accepts

How does $D_{M,w}$ work?



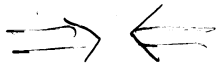
- 1) First read head verifies C_1 is correct starting config of M on w
- 2) Two heads move in lock-step to verify each C_i and C_{i+1} correctly line up according to M 's transition function
- 3) First head checks that q_{accept} is encountered. If so, and all configs seen so far are valid, accept.

Q] Define the "Busy Beaver" function $f(k) =$
 to be the maximum length string printed by a k -state machine
 which halts on a blank input.
 Show f is not computable.
 on a blank input

Proof #1 | Suppose it were computable. Then define machine

$G =$ "on blank input

- 1) Get own description $\langle G \rangle$, let $K = \#$ states in $\langle G \rangle$
- 2) Compute $f(K)$
- 3) Print $f(K) + 1$ symbols on the tape"



Alternative solution : Show how to decide the halting
 problem if we can compute the Busy Beaver function.

10/21/05

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