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Problem Set No.2

Issued: Sep 16, 2008
Due: Sep 23, 2008

Problem 1: text problem 2.18

Problem 2: text problem 2.20

Problem 3: text problem 2.22

Problem 4: text problem 2.23* (Bonus Problem)

Problem 5: text problem 2.26

Problem 6: text problem 3.1

Problem 7: text problem 3.7

Problem 8: TCP Window Evolution: Additive Increase and Multiplicative Decrease (AIMD)

Consider the following variation of Go-Back-N where the sender adjusts its window size dynamically as follows:

Sender (for current window size N)

- Upon receipt of an ACK, $N = N + 1/N$.
- Upon receipt of a NACK, $N = N/2$

Receiver

- Upon detection of a packet loss, send a NACK
- Upon successful reception of a packet, send an ACK

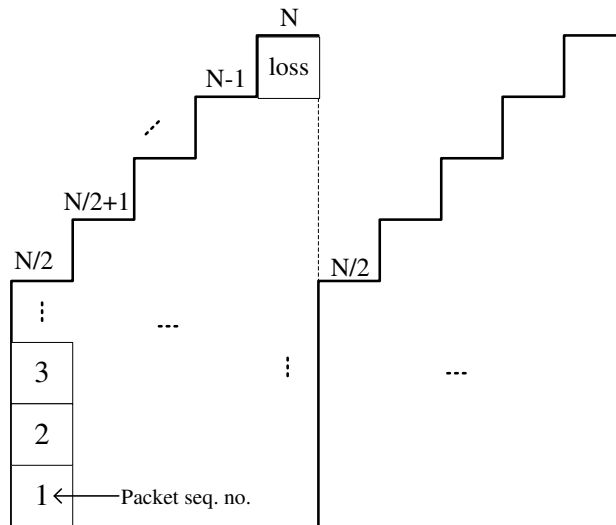
The following is assumed:

- Sender adjusts its window immediately when it receives an ACK/NACK.
- For given window size M , define a *successful round* as a period during which M packets are successfully transmitted and all their ACKs are received. Assume for simplicity that the window size increases by 1 after each successful round.

- Packets losses are independent and occur with probability $p(\ll 1)$.

Under these assumptions, we are interested in finding an average equilibrium window size in terms of p .

For simplicity, consider the following steady state. The packet loss occurs when the window size is N and the lost packet is the last packet sent (in the current window of N). Assume that the sender receives a NACK immediately (without delay). In this case, it will decrease its window size to $N/2$. Again, the window size increases up to N and decreases by half. This AIMD behavior continues to happen. A *cycle* is defined as the period from the time when the transmission with window size of $N/2$ is started to the time when a NACK is received (i.e., the period of time during which the window size increased from $N/2$ to N). The following figure depicts the evolution of the window size:



- Show that the number of packets transmitted in a single cycle is $\frac{3}{8}N^2 + \frac{3}{4}N$.
- Let α be the sequence number of the first lost packet. This is obviously a random variable as the packet loss occurs randomly. Show $E[\alpha] = \frac{1}{p}$ where $E[X]$ is the expectation on a random variable X .
- Show that the *average* window size in the steady state can be approximated by $\sqrt{\frac{3}{2p}}$.