A General Framework for Designing Approximation Schemes for Combinatorial Optimization Problems with Many Objectives Combined into One

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Outline

1. A general framework for designing approximation schemes
2. Applying the general framework to the following problems
   - Scheduling on unrelated parallel machines
   - Santa Claus problem
   - A resource allocation problem

The General Framework

Designing FPTAS for the following general problem:

$$\text{minimize } g(x) = h(f_1(x), \ldots, f_m(x)), \qquad x \in X.$$

- Each f_i non-negative function computable in polynomial time.
- h is a general norm for combining the functions together.
- m **is fixed.**

Scheduling on Unrelated Parallel Machines ($R||C_{\max}$)

- m machines, n jobs.
- Each job is to be processed on one of the machines.
- Processing time of job k on machine i is p_{ik}.
- **Objective**: Minimize the maximum processing time over all the machines.

The Santa Claus/Max-Min Fair Allocation Problem

- m agents, n resources.
- Each resource can be given to only one of the agents.
- The utility of agent i for resource k is p_{ik}.
- **Objective**: Maximize the minimum utility over all the agents.

Another Resource Allocation Problem

- m agents, n resources.
- More than one agent may need to use a resource, but a resource can be allocated to only one of the agents.
- The cost of accessing resource k for agent i, if it is allocated to agent j, is c_{ij}^k ($c_{ii}^k = 0$).
- **Objective**: Allocate the resources to agents so that the maximum cost over all the agents is minimized.

Our Results

- FPTAS for any norm used for combining the different objective functions.
- FPTAS for the Santa Claus problem and the resource allocation problem with fixed number of agents.
  - First FPTAS for these two problems.
- FPTAS for the $R||C_{max}$ problem with fixed number of machines having space complexity polynomial in m.
  - All previous FPTAS had space complexity exponential in m.

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Special Case: l_∞ norm

$$\text{minimize } g(x) = \max(f_1(x), \ldots, f_m(x)), \qquad x \in X.$$

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- X is a discrete set.
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π : specific instance of a problem.

$|\pi|$: the input problem size.

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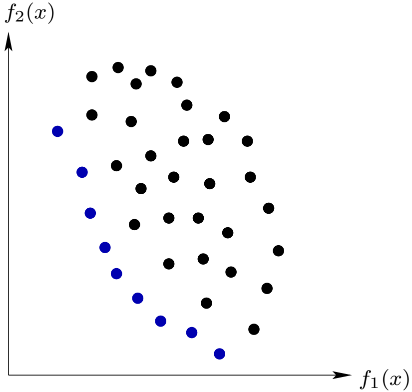
Pareto-optimal Front

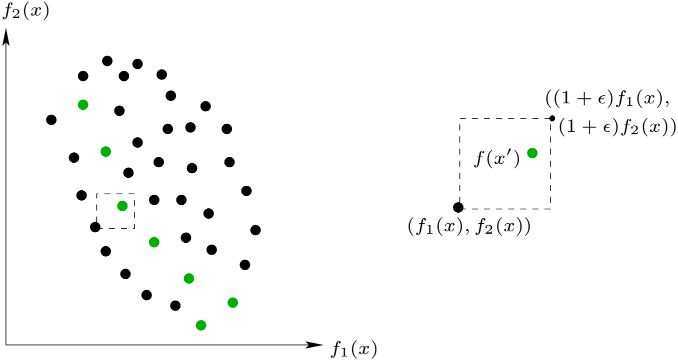
Pareto-optimal front ($P(\pi)$) is the set of all non-dominated solution points.



Approximate Pareto-optimal Front

Set of solutions $P_\epsilon(\pi)$ such that:

for all $x \in X$, there is $x' \in P_\epsilon(\pi)$ such that $f_i(x') \leq (1+\epsilon) f_i(x)$.



Why Approximate Pareto Front?

For many multi-objective combinatorial problems,

- It may not be tractable to compute $P(\pi)$.
- $P(\pi)$ may contain exponentially many points.

Example: two-objective shortest path problem (Hansen 1979).

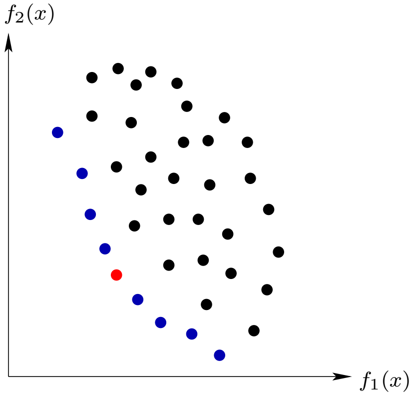
However:

For a **fixed number** m **of objectives**, there is a $P_\epsilon(\pi)$ whose cardinality is bounded by a polynomial in $|\pi|$ and $1/\epsilon$.

(Papadimitriou and Yannakakis 2000).

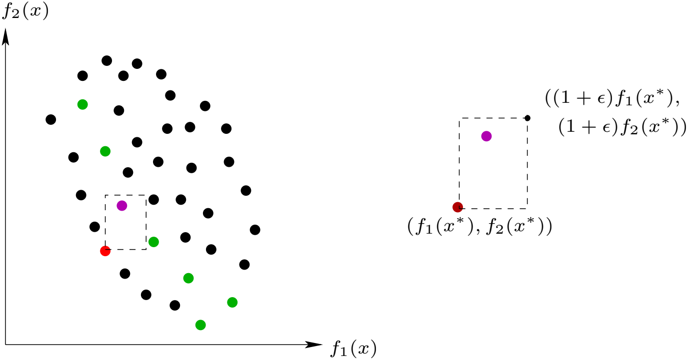
A Lemma

An optimal solution of the problem $\min g(x) = \max_{i=1,\ldots,m} f_i(x)$ lies on $P(\pi)$.



Another Lemma

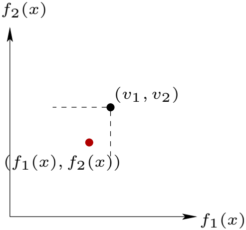
Let $\hat{x}$ be the solution in $P_\epsilon(\pi)$ that minimizes $g(x)$ over all the points in $P_\epsilon(\pi)$. Then $\hat{x}$ is a $(1+\epsilon)$-approximate solution of the optimization problem.



The Gap Theorem [PY'00]

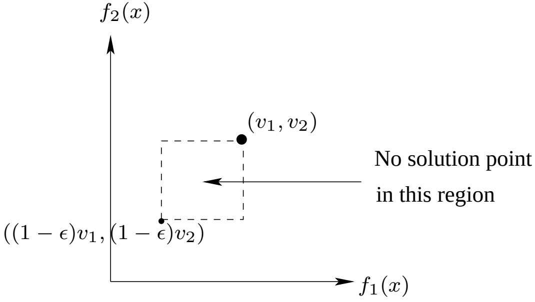
For a fixed m, it is possible to find a $P_\epsilon(\pi)$ in time polynomial in $|\pi|$ and $1/\epsilon$ *if and only if* the following "gap problem" can be solved in polynomial time.

Gap problem: Given an m vector of values $(v_1, \ldots, v_m)$, either (i) return a solution $x \in X$ such that $f_i(x) \leq v_i$ for all $i = 1, \ldots, m$, or .....



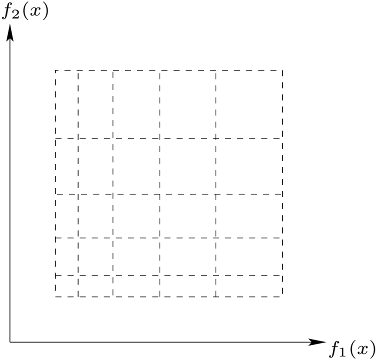
The Gap Theorem (contd.)

(ii) assert that there is no $x' \in X$ such that $f_i(x') \leq (1-\epsilon)v_i$ for all $i = 1, \ldots, m$.



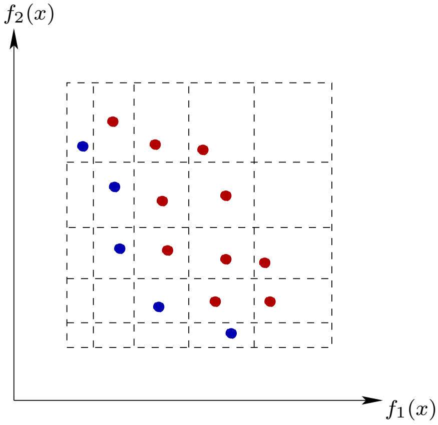
Proof for the 'If' Part

Divide the solution space into small hyperrectangles, such that in each dimension, the ratio of successive divisions is equal to $1 + \epsilon'$ (ϵ' depends on ϵ).



Proof for the 'If' Part (contd.)

For each corner point, solve the gap problem, and keep only the undominated set of solutions.



A Sufficient Condition for Solving Gap Problem

The setting:

- the functions $f_i(x)$ are linear functions
- X is a discrete set ($\{0,1\}^d$)

then the gap problem can be solved in polynomial time, *if* the following *exact problem* can be solved in *pseudo-polynomial* time:

Given a non-negative integer C and a vector $(c_1, \ldots, c_d) \in \mathbb{Z}_+$, does there exist a solution $x \in X$ such that

$$\sum_{i=1}^{d} c_i x_i = C?$$

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Proof Outline for Sufficiency of Given Condition

- Scale the coefficients in f_i to get f_i', in which the maximum magnitude of coefficients is $r = \lceil d/\epsilon \rceil$.
- Find x such that $f_i' \leq r$ for all i, or assert no such x exists.
- $(r+1)^m$ ways of having $f_i'(x) = b_i$, $b_i \leq r$
- Check for each case by finding an x such that $\sum_{i=1}^{m} M^{i-1} f_i'(x) = \sum_{i=1}^{m} M^{i-1} b_i$, M is maximum value f' can take.

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Exact Problem for the 3 Problems

For the Resource Allocation, the Santa Claus and the $R||C_{\max}$ problem, the exact problem is:

Given an integer C, does there exist a $0/1$-vector such that

$$\begin{aligned} \sum_{k=1}^{n}\sum_{j=1}^{m} c_{jk}x_{jk} &= C, \\ \text{subject to} \sum_{j=1}^{m} x_{jk} &= 1, \quad \text{for} \quad k = 1, \ldots, n. \end{aligned}$$

Solving the Exact Problem

- This exact problem can be formulated as a reachability problem in a directed graph.
- By using depth first search, and generating the graph "on the fly", we can solve the exact problem in *pseudo-polynomial time*, and in *polynomial space*.

Overview of the FPTAS

To summarize, the proposed FPTAS is:

- Divide the solution space into (polynomially many) smaller hyperrectangles.
- For each corner point, solve the gap problem and keep only the undominated set of solutions.
- Each gap problem can be solved by making (polynomially many) calls to the (pseudo-polynomial time) algorithm for exact problem. Each input to this algorithm has numbers whose magnitude is polynomial in input size.

Thank You!

Questions??