

Introduction

The Xilinx LogiCORE™ Block Memory Generator is an advanced memory constructor, generating area and performance optimized memories using embedded block RAM resources in Xilinx FPGAs. Available through the CORE Generator™ system, the core allows users to quickly create optimized memories to leverage the performance and features of block RAMs in Xilinx FPGAs.

Features

- Generates single-port RAM, simple dual-port RAM, true dual-port RAM, single-port ROM, dual-port ROM
- Performance up to 450 MHz
- Supports data widths from 1 to 1152 bits
- Supports memory depths from 8 to 9M words (limited only by memory resources on selected part)
- Supports configurable port aspect ratios for dual-port configurations
- Supports read-to-write aspect ratios in Virtex™-4 FPGAs
- Optimized algorithm for minimum block RAM resource utilization
- Configurable memory initialization
- Supports individual write enable per byte in Virtex-4 devices
- Optimized VHDL and Verilog behavioral models for fast simulation times
- Structural simulation models for precise simulation of memory behaviors
- Selectable operating mode per port: WRITE_FIRST, READ_FIRST, or NO_CHANGE

| LogiCORE Facts | |
|---------------------------------|--|
| Core Specifics | |
| Supported Device Family | Virtex-4, Virtex-II Pro, Virtex-II, Spartan™-3E, Spartan-3 |
| Package | All |
| Speed Grade | All |
| Performance | Varied, based on core parameters |
| Core Resources | |
| Block RAM | Varied, based on core parameters |
| DCM | None |
| BUFG | None |
| IOBs/RocketIO™ | None |
| PPC | None |
| IOB-FF/TBUFs | None |
| Provided with Core | |
| Documentation | Product Specification |
| Design File Formats | NGC netlist |
| Design Tool Requirements | |
| Xilinx Implementation Tools | ISE 8.1i |
| Supported Behavioral Models | VHDL, Verilog |
| Synthesis | XST 8.1i |
| Support | |
| Provided by Xilinx, Inc. | |

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Overview

The Block Memory Generator core uses the embedded Block Memory primitives available in Xilinx FPGAs, extending the functionality and capabilities of a single primitive to memories of arbitrary widths and depths. The sophisticated algorithms behind the Block Memory Generator core produce optimized solutions for a wide range of configurations, providing convenient access to memories optimized for your needs.

The Block Memory Generator has two fully independent ports. The A port has a write and a read interface, as does the B port. These four interfaces access a shared memory space, and can be configured uniquely for the Virtex-4 architecture. When not using all four interfaces, the user can select a simplified memory configuration such as Single-Port Memory or Simple Dual-Port Memory, which allows the core to more efficiently use the available resources.

For users familiar with the Single-Port Block Memory and Dual-Port Block Memory cores, the Block Memory Generator is not fully backwards compatible with these cores. The differences are discussed in "[Revision History](#)" on page 46.

Applications

The Block Memory Generator core is used to create customized memories to suit any application. Typical applications include:

- **Single-Port Memory:** Processor scratch RAM, very large look-up tables.
- **Dual-Port Memory:** Content addressable memories, FIFOs, multi-processor storage.
- **Read-Only Memory:** Program code storage, initialization ROM.

Feature Summary

Memory Types

The Block Memory Generator core uses embedded block RAM to generate the following types of memories:

- Single-port RAM
- Simple dual-port RAM
- True dual-port RAM
- Single-port ROM
- Dual-port ROM

For dual-port memories, each port operates completely independently. Operating mode, clock frequency, and optional pins are selectable per port.

Selectable Memory Algorithm

The core concatenates block RAM primitives according to one of the following algorithms:

- **Minimum Area Algorithm:** The memory is generated using the minimum number of block RAM primitives. Both data and parity bits are utilized.
- **Selectable Primitive Algorithm:** The memory is generated using only one type of block RAM primitive. For a complete list of primitives available for each device family, see the data sheet for that family.

Configurable Width and Depth

The Block Memory Generator can generate memory structures from 1 to 1152 bits wide, and at least eight locations deep. The maximum depth of the memory is limited only by the number of block RAM primitives in the target device.

Selectable Port Aspect Ratios

The core supports the same port aspect ratios as the block RAM primitives:

- In all supported device families, the A port width may differ from the B port width by a factor of 1, 2, 4, 8, 16, or 32.
- In Virtex-4 FPGA-based memories, the read width may differ from the write width by a factor of 1, 2, 4, 8, 16, or 32 for each port. The maximum ratio between any two of the data widths (DINA, DOUTA, DINB, and DOUTB) is 32:1.

Memory Initialization

The memory contents can be optionally initialized using a memory coefficient (COE) file or by using the default data option. A COE file can define the initial contents of each individual memory location, while the default data option defines the initial content of all locations.

Selectable Operating Mode per Port

The Block Memory Generator supports the block RAM primitive operating modes: WRITE FIRST, READ FIRST, and NO CHANGE. Each port may be assigned an operating mode.

Optional Output Registers

The Block Memory Generator provides two optional stages of output registering to increase memory performance. The core supports the Virtex-4 embedded block RAM registers as well as registers implemented in the FPGA fabric.

Optional Synchronous Set/Reset Pin

The core provides optional set/reset pins (SSRA and SSRB) pin per port that synchronously initialize the read output to a programmable value.

Optional Enable Pin

The core provides optional port enable pins (ENA and ENB) to control the operation of the memory. When deasserted, no read, write, or reset operations are performed on the respective port. If the enable pins are not used, it is assumed that the port is always enabled.

Optional Byte-Write Enable

In Virtex-4 FPGA-based memories, the Block Memory Generator core provides byte-write support for memory widths of 9-bit multiples.

Functional Description

The Block Memory Generator is used to build custom memory modules from block RAM primitives in Xilinx FPGAs. The core implements an optimal memory by arranging block RAM primitives based on user selections, automating the process of primitive instantiation and concatenation. Using the CORE

Generator graphical user interface (GUI), the user can configure the core and rapidly generate a highly optimized custom memory solution.

Memory Type

The Block Memory Generator core can create five types of memories: Single-port RAM, Simple dual-port RAM, True dual-port RAM, Single-port ROM, and Dual-port ROM. Figures 1 through 5 illustrate the signals available for each memory type. Optional pins are shown in italics.

The single-port ROM allows read access to the memory space through a single port, as shown in Figure 1.

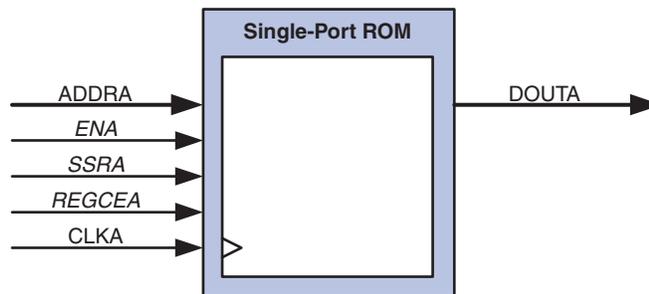


Figure 1: Single-port ROM

The dual-port ROM allows read access to the memory space through two ports, as shown in Figure 2.

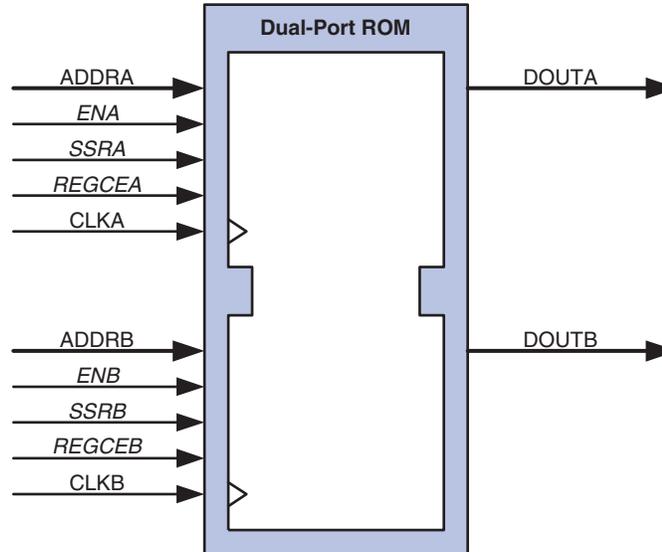


Figure 2: Dual-port ROM

The single-port RAM allows read and write access to the memory through a single port, as shown in Figure 3.

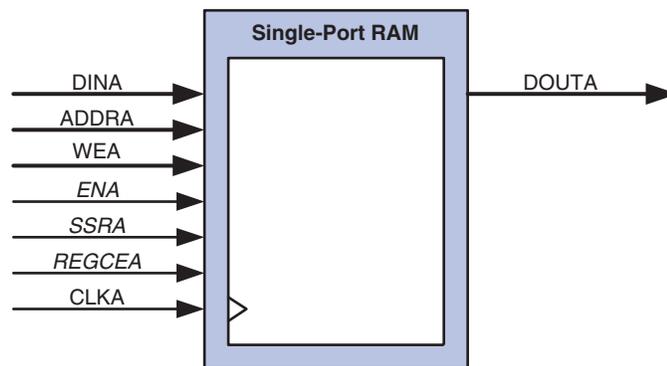


Figure 3: Single-port RAM

The simple dual-port RAM shown in Figure 4 provides two ports, A and B. Write access to the memory is allowed via port A, and read access is allowed via port B.

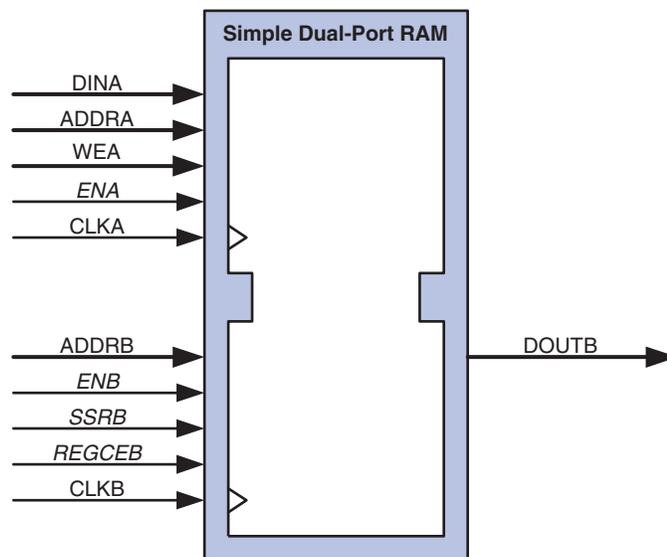


Figure 4: Simple Dual-port RAM

The true dual-port RAM shown in **Figure 5** provides two ports, A and B. Read and write accesses to the memory are allowed on either port.

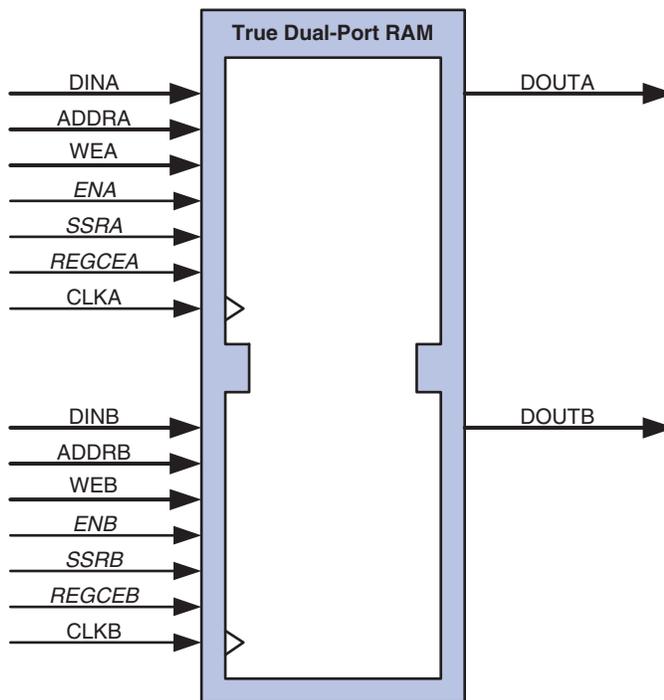


Figure 5: True Dual-port RAM

Selectable Memory Algorithm

The Block Memory Generator core arranges block RAM primitives according to one of two algorithms: the minimum area algorithm and the selectable primitive algorithm.

Minimum Area Algorithm

The minimum area algorithm provides a highly optimized solution, resulting in a minimum number of block RAM primitives used, while reducing output multiplexing. Figure 6 depicts two examples of memories built using the minimum area algorithm.

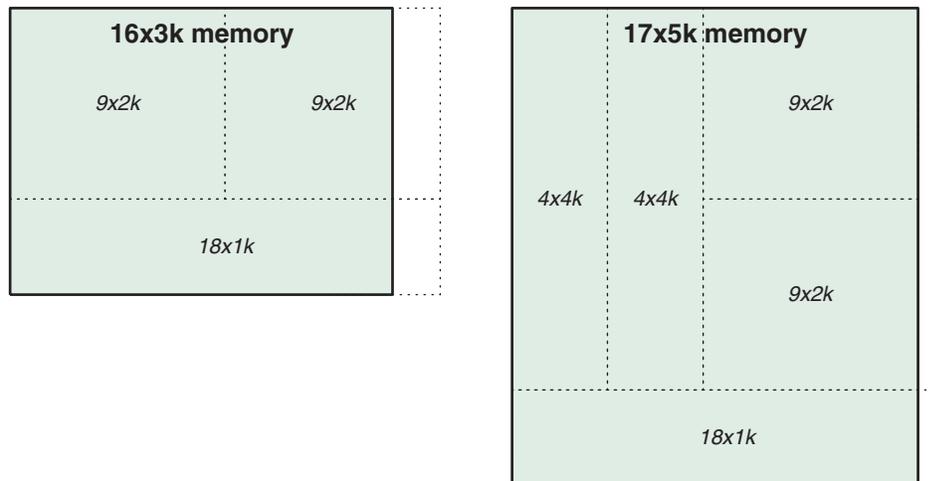


Figure 6: Examples of the Minimum Area Algorithm

In the first example, a 16x3k memory is implemented using three block RAMs. While it may have been possible to concatenate three 18x1k block RAMs in depth, this would require more output multiplexing. The minimum area algorithm maximizes performance in this way while maintaining minimum block RAM usage.

The second example, a 17x5k memory, further demonstrates how the algorithm can pack block RAMs efficiently to use the fewest resources while maximizing performance by reducing output multiplexing.

Selectable Primitive Algorithm

The selectable primitive algorithm allows the user to select a single block RAM primitive type. The core will build the memory by concatenating this single primitive type in width and depth. It is useful in

systems that require a fixed primitive type. Figure 7 depicts two 16x3k memories, one built using the 9x2k primitive type, the other built using the 4x4k primitive type.

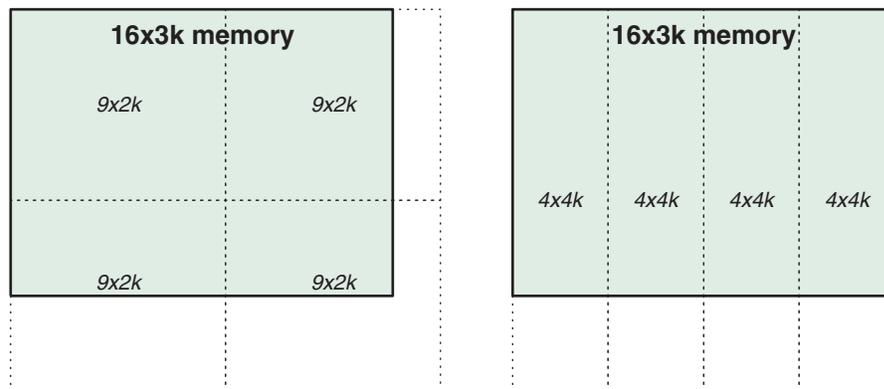


Figure 7: Examples of the Selectable Primitive Algorithms

Note that both implementations use four block RAMs, and that some of the resources utilized extend beyond the usable memory space. It is up to the user to decide which primitive type is best for their application.

The selectable primitive algorithm supports all of the block RAM primitive types: 1x16k, 2x8k, 4x4k, 9x2k, 18x1k, and 36x512. In Virtex-4 based memories, the 1x32k cascaded primitive type is also supported, which cascades two block RAMs primitives in depth using an embedded mux for faster operation.

When using data width aspect ratios, the primitive type dimensions are chosen with respect to the A port write width. Note that primitive selection may limit port aspect ratios as described in "[Aspect Ratio Limitations](#)" on page 12. When using the byte write feature, only the 9x2k, 18x1k, and 36x512 primitive types are available.

Selectable Width and Depth

The Block Memory Generator generates memories with widths from 1 to 1152 bits, and with depths of 8 or more words. The memory is built by concatenating block RAM primitives, and total memory size is limited only by the number of block RAMs on the target device.

Write operations to out of range addresses are guaranteed not to corrupt data in the memory, while read operations to out of range addresses will return invalid data. Note that the set/reset function should not be asserted while accessing an out of range address – it will also result in invalid data on the output.

Operating Mode

The operating mode determines which data word is presented on the output during a write operation. Each port is assigned an operating mode. The Block Memory Generator supports the three operating modes supported by the block RAM primitives. These operating modes are described in detail below. For more information, refer to the Block RAM section of the user guide for the specific device family.

- **Write First Mode:** In WRITE_FIRST mode, the input data is simultaneously written into memory and driven on the data output, as shown in Figure 8. This transparent mode offers the flexibility of

using the data output bus during a write operation on the same port.

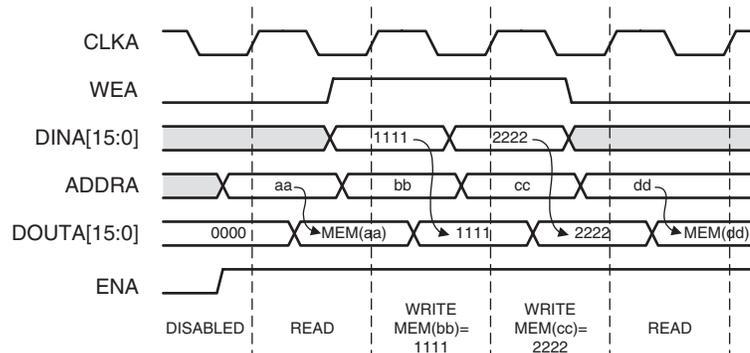


Figure 8: Write First Mode Example

Note that the WRITE_FIRST operation for Virtex-4 is affected by the optional read-to-write aspect ratio and byte-write features. For more details, see "Virtex-4 Write First Mode" on page 12.

- **Read First Mode:** In READ_FIRST mode, data previously stored at the write address appears on the data output, while the input data is being stored in memory. This read-before-write behavior is illustrated in Figure 9.

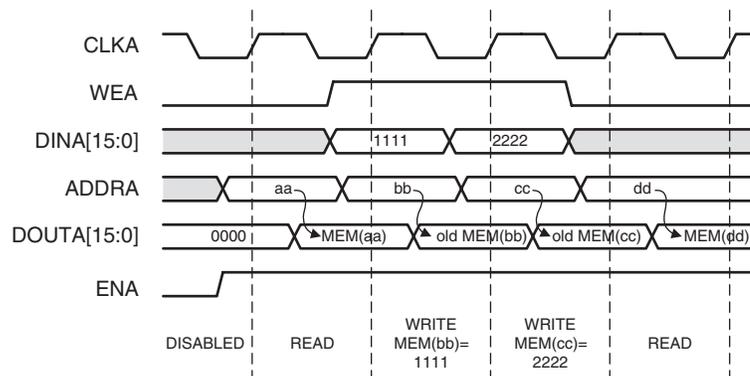


Figure 9: Read First Mode Example

- **No Change Mode:** In NO_CHANGE mode, the output latches remain unchanged during a write operation. As shown in Figure 10, the data output is still the previous read data and is unaffected

by a write operation on the same port.

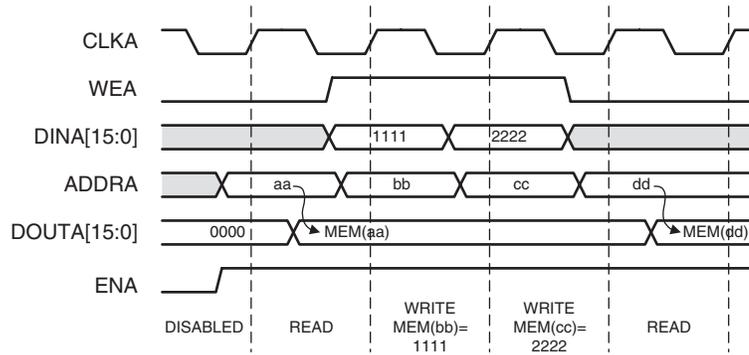


Figure 10: No Change Mode Example

Data Width Aspect Ratios

The Block Memory Generator supports data width aspect ratios. This allows the port A data width to be different than the port B data width, as described in Port Aspect Ratios. In Virtex-4 FPGA-based memories, all four data busses (DINA, DOUTA, DINB, and DOUTB) can have different widths, as described in "Virtex-4 Read-to-Write Aspect Ratios" on page 11.

The limitations of the data width aspect ratio feature (some of which are imposed by other optional features) are described in "Aspect Ratio Limitations" on page 12. The CORE Generator GUI ensures only valid aspect ratios as selected.

Port Aspect Ratios

The Block Memory Generator supports port aspect ratios of 1:32, 1:16, 1:8, 1:4, 1:2, 1:1, 2:1, 4:1, 8:1, 16:1, and 32:1. The port A data width can be up to 32 times larger than the port B data width, or vice versa. The smaller data words are arranged in little-endian format, as illustrated in Figure 11.

Port Aspect Ratio Example

Consider a true dual-port RAM of 32x2048, which is the A port width and depth. From the perspective of an 8-bit B port, the depth would be 8192. The ADDRA bus is 11 bits, while the ADDR B bus is 13 bits. The data is stored little-endian, as shown in Figure 11. Note that A_n is the data word at address n, with respect to the A port. B_n is the data word at address n with respect to the B port. A_0 is comprised of $B_3, B_2, B_1,$ and B_0 .

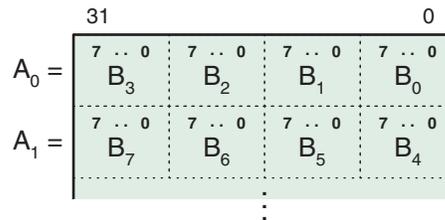


Figure 11: Port Aspect Ratio Example Memory Map

Virtex-4 Read-to-Write Aspect Ratios

When implementing RAMs targeting Virtex-4 FPGAs, the Block Memory Generator allows read and write aspect ratios. On each port, the read to write data width ratio can be 1:32, 1:16, 1:8, 1:4, 1:2, 1:1, 2:1, 4:1, 8:1, 16:1, or 32:1.

For true dual-port RAMs, the four data busses, DINA, DOUTA, DINB, and DOUTB, can have different widths. For single-port RAMs, DINA and DOUTA widths can be independent. The maximum ratio between any two data buses is 32:1. The widest data bus can be no larger than 1152 bits.

If the read and write data widths are different, the memory depth is different with respect to read and write accesses. The depth ratio is the inverse of the width ratio. The address bus must be large enough to address the deeper of the two depths, since it is shared by read and write accesses. For the shallower, the least significant bits of the address bus are ignored. The smaller data words are arranged in little-endian format, as illustrated in Figure 12.

Virtex-4 Read-to-Write Aspect Ratio Example

Consider a true dual-port RAM of 64x512, which is the A port write width and depth. Table 1 lists the four data port widths and their respective depths for this example.

Table 1: Read-to-Write Aspect Ratio Example Ports

| Interface | Data Width | Memory Depth |
|--------------|------------|--------------|
| Port A Write | 64 | 512 |
| Port A Read | 16 | 2048 |
| Port B Write | 256 | 128 |
| Port B Read | 32 | 1024 |

The ADDRA width is determined by the larger of the A port depths (2048). Therefore, ADDRA is 11 bits wide. On port A, read operations utilize the entire ADDRA bus, while write operations ignore the least significant 2 bits.

In the same way, the ADDR B width is determined by the larger of the B port depths (1024). Therefore, ADDR B is 10 bits wide. On port B, read operations utilize the entire ADDR B bus, while write operations ignore the least significant 3 bits.

The memory map in Figure 12 shows how port B write words are related to port A write words, in a little-endian arrangement. Note that AW_n is the write data word at address 'n' with respect to port A, while BW_n is the write data word at address 'n' with respect to port B.

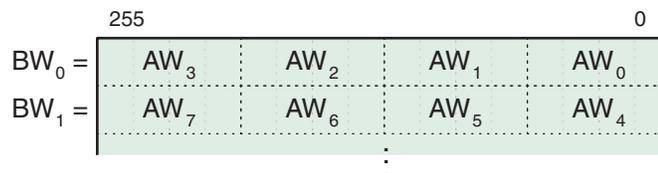


Figure 12: Read-to-Write Aspect Ratio Example Memory Map

BW_0 is made up of AW_3 , AW_2 , AW_1 , and AW_0 . In the same way, BR_0 is made up of AR_1 and AR_0 , and AW_0 is made up of BR_1 and BR_0 .

In the above example, the largest data width ratio is port B write words (256 bits) to port A read words (16 bits). This ratio is 16:1.

Aspect Ratio Limitations

In general, no port data width may be wider than 1152 bits, and no two data widths can have a ratio greater than 32:1. However, some optional features further limit data width aspect ratios:

- **Byte-writes:** When using byte-writes, no two data widths can have a ratio greater than 4:1.
- **Selectable primitive algorithm:** When using the selectable primitive algorithm with an N-bit wide primitive, aspect ratios are limited to 32:N and 1:N from the port A write width. For example, using the 4x4k primitive type, the other ports may be no more than 8 times (32:4) larger than port A write width and no less than 4 times (1:4) smaller.

Virtex-4 Byte-Writes

The Block Memory Generator provides byte-write support in Virtex-4 devices. Byte-writes are available for memories of 9-bit multiples — parity bits are included.

When byte-writes are enabled, the WE [A | B] bus is N bits wide, where N is the number of bytes in DIN [A | B]. The most significant bit in the write enable bus corresponds to the most significant byte in the input word. Bytes will be stored in memory only if the corresponding bit in the write enable bus is asserted during the write operation.

To remove parity and use 8-bit bytes, the user may tie off the 9th bit of each byte in the data words. The byte-write feature may be used in conjunction with the Virtex-4 data width aspect ratios. However, it may not be used with the NO_CHANGE operating mode. The byte-write feature also affects the operation of WRITE_FIRST mode, as described in "[Virtex-4 Write First Mode](#)" on page 12.

Byte-Write Example

Consider a single-port RAM with a data width of 24 bits, or 3 bytes with parity bits tied off. The write enable bus, WEA, consists of 3 bits. [Figure 13](#) illustrates the use of byte-writes, and shows the contents of the RAM at address 0. Assume all memory locations are initialized to 0.

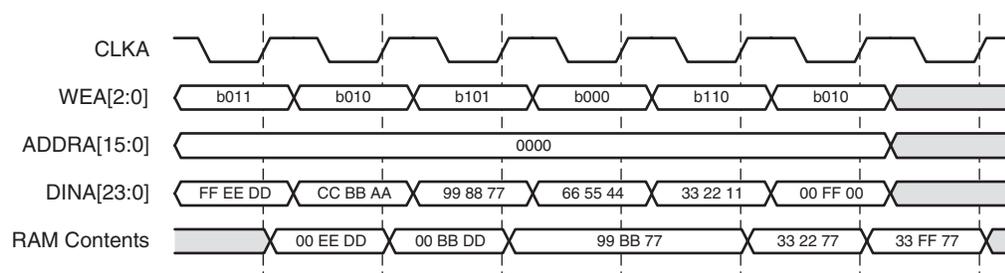


Figure 13: Byte-write Example

Virtex-4 Write First Mode

The byte-write feature and the read-to-write data width aspect ratio have special consideration when using WRITE_FIRST mode in Virtex-4 FPGAs. In general, when performing a write in WRITE_FIRST mode, the concurrent read operation shows the newly written data on the output of the core.

When writing a full data word (all bits of WE are asserted), data on the output will be valid. If the read width is greater than the write width, the output word shows newly written data along with old data which the write did not affect.

When writing a partial data word (some bits of WE are asserted), data on the output will not be valid in WRITE_FIRST mode. Subsequent reads will return valid data, with byte-writes applied appropriately.

Virtex-4 Write First Mode Examples

The waveform in Figure 14 illustrates when the read width (32 bits) is greater than the write width (16 bits) in WRITE_FIRST mode. A full word write and a partial word write are shown. Assume all memory locations are initialized to 0. Since the read data width is twice as wide, the relative memory space will be half as deep. Thus, read operations ignore the least significant bit of ADDRA.

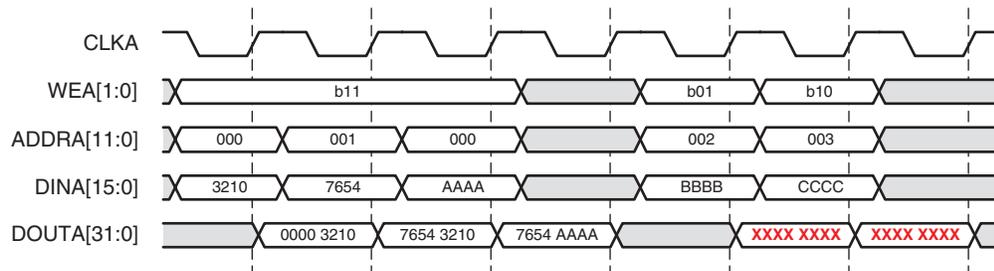


Figure 14: Virtex-4 Write-First Mode Example 1

The waveform in Figure 15 illustrates when the read width (18 bits) is equal to the write width (18 bits) in WRITE_FIRST mode. A full word write and a partial word write are shown. Note that the 18-bit signals are shown as two separate 9-bit bytes.

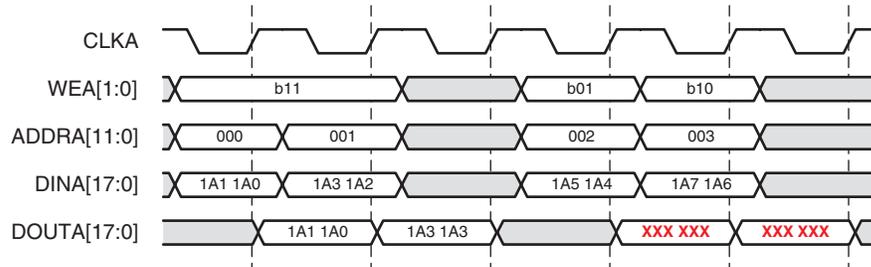


Figure 15: Virtex-4 Write-first Mode Example 2

The waveform in Figure 16 illustrates when the read width (16 bits) is less than the write width (32 bits) in WRITE_FIRST mode. A full word write and a partial word write are shown. Since the write data

width is twice as wide, the relative memory space will be half as deep. Thus, write operations ignore the least significant bit of ADDRA.

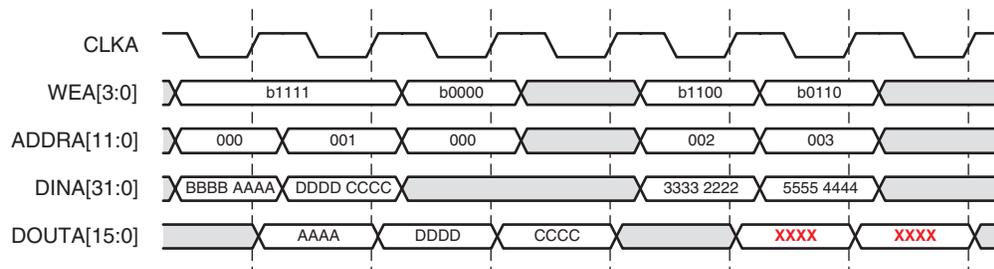


Figure 16: Virtex-4 Write-first Mode Example 3

In Figures 14, 15, and 16, 'X' indicates undefined data. Note that the structural simulation model generates these 'X's' purposefully to match the core behavior, indicating that data is undefined. The behavioral model never generates 'X's' though output data is invalid and should be ignored in these cases. For more information, see "Simulation Models" on page 18.

Collision Behavior

The Block Memory Generator core supports dual-port RAM implementations. Each port is equivalent and independent, yet they access the same memory space. In such an arrangement, is it possible to have data collisions. The ramifications of this behavior are described for both asynchronous and synchronous clocks below.

Collisions and Asynchronous Clocks

Using asynchronous clocks, when one port writes data to a memory location, the other port must not read or write that location for a specified amount of time. This clock-to-clock setup time is defined in the device data sheet, along with other block RAM switching characteristics.

Collisions and Synchronous Clocks

Synchronous clocks cause a number of special case collision scenarios, described below.

Synchronous Write-Write Collisions

A write-write collision occurs if both ports attempt to write to the same location in memory. The resulting contents of the memory location are unknown. Note that write-write collisions affect memory content, as opposed to write-read collisions which only affect data output.

Synchronous Write-Write Collisions and Byte-Writes

When using byte-writes, memory contents are not corrupted when separate bytes are written in the same data word. RAM contents are corrupted only when both ports attempt to write the same byte. Figure 17 illustrates this case. Assume ADDR_A = ADDR_B = 0.

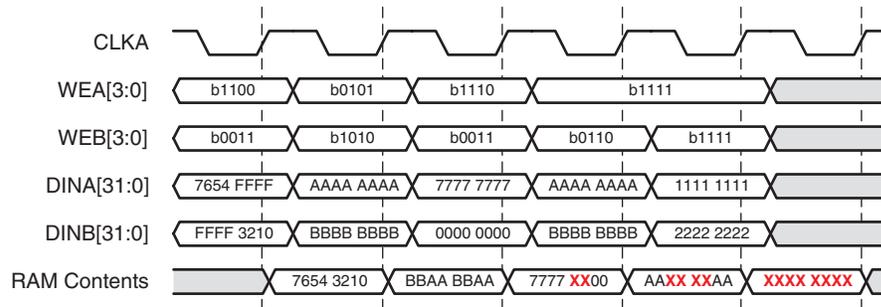


Figure 17: Write-Write Collision Example

Synchronous Write-Read Collisions

A synchronous write-read collision may occur if a port attempts to write a memory location and the other port reads the same location. While memory contents are not corrupted in write-read collisions, the validity of the output data depends on the write port operating mode.

- If the write port is in READ_FIRST mode, the other port can reliably read the old memory contents.
- If the write port is in WRITE_FIRST or NO_CHANGE mode, data on the output of the read port is invalid.
- In the case of byte-writes, only bytes which are updated will be invalid on the read port output.

Figure 18 illustrates write-read collisions and the effects of byte-writes. DOUTB is shown for when port A is in WRITE_FIRST mode and READ_FIRST mode. Assume ADDR_A = ADDR_B = 0, port B is always reading, and all memory locations are initialized to 0. Recall that the RAM contents are never corrupted in write-read collisions.

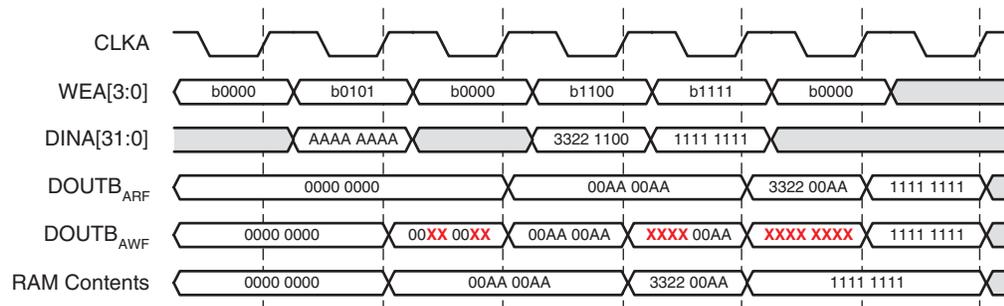


Figure 18: Write-read Collision Example

Optional Output Registers

The Block Memory Generator allows optional output registers to improve the performance of the core. The user may choose to include register stages at two places: at the output of the block RAM primitives and at the output of the core.

Registers at the output of the block RAM primitives reduce the impact of the clock to out delay of the primitives. Registers at the output of the core isolate the delay from the block RAMs through the output multiplexers. Note that each optional register stage used adds an additional clock cycle of latency to the read operation.

For Virtex-4 FPGAs, the primitive output register stage may be implemented using the Virtex-4 embedded block RAM registers (requiring no further FPGA resources). All other register stages are implemented in FPGA fabric. **Figure 19** shows an example of a Virtex-4 based memory that has been configured using both output register stages.

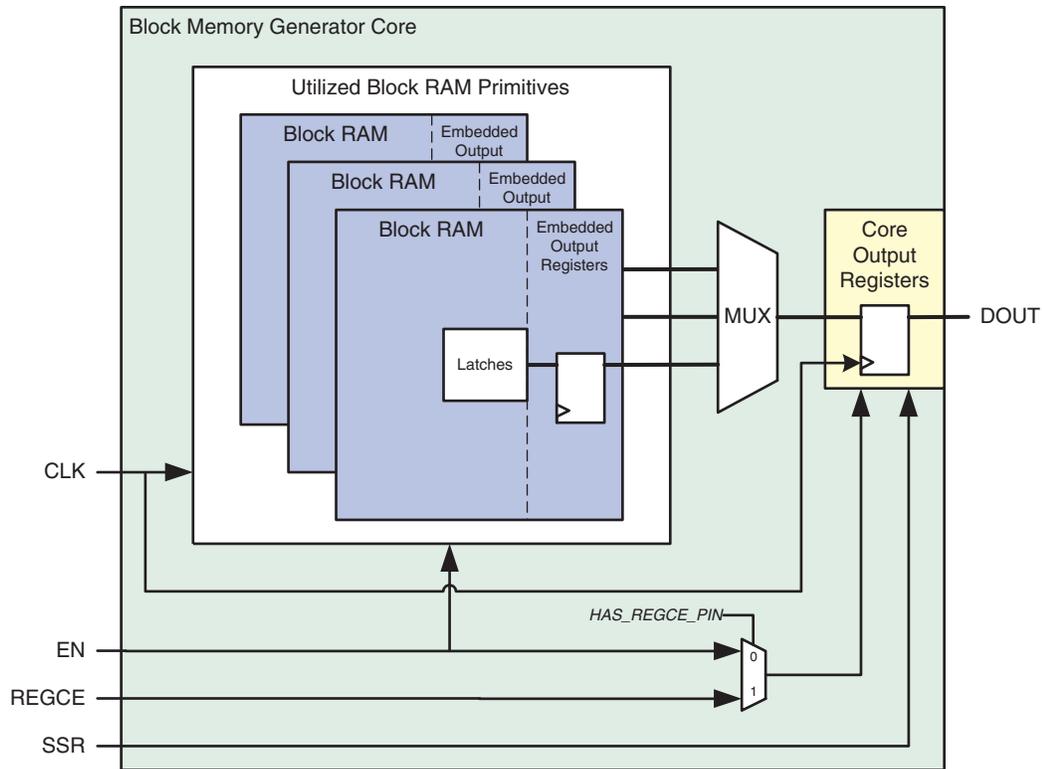


Figure 19: Virtex-4 Block Memory Generated With Both Embedded and Core Output Registers Enabled

For a complete description of the output options supported, see it.

Optional Register Clock Enable Pins

The optional output registers are enabled by the EN signal by default. However, the optional register clock enable pins (REGCEA and REGCEB) allow the user to control the clock enable pin for the last register in the memory output block.

Thus, the data output from the core can be controlled independent of the flow of data through the rest of the core. When using the REGCE pin, the last output register operates independently of the EN signal. For a complete description of the output options supported, see ["Output Register Configurations" on page 37](#).

Optional Synchronous Set/Reset Pins

The synchronous set/reset pins (SSRA and SSRB) control the reset operation of the last register in the output stage. For memories with no output registers, the reset pins control the memory output latches.

When SSR and EN are asserted on a given port, the data on the output of that port is driven to the reset value defined in the CORE Generator GUI.

For Virtex-4 FPGAs, if the option to use the synchronous set/reset pin is selected in conjunction with memory primitive registers and without core output registers, then the Virtex-4 embedded block RAM registers may not be utilized. For more information, refer to "Output Register Configurations" on page 37.

Memory Output Flow Control Examples

The combination of the enable (EN), reset (SSR), and register enable (REGCE) pins allow a wide range of data flows in the output stage. Figure 20 and Figure 21 are examples on how this can be accomplished. Keep in mind that the SSR and REGCE pins apply only to the last register stage.

Figure 20 depicts how SSR can be used to control the data output to allow only intended data through. Assume that both output registers are used, the port A reset value is 0xFFFF, and that EN and REGCE are always asserted. The data on the block RAM memory latch is labeled LATCH, while the output of the block RAM embedded register is labeled REG1. The output of the last register is the output of the core, DOUT.

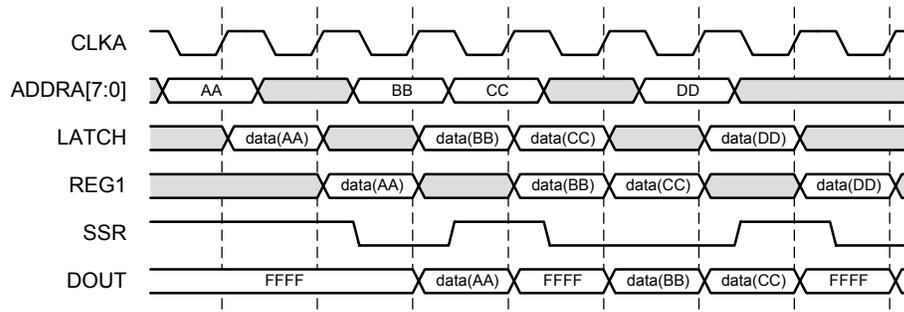


Figure 20: Flow Control Using SSR

Figure 21 depicts how REGCE can be used to latch the data output to allow only intended data through. Assume that only the memory primitive registers are used, and that EN is always asserted and SSR is

always deasserted. The data on the block RAM memory latch is labeled latch, while the output of the last register, the block RAM embedded register, is the core output, DOUT.

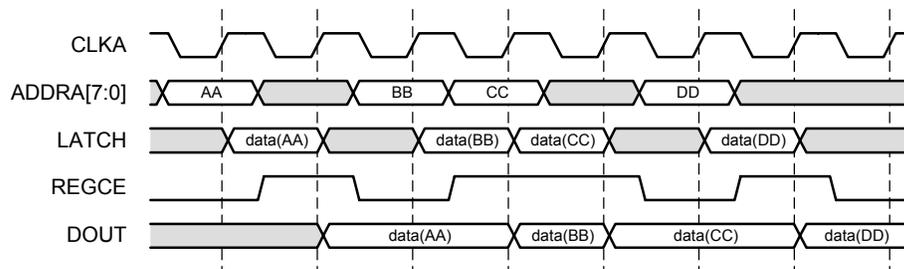


Figure 21: Flow Control Using REGCE

Simulation Models

The Block Memory Generator core provides two types of functional simulation models to its users:

- Behavioral Simulation Models (VHDL and Verilog)
- Structural/Unisim based Simulation Models (VHDL and Verilog)

The behavioral simulation models is a simplified model of the core while the structural (unisim) based simulation models is an accurate modeling of the internals of the core. The behavioral simulation models is written purely in RTL and simulates faster than the structural based simulation models and is ideal for debugging. Moreover, the memory is modeled in a 2-dimensional array which makes it easier to probe contents of the memory.

The structural simulation model uses primitive instantiations to more accurately model the behavior of the core. Collision detection and 'X' generation are better modeled using the structural simulation model. However, simulation time is longer and debug may be more difficult.

The "Simulation Files" options in the CORE Generator Project Options determines which type of functional simulation models are generated.

Table 2 lists the differences between the two functional simulation models:

Table 2: Differences Between Two Simulation Models

| | Behavioral Models | Structural (Unisim) based Models |
|--------------------------------|---|---|
| When core output is undefined. | Will never generate 'X'. | Will generate 'X' to match core. |
| Out-of-range address access | Optionally flags a warning message. | Will generate 'X'. |
| Collision behavior | Will not generate 'X' on output and flag a warning message. | Will generate 'X' to match core. |
| Byte-write collision behavior | Will flag all byte-write collisions. | Will not flag collisions if byte-writes do not overlap. |

Signal List

Table 3 provides a description of the Block Memory Generator core signals. The widths of the data ports (DINA, DOUTA, DINB, and DOUTB) are selected by the user in the CORE Generator GUI. The address port (ADDRA and ADDR B) widths are determined by the memory depth with respect to each port, as selected by the user in the GUI. The write enable ports (WEA and WEB) are busses of width 1 when byte-writes are disabled. When byte-writes are enabled, WEA and WEB widths depend on the byte size and write data widths selected in the GUI.

Table 3: Core Signal Pinout

| Name | Direction | Description |
|--------|-----------|---|
| CLKA | Input | Port A clock: Port A operations are synchronous to this clock. For synchronous operation, this must be driven by the same signal as CLKB. |
| ADDRA | Input | Port A address: Addresses the memory space for Port A read and write operations. Available in all configurations. |
| DINA | Input | Port A Data Input: Data input to be written into the memory via Port A. Available in all RAM configurations. |
| DOUTA | Output | Port A Data Output: Data output from read operations via Port A. Available in all configurations except simple dual-port RAM. |
| ENA | Input | Port A Clock Enable: Enables read, write, and reset operations via Port A. Optional in all configurations. |
| WEA | Input | Port A Write Enable: Enables write operations via Port A. Available in all RAM configurations. |
| SSRA | Input | Port A Synchronous Set/Reset: Resets the Port A memory output latch or output register. Optional in all configurations. |
| REGCEA | Input | Port A Register Enable: Enables the last output register. Optional in all configurations with output registers. |
| CLKB | Input | Port B Clock: Port B operations are synchronous to this clock. Available in dual-port configurations. For synchronous operation, this must be driven by the same signal as CLKA. |
| ADDRB | Input | Port B address: Addresses the memory space for Port B read and write operations. Available in dual-port configurations. |
| DINB | Input | Port B Data Input: Data input to be written into the memory via Port B. Available in true dual-port RAM configurations. |
| DOUTB | Output | Port B Data Output: Data output from read operations via Port B. Available in dual-port configurations. |
| ENB | Input | Port B Clock Enable: Enables read, write, and reset operations via Port B. Optional in dual-port configurations. |
| WEB | Input | Port B Write Enable: Enables write operations via Port B. Available in dual-port RAM configurations. |
| SSRB | Input | Port B Synchronous Set/Reset: Resets the Port B memory output latch or output register. Optional in dual-port configurations. |
| REGCEB | Input | Port B Register Enable: Enables the last output register. Optional in dual-port configurations with output registers. |

Generating the Core

The Block Memory Generator can be found in the CORE Generator GUI View by Function pane: Under Memories & Storage Elements > RAMs & ROMs. This section describes the options available in the Block Memory Generator GUI.

CORE Generator Parameter Screens

The Block Memory Generator GUI includes five screens:

- [Block Memory Generator Main Screen](#)
- [Port A Options Screen](#)
- [Port B Options Screen](#)
- [Output Registers and Memory Initialization Screen](#)
- [Simulation Model Options and Information Screen](#)

In addition, all the screens share common tabs and buttons to provide information about the core and to navigate the Block Memory Generator GUI.

Block Memory Generator Main Screen

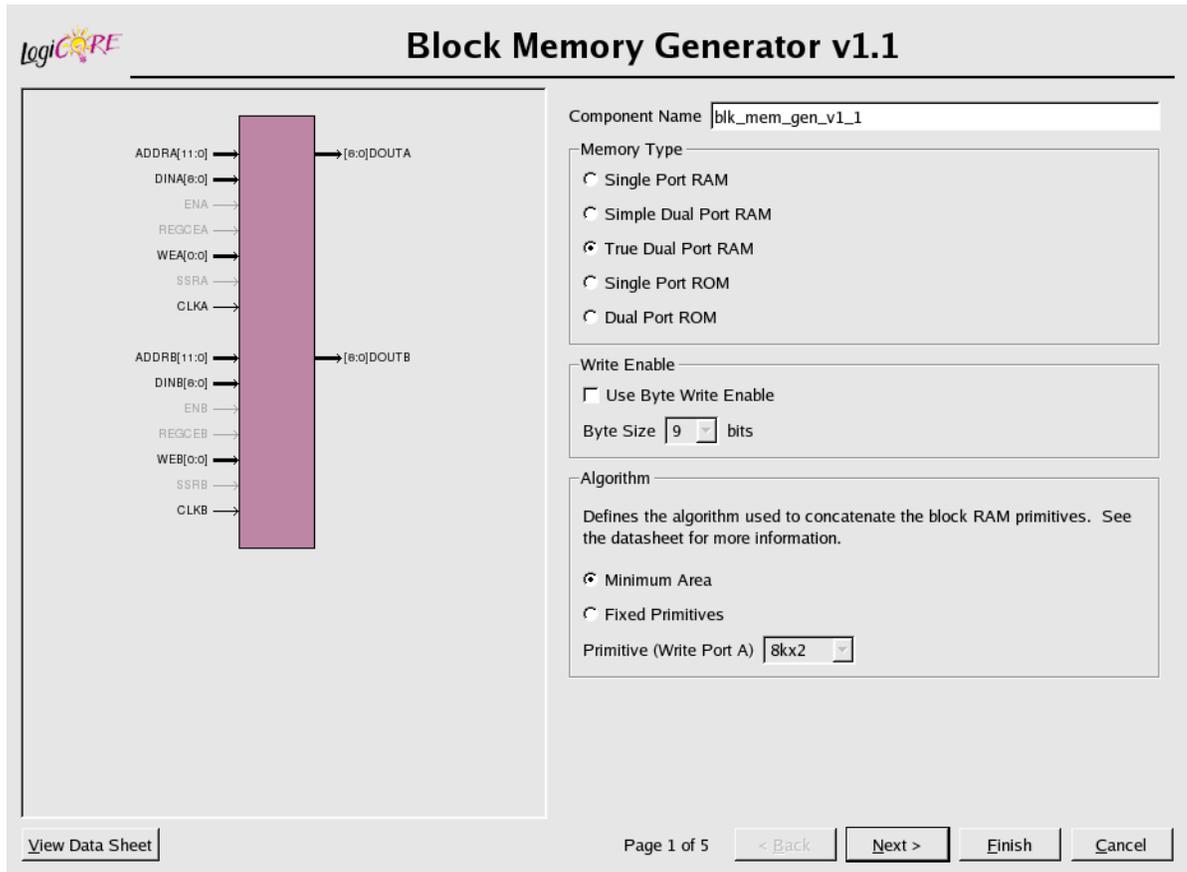


Figure 22: Block Memory Generator Main Screen

Component Name

The base name of the output files generated for the core. Names must begin with a letter and be composed of any of the following characters: a to z, 0 to 9, and “_”. Names can not be Verilog or VHDL reserved words.

Memory Type

Select the type of memory to be generated.

- Single-port RAM
- Simple dual-port RAM
- True dual-port RAM
- Single-port ROM
- Dual-port ROM configurations.

Write Enable

When targeting Virtex-4, select whether to use the byte-write enable feature. Byte size is 9 bits, though it may be optional in future versions.

Algorithm

Select the algorithm used to implement the memory.

- **Minimum Area Algorithm:** Generates a core using the least number of primitives.
- **Selectable Primitive Algorithm:** Generates a core that concatenates a single primitive type to implement the memory. Choose which primitive type to use in the drop down box.

Port A Options Screen

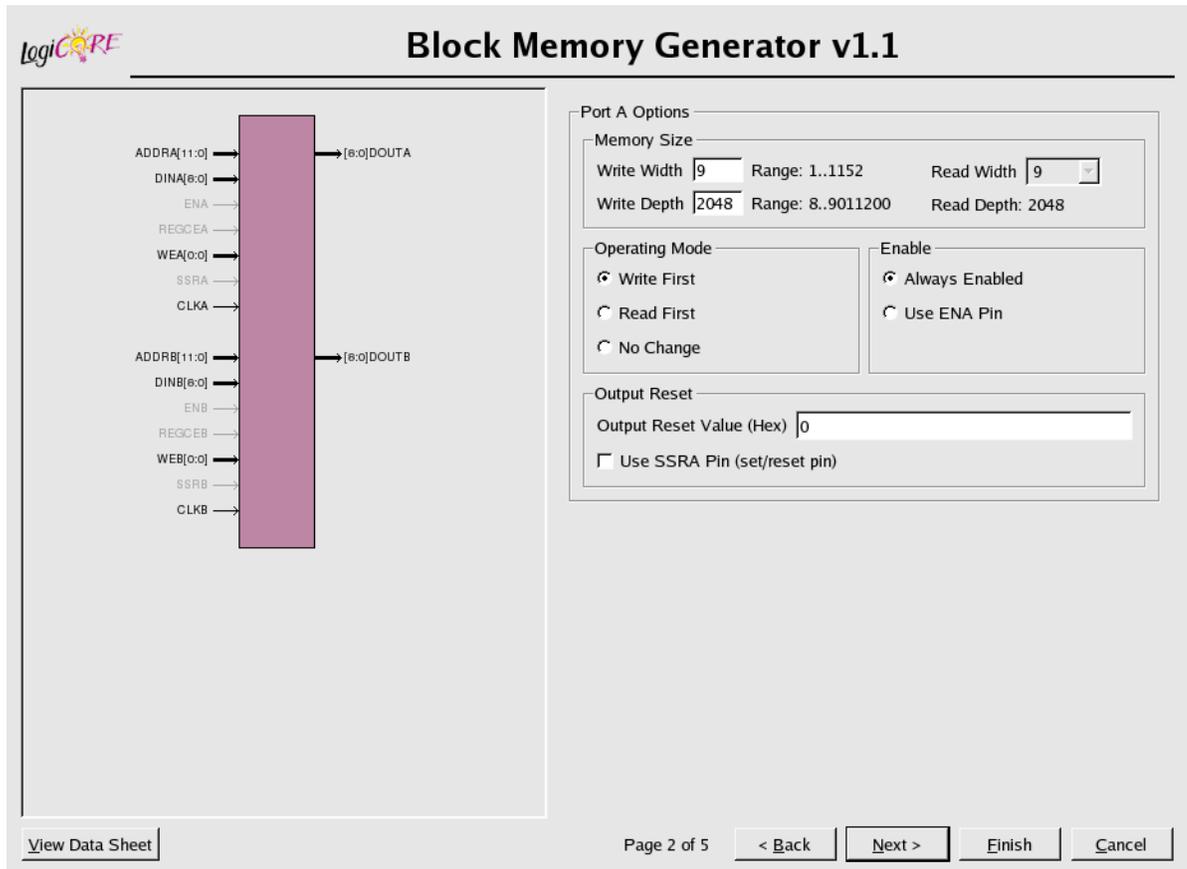


Figure 23: Port A Options

Memory Size (Port A)

Specify the port A write width and depth. Select the port A read width from the drop-down menu of valid choices. The read depth is calculated automatically.

Operating Mode (Port A)

Specify the port A operating mode.

- READ_FIRST
- WRITE_FIRST
- NO_CHANGE

Enable (Port A)

Select the enable type:

- Always enabled (no ENA pin available)
- Use ENA pin

Output Reset (Port A)

Specify the reset value of the memory output latch and output registers. These values are with respect to the read port widths. Choose whether a set/reset pin (SSRA) is needed.

Port B Options Screen

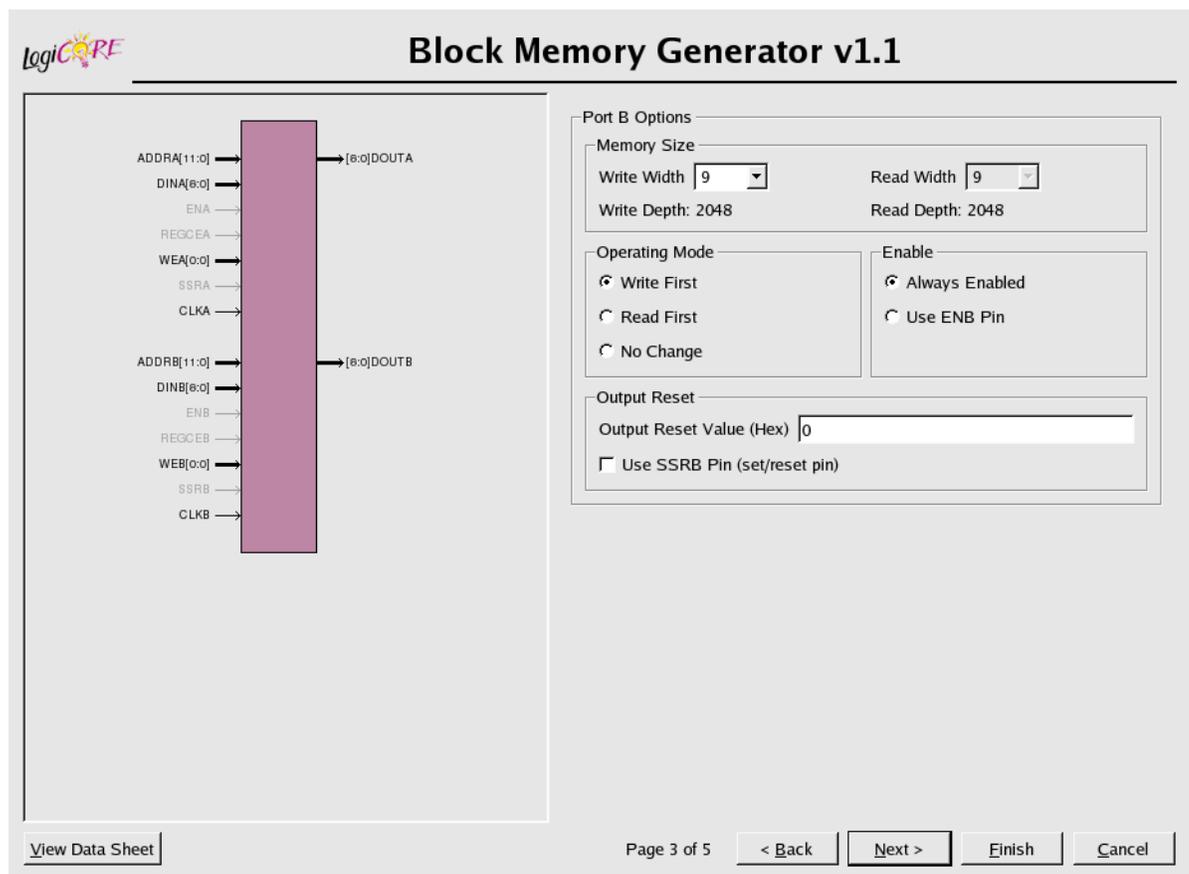


Figure 24: Port B Options

Memory Size (Port B)

Select the port B write and read widths from the drop-down menu of valid choices. The read depth is calculated automatically.

Operating Mode (Port B)

Specify the port B write mode.

- READ_FIRST
- WRITE_FIRST

- NO_CHANGE

Enable (Port B)

Select the enable type:

- Always enabled (no ENB pin available)
- Use ENB pin

Output Reset (Port B)

Specify the reset value of the memory output latch and output registers. These values are with respect to the read port widths. Choose whether a set/reset pin (SSRB) is needed.

Output Registers and Memory Initialization Screen

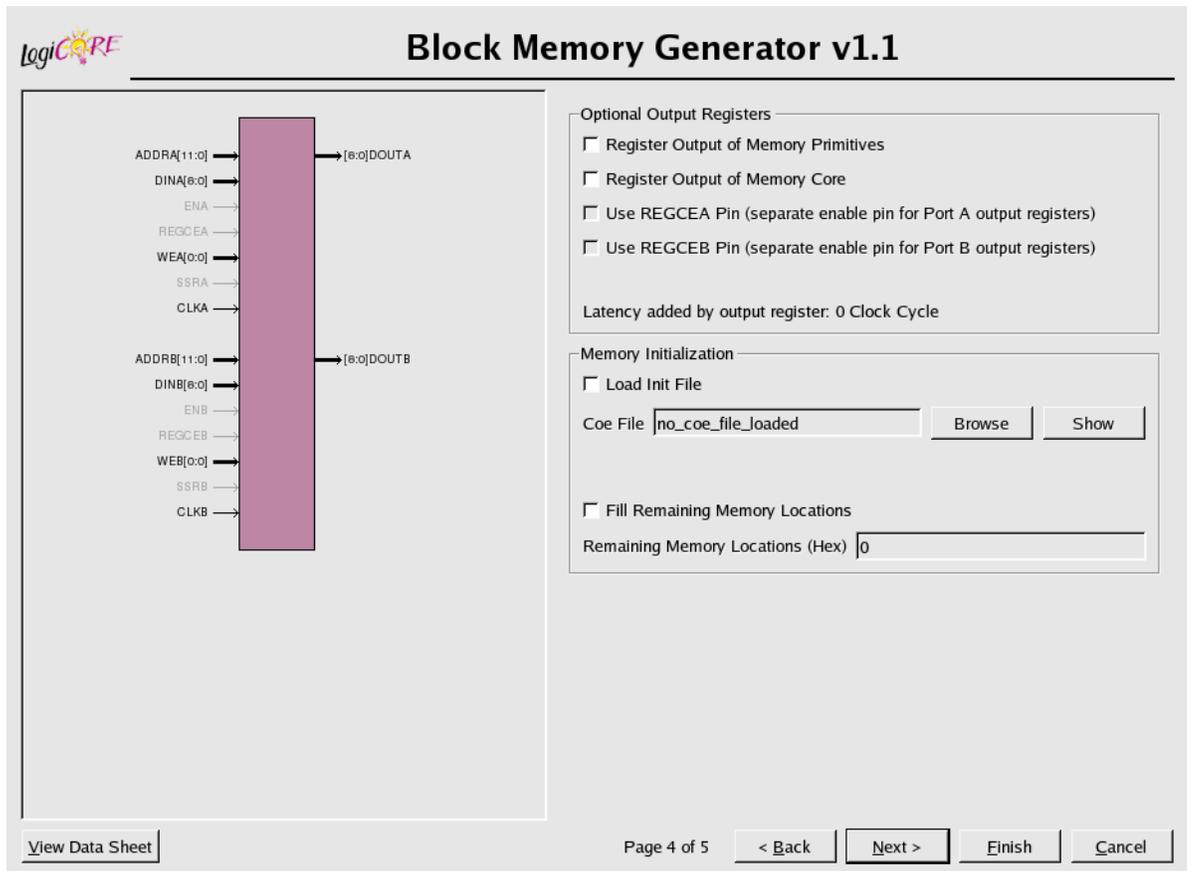


Figure 25: Output Registers and Memory Initialization Screen

Optional Output Registers

Select which output register stages to include:

- **Register Output of Memory Primitives.** Select to insert output register after the memory primitives. When targeting Virtex-4 FPGAs, the embedded output registers in the block RAM primitives will be used if the user chooses to registers the output of the memory primitives. For other architectures, the registers in the FPGA slices will be used.
- **Register Output of Memory Core.** Select to insert register output of the memory core. The registers

in the FPGA slices will be used.

Select whether to have separate enable pins for the last register stage if output registers are selected.

Memory Initialization

Select whether to initialize the memory contents using a COE file, and whether to initialize the remaining memory contents with a default value. When using asymmetric port widths or data widths, the COE file and the default value are with respect to the port A write width.

Simulation Model Options and Information Screen

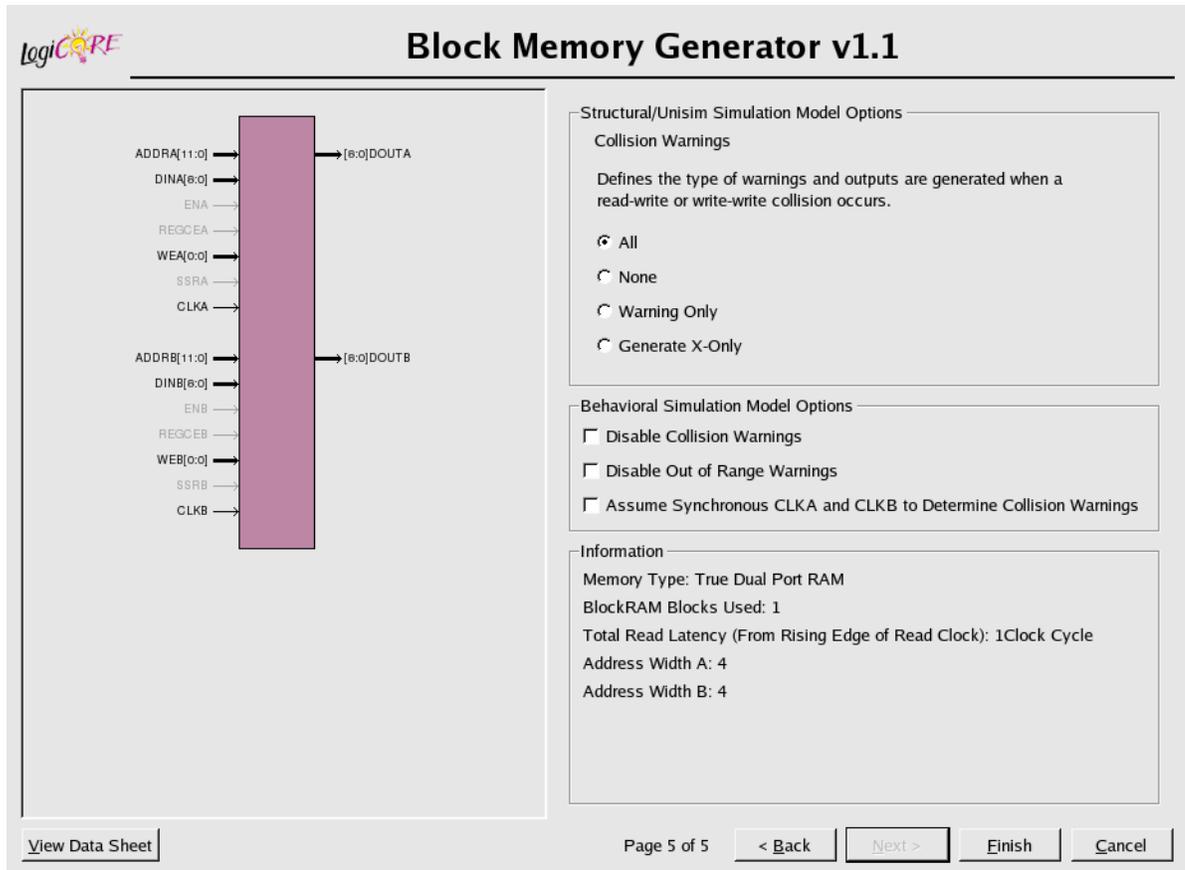


Figure 26: Simulation Model Options and Information Screen

Structural/UNISIM Simulation Model Options

Select the type of warning messages and outputs generated by the structural simulation model in the event of collisions.

Behavioral Simulation Model Options

Select the type of warning messages generated by the behavioral simulation model. Select whether the model should assume synchronous clocks for collision warnings.

Information Panel

This pane displays an informational summary of the selected core options.

Specifying Initial Memory Contents

The Block Memory Generator core supports memory initialization using a memory coefficient (COE) file, the default data option in the CORE Generator GUI, or a combination of both.

The COE file can specify the initial contents of each memory location, while default data specifies the contents of all memory locations. When used in tandem, the COE file can specify a portion of the memory space, while default data fills the rest the remaining memory space. COE files and default data shall be formatted with respect to the port A write width (or port A read width for ROM's).

A COE a text file which specifies two parameters:

- **memory_initialization_radix:** The radix of the values in the memory_initialization_vector. Valid choices are 2, 10, or 16.
- **memory_initialization_vector:** Defines the contents of each memory element. Each value is LSB-justified, separated by a space, and assumed to be in the radix defined by memory_initialization_radix.

The following is an example COE file. Note that semi-colon is the end of line character.

```
; Sample initialization file for a
; 32-bit wide by 16 deep RAM
memory_initialization_radix = 16;
memory_initialization_vector =
0 1 2 3 4 5 6 7
8 9 A B C D E F;
```

Verification

The Block Memory Generator core and the simulation models delivered with it are rigorously verified using advanced verification techniques, including a constrained-random configuration generator and a cycle-accurate bus functional model.

Resource Utilization and Performance

The resource utilization and performance of the Block Memory Generator core is highly dependent on user selections, such as algorithm, optional output registers, and memory size.

Block RAM Usage

When using the fixed primitive algorithm, the number of block RAM primitives used is equal to the width ratio (rounded up) multiplied by the depth ratio (rounded up), where the width ratio is the width of the memory divided by the width of the selected primitive, and the depth ratio is the depth of the memory divided by the depth of the primitive selected.

When using the minimum area algorithm, it is not as easy to determine the exact block RAM count. This is because the actual algorithm performs complex manipulations to produce an optimal solution. The optimistic estimate is total memory bits divided by 18k (the total number of bits per primitive) rounded up. Given that the algorithm packs block RAM's very efficiently, this estimate is often very accurate for most memories.

LUT Utilization and Performance

The LUT utilization and performance of the core are directly related to the arrangement of primitives and the selection of output registers. Particularly, the number of primitives cascaded in depth to imple-

ment a memory determines the size of the output multiplexer and the size of the input decoder, which are implemented in the FPGA fabric.

It should be noted that while the minimum area algorithm's primary goal is to use the minimum number of block RAM primitives, it has a secondary goal of maximizing performance – as long as block RAM usage does not increase.

Resource Utilization and Performance Examples

The following tables provide examples of actual resource utilization and performance for Block Memory Generator implementations. Each section highlights the effects of a specific feature on resource utilization and performance.

Benchmarks were performed targeting a Virtex-II Pro FPGA in the -5 speed grade (2vp30-ff1152-5) and a Virtex-4 FPGA in the -10 speed grade (4vlx60-ff1148-10). Better performance may be possible with higher speed grades.

In the benchmark designs below, the core was encased in a wrapper with input and output registers to remove the effects of IO delays from the results; performance may vary depending on the user design. The minimum area algorithm was used unless otherwise noted. The examples below highlight the use of embedded registers in Virtex-4 and the subsequent significant performance improvement.

Single Primitive

The Block Memory Generator does not add additional logic if the memory can be implemented in a single Block RAM primitive. [Table 4](#) shows performance data for single-primitive memories.

Table 4: Single Primitive Examples

| Memory Type | Options | Width x Depth | Resource Utilization | | | Performance | |
|--------------------|---------|---------------|----------------------|------|-----|---------------|----------|
| | | | Block RAMs | LUTs | FFs | Virtex-II Pro | Virtex-4 |
| True Dual-Port RAM | | 36x512 | 1 | 0 | 0 | 331 | 307 |
| | | 9x2k | 1 | 0 | 0 | 350 | 351 |
| | 1 | 36x512 | 1 | 0 | 0 | N/A | 365 |
| | | 9x2k | 1 | 0 | 0 | N/A | 400 |

Options

1. Virtex-4 embedded output registers

Output Registers

The Block Memory Generator optional output registers increase the performance of memories by isolating the Block RAM primitive clock-to-out delays and the data output multiplexer delays.

The output registers are only implemented for output ports. For this reason, when output registers are used, a single port RAM requires fewer resources than a true dual port RAM. Note that the effects of the core output registers are not fully illustrated due to the simple register wrapper used. In a full-scale user design, core output registers may improve performance notably.

Table 5: Virtex-II Pro Output Register Examples

| Memory Type | Width x Depth | Options | Virtex-II Pro | | | |
|--------------------|---------------|---------|---------------|-----|------|-------------------|
| | | | Block RAMs | FFs | LUTs | Performance (MHz) |
| True Dual-Port RAM | 17x5k | | 5 | 6 | 42 | 290 |
| | | 1 | 5 | 98 | 110 | 288 |
| | | 2 | 5 | 40 | 42 | 291 |
| | | 1,2 | 5 | 134 | 76 | 344 |
| Single-Port RAM | 17x5k | | 5 | 3 | 21 | 297 |
| | | 1 | 5 | 49 | 55 | 312 |
| | | 2 | 5 | 20 | 21 | 299 |
| | | 1,2 | 5 | 67 | 38 | 350 |

Options

1. Primitive output registers
2. Core output registers

In Virtex-4 architectures, the embedded Block RAM may be utilized, reducing the FPGA fabric resources required to create the registers.

Table 6: Virtex-4 Output Register Examples

| Memory Type | Width x Depth | Options | Virtex-4 | | | |
|--------------------|---------------|---------|------------|-----|------|-------------------|
| | | | Block RAMs | FFs | LUTs | Performance (MHz) |
| True Dual-Port RAM | 17x5k | | 5 | 6 | 42 | 249 |
| | | 1 | 5 | 12 | 42 | 327 |
| | | 2 | 5 | 40 | 42 | 253 |
| | | 1,2 | 5 | 52 | 42 | 370 |
| Single-Port RAM | 17x5k | | 5 | 3 | 21 | 263 |
| | | 1 | 5 | 6 | 21 | 375 |
| | | 2 | 5 | 20 | 21 | 263 |
| | | 1,2 | 5 | 26 | 21 | 382 |

Options

1. Embedded primitive output registers
2. Core output registers

Aspect Ratios

The Block Memory Generator selectable port and data width aspect ratios may increase Block RAM usage and affect performance, because aspect ratios limit the primitive types available to the algorithm, which can reduce packing efficiency. Large aspect ratios, such as 1:32, have a greater impact than small aspect ratios. Note that width and depth are reported with respect to the port A write interface.

Table 7: Virtex-II Pro Aspect Ratio Examples

| Memory Type | Width x Depth | Options | Virtex-II Pro | | | |
|--------------------|---------------|---------|---------------|-----|------|-------------------|
| | | | Block RAMs | FFs | LUTs | Performance (MHz) |
| True Dual-Port RAM | 17x5k | 1 | 5 | 6 | 42 | 290 |
| | 136x640 | 2 | 9 | 0 | 0 | 259 |

Options

1. Port aspect ratio = 1:1
2. Port aspect ratio = 1:8 (A port is 136x640; B port is 17x5k)

Table 8: Virtex-4 Aspect Ratio Examples

| Memory Type | Width x Depth | Options | Virtex-4 | | | |
|-----------------|---------------|---------|------------|-----|------|-------------------|
| | | | Block RAMs | FFs | LUTs | Performance (MHz) |
| Single-Port RAM | 17x5k | 1 | 5 | 3 | 21 | 263 |
| | | 2 | 9 | 0 | 0 | 301 |
| | 136x640 | 3 | 9 | 0 | 0 | 319 |

Options

1. Data width aspect ratio = 1:1
2. Data width aspect ratio = 1:8 (read port is 136x640, write port is 17x5k)
3. Data width aspect ratio = 8:1 (read port is 17x5k, write port is 136x640)

Algorithm

The differences between the minimum area algorithm and the selectable primitive algorithm are discussed in detail in "[Selectable Memory Algorithm](#)" on page 2. Table 9 shows examples of the resource utilization and the performance difference between them for two selected configurations.

Table 9: Memory Algorithm Examples

| Memory Type | Width x Depth | Options | Resource Utilization | | | Performance (MHz) | |
|-----------------|---------------|---------|----------------------|-----|------|-------------------|----------|
| | | | Block RAMs | FFs | LUTs | Virtex-II Pro | Virtex-4 |
| Single-Port RAM | 17x5k | 1 | 5 | 3 | 21 | 297 | 263 |
| | | 2 | 5 | 3 | 56 | 292 | 247 |
| | 36x4k | 1 | 8 | 1 | 38 | 285 | 291 |
| | | 3 | 8 | 4 | 152 | 252 | 217 |

Options

1. Minimum area algorithm
2. Selectable Primitive Algorithm using 18x1k Block RAMs
3. Selectable Primitive Algorithm using 36x512 Block RAMs

Supplemental Information

The following sections provide additional information about working with the Block Memory Generator core.

- **Interfacing with Older Memory Cores.** Defines the differences between older memory cores and the Block Memory Generator core.
- **SIM Parameters.** Defines the SIM parameters used to specify the core configuration.
- **Output Register Configurations.** Provides information optional output registers used to improve core performance.

Interfacing with Older Memory Cores

This section describes the differences between the LogiCORE Single and Dual Port Block Memory cores and the Block Memory Generator. The new Block Memory Generator is not backward compatible with the Single and Dual Port Block Memory core in several aspects:

- Single Port Block Memory core and the Block Memory Generator do not have the same port names for a single-port memory configuration
- Dual Port Block Memory core and the Block Memory Generator do not have the same reset pin name
- XCO files for the previous memory cores are NOT compatible with the Block Memory Generator
- Pin polarity, handshaking, and input register features supported by the previous memory cores are not supported in the Block Memory Generator
- The behavior of set/reset pin (SSR) and the enable pin (EN) in the Block Memory Generator will differ from previous cores when using optional output registers

Port Memory Pin Names

In the Block Memory Generator, the port names have been changed to reflect the actual ports on the block RAM primitives. [Table 10](#) and [Table 11](#) reflect the changes in port names.

Table 10: Port Name Changes from Single Port Block Memory Core

| Single Port Block Memory v6.2 (Old core) | Block Memory Generator v1.1 (New Core, Single Port Configuration) |
|---|--|
| DIN | DINA |
| ADDR | ADDRA |
| EN | ENA |
| WE | WEA |
| SINIT | SSRA |
| CLK | CLKA |
| DOUT | DOUTA |
| ND | Not supported |
| RFD | Not supported |
| RDY | Not supported |

Table 11: Port Name Changes from Dual Port Block Memory Core

| Dual Port Block Memory v6.2 (Old core) | Block Memory Generator v1.1 (New Core, Single Port Configuration) |
|---|--|
| SINITA | SSRA |
| SINITB | SSRB |
| ND [A B] | Not supported |
| RFD [A B] | Not supported |
| RDY [A B] | Not supported |

Obsolete Features

Minimum Parameterizable Depth

The minimum depth for the previous cores is 2. The Block Memory Generator core has a minimum depth of 8.

Pin Polarity Option

The Block Memory Generator does not support pin polarity options on the clock, enable, write enable and set/reset input pins. When active low signaling is desired, users can invert the signals prior to the input of the core.

Handshaking Pins

The Block Memory Generator does not support handshaking pins: ND [A | B] , RFD [A | B] and RDY [A | B] .

Input Registers

The Block Memory Generator does not support input registers on port DIN, ADDR, and WE prior to accessing the memory primitives.

Modified Behaviors

Enable Pin

Table 12 illustrates the difference in enable behavior between the previous cores and the Block Memory Generator.

Table 12: Old and New Core Differences

| | Single (or Dual) Port Block Memory v6.2 | Block Memory Generator v1.1 |
|---|--|---|
| 1 | Enable signal is pipelined and each pipeline register enables the registers or memories in that pipe stage | Enable signal in NOT pipelined and is tied directly to the enables for all registers and memories |
| 2 | Single enable pin that controls the enables of all registers and memories | Two optional enable pins provided: <ul style="list-style-type: none"> REGCE pin that only controls the enables of the registers in the last output stage EN pin that controls the enables of all other registers and memory |

Synchronous Reset Pin

In the previous memory cores, the synchronous reset (SINIT pin) initializes all memory latches and output registers. In the Block Memory Generator, the synchronous set/reset (SSR pin) only resets the registers in the last output stage. When there are no output register stages present, it will initialize the memory latches.

SIM Parameters

Table 13 defines the SIM parameters used to specify the configuration of the core. These parameters are only used to manually instantiate the core in HDL, calling the CORE Generator dynamically. This parameter list does not apply to users that generate the core using the CORE Generator GUI.

Table 13: SIM Parameters

| | SIM Parameter | Type | Range | Description |
|----|--------------------------|---------|---|---|
| 1 | C_FAMILY | String | Spartan-3, Spartan-3E, Virtex-II, Virtex-II Pro, Virtex-4, | Target device family |
| 2 | C_MEM_TYPE | Integer | 0: Single Port RAM 1: Simple Dual Port RAM 2: True Dual Port RAM 3: Single Port ROM 4: Dual Port RAM | Type of memory |
| 3 | C_ALOGRITHM | | 0 (selectable primitive), 1 (minimum area) | Type of algorithm |
| 4 | C_PRIM_TYPE | Integer | 0 (1-bit wide) 1 (2-bit wide) 2 (4-bit wide) 3 (9-bit wide) 4 (18-bit wide) 5 (36-bit wide) 6 (72-bit wide, single-port only) | If selectable primitive algorithm is chosen, defines which type of primitive to use to build memory |
| 5 | C_BYTE_SIZE | Integer | 8, 9 | Defines size of a byte; 8-bit or 9-bit |
| 6 | C_SIM_COLLISION_CHECK | String | None, Generate_X, All, Warnings_only | Defines warning collision checks in structural/unisim simulation model |
| 7 | C_COMMON_CLOCK | Integer | 0, 1 | Defines whether to optimize behavioral model collision check by assume clocks are synchronous |
| 8 | C_DISABLE_WARN_BHV_COLL | Integer | 0, 1 | Disables the behavioral model from generating warnings due to read-write collisions |
| 9 | C_DISABLE_WARN_BHV_RANGE | Integer | 0, 1 | Disables the behavioral model from generating warnings due to address out of range |
| 10 | C_LOAD_INIT_FILE | Integer | 0, 1 | Defined whether to load initialization file |
| 11 | C_INIT_FILE_NAME | String | "..." | Name of initialization file (MIF format) |
| 12 | C_USE_DEFAULT_DATA | Integer | 0, 1 | Defines whether to use default data for the memory |

Table 13: SIM Parameters (Continued)

| | SIM Parameter | Type | Range | Description |
|----|-----------------------|---------|------------------------------------|---|
| 13 | C_DEFAULT_DATA | String | "..." | Defines a default data for the memory |
| 14 | C_HAS_MEM_OUTPUT_REGS | Integer | 0, 1 | Defines whether a register stage is added at the output of the memory latch |
| 15 | C_HAS_MUX_OUTPUT_REGS | Integer | 0, 1 | Defines whether a register stage is added at the output of the memory core |
| 16 | C_WRITE_WIDTHA | Integer | 1 to 1152 | Defines width of write port A |
| 17 | C_READ_WIDTHA | Integer | 1 to 1152 | Defines width of read port A |
| 18 | C_WRITE_DEPTHA | Integer | 8 to 9011200 | Defines depth of write port A |
| 19 | C_READ_DEPTHA | Integer | 8 to 9011200 | Defines depth of read port A |
| 20 | C_ADDRA_WIDTH | Integer | 3 to 24 | Defines the width of address A |
| 21 | C_WRITE_MODEA | String | Write_First, Read_first, No_change | Defines the write mode for port A |
| 22 | C_HAS_ENA | Integer | 0, 1 | Defines whether port A has an enable pin |
| 23 | C_HAS_REGCEA | Integer | 0, 1 | Defines whether port A has an enable pin for its output register |
| 24 | C_HAS_SSRA | Integer | 0, 1 | Defines whether port A has an synchronous reset pin |
| 25 | C_SINITA_VAL | String | "..." | Defines initialization/power-on value for port A output |
| 26 | C_USE_BYTE_WEA | Integer | 0, 1 | Defines whether byte-write feature is used on port A |
| 27 | C_WEA_WIDTH | Integer | 1 to 128 | Defines width of WEA pin for port A |
| 28 | C_WRITE_WIDTHB | Integer | 1 to 1152 | Defines width of write port B |
| 29 | C_READ_WIDTHB | Integer | 1 to 1152 | Defines width of read port B |
| 30 | C_WRITE_DEPTHB | Integer | 8 to 9011200 | Defines depth of write port B |

Table 13: SIM Parameters (Continued)

| | SIM Parameter | Type | Range | Description |
|----|----------------|---------|------------------------------------|--|
| 31 | C_READ_DEPTHB | Integer | 8 to 9011200 | Defines depth of read port B |
| 32 | C_ADDRB_WIDTH | Integer | 3 to 24 | Defines the width of address B |
| 33 | C_WRITE_MODEB | String | Write_First, Read_first, No_change | Defines the write mode for port B |
| 34 | C_HAS_ENB | Integer | 0, 1 | Defines whether port B has an enable pin |
| 35 | C_HAS_REGCEB | Integer | 0, 1 | Defines whether port B has an enable pin for its output register |
| 36 | C_HAS_SSRB | Integer | 0, 1 | Defines whether port B has an synchronous reset pin |
| 37 | C_SINITB_VAL | String | "..." | Defines initialization/power-on value for port B output |
| 38 | C_USE_BYTE_WEB | Integer | 0, 1 | Defines whether byte-write feature is used on port B |
| 39 | C_WEB_WIDTH | Integer | 1 to 128 | Defines width of WEA pin for port B |

Output Register Configurations

The Block Memory Generator core allows optional output registers to improve the performance of the core. Since Virtex-4 block RAM primitives have embedded output registers that the previous architectures did not support, the configurations described below are separated into Virtex-4 and non-Virtex-4 implementations. Note that each optional register used adds an additional clock cycle of latency to the read operation.

Virtex-4 Output Register Configurations

To tailor the register options, two check boxes are provided in the CORE Generator GUI under the Output Pipelining section. The embedded output registers are enabled when the box labeled “*Register output of memory primitive*” is checked. Likewise, registers at the output of the core are enabled by checking the box labeled “*Register output of memory core.*” Illustrated below are the resulting implementations of the four permutations supported by these configuration options.

Virtex-4 Memory with Embedded and Core Output Registers

When both options are selected, then a memory core is generated with both the Virtex-4 embedded registers enabled and a set of registers on the output, as shown in [Figure 27](#). If you are building a large memory construct, selecting this configuration can provide improved performance.

- Register output of memory primitive
- Register output of memory core

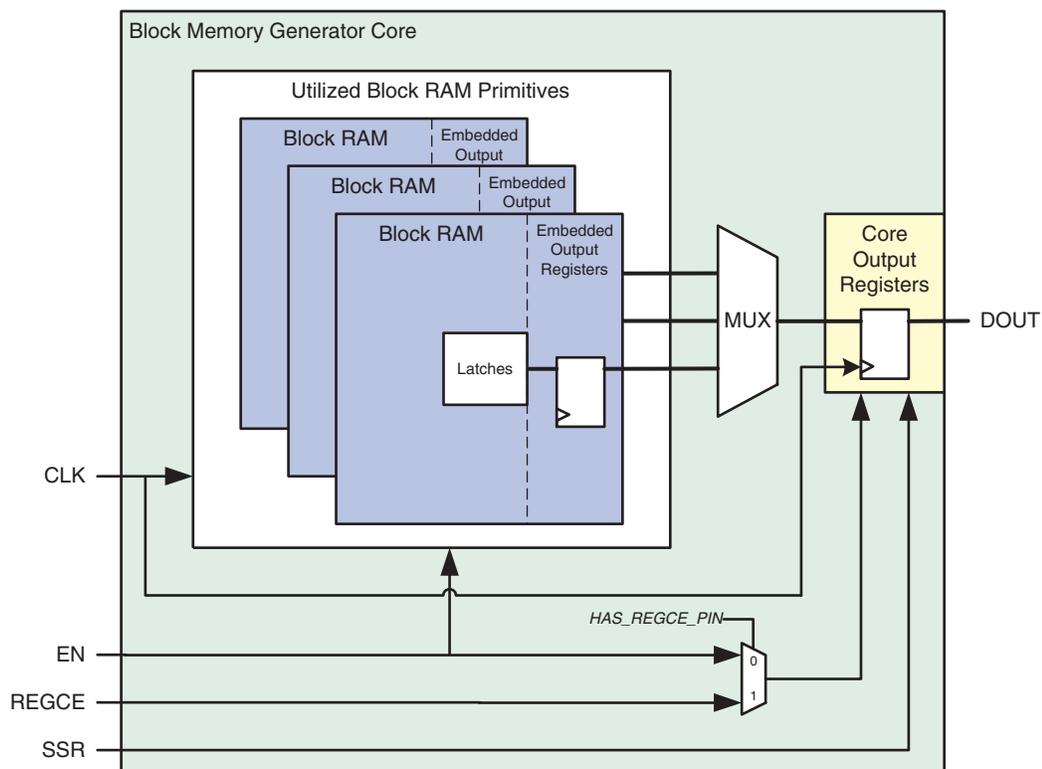


Figure 27: Virtex-4 Block Memory Generated With Both Embedded and Core Output Registers Enabled

Virtex-4 Memory with Embedded Registers Only

When only the option to use the primitive registers is selected, then a memory core is generated that registers the output of the block RAM primitives. Note that the output of any multiplexing which may be required to combine multiple primitives, will not be registered in this configuration, as shown in Figure 28.

- Register output of memory primitive
- Register output of memory core

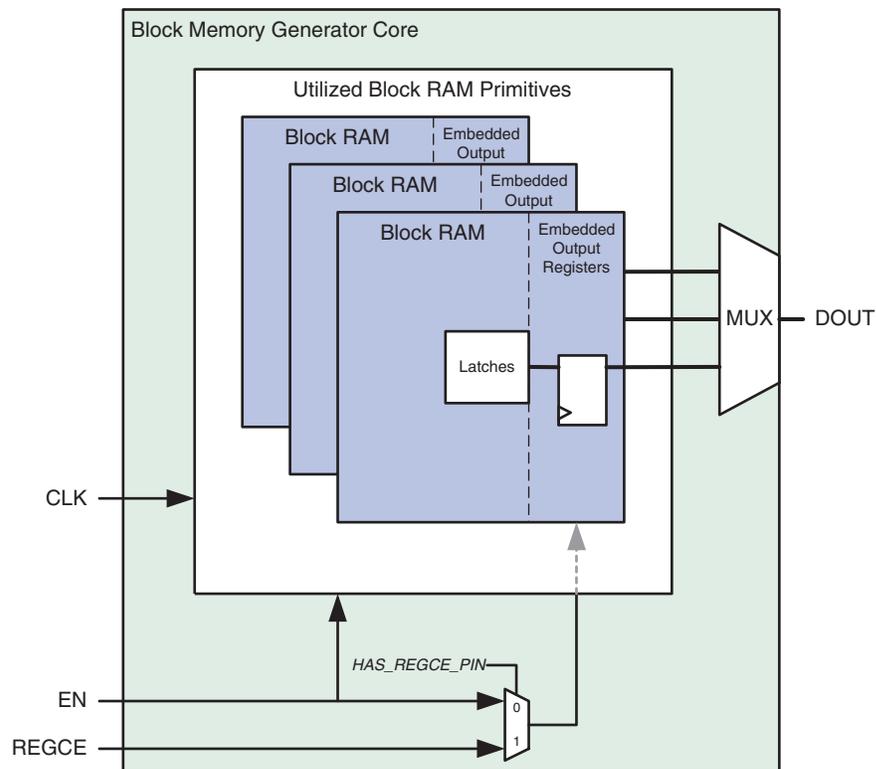


Figure 28: Virtex-4 Block Memory Generated With Embedded Registers

If the option to use the synchronous reset pin (SSRA, SSRB) is selected, then the Virtex-4 embedded block RAM registers can not be utilized. The primitive registers will be built from FPGA fabric as illustrated in Figure 29. If an extra cycle of latency is acceptable, then select both register stages as shown above in Figure 27 to enable the Virtex-4 embedded block RAM registers.

- Register output of memory primitive
- Register output of memory core
- Use SINITA (or SINITB) pin for resetting Port A DOUT bus

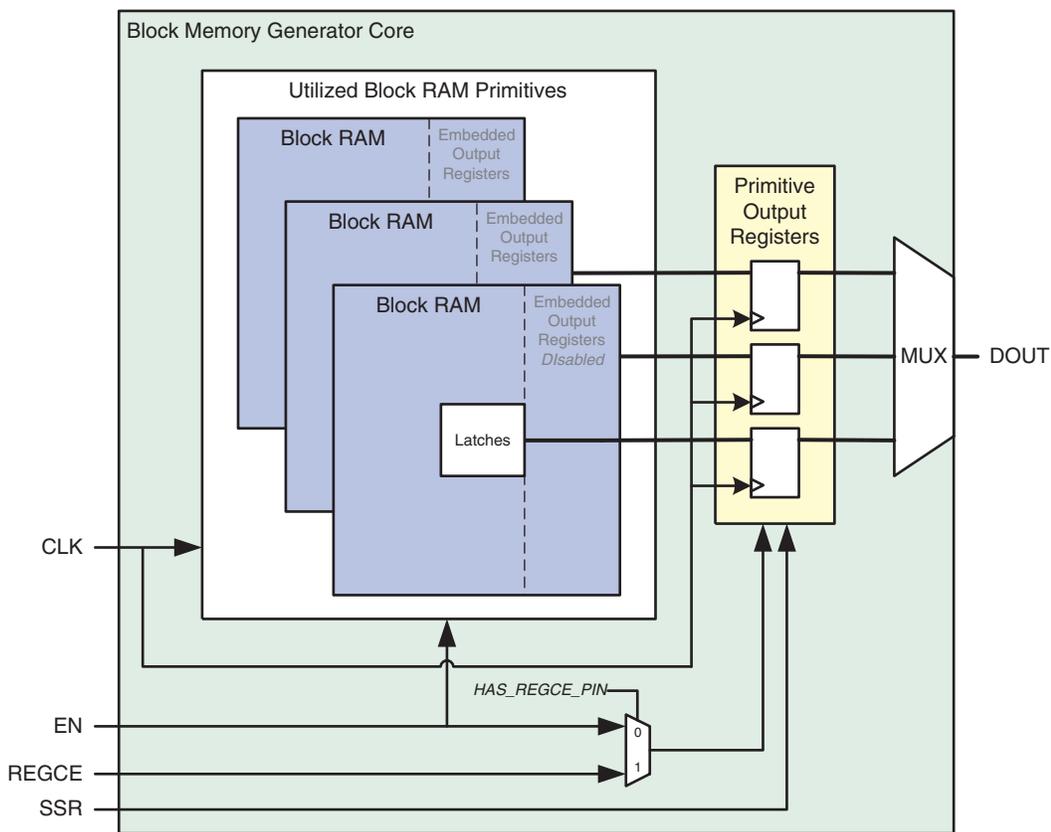


Figure 29: Virtex-4 Block Memory Generated With Embedded Registers and SSR

Virtex-4 Memory with Core Output Registers Only

When only the output is selected to be registered, then the core that is generated has the Virtex-4 embedded registers disabled, as shown in Figure 30.

- Register output of memory primitive
- Register output of memory core

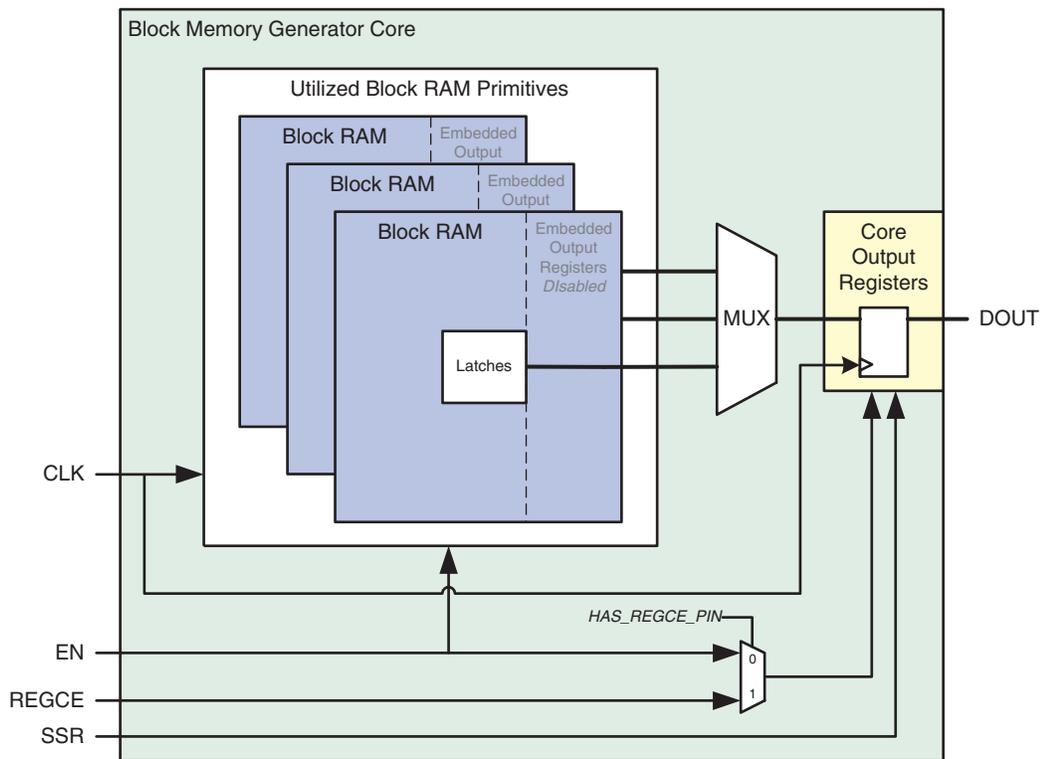


Figure 30: Virtex-4 Block Memory Generated With Core Output Registers

Virtex-4 Memory with No Optional Registers

When the no optional registers are selected, then the output of memory primitive is driven directly from the RAM primitive latches, and for cores that use more than one primitive, the output is driven from any the multiplexers required, and shown in [Figure 31](#).

- Register output of memory primitive
- Register output of memory core

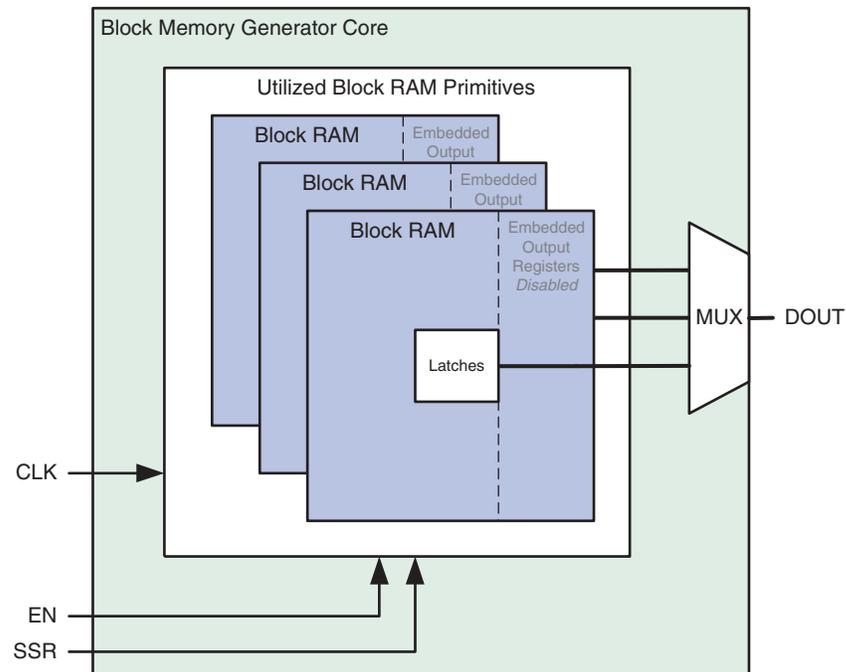


Figure 31: Virtex-4 Block Memory Generated With No Optional Registers

Non-Virtex-4: Output Register Configurations

To tailor the register options of non-Virtex-4 architectures, two check boxes are provided in the CORE Generator GUI under the Output Pipelining section. For implementing registers on the outputs of the individual block RAM primitives, the box labeled “Register output of memory primitive” is checked. Likewise, registering the output of the core is enabled by checking the box labeled “Register output of memory core.” Illustrated below are the resulting implementations of the four permutations supported by these configuration options.

Non-Virtex-4: Memory with Primitive and Core Output Registers

When both options are selected, then a memory core is generated with registers on the outputs of the individual RAM primitives and on the core output, as shown in [Figure 32](#). If you are building a large memory construct, selecting this configuration can provide improved performance.

- Register output of memory primitive
- Register output of memory core

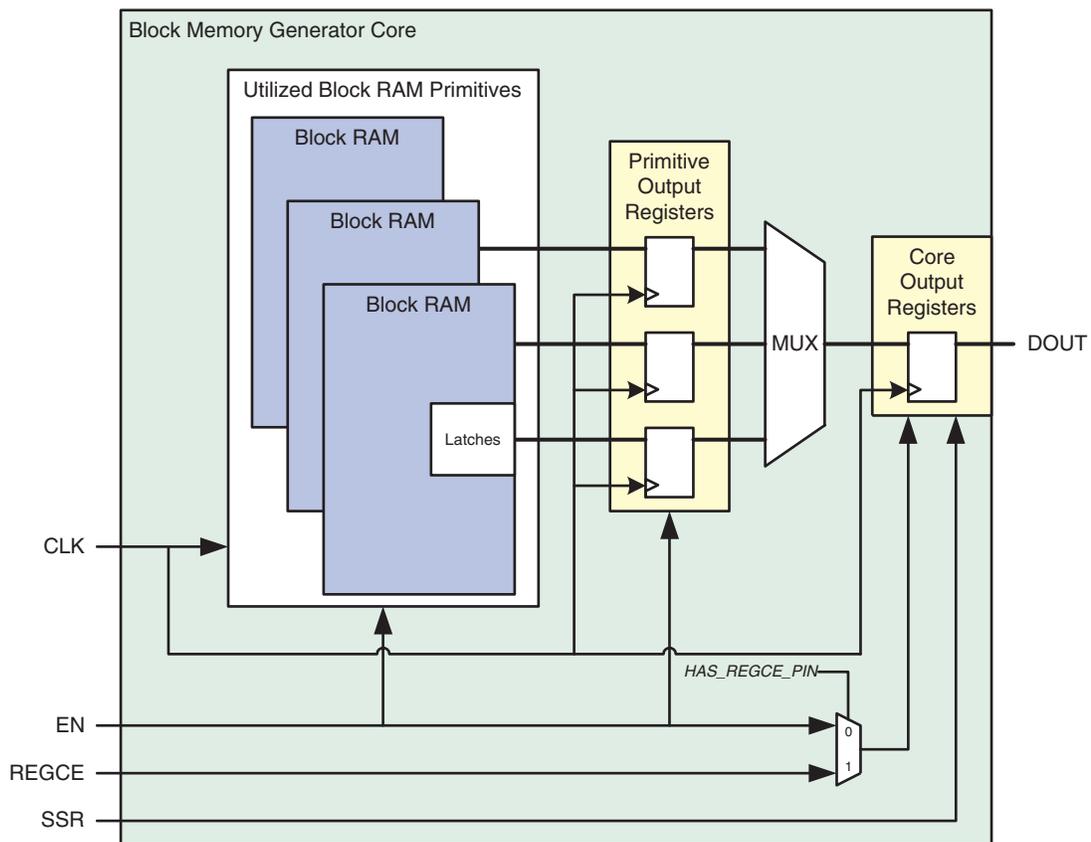


Figure 32: Non-Virtex-4 Memory With Both Primitive and Core Output Registers Enabled

Non-Virtex-4: Memory with Embedded Registers Only

When the option to register the primitive is selected, then a core is generated that only registers the output of the RAM primitives. Note that the output of any multiplexing required to combine multiple primitives will not be registered in this configuration, as shown in **Figure 33**.

- Register output of memory primitive
- Register output of memory core

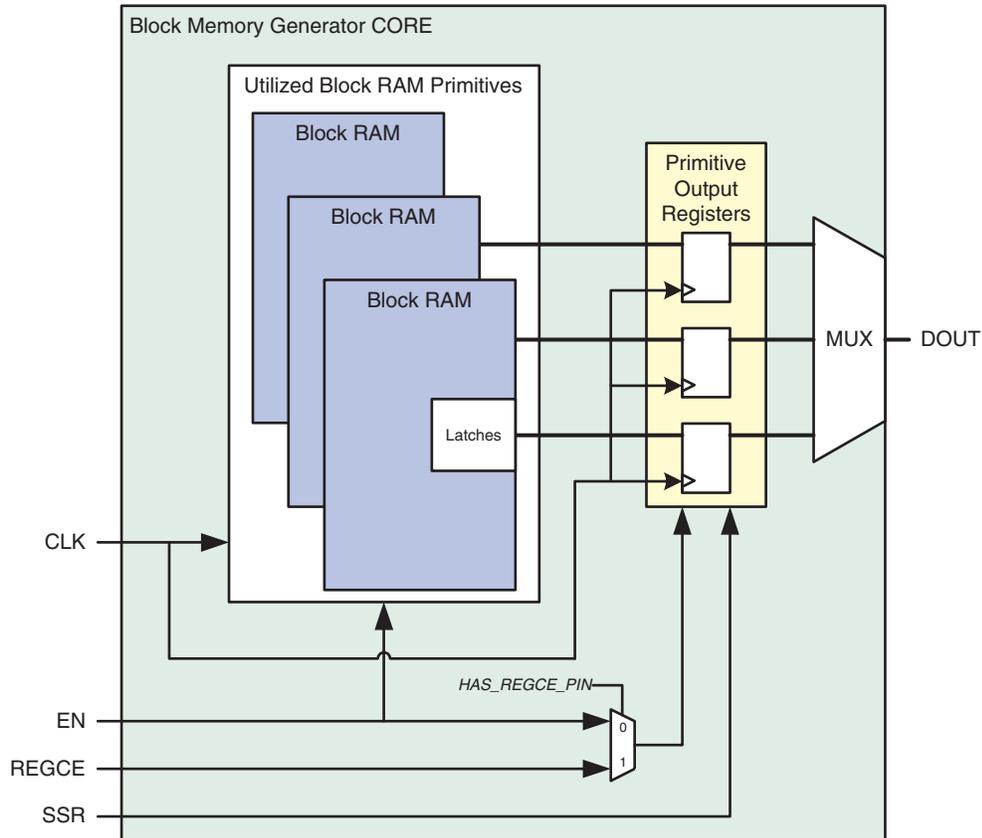


Figure 33: Non-Virtex-4 Block Memory Generated With Primitive Registers

Non-Virtex-4: Memory with Core Output Registers Only

A memory that is configured with only the output registered selected is shown in **Figure 34**.

- Register output of memory primitive
- Register output of memory core

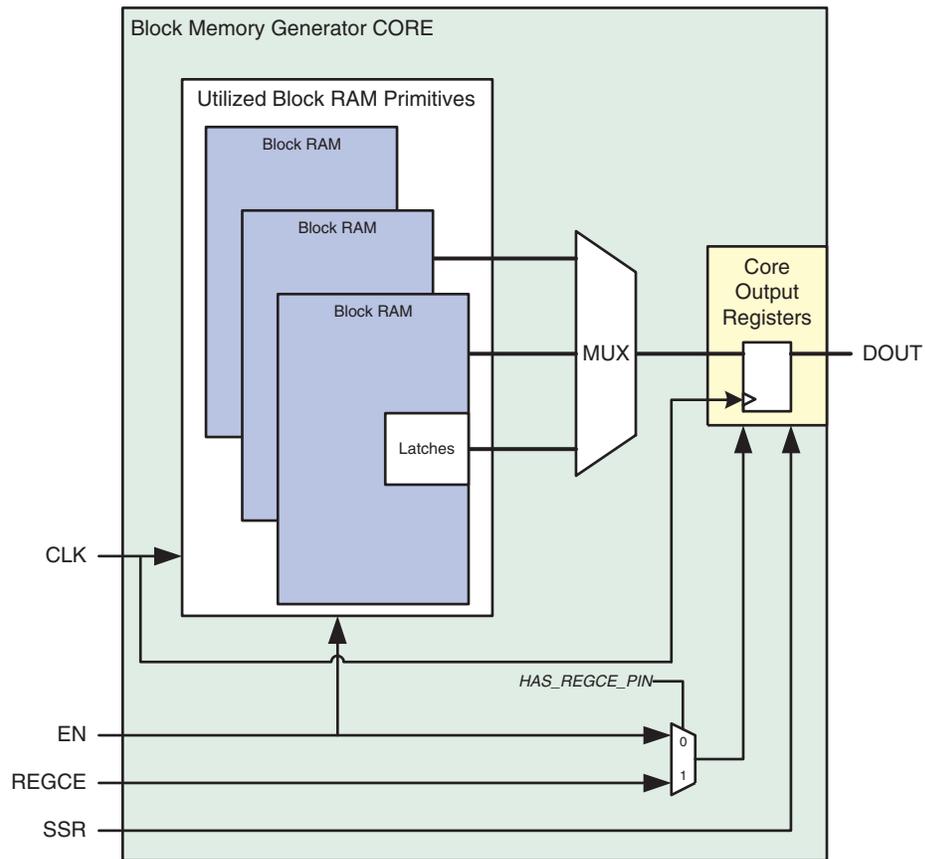


Figure 34: Non-Virtex-4 Block Memory Generated With Core Output Registers

Non-Virtex-4: Memory with No Optional Registers

When no output register option is selected, then the output of the memory primitive is driven directly from the memory latches, and for cores that use more than one primitive, the core output is driven from any multiplexers required, and shown in [Figure 35](#).

- Register output of memory primitive
- Register output of memory core

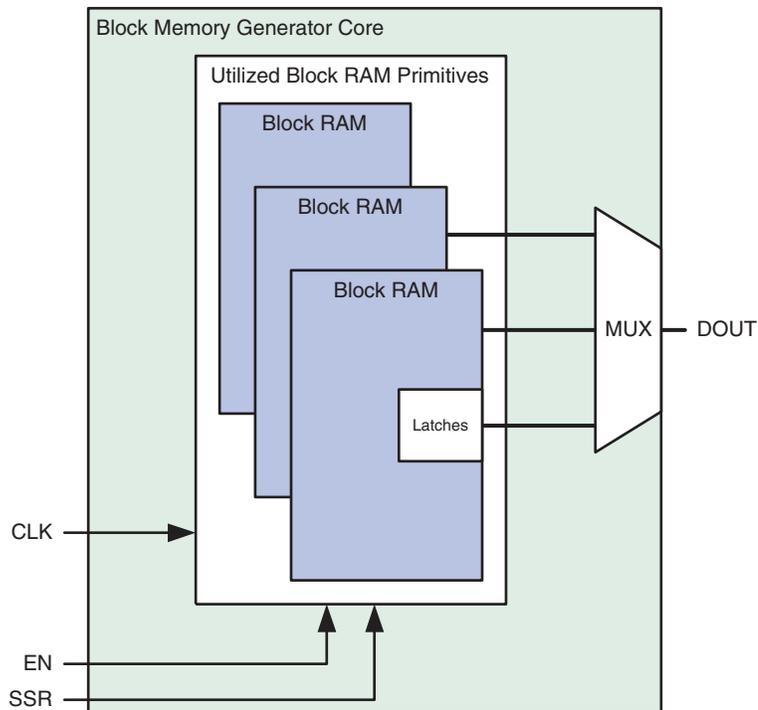


Figure 35: Non-Virtex-4 Block Memory Generated With No Optional Registers

Ordering Information

This free core is included with the Xilinx ISE CORE Generator system. Updates to the core are bundled with ISE IP Updates, which are accessible from the Xilinx [Download Center](#). To order Xilinx software, please contact your local Xilinx sales representative. Information on additional Xilinx LogiCORE modules is available at the Xilinx [IP Center](#).

Related Information

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Revision History

| Date | Version | Revision |
|----------|---------|-------------------------|
| 01/11/06 | 1.0 | Initial Xilinx release. |