



# Efficient calibration techniques for large-scale traffic simulators



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## ARTICLE INFO

### Article history:

Received 15 May 2016

Revised 2 December 2016

Accepted 3 December 2016

## ABSTRACT

Road transportation simulators are increasingly used by transportation stakeholders around the world for the analysis of intricate transportation systems. Model calibration is a crucial prerequisite for transportation simulators to reliably reproduce and predict traffic conditions. This paper considers the calibration of transportation simulators. The methodology is suitable for a broad family of simulators. Its use is illustrated with stochastic and computationally costly simulators. The calibration problem is formulated as a simulation-based optimization (SO) problem. We propose a metamodel approach. The analytical metamodel combines information from the simulator with information from an analytical differentiable and tractable network model that relates the calibration parameters to the simulation-based objective function. The proposed algorithm is validated by considering synthetic experiments on a toy network. It is then used to address a calibration problem with real data for a large-scale network: the Berlin metropolitan network with over 24300 links and 11300 nodes. The performance of the proposed approach is compared to a traditional benchmark method. The proposed approach significantly improves the computational efficiency of the calibration algorithm with an average reduction in simulation runtime until convergence of more than 80%. The results illustrate the scalability of the approach and its suitability for the calibration of large-scale computationally inefficient network simulators.

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## 1. Introduction

This paper focuses on the calibration (i.e., the estimation of the input parameters) of simulation-based road transportation models. We use the term “traffic simulator” to denote any simulation-based transportation model, whether macroscopic, mesoscopic or microscopic. The problem of model calibration has been extensively studied by the transportation community. A survey of both analytical and simulation-based calibration problems and algorithms is given in Balakrishna (2006). The most extensively studied calibration problem is, arguably, that of the calibration of origin-destination (OD) matrices, with seminal work such as Cascetta and Nguyen (1988); Cascetta et al. (1993) and more recent work such as Zhou and Mahmassani (2007; 2006); Zhou (2004); Ashok and Ben-Akiva (2002) and Ashok (1996), as well as the deployment of path flow estimators (PFEs), with the seminal work of Bell et al. (1997). Recent PFE reviews include Flötteröd (2008) and Buisson et al. (2012).

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This paper focuses on the calibration of simulators that are computationally costly to evaluate. The high computational cost can be due to: (i) the simulation of high levels of demand along with a high-resolution representation of demand (e.g., disaggregate representation of travelers), (ii) the simulation of a large-scale network, (iii) the use of a stochastic simulator requiring the evaluation of numerous simulation replications, (iv) the desire to evaluate performance under equilibrium conditions, which requires running sequentially multiple simulation-based assignment iterations (Nagel and Flötteröd, 2012).

Given the high computational costs involved in evaluating the simulation models, there is a need for calibration algorithms that can identify solutions with good performance at a low computational cost. That is, algorithms that can identify good solutions within few algorithmic iterations, i.e., with few simulated (near-)equilibrium responses.

We first discuss the main challenges of addressing calibration problems. We then state how the approach proposed in this paper addresses these challenges. The traffic model is often stochastic. It can involve sampling for every traveler from a variety of disaggregate behavioral models (e.g., choice models such as departure time, mode, route, lane-changing, etc.). Thus, a single run of the simulator involves drawing, for each of the thousands or hundreds of thousands of travelers, from a set of behavioral distributions. Given a sample of behavioral choices, a traffic flow model is used to propagate the travelers throughout the network. A review of state-of-the-art simulation models is presented in Barceló (2010). Thus, the mapping between the calibration parameters and the probability distribution of a given performance metric (e.g., the objective function of a calibration problem) is an intricate function.

This mapping is often non-convex and may contain multiple local minima. The stochasticity of the simulator requires the use of optimization algorithms that account for the lack of both: (i) a closed-form expression of the objective function and, (ii) exact function evaluations (since the functions can only be estimated via simulation). The traditional statistical and numerical algorithms are of limited use for calibration problems, since the underlying simulators often lack the strong assumptions required by these methods (e.g., normality, ergodicity) (Buisson et al., 2012). Hence, the traditional approach to calibration has been the use of simulation-based optimization (SO) algorithms. Most SO algorithms are general-purpose algorithms, they are not tailored to the intricacies of transportation simulators. They have been used extensively for calibration, several reviews include Ben-Akiva et al. (2012); Balakrishna (2006) and Antoniou (2004). Algorithms frequently used for calibration include simultaneous perturbation stochastic approximation (SPSA) (Spall, 1992), genetic algorithms (GA) (Holland, 1975), particle filters and Kalman filters.

The generality of these SO algorithms stems from the fact that they treat the simulator as a black-box. The main implication of this is that they are designed to achieve asymptotic (i.e., large-sample size) convergence properties. They are not designed to identify good solutions within few algorithmic iterations, i.e., they are not computationally efficient. Yet they are used by the transportation community under tight computational budgets, i.e., under small-sample size conditions. One approach to derive computationally efficient algorithms, is to exploit the structure of the underlying calibration problem. General-purpose SO algorithms exploit limited problem-structure (e.g., at most they are based on numerical linearizations). One recent work that does exploit problem structure within the calibration algorithm is that of Flötteröd et al. (2011). It formulates and embeds within the algorithm an analytical approximation of the first-order derivative of the simulator's measurement equation. This leads to significant reductions in the computational requirements of the algorithm.

Traditional SO algorithms, although designed to guarantee asymptotic properties, are typically used for calibration under tight computational budgets. This makes them sensitive to the initial points. Given the difficulty of identifying good initial points to the calibration problem, there is a need for algorithms that perform well under tight computational budgets while being robust to the quality of the initial points. Additionally, the lack of computationally efficient calibration algorithms has led recent calibration research to focus on the design of dimensionality reduction methods (e.g., sensitivity analysis methods) (Ge et al., 2014; Ge and Menendez, 2016; 2014; Ciuffo and Azevedo, 2014).

This paper addresses the following challenges.

**Computational efficiency** We propose an algorithm that can identify points with good performance within few algorithmic iterations. Therefore, it is a computationally efficient algorithm that reflects well the computational conditions under which calibration problems are addressed by both the transportation research and practice communities. This is achieved by designing an algorithm that exploits the transportation-specific structure of the calibration problem. More specifically, the proposed approach solves at every iteration of the calibration algorithm, an analytical (i.e., not simulation-based) approximate calibration problem. This analytical problem is solved by using information from an analytical traffic model. The latter is highly efficient. It is formulated as a system of nonlinear equations. Hence, it can be evaluated with a variety of efficient solvers. This is key for the efficiency of the calibration algorithm.

The proposed approach resorts to the use of a derivative-free algorithm. In other words, it does not rely on estimates of the derivatives of the simulation-based objective function. This further contributes to the efficiency of the algorithm.

**Analytical structural information** The algorithm embeds an analytical approximation of the simulator. This contributes to a largely unresolved methodological challenge which is the formulation of tractable measurement equations that link available surveillance field data to the simulator's calibration parameters.

**Robustness to initial conditions** The algorithm can identify good solutions within few iterations regardless of the initial points. It is robust to the quality of the initial points.

**Stochasticity** The algorithm is a simulation-based optimization algorithm that accounts for the simulator's stochasticity.

**Large-scale networks** The algorithm is applied to a calibration problem for the Berlin metropolitan network (over 24300 links and 11300 nodes). Compared to a black-box algorithm, the proposed approach reduces simulation runtime until convergence by more than 80%, on average. This illustrates the efficiency of the approach for the calibration of large-scale networks. The design of a calibration algorithm suitable for large-scale networks is achieved with the formulation of a scalable analytical model. In particular, the model is defined as a system of nonlinear equations with a dimension that scales linearly with the number of links, and linearly with the number of OD pairs.

This paper formulates the calibration problem as a simulation-based optimization problem and uses a metamodel SO algorithm. The paper formulates a novel metamodel suitable for demand calibration problems. The algorithm is used to address a calibration problem for a large-scale Berlin metropolitan network. [Section 2](#) presents the proposed methodology, followed by case studies on both a toy network and a Berlin metropolitan network ([Section 3](#)). We conclude with a brief discussion ([Section 4](#)). [Appendix A](#) details the SO algorithm. [Appendix B](#) contains a list of the notation used throughout the paper. [Appendix C](#) contains additional details of the numerical results of [Section 3](#).

With the increase in the availability, diversity and quality of travel data, comes an increasing interest and relevance of the joint calibration of demand and supply parameters (e.g., [Balakrishna, 2006](#); [Antoniou et al., 2007](#); [Vaze et al., 2009](#)). The formulation of the proposed framework for the calibration of supply parameters has been recently derived ([Zhang et al., 2016a](#)). Thus, we expect the general ideas presented in this paper to be suitable for the joint calibration of demand and supply parameters.

This paper illustrates the use of the proposed algorithm with the calibration of a single demand parameter. Hence, the results do not indicate the suitability of the approach for high-dimensional problems. Nonetheless, the metamodel ideas proposed in this paper have shown to be scalable for traffic management problems (e.g., [Osorio and Chong, 2015](#); [Chen et al., 2012](#)). We expect these ideas to also scale well for calibration problems.

The case studies show that the proposed algorithm identifies good solutions within a couple of iterations. Hence, it can also be used as a technique to identify good initial solutions to launch or initialize traditional general-purpose calibration algorithms. This is particularly important when using the general-purpose algorithms under tight computational budgets, which are by design sensitive to the initial conditions.

## 2. Methodology

The general framework discussed in this paper is suitable to address a broad class of calibration problems (e.g., demand, supply, time-independent, time-dependent, etc.). [Section 2.1](#) formulates a general calibration problem. [Section 2.2](#) presents a framework suitable to address this general problem. More specifically, the framework is based on the use of a family of optimization algorithms known as metamodel simulation-based optimization (SO) algorithms. In order to provide a more detailed formulation of the proposed framework, [Section 2.3](#) focuses on one specific calibration problem which is a time-independent one-dimensional demand calibration problem. [Section 2.3](#) formulates the calibration problem as a metamodel SO problem. It then formulates a suitable metamodel.

### 2.1. Formal problem statement

We consider the calibration of travel demand parameters from link flows. Origin-destination (OD) pairs are trip production and attraction points connected by a set of routes in an urban network. Let  $\tau$  and  $i$  index the simulation time intervals of duration  $T$  and the network links, respectively, and let  $n = 1 \dots N$  index the individual trip-makers (i.e., simulated travelers or agents) in the system. Denoting by  $\Delta_{ni\tau} \in \{0, 1\}$  the stochastic binary indicator of traveler  $n$  crossing link  $i$  in time interval  $\tau$ , the stochastic simulated link flow rate  $F_{i,\tau}$  on link  $i$  in time interval  $\tau$  is

$$F_{i,\tau} = \frac{1}{T} \sum_{n=1}^N \Delta_{ni\tau} \quad (1)$$

with the expectation

$$E[F_{i,\tau}] = \frac{1}{T} \sum_{n=1}^N \delta_{ni\tau}, \quad (2)$$

where  $\delta_{ni\tau} = E[\Delta_{ni\tau}]$  is the probability that traveler  $n$  crosses link  $i$  in time interval  $\tau$  ([Flötteröd and Bierlaire, 2009](#)). This probability is in turn dependent on the network conditions  $x$ , in particular travel times of the various routes, which in turn depend on the travel behavior (e.g., departure time, route choice), requiring typically to solve (2) iteratively. This is subsequently expressed by writing  $\delta_{ni\tau} = \delta_{ni\tau}(\theta; x)$  with  $\theta$  a parameter vector of the underlying behavioral model. Here,  $\theta$  represents a general calibration parameter vector. It can include, for instance, coefficients of attributes of behavioral models, such as the travel time coefficient of a route choice model or of a departure time choice model. In a transportation simulator that simulates the movement of individual travelers (microscopic and often mesoscopic), these iterations can be interpreted as a learning process over subsequent days, where in each day some travelers update their travel decisions (typically route

choice, in some models also time and mode choice) based on the most recent network conditions  $x$ , followed by a simulation of the corresponding vehicle flows through the network, which in turn yields updated network conditions.

Let  $y_{i,\tau}$  be the number of vehicles counted in the field on link  $i$  in time interval  $\tau$ . A traditional nonlinear least squares formulation of the calibration problem is then to minimize the following objective function:

$$f(\theta) = \min_{\theta \in \Theta} \sum_{i \in \mathcal{I}} \sum_{\tau} (y_{i,\tau} - E[F_{i,\tau}(\theta; x)])^2. \quad (3)$$

The summation considers all time intervals  $\tau$  and all links  $i$  that belong to the set of links with measurements available, denoted  $\mathcal{I}$ . The feasible region  $\Theta$  is defined analytically, and often consists of simple bound constraints.

Formulation (3) illustrates the main challenges of the calibration problem described in Section 1. The function  $E[F_{i,\tau}(\theta; x)]$  has no closed-form expression available, since  $x$  (e.g., the travel times) and  $E[F_{i,\tau}(\theta; x)]$  (i.e., the expected link flows) can only be estimated by evaluating the stochastic simulation model. Thus, Problem (3) cannot be solved using traditional analytical and deterministic optimization approaches. The function  $E[F_{i,\tau}(\theta; x)]$  is a nonlinear function that describes intricate spatial-temporal traffic phenomena in the stochastic traffic simulator and lacks sound mathematical properties such as convexity.

The purpose of this paper is to propose an efficient calibration algorithm for such difficult problems. The main idea is to embed within the algorithm analytical structural problem-specific information. In particular, we propose to formulate and provide the algorithm with an analytical approximation of the relationship between  $\theta$  and  $E[F_{i,\tau}(\theta; x)]$ . We expect this analytical information to enable computational efficiency.

Developing such an analytical approximation is a challenging problem because the approximated mapping involves the highly nonlinear and stochastic *network loading map* of path flows on network conditions, comprising in the simulation context all difficulties that come along with real traffic flow dynamics in urban networks (including, e.g., multi-lane flows, spillbacks, flow interactions in intricate intersections).

## 2.2. Metamodel simulation-based optimization methods

The following summary of ideas developed by Osorio and Bierlaire (2013) is needed to make the proposed calibration approach concrete. The broad family of SO problems considered can be formulated as follows:

$$\min_{\theta \in \Theta} f(\theta, z; \hat{q}). \quad (4)$$

Problem (4) consists of two components: a simulation-based objective function  $f$ , and a feasible region  $\Theta$ . The objective function  $f$  is not available in closed-form, it can only be estimated via simulation. It depends on: the calibration parameters  $\theta$ , endogenous simulation variables  $z$  (e.g., link flows, travel times) and exogenous simulation parameters  $\hat{q}$  (e.g., network topology). The feasible region  $\Theta$  is defined by a set of constraints assumed analytical (rather than simulation-based), differentiable and of general-form (e.g., non-convex).

For a review of metamodel SO methods, we refer the reader to Osorio (Chap. 5, 2010). The main idea underlying metamodel SO algorithms is to address the simulation-based Problem (4) by iteratively solving a set of analytical problems. At iteration  $k$  the SO algorithm solves an analytical problem with the following form:

$$\min_{\theta \in \Theta} m_k(\theta; \beta_k). \quad (5)$$

The main idea is to replace the simulation-based objective function ( $f$  in (4)) with an analytical approximation of it ( $m_k$  in (5)). The function  $m_k$  is known as the metamodel. It is often a parametric function, with the iteration-specific parameter vector,  $\beta_k$ , often being fitted based on simulated observations. Metamodel SO techniques iterate over two main steps. First, the metamodel is constructed based on the sample of simulated observations. Second, it is used to perform optimization and derive a trial point (e.g., a calibration parameter value). The performance of the trial point can be evaluated by the simulator, which leads to new simulation observations. As new observations become available the accuracy of the metamodel can be improved (Step 1), leading ultimately to better trial points (Step 2). In this paper, we use the derivative-free trust region algorithm of Conn et al. (2009). A full description of the algorithm is provided in Section 4.2 of Osorio and Bierlaire (2013). In order to use this framework for calibration, the main challenge is to formulate a metamodel that: (i) provides a good approximation of the mapping of the decision vector to the objective function, and (ii) is also computationally efficient such that Problem (5), which needs to be solved at every iteration of the algorithm, can be solved efficiently. This paper contributes by formulating a metamodel that achieves these two goals.

Metamodels can be classified as either: (i) functional models, these are general-purpose functions suitable to approximate an arbitrary function  $f$  (e.g., polynomials); (ii) physical models, these are problem-specific functions. In other words, their functional form depends on the specific problem.

We use a metamodel that combines ideas from functional models and physical models. At iteration  $k$  it is defined by:

$$m_k(\theta; \beta_k) = \beta_{k,0} f_A(\theta) + \phi(\theta; \beta_k). \quad (6)$$

Eq. (6) defines  $m_k$  as a linear combination of a general-purpose parametric function (e.g., a polynomial) denoted  $\phi$  and a problem-specific approximation of  $f$  (defined in (4)) denoted  $f_A$ . The first element of the vector  $\beta_k$  is denoted  $\beta_{k,0}$ . Based

on (6), the metamodel can be interpreted as a problem-specific approximation ( $f_A$ ) which is corrected by a scaling factor ( $\beta_{k,0}$ ) and an additive general-purpose correction term ( $\phi$ ).

The problem-specific approximation ( $f_A$ ) is derived from an analytical macroscopic traffic model. Hence, the analytical problem solved at every iteration is defined by:

$$\min_{\theta \in \Theta} m_k(\theta; \beta_k) \quad (7)$$

$$h(\theta, \tilde{v}; \tilde{q}) = 0, \quad (8)$$

where  $h$  denotes the analytical macroscopic traffic model, with endogenous variables  $\theta$  and  $\tilde{v}$  (e.g., expected link queue-lengths) and exogenous parameters  $\tilde{q}$  (e.g., network topology). This problem differs from Problem (5) in the presence of an additional set of constraints (8). This set of constraints represents the analytical macroscopic traffic model used.

The key to achieving computational efficiency is in formulating a problem-specific approximation  $f_A$  that is a good approximation of the true, unknown, simulation-based objective function  $f$ . Hence, for a given transportation problem the main challenge is in the formulation of a suitable traffic model (function  $h$  of (8)) that satisfies the following requirements.

- (i) It leads to a good analytical approximation ( $f_A$ ) of the simulation-based objective function ( $f$ ).
- (ii) It is a scalable traffic model, such that large-scale networks can be addressed.
- (iii) It is computationally efficient to solve. Every iteration of the SO algorithm requires solving Problem (7)–(8), which contains the analytical traffic model as a set of constraints. Hence, the traffic model needs to be computationally inexpensive to evaluate.

The broad family of metamodels defined by (6) have been used to efficiently address large-scale urban traffic management problems while using inefficient yet detailed stochastic microscopic simulators (Osorio and Selvam, 2016; Osorio and Chong, 2015; Osorio and Nanduri, 2015a; 2015b; Chen et al., 2012; Chong and Osorio, 2016). Metamodel approaches have also been used recently for addressing various transportation problems, such as in Chen et al. (2016), where a pricing optimization problem is addressed based on a large-scale mesoscopic network model. As in traditional simulation literature, the metamodel of Chen et al. (2016) is a general-purpose (also known as a functional) metamodel. This comes with the advantage of being a general-purpose methodology, which can be directly applied to a problem regardless of its formulation (e.g., the choice of the objective function, the decision variables, the underlying network structure). Nonetheless, general-purpose metamodels lack problem-specific structural information, and hence are not designed to be computationally efficient.

### 2.3. Metamodel formulation

The novel methodology proposed in this paper is valid for a general class of calibration problems. In order to illustrate and implement a specific instance of it, we focus, hereafter, on: (i) a time-independent calibration problem (i.e., we consider a single time interval), (ii) the calibration of a single behavioral parameter, a scalar  $\theta$ , which governs route choice. Recall, that Section 2.2 presented a general calibration problem with a general calibration vector denoted  $\theta$ . This section considers a specific calibration problem. Hence,  $\theta$  now represents a scalar that denotes the travel time coefficient of a route choice model.

The specific SO problem is then formulated as:

$$f(\theta) = \min_{\theta} \sum_{i \in \mathcal{I}} (y_i - E[F_i(\theta; x)])^2 \quad (9)$$

$$\theta_L \leq \theta \leq \theta_U, \quad (10)$$

where  $\theta_L$  (resp.  $\theta_U$ ) denotes a lower (resp. upper) bound. The purpose of the metamodel is to approximate the simulation-based performance metric  $E[F_i(\theta; x)]$ , which denotes the expected flow of link  $i$ . Let  $m_{i,k}$  denote the analytical approximation of  $E[F_i(\theta; x)]$  at iteration  $k$  of the SO algorithm. Given an expression for  $m_{i,k}$ , the SO algorithm would then solve a series of problems of the form:

$$\min_{\theta} \sum_{i \in \mathcal{I}} (y_i - m_{i,k}(\theta; \beta_{i,k}))^2 \quad (11)$$

$$\theta_L \leq \theta \leq \theta_U \quad (12)$$

$$h(\theta, \tilde{v}; \tilde{q}) = 0, \quad (13)$$

where  $h$  denotes the analytical traffic model used to analytically approximate the expected link demands.

The metamodel is formulated for link  $i$  and iteration  $k$  as:

$$m_{i,k}(\theta; \beta_{i,k}) = \beta_{i,k,0} \lambda_i(\theta) + \beta_{i,k,1} + \beta_{i,k,2} \theta, \quad (14)$$

where  $\theta$  is the (scalar) behavioral parameter,  $\lambda_i(\theta)$  is the expected demand for link  $i$  approximated by the analytical traffic model, and  $\beta_{i,k}$  is a three-dimensional vector of metamodel parameters:  $\beta_{i,k} = [\beta_{i,k,0}, \beta_{i,k,1}, \beta_{i,k,2}]$ . The metamodel can be interpreted as the analytical approximation ( $\lambda_i(\theta)$ ) corrected by a scaling factor ( $\beta_{i,k,0}$ ) and an additive linear in  $\theta$  error term (represented by  $\beta_{i,k,1} + \beta_{i,k,2}\theta$ ). Flötteröd et al. (2011) formulate a related calibration problem using a metamodel of the form (14) with  $\beta_{i,k,0}$  set to zero. In other words, it is a general-purpose (i.e., purely functional) metamodel that is also scalable and tractable, yet it does not embed an analytical problem-specific network model that provides structural information to the algorithm.

We now present the analytical traffic model that will yield the analytical approximation of expected link demands ( $\lambda_i(\theta)$ ). The model is a probabilistic and differentiable network model. We map the road network as a probabilistic queueing network. Each link is modeled as a single queue. Hereafter, the terms “link” and “queue” are used interchangeably. We first introduce notation, we then present the model formulation and comment on its derivation.

$d_s$	expected travel demand for OD pair $s$ ;
$\ell_i$	space capacity of queue $i$ ;
$l_i$	length of link $i$ (road length);
$v_i$	maximum speed of link $i$ ;
$\mu_i$	service rate of queue $i$ ;
$\gamma_i$	expected external demand of queue $i$ ;
$\lambda_i$	expected demand of queue $i$ ;
$\tilde{t}_i$	expected travel time of queue $i$ ;
$\tilde{n}_i$	expected number of vehicles in queue $i$ ;
$p_{ij}$	turning probability from queue $i$ to queue $j$ ;
$t_r$	expected travel time of route $r$ ;
$f_r$	expected demand on route $r$ ;
$\tilde{p}_{sr}$	probability that a traveler of OD pair $s$ takes route $r$ ;
$\theta$	travel time coefficient of the route choice model;
$\mathcal{S}$	set of OD pairs;
$\mathcal{Q}$	set of queues;
$\mathcal{R}$	set of routes;
$\mathcal{R}_s$	set of routes of OD pair $s$ ;
$\mathcal{G}_{ij}$	set of routes that consecutively go through queues $i$ and $j$ ;
$\mathcal{H}_i$	set of routes that go through queue $i$ ;
$\mathcal{T}_i$	set of routes that start with queue $i$ ;
$\Psi_r$	set of links of route $r$ .

$$\begin{cases}
 f_r = \sum_{s \in \mathcal{S}} d_s \tilde{p}_{sr} & \forall r \in \mathcal{R} & (a) \\
 \tilde{p}_{sr} = \frac{e^{\theta t_r}}{\sum_{j \in \mathcal{R}_s} e^{\theta t_j}} & \forall s \in \mathcal{S}, \forall r \in \mathcal{R}_s & (b) \\
 t_r = \sum_{i \in \Psi_r} \tilde{t}_i & \forall r \in \mathcal{R} & (c) \\
 \tilde{t}_i = \frac{l_i}{v_i} + \frac{\tilde{n}_i}{\lambda_i} & \forall i \in \mathcal{Q} & (d) \\
 \tilde{n}_i = \frac{\tilde{\rho}_i}{1 - \tilde{\rho}_i} - \frac{(\ell_i + 1) \tilde{\rho}_i^{\ell_i + 1}}{1 - \tilde{\rho}_i^{\ell_i + 1}} & \forall i \in \mathcal{Q} & (e) \\
 \tilde{\rho}_i = \frac{\lambda_i}{\mu_i} & \forall i \in \mathcal{Q} & (f) \\
 \lambda_i = \gamma_i + \sum_{j \in \mathcal{Q}} p_{ji} \lambda_j & \forall i \in \mathcal{Q} & (g) \\
 p_{ij} = \frac{\sum_{r \in \mathcal{G}_{ij}} f_r}{\sum_{w \in \mathcal{H}_i} f_w} & \forall i \in \mathcal{Q}, \forall j \in \mathcal{Q} & (h) \\
 \gamma_i = \sum_{r \in \mathcal{T}_i} f_r & \forall i \in \mathcal{Q} & (i)
 \end{cases} \quad (15)$$

Eq. (15)a describes how expected OD demand is mapped to expected route demand through the route choice model. In other words, it defines the expected route demand as a weighted sum of expected OD demands. Eq. (15)b is the route choice model. It is a multinomial logit (MNL) model with a single attribute: expected route travel time. This is a simplification of the route choice model used in the simulation models of the case studies of Section 3. For a detailed description of the route choice models of the simulator, see Zhang et al. (2016b). Eq. (15)c defines the expected route travel time as the sum of expected link travel times.

Eq. (15)d approximates the expected link travel time as the sum of expected free-flow travel time and expected delay. The analytical approximation model used here is based on stochastic point queues, meaning that it does not capture spill-backs and that all link outflow constraints result from the link's downstream bottleneck capacity. Recall that the structural



metamodel is, by design, a simplified approximation of the simulator, which may itself use space queues and capture spillbacks. The expected time it takes a vehicle to traverse a spatial link with a point queue downstream bottleneck is the sum of: (i) the time to join the point queue, which corresponds to term  $\frac{l_i}{v_i}$  and represents a constant free-flow travel time, and (ii) the expected delay or time needed to pass the bottleneck, which corresponds to term  $\frac{\bar{n}_i}{\lambda_i}$  and is obtained by modeling the downstream bottleneck capacity as a distributed quantity in order to capture variability in link discharge flows.

The expected delay is based on the use of Little's law (Little, 2011; 1961), assuming an infinite space capacity queue. The expected free-flow travel time is defined as the travel time to travel the link at maximum speed. Eq. (15)e approximates the expected queue-length of a given queue. It is obtained by considering the queue as an isolated M/M/1/ℓ queue, with finite space capacity ℓ. A derivation of this expression can be found in Appendix A of Osorio and Chong (2015). Eq. (15)f defines  $\bar{\rho}_i$ , which is known as the traffic intensity of the queue, as the ratio of expected demand to expected supply. Eq. (15)g is a flow conservation equation that relates the expected demand on queue *i* (denoted  $\lambda_i$ ) to the sum of expected external demand to queue *i* (denoted  $\gamma_i$ ) and expected demand arising from upstream queues. The expected demand,  $\lambda_i$ , is also referred to as the arrival rate of the queue. The expected external demand,  $\gamma_i$ , represents demand that arises from outside the network, i.e., trips that start at queue *i*. The turning probability from queue *i* to queue *j* (denoted  $p_{ij}$ ) is defined by Eq. (15)h as the ratio of the expected demand from *i* to *j* and the expected demand of queue *i*. The expected external demand for queue *i* is defined in Eq. (15)i as the sum of the expected demand over all routes that start at queue *i*.

For the simulation model of the case studies, the set of route alternatives for a given OD pair is endogenous (i.e., it varies with  $\theta$  and across assignment iterations). Nonetheless, for the purpose of tractability a fixed (i.e., exogenous) route choice is considered for the analytical model. For every OD pair, a set of 10 routes is constructed. Details on the derivation of this choice set are given in Zhang et al. (2016b). Since the analytical model does not capture congestion-dependent route choice changes, these are captured by the polynomial component ( $\phi$  of Eq. (6)) of the metamodel. The use of an endogenous, iteration-dependent, route choice set could yield more accurate analytical results, yet at the cost of a higher computational burden.

In summary, the analytical approximation of  $\lambda_i(\theta)$  of Eq. (14) is obtained by evaluating the analytical network model defined by the system of nonlinear differentiable Eqs. (15). The exogenous parameters of this system of equations are:  $\theta, d_s, \mu_i, \gamma_i, \ell_i, l_i, v_i, S, Q, R, R_s, G_{ij}, H_i, T_i, \Psi_r$ . All other variables are endogenous variables that are obtained when solving the above system of equations.

As is illustrated with the case studies of Section 3, the proposed model (system of Eqs. (15)) works well for scenarios with various levels of congestion, including congested scenarios. This is remarkable given that the model does not account for the occurrence of spillbacks and their impact on the performance of upstream links. Nonetheless, the model accounts for the impact of the link's finite space capacity on the expected link delay. More specifically, the expected link travel time equation (Eq. (15)d) consists of the summation of expected link free-flow travel time and expected link delay. The delay term is based on the approximation of the expected number of vehicles on the link (Eq. (15)e) which assumes that each link has a finite space capacity (denoted  $\ell_i$ ). Hence, the impact of finite space capacity on the expected delay is accounted for.

Eq. (13) represents the system of Eqs. (15). The function  $h$  is of class  $C^\infty$ . There exists a variety of algorithms to efficiently solve this differentiable system of nonlinear equations. The dimension of the system of equations scales linearly with the number of links in the network and linearly with the number of OD pairs. This makes it a scalable model suitable for the calibration of large-scale networks.

This metamodel framework is not constrained to the use of the concrete analytical queueing-theoretic network model (15) but is compatible with a variety of other analytical network models. We consider this flexibility a strength of the proposed framework. The proposed analytical model is particularly efficient, since its evaluation consists of solving a system of nonlinear equations, the dimensionality of which scales linearly with the number of links in the network and linearly with the number of OD pairs. This makes it a scalable model suitable for the calibration of large-scale networks. The proposed framework can be used with other analytical network models. Ongoing work studies the use of traffic-theoretic network loading models for simulation-based optimization, such as the model of Osorio and Flötteröd (2014), which is a stochastic model consistent with Newell's deterministic simplified theory of kinematic waves (Newell, 1993). The use of approximate expressions of local, and path marginal cost functions (e.g., Ghali and Smith, 1995; Shen et al., 2007; Qian and Zhang, 2011; Lu et al., 2013) could be of interest. The main challenge in using more traditional traffic-theoretic models is to develop formulations that both: (i) have endogenous user-equilibrium assignment, and (ii) are computationally efficient for large-scale networks.

The modeling of route choice sets is an in general unresolved problem; this is so because route choice sets are not directly observable (e.g., Frejinger et al., 2009; Flötteröd and Bierlaire, 2013). We hence consider it adequate to deploy the simplest possible approach to route choice set generation (where the choice set is exogenously given) and to rely on the general-purpose polynomial term in the metamodel to absorb the resulting modeling error.

We briefly describe how the exogenous route choice set is derived. A more detailed description is provided in Zhang et al. (2016b). The exogenous set consists of 10 route alternatives per OD pair. We consider different behavioral parameter values. For a given value, we run a set of sequential assignment iterations, and extract the set of routes used by the simulator in the last assignment iteration. We group the set of routes extracted from the various behavioral parameter values. Then for a given OD pair, the final set of 10 routes is determined by selecting the set (of 10 routes) with maximal distance-based

**Table 1**  
Network attributes.

	Toy network	Berlin network
Number of links	6	24335
Number of nodes	6	11345
Number of OD pairs	1	3635
Expected demand (veh/h)	1400	172900

**Table 2**  
Experimental design.

	Toy network	Berlin network
Bounds for $\theta$ values (1/h), $[\theta_L, \theta_U]$	$[-60, 0]$	$[-60, 0]$
True $\theta$ values (1/h), $\theta^*$	$\{-5, -20, -55\}$	N.A.
Initial $\theta$ values (1/h), $\theta_0$	$\{0, -40, -60\}$	$\{0, -40, -60\}$
Computational budget	30	20
Simulation replications	5	10
Simulation assignment iterations	50	100
Total simulation assignment iterations per algorithmic run	7500	20000

overlap with the entire set of extracted routes. For a given network, this process is carried out once, prior to calibration. The route choice set is then kept fixed throughout the entire calibration process.

The use of an exogenous choice set contributes to the computational efficiency of the proposed approach. As will be discussed in Section 3.3.2, for the Berlin metropolitan network, the analytical model with exogenous route choice set yields an accurate approximation of the form of the simulation-based objective function. This highlights the negligible effect that the exogenous route choice set has on the analytical model's accuracy. Nonetheless, this observation is network- and problem-specific. A discussion on extensions of this framework to allow for the use of endogenous route choice sets is given in Section 4.

### 3. Case studies

#### 3.1. Experimental design

We apply the proposed approach to two case studies: a hypothetical toy network and a Berlin metropolitan network. The simulator used is MATSim (Horni et al., 2016). The main purpose of both case studies is to evaluate the added value of embedding within the calibration algorithm the problem-specific analytical structural information, which is provided by the analytical traffic model. For each case study, we compare the performance of two calibration approaches that only differ in the use or not of the analytical traffic model. All other algorithmic details are identical. The first approach is the proposed approach (denoted algorithm *Am*). It uses the metamodel defined by (14). The second approach (denoted algorithm *A $\phi$* ) considers a metamodel defined for iteration  $k$  and link  $i$  as:

$$\phi_{i,k}(\theta; \beta_{i,k}) = \beta_{i,k,1} + \beta_{i,k,2}\theta. \quad (16)$$

This metamodel differs from that of (14) in the absence of the macroscopic traffic model. In other words, compared to *Am*, *A $\phi$*  uses the same general-purpose metamodel component but has no problem-specific metamodel component.

The network topology characteristics of both networks are summarized in Table 1. The main details of the experimental design for each network are displayed in Table 2. The first row of Table 2 considers the lower and upper bounds for  $\theta$ , this defines the feasible region,  $\Theta$ . For the toy network, we consider a set of hypothetical  $\theta$  values, based on which we simulate synthetic traffic counts. The second row of the table displays these hypothetical values, which we refer to as the true values and denote  $\theta^*$ . For the Berlin network,  $\theta^*$  is unknown. Recall that  $\theta$  is the travel time coefficient of the route choice model. Hence its unit is the inverse of that of the travel time. In other words, the concrete value of  $\theta$  depends on the time unit used for travel time. For both networks, the travel times are computed in hours, hence the unit of  $\theta$  is  $\text{h}^{-1}$ .

For a given  $\theta^*$ , we initialize the SO algorithms with three different initial points, denoted  $\theta_0$ . The initial values used are displayed in row 3. Therefore, in total there are 9 different experiments for the toy network and 3 different experiments for the Berlin network. For each experiment, we run each SO algorithm (*Am* or *A $\phi$* ) 3 times. The need to run an algorithm multiple times for a given experiment is due to the stochasticity of the traffic simulator. Each algorithmic run consists of a maximum number of points ( $\theta$  values) to be evaluated. This is known as the computational budget or the sampling budget (displayed in row 4). Once this computational budget is reached, the algorithm is terminated.

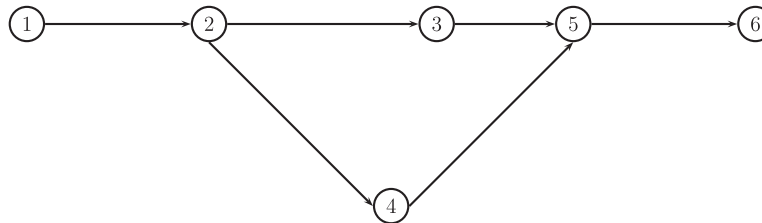
For a given  $\theta$  value, an estimate of the simulation-based objective function (Eq. (9)) is obtained by averaging over a set of independent simulation replications. The number of independent simulation replications is displayed in row 5. Each simulation replication consists of a sequential set of simulation assignment iterations (displayed in row 6). For a given simulation replication, the estimate of  $E[F_i(\theta; x)]$  is obtained by averaging the observations from the last 5 assignment iterations. Details



**Table 3**

Simulation runtime statistics in minutes.

	Toy network	Berlin network
Average runtime per assignment iteration	0.06	1.2
Average runtime per SO iteration	3	120
Average runtime per experiment	90	2400

**Fig. 1.** Toy network topology.**Table 4**

Toy network link properties.

Link index	Nodes connected	Length (km)	Maximum speed (km/h)	Signalized
1	1 → 2	2.5	72	un-signalized
2	2 → 3	7.5	54	signalized
3	3 → 5	2.5	54	un-signalized
4	2 → 4	7.1	72	un-signalized
5	4 → 5	7.1	72	un-signalized
6	5 → 6	2	72	un-signalized

on how the assignment iterations are initialized are given in Zhang et al. (2016b). The last row of the table indicates the total number of simulation assignment iterations per SO algorithmic run. For example, for the toy network, the computation of a single experiment (i.e., one run of the algorithm) requires a total of 7500 simulation assignment iterations: the performance of a total of 30  $\theta$  values are estimated, each estimation involves 5 independent replications, each of which requires 50 sequential assignment iterations. This leads to a total of  $30 \times 5 \times 50 = 7500$  simulation assignment iterations. Similarly, for the Berlin network each SO run involves  $20 \times 10 \times 100 = 20000$  simulation assignment iterations. In other words, we allow for a tight computational budget, which is defined as a small number of iterations of the calibration algorithm. This number is 30 (resp. 20) for the toy (resp. Berlin) network. Each iteration of the calibration algorithm, involves an estimation of the simulation-based optimization objective function ( $f$  of Eq. (9)). For each estimation, we carry out a set of sequential assignment iterations. Hence, each estimation involves calling the simulator  $5 \times 50 = 250$  (resp.  $10 \times 100 = 1000$ ) times.

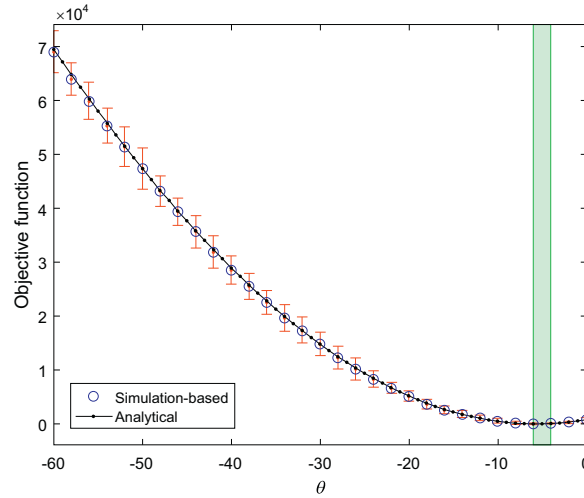
Table 3 displays the runtime statistics considering the 9 (resp. 3) experiments for the toy (resp. Berlin) network. For the toy network, a single assignment iteration takes an average of 0.06 min, leading to an average of 3 min per SO iteration and 90 min per SO algorithmic run (i.e., one experiment). For the Berlin network an assignment iteration averages 1.2 min, an SO iteration 120 min, and an algorithmic run 2400 min (i.e., 40 h). The toy network experiments are carried out on a standard laptop with a 4-core Intel i7-3740QM processor and 8GB RAM. The Berlin network experiments are carried out on a server with a 40-core Intel Xeon E5-2660 processor and 64GB RAM.

### 3.2. Toy network

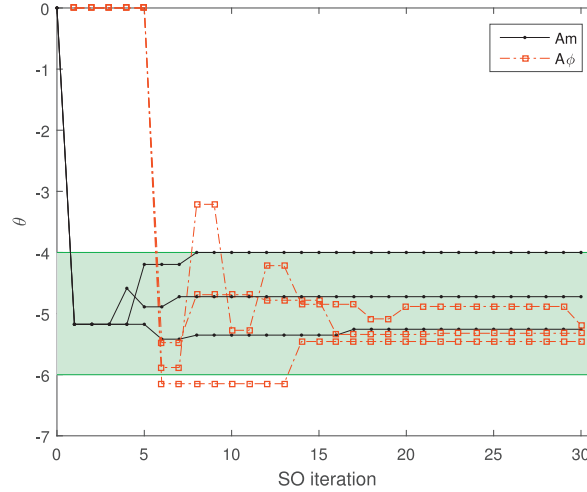
#### 3.2.1. Network attributes

For the hypothetical toy network, we pick a set of true  $\theta$  values,  $\theta^*$  (row 2 of Table 2). We use these values to generate the “real” traffic counts via simulation. The topology of the network is displayed in Fig. 1. Each link consists of a single lane road. Table 4 details the link properties.

The network has one OD pair (node 1 to node 6) and an expected demand of 1400 vehicles per hour. There are two alternative routes connecting the OD pair, a route to the north which goes through node 3 and a route to the south which goes through node 4. The northern route has a signal controlled intersection at node 3, whereas the southern route is un-signalized. The traffic signal control at node 3 is green for 75 seconds out of the 100 seconds cycle time. The free-flow travel time on the route to the north (resp. south) is approximately 14.9 (resp. 15.5) min. The free-flow travel time on the route to the north is shorter and hence it is preferred when there is no congestion. As congestion increases, the route to the south becomes increasingly attractive.



**Fig. 2.** Simulation-based and analytical objective functions for  $\theta^* = -5h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 3.** Algorithmic solutions versus iterations, for  $\theta^* = -5h^{-1}$  and  $\theta_0 = 0h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

### 3.2.2. Results

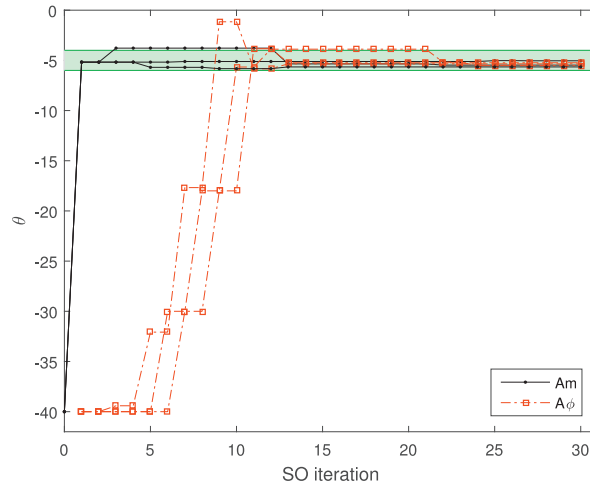
We first consider the experiment with  $\theta^* = -5h^{-1}$ . Fig. 2 displays the simulation-based objective function and the objective function approximation provided by the analytical traffic model (Eqs. (15)). The simulation-based objective function is defined by Eq. (9). The simulation-based estimates are displayed as blue circles. For a given point, the estimate is based on 5 simulation replications. The red error bars displayed are 95% confidence intervals. The black curve is the analytically approximated objective function, which is defined as:

$$\sum_{i \in \mathcal{I}} (y_i - \lambda_i(\theta))^2. \quad (17)$$

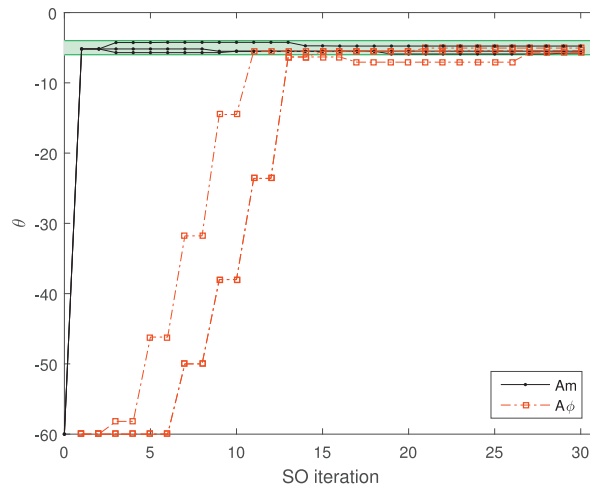
In other words, the analytical approximation is obtained by replacing the simulation-based metric (expected link flow) with the analytical metric. The black curve appears as an excellent approximation of the simulation-based objective function.

Fig. 2 also displays a green range for  $\theta$ . This range is a set of  $\theta$  values that have statistically equivalent objective function values to that of  $\theta^*$ . If an SO method yields  $\theta$  values within this range, we consider it to have converged. Statistical equivalence is tested with a paired  $t$ -test where the null hypothesis assumes equal expectations, while the alternative hypothesis assumes unequal expectations. For this experiment, the equivalent region is  $[-6, -4]$  (in units  $h^{-1}$ ).

Figs. 3–5 each considers a given initial point,  $\theta_0$ . The x-axis displays the iteration of the SO algorithm, the y-axis displays the current iterate (i.e., the best  $\theta$  value found so far by the algorithm). The solid black lines correspond to the proposed method,  $Am$ , while the dashed red lines correspond to the benchmark method,  $A\phi$ . Recall that the only difference in the



**Fig. 4.** Algorithmic solutions versus iterations, for  $\theta^* = -5h^{-1}$  and  $\theta_0 = -40h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 5.** Algorithmic solutions versus iterations, for  $\theta^* = -5h^{-1}$  and  $\theta_0 = -60h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

methods is their metamodel formulation, all other algorithmic details are identical for both methods. These figures also display the aforementioned equivalent region (in green).

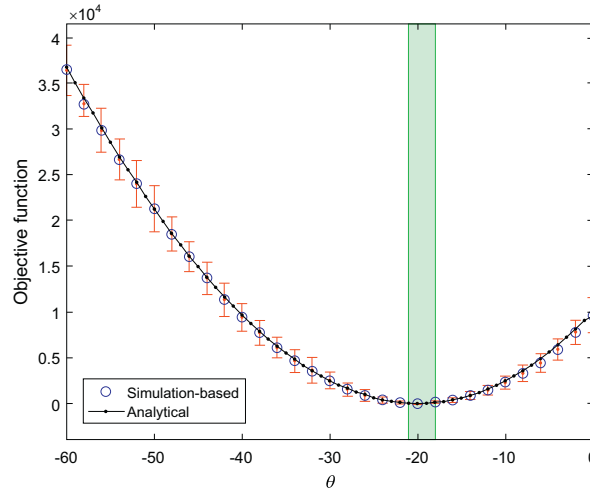
Note that in Fig. 5, 2 of the red curves overlap from iteration 1 until iteration 14, this occurs for the right-most red curve. For all 3 figures (Figs. 3–5), the following observations hold.

- All  $Am$  and all  $A\phi$  runs converge.
- All  $Am$  runs converge faster than the  $A\phi$  runs.

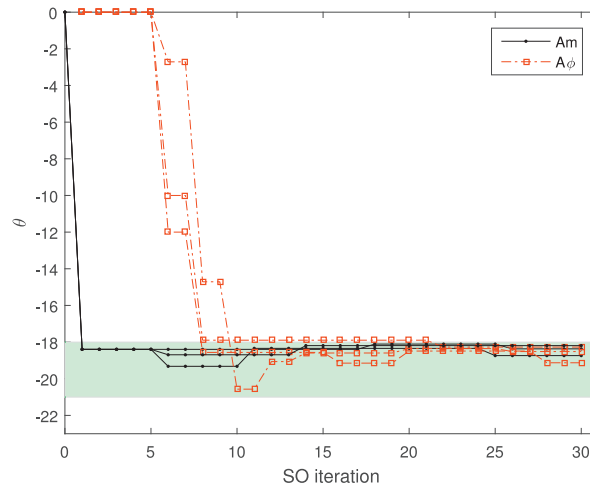
We now consider a true value of  $\theta^* = -20h^{-1}$ . The different objective functions are displayed in Fig. 6. Again, the analytical approximation of the objective function provided by the analytical network model (Eqs. (15)) is an excellent approximation of its simulation-based counterpart. The statistically equivalent region for  $\theta$  is  $[-21, -18]$  (in units  $h^{-1}$ ). The corresponding experimental results are shown in Figs. 7, 8 and 9. Note that in Fig. 9, all 3 red curves overlap from iteration 1 until iteration 8. For Figs. 7, 8 and 9, all  $Am$  runs converge, while 7 out of 9  $A\phi$  runs converge. Of the 2  $A\phi$  runs that do not converge, they both had current iterates within the green region, yet exited the region. For the converged runs, convergence tends to be faster for  $Am$  than for  $A\phi$ .

The different objective functions for the experiments with  $\theta^* = -55h^{-1}$  are displayed in Fig. 10. As before, the analytical objective function approximation is almost identical to the simulation-based objective function. The statistically equivalent region for  $\theta$  is  $[-57, -54]$  (in units  $h^{-1}$ ). The results for these experiments are displayed in Figs. 11–13.

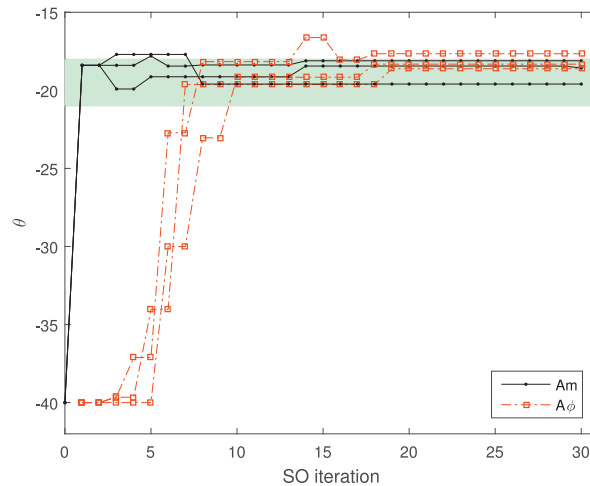
Note that in Fig. 11, 2 of the red curves that overlap from iteration 1 until iteration 11, this occurs for the left-most red curve. As for the previous experiments, all  $Am$  runs converge. Eight out of the nine  $A\phi$  runs converge.



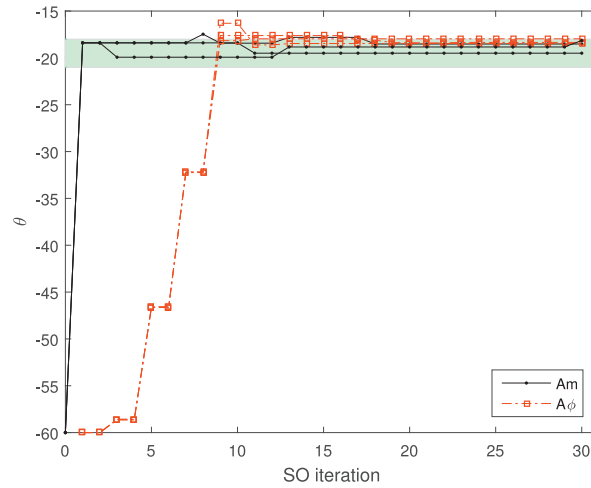
**Fig. 6.** Simulation-based and analytical objective functions for  $\theta^* = -20h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



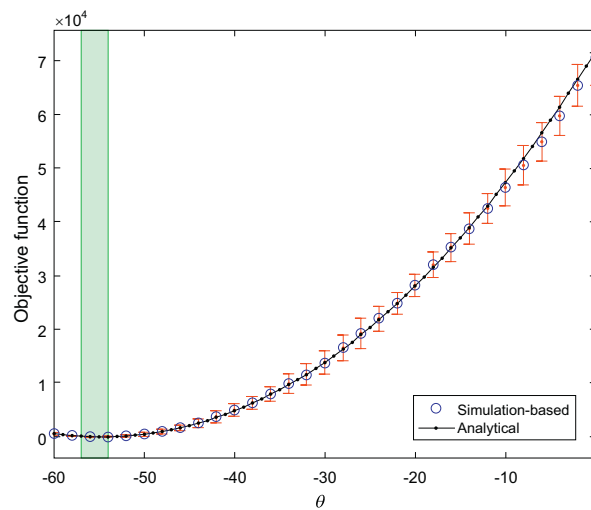
**Fig. 7.** Algorithmic solutions versus iterations, for  $\theta^* = -20h^{-1}$  and  $\theta_0 = 0h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



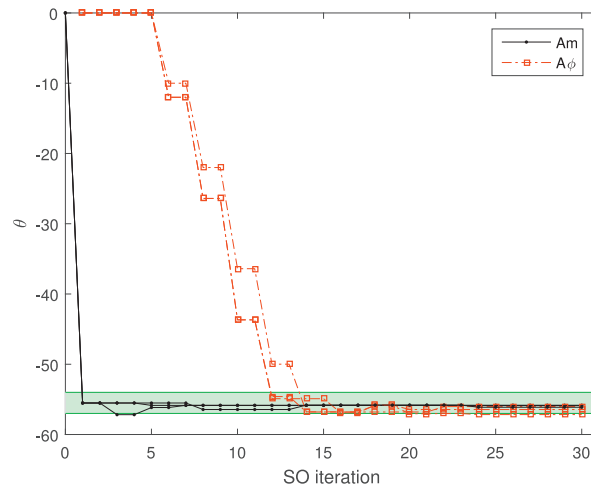
**Fig. 8.** Algorithmic solutions versus iterations, for  $\theta^* = -20h^{-1}$  and  $\theta_0 = -40h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



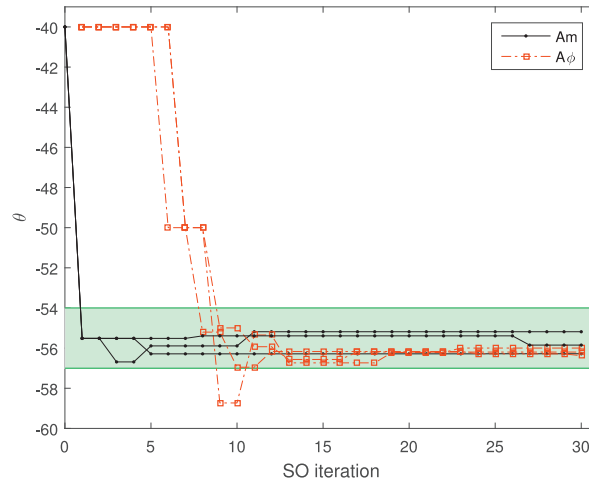
**Fig. 9.** Algorithmic solutions versus iterations, for  $\theta^* = -20h^{-1}$  and  $\theta_0 = -60h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



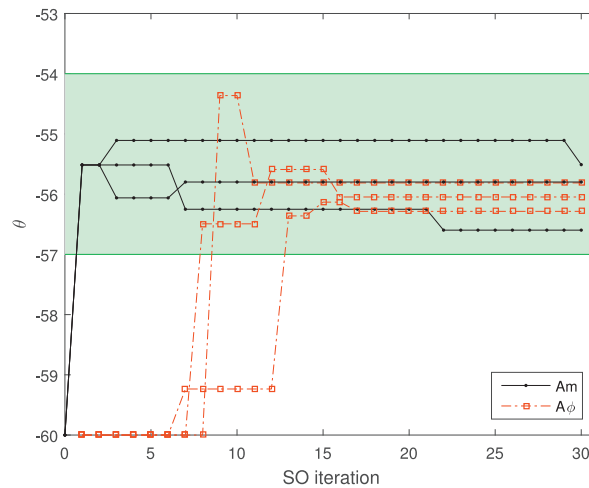
**Fig. 10.** Simulation-based and analytical objective functions for  $\theta^* = -55h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 11.** Algorithmic solutions versus iterations, for  $\theta^* = -55h^{-1}$  and  $\theta_0 = 0h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 12.** Algorithmic solutions versus iterations, for  $\theta^* = -55h^{-1}$  and  $\theta_0 = -40h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 13.** Algorithmic solutions versus iterations, for  $\theta^* = -55h^{-1}$  and  $\theta_0 = -60h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

**Table 5**

Number of algorithmic iterations until convergence for the toy network.

		<i>Am</i>			<i>Aφ</i>		
$\theta^* = -5$	$\theta_0 = 0$	1	1	1	10	6	14
	$\theta_0 = -40$	1	13	1	13	22	11
	$\theta_0 = -60$	1	1	1	11	27	15
$\theta^* = -20$	$\theta_0 = 0$	1	1	1	22	8	10
	$\theta_0 = -40$	8	6	1	7	30	10
	$\theta_0 = -60$	1	18	1	17	11	30
$\theta^* = -55$	$\theta_0 = 0$	1	1	5	30	12	14
	$\theta_0 = -40$	1	1	1	11	9	8
	$\theta_0 = -60$	1	1	1	9	13	8

From the perspective of the values proposed by the algorithms for the calibration parameter,  $\theta$ , both SO methods yield values with good performance and can do so within few algorithmic iterations. Overall the performance of both methods is similar and good. Overall, *Am* identifies good solutions faster than *Aφ*, and systematically converges.

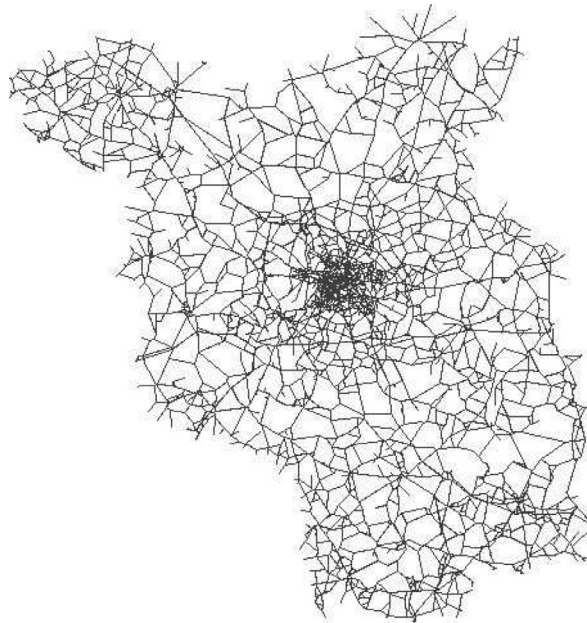
Let us compare the performance of both methods in terms of their computational efficiency. Table 5 considers all experiments mentioned above. It displays for each experiment and each method, the number of algorithmic iterations until convergence (i.e., the first time the equivalent region is entered and not exited thereafter). If a method upon termination



**Table 6**

Simulation runtimes until convergence for the toy network (minutes).

		<i>Am</i>			<i>A<math>\phi</math></i>		
$\theta^* = -5$	$\theta_0 = 0$	3	3	3	30	18	42
	$\theta_0 = -40$	3	39	3	39	66	33
	$\theta_0 = -60$	3	3	3	33	81	45
$\theta^* = -20$	$\theta_0 = 0$	3	3	3	66	24	30
	$\theta_0 = -40$	24	18	3	21	90	30
	$\theta_0 = -60$	3	54	3	51	33	90
$\theta^* = -55$	$\theta_0 = 0$	3	3	15	90	36	42
	$\theta_0 = -40$	3	3	3	33	27	24
	$\theta_0 = -60$	3	3	3	27	39	24

**Fig. 14.** Metropolitan Berlin simulation network.

of the algorithm (i.e., iteration 30) has not converged, then we indicate a value equal to the computational budget (i.e., the maximum number of iterations) of 30. This underestimates the convergence statistics for the non-converged experiments.

This table indicates that for *Am* and for  $\theta^* = -5h^{-1}$ , 8 out of the 9 experiments converge after 1 iteration. For  $\theta^* = -20h^{-1}$ , this happens 6 out of 9 times, and for  $\theta^* = -55h^{-1}$ , this happens 8 out of 9 times. At the first algorithmic iteration, only a single objective function estimate is available (the estimate obtained at the initial value  $\theta_0$ ). Hence, this instantaneous convergence is due to the information provided by the analytical traffic model. When considering all true values (i.e., all 27 experiments), *Am* converges on average at iteration 2.7, and in the worst case at iteration 18. The method *A $\phi$*  converges on average at iteration 14.4 (this average includes the non-converged cases, where we input an iteration value of 30), at best at iteration 6, and does not converge for 3 experiments. The corresponding numerical values of the solutions derived by each method for each experiment are presented in Table 11 of the Appendix C.

Table 6 analyzes the convergence of the methods in terms of their simulation runtimes until convergence. The runtime is in minutes. As in the above analysis, if a method has not converged, we indicate the total simulation runtime used until the algorithm was terminated. This underestimates the runtime needed until convergence. The method *Am* converges on average within 8 min of simulation, while *A $\phi$*  converges on average in 43.1 min. By providing the algorithm with the analytical information, we can converge with an average 81.4% reduction in simulation runtime. This table highlights the computational efficiencies that are achieved when using *Am*.

### 3.3. Berlin metropolitan network

#### 3.3.1. Network attributes

Fig. 14 displays the metropolitan Berlin network topology, as modeled in the simulator. As a reference, we also display here the road map of the corresponding network (Fig. 15). This network represents the metropolitan area of Berlin, Germany.



Fig. 15. Metropolitan Berlin road network. Map source: <http://www.mapsofworld.com/germany/states/brandenburg>, downloaded on 10/30/2014.

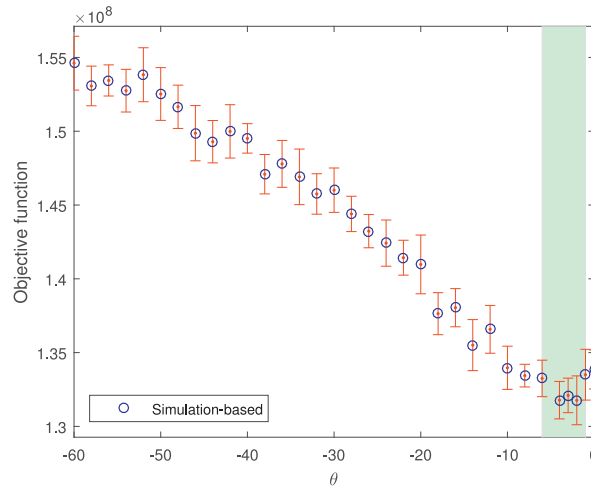


Fig. 16. Simulation-based objective function for the Berlin network. (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

The area includes the city of Berlin and the broader federal state of Brandenburg. It consists of 24335 links, 11345 nodes and 3635 OD pairs. Real traffic counts from sensors on 346 links are available per hour for 24 h. We focus on the morning peak hour: 8–9am, during which expected demand is 172900 vehicles. For more data and model details, see Ziemke et al. (page 120, Section “Counts”, 2015) and Ziemke (pages 57–59, 2013).

### 3.3.2. Results

Fig. 16 displays the simulation-based objective function estimates with corresponding 95% confidence intervals obtained from 20 simulation replications. As for the toy network, the green region represents the range of statistically equivalent  $\theta$  values, which is  $[-6, -1]$  (in units  $h^{-1}$ ). Based on the finite and small set of simulated points, the  $\theta$  value at which the minimum objective function is obtained is near  $-2h^{-1}$ . The analytical approximation of the objective function derived from the analytical network model (Eqs. (15)) is displayed in Fig. 17. Both functions have similar form and seem to be non-convex. Note that both figures differ in the y-axis limits. Hence, the analytical function approximates well the form of the simulation-based function, yet is not scaled properly. The scaling is corrected by the metamodel (term  $\beta_{i,k,0}$  of Eq. (14)). The minimum of the analytical objective function is obtained at  $-5h^{-1}$ , which is in the green statistically equivalent region.

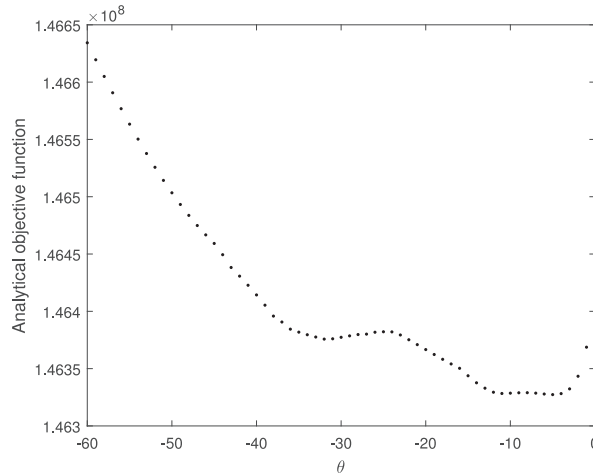


Fig. 17. Analytical objective function for the Berlin network.

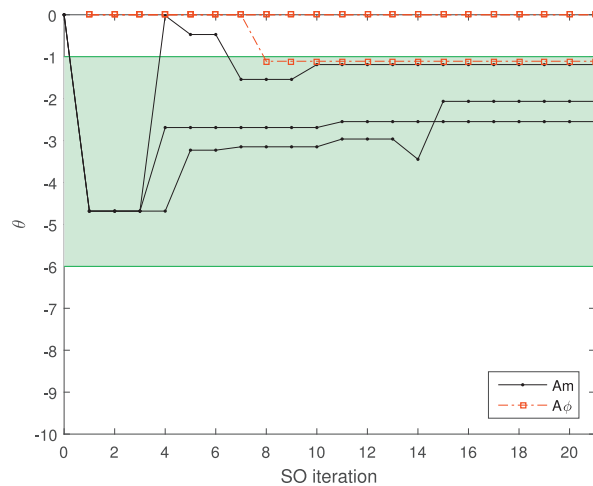


Fig. 18. Algorithmic solutions versus iterations, for  $\theta_0 = 0h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

We proceed as for the toy network: we run experiments for 3 different initial values  $\theta_0$ . For each initial value, we plot the current iterate (i.e., best  $\theta$  value identified so far by the algorithm) versus algorithmic iteration. We do this for both SO methods:  $Am$  and  $A\phi$ . The results are presented in Figs. 18, 19 and 20.

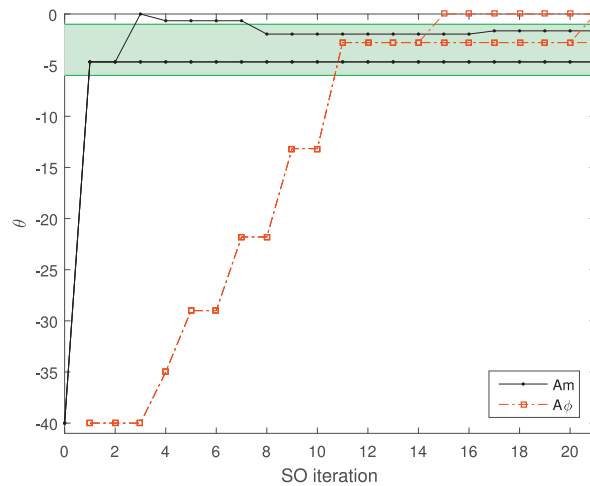
In these figures, some curves overlap: in Fig. 18, two dashed red lines overlap at  $\theta = 0h^{-1}$  for all iterations; in Fig. 19, two solid black lines overlap at  $\theta = -4.7h^{-1}$ .

In Fig. 18, all  $Am$  runs converge, 1 out of 3  $A\phi$  runs converges, the remaining 2 stay at the initial  $\theta$  value of  $0h^{-1}$ . In Fig. 19, all  $Am$  runs converge, 1 out of the 3  $A\phi$  runs converges, the remaining 2 reach and stay at the lower bound value for  $\theta$  equal to  $0h^{-1}$ . In Fig. 20, all  $Am$  runs converge, none of the  $A\phi$  runs converges, 2 reach the lower bound and the third stays at the value of  $-0.4h^{-1}$ .

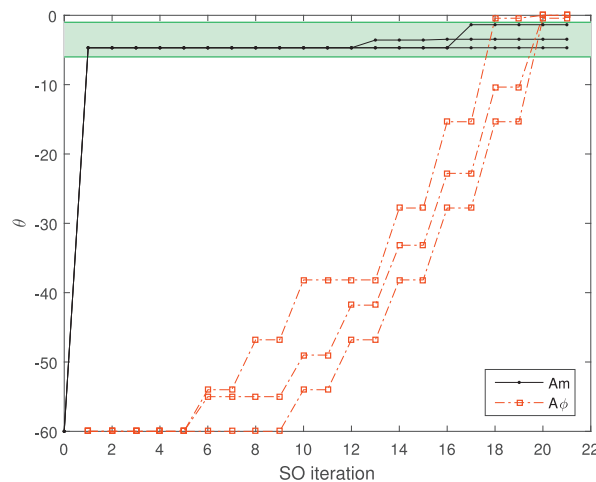
For all 3 figures, the following observations hold.

- All runs of  $Am$  converge. Seven out of the nine runs converge at the first iteration. This is thanks to the analytical network model approximating well the simulation-based objective function.
- Only 2 out of the 9 runs of  $A\phi$  converge. Of the non-converged runs, two stay at the initial value of  $\theta = 0h^{-1}$  (Fig. 18), 4 leave the initial value, reach and stay at the lower bound value of  $\theta = 0h^{-1}$  (Figs. 19 and 20).
- When comparing the converged runs: all  $Am$  runs converge faster than the  $A\phi$  runs.

We analyze the convergence statistics of the two methods. We proceed as for the toy network. Table 7 considers all Berlin experiments mentioned above. It displays for each experiment and each method, the number of algorithmic iterations until convergence (i.e., the first time the equivalent region is entered and not exited thereafter). If a method upon termination of the algorithm (i.e., iteration 20) has not converged, then we indicate a value equal to the computational budget, which is



**Fig. 19.** Algorithmic solutions versus iterations, for  $\theta_0 = -40\text{h}^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 20.** Algorithmic solutions versus iterations, for  $\theta_0 = -60\text{h}^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

**Table 7**

Number of algorithmic iterations until convergence for the Berlin network.

	<i>Am</i>			<i>A<math>\phi</math></i>		
$\theta_0 = 0$	1	7	1	20	20	8
$\theta_0 = -40$	1	1	8	20	20	11
$\theta_0 = -60$	1	1	1	20	20	20

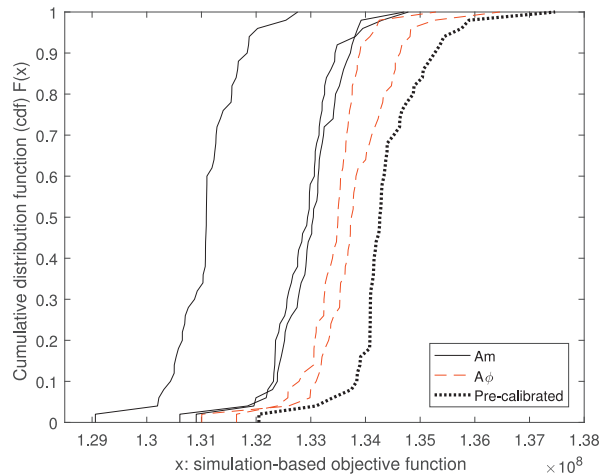
20. This underestimates the convergence statistics for the non-converged experiments. This table indicates that method *Am* converges at iteration 1 for 7 out of the 9 experiments. On average it converges at iteration 2.4 and at worst at iteration 8. Method *A $\phi$*  only converges 2 out of the 9 runs. For those two converged experiments, convergence is achieved at iterations 8 and 11, respectively. The corresponding numerical values of the solutions derived by each method for each experiment are presented in Table 12 of the Appendix C.

The corresponding statistics in terms of simulation runtimes are displayed in Table 8. If a given method has not converged upon termination of the algorithm, then we indicate the total simulation runtime upon termination (2400 min). As before, this underestimates convergence runtime for method *A $\phi$*  which often does not converge. On average method *Am* converges within 293.3 min (i.e., 4.9 h), while *A $\phi$*  does so within 2120 min (i.e., 35.3 h). Method *Am* achieves on average a 86.2% reduction in runtime until convergence, which corresponds in this case study to average savings of 30.4 h of simulation per experiment.

**Table 8**

Simulation runtimes until convergence for the Berlin network (minutes).

	$Am$		$A\phi$			
$\theta_0 = 0$	120	840	120	2400	2400	960
$\theta_0 = -40$	120	120	960	2400	2400	1320
$\theta_0 = -60$	120	120	120	2400	2400	2400



**Fig. 21.** Objective function distributions for each solution of each method, when initialized with  $\theta_0 = 0h^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

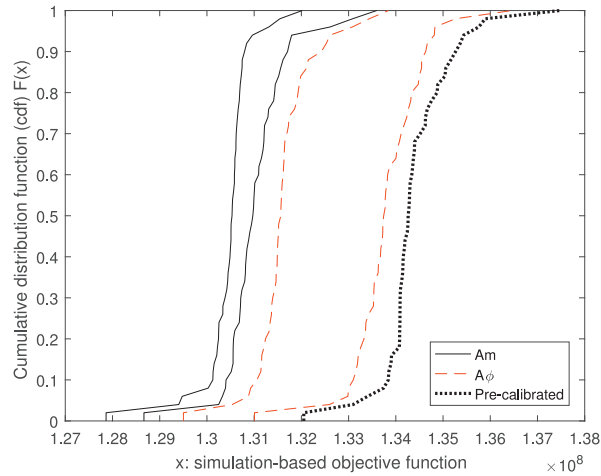
We now compare the performance of the best solutions proposed by each method with the performance of the value currently used as part of the Berlin model, the latter was obtained from prior calibration efforts (Ziemke et al., 2015). Let  $\tilde{\theta}$  denote this pre-calibrated value. It has a value of  $-6h^{-1}$ .

Recall, that for a given initial value,  $\theta_0$ , we ran each method 3 times, this leads to 3 solutions. We now compare the performance of these solutions with that of  $\tilde{\theta}$ . In order to evaluate the performance of a given  $\theta$  value, we run 50 simulation replications and obtain 50 observations of the objective function. More specifically, for a given  $\theta$  value, we run 100 assignment iterations to obtain 1 objective function estimate. We repeat this process 50 times to obtain 50 independent simulation observations (or estimates).

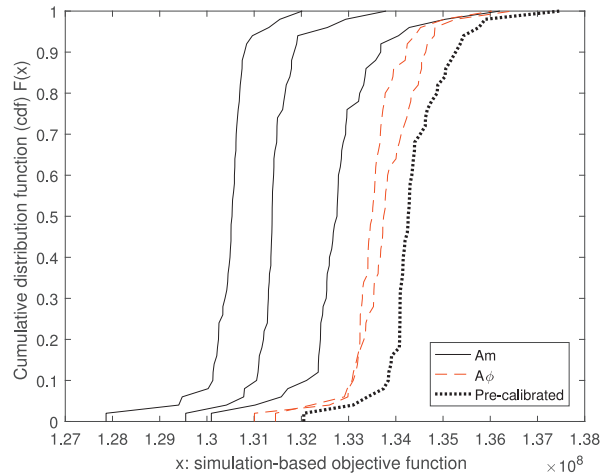
Fig. 21 displays the cumulative distribution function (cdf) of these 50 observations. Each curve corresponds to a given  $\theta$  value. The solid black (resp. dashed red) curves correspond to values derived by  $Am$  (resp.  $A\phi$ ). The dotted black curve corresponds to the pre-calibrated value  $\tilde{\theta}$ . The more a curve is shifted to the left, the lower the objective function estimates, i.e., the better its performance. More specifically, for a given  $x$ -value, the corresponding  $y$ -value on the cdf curve gives the proportion of observations (out of the 50) that have objective function values smaller than  $x$ . Fig. 21 considers the solutions obtained with an initial value of  $\theta_0 = 0h^{-1}$ . 2 out of the 3 runs of  $A\phi$  yield the same solution (i.e., the algorithm considers for all iterations the initial value,  $\theta_0 = 0h^{-1}$ , as the best value). Thus, only 2 dashed red curves are visible. The right-most dashed red curve represents the 2 identical solutions (i.e.,  $\theta = 0h^{-1}$ ). All 3 values derived by  $Am$  outperform the 3 values obtained by  $A\phi$ . They also outperform the pre-calibrated value.

Fig. 22 displays the cdf of each solution obtained when initializing the algorithms with  $\theta_0 = -40h^{-1}$ . Here as well, 2 out of the 3 runs of  $A\phi$  yield the same solution of  $0h^{-1}$ . Hence, only 2 dashed red curves are visible, one of which represents 2 identical solutions. The proposed method  $Am$  also yields 2 solutions that are similar (both are approximately  $-4.7h^{-1}$ ). The same conclusions as above hold: all solutions derived by  $Am$  outperform all solutions derived by  $A\phi$  and outperform the pre-calibrated value. Fig. 23 displays the cdf of each solution obtained when initializing the algorithms with  $\theta_0 = -60h^{-1}$ . The method  $A\phi$  yields a solution of  $0h^{-1}$  for 2 out of the 3 runs. Hence, only 2 distinct dashed red curves are visible. For this figure, the same conclusions as above hold.

For each initial value  $\theta_0$  and each method, we choose only one proposed solution which is defined as that with the smallest objective function average (the average is over the 50 simulation replications). In Table 9, we give the numerical values of the best proposed solutions for the three sets of experiments. Fig. 24 presents the cdf of the best solutions for all three initial values. It also displays the cdf of pre-calibrated value and of the value obtained by solving the problem with the analytical traffic model only (i.e., no simulation). The solid black curves correspond to  $Am$  solutions. Only two curves appear, because  $Am$  yields the same solution ( $\theta = -4.7h^{-1}$ ) for two initial values ( $\theta_0 = -40h^{-1}$  or  $\theta_0 = -60h^{-1}$ ). Hence, the left-most solid black curve represents two solutions. The dashed red curves correspond to the three solutions of  $A\phi$ . The



**Fig. 22.** Objective function distributions for each solution of each method, when initialized with  $\theta_0 = -40\text{h}^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)



**Fig. 23.** Objective function distributions for each solution of each method, when initialized with  $\theta_0 = -60\text{h}^{-1}$ . (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

**Table 9**

Best solution of each method and each initial  $\theta$  value.

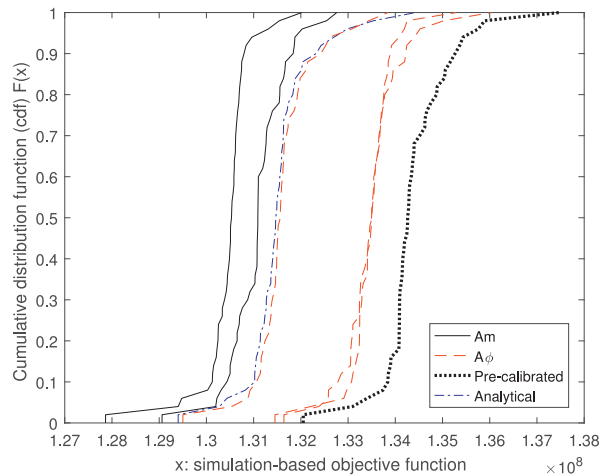
	$Am$	$A\phi$
$\theta_0 = 0$	-2.1	-1.1
$\theta_0 = -40$	-4.7	-2.8
$\theta_0 = -60$	-4.7	-0.4

dash-dotted blue curve corresponds to the solution of the analytical traffic model (i.e., no simulation) and the dotted black curve corresponds to the pre-calibrated value.

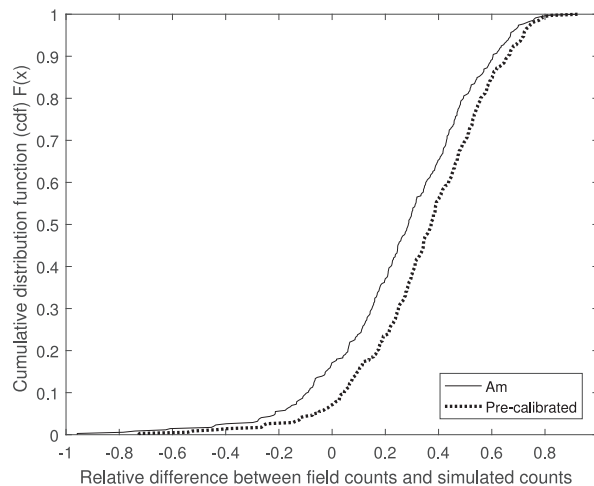
Fig. 24 indicates that the solutions proposed by  $Am$ , for each of the three initial values, outperform all 3 solutions proposed by  $A\phi$ . It also outperforms the pre-calibrated value. This figure also shows that the analytical solution outperforms all  $A\phi$  solutions. This indicates the added value of the analytical network model information. All  $Am$  solutions outperform the analytical solution. This indicates the added value of combining analytical and simulation-based information, rather than merely solving the problem with an analytical-only approach.

Recall that the traffic field data consists of traffic counts from sensors on 346 links. Fig. 25 compares the performance of the best  $Am$  solution ( $\theta = -4.7\text{h}^{-1}$ , the left-most curve in Fig. 24) and the pre-calibrated solution. For each solution and each sensor location, we compute the relative error of the count:  $(y_i - \hat{E}[F_i(\theta; x)])/y_i$ , where  $y_i$  represents the field count at location  $i$  and  $\hat{E}[F_i(\theta; x)]$  represents the simulated estimate of the count at that location. Fig. 25 displays two curves, one for





**Fig. 24.** Objective function distributions for the best solutions of each method and each initial  $\theta$  value. (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

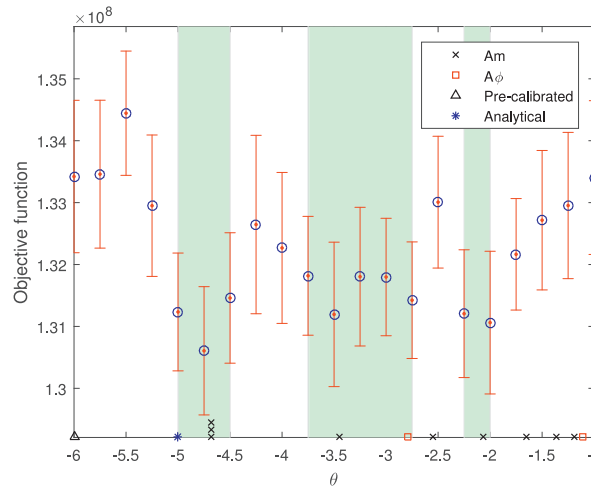


**Fig. 25.** Empirical cdf of the relative difference of traffic counts.

each solution ( $Am$  or pre-calibrated). The curve is the cumulative distribution function of the relative error. The distribution is over all 346 links. Let us illustrate how to interpret these curves. The solid black curve corresponds to method  $Am$ . A vertical line through  $x = 0$  intersects the  $Am$  curve at  $y = 0.17$ . This means that under the  $Am$  solution, 17% of the sensor locations yield a negative relative difference, i.e., they overestimate 17% of the counts. The pre-calibrated curve intersects this vertical line at  $y = 0.07$ , meaning that 7% of the locations overestimate the counts. Similarly, a horizontal line through  $y = 0.6$  intersects the  $Am$  curve at  $x = 0.35$ . This indicates that 60% of the relative errors are below 0.35. This horizontal line intersects the pre-calibrated curve at  $x = 0.44$ , which indicates that 60% of the relative errors are below the  $x = 0.44$ . From this figure we can deduce that 31% of the counts have a relative error within  $[-0.2, 0.2]$  under the  $Am$  value, while this is the case of only 21% of the counts under the pre-calibrated value.

We now carry out a more detailed analysis of the performance of the different points within the green equivalent region. Fig. 26 considers the equivalent region  $[-6, -1]$  (in units  $h^{-1}$ ). It displays a more detailed simulation-based estimate of the objective function. The function is now sampled with a step size of 0.25 (i.e., a smaller step size than in Fig. 16). Each estimate is obtained as an average over 30 simulation replications. Its corresponding 95% confidence interval is displayed. Within this interval  $[-6, -1]$  (in units  $h^{-1}$ ), we compute a new green equivalent region. We identify the point with the smallest objective function estimate. This is  $\hat{\theta} = -4.75h^{-1}$ . We conduct for each point, a paired  $t$ -test to test whether its performance (i.e., objective function value) is equivalent to that of  $\hat{\theta}$ . If two adjacent points have statistically equivalent performance to that of  $\hat{\theta}$ , then both points belong to the new equivalent region. The resulting green region is now defined (in units  $h^{-1}$ ) as:  $[-5, -4.5]$ ,  $[-3.75, -2.75]$ , and  $[-2.25, -2]$ . This new region is displayed in green in Fig. 26.

Fig. 26 indicates along the  $x$ -axis the location of: the 9 solutions proposed by  $Am$  (all 9 solutions fall in the initial green equivalent region of  $[-6, -1]$  (in units  $h^{-1}$ ); they are represented by black crosses), 2 solutions proposed by  $A\phi$  (2 solutions fall in  $[-6, -1]$  (in units  $h^{-1}$ ); they are represented by red squares), the pre-calibrated value  $\hat{\theta}$  (represented by a



**Fig. 26.** More detailed estimation of the simulation-based objective function in the region  $[-6, -1]$  (in units  $\text{h}^{-1}$ ). (For interpretation of the references to colour in this figure legend, the reader is referred to the web version of this article.)

black triangle), and the solution obtained using only the analytical traffic model (i.e., no simulation-based optimization; it is represented by a blue asterisk). Note that 3 of the  $Am$  solutions overlap at the value  $-4.7\text{h}^{-1}$ , hence they are displayed as stacked upon each other on the x-axis.

Considering this more accurate estimation of the equivalent region, we observe that 5 out of the 9 solutions of  $Am$  fall within this region, while only 1 solution of  $A\phi$  is in the region. The solution using only the analytical traffic model also falls in this region.

The scenarios of Section 3 all consider high levels of congestion. For instance, in the toy network the expected demand is 1400 veh/h, and the expected supply of links 1 and 6 is 1200 veh/h. Hence, the ratio of expected demand to expected supply is 1.17. The analytical model describes the within link build-up of congestion but does not capture the occurrence and impact of vehicular spillbacks. Nonetheless, as shown in Figs. 2, 6 and 10, it provides an accurate approximation of the simulation-based objective function. The Berlin network consists of a set of links with varying levels of congestion: ranging from uncongested to highly congested. For instance, the city center contains approximately 7% of links with a ratio of expected demand to expected supply that is greater than 1. For the entire metropolitan network, 2% of the links have a ratio greater than 1.

#### 4. Conclusions and discussion

In this paper, we propose a computationally efficient calibration algorithm. Efficiency is achieved by providing analytical problem-specific information to the algorithm. We formulate a metamodel that embeds information from an analytical differentiable and tractable network model. The analytical network model provides an analytical description of how the calibration parameter is related to the objective function. The performance of the method is evaluated with case studies for both a hypothetical toy network and for the Berlin metropolitan area network. The performance of the proposed method is compared to that of an algorithm that differs only in that the information from the analytical network model is not provided to it. For both networks, the proposed approach significantly improves the computational efficiency of the calibration algorithm with an average reduction in the simulation runtime until convergence of more than 80%. The simulator used in both case studies of this paper, MATSim, is computationally efficient compared to other higher-resolution simulators. Inefficient simulators are typically calibrated under tight computational budgets. By accelerating the convergence of traditional black-box calibration algorithms, the proposed approach is of particular interest for the calibration of inefficient simulators.

This metamodel framework can be extended to address any continuous calibration problem. It can be extended to allow for other types of demand or supply calibration parameters, as well as to allow for other types of traffic measurements to be used (e.g., speeds or occupancies). For a given calibration problem formulation, one of two types of approaches can be followed. One approach is to extend the analytical model (system of Eqs. (15)) to describe the relationship between any new calibration parameters and any traffic measurements. This involves formulating an analytical and differentiable approximation of the mapping between the calibration parameters and the traffic measurements. As part of ongoing work, this has been formulated and is being tested for several problems, including: (i) the supply calibration of various types of traffic simulators (Zhang et al., 2016a), (ii) calibration of origin-destination matrices. For instance, the latter uses an endogenous representation of the external travel demand (i.e.,  $\gamma$  of Eq. (15)g becomes endogenous). Second, if the extension of the analytical model deems computationally inefficient, then the current analytical traffic model can be used as is, and the additional parameters can be included in the polynomial error term of the metamodel. In this second case, since the dimension

of the system of equations scales linearly with the number of links in the network and linearly with the number of OD pairs, higher-dimensional problems can be readily addressed. In the first case, the use for higher-dimensional problems will depend on the scalability of the new traffic model formulation.

A generalization to other data sources and parameters to be calibrated is part of the planned research effort. Ongoing work is also investigating the use of these metamodel ideas to enhance the efficiency and robustness of traditional data-driven calibration algorithms (e.g., Kalman and particle filters).

The analytical route choice model (Eq. (15)b) is a multinomial logit (MNL) model. This calibration algorithm can be used for a simulator that embeds a different route choice model specification by replacing Eq. (15)b with an analytical and differentiable approximation of the simulator's route choice model. If such an approximation is not available or is not computationally efficient, then one can continue to use the MNL formulation of this paper. In this case, the polynomial error term of the metamodel will capture the effects of this specification difference between the analytical and the simulation-based models.

In this paper, the simulation-based optimization algorithm used is derivative-free, i.e., it does not require the evaluation of derivatives of the simulation-based objective function. Derivative-based algorithms have been proposed for calibration problems (e.g., Yang, 1995; Antoniou et al., 2011; Balakrishna, 2006). Such algorithms have been traditionally designed to achieve asymptotic properties. There is now extensive ongoing and recent work that formulates their extensions for use under tight computational budgets (Lu et al., 2015; Tympakianaki et al., 2015). Another area of ongoing research aims to achieve efficiency through the combined use of multiple models with varying efficiency-accuracy trade-offs (e.g., Corthout et al., 2014; Osorio and Selvam, 2016).

The analytical model of the proposed approach is based on the use of an exogenous route choice set. This set is computed once prior to calibration. This contributes to the computational efficiency of the calibration algorithm. For cases where the use of an endogenous route choice set is desirable, one efficient approach would be to use at each iteration of the calibration algorithm, the route choice set used by the simulator for the current iterate (i.e., the calibration vector value considered as best so far). This would require, at every iteration of the calibration algorithm, to compute a smaller set of routes to be used by the analytical model (this can be done as described in Section 2.3), and then to compute the exogenous parameters of the analytical model that depend on the route choice set. This can be done efficiently. This would allow the route choice set used to solve Problem (7)–(8) to vary across iterations of the calibration algorithm. However, for a given iteration, Problem (7)–(8) would still be solved assuming a fixed route choice set.

## Acknowledgments

The work of C. Zhang and C. Osorio is partially supported by the U.S. National Science Foundation under Grant No. 1334304. Any opinions, findings, and conclusions or recommendations expressed in this material are those of the authors and do not necessarily reflect the views of the National Science Foundation. The authors thank Prof. Kai Nagel and his group at TU Berlin for the data. In particular, Dominik Ziemke provided modeling and data help. The iMobility laboratory of the KTH Royal Institute of Technology provided access to computing resources used to generate parts of the experimental results.

## Appendix A

This appendix presents the SO algorithm. The algorithm is described using notation of Osorio and Bierlaire (2013). The following notations are defined for a given iteration  $k$ : current iterate as  $x_k$ , trust region radius as  $\Delta_k$ , metamodel coefficient vector as  $v_k$ , metamodel as  $m_k$ , total number of simulation runs carried out up until and including iteration  $k$  as  $n_k$ , and total number of successive trial points rejected as  $\mu_k$ .

The constants  $\eta_1$ ,  $\bar{\gamma}$ ,  $\gamma_{inc}$ ,  $\bar{\tau}$ ,  $\bar{d}$ ,  $\bar{\mu}$ ,  $\Delta_{max}$  are given such that:  $0 < \eta_1 < 1$ ,  $0 < \bar{\gamma} < 1 < \gamma_{inc}$ ,  $0 < \bar{\tau} < 1$ ,  $0 < \bar{d} < \Delta_{max}$ ,  $\bar{\mu} \in \mathbb{N}^*$ . Set the total number of simulation runs permitted,  $n_{max}$ , this determines the computational budget. Set the number of simulation replications per point  $\bar{r}$ .

### 0. Initialization.

Set  $k = 0$ ,  $n_0 = 0$ ,  $\mu_0 = 0$ . Determine  $x_0$  and  $\Delta_0$  ( $\Delta_0 \in (0, \Delta_{max}]$ ). Given the initial point  $x_0$ , compute  $f_A(x_0)$  (analytical approximation of Eq. (4) or equivalently of Eq. (9)) and  $\hat{f}(x_0)$  (simulated estimate of Eq. (4) or equivalently of Eq. (9)).

### 1. Analytical-only calibration.

Solve a Problem (11)–(13) using only the analytical network model and without using any simulation information, i.e., set the metamodel equal to  $f_A$ . Let  $x_1$  denote the solution to this problem, compute  $f_A(x_1)$  and  $\hat{f}(x_1)$ . Fit the metamodel  $m_1$ .

### 2. Step calculation.

Solve Problem (11)–(13), and find a solution, denoted  $x_k + s_k$  and referred to as the trial point, that lies in the trust region (i.e.,  $\|s_k\| \leq \Delta_k$ ).

### 3. Acceptance or rejection of the trial point.

Compute  $\hat{f}(x_k + s_k)$  and

$$\rho_k = \frac{\hat{f}(x_k) - \hat{f}(x_k + s_k)}{m_k(x_k) - m_k(x_k + s_k)} \quad (18)$$

- If  $\rho_k \geq \eta_1$  and  $\hat{f}(x_k) - \hat{f}(x_k + s_k) > 0$ , then accept the trial point:  $x_{k+1} = x_k + s_k$ ,  $\mu_k = 0$ ;
- Otherwise, reject the trial point:  $x_{k+1} = x_k$ ,  $\mu_k = \mu_k + 1$ .

Include the new observation in the set of sampled points ( $n_k = n_k + \bar{r}$ ), and fit the new metamodel  $m_{k+1}$ .

#### 4. Model improvement.

Compute  $\tau_{k+1} = \frac{\|v_{k+1} - v_k\|}{\|v_k\|}$ . If  $\tau_{k+1} < \bar{\tau}$ , then improve the model by simulating the performance of a new point  $x$ , which is uniformly and randomly drawn from the feasible space. Evaluate  $f_A$  and  $\hat{f}$  at  $x$ . Include this new observation in the set of sampled points ( $n_k = n_k + \bar{r}$ ). Update metamodel  $m_{k+1}$ .

#### 5. Trust region radius update.

$$\Delta_{k+1} = \begin{cases} \min\{\gamma_{inc} \Delta_k, \Delta_{max}\}, & \text{if } \rho_k > \eta_1 \\ \max\{\bar{\gamma} \Delta_k, \bar{d}\}, & \text{if } \rho_k \leq \eta_1 \text{ and } \mu_k \geq \bar{\mu} \\ \Delta_k, & \text{otherwise.} \end{cases}$$

If  $\rho \leq \eta_1$  and  $\mu_k \geq \bar{\mu}$ , then set  $\mu_k = 0$ . Set  $n_{k+1} = n_k$ ,  $\mu_{k+1} = \mu_k$ ,  $k = k + 1$ .

If  $n_k < n_{max}$ , then go to step 2. Otherwise, stop.

The general-purpose method  $A\phi$  differs from the proposed method  $Am$  as follows: (i) the metamodel is formulated as Eq. (16) instead of Eq. (14); (ii) step 1 is not carried out.

## Appendix B

Table 10 displays a list of the notation used in the paper. The variables and parameters in the table are listed in the same order as they appear in the paper.

**Table 10**  
List of all variables and parameters.

Notation	Description
$T$	simulation time interval size
$\tau$	simulation time interval index
$i$	network link index
$n$	traveler index
$N$	total number of travelers in the system
$\Delta_{nir}$	stochastic binary indicator of traveler $n$ crossing link $i$ in time interval $\tau$
$F_{i,\tau}$	stochastic simulated link flow rate on link $i$ in time interval $\tau$
$\delta_{nir}$	probability that traveler $n$ crosses link $i$ in time interval $\tau$
$\theta$	general calibration parameter vector
$x$	network conditions
$y_{i,\tau}$	number of vehicles counted in the field on link $i$ in time interval $\tau$
$\mathcal{I}$	set of links with measurements available
$\Theta$	feasible region of the behavioral parameter
$k$	SO algorithm iteration index
$z$	endogenous simulation model variables
$\hat{q}$	exogenous simulation model variables
$f(\theta, z; \hat{q})$	simulation-based objective function
$\beta_{ik}$	vector of metamodel parameters for link $i$ at iteration $k$
$m_{i,k}(\theta; \beta_{i,k})$	metamodel for link $i$ at iteration $k$
$f_A(\theta)$	analytical problem-specific approximation of $f(\theta)$
$\tilde{v}$	endogenous analytical model variables
$\tilde{q}$	exogenous analytical model variables
$h(\theta, \tilde{v}; \tilde{q})$	analytical macroscopic traffic model
$\theta_L$	lower bound for $\theta$
$\theta_U$	upper bound for $\theta$
$\lambda_i(\theta)$	expected demand for link $i$
$s$	OD pair index
$d_s$	expected travel demand for OD pair $s$
$\ell_i$	space capacity of queue $i$
$l_i$	length of link $i$ (road length)
$v_i$	maximum speed of link $i$
$\mu_i$	service rate of queue $i$
$\gamma_i$	expected external demand of link $i$
$\bar{t}_i$	expected travel time of queue $i$
$\bar{n}_i$	expected number of vehicles in queue $i$

(continued on next page)

**Table 10** (continued)

Notation	Description
$p_{ij}$	turning probability from queue $i$ to queue $j$
$t_r$	expected travel time of route $r$
$f_r$	expected demand on route $r$
$\bar{p}_{sr}$	probability that a traveler of OD pair $s$ takes route $r$
$S$	set of OD pairs
$\mathcal{Q}$	set of queues
$\mathcal{R}$	set of routes
$\mathcal{R}_s$	set of routes of OD pair $s$
$\mathcal{G}_{ij}$	set of routes that consecutively go through queues $i$ and $j$
$\mathcal{H}_i$	set of routes that go through queue $i$
$\mathcal{T}_i$	set of routes that start with queue $i$
$\Psi_r$	set of links of route $r$
$Am$	proposed metamodel approach
$A\phi$	general-purpose metamodel approach
$\phi_{i,k}(\theta; \beta_{i,k})$	general-purpose metamodel for link $i$ at iteration $k$
$\theta^*$	true value for $\theta$
$\theta_0$	initial value for $\theta$

**Table 11**

Numerical values of the solutions (in units  $\text{h}^{-1}$ ) derived by each method for each experiment of the toy network.

		$Am$		$A\phi$			
$\theta^* = -5$	$\theta_0 = 0$	-4.7	-5.3	-4.0	-5.2	-5.3	-5.5
	$\theta_0 = -40$	-5.7	-5.5	-5.0	-5.2	-5.5	-5.3
	$\theta_0 = -60$	-4.8	-5.5	-5.8	-5.1	-5.0	-5.4
$\theta^* = -20$	$\theta_0 = 0$	-18.1	-18.4	-18.7	-18.3	-19.1	-18.5
	$\theta_0 = -40$	-19.6	-18.1	-18.6	-18.6	-17.7	-18.3
	$\theta_0 = -60$	-18.1	-18.5	-19.5	-18.3	-18.5	-18.0
$\theta^* = -55$	$\theta_0 = 0$	-56.1	-55.9	-55.9	-57.2	-56.4	-56.0
	$\theta_0 = -40$	-55.8	-56.3	-55.2	-56.0	-56.2	-56.4
	$\theta_0 = -60$	-55.8	-55.5	-56.6	-55.8	-56.3	-56.0

**Table 12**

Numerical values of the solutions (in units  $\text{h}^{-1}$ ) derived by each method for each experiment of the Berlin network.

		$Am$		$A\phi$		
$\theta_0 = 0$		-2.5	-1.2	-2.1	0.0	-1.1
$\theta_0 = -40$		-4.7	-4.7	-1.6	0.0	-2.8
$\theta_0 = -60$		-1.4	-4.7	-3.5	0.0	-0.4

## Appendix C

### Numerical values of the solutions derived by each method

Table 11 (resp. 12) displays the numerical values of the solutions derived by each method and each experiment of the toy (resp. Berlin) network.

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