

INFORMATION AND MINIMAL LENGTH CODING

6.02 / Nov. 30, 2010

BASICS:

GIVEN: A PROBABILITY DISTRIBUTION
ON A FINITE ALPHABET

(EX.: $\{ "A", "B", "C" \}$, $P_A = \frac{1}{4}$, $P_B = \frac{1}{3}$, $P_C = \frac{5}{12}$)

WANT: A VARIABLE LENGTH ^{BIT} CODE, i.e.

A CORRESPONDENCE ALPHABET $\rightarrow$ BIT SEQUENCES
SUCH THAT THE SEQUENCE CORRESPONDING TO
ONE ~~AND~~ SYMBOL IS NEVER A BEGINNING
OF A SEQUENCE CORRESPONDING TO ANOTHER SYMBOL

Ex.:

$A \rightarrow 0$
 $B \rightarrow 10$
 $C \rightarrow 11$ } is a
VALID
ENCODING
CODE

$A \rightarrow 0$
 $B \rightarrow 1$
 $C \rightarrow 01$ } is NOT A
VALID CODE

PERFORMANCE ("AVERAGE LENGTH"):

$$A.L. = \sum_{\text{SYMBOLS}} P(\text{SYMBOL}) \cdot \text{LENGTH}(\text{SYMBOL ENCODING})$$

(EX.: $1 \cdot \frac{1}{4} + 2 \cdot \frac{1}{3} + 2 \cdot \frac{5}{12} = \frac{7}{4}$ FOR THE CODE ABOVE)

USEFUL RESULTS:

$$* A.L. \geq \underbrace{\sum_{\text{SYMBOLS}} P(\text{SYMBOL}) \cdot \log_2 \left(\frac{1}{P(\text{SYMBOL})} \right)}_{\text{"INFORMATION ENTROPY" / "SHANNON'S ENTROPY"}} \quad \text{"INFORMATION" DELIVERED BY THE SYMBOL}$$

$$* \underbrace{\text{MINIMAL}}_{A.L.} \leq 1 + \sum_s P(s) \log_2 (1/P(s))$$

* MINIMAL A.L. IS DELIVERED BY HUFFMAN CODES

* FOR THE ALPHABET OF n -LENGTH WORDS
 $W = (s_1, s_2, \dots, s_n)$, WITH $P((s_1, s_2, \dots, s_n)) = P(s_1)P(s_2) \dots P(s_n)$

THE ENTROPY IS $n \cdot (\text{ORIGINAL ENTROPY})$

i.e. THE ENTROPY IS OPTIMAL AVERAGE NUMBER OF BITS ($n \rightarrow \infty$)

6.02/Nov 30, 2010Ex. 1

YOU ARE TOLD THAT TWO INDEPENDENT
TOSSES OF A FAIR DIE (SIDES 1, 2, 3, 4, 5, 6)
HAVE YIELDED A SUM OF 4.
HOW MUCH INFO IS THIS?

3 possibilities LEFT OUT OF 36
(1+3, 2+2, 3+1)

$$\text{INFO} = \log_2 (36/3) \approx 3.58$$

$$(\text{IF THE SUM} = 7, \text{ INFO} = \log_2 (36/6) \approx \frac{11.72}{3.58-1} = 2.58)$$

Ex. 2

UNFAIR COIN: ($P(\text{TAIL}) = 0.02$)

NEED: TO ENCODE THE OUTCOME OF 100 TOSSES,
TO MINIMIZE A.L. MINIMAL A.L. = ?

THE ALPHABET HAS 2^{100} SYMBOLS!

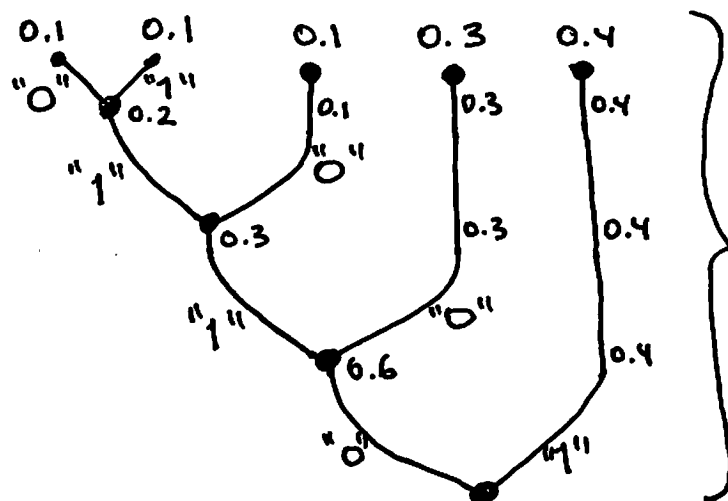
$$\text{ENTROPY} = 100 \cdot [0.02 \cdot \log \frac{1}{0.02} + 0.98 \cdot \log \frac{1}{0.98}] \approx 14.144$$

HENCE:

$$14.1 < \text{MINIMAL A.L.} < 15.2$$

Ex. 3

HUFFMAN CODE FOR $(P_1, P_2, P_3, P_4, P_5) = (0.1, 0.1, 0.1, 0.3, 0.4)$

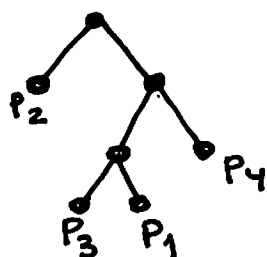


ASSIGN 0/1
ARBITRARILY
AT EACH BRANCHING

Ex. 4

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③



THIS IS A
HUFFMAN CODE TREE.
WHAT CAN BE SAID ABOUT
 P_1, P_2, P_3, P_4 ?

$$\left. \begin{array}{l} P_1 \leq P_2 \\ P_1 \leq P_4 \\ P_3 \leq P_2 \\ P_3 \leq P_4 \end{array} \right\} \text{BECAUSE } P_1, P_3 \text{ ARE THE TWO LONGEST}$$

$$\left. \begin{array}{l} P_1 + P_3 \leq P_2 \\ P_4 \leq P_2 \end{array} \right\} \text{BECAUSE } (P_1, P_3) \text{ IS JOINED WITH } P_4$$

IS THIS CONSISTENT WITH $P_2 = 0.3$?

NO, BECAUSE THEN $1 = P_1 + P_2 + P_3 + P_4 \leq 3P_2 = 0.9$!

IS THIS CONSISTENT WITH $P_2 = \frac{1}{3}$?

YES, HERE'S AN EXAMPLE: $P_2 = \frac{1}{3}, P_4 = \frac{1}{3}, P_1 = \frac{1}{6}, P_3 = \frac{1}{6}$

Ex. 5

$$P(S_1) = \frac{1}{2}$$

$$P(S_2) = P(S_3) = \dots = P(S_{100}) = \frac{1}{198}$$

WHAT IS THE LENGTH OF ENCODING S_1
IN THE HUFFMAN CODE?

ANSWER: 1

Ex. 6

$$P_1 = 0.24, P_2 = 0.25, P_3 = 0.28, P_4 = 0.23$$

HOW MANY OPTIMAL LENGTH ENCODINGS?

ANSWER: $4! = 24$