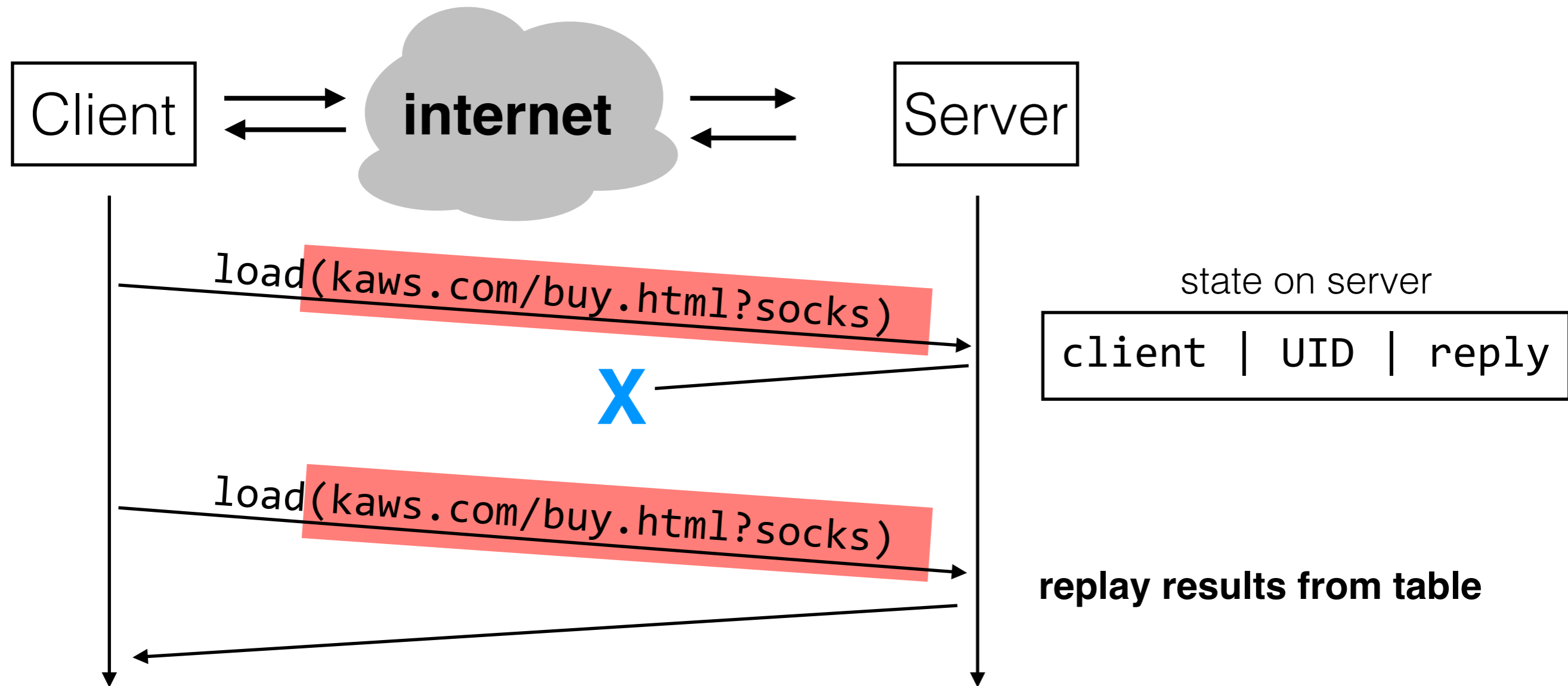


6.033 Spring 2018

Lecture #2

- Naming in systems
- Case study: DNS

Last Time: Enforced Modularity via Client/Server Model



Today: Naming

allows modules to interact

Examples of Names

mit.edu

lacurts@mit.edu

lacurts

R0

main

WebBrowser

/mit/6.033/www/schedule.shtml

http://web.mit.edu/about

617-253-7341

128.30.2.121

hostname

email

username

x86 register name

function name

class name

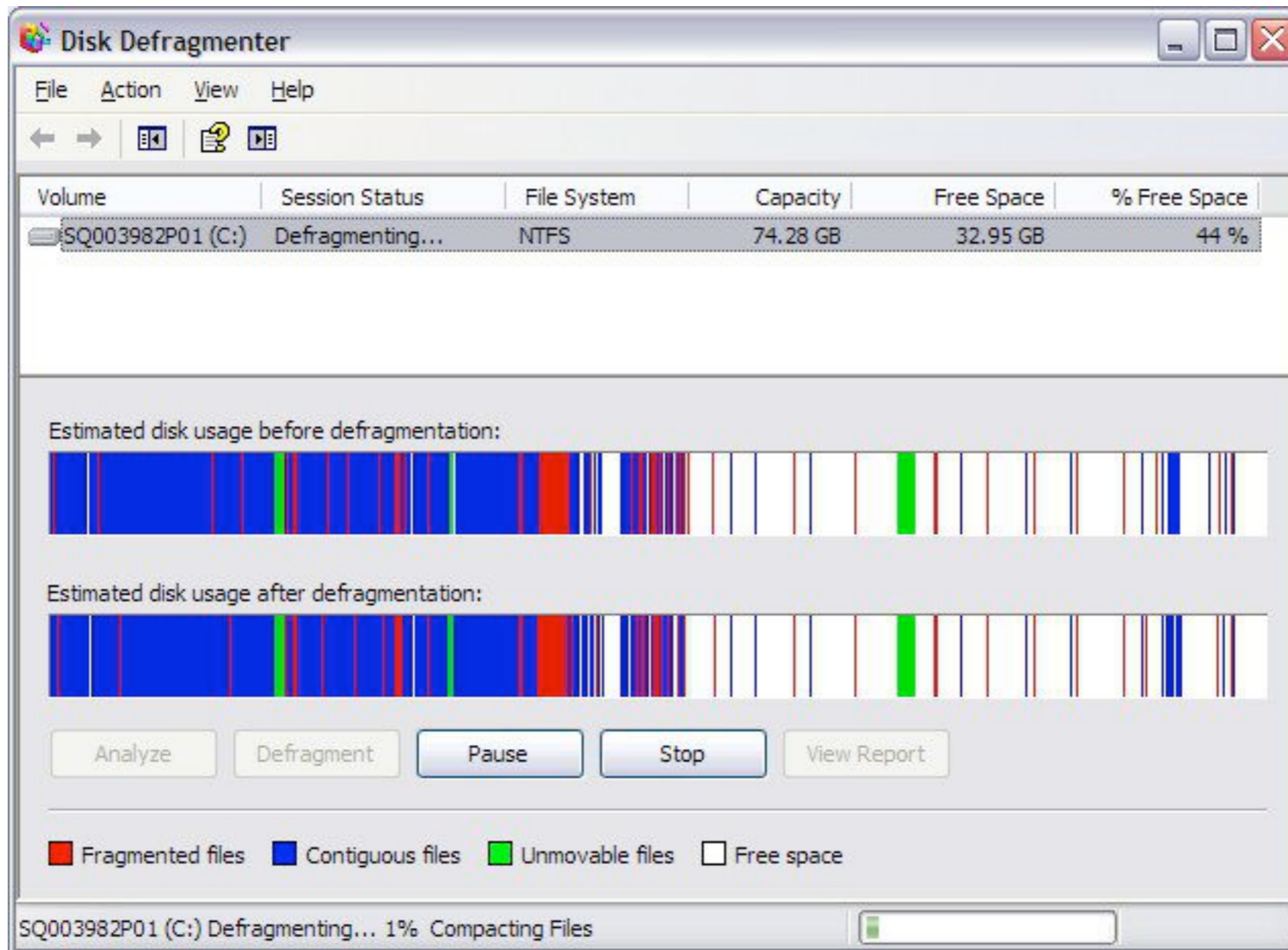
path name

URL

phone number

IP Address

why use names?



why use names?

Naming Schemes

1. Set of all possible **names**
2. Set of all possible **values**
3. **Look-up algorithm** to translate a name into a value (or set of values, or “none”)

Domain Name System

1. **names:** hostnames (`web.mit.edu`)

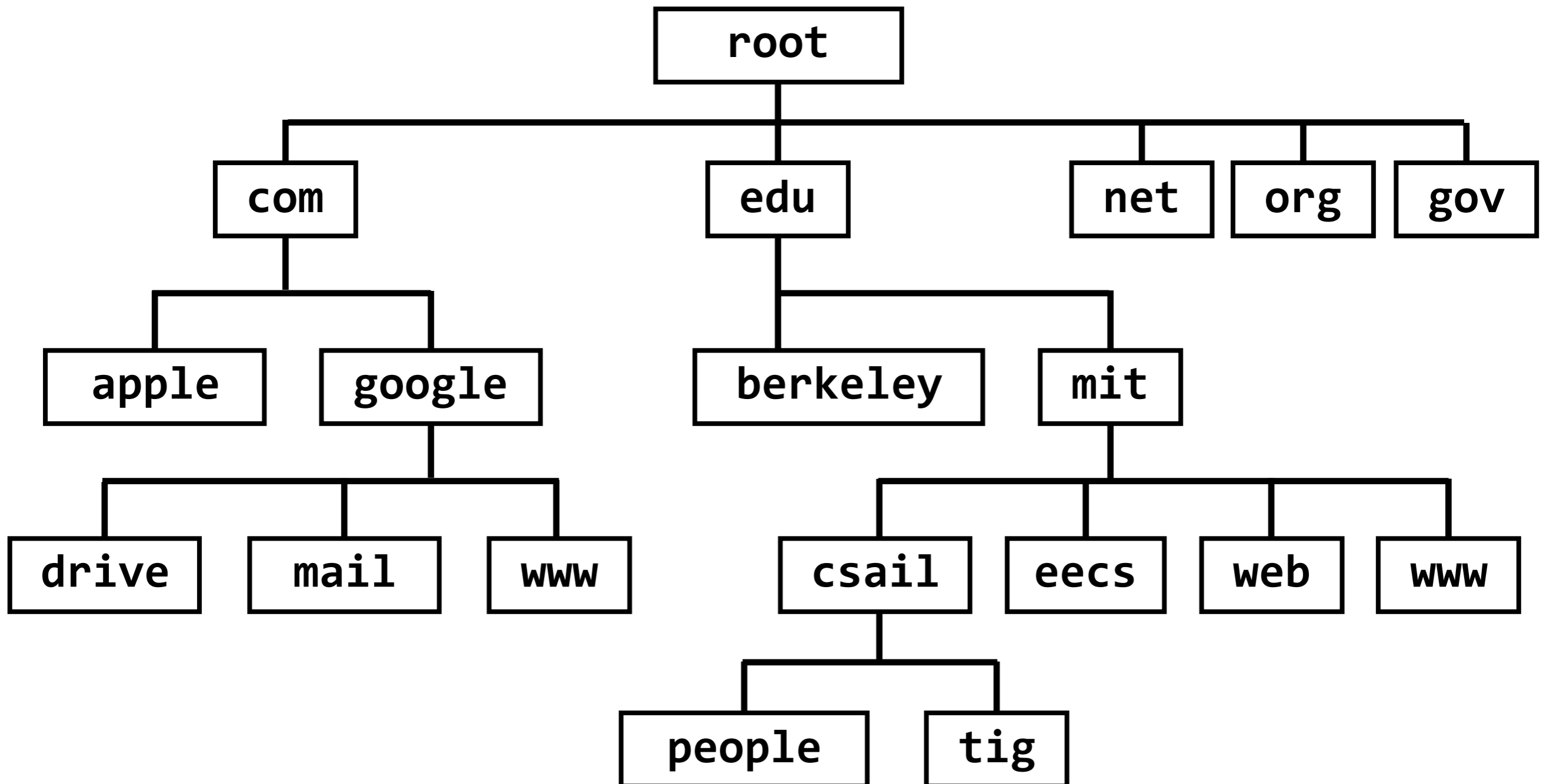
2. **values:** IP addresses (`18.9.22.69`)

IP addresses are imbued with location information: routers can send packets to an IP address, but not to a hostname

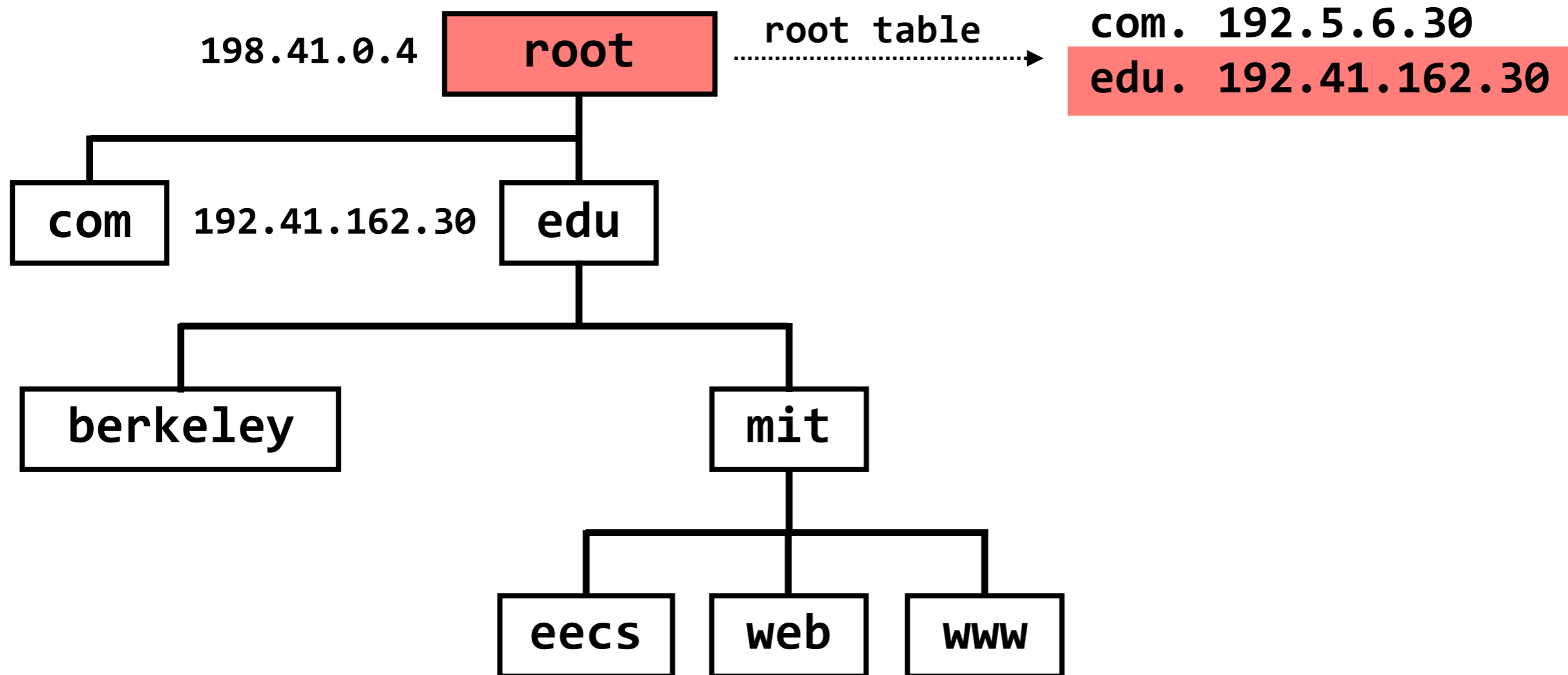
3. **look-up algorithm:** resolves a hostname to an IP address so that your machine knows where to send data

DNS Hierarchy

(a partial view)



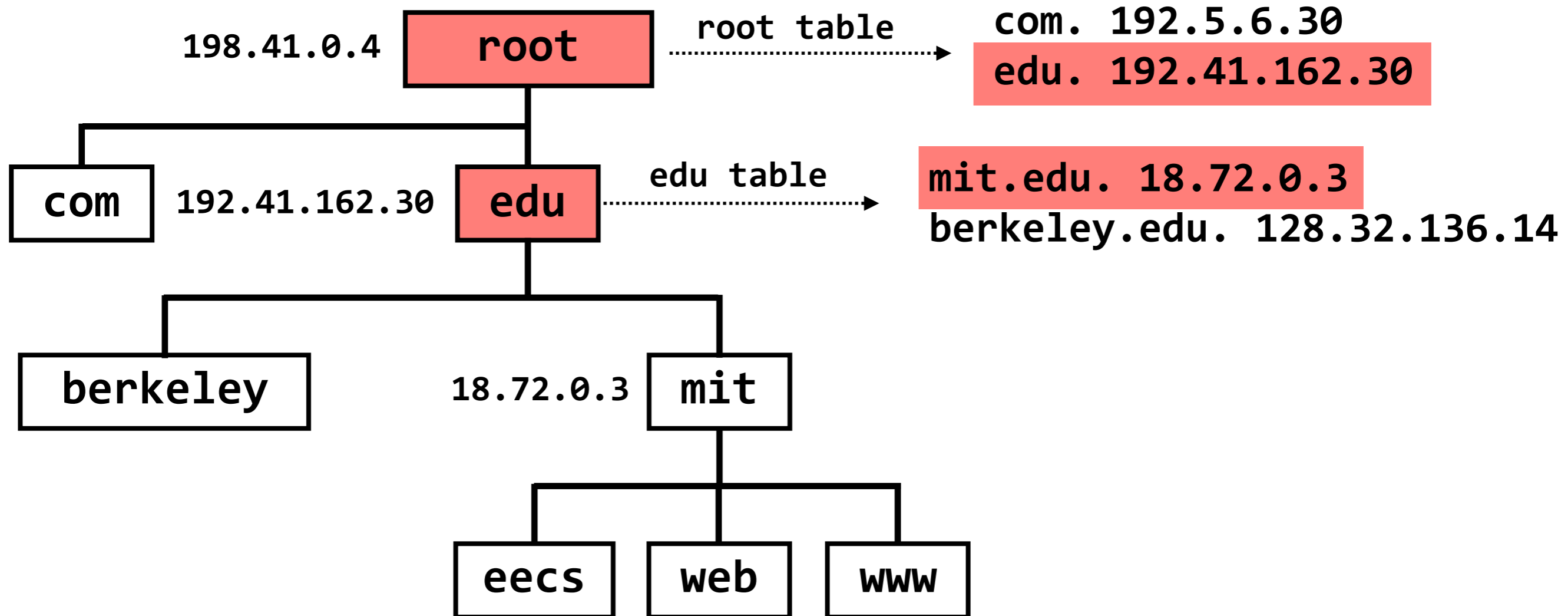
DNS Look-up for web.mit.edu



query to: 198.41.0.4

result: edu. 192.41.162.30

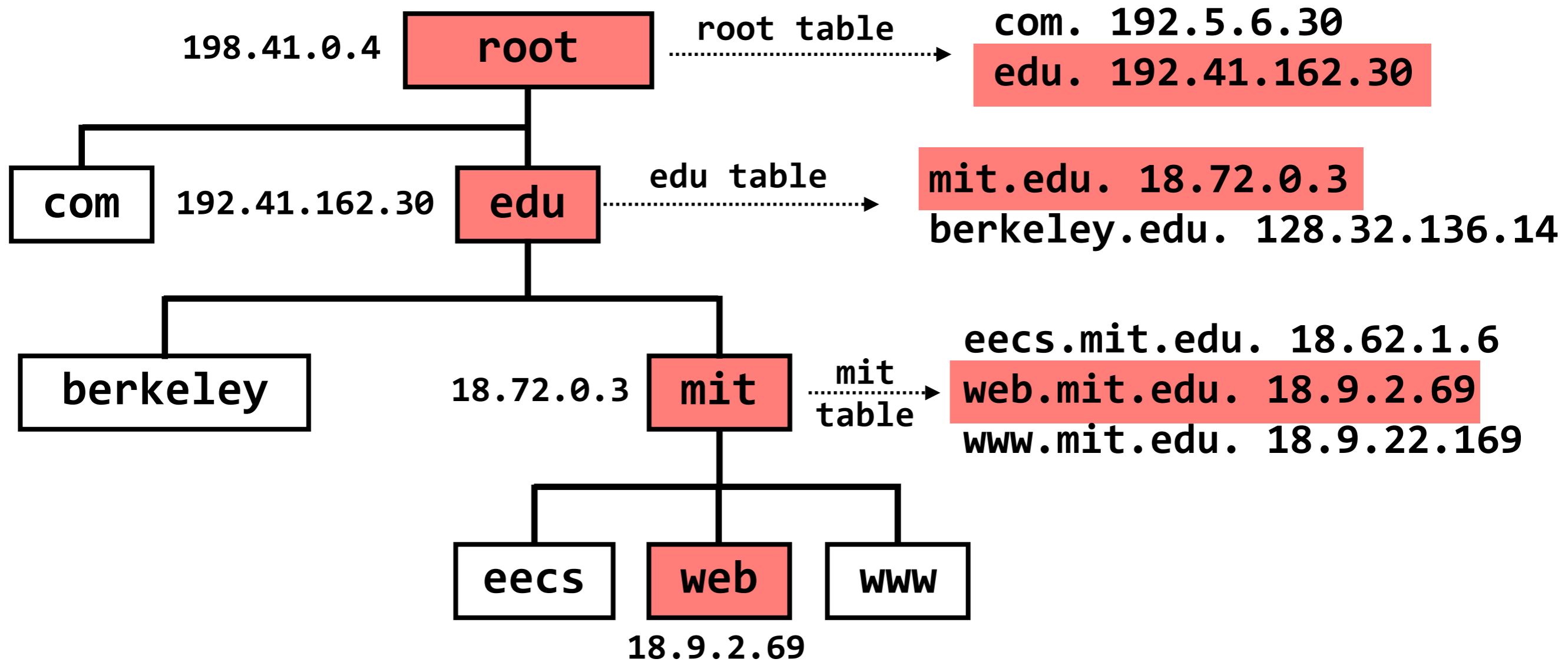
DNS Look-up for web.mit.edu



query to: 192.41.162.30

result: mit.edu. 18.72.0.3

DNS Look-up for web.mit.edu

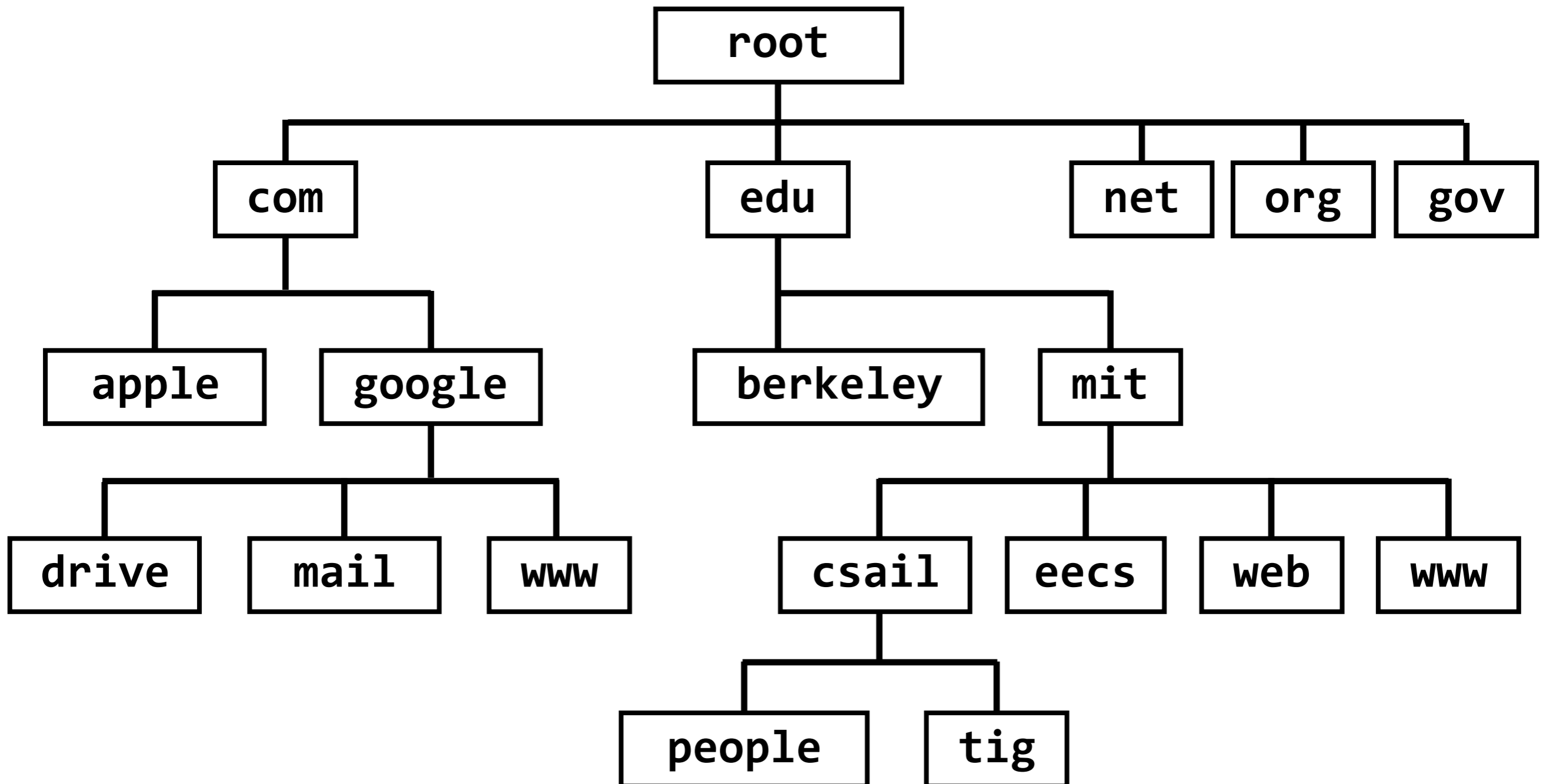


query to: 18.72.0.3

result: web.mit.edu. 18.9.2.69

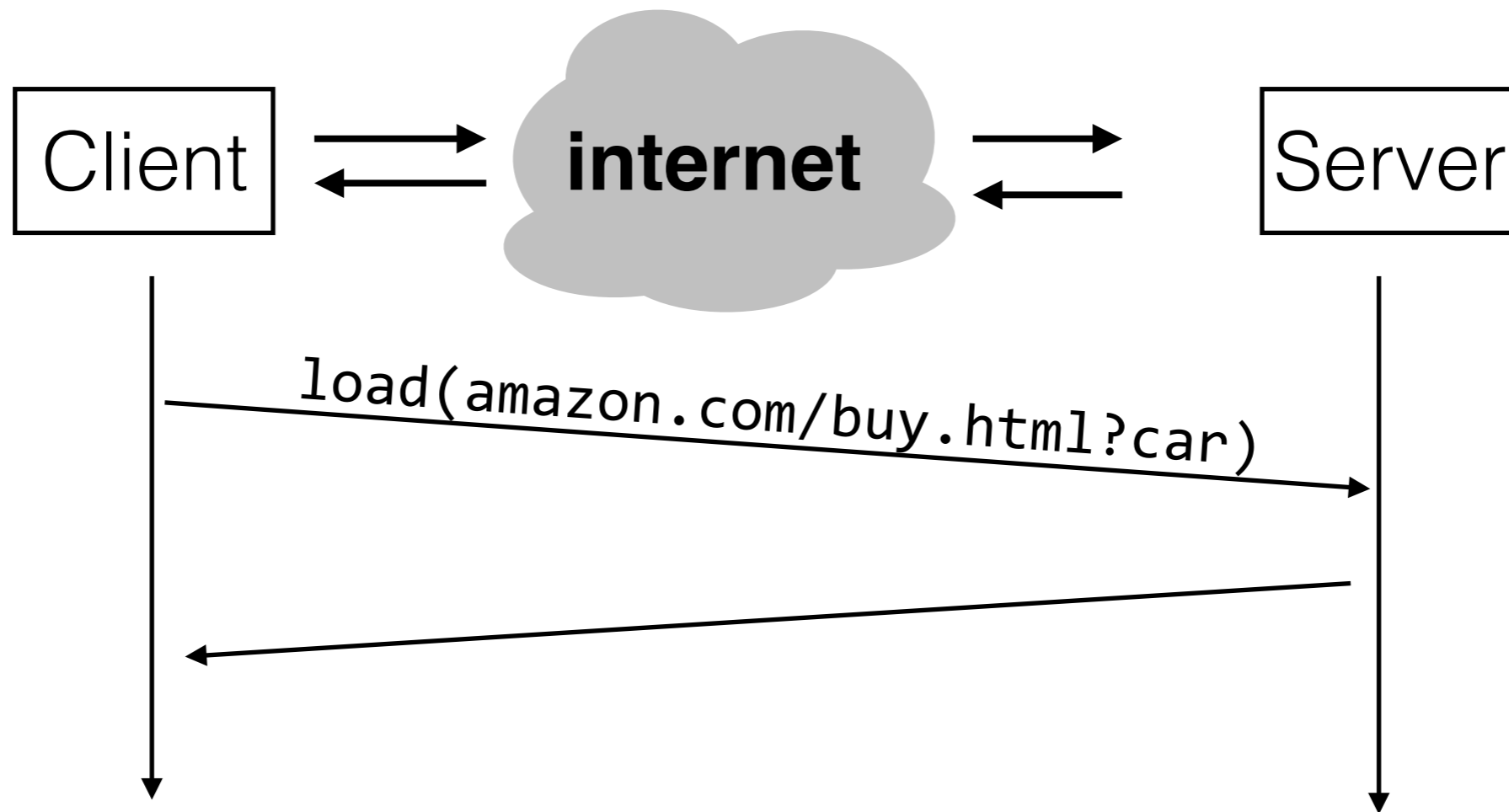
DNS Hierarchy

(a partial view)



- **Modularity** (and abstraction) limit complexity. One way to enforce modularity is to use a client/server design
- **Naming** is what allows modules — for example, a client and a server — to communicate; it is pervasive across systems
- **DNS** maps hostnames to IP addresses. It is also a good example of **hierarchy**.

Lingering Problem



what if we don't want our modules to be on entirely separate machines? how can we **enforce modularity on a single machine?**