Fall 2023

# Special structures

What makes some sequential decision-making problems easy?

### Cathy Wu

6.7920 Reinforcement Learning: Foundations and Methods



- 1. 6.231 Sp22 Lecture 3 notes, Section 2 [N3 §2]
- 2. DPOC vol 1, 3.1 (LQR), 3.3-3.4
- (Optional) MIT Underactuated Robotics, Chapter 10 Trajectory Optimization [link]

## Outline

- 1. Recap & roadmap
- 2. Template for structural DP arguments
- 3. Example: optimal stopping
- 4. Linear quadratic control (LQR)

## Outline

### 1. Recap & roadmap

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### So far: sequential decision making is hard

"Roadmap"

This time: What makes *some* sequential decision problems easy?

Next time [3x]: Why is there still *hope* of solving sequential decision problems? (general solutions for *small-state* problems)

Next next time [8x]: Why is there still hope of solving *large-state* problems?

## Outline

1. Recap & roadmap

### 2. Template for structural DP arguments

- a. Convexity, monotonicity
- 3. Example: optimal stopping
- 4. Linear quadratic control (LQR)

### Template for Structural DP Arguments

- 1. Recognize that the **terminal** reward/cost-to-go function  $V_T^*$  has a **nice property** (base case in induction proof).
  - Example: convexity or monotonicity
- 2. Then, argue that this property implies that the policy  $\pi^*_{T-1}$  has some nice structure.
  - Example: a threshold policy is optimal
- 3. Extend this with an **induction step**: we show that if a reward-to-go function V satisfies the property, then the "next" reward-to-go function:

$$V^{-}(x) = \max_{a \in A(s)} \mathbb{E}\left[g(s, a, w) + V(f(s, a, w))\right]$$

that is generated by a step of the DP algorithm will also satisfy this property.

### **Operations that Preserve Convexity**

- Comes in handy in showing the convexity of reward-to-go functions.
- Non-negative weighted sums:
  - If  $f_1, ..., f_m: \mathcal{D} \to \mathbb{R}$  are convex and  $w_1, ..., w_m \ge 0$ , then  $w_1 f_1 + ... + w_m f_m$  is convex.
  - For some  $f: \mathcal{X} \times \mathcal{Y} \to \mathbb{R}$ , the expectation  $g: \mathcal{X} \to \mathbb{R}$  defined as  $g(x) = \int f(x, y) w(y) dy$

is convex if  $w(y) \ge 0$  and the mapping  $x \mapsto f(x, y)$  is convex for all  $y \in \mathcal{Y}$ .

- Composition with an affine map:
  - g(x) = f(Ax + b) is convex if f is convex.
- Point-wise supremum:
  - $g(x) = \sup_{y \in \mathcal{Y}} f(x, y)$  is convex if  $x \mapsto f(x, y)$  is convex for all  $y \in \mathcal{Y}$ .

*Further reading: For a detailed treatment, please refer to the book Convex Optimization by Boyd and Vandenberghe.* 

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- **3.** Example: optimal stopping
- 4. Linear quadratic control (LQR)

### Asset Selling With Irrevocable Decisions

- Discrete time setting,  $t = \{0, 1, \dots, T-1\}$
- Problem: you have an asset to sell by time T.
  - At each epoch
    - You receive an offer  $w_t$  drawn independently from some distribution W (bounded).
    - You must either accept the offer and invest the money at a fixed interest rate r > 0 or reject and wait for the next offer.
  - Goal: maximize the expected final revenue.
- Notes:
  - Continuous state problem!
  - Assume that a rejected offer is lost.

### Asset Selling With Irrevocable Decisions

State s<sub>t</sub>

$$s_{t+1} = \begin{cases} \text{sold} & \text{if } A_t = \text{Accept or } s_t = \text{sold} \\ w_t & o.w. \end{cases}$$

 $\forall \{t = 0, \dots T - 1\}.$ 

- Set  $s_0 = 0$  as a dummy variable.
- The state space is  $S \subset \mathbb{R} \cup \{\text{sold}\}$ .

Action space:

$$A_t(s_t) = \begin{cases} \emptyset & \text{if } s_t = \text{sold} \\ \{\text{Accept, Reject}\} & o.w. \end{cases}$$

The revenue for each period is defined as:

$$g_t(s_t, u_t, w_t) = \begin{cases} 0 & \text{if } u_t \neq \text{Accept} \\ (1+r)^{T-t} s_t & \text{if } u_t = \text{Accept} \end{cases}$$

with the revenue for the final state being:

$$g_T(s_T) = \begin{cases} 0 & \text{if } s_T = \text{sold} \\ s_T & o.w. \end{cases}$$

### DP recursion & optimal policy

• Following the DP algorithm described in the previous section, set  $V_T^*(s) = g_T(s)$ . For  $t = \{T - 1, T - 2, ..., 0\}$ , set:  $V_t^*(s) = \begin{cases} \max\{(1 + r)^{T-t}s, \mathbb{E}[V_{t+1}^*(w_t)]\} & \text{if } s \neq \text{sold} \\ 0 & \text{if } s = \text{sold} \end{cases}$ 

Given the structure of the value-to-go functions, V<sup>\*</sup><sub>t</sub>(s),
 the optimal policy can be easily computed as the following threshold policy:

$$\pi_t^*(s_t) | (s_t \neq \text{sold}) = \begin{cases} \text{Accept} & \text{if } s_t \ge \alpha_t \\ \text{Reject} & \text{if } s_t \le \alpha_t \end{cases}$$

where the thresholds,  $\alpha_t = \frac{\mathbb{E}[V_{t+1}^*(w_t)]}{(1+r)^{T-t}}$ , depend on time *t*.

• We remark that  $\alpha_t$ 's obey their own recursion. Since we must accept the last offer,  $\alpha_T = -\infty$ . For  $t = \{T - 1, ..., 0\}$ 

$$\alpha_t = \frac{1}{1+r} \mathbb{E}[\max\{w_t, \alpha_{t+1}\}]$$
Proof: By induction.

• **Remark (infinite horizon version)**: With i.i.d. offers, the optimal policy is stationary and the optimal threshold  $\alpha^*$  is the solution to the fixed point equation:

$$\alpha = \frac{1}{1+r} \mathbb{E}[\max\{w, \alpha\}]$$

In plain English: Accept the offer if it's better to invest now than wait for a slightly better offer in the future, which loses out on factor(s) of the interest rate,

### Asset Selling With Offers Retained

- Now consider the setting:
  - The offers  $w_0, \ldots, w_{T-1}$  are i.i.d., non-negative, bounded.
  - The rejected offers are not lost. At any period *t*, we can choose the highest offer received so far.
- To accommodate this setting, we define the state such that

 $s_{t+1} = \begin{cases} \text{sold} & \text{if } A_t = \text{Accept or } s_t = \text{sold} \\ \max\{s_t, w_t\} & o.w. \end{cases}$  $\forall t = \{0, \dots, T-1\}.$ 

• The action space and functions  $g_t$ 's stay the same.

### Optimal policy

#### Proposition

An optimal policy for asset selling with offers retained is a stationary policy  $\pi^* = (\mu^*, \mu^*, ..., \mu^*)$ , where for  $s \neq \text{sold}$ ,  $\pi_t^*(s) = \begin{cases} \text{Accept if } s \geq \frac{1}{1+r} \mathbb{E}_w[\max\{s, w\}] \\ \text{Reject } o.w. \end{cases}$ 

### Proof (Proposition)

- 1. Monotonicity: For  $s \neq \text{sold}$ , we can set  $V_T^*(s) = s$ . For t = T 1 and  $s \neq \text{sold}$ ,  $V_{T-1}^*(s) = \max\{(1+r)s, \mathbb{E}[\max\{w_{T-1}, s\}]\}$   $\geq (1+r)s$  $= (1+r)V_T^*(s)$
- 2. By induction, assume that  $V_{t+1}^*(s) \ge (1+r)V_{t+2}^*(s)$ . Then  $V_t^*(s) = \max\{(1+r)^{T-t}s, \mathbb{E}[V_{t+1}^*(\max\{s, w_t\})]\}$   $\ge \max\{(1+r)^{T-t}s, (1+r)\mathbb{E}[V_{t+2}^*(\max\{s, w_t\})]\}$   $= (1+r)\max\{(1+r)^{T-(t+1)}s, \mathbb{E}[V_{t+2}^*(\max\{s, w_t\})]\}$  $= (1+r)V_{t+1}^*(s)$
- 3. Optimal stopping set:  $S_t^* := \{s \mid s \ge \alpha_t \coloneqq (1+r)^{-(T-t)} \mathbb{E}[V_{t+1}^*(\max\{s, w_t\})]\}$
- **4.** Convergence: thresholds  $\alpha_t$  converge (backwards) because:
  - Thresholds  $\alpha_t$  are monotonically increasing (backwards)

$$a_t \ge \alpha_{t+1} \to S_t^* \subseteq S_{t+1}^*$$

- Thresholds  $\alpha_t$  are bounded above (bounded offers)
- Thresholds  $\alpha_t \rightarrow \frac{1}{1+r} \mathbb{E}_w[\max\{s, w\}]$ , since 1)  $S_t^* \supseteq S_{t+1}^*$ , 2)  $a_{T-1} = \frac{1}{1+r} \mathbb{E}_w[\max\{s, w\}]$

Wu

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- 1. Recap & roadmap
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- 3. Example: optimal stopping
- 4. Linear quadratic control (LQR)
  - a. Finite horizon LQR
  - b. Linear quadratic Gaussian & Certainty equivalence
  - c. Infinite horizon LQR & Algebraic Riccati Equations

### Notation "break"

In the following section, and in deference to the rich tradition in control theory, we will be using standard control theory notation

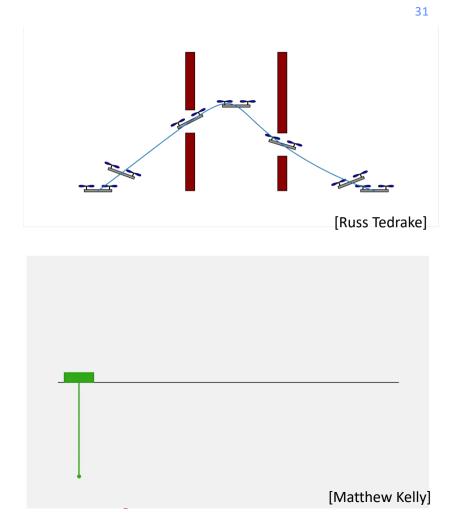
- x and u, in place of s and a, to denote state and the control
- c(x, u) in place of r(s, a), to denote immediate cost or reward
- r(s,a) = -c(x,u)

### Trajectory optimization

Synthesis of Complex Behaviors with Online Trajectory Optimization

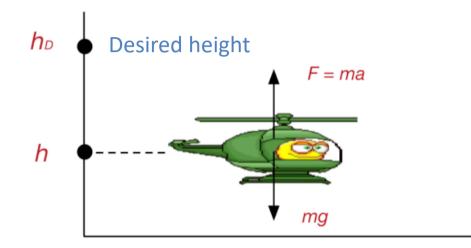
Yuval Tassa, Tom Erez & Emo Todorov

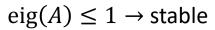
IEEE International Conference on Intelligent Robots and Systems 2012



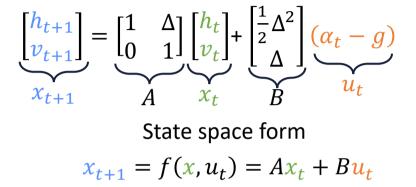
### Linear quadratic control

Assumptions: deterministic, finite horizon, discrete time





Further reading: Chen, Chi-Tsong. Linear system theory and design. 1984. 33



Linear time-invariant (LTI) system

w.l.o.g.

The dynamics (discrete form) are governed by the equations of motion is:

$$h_{t+1} = h_t + \Delta v_t + \frac{1}{2}\Delta^2(\alpha_t - g)$$
$$v_{t+1} = v_t + \Delta(\alpha_t - g)$$

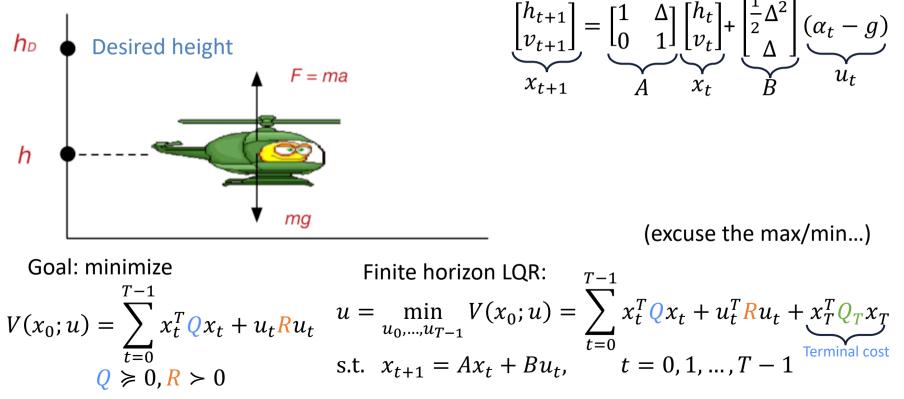
where  $\Delta$  = time step (sec)

Adapted from Kevin Jamieson

 $\begin{aligned} \mathbf{x}_t &\coloneqq \begin{bmatrix} h_t \\ v_t \end{bmatrix} - \mathbf{x}_D \\ \mathbf{x}_D &\coloneqq \begin{bmatrix} h_D \\ 0 \end{bmatrix} \end{aligned}$ 

### Linear quadratic control

Assumptions: deterministic, finite horizon, discrete time

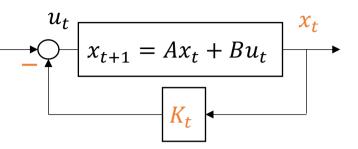


Adapted from Kevin Jamieson

### Linear quadratic control

Finite horizon LQR

$$u = \min_{u_0, \dots, u_{T-1}} V(x_0; u) = \sum_{t=0}^{T-1} x_t^T Q x_t + u_t^T R u_t + x_T^T Q_T x_T$$
  
s.t.  $x_{t+1} = A x_t + B u_t$ ,  $t = 0, 1, \dots, T-1$ 



 $eig(A - BK_t) \le 1 \rightarrow stable$ 

Optimal control law is a linear feedback controller:  $x_{t+1} = Ax_t + Bu_t = (A - BK_t)x_t$ 

#### Theorem (Finite horizon LQR)

The optimal cost-to-go and optimal control at time t are given by:

$$V^*(x_t) = x_t^T P_t x_t$$
$$u_t^* = -K_t x_t$$

where

$$P_{t} = Q + K_{t}^{T} R K_{t} + (A - B K_{t})^{T} P_{t+1} (A - B K_{t}), \qquad P_{T} = Q_{T}$$
  

$$K_{t} = (R + B^{T} P_{t+1} B)^{-1} B^{T} P_{t+1} A, \qquad t \in \{0, ..., T - 1\}$$

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### Proof (induction)

Base case (stage T):

$$V^*(x_t) = x_t^T P_t x_t$$
  
$$\Rightarrow P_T = Q_T$$

Finite horizon LQR:

Theorem (Finite horizon LQR)

T = 1

s.t.  $x_{t+1} = Ax_t + Bu_t$ , t = 0, 1, ..., T - 1

 $u = \min_{u_0, \dots, u_{T-1}} V(x_0; u) = \sum_{t=0}^{T} x_t^T Q x_t + u_t^T R u_t + x_T^T Q_T x_T$ 

The optimal cost-to-go and optimal control at time t are given by:

 $V^*(x_t) = x_t^T P_t x_t$ • Special structure:  $V^*(x_T) = x_T^T Q_T x_T$  is convex.  $u_t^* = -K_t x_t$ • Induction: assume  $P_t$  holds &  $V^*(x_t)$  convex, show for t-1 where  $\overline{P_t} = Q + K_t^T R K_t + (A - B K_t)^T P_{t+1} (A - B K_k), \qquad P_T = Q_T$ Recall:  $r_t(x_t, u_t) \coloneqq x_t^T O x_t + u_t R u_t$  $K_t = (R + B^T P_{t+1} B)^{-1} B^T P_{t+1} A_t$  $t \in \{0, \dots, T-1\}$  $V^*(x_{t-1}) = \min_{u_{t-1}} [x_{t-1}^T Q x_{t-1} + u_{t-1}^T R u_{t-1} + V^*(x_t)]$ (principle of optimality)  $= \min[x_{t-1}^T Q x_{t-1} + u_{t-1}^T R u_{t-1} + x_t^T P_t x_t]$ (induction hypothesis)  $= \min_{u_{t-1}} [x_{t-1}^T Q x_{t-1} + u_{t-1}^T R u_{t-1} + (A x_{t-1} + B u_{t-1})^T P_t (A x_{t-1} + B u_{t-1})] \quad \text{(system equations)}$  $\nabla_{u_{t-1}} V^*(x_{t-1}) = 2u_{t-1}^T R + 2(Ax_{t-1} + Bu_{t-1})^T P_t B = 0$ (convexity)  $u_{t-1}^* = (R + B^T P_t B)^{-1} B^T P_t A x_{t-1} = -K_{t-1} x_{t-1}$  $(R > 0, \text{ derives } K_t \text{ for any } t)$  $V^{*}(x_{t-1}) = x_{t-1}^{T} Q x_{t-1} + u_{t-1}^{*T} R u_{t-1}^{*} + (A x_{t-1} + B u_{t-1}^{*})^{T} P_{t}(A x_{t-1} + B u_{t-1}^{*})$  $= x_{t-1}^{T} \left( Q + K_{t-1}^{T} R K_{t-1} + (A - B K_{t-1})^{T} P_{t} (A - B K_{t-1}) \right) x_{t-1}$ (derives  $P_{t-1}$ )  $= x_{t-1}^T P_{t-1} x_{t-1}$ 

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### Linear quadratic control (stochastic)

- Assumptions: deterministic, finite horizon, discrete time
- Gaussian noise → Linear quadratic Gaussian (LQG) problem

$$x_{t+1} = f(x_t, u_t, \epsilon_t) = Ax_t + Bu_t + \epsilon_t \quad \epsilon_t \sim \mathcal{N}(0, \Sigma)$$

• Revised optimization problem:

$$u = \min_{u_0, \dots, u_{T-1}} V(x_0; u) = \mathbb{E} \left[ \sum_{t=0}^{T-1} x_t^T Q x_t + u_t R u_t + x_T^T Q_f x_T \right]$$
  
subject to  $x_{t+1} = A x_t + B u_t + \epsilon_t$ 

#### Theorem (LQG)

The optimal cost-to-go and optimal control at time t are given by:

$$V^*(x_t) = x_t^T P_t x_t + \Sigma_t$$
$$u_t^* = -K_t x_t$$

**certainty equivalence**: control as if disturbances were known (deterministic)!

where

$$\begin{split} P_t &= Q + K_t^T R K_t + (A - B K_t)^T P_{t+1} (A - B K_t), \quad P_T = Q_f \\ K_t &= (R + B^T P_{t+1} B)^{-1} B^T P_{t+1} A, \quad \Sigma_{t-1} = Tr \left( \Sigma P_t \right) + \Sigma_t, \quad \Sigma_T = 0 \\ t \in \{0, \dots, T-1\} \end{split}$$

- Intuition: noise terms are independent of actions → optimal actions don't change.
- Exercise: complete the proof!

### Linear quadratic control (towards infinite horizon)

- Assumptions: deterministic, finite horizon, discrete time
- Revised optimization problem:

$$u^* = \min_{u_0, \dots, u_{T-1}} V(x_0; u) = \lim_{T \to \infty} \sum_{t=0}^{T-1} x_t^T Q x_t + u_t R u_t$$
  
Subject to  $x_{t+1} = A x_t + B u_t$  Later: infinite horizon problems

TT 1

- Before (finite horizon): finite horizon  $\rightarrow$  finite sum.
- Now, need some condition to keep sum finite.
  - System (A, B) is **controllable** if A is full rank &  $\overline{A} := [B \ AB \ A^2B \ ... \ A^{n-1}B]$  is full rank (n).

#### Theorem (infinite horizon LQR)

If the system (A, B) is controllable, the optimal cost-to-go and optimal control converges to

$$V^*(x) = x^T P x$$
$$u^* = -K x$$

Algebraic Riccati Equation (ARE)

No "final ston"

where

$$P = Q + A^{T}PA - A^{T}PB(R + B^{T}PB)^{-1}B^{T}PA$$
$$K = (R + B^{T}PB)^{-1}B^{T}PA$$

Exercise: show that the expression is equivalent to before (in the limit).

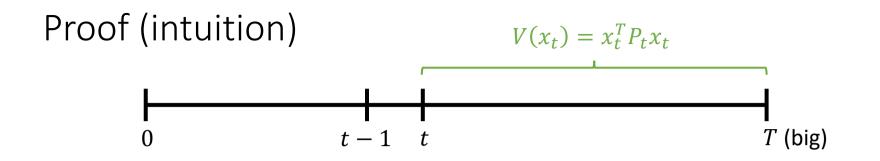
Controllability (for linear systems)

- System is **controllable** if A is full rank &  $\overline{A} = \begin{bmatrix} B & AB & A^2B & \dots & A^{n-1}B \end{bmatrix}$  is full rank (n).
- Intuition: Can s' be reached within n steps from any s?

$$\begin{aligned} x_{t+1} &= Ax_t + Bu_t \\ &= A(Ax_{t-1} + Bu_{t-1}) + Bu_t \\ &= A^2 x_{t-1} + ABu_{t-1} + Bu_t \\ &= A^3 x_{t-2} + A^2 Bu_{t-2} + ABu_{t-1} + Bu_t \end{aligned}$$

...

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For simplicity, take  $P_T = Q_T = 0$ 

$$x^T P_t x \le x^T P_{t-1} x$$
 (PSD)

• As  $T \to \infty$ ,  $x^T P_0 x$  must converge or go to infinity

- Controllability  $\rightarrow$  For every x, there is a sequence  $u_0, \dots, u_{n-1}$  (where  $x \in \mathbb{R}^n$ ) that drives x to 0.
- After n steps, can set  $u_k = 0$  for  $k \ge n$ .
- Controllability
  - $\rightarrow x^T P_0 x$  is bounded above, for any x
  - $\rightarrow P_0$  converges to finite limit.

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### LQR example (implementation)

$$\begin{aligned} x(k+1) &= \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} x(k) + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u(k) \\ y(k) &= \begin{bmatrix} 1 & 0 \end{bmatrix} x(k) \end{aligned}$$

Let  $Q = C^T C = I, R = 0.3$ . Infinite horizon. Solve the optimal control. Solve DARE

$$S = A^{T}SA - A^{T}SB(R + B^{T}SB)^{-1}B^{T}SA + Q, K = -(R + B^{T}SB)^{-1}B^{T}SA$$
$$\Rightarrow S = \begin{bmatrix} 2.751.91\\ 1.913.34 \end{bmatrix}$$
$$K = [-0.524, -1.44]$$

Courtesy Ding Zhao (CMU)

### LQR example (implementation)

```
1 from future import division, print function
 2 import numpy as np
 3 import scipy.linalq
 4 def dlgr(A,B,Q,R):
       """Solve the discrete time lqr controller.
 5
       x[k+1] = A x[k] + S B u[k]
 6
 7
       cost = sum x[k].T*Q*x[k] + u[k].T*R*u[k]
       ....
 8
 9
       #ref Bertsekas, p.151
      #first, try to solve the ricatti equation
10
11
       S = np.matrix(scipy.linalq.solve discrete are(A, B, Q, R))
12
       #compute the LOR gain
       K = -np.matrix(scipy.linalg.inv(B.T*S*B+R)*(B.T*S*A))
13
14
       eigVals, eigVecs = scipy.linalg.eig(A+B*K)
15
       return K, S, eigVals
```

```
1 A = np.array([[1,1],[0,1]])
2 B = np.array([[0],[1]])
3 Q = np.eye(2)
4 R = 0.3
5
6 K,S,_ = dlqr(A,B,Q,R)
7 print("S:", S)
8 print("K:", K)
```

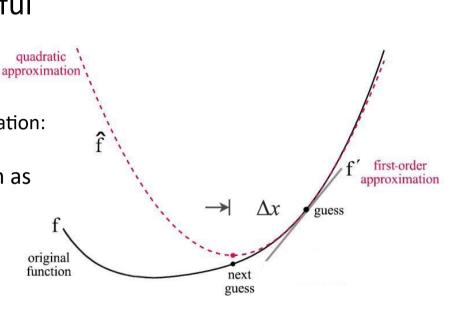
```
S: [[2.75078485 1.90801622]
[1.90801622 3.34052588]]
K: [[-0.52410456 -1.44169888]]
```

### Further reading & extensions

- Further reading: 6.231 Sp22 Lecture 3 notes, Section 2 [N3 §2]
  - Excellent exposition, generally
  - Further discussion on observability
  - Stability
  - Loose ends (connections to other topics
- Lots of extensions
  - Continuous time (Callier & Desoer)
  - Model estimation, via LS & recursive LS
  - Adaptive control (Abbasi-Yadkori, 2011)
  - Unknown models, robust LQR (Dean, 2017)
  - Time Varying Regression with Hidden Linear Dynamics (Mania, 2022)

### LQR – final notes

- Iterative LQR remains a powerful approach, e.g. in robotics.
- Extensions
  - Iterative LQR (iLQR) (full implementation: <u>https://github.com/anassinator/ilqr</u>)
    - 1. Approximate a nonlinear system as LQR using Taylor expansion
    - 2. Take a step or three
    - Rinse and repeat to update the model & objective



[Jonathan Hui, 2018]

### And there you have it - iLQR

Synthesis of Complex Behaviors with Online Trajectory Optimization

Yuval Tassa, Tom Erez & Emo Todorov

IEEE International Conference on Intelligent Robots and Systems 2012

### Summary & takeaways

- Certain DP problems admit closed form solutions, such as optimal stopping and linear quadratic control (LQR).
- DP for problems with special structures can be analyzed by induction, by showing that the special structure holds from one step to the previous, as well as for the terminal case. Special structures include convexity and monotonicity.
- LQR exhibits certainty equivalence: the optimal policy remains the same when random disturbances are replaced with their means (conditional expectation).

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- 1. 6.231 Sp22 Lecture 3 notes, Section 2 [N3 §2]
- 2. DPOC vol 1, 3.1 (LQR), 3.3-3.4
- 3. Some material adapted from:
  - Daniel Russo (Columbia)