

Dimitri P. Bertsekas



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Convex Optimization Theory

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Dimitri Bertsekas studied Mechanical and Electrical Engineering at the National Technical University of Athens, Greece, and obtained his Ph.D. in system science from the Massachusetts Institute of Technology. He has held faculty positions with the Engineering-Economic Systems Department, Stanford University, and the Electrical Engineering Department of the University of Illinois, Urbana. Since 1979 he has been teaching at the Electrical Engineering and Computer Science Department of the Massachusetts Institute of Technology (M.I.T.), where he is McAfee Professor of Engineering. In 2019, he joined the School of Computing, Informatics, and Decision Systems Engineering at the Arizona State University, Tempe, AZ, as Fulton Professor of Computational Decision Making.

Professor Bertsekas' teaching and research have spanned several fields, including deterministic optimization, dynamic programming and stochastic control, large-scale and distributed computation, artificial intelligence, and data communication networks. He has authored or coauthored numerous research papers and eighteen books, several of which are currently used as textbooks in MIT classes, including "Dynamic Programming and Optimal Control," "Data Networks," "Introduction to Probability," and "Nonlinear Programming."

Professor Bertsekas was awarded the INFORMS 1997 Prize for Research Excellence in the Interface Between Operations Research and Computer Science for his book "Neuro-Dynamic Programming" (co-authored with John Tsitsiklis), the 2001 AACC John R. Ragazzini Education Award, the 2009 INFORMS Expository Writing Award, the 2014 AACC Richard Bellman Heritage Award, the 2014 INFORMS Khachiyan Prize for Life-Time Accomplishments in Optimization, the 2015 MOS/SIAM George B. Dantzig Prize, and the 2022 IEEE Control Systems Award. In 2018 he shared with his coauthor, John Tsitsiklis, the 2018 INFORMS John von Neumann Theory Prize for the contributions of the research monographs "Parallel and Distributed Computation" and "Neuro-Dynamic Programming." Professor Bertsekas was elected in 2001 to the United States National Academy of Engineering for "pioneering contributions to fundamental research, practice and education of optimization/control theory, and especially its application to data communication networks."

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Preface

This book aims at an accessible, concise, and intuitive exposition of two related subjects that find broad practical application:

- (a) Convex analysis, particularly as it relates to optimization.
- (b) Duality theory for optimization and minimax problems, mainly within a convexity framework.

The focus on optimization is to derive conditions for existence of primal and dual optimal solutions for constrained problems such as

minimize
$$f(x)$$

subject to $x \in X$, $g_j(x) \le 0$, $j = 1, ..., r$.

Other types of optimization problems, such as those arising in Fenchel duality, are also part of our scope. The focus on minimax is to derive conditions guaranteeing the equality

$$\inf_{x \in X} \sup_{z \in Z} \phi(x, z) = \sup_{z \in Z} \inf_{x \in X} \phi(x, z),$$

and the attainment of the "inf" and the "sup."

The treatment of convexity theory is fairly detailed. It touches upon nearly all major aspects of the subject, and it is sufficient for the development of the core analytical issues of convex optimization. The mathematical prerequisites are a first course in linear algebra and a first course in real analysis. A summary of the relevant material is provided in an appendix. Prior knowledge of linear and nonlinear optimization theory is not assumed, although it will undoubtedly be helpful in providing context and perspective. Other than this modest background, the development is self-contained, with rigorous proofs provided throughout.

We have aimed at a unified development of the strongest possible forms of duality with the most economic use of convexity theory. To this end, our analysis often departs from the lines of Rockafellar's 1970 book and other books that followed the Fenchel/Rockafellar formalism. For example, we treat differently closed set intersection theory and preservation of closure under linear transformations (Sections 1.4.2 and 1.4.3); we develop subdifferential calculus by using constrained optimization duality (Section 5.4.2); and we do not rely on concepts such as infimal convolution, image, polar sets and functions, bifunctions, and conjugate saddle functions. Perhaps our greatest departure is in duality theory itself: similar to Fenchel/Rockafellar, our development rests on Legendre/Fenchel conjugacy ideas, but is far more geometrical and visually intuitive.

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Our duality framework is based on two simple geometrical problems: the min common point problem and the max crossing point problem. The salient feature of the min common/max crossing (MC/MC) framework is its highly visual geometry, through which all the core issues of duality theory become apparent and can be analyzed in a unified way. Our approach is to obtain a handful of broadly applicable theorems within the MC/MC framework, and then specialize them to particular types of problems (constrained optimization, Fenchel duality, minimax problems, etc). We address all duality questions (existence of duality gap, existence of dual optimal solutions, structure of the dual optimal solution set), and other issues (subdifferential theory, theorems of the alternative, duality gap estimates) in this way.

Fundamentally, the MC/MC framework is closely connected to the conjugacy framework, and owes its power and generality to this connection. However, the two frameworks offer complementary starting points for analysis and provide alternative views of the geometric foundation of duality: conjugacy emphasizes functional/algebraic descriptions, while MC/MC emphasizes set/epigraph descriptions. The MC/MC framework is simpler, and seems better suited for visualizing and investigating questions of strong duality and existence of dual optimal solutions. The conjugacy framework, with its emphasis on functional descriptions, is more suitable when mathematical operations on convex functions are involved, and the calculus of conjugate functions can be brought to bear for analysis or computation.

The book evolved from the earlier book of the author [BNO03] on the subject (coauthored with A. Nedić and A. Ozdaglar), but has different character and objectives. The 2003 book was quite extensive, was structured (at least in part) as a research monograph, and aimed to bridge the gap between convex and nonconvex optimization using concepts of nonsmooth analysis. By contrast, the present book is organized differently, has the character of a textbook, and concentrates exclusively on convex optimization. Despite the differences, the two books have similar style and level of mathematical sophistication, and share some material.

The chapter-by-chapter description of the book follows:

Chapter 1: This chapter develops all of the convex analysis tools that are needed for the development of duality theory in subsequent chapters. It covers basic algebraic concepts such as convex hulls and hyperplanes, and topological concepts such as relative interior, closure, preservation of closedness under linear transformations, and hyperplane separation. In addition, it develops subjects of special interest in duality and optimization, such as recession cones and conjugate functions.

Chapter 2: This chapter covers polyhedral convexity concepts: extreme points, the Farkas and Minkowski-Weyl theorems, and some of their applications in linear programming. It is not needed for the developments of subsequent chapters, and may be skipped at first reading.

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Chapter 3: This chapter focuses on basic optimization concepts: types of minima, existence of solutions, and a few topics of special interest for duality theory, such as partial minimization and minimax theory.

Chapter 4: This chapter introduces the MC/MC duality framework. It discusses its connection with conjugacy theory, and it charts its applications to constrained optimization and minimax problems. It then develops broadly applicable theorems relating to strong duality and existence of dual optimal solutions.

Chapter 5: This chapter specializes the duality theorems of Chapter 4 to important contexts relating to linear programming, convex programming, and minimax theory. It also uses these theorems as an aid for the development of additional convex analysis tools, such as a powerful nonlinear version of Farkas' Lemma, subdifferential theory, and theorems of the alternative. A final section is devoted to nonconvex problems and estimates of the duality gap, with special focus on separable problems.

In aiming for brevity, I have omitted a number of topics that an instructor may wish for. One such omission is applications to specially structured problems; the book by Boyd and Vanderbergue [BoV04], as well as my book on parallel and distributed computation with John Tsitsiklis [BeT89] cover this material extensively (both books are available on line).

Another important omission is computational methods. However, I have written a long supplementary sixth chapter (over 100 pages), which covers the most popular convex optimization algorithms (and some new ones), and can be downloaded from the book's web page

http://www.athenasc.com/convexduality.html.

This chapter, together with a more comprehensive treatment of convex analysis, optimization, duality, and algorithms will be part of a more extensive textbook that I am currently writing. Until that time, the chapter will serve instructors who wish to cover convex optimization algorithms in addition to duality (as I do in my M.I.T. course). This is a "living" chapter that will be periodically updated. Its current contents are as follows:

Chapter 6 on Algorithms: 6.1. Problem Structures and Computational Approaches; 6.2. Algorithmic Descent; 6.3. Subgradient Methods; 6.4. Polyhedral Approximation Methods; 6.5. Proximal and Bundle Methods; 6.6. Dual Proximal Point Algorithms; 6.7. Interior Point Methods; 6.8. Approximate Subgradient Methods; 6.9. Optimal Algorithms and Complexity.

While I did not provide exercises in the text, I have supplied a substantial number of exercises (with detailed solutions) at the book's web page. The reader/instructor may also use the end-of-chapter problems (a total of 175) given in [BNO03], which have similar style and notation to the present book. Statements and detailed solutions of these problems can be downloaded from the book's web page and are also available on line at

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http://www.athenasc.com/convexity.html.

The book may be used as a text for a theoretical convex optimization course; I have taught several variants of such a course at MIT and elsewhere over the last ten years. It may also be used as a supplementary source for nonlinear programming classes, and as a theoretical foundation for classes focused on convex optimization models (rather than theory).

The book has been structured so that the reader/instructor can use the material selectively. For example, the polyhedral convexity material of Chapter 2 can be omitted in its entirety, as it is not used in Chapters 3-5. Similarly, the material on minimax theory (Sections 3.4, 4.2.5, and 5.5) may be omitted; and if this is done, Sections 3.3 and 5.3.4, which use the tools of partial minimization, may be omitted. Also, Sections 5.4-5.7 are "terminal" and may each be omitted without effect on other sections.

A "minimal" self-contained selection, which I have used in my nonlinear programming class at MIT (together with the supplementary web-based Chapter 6 on algorithms), consists of the following:

- Chapter 1, except for Sections 1.3.3 and 1.4.1.
- Section 3.1.
- Chapter 4, except for Section 4.2.5.
- Chapter 5, except for Sections 5.2, 5.3.4, and 5.5-5.7.

This selection focuses on nonlinear convex optimization, and excludes all the material relating to polyhedral convexity and minimax theory.

I would like to express my thanks to several colleagues for their contributions to the book. My collaboration with Angelia Nedić and Asuman Ozdaglar on our 2003 book was important in laying the foundations of the present book. Huizhen (Janey) Yu read carefully early drafts of portions of the book, and offered several insightful suggestions. Paul Tseng contributed substantially through our joint research on set intersection theory, given in part in Section 1.4.2 (this research was motivated by earlier collaboration with Angelia Nedić). Feedback from students and colleagues, including Dimitris Bisias, Vivek Borkar, John Tsitsiklis, Mengdi Wang, and Yunjian Xu, is highly appreciated. Finally, I wish to thank the many outstanding students in my classes, who have been a continuing source of motivation and inspiration.

Basic Concepts of Convex Analysis

Convex sets and functions are very useful in optimization models, and have a rich mathematical structure that is convenient for analysis and algorithms. Much of this structure can be traced to a few fundamental properties. For example, each closed convex set can be described in terms of the hyperplanes that support the set, each point on the boundary of a convex set can be approached through the relative interior of the set, and each halfline belonging to a closed convex set still belongs to the set when translated to start at any point in the set.

Yet, despite their favorable structure, convex sets and their analysis are not free of anomalies and exceptional behavior, which cause serious difficulties in theory and applications. For example, contrary to affine and compact sets, some basic operations such as linear transformation and vector sum may not preserve the closedness of closed convex sets. This in turn complicates the treatment of some fundamental optimization issues, including the existence of optimal solutions and duality.

For this reason, it is important to be rigorous in the development of convexity theory and its applications. Our aim in this first chapter is to establish the foundations for this development, with a special emphasis on issues that are relevant to optimization.

1.1 CONVEX SETS AND FUNCTIONS

We introduce in this chapter some of the basic notions relating to convex sets and functions. This material permeates all subsequent developments in this book. Appendix A provides the definitions, notational conventions, and results from linear algebra and real analysis that we will need. We first define convex sets (cf. Fig. 1.1.1).

Definition 1.1.1: A subset C of \Re^n is called *convex* if

$$\alpha x + (1 - \alpha)y \in C$$
, $\forall x, y \in C, \forall \alpha \in [0, 1]$.

Note that the empty set is by convention considered to be convex. Generally, when referring to a convex set, it will usually be apparent from the context whether this set can be empty, but we will often be specific in order to minimize ambiguities. The following proposition gives some operations that preserve convexity.

Proposition 1.1.1:

(a) The intersection $\cap_{i \in I} C_i$ of any collection $\{C_i \mid i \in I\}$ of convex sets is convex.

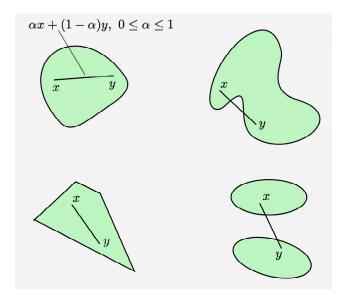


Figure 1.1.1. Illustration of the definition of a convex set. For convexity, linear interpolation between any two points of the set must yield points that lie within the set. Thus the sets on the left are convex, but the sets on the right are not.

- (b) The vector sum $C_1 + C_2$ of two convex sets C_1 and C_2 is convex.
- (c) The set λC is convex for any convex set C and scalar λ . Furthermore, if C is a convex set and λ_1 , λ_2 are positive scalars,

$$(\lambda_1 + \lambda_2)C = \lambda_1C + \lambda_2C.$$

- (d) The closure and the interior of a convex set are convex.
- (e) The image and the inverse image of a convex set under an affine function are convex.

Proof: The proof is straightforward using the definition of convexity. To prove part (a), we take two points x and y from $\bigcap_{i \in I} C_i$, and we use the convexity of C_i to argue that the line segment connecting x and y belongs to all the sets C_i , and hence, to their intersection.

Similarly, to prove part (b), we take two points of $C_1 + C_2$, which we represent as $x_1 + x_2$ and $y_1 + y_2$, with $x_1, y_1 \in C_1$ and $x_2, y_2 \in C_2$. For any $\alpha \in [0, 1]$, we have

$$\alpha(x_1 + x_2) + (1 - \alpha)(y_1 + y_2) = (\alpha x_1 + (1 - \alpha)y_1) + (\alpha x_2 + (1 - \alpha)y_2).$$

By convexity of C_1 and C_2 , the vectors in the two parentheses of the right-

hand side above belong to C_1 and C_2 , respectively, so that their sum belongs to $C_1 + C_2$. Hence $C_1 + C_2$ is convex. The proof of part (c) is left as an exercise for the reader. The proof of part (e) is similar to the proof of part (b).

To prove part (d), let C be a convex set. Choose two points x and y from the closure of C, and sequences $\{x_k\} \subset C$ and $\{y_k\} \subset C$, such that $x_k \to x$ and $y_k \to y$. For any $\alpha \in [0, 1]$, the sequence

$$\{\alpha x_k + (1-\alpha)y_k\},\$$

which belongs to C by the convexity of C, converges to $\alpha x + (1-\alpha)y$. Hence $\alpha x + (1-\alpha)y$ belongs to the closure of C, showing that the closure of C is convex. Similarly, we choose two points x and y from the interior of C, and we consider open balls that are centered at x and y, and have sufficiently small radius r so that they are contained in C. For any $\alpha \in [0,1]$, consider the open ball of radius r that is centered at $\alpha x + (1-\alpha)y$. Any point in this ball, say $\alpha x + (1-\alpha)y + z$, where $\|z\| < r$, belongs to C, because it can be expressed as the convex combination $\alpha(x+z) + (1-\alpha)(y+z)$ of the vectors x+z and y+z, which belong to C. Hence the interior of C contains $\alpha x + (1-\alpha)y$ and is therefore convex. Q.E.D.

Special Convex Sets

We will often consider some special sets, which we now introduce. A *hyperplane* is a set specified by a single linear equation, i.e., a set of the form $\{x \mid a'x = b\}$, where a is a nonzero vector and b is a scalar. A *halfspace* is a set specified by a single linear inequality, i.e., a set of the form $\{x \mid a'x \leq b\}$, where a is a nonzero vector and b is a scalar. It is clearly closed and convex. A set is said to be *polyhedral* if it is nonempty and it is the intersection of a finite number of halfspaces, i.e., if it has the form

$$\{x \mid a'_j x \le b_j, j = 1, \dots, r\},\$$

where a_1, \ldots, a_r and b_1, \ldots, b_r are some vectors in \Re^n and scalars, respectively. A polyhedral set is convex and closed, being the intersection of halfspaces [cf. Prop. 1.1.1(a)].

A set C is said to be a *cone* if for all $x \in C$ and $\lambda > 0$, we have $\lambda x \in C$. A cone need not be convex and need not contain the origin, although the origin always lies in the closure of a nonempty cone (see Fig. 1.1.2). A polyhedral cone is a set of the form

$$C = \{x \mid a_i' x \le 0, j = 1, \dots, r\},\$$

where a_1, \ldots, a_r are some vectors in \Re^n . A subspace is a special case of a polyhedral cone, which is in turn a special case of a polyhedral set.

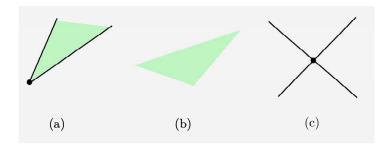


Figure 1.1.2. Illustration of convex and nonconvex cones. Cones (a) and (b) are convex, while cone (c), which consists of two lines passing through the origin, is not convex. Cone (a) is polyhedral. Cone (b) does not contain the origin.

1.1.1 Convex Functions

We now define a real-valued convex function (cf. Fig. 1.1.3).

Definition 1.1.2: Let C be a convex subset of \Re^n . We say that a function $f:C\mapsto\Re$ is *convex* if

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y), \qquad \forall \ x, y \in C, \ \forall \ \alpha \in [0, 1].$$

$$(1.1)$$

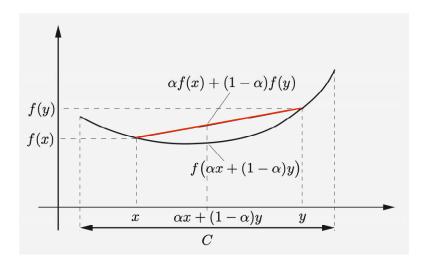


Figure 1.1.3. Illustration of the definition of a function $f: C \mapsto \Re$ that is convex. The linear interpolation $\alpha f(x) + (1-\alpha)f(y)$ overestimates the function value $f(\alpha x + (1-\alpha)y)$ for all $\alpha \in [0,1]$.

Note that, according to our definition, convexity of the domain C is a prerequisite for convexity of a function $f: C \to \Re$. Thus when calling a function convex, we imply that its domain is convex.

We introduce some variants of the basic definition of convexity. A convex function $f: C \mapsto \Re$ is called *strictly convex* if the inequality (1.1) is strict for all $x, y \in C$ with $x \neq y$, and all $\alpha \in (0,1)$. A function $f: C \mapsto \Re$, where C is a convex set, is called *concave* if the function (-f) is convex.

An example of a convex function is an affine function, one of the form f(x) = a'x + b, where $a \in \mathbb{R}^n$ and $b \in \mathbb{R}$; this is straightforward to verify using the definition of convexity. Another example is a norm $\|\cdot\|$ (any norm not just the Euclidean norm), since by the triangle inequality, we have

$$\|\alpha x + (1 - \alpha)y\| \le \|\alpha x\| + \|(1 - \alpha)y\| = \alpha \|x\| + (1 - \alpha)\|y\|,$$

for any $x, y \in \Re^n$, and $\alpha \in [0, 1]$ (see Definition A.2.1 in Appendix A).

If $f: C \mapsto \Re$ is a function and γ is a scalar, the sets $\{x \in C \mid f(x) \leq \gamma\}$ and $\{x \in C \mid f(x) < \gamma\}$, are called *level sets* of f. If f is a convex function, then all its level sets are convex. To see this, note that if $x, y \in C$ are such that $f(x) \leq \gamma$ and $f(y) \leq \gamma$, then for any $\alpha \in [0,1]$, we have $\alpha x + (1 - \alpha)y \in C$, by the convexity of C, so

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y) \le \gamma,$$

by the convexity of f. A similar proof also shows that the level sets $\{x \in C \mid f(x) < \gamma\}$ are convex when f is convex. Note, however, that convexity of the level sets does not imply convexity of the function; for example, the scalar function $f(x) = \sqrt{|x|}$ has convex level sets but is not convex.

Extended Real-Valued Convex Functions

We generally prefer to deal with convex functions that are real-valued and are defined over the entire space \Re^n (rather than over just a convex subset), because they are mathematically simpler. However, in some situations, prominently arising in the context of optimization and duality, we will encounter operations resulting in functions that can take infinite values. For example, the function

$$f(x) = \sup_{i \in I} f_i(x),$$

where I is an infinite index set, can take the value ∞ even if the functions f_i are real-valued, and the conjugate of a real-valued function often takes infinite values (cf. Section 1.6).

Furthermore, we will encounter functions f that are convex over a convex subset C and cannot be extended to functions that are real-valued and convex over the entire space \Re^n [e.g., the function $f:(0,\infty) \mapsto \Re$

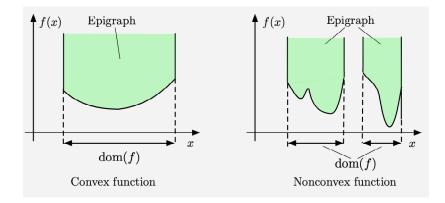


Figure 1.1.4. Illustration of the epigraphs and effective domains of extended real-valued convex and nonconvex functions.

defined by f(x) = 1/x]. In such situations, it may be convenient, instead of restricting the domain of f to the subset C where f takes real values, to extend the domain to all of \Re^n , but allow f to take infinite values.

We are thus motivated to introduce extended real-valued functions that can take the values $-\infty$ and ∞ at some points. Such functions can be characterized using the notion of epigraph, which we now introduce.

The *epigraph* of a function $f: X \mapsto [-\infty, \infty]$, where $X \subset \mathbb{R}^n$, is defined to be the subset of \mathbb{R}^{n+1} given by

$$epi(f) = \{(x, w) \mid x \in X, w \in \Re, f(x) \le w\}.$$

The effective domain of f is defined to be the set

$$dom(f) = \{x \in X \mid f(x) < \infty\}$$

(see Fig. 1.1.4). It can be seen that

$$dom(f) = \{x \mid \text{there exists } w \in \Re \text{ such that } (x, w) \in epi(f)\},\$$

i.e., $\operatorname{dom}(f)$ is obtained by a projection of $\operatorname{epi}(f)$ on \Re^n (the space of x). Note that if we restrict f to its effective domain, its epigraph remains unaffected. Similarly, if we enlarge the domain of f by defining $f(x) = \infty$ for $x \notin X$, the epigraph and the effective domain remain unaffected.

It is often important to exclude the degenerate case where f is identically equal to ∞ [which is true if and only if $\operatorname{epi}(f)$ is empty], and the case where the function takes the value $-\infty$ at some point [which is true if and only if $\operatorname{epi}(f)$ contains a vertical line]. We will thus say that f is proper if $f(x) < \infty$ for at least one $x \in X$ and $f(x) > -\infty$ for all $x \in X$, and we will say that f improper if it is not proper. In words, a function is proper if and only if its epigraph is nonempty and does not contain a vertical line.

A difficulty in defining extended real-valued convex functions f that can take both values $-\infty$ and ∞ is that the term $\alpha f(x) + (1-\alpha)f(y)$ arising in our earlier definition for the real-valued case may involve the forbidden sum $-\infty + \infty$ (this, of course, may happen only if f is improper, but improper functions arise on occasion in proofs or other analyses, so we do not wish to exclude them $a\ priori$). The epigraph provides an effective way of dealing with this difficulty.

Definition 1.1.3: Let C be a convex subset of \Re^n . We say that an extended real-valued function $f: C \mapsto [-\infty, \infty]$ is *convex* if $\operatorname{epi}(f)$ is a convex subset of \Re^{n+1} .

It can be easily verified that, according to the above definition, convexity of f implies that its effective domain $\mathrm{dom}(f)$ and its level sets $\{x \in C \mid f(x) \leq \gamma\}$ and $\{x \in C \mid f(x) < \gamma\}$ are convex sets for all scalars γ . Furthermore, if $f(x) < \infty$ for all x, or $f(x) > -\infty$ for all x, then

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y), \qquad \forall x, y \in C, \ \forall \ \alpha \in [0, 1], \ (1.2)$$

so the preceding definition is consistent with the earlier definition of convexity for real-valued functions.

By passing to epigraphs, we can use results about sets to infer corresponding results about functions (e.g., proving convexity). The reverse is also possible, through the notion of indicator function $\delta: \Re^n \mapsto (-\infty, \infty]$ of a set $X \subset \Re^n$, defined by

$$\delta(x\mid X) = \begin{cases} 0 & \text{if } x\in X,\\ \infty & \text{otherwise.} \end{cases}$$

In particular, a set is convex if and only if its indicator function is convex, and it is nonempty if and only if its indicator function is proper.

A convex function $f: C \mapsto (-\infty, \infty]$ is called *strictly convex* if the inequality (1.2) is strict for all $x, y \in \text{dom}(f)$ with $x \neq y$, and all $\alpha \in (0, 1)$. A function $f: C \mapsto [-\infty, \infty]$, where C is a convex set, is called *concave* if the function $(-f): C \mapsto [-\infty, \infty]$ is convex as per Definition 1.1.3.

Sometimes we will deal with functions that are defined over a (possibly nonconvex) domain C but are convex when restricted to a convex subset of their domain. The following definition formalizes this case.

Definition 1.1.4: Let C and X be subsets of \Re^n such that C is nonempty and convex, and $C \subset X$. We say that an extended real-valued function $f: X \mapsto [-\infty, \infty]$ is *convex over* C if f becomes convex when the domain of f is restricted to C, i.e., if the function $\tilde{f}: C \mapsto [-\infty, \infty]$, defined by $\tilde{f}(x) = f(x)$ for all $x \in C$, is convex.

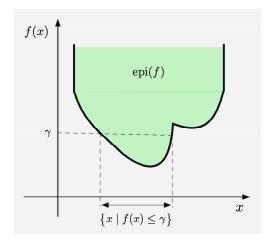


Figure 1.1.5. Visualization of the epigraph of a function in relation to its level sets. It can be seen that the level set $\{x \mid f(x) \leq \gamma\}$ can be identified with a translation of the intersection of $\operatorname{epi}(f)$ and the "slice" $\{(x,\gamma) \mid x \in \Re^n\}$, suggesting that $\operatorname{epi}(f)$ is closed if and only if all the level sets are closed.

By replacing the domain of an extended real-valued proper convex function with its effective domain, we can convert it to a real-valued function. In this way, we can use results stated in terms of real-valued functions, and we can also avoid calculations with ∞ . Thus, nearly all the theory of convex functions can be developed without resorting to extended real-valued functions. The reverse is also true, namely that extended real-valued functions can be adopted as the norm; for example, this approach is followed by Rockafellar [Roc70]. We will adopt a flexible approach, and use both real-valued and extended real-valued functions, depending on the context.

1.1.2 Closedness and Semicontinuity

If the epigraph of a function $f: X \mapsto [-\infty, \infty]$ is a closed set, we say that f is a closed function. Closedness is related to the classical notion of lower semicontinuity. Recall that f is called *lower semicontinuous* at a vector $x \in X$ if

$$f(x) \le \liminf_{k \to \infty} f(x_k)$$

for every sequence $\{x_k\} \subset X$ with $x_k \to x$. We say that f is lower semicontinuous if it is lower semicontinuous at each point x in its domain X. We say that f is upper semicontinuous if -f is lower semicontinuous. These definitions are consistent with the corresponding definitions for real-valued functions [cf. Definition A.2.4(c)].

The following proposition connects closedness, lower semicontinuity, and closedness of the level sets of a function; see Fig. 1.1.5.

Proposition 1.1.2: For a function $f: \Re^n \mapsto [-\infty, \infty]$, the following are equivalent:

- (i) The level set $V_{\gamma} = \{x \mid f(x) \leq \gamma\}$ is closed for every scalar γ .
- (ii) f is lower semicontinuous.
- (iii) epi(f) is closed.

Proof: If $f(x) = \infty$ for all x, the result trivially holds. We thus assume that $f(x) < \infty$ for at least one $x \in \Re^n$, so that $\operatorname{epi}(f)$ is nonempty and there exist level sets of f that are nonempty.

We first show that (i) implies (ii). Assume that the level set V_{γ} is closed for every scalar γ . Suppose, to arrive at a contradiction, that

$$f(\overline{x}) > \liminf_{k \to \infty} f(x_k)$$

for some \overline{x} and sequence $\{x_k\}$ converging to \overline{x} , and let γ be a scalar such that

$$f(\overline{x}) > \gamma > \liminf_{k \to \infty} f(x_k).$$

Then, there exists a subsequence $\{x_k\}_{\mathcal{K}}$ such that $f(x_k) \leq \gamma$ for all $k \in \mathcal{K}$, so that $\{x_k\}_{\mathcal{K}} \subset V_{\gamma}$. Since V_{γ} is closed, \overline{x} must also belong to V_{γ} , so $f(\overline{x}) \leq \gamma$, a contradiction.

We next show that (ii) implies (iii). Assume that f is lower semicontinuous over \Re^n , and let $(\overline{x}, \overline{w})$ be the limit of a sequence

$$\{(x_k, w_k)\} \subset \operatorname{epi}(f).$$

Then we have $f(x_k) \leq w_k$, and by taking the limit as $k \to \infty$ and by using the lower semicontinuity of f at \overline{x} , we obtain

$$f(\overline{x}) \le \liminf_{k \to \infty} f(x_k) \le \overline{w}.$$

Hence, $(\overline{x}, \overline{w}) \in \operatorname{epi}(f)$ and $\operatorname{epi}(f)$ is closed.

We finally show that (iii) implies (i). Assume that $\operatorname{epi}(f)$ is closed, and let $\{x_k\}$ be a sequence that converges to some \overline{x} and belongs to V_{γ} for some scalar γ . Then $(x_k, \gamma) \in \operatorname{epi}(f)$ for all k and $(x_k, \gamma) \to (\overline{x}, \gamma)$, so since $\operatorname{epi}(f)$ is closed, we have $(\overline{x}, \gamma) \in \operatorname{epi}(f)$. Hence, \overline{x} belongs to V_{γ} , implying that this set is closed. **Q.E.D.**

For most of our development, we prefer to use the closedness notion, rather than lower semicontinuity. One reason is that contrary to closedness, lower semicontinuity is a domain-specific property. For example, the function $f: \Re \mapsto (-\infty, \infty]$ given by

$$f(x) = \begin{cases} 0 & \text{if } x \in (0,1), \\ \infty & \text{if } x \notin (0,1), \end{cases}$$

is neither closed nor lower semicontinuous; but if its domain is restricted to (0,1) it becomes lower semicontinuous.

On the other hand, if a function $f: X \mapsto [-\infty, \infty]$ has a closed effective domain dom(f) and is lower semicontinuous at every $x \in \text{dom}(f)$, then f is closed. We state this as a proposition. The proof follows from the argument we used to show that (ii) implies (iii) in Prop. 1.1.2.

Proposition 1.1.3: Let $f: X \mapsto [-\infty, \infty]$ be a function. If dom(f) is closed and f is lower semicontinuous at each $x \in dom(f)$, then f is closed.

As an example of application of the preceding proposition, the indicator function of a set X is closed if and only if X is closed (the "if" part follows from the proposition, and the "only if" part follows using the definition of epigraph). More generally, if f_X is a function of the form

$$f_X(x) = \begin{cases} f(x) & \text{if } x \in X, \\ \infty & \text{otherwise,} \end{cases}$$

where $f: \mathbb{R}^n \to \mathbb{R}$ is a continuous function, then it can be shown that f_X is closed if and only if X is closed.

We finally note that an improper closed convex function is very peculiar: it cannot take a finite value at any point, so it has the form

$$f(x) = \begin{cases} -\infty & \text{if } x \in \text{dom}(f), \\ \infty & \text{if } x \notin \text{dom}(f). \end{cases}$$

To see this, consider an improper closed convex function $f: \mathbb{R}^n \mapsto [-\infty, \infty]$, and assume that there exists an x such that f(x) is finite. Let \overline{x} be such that $f(\overline{x}) = -\infty$ (such a point must exist since f is improper and f is not identically equal to ∞). Because f is convex, it can be seen that every point of the form

$$x_k = \frac{k-1}{k}x + \frac{1}{k}\overline{x}, \quad \forall \ k = 1, 2, \dots$$

satisfies $f(x_k) = -\infty$, while we have $x_k \to x$. Since f is closed, this implies that $f(x) = -\infty$, which is a contradiction. In conclusion, a closed convex function that is improper cannot take a finite value anywhere.

1.1.3 Operations with Convex Functions

We can verify the convexity of a given function in a number of ways. Several commonly encountered functions, such as affine functions and norms, are convex. An important type of convex function is a *polyhedral function*,

which by definition is a proper convex function whose epigraph is a polyhedral set. Starting with some known convex functions, we can generate other convex functions by using some common operations that preserve convexity. Principal among these operations are the following:

- (a) Composition with a linear transformation.
- (b) Addition, and multiplication with a nonnegative scalar.
- (c) Taking supremum.
- (d) Taking partial minimum, i.e., minimizing with respect to z a function that is (jointly) convex in two vectors x and z.

The following three propositions deal with the first three cases, and Section 3.3 deals with the fourth.

Proposition 1.1.4: Let $f: \mathbb{R}^m \mapsto (-\infty, \infty]$ be a given function, let A be an $m \times n$ matrix, and let $F: \mathbb{R}^n \mapsto (-\infty, \infty]$ be the function

$$F(x) = f(Ax), \qquad x \in \Re^n.$$

If f is convex, then F is also convex, while if f is closed, then F is also closed.

Proof: Let f be convex. We use the definition of convexity to write for any $x, y \in \mathbb{R}^n$ and $\alpha \in [0, 1]$,

$$F(\alpha x + (1 - \alpha)y) = f(\alpha Ax + (1 - \alpha)Ay)$$

$$\leq \alpha f(Ax) + (1 - \alpha)f(Ay)$$

$$= \alpha F(x) + (1 - \alpha)F(y).$$

Hence F is convex.

Let f be closed. Then f is lower semicontinuous at every $x \in \Re^n$ (cf. Prop. 1.1.2), so for every sequence $\{x_k\}$ converging to x, we have

$$f(Ax) \le \liminf_{k \to \infty} f(Ax_k),$$

or

$$F(x) \le \liminf_{k \to \infty} F(x_k)$$

for all k. It follows that F is lower semicontinuous at every $x \in \Re^n$, and hence is closed by Prop. 1.1.2. **Q.E.D.**

The next proposition deals with sums of function and it is interesting to note that it can be viewed as a special case of the preceding proposition, which deals with compositions with linear transformations. The reason is that we may write a sum $F = f_1 + \cdots + f_m$ in the form F(x) = f(Ax), where A is the matrix defined by $Ax = (x, \dots, x)$, and $f : \Re^{mn} \mapsto (-\infty, \infty]$ is the function given by

$$f(x_1,\ldots,x_m) = f_1(x_1) + \cdots + f_m(x_m).$$

Proposition 1.1.5: Let $f_i: \Re^n \mapsto (-\infty, \infty], i = 1, ..., m$, be given functions, let $\gamma_1, ..., \gamma_m$ be positive scalars, and let $F: \Re^n \mapsto (-\infty, \infty]$ be the function

$$F(x) = \gamma_1 f_1(x) + \dots + \gamma_m f_m(x), \qquad x \in \Re^n.$$

If f_1, \ldots, f_m are convex, then F is also convex, while if f_1, \ldots, f_m are closed, then F is also closed.

Proof: The proof follows closely the one of Prop. 1.1.4. **Q.E.D.**

Proposition 1.1.6: Let $f_i: \Re^n \to (-\infty, \infty]$ be given functions for $i \in I$, where I is an arbitrary index set, and let $f: \Re^n \mapsto (-\infty, \infty]$ be the function given by

$$f(x) = \sup_{i \in I} f_i(x).$$

If f_i , $i \in I$, are convex, then f is also convex, while if f_i , $i \in I$, are closed, then f is also closed.

Proof: A pair (x, w) belongs to $\operatorname{epi}(f)$ if and only if $f(x) \leq w$, which is true if and only if $f_i(x) \leq w$ for all $i \in I$, or equivalently $(x, w) \in \cap_{i \in I} \operatorname{epi}(f_i)$. Therefore,

$$\operatorname{epi}(f) = \bigcap_{i \in I} \operatorname{epi}(f_i).$$

If the functions f_i are convex, the epigraphs $epi(f_i)$ are convex, so epi(f) is convex, and f is convex. If the functions f_i are closed, the epigraphs $epi(f_i)$ are closed, so epi(f) is closed, and f is closed. **Q.E.D.**

1.1.4 Characterizations of Differentiable Convex Functions

For once or twice differentiable functions, there are some additional criteria for verifying convexity, as we will now discuss. A useful alternative characterization of convexity for differentiable functions is given in the following proposition and is illustrated in Fig. 1.1.6.

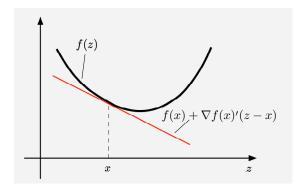


Figure 1.1.6. Characterization of convexity in terms of first derivatives. The condition $f(z) \ge f(x) + \nabla f(x)'(z-x)$ states that a linear approximation, based on the gradient, underestimates a convex function.

Proposition 1.1.7: Let C be a nonempty convex subset of \mathbb{R}^n and let $f: \mathbb{R}^n \to \mathbb{R}$ be differentiable over an open set that contains C.

(a) f is convex over C if and only if

$$f(z) \ge f(x) + \nabla f(x)'(z - x), \quad \forall x, z \in C.$$
 (1.3)

(b) f is strictly convex over C if and only if the above inequality is strict whenever $x \neq z$.

Proof: The ideas of the proof are geometrically illustrated in Fig. 1.1.7. We prove parts (a) and (b) simultaneously. Assume that the inequality (1.3) holds. Choose any $x, y \in C$ and $\alpha \in [0, 1]$, and let $z = \alpha x + (1 - \alpha)y$. Using the inequality (1.3) twice, we obtain

$$f(x) \ge f(z) + \nabla f(z)'(x - z),$$

$$f(y) > f(z) + \nabla f(z)'(y-z).$$

We multiply the first inequality by α , the second by $(1-\alpha)$, and add them to obtain

$$\alpha f(x) + (1 - \alpha)f(y) \ge f(z) + \nabla f(z)'(\alpha x + (1 - \alpha)y - z) = f(z),$$

which proves that f is convex. If the inequality (1.3) is strict as stated in part (b), then if we take $x \neq y$ and $\alpha \in (0,1)$ above, the three preceding inequalities become strict, thus showing the strict convexity of f.

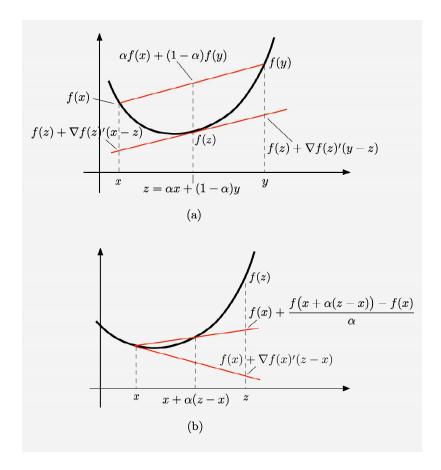


Figure 1.1.7. Geometric illustration of the ideas underlying the proof of Prop. 1.1.7. In figure (a), we linearly approximate f at $z = \alpha x + (1-\alpha)y$. The inequality (1.3) implies that

$$f(x) \ge f(z) + \nabla f(z)'(x-z),$$

$$f(y) \ge f(z) + \nabla f(z)'(y-z).$$

As can be seen from the figure, it follows that $\alpha f(x) + (1-\alpha)f(y)$ lies above f(z), so f is convex.

In figure (b), we assume that f is convex, and from the figure's geometry, we note that

$$f(x) + \frac{f(x + \alpha(z - x)) - f(x)}{\alpha}$$

lies below f(z), is monotonically nonincreasing as $\alpha \downarrow 0,$ and converges to

$$f(x) + \nabla f(x)'(z - x).$$

It follows that

$$f(z) \ge f(x) + \nabla f(x)'(z - x).$$

Conversely, assume that f is convex, let x and z be any vectors in C with $x \neq z$, and consider the function $g:(0,1] \mapsto \Re$ given by

$$g(\alpha) = \frac{f(x + \alpha(z - x)) - f(x)}{\alpha}, \quad \alpha \in (0, 1].$$

We will show that $g(\alpha)$ is monotonically increasing with α , and is strictly monotonically increasing if f is strictly convex. This will imply that

$$\nabla f(x)'(z-x) = \lim_{\alpha \downarrow 0} g(\alpha) \le g(1) = f(z) - f(x),$$

with strict inequality if g is strictly monotonically increasing, thereby showing that the desired inequality (1.3) holds, and holds strictly if f is strictly convex. Indeed, consider any α_1, α_2 , with $0 < \alpha_1 < \alpha_2 < 1$, and let

$$\overline{\alpha} = \frac{\alpha_1}{\alpha_2}, \quad \overline{z} = x + \alpha_2(z - x).$$
 (1.4)

We have

$$f(x + \overline{\alpha}(\overline{z} - x)) \le \overline{\alpha}f(\overline{z}) + (1 - \overline{\alpha})f(x),$$

or

$$\frac{f(x + \overline{\alpha}(\overline{z} - x)) - f(x)}{\overline{\alpha}} \le f(\overline{z}) - f(x), \tag{1.5}$$

and the above inequalities are strict if f is strictly convex. Substituting the definitions (1.4) in Eq. (1.5), we obtain after a straightforward calculation

$$\frac{f(x+\alpha_1(z-x))-f(x)}{\alpha_1} \le \frac{f(x+\alpha_2(z-x))-f(x)}{\alpha_2},$$

or

$$g(\alpha_1) \leq g(\alpha_2),$$

with strict inequality if f is strictly convex. Hence g is monotonically increasing with α , and strictly so if f is strictly convex. **Q.E.D.**

Note a simple consequence of Prop. 1.1.7(a): if $f: \Re^n \to \Re$ is a differentiable convex function and $\nabla f(x^*) = 0$, then x^* minimizes f over \Re^n . This is a classical sufficient condition for unconstrained optimality, originally formulated (in one dimension) by Fermat in 1637. Similarly, from Prop. 1.1.7(a), we see that the condition

$$\nabla f(x^*)'(z - x^*) > 0, \qquad \forall \ z \in C,$$

implies that x^* minimizes a differentiable convex function f over a convex set C. This sufficient condition for optimality is also necessary. To see this,

assume to arrive at a contradiction that x^* minimizes f over C and that $\nabla f(x^*)'(z-x^*)<0$ for some $z\in C$. By differentiation, we have

$$\lim_{\alpha \downarrow 0} \frac{f\left(x^* + \alpha(z - x^*)\right) - f(x^*)}{\alpha} = \nabla f(x^*)'(z - x^*) < 0,$$

so $f(x^* + \alpha(z - x^*))$ decreases strictly for sufficiently small $\alpha > 0$, contradicting the optimality of x^* . We state the conclusion as a proposition.

Proposition 1.1.8: Let C be a nonempty convex subset of \Re^n and let $f: \Re^n \to \Re$ be convex and differentiable over an open set that contains C. Then a vector $x^* \in C$ minimizes f over C if and only if

$$\nabla f(x^*)'(z - x^*) \ge 0, \quad \forall z \in C.$$

Let us use the preceding optimality condition to prove a basic theorem of analysis and optimization.

Proposition 1.1.9: (Projection Theorem) Let C be a nonempty closed convex subset of \Re^n , and let z be a vector in \Re^n . There exists a unique vector that minimizes the Euclidean norm distance ||z-x|| over $x \in C$, called the projection of z on C. Furthermore, a vector x^* is the projection of z on C if and only if

$$(z - x^*)'(x - x^*) \le 0, \quad \forall \ x \in C.$$
 (1.6)

Proof: Minimizing ||z - x|| is equivalent to minimizing the convex and differentiable function

$$f(x) = \frac{1}{2}||z - x||^2$$
.

By Prop. 1.1.8, x^* minimizes f over C if and only if

$$\nabla f(x^*)'(x - x^*) \ge 0, \quad \forall \ x \in C.$$

Since $\nabla f(x^*) = x^* - z$, this condition is equivalent to Eq. (1.6).

Minimizing f over C is equivalent to minimizing f over the compact set $C \cap \{\|z - x\| \le \|z - w\|\}$, where w is any vector in C. By Weierstrass' Theorem (Prop. A.2.7), it follows that there exists a minimizing vector. To show uniqueness, let x_1^* and x_2^* be two minimizing vectors. Then by Eq. (1.6), we have

$$(z-x_1^*)'(x_2^*-x_1^*) \le 0,$$
 $(z-x_2^*)'(x_1^*-x_2^*) \le 0.$

Adding these two inequalities, we obtain

$$(x_2^* - x_1^*)'(x_2^* - x_1^*) = ||x_2^* - x_1^*||^2 \le 0,$$

so $x_2^* = x_1^*$. **Q.E.D.**

For twice differentiable convex functions, there is another characterization of convexity, given by the following proposition.

Proposition 1.1.10: Let C be a nonempty convex subset of \Re^n and let $f: \Re^n \to \Re$ be twice continuously differentiable over an open set that contains C.

- (a) If $\nabla^2 f(x)$ is positive semidefinite for all $x \in C$, then f is convex over C.
- (b) If $\nabla^2 f(x)$ is positive definite for all $x \in C$, then f is strictly convex over C.
- (c) If C is open and f is convex over C, then $\nabla^2 f(x)$ is positive semidefinite for all $x \in C$.

Proof: (a) Using the mean value theorem (Prop. A.3.1), we have for all $x, y \in C$,

$$f(y) = f(x) + \nabla f(x)'(y - x) + \frac{1}{2}(y - x)'\nabla^2 f(x + \alpha(y - x))(y - x)$$

for some $\alpha \in [0,1]$. Therefore, using the positive semidefiniteness of $\nabla^2 f$, we obtain

$$f(y) \ge f(x) + \nabla f(x)'(y - x), \quad \forall x, y \in C.$$

From Prop. 1.1.7(a), we conclude that f is convex over C.

- (b) Similar to the proof of part (a), we have $f(y) > f(x) + \nabla f(x)'(y-x)$ for all $x, y \in C$ with $x \neq y$, and the result follows from Prop. 1.1.7(b).
- (c) Assume, to obtain a contradiction, that there exist some $x \in C$ and some $z \in \Re^n$ such that $z' \nabla^2 f(x) z < 0$. Since C is open and $\nabla^2 f$ is continuous, we can choose z to have small enough norm so that $x + z \in C$ and $z' \nabla^2 f(x + \alpha z) z < 0$ for every $\alpha \in [0,1]$. Then, using again the mean value theorem, we obtain $f(x + z) < f(x) + \nabla f(x)' z$, which, in view of Prop. 1.1.7(a), contradicts the convexity of f over C. Q.E.D.

If f is convex over a convex set C that is not open, $\nabla^2 f(x)$ may not be positive semidefinite at any point of C [take for example n=2, $C=\{(x_1,0)\mid x_1\in\Re\}$, and $f(x)=x_1^2-x_2^2$]. However, it can be shown that the conclusion of Prop. 1.1.10(c) also holds if C has nonempty interior instead of being open.

1.2 CONVEX AND AFFINE HULLS

We now discuss issues relating to the convexification of nonconvex sets. Let X be a nonempty subset of \Re^n . The convex hull of a set X, denoted $\operatorname{conv}(X)$, is the intersection of all convex sets containing X, and is a convex set by Prop. 1.1.1(a). A convex combination of elements of X is a vector of the form $\sum_{i=1}^m \alpha_i x_i$, where m is a positive integer, x_1, \ldots, x_m belong to X, and $\alpha_1, \ldots, \alpha_m$ are scalars such that

$$\alpha_i \ge 0, \quad i = 1, \dots, m, \qquad \sum_{i=1}^{m} \alpha_i = 1.$$

Note that a convex combination belongs to conv(X) (see the construction of Fig. 1.2.1). For any convex combination and function $f: \mathbb{R}^n \to \mathbb{R}$ that is $convex \ over \ conv(X)$, we have

$$f\left(\sum_{i=1}^{m} \alpha_i x_i\right) \le \sum_{i=1}^{m} \alpha_i f(x_i). \tag{1.7}$$

This follows by using repeatedly the definition of convexity together with the construction of Fig. 1.2.1. The preceding relation is a special case of a relation known as *Jensen's inequality*, which finds wide use in applied mathematics and probability theory.

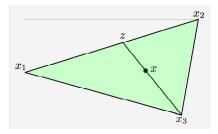


Figure 1.2.1. Construction of a convex combination of m vectors by forming a sequence of m-1 convex combinations of pairs of vectors (first combine two vectors, then combine the result with a third vector, etc). For example,

$$x = \alpha_1 x_1 + \alpha_2 x_2 + \alpha_3 x_3 = (\alpha_1 + \alpha_2) \left(\frac{\alpha_1}{\alpha_1 + \alpha_2} x_1 + \frac{\alpha_2}{\alpha_1 + \alpha_2} x_2 \right) + \alpha_3 x_3,$$

so the convex combination $\alpha_1x_1 + \alpha_2x_2 + \alpha_3x_3$ can be obtained by forming the convex combination

$$z = \frac{\alpha_1}{\alpha_1 + \alpha_2} x_1 + \frac{\alpha_2}{\alpha_1 + \alpha_2} x_2,$$

and then by forming the convex combination $x = (\alpha_1 + \alpha_2)z + \alpha_3x_3$ as shown in the figure. This shows that a convex combination of vectors from a convex set belongs to the set, and that a convex combination of vectors from a nonconvex set belongs to the convex hull of the set.

It is straightforward to verify that the set of all convex combinations of elements of X is equal to conv(X). In particular, if X consists of a finite number of vectors x_1, \ldots, x_m , its convex hull is

$$conv({x_1, ..., x_m}) = \left\{ \sum_{i=1}^m \alpha_i x_i \mid \alpha_i \ge 0, i = 1, ..., m, \sum_{i=1}^m \alpha_i = 1 \right\}.$$

Also, for any set S and linear transformation A, we have conv(AS) = A conv(S). From this it follows that for any sets S_1, \ldots, S_m , we have $conv(S_1 + \cdots + S_m) = conv(S_1) + \cdots + conv(S_m)$.

We recall that an affine set M in \Re^n is a set of the form x+S, where x is some vector, and S is a subspace uniquely determined by M and called the subspace parallel to M. Alternatively, a set M is affine if it contains all the lines that pass through pairs of points $x,y\in M$ with $x\neq y$. If X is a subset of \Re^n , the affine hull of X, denoted aff(X), is the intersection of all affine sets containing X. Note that aff(X) is itself an affine set and that it contains $\operatorname{conv}(X)$. The dimension of $\operatorname{aff}(X)$ is defined to be the dimension of the subspace parallel to $\operatorname{aff}(X)$. It can be shown that

$$\operatorname{aff}(X) = \operatorname{aff}(\operatorname{conv}(X)) = \operatorname{aff}(\operatorname{cl}(X)).$$

For a convex set C, the *dimension* of C is defined to be the dimension of aff(C).

Given a nonempty subset X of \Re^n , a nonnegative combination of elements of X is a vector of the form $\sum_{i=1}^m \alpha_i x_i$, where m is a positive integer, x_1, \ldots, x_m belong to X, and $\alpha_1, \ldots, \alpha_m$ are nonnegative scalars. If the scalars α_i are all positive, $\sum_{i=1}^m \alpha_i x_i$ is said to be a positive combination. The cone generated by X, denoted $\operatorname{cone}(X)$, is the set of all nonnegative combinations of elements of X. It is easily seen that $\operatorname{cone}(X)$ is a convex cone containing the origin, although it need not be closed even if X is compact, as shown in Fig. 1.2.2 [it can be proved that $\operatorname{cone}(X)$ is closed in special cases, such as when X is finite; see Section 1.4.3].

The following is a fundamental characterization of convex hulls (see Fig. 1.2.3).

Proposition 1.2.1: (Caratheodory's Theorem) Let X be a non-empty subset of \Re^n .

- (a) Every nonzero vector from cone(X) can be represented as a positive combination of linearly independent vectors from X.
- (b) Every vector from conv(X) can be represented as a convex combination of no more than n+1 vectors from X.

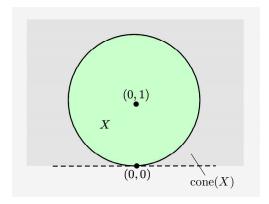


Figure 1.2.2. An example in \Re^2 where X is convex and compact, but cone(X) is not closed. Here

$$X = \left\{ (x_1, x_2) \mid x_1^2 + (x_2 - 1)^2 \le 1 \right\}, \quad \operatorname{cone}(X) = \left\{ (x_1, x_2) \mid x_2 > 0 \right\} \cup \left\{ (0, 0) \right\}.$$

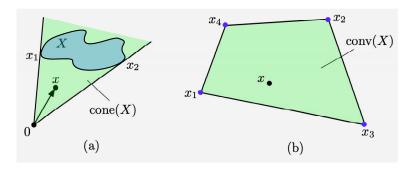


Figure 1.2.3. Illustration of Caratheodory's Theorem. In (a), X is a nonconvex set in \Re^2 , and a point $x \in \text{cone}(X)$ is represented as a positive combination of the two linearly independent vectors $x_1, x_2 \in X$. In (b), X consists of four points x_1, x_2, x_3, x_4 in \Re^2 , and the point $x \in \text{conv}(X)$ shown in the figure can be represented as a convex combination of the three vectors x_1, x_2, x_3 . Note also that x can alternatively be represented as a convex combination of the vectors x_1, x_3, x_4 , so the representation is not unique.

Proof: (a) Consider a vector $x \neq 0$ from $\operatorname{cone}(X)$, and let m be the smallest integer such that x has the form $\sum_{i=1}^m \alpha_i x_i$, where $\alpha_i > 0$ and $x_i \in X$ for all $i = 1, \ldots, m$. We argue by contradiction. If the vectors x_i are linearly dependent, there exist scalars $\lambda_1, \ldots, \lambda_m$, with $\sum_{i=1}^m \lambda_i x_i = 0$ and at least one λ_i is positive. Consider the linear combination $\sum_{i=1}^m (\alpha_i - \overline{\gamma}\lambda_i)x_i$, where $\overline{\gamma}$ is the largest γ such that $\alpha_i - \gamma\lambda_i \geq 0$ for all i. This combination provides a representation of x as a positive combination of fewer than x vectors of x a contradiction. Therefore, x_1, \ldots, x_m are linearly independent.

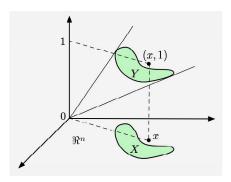


Figure 1.2.4. Illustration of the proof of Caratheodory's Theorem for convex hulls using the version of the theorem for generated cones. We consider the set

$$Y = \left\{ (y, 1) \mid y \in X \right\} \subset \Re^{n+1}$$

and apply Prop. 1.2.1(a).

(b) We apply part (a) to the following subset of \Re^{n+1} :

$$Y = \{(y,1) \mid y \in X\}$$

(cf. Fig. 1.2.4). If $x \in \operatorname{conv}(X)$, we have $x = \sum_{i=1}^{I} \gamma_i x_i$ for an integer I > 0 and scalars $\gamma_i > 0$, $i = 1, \ldots, I$, with $1 = \sum_{i=1}^{I} \gamma_i$, so that $(x, 1) \in \operatorname{cone}(Y)$. By part (a), we have $(x, 1) = \sum_{i=1}^{m} \alpha_i(x_i, 1)$ for some scalars $\alpha_1, \ldots, \alpha_m > 0$ and (at most n+1) linearly independent vectors $(x_1, 1), \ldots, (x_m, 1)$. Thus, $x = \sum_{i=1}^{m} \alpha_i x_i$ and $1 = \sum_{i=1}^{m} \alpha_i$. Q.E.D.

Note that the proof of part (b) of Caratheodory's Theorem shows that if $m \geq 2$, the m vectors $x_1, \ldots, x_m \in X$ used to represent a vector in $\operatorname{conv}(X)$ may be chosen so that $x_2 - x_1, \ldots, x_m - x_1$ are linearly independent [if $x_2 - x_1, \ldots, x_m - x_1$ were linearly dependent, there exist $\lambda_2, \ldots, \lambda_m$, not all 0, with $\sum_{i=2}^m \lambda_i(x_i - x_1) = 0$ so that by defining $\lambda_1 = -(\lambda_2 + \cdots + \lambda_m)$,

$$\sum_{i=1}^{m} \lambda_i(x_i, 1) = 0,$$

contradicting the linear independence of $(x_1, 1), \ldots, (x_m, 1)$.

Caratheodory's Theorem can be used to prove several other important results. An example is the following proposition.

Proposition 1.2.2: The convex hull of a compact set is compact.

Proof: Let X be a compact subset of \Re^n . To show that $\operatorname{conv}(X)$ is compact, we take a sequence in $\operatorname{conv}(X)$ and show that it has a convergent subsequence whose limit is in $\operatorname{conv}(X)$. Indeed, by Caratheodory's Theorem, a sequence in $\operatorname{conv}(X)$ can be expressed as $\left\{\sum_{i=1}^{n+1} \alpha_i^k x_i^k\right\}$, where for all k and $i, \alpha_i^k \geq 0, x_i^k \in X$, and $\sum_{i=1}^{n+1} \alpha_i^k = 1$. Since the sequence

$$\left\{ (\alpha_1^k, \dots, \alpha_{n+1}^k, x_1^k, \dots, x_{n+1}^k) \right\}$$

is bounded, it has a limit point $\{(\alpha_1, \ldots, \alpha_{n+1}, x_1, \ldots, x_{n+1})\}$, which must satisfy $\sum_{i=1}^{n+1} \alpha_i = 1$, and $\alpha_i \geq 0$, $x_i \in X$ for all i. Thus, the vector $\sum_{i=1}^{n+1} \alpha_i x_i$, which belongs to $\operatorname{conv}(X)$, is a limit point of the sequence $\{\sum_{i=1}^{n+1} \alpha_i^k x_i^k\}$, showing that $\operatorname{conv}(X)$ is compact. **Q.E.D.**

Note that the convex hull of an unbounded closed set need not be closed. As an example, for the closed subset of \Re^2

$$X = \{(0,0)\} \cup \{(x_1,x_2) \mid x_1x_2 \ge 1, x_1 \ge 0, x_2 \ge 0\},\$$

the convex hull is

$$conv(X) = \{(0,0)\} \cup \{(x_1,x_2) \mid x_1 > 0, x_2 > 0\},\$$

which is not closed.

We finally note that just as one can convexify nonconvex sets through the convex hull operation, one can also convexify a nonconvex function by convexification of its epigraph. In fact, this can be done in a way that the optimality of the minima of the function is maintained (see Section 1.3.3).

1.3 RELATIVE INTERIOR AND CLOSURE

We now consider some generic topological properties of convex sets and functions. Let C be a nonempty convex subset of \Re^n . The closure of C, denoted $\operatorname{cl}(C)$, is also a nonempty convex set [Prop. 1.1.1(d)]. The interior of C is also convex, but it may be empty. It turns out, however, that convexity implies the existence of interior points relative to the affine hull of C. This is an important property, which we now formalize.

Let C be a nonempty convex set. We say that x is a relative interior point of C if $x \in C$ and there exists an open sphere S centered at x such that $S \cap \operatorname{aff}(C) \subset C$, i.e., x is an interior point of C relative to the affine hull of C. The set of all relative interior points of C is called the relative interior of C, and is denoted by $\operatorname{ri}(C)$. The set C is said to be relatively open if $\operatorname{ri}(C) = C$. The vectors in $\operatorname{cl}(C)$ that are not relative interior points are said to be relative boundary points of C, and their collection is called the relative boundary of C.

For an example, let C be a line segment connecting two distinct points in the plane. Then $\mathrm{ri}(C)$ consists of all points of C except for the two end points, and the relative boundary of C consists of the two end points. For another example, let C be an affine set. Then $\mathrm{ri}(C)=C$ and the relative boundary of C is empty.

The most fundamental fact about relative interiors is given in the following proposition.

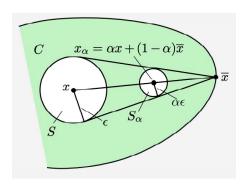


Figure 1.3.1. Proof of the Line Segment Principle for the case where $\overline{x} \in C$. Since $x \in \text{ri}(C)$, there exists an open sphere $S = \{z \mid ||z - x|| < \epsilon\}$ such that $S \cap \text{aff}(C) \subset C$. For all $\alpha \in (0,1]$, let $x_{\alpha} = \alpha x + (1-\alpha)\overline{x}$ and let $S_{\alpha} = \{z \mid ||z - x_{\alpha}|| < \alpha \epsilon\}$. It can be seen that each point of $S_{\alpha} \cap \text{aff}(C)$ is a convex combination of \overline{x} and some point of $S \cap \text{aff}(C)$. Therefore, by the convexity of $S_{\alpha} \cap \text{aff}(C) \subset C$, implying that $x_{\alpha} \in \text{ri}(C)$.

Proposition 1.3.1: (Line Segment Principle) Let C be a nonempty convex set. If $x \in ri(C)$ and $\overline{x} \in cl(C)$, then all points on the line segment connecting x and \overline{x} , except possibly \overline{x} , belong to ri(C).

Proof: For the case where $\overline{x} \in C$, the proof is given in Fig. 1.3.1. Consider the case where $\overline{x} \notin C$. To show that for any $\alpha \in (0,1]$ we have $x_{\alpha} = \alpha x + (1-\alpha)\overline{x} \in \text{ri}(C)$, consider a sequence $\{x_k\} \subset C$ that converges to \overline{x} , and let $x_{k,\alpha} = \alpha x + (1-\alpha)x_k$. Then as in Fig. 1.3.1, we see that $\{z \mid ||z-x_{k,\alpha}|| < \alpha \epsilon\} \cap \text{aff}(C) \subset C$ for all k, where ϵ is such that the open sphere $S = \{z \mid ||z-x|| < \epsilon\}$ satisfies $S \cap \text{aff}(C) \subset C$. Since $x_{k,\alpha} \to x_{\alpha}$, for large enough k, we have

$$\{z \mid ||z - x_{\alpha}|| < \alpha \epsilon/2\} \subset \{z \mid ||z - x_{k,\alpha}|| < \alpha \epsilon\}.$$

It follows that $\{z \mid ||z - x_{\alpha}|| < \alpha \epsilon/2\} \cap \operatorname{aff}(C) \subset C$, which shows that $x_{\alpha} \in \operatorname{ri}(C)$. Q.E.D.

A major consequence of the Line Segment Principle is given in the following proposition.

Proposition 1.3.2: (Nonemptiness of Relative Interior) Let C be a nonempty convex set. Then:

(a) ri(C) is a nonempty convex set, and has the same affine hull as C.

(b) If m is the dimension of $\operatorname{aff}(C)$ and m > 0, there exist vectors $x_0, x_1, \ldots, x_m \in \operatorname{ri}(C)$ such that $x_1 - x_0, \ldots, x_m - x_0$ span the subspace parallel to $\operatorname{aff}(C)$.

Proof: (a) Convexity of $\operatorname{ri}(C)$ follows from the Line Segment Principle (Prop. 1.3.1). By using a translation argument if necessary, we assume without loss of generality that $0 \in C$. Then $\operatorname{aff}(C)$ is a subspace whose dimension will be denoted by m. To show that $\operatorname{ri}(C)$ is nonempty, we will use a basis for $\operatorname{aff}(C)$ to construct a relatively open set.

If the dimension m is 0, then C and $\operatorname{aff}(C)$ consist of a single point, which is a unique relative interior point. If m>0, we can find m linearly independent vectors z_1,\ldots,z_m in C that span $\operatorname{aff}(C)$; otherwise there would exist r< m linearly independent vectors in C whose span contains C, contradicting the fact that the dimension of $\operatorname{aff}(C)$ is m. Thus z_1,\ldots,z_m form a basis for $\operatorname{aff}(C)$.

Consider the set

$$X = \left\{ x \mid x = \sum_{i=1}^{m} \alpha_i z_i, \sum_{i=1}^{m} \alpha_i < 1, \, \alpha_i > 0, \, i = 1, \dots, m \right\}$$

(see Fig. 1.3.2), and note that $X \subset C$ since C is convex. We claim that this set is open relative to $\operatorname{aff}(C)$, i.e., for every vector $\overline{x} \in X$, there exists an open ball B centered at \overline{x} such that $\overline{x} \in B$ and $B \cap \operatorname{aff}(C) \subset X$. To see this, fix $\overline{x} \in X$ and let x be another vector in $\operatorname{aff}(C)$. We have $\overline{x} = Z\overline{\alpha}$ and $x = Z\alpha$, where Z is the $n \times m$ matrix whose columns are the vectors z_1, \ldots, z_m , and $\overline{\alpha}$ and α are suitable m-dimensional vectors, which are unique since z_1, \ldots, z_m form a basis for $\operatorname{aff}(C)$. Since Z has linearly independent columns, the matrix Z'Z is symmetric and positive definite, so for some positive scalar γ , which is independent of x and \overline{x} , we have

$$||x - \overline{x}||^2 = (\alpha - \overline{\alpha})' Z' Z(\alpha - \overline{\alpha}) \ge \gamma ||\alpha - \overline{\alpha}||^2.$$
 (1.8)

Since $\overline{x} \in X$, the corresponding vector $\overline{\alpha}$ lies in the open set

$$A = \left\{ (\alpha_1, \dots, \alpha_m) \mid \sum_{i=1}^m \alpha_i < 1, \, \alpha_i > 0, \, i = 1, \dots, m \right\}.$$

From Eq. (1.8), we see that if x lies in a suitably small ball centered at \overline{x} , the corresponding vector α lies in A, implying that $x \in X$. Hence X contains the intersection of $\operatorname{aff}(C)$ and an open ball centered at \overline{x} , so X is open relative to $\operatorname{aff}(C)$. It follows that all points of X are relative interior points of C, so that $\operatorname{ri}(C)$ is nonempty. Also, since by construction,

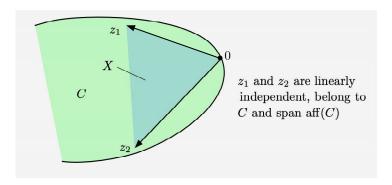


Figure 1.3.2. Construction of the relatively open set X in the proof of nonemptiness of the relative interior of a convex set C that contains the origin, assuming that m > 0. We choose m linearly independent vectors $z_1, \ldots, z_m \in C$, where m is the dimension of aff(C), and we let

$$X = \left\{ \sum_{i=1}^{m} \alpha_i z_i \mid \sum_{i=1}^{m} \alpha_i < 1, \ \alpha_i > 0, \ i = 1, \dots, m \right\}.$$

Any point in X is shown to be a relative interior point of C.

 $\operatorname{aff}(X) = \operatorname{aff}(C)$ and $X \subset \operatorname{ri}(C)$, we see that $\operatorname{ri}(C)$ and C have the same affine hull.

(b) Let x_0 be a relative interior point of C [there exists such a point by part (a)]. Translate C to $C - x_0$ (so that x_0 is translated to the origin), and consider vectors $z_1, \ldots, z_m \in C - x_0$ that span aff $(C - x_0)$, as in the proof of part (a). Let $\alpha \in (0,1)$. Since $0 \in \text{ri}(C - x_0)$, by the Line Segment Principle (Prop. 1.3.1), we have $\alpha z_i \in \text{ri}(C - x_0)$ for all $i = 1, \ldots, m$. It follows that the vectors

$$x_i = x_0 + \alpha z_i, \qquad i = 1, \dots, m,$$

are such that x_1-x_0,\ldots,x_m-x_0 belong to ri(C) and span aff(C). Q.E.D.

Here is another useful consequence of the Line Segment Principle.

Proposition 1.3.3: (Prolongation Lemma) Let C be a nonempty convex set. A vector x is a relative interior point of C if and only if every line segment in C having x as one endpoint can be prolonged beyond x without leaving C [i.e., for every $\overline{x} \in C$, there exists a $\gamma > 0$ such that $x + \gamma(x - \overline{x}) \in C$].

Proof: If $x \in ri(C)$, the given condition clearly holds, using the definition of relative interior point. Conversely, let x satisfy the given condition, and

let \overline{x} be a point in ri(C) (by Prop. 1.3.2, there exists such a point). If $x = \overline{x}$, we are done, so assume that $x \neq \overline{x}$. By the given condition, there is a $\gamma > 0$ such that $y = x + \gamma(x - \overline{x}) \in C$, so that x lies strictly within the line segment connecting \overline{x} and y. Since $\overline{x} \in \text{ri}(C)$ and $y \in C$, by the Line Segment Principle (Prop. 1.3.1), it follows that $x \in \text{ri}(C)$. Q.E.D.

We will see in the following chapters that the notion of relative interior is pervasive in convex optimization and duality theory. As an example, we provide an important characterization of the set of optimal solutions in the case where the cost function is concave.

Proposition 1.3.4: Let X be a nonempty convex subset of \Re^n , let $f: X \mapsto \Re$ be a concave function, and let X^* be the set of vectors where f attains a minimum over X, i.e.,

$$X^* = \left\{ x^* \in X \mid f(x^*) = \inf_{x \in X} f(x) \right\}.$$

If X^* contains a relative interior point of X, then f must be constant over X, i.e., $X^* = X$.

Proof: Let x^* belong to $X^* \cap \text{ri}(X)$, and let x be any vector in X. By the Prolongation Lemma (Prop. 1.3.3), there exists a $\gamma > 0$ such that

$$\hat{x} = x^* + \gamma(x^* - x)$$

belongs to X, implying that

$$x^* = \frac{1}{\gamma + 1}\hat{x} + \frac{\gamma}{\gamma + 1}x$$

(see Fig. 1.3.3). By the concavity of f, we have

$$f(x^*) \ge \frac{1}{\gamma + 1} f(\hat{x}) + \frac{\gamma}{\gamma + 1} f(x),$$

and using $f(\hat{x}) \ge f(x^*)$ and $f(x) \ge f(x^*)$, this shows that $f(x) = f(x^*)$. Q.E.D.

One consequence of the preceding proposition is that a linear cost function f(x) = c'x, with $c \neq 0$, cannot attain a minimum at some interior point of a convex constraint set, since such a function cannot be constant over an open sphere. This will be further discussed in Chapter 2, after we introduce the notion of an extreme point.

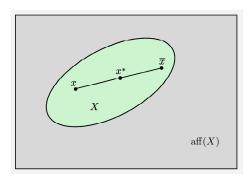


Figure 1.3.3. The idea of the proof of Prop. 1.3.4. If $x^* \in \operatorname{ri}(X)$ minimizes f over X and f is not constant over X, then there exists $x \in X$ such that $f(x) > f(x^*)$. By the Prolongation Lemma (Prop. 1.3.3), there exists $\overline{x} \in X$ such that x^* lies strictly between x and \overline{x} . Since f is concave and $f(x) > f(x^*)$, we must have $f(\overline{x}) < f(x^*)$ - a contradiction of the optimality of x^* .

1.3.1 Calculus of Relative Interiors and Closures

To deal with set operations such as intersection, vector sum, and linear transformation in convex analysis, we need tools for calculating the corresponding relative interiors and closures. These tools are provided in the next five propositions. Here is an informal summary of their content:

- (a) Two convex sets have the same closure if and only if they have the same relative interior.
- (b) Relative interior and closure commute with Cartesian product and inverse image under a linear transformation.
- (c) Relative interior commutes with image under a linear transformation and vector sum, but closure does not.
- (d) Neither closure nor relative interior commute with set intersection, unless the relative interiors of the sets involved have a point in common.

Proposition 1.3.5: Let C be a nonempty convex set.

- (a) $\operatorname{cl}(C) = \operatorname{cl}(\operatorname{ri}(C)).$
- (b) $\operatorname{ri}(C) = \operatorname{ri}(\operatorname{cl}(C))$.
- (c) Let \overline{C} be another nonempty convex set. Then the following three conditions are equivalent:
 - (i) C and \overline{C} have the same relative interior.
 - (ii) C and \overline{C} have the same closure.
 - (iii) $\operatorname{ri}(C) \subset \overline{C} \subset \operatorname{cl}(C)$.

Proof: (a) Since $ri(C) \subset C$, we have $cl(ri(C)) \subset cl(C)$. Conversely, let

 $\overline{x} \in \text{cl}(C)$. We will show that $\overline{x} \in \text{cl}(\text{ri}(C))$. Let x be any point in ri(C) [there exists such a point by Prop. 1.3.2(a)], and assume that $\overline{x} \neq x$ (otherwise we are done). By the Line Segment Principle (Prop. 1.3.1), we have $\alpha x + (1 - \alpha)\overline{x} \in \text{ri}(C)$ for all $\alpha \in (0, 1]$. Thus, \overline{x} is the limit of the sequence

$$\{(1/k)x + (1-1/k)\overline{x} \mid k \ge 1\}$$

that lies in ri(C), so $\overline{x} \in cl(ri(C))$.

- (b) The inclusion $\operatorname{ri}(C) \subset \operatorname{ri}(\operatorname{cl}(C))$ follows from the definition of a relative interior point and the fact $\operatorname{aff}(C) = \operatorname{aff}(\operatorname{cl}(C))$ (the proof of this is left for the reader). To prove the reverse inclusion, let $z \in \operatorname{ri}(\operatorname{cl}(C))$. We will show that $z \in \operatorname{ri}(C)$. By Prop. 1.3.2(a), there exists an $x \in \operatorname{ri}(C)$. We may assume that $x \neq z$ (otherwise we are done). We use the Prolongation Lemma [Prop. 1.3.3, applied within the set $\operatorname{cl}(C)$] to choose $\gamma > 0$, with γ sufficiently close to 0 so that the vector $y = z + \gamma(z x)$ belongs to $\operatorname{cl}(C)$. Then we have $z = (1 \alpha)x + \alpha y$ where $\alpha = 1/(\gamma + 1) \in (0, 1)$, so by the Line Segment Principle (Prop. 1.3.1, applied within the set C), we obtain $z \in \operatorname{ri}(C)$.
- (c) If $\operatorname{ri}(C)=\operatorname{ri}(\overline{C})$, part (a) implies that $\operatorname{cl}(C)=\operatorname{cl}(\overline{C})$. Similarly, if $\operatorname{cl}(C)=\operatorname{cl}(\overline{C})$, part (b) implies that $\operatorname{ri}(C)=\operatorname{ri}(\overline{C})$. Thus, (i) and (ii) are equivalent. Also, (i), (ii), and the relation $\operatorname{ri}(\overline{C})\subset \overline{C}\subset\operatorname{cl}(\overline{C})$ imply condition (iii). Finally, let condition (iii) hold. Then by taking closures, we have $\operatorname{cl}(\operatorname{ri}(C))\subset\operatorname{cl}(\overline{C})\subset\operatorname{cl}(C)$, and by using part (a), we obtain $\operatorname{cl}(C)\subset\operatorname{cl}(\overline{C})\subset\operatorname{cl}(C)$. Hence $\operatorname{cl}(\overline{C})=\operatorname{cl}(C)$, i.e., (ii) holds. Q.E.D.

We now consider the image of a convex set C under a linear transformation A. Geometric intuition suggests that $A \cdot \operatorname{ri}(C) = \operatorname{ri}(A \cdot C)$, since spheres within C are mapped onto ellipsoids within the image $A \cdot C$ (relative to the corresponding affine hulls). This is shown in part (a) of the following proposition. However, the image of a closed convex set under a linear transformation is not closed [see part (b) of the following proposition], and this is a major source of analytical difficulty in convex optimization.

Proposition 1.3.6: Let C be a nonempty convex subset of \mathbb{R}^n and let A be an $m \times n$ matrix.

- (a) We have $A \cdot ri(C) = ri(A \cdot C)$.
- (b) We have $A \cdot \operatorname{cl}(C) \subset \operatorname{cl}(A \cdot C)$. Furthermore, if C is bounded, then $A \cdot \operatorname{cl}(C) = \operatorname{cl}(A \cdot C)$.

Proof: (a) For any set X, we have $A \cdot \operatorname{cl}(X) \subset \operatorname{cl}(A \cdot X)$, since if a sequence $\{x_k\} \subset X$ converges to some $x \in \operatorname{cl}(X)$ then the sequence $\{Ax_k\}$, which belongs to $A \cdot X$, converges to Ax, implying that $Ax \in \operatorname{cl}(A \cdot X)$. We use

this fact and Prop. 1.3.5(a) to write

$$A \cdot ri(C) \subset A \cdot C \subset A \cdot cl(C) = A \cdot cl(ri(C)) \subset cl(A \cdot ri(C)).$$

Thus the convex set $A \cdot C$ lies between the convex set $A \cdot \text{ri}(C)$ and the closure of that set, implying that the relative interiors of the sets $A \cdot C$ and $A \cdot \text{ri}(C)$ are equal [Prop. 1.3.5(c)]. Hence $\text{ri}(A \cdot C) \subset A \cdot \text{ri}(C)$.

To show the reverse inclusion, we take any $z \in A \cdot \mathrm{ri}(C)$ and we show that $z \in \mathrm{ri}(A \cdot C)$. Let x be any vector in $A \cdot C$, and let $\overline{z} \in \mathrm{ri}(C)$ and $\overline{x} \in C$ be such that $A\overline{z} = z$ and $A\overline{x} = x$. By the Prolongation Lemma (Prop. 1.3.3), there exists a $\gamma > 0$ such that the vector $\overline{y} = \overline{z} + \gamma(\overline{z} - \overline{x})$ belongs to C. Thus we have $A\overline{y} \in A \cdot C$ and $A\overline{y} = z + \gamma(z - x)$, so by the Prolongation Lemma, it follows that $z \in \mathrm{ri}(A \cdot C)$.

(b) By the argument given in part (a), we have $A \cdot \operatorname{cl}(C) \subset \operatorname{cl}(A \cdot C)$. To show the converse, assuming that C is bounded, choose any $z \in \operatorname{cl}(A \cdot C)$. Then, there exists a sequence $\{x_k\} \subset C$ such that $Ax_k \to z$. Since C is bounded, $\{x_k\}$ has a subsequence that converges to some $x \in \operatorname{cl}(C)$, and we must have Ax = z. It follows that $z \in A \cdot \operatorname{cl}(C)$. Q.E.D.

Note that if C is closed and convex but unbounded, the set $A \cdot C$ need not be closed [cf. part (b) of the preceding proposition]. For example, projection on the horizontal axis of the closed convex set

$$\{(x_1, x_2) \mid x_1 > 0, x_2 > 0, x_1 x_2 \ge 1\},\$$

shown in Fig. 1.3.4, is the (nonclosed) halfline $\{(x_1, x_2) \mid x_1 > 0, x_2 = 0\}$. Generally, the vector sum of sets C_1, \ldots, C_m can be viewed as the result of the linear transformation $(x_1, \ldots, x_m) \mapsto x_1 + \cdots + x_m$ on the Cartesian product $C_1 \times \cdots \times C_m$. Thus, results involving linear transformations, such as the one of the preceding proposition, yield corresponding

results for vector sums, such as the one of the following proposition.

Proposition 1.3.7: Let C_1 and C_2 be nonempty convex sets. We have

$$ri(C_1 + C_2) = ri(C_1) + ri(C_2), \quad cl(C_1) + cl(C_2) \subset cl(C_1 + C_2).$$

Furthermore, if at least one of the sets C_1 and C_2 is bounded, then

$$cl(C_1) + cl(C_2) = cl(C_1 + C_2).$$

Proof: Consider the linear transformation $A: \mathbb{R}^{2n} \to \mathbb{R}^n$ given by

$$A(x_1, x_2) = x_1 + x_2, \qquad x_1, x_2 \in \Re^n.$$

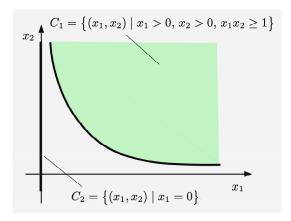


Figure 1.3.4. An example where the sum of two closed convex sets C_1 and C_2 is not closed. Here

$$C_1 = \left\{ (x_1, x_2) \mid x_1 > 0, \, x_2 > 0, \, x_1 x_2 \geq 1 \right\}, \qquad C_2 = \left\{ (x_1, x_2) \mid x_1 = 0 \right\},$$

and $C_1 + C_2$ is the open halfspace $\{(x_1, x_2) \mid x_1 > 0\}$. Also the projection of the set C_1 on the horizontal axis is not closed.

The relative interior of the Cartesian product $C_1 \times C_2$ (viewed as a subset of \Re^{2n}) is $\operatorname{ri}(C_1) \times \operatorname{ri}(C_2)$ (the easy proof of this is left for the reader). Since

$$A(C_1 \times C_2) = C_1 + C_2,$$

from Prop. 1.3.6(a), we obtain $ri(C_1 + C_2) = ri(C_1) + ri(C_2)$.

Similarly, the closure of $C_1 \times C_2$ is $\mathrm{cl}(C_1) \times \mathrm{cl}(C_2)$. From Prop. 1.3.6(b), we have

$$A \cdot \operatorname{cl}(C_1 \times C_2) \subset \operatorname{cl}(A \cdot (C_1 \times C_2)),$$

or equivalently, $\operatorname{cl}(C_1) + \operatorname{cl}(C_2) \subset \operatorname{cl}(C_1 + C_2)$.

To show the reverse inclusion, assuming that C_1 is bounded, let $x \in \operatorname{cl}(C_1+C_2)$. Then there exist sequences $\{x_k^1\} \subset C_1$ and $\{x_k^2\} \subset C_2$ such that $x_k^1+x_k^2 \to x$. Since $\{x_k^1\}$ is bounded, it follows that $\{x_k^2\}$ is also bounded. Thus, $\{(x_k^1,x_k^2)\}$ has a subsequence that converges to a vector (x^1,x^2) , and we have $x^1+x^2=x$. Since $x^1\in\operatorname{cl}(C_1)$ and $x^2\in\operatorname{cl}(C_2)$, it follows that $x\in\operatorname{cl}(C_1)+\operatorname{cl}(C_2)$. Hence $\operatorname{cl}(C_1+C_2)\subset\operatorname{cl}(C_1)+\operatorname{cl}(C_2)$. Q.E.D.

The requirement that at least one of the sets C_1 and C_2 be bounded is essential in the preceding proposition. This is illustrated by the example of Fig. 1.3.4.

Proposition 1.3.8: Let C_1 and C_2 be nonempty convex sets. We have

$$\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2) \subset \operatorname{ri}(C_1 \cap C_2), \qquad \operatorname{cl}(C_1 \cap C_2) \subset \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2).$$

Furthermore, if the sets $ri(C_1)$ and $ri(C_2)$ have a nonempty intersection, then

$$\mathrm{ri}(C_1 \cap C_2) = \mathrm{ri}(C_1) \cap \mathrm{ri}(C_2), \qquad \mathrm{cl}(C_1 \cap C_2) = \mathrm{cl}(C_1) \cap \mathrm{cl}(C_2).$$

Proof: Take any $x \in \text{ri}(C_1) \cap \text{ri}(C_2)$ and any $y \in C_1 \cap C_2$. By the Prolongation Lemma (Prop. 1.3.3), it can be seen that the line segment connecting x and y can be prolonged beyond x by a small amount without leaving C_1 and also by another small amount without leaving C_2 . Thus, by using the lemma again, it follows that $x \in \text{ri}(C_1 \cap C_2)$, so that

$$ri(C_1) \cap ri(C_2) \subset ri(C_1 \cap C_2).$$

Also, since the set $C_1 \cap C_2$ is contained in the closed set $cl(C_1) \cap cl(C_2)$, we have

$$\operatorname{cl}(C_1 \cap C_2) \subset \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2).$$

To show the reverse inclusions assuming that $\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2) \neq \emptyset$, let $y \in \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2)$, and let $x \in \operatorname{ri}(C_1) \cap \operatorname{ri}(C_2)$. By the Line Segment Principle (Prop. 1.3.1), $\alpha x + (1 - \alpha)y \in \operatorname{ri}(C_1) \cap \operatorname{ri}(C_2)$ for all $\alpha \in (0, 1]$ (see Fig. 1.3.5). Hence, y is the limit of a sequence $\alpha_k x + (1 - \alpha_k)y \subset \operatorname{ri}(C_1) \cap \operatorname{ri}(C_2)$ with $\alpha_k \to 0$, implying that $y \in \operatorname{cl}(\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2))$. Thus,

$$\operatorname{cl}(C_1) \cap \operatorname{cl}(C_2) \subset \operatorname{cl}(\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2)) \subset \operatorname{cl}(C_1 \cap C_2).$$

We showed earlier that $\operatorname{cl}(C_1 \cap C_2) \subset \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2)$, so equality holds throughout in the preceding relation, and therefore $\operatorname{cl}(C_1 \cap C_2) = \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2)$. Furthermore, the sets $\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2)$ and $C_1 \cap C_2$ have the same closure. Therefore, by Prop. 1.3.5(c), they have the same relative interior, so that

$$ri(C_1 \cap C_2) = ri(ri(C_1) \cap ri(C_2)) \subset ri(C_1) \cap ri(C_2).$$

We showed earlier the reverse inclusion, so $ri(C_1 \cap C_2) = ri(C_1) \cap ri(C_2)$. **Q.E.D.**

The requirement that $ri(C_1) \cap ri(C_2) \neq \emptyset$ is essential in part (a) of the preceding proposition. As an example, consider the following subsets of the real line:

$$C_1 = \{x \mid x \ge 0\}, \qquad C_2 = \{x \mid x \le 0\}.$$

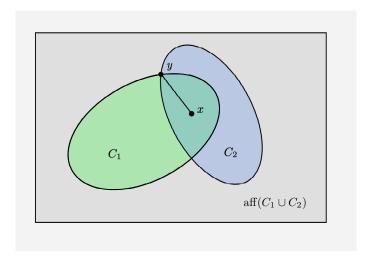


Figure 1.3.5. Construction used to show that

$$\operatorname{cl}(C_1) \cap \operatorname{cl}(C_2) \subset \operatorname{cl}(C_1 \cap C_2),$$

assuming that there exists $x \in ri(C_1) \cap ri(C_2)$ (cf. Prop. 1.3.8). Any $y \in cl(C_1) \cap cl(C_2)$ can be approached along the line segment of $ri(C_1) \cap ri(C_2)$ connecting it with x, so it belongs to the closure of $ri(C_1) \cap ri(C_2)$ and hence also to $cl(C_1 \cap C_2)$.

Then we have $ri(C_1 \cap C_2) = \{0\} \neq \emptyset = ri(C_1) \cap ri(C_2)$. Also, consider the following subsets of the real line:

$$C_1 = \{x \mid x > 0\}, \qquad C_2 = \{x \mid x < 0\}.$$

Then we have $\operatorname{cl}(C_1 \cap C_2) = \emptyset \neq \{0\} = \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2)$.

Proposition 1.3.9: Let C be a nonempty convex subset of \Re^m , and let A be an $m \times n$ matrix. If $A^{-1} \cdot \text{ri}(C)$ is nonempty, then

$$ri(A^{-1} \cdot C) = A^{-1} \cdot ri(C), \qquad cl(A^{-1} \cdot C) = A^{-1} \cdot cl(C),$$

where A^{-1} denotes inverse image of the corresponding set under A.

Proof: Define the sets

$$D = \Re^n \times C, \qquad S = \{(x, Ax) \mid x \in \Re^n\},\$$

and let T be the linear transformation that maps $(x, y) \in \mathbb{R}^{n+m}$ into $x \in \mathbb{R}^n$. We have

$$A^{-1} \cdot C = \{x \mid Ax \in C\} = T \cdot \{(x, Ax) \mid Ax \in C\} = T \cdot (D \cap S),$$

from which

$$ri(A^{-1} \cdot C) = ri(T \cdot (D \cap S)). \tag{1.9}$$

Similarly, we have

$$A^{-1} \cdot \operatorname{ri}(C) = \left\{ x \mid Ax \in \operatorname{ri}(C) \right\} = T \cdot \left\{ (x, Ax) \mid Ax \in \operatorname{ri}(C) \right\} = T \cdot \left(\operatorname{ri}(D) \cap S \right), \tag{1.10}$$

where the last equality holds because $ri(D) = \Re^n \times ri(C)$ (cf. Prop. 1.3.8). Since by assumption, $A^{-1} \cdot ri(C)$ is nonempty, we see that $ri(D) \cap S$ is nonempty. Therefore, using the fact ri(S) = S, and Props. 1.3.6(a) and 1.3.8, it follows that

$$ri(T \cdot (D \cap S)) = T \cdot ri(D \cap S) = T \cdot (ri(D) \cap S). \tag{1.11}$$

Combining Eqs. (1.9)-(1.11), we obtain

$$ri(A^{-1} \cdot C) = A^{-1} \cdot ri(C).$$

To show the second relation, note that

$$A^{-1} \cdot \operatorname{cl}(C) = \{ x \mid Ax \in \operatorname{cl}(C) \} = T \cdot \{ (x, Ax) \mid Ax \in \operatorname{cl}(C) \} = T \cdot (\operatorname{cl}(D) \cap S),$$

where the last equality holds because $cl(D) = \Re^n \times cl(C)$. Since $ri(D) \cap S$ is nonempty and ri(S) = S, it follows from Prop. 1.3.8 that

$$cl(D) \cap S = cl(D \cap S).$$

Using the last two relations and the continuity of T, we obtain

$$A^{-1} \cdot \operatorname{cl}(C) = T \cdot \operatorname{cl}(D \cap S) \subset \operatorname{cl}(T \cdot (D \cap S)),$$

which combined with Eq. (1.9) yields

$$A^{-1} \cdot \operatorname{cl}(C) \subset \operatorname{cl}(A^{-1} \cdot C).$$

To show the reverse inclusion, let \overline{x} be a vector in $\operatorname{cl}(A^{-1} \cdot C)$. Then there exists a sequence $\{x_k\}$ converging to \overline{x} such that $Ax_k \in C$ for all k. Since $\{x_k\}$ converges to \overline{x} , we see that $\{Ax_k\}$ converges to $A\overline{x}$, so that $A\overline{x} \in cl(C)$, or equivalently, $\overline{x} \in A^{-1} \cdot cl(C)$. Q.E.D.

We finally show a useful characterization of the relative interior of sets involving two variables. It generalizes the Cartesian product formula

$$ri(C_1 \times C_2) = ri(C_1) \times ri(C_2)$$

for two convex sets $C_1 \subset \Re^n$ and $C_2 \in \Re^m$.

Proposition 1.3.10: Let C be a convex subset of \Re^{n+m} . For $x \in \Re^n$, denote

$$C_x = \{ y \mid (x, y) \in C \},\$$

and let

$$D = \{x \mid C_x \neq \emptyset\}.$$

Then

$$ri(C) = \{(x,y) \mid x \in ri(D), y \in ri(C_x)\}.$$

Proof: Since D is the projection of C on the x-axis, from Prop. 1.3.6,

$$ri(D) = \{x \mid \text{there exists } y \in \Re^m \text{ with } (x, y) \in ri(C)\},$$

so that

$$\operatorname{ri}(C) = \bigcup_{x \in \operatorname{ri}(D)} \Big(M_x \cap \operatorname{ri}(C) \Big),$$

where $M_x = \{(x, y) \mid y \in \Re^m \}$. For every $x \in ri(D)$, we have

$$M_x \cap \operatorname{ri}(C) = \operatorname{ri}(M_x \cap C) = \{(x, y) \mid y \in \operatorname{ri}(C_x)\},\$$

where the first equality follows from Prop. 1.3.8. By combining the preceding two equations, we obtain the desired result. Q.E.D.

1.3.2 Continuity of Convex Functions

We now derive a basic continuity property of convex functions.

Proposition 1.3.11: If $f: \Re^n \mapsto \Re$ is convex, then it is continuous. More generally, if $f: \Re^n \mapsto (-\infty, \infty]$ is a proper convex function, then f, restricted to dom(f), is continuous over the relative interior of dom(f).

Proof: Restricting attention to the affine hull of dom(f) and using a transformation argument if necessary, we assume without loss of generality that the origin is an interior point of dom(f) and that the unit cube

$$X = \{x \mid ||x||_{\infty} \le 1\}$$

is contained in dom(f) (we use the norm $||x||_{\infty} = \max_{j \in \{1,...,n\}} |x_j|$). It will suffice to show that f is continuous at 0, i.e., that for any sequence $\{x_k\} \subset \mathbb{R}^n$ that converges to 0, we have $f(x_k) \to f(0)$.

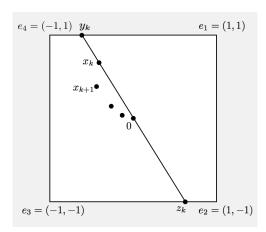


Figure 1.3.6. Construction for the proof of continuity of a real-valued convex function (cf. Prop. 1.3.11).

Let e_i , $i=1,\ldots,2^n$, be the corners of X, i.e., each e_i is a vector whose entries are either 1 or -1. It can be seen that any $x\in X$ can be expressed in the form $x=\sum_{i=1}^{2^n}\alpha_ie_i$, where each α_i is a nonnegative scalar and $\sum_{i=1}^{2^n}\alpha_i=1$. Let $A=\max_i f(e_i)$. From Jensen's inequality [Eq. (1.7)], it follows that $f(x)\leq A$ for every $x\in X$.

For the purpose of proving continuity at 0, we can assume that $x_k \in X$ and $x_k \neq 0$ for all k. Consider the sequences $\{y_k\}$ and $\{z_k\}$ given by

$$y_k = \frac{x_k}{\|x_k\|_{\infty}}, \qquad z_k = -\frac{x_k}{\|x_k\|_{\infty}};$$

(cf. Fig. 1.3.6). Using the definition of a convex function for the line segment that connects y_k , x_k , and 0, we have

$$f(x_k) \le (1 - ||x_k||_{\infty}) f(0) + ||x_k||_{\infty} f(y_k).$$

Since $||x_k||_{\infty} \to 0$ and $f(y_k) \le A$ for all k, by taking the limit as $k \to \infty$, we obtain

$$\limsup_{k \to \infty} f(x_k) \le f(0).$$

Using the definition of a convex function for the line segment that connects x_k , 0, and z_k , we have

$$f(0) \le \frac{\|x_k\|_{\infty}}{\|x_k\|_{\infty} + 1} f(z_k) + \frac{1}{\|x_k\|_{\infty} + 1} f(x_k)$$

and letting $k \to \infty$, we obtain

$$f(0) \le \liminf_{k \to \infty} f(x_k).$$

Thus, $\lim_{k\to\infty} f(x_k) = f(0)$ and f is continuous at zero. Q.E.D.

Among other things, the proposition implies that a real-valued convex function is continuous and hence closed (cf. Prop. 1.1.2). We also have the following stronger result for the case of a function of one variable.

Proposition 1.3.12: If C is a closed interval of the real line, and $f: C \mapsto \Re$ is closed and convex, then f is continuous over C.

Proof: By the preceding proposition, f is continuous in the relative interior of C. To show continuity at a boundary point \overline{x} , let $\{x_k\} \subset C$ be a sequence that converges to \overline{x} , and write

$$x_k = \alpha_k x_0 + (1 - \alpha_k) \overline{x}, \quad \forall k$$

where $\{\alpha_k\}$ is a nonnegative sequence with $\alpha_k \to 0$. By convexity of f, we have for all k such that $\alpha_k \le 1$,

$$f(x_k) \le \alpha_k f(x_0) + (1 - \alpha_k) f(\overline{x}),$$

and by taking the limit as $k \to \infty$, we obtain

$$\limsup_{k \to \infty} f(x_k) \le f(\overline{x}).$$

Consider the function $\tilde{f}: \Re \mapsto (-\infty, \infty]$, which takes the value f(x) for $x \in C$ and ∞ for $x \notin C$, and note that it is closed (since it has the same epigraph as f), and hence lower semicontinuous (cf. Prop. 1.1.2). It follows that $f(\overline{x}) \leq \liminf_{k \to \infty} f(x_k)$, thus implying that $f(x_k) \to f(\overline{x})$, and that f is continuous at \overline{x} . Q.E.D.

1.3.3 Closures of Functions

In this section, we study operations that can transform a given function to a closed and/or convex function, while preserving much of its essential character. These operations play an important role in optimization and other contexts.

A nonempty subset E of \Re^{n+1} is the epigraph of some function if for every $(\overline{x}, \overline{w}) \in E$, the set $\{w \mid (\overline{x}, w) \in E\}$ is either the real line or else it is a halfline that is bounded below and contains its (lower) endpoint. Then E is the epigraph of the function $f: D \mapsto [-\infty, \infty]$, where

$$D = \{x \mid \text{there exists } w \in \Re \text{ with } (x, w) \in E\},\$$

and

$$f(x) = \inf\{w \mid (x, w) \in E\}, \quad \forall x \in D$$

[the infimum is actually attained if f(x) is finite]. Note that E is also the epigraph of other functions with different domain than f (but the same effective domain); for example, $\tilde{f}: \Re^n \mapsto [-\infty, \infty]$, where $\tilde{f}(x) = f(x)$ for $x \in D$ and $\tilde{f}(x) = \infty$ for $x \notin D$. If E is the empty set, it is the epigraph of the function that is identically equal to ∞ .

The closure of the epigraph of a function $f: X \mapsto [-\infty, \infty]$ can be seen to be a legitimate epigraph of another function. This function, called the *closure of* f and denoted closure of f and closure

$$(\operatorname{cl} f)(x) = \inf\{w \mid (x, w) \in \operatorname{cl}(\operatorname{epi}(f))\}, \qquad x \in \Re^n.$$

When f is convex, the set cl(epi(f)) is closed and convex [since the closure of a convex set is convex by Prop. 1.1.1(d)], implying that cl f is closed and convex since epi(cl f) = cl(epi(f)) by definition.

The closure of the convex hull of the epigraph of f is the epigraph of some function, denoted $\check{\operatorname{cl}} f$ called the *convex closure of* f. It can be seen that $\check{\operatorname{cl}} f$ is the closure of the function $F: \Re^n \mapsto [-\infty, \infty]$ given by

$$F(x) = \inf\{w \mid (x, w) \in \operatorname{conv}(\operatorname{epi}(f))\}, \qquad x \in \Re^n.$$
 (1.12)

It is easily shown that F is convex, but it need not be closed and its domain may be strictly contained in $\operatorname{dom}(\check{\operatorname{cl}} f)$ (it can be seen though that the closures of the domains of F and $\check{\operatorname{cl}} f$ coincide).

From the point of view of optimization, an important property is that the minimal values of f, $\operatorname{cl} f$, F, and $\operatorname{\check{cl}} f$ coincide, as stated in the following proposition:

Proposition 1.3.13: Let $f: X \mapsto [-\infty, \infty]$ be a function. Then

$$\inf_{x\in X}f(x)=\inf_{x\in X}(\operatorname{cl} f)(x)=\inf_{x\in\Re^n}(\operatorname{cl} f)(x)=\inf_{x\in\Re^n}F(x)=\inf_{x\in\Re^n}(\operatorname{\check{cl}} f)(x),$$

where F is given by Eq. (1.12). Furthermore, any vector that attains the infimum of f over X also attains the infimum of $\operatorname{cl} f$, F, and $\operatorname{cl} f$.

[†] A note regarding the definition of closure: in Rockafellar [Roc70], p. 52, what we call "closure" of f is called the "lower semi-continuous hull" of f, and "closure" of f is defined somewhat differently (but denoted $\operatorname{cl} f$). Our definition of "closure" of f works better for our purposes, and results in a more streamlined analysis. It coincides with the one of [Roc70] when f is proper convex. For this reason the results of this section correspond to results in [Roc70] only in the case where the functions involved are proper convex. In Rockafellar and Wets [RoW98], p. 14, our "closure" of f is called the "lsc regularization" or "lower closure" of f, and is denoted by $\operatorname{cl} f$. Thus our notation is consistent with the one of [RoW98].

Proof: If $\operatorname{epi}(f)$ is empty, i.e., $f(x) = \infty$ for all x, the results trivially hold. Assume that $\operatorname{epi}(f)$ is nonempty, and let $f^* = \inf_{x \in \Re^n}(\operatorname{cl} f)(x)$. For any sequence $\{(\overline{x}_k, \overline{w}_k)\} \subset \operatorname{cl}(\operatorname{epi}(f))$ with $\overline{w}_k \to f^*$, we can construct a sequence $\{(x_k, w_k)\} \subset \operatorname{epi}(f)$ such that $|w_k - \overline{w}_k| \to 0$, so that $w_k \to f^*$. Since $x_k \in X$, $f(x_k) \leq w_k$, we have

$$\limsup_{k \to \infty} f(x_k) \le f^* \le (\operatorname{cl} f)(x) \le f(x), \qquad \forall \ x \in X,$$

so that

$$\inf_{x\in X}f(x)=\inf_{x\in X}(\operatorname{cl} f)(x)=\inf_{x\in\Re^n}(\operatorname{cl} f)(x).$$

Choose $\{(x_k, w_k)\}\subset \operatorname{conv}(\operatorname{epi}(f))$ with $w_k\to \inf_{x\in\Re^n}F(x)$. Each (x_k, w_k) is a convex combination of vectors from $\operatorname{epi}(f)$, so that $w_k\geq \inf_{x\in X}f(x)$. Hence $\inf_{x\in\Re^n}F(x)\geq \inf_{x\in X}f(x)$. On the other hand, we have $F(x)\leq f(x)$ for all $x\in X$, so it follows that $\inf_{x\in\Re^n}F(x)=\inf_{x\in X}f(x)$. Since $\operatorname{cl} f$ is the closure of F, it also follows (based on what was shown in the preceding paragraph) that $\inf_{x\in\Re^n}(\operatorname{cl} f)(x)=\inf_{x\in\Re^n}F(x)$.

We have $f(x) \ge (\operatorname{cl} f)(x)$ for all x, so if x^* attains the infimum of f,

$$\inf_{x \in \Re^n} (\operatorname{cl} f)(x) = \inf_{x \in X} f(x) = f(x^*) \ge (\operatorname{cl} f)(x^*),$$

showing that x^* attains the infimum of cl f. Similarly, x^* attains the infimum of F and cl f. Q.E.D.

The following is a characterization of closures and convex closures.

Proposition 1.3.14: Let $f: \mathbb{R}^n \mapsto [-\infty, \infty]$ be a function.

- (a) cl f is the greatest closed function majorized by f, i.e., if $g: \Re^n \mapsto [-\infty, \infty]$ is closed and satisfies $g(x) \leq f(x)$ for all $x \in \Re^n$, then $g(x) \leq (\operatorname{cl} f)(x)$ for all $x \in \Re^n$.
- (b) cl f is the greatest closed and convex function majorized by f, i.e., if $g: \Re^n \mapsto [-\infty, \infty]$ is closed and convex, and satisfies $g(x) \leq f(x)$ for all $x \in \Re^n$, then $g(x) \leq (\operatorname{cl} f)(x)$ for all $x \in \Re^n$.

Proof: (a) Let $g: \Re^n \mapsto [-\infty, \infty]$ be closed and such that $g(x) \leq f(x)$ for all x. Then $\operatorname{epi}(f) \subset \operatorname{epi}(g)$. Since $\operatorname{epi}(\operatorname{cl} f) = \operatorname{cl}(\operatorname{epi}(f))$, we have that $\operatorname{epi}(\operatorname{cl} f)$ is the intersection of all closed sets $E \subset \Re^{n+1}$ with $\operatorname{epi}(f) \subset E$, so that $\operatorname{epi}(\operatorname{cl} f) \subset \operatorname{epi}(g)$. It follows that $g(x) \leq (\operatorname{cl} f)(x)$ for all $x \in \Re^n$.

(b) Similar to the proof of part (a). Q.E.D.

Working with the closure of a convex function is often useful because in some sense the closure "differs minimally" from the original. In particular, we can show that a convex function coincides with its closure on the relative interior of its domain. This and other properties of closures are derived in the following proposition.

Proposition 1.3.15: Let $f: \Re^n \mapsto [-\infty, \infty]$ be a convex function. Then:

(a) We have

$$cl(dom(f)) = cl(dom(cl f)), \quad ri(dom(f)) = ri(dom(cl f)),$$

$$(\operatorname{cl} f)(x) = f(x), \quad \forall \ x \in \operatorname{ri}(\operatorname{dom}(f)).$$

Furthermore, $\operatorname{cl} f$ is proper if and only if f is proper.

(b) If $x \in ri(dom(f))$, we have

$$(\operatorname{cl} f)(y) = \lim_{\alpha \downarrow 0} f(y + \alpha(x - y)), \quad \forall y \in \Re^n.$$

Proof: (a) From Prop. 1.3.10, we have

$$\operatorname{ri}(\operatorname{epi}(f)) = \{(x, w) \mid x \in \operatorname{ri}(\operatorname{dom}(f)), f(x) < w\}, \tag{1.13}$$

$$\operatorname{ri}(\operatorname{epi}(\operatorname{cl} f)) = \{(x, w) \mid x \in \operatorname{ri}(\operatorname{dom}(\operatorname{cl} f)), (\operatorname{cl} f)(x) < w\}. \tag{1.14}$$

Since $\operatorname{epi}(f)$ and $\operatorname{epi}(\operatorname{cl} f)$ have the same closure, they have the same relative interior [Prop. 1.3.5(c)], i.e., the sets of Eqs. (1.13) and (1.14) are equal. Hence $\operatorname{dom}(f)$ and $\operatorname{dom}(\operatorname{cl} f)$ have the same relative interior and therefore also the same closure. Thus, the equality of the sets (1.13) and (1.14) yields

$$\begin{split} \big\{(x,w) \mid x \in \mathrm{ri}\big(\mathrm{dom}(f)\big), f(x) < w \big\} \\ &= \big\{(x,w) \mid x \in \mathrm{ri}\big(\mathrm{dom}(f)\big), \, (\mathrm{cl}\, f)(x) < w \big\}, \end{split}$$

from which it follows that $f(x) = (\operatorname{cl} f)(x)$ for all $x \in \operatorname{ri}(\operatorname{dom}(f))$.

If cl f is proper, clearly f is proper. Conversely, if cl f is improper, then $(\operatorname{cl} f)(x) = -\infty$ for all $x \in \operatorname{dom}(\operatorname{cl} f)$ (cf. the discussion at the end of Section 1.1.2). Hence $(\operatorname{cl} f)(x) = -\infty$ for all $x \in \operatorname{ri}(\operatorname{dom}(\operatorname{cl} f)) = \operatorname{ri}(\operatorname{dom}(f))$. Using what was just proved, it follows that $f(x) = (\operatorname{cl} f)(x) = -\infty$ for all $x \in \operatorname{ri}(\operatorname{dom}(f))$, implying that f is improper.

(b) Assume first $y \notin \text{dom}(\text{cl } f)$, i.e., $(\text{cl } f)(y) = \infty$. Then, by the lower semicontinuity of cl f, we have $(\text{cl } f)(y_k) \to \infty$ for all sequences $\{y_k\}$ with $y_k \to y$, from which $f(y_k) \to \infty$, since $(\text{cl } f)(y_k) \le f(y_k)$. Hence $(\text{cl } f)(y) = \lim_{\alpha \downarrow 0} f(y + \alpha(x - y)) = \infty$.

Assume next that $y \in \text{dom}(\operatorname{cl} f),$ and consider the function $g:[0,1] \mapsto \Re$ given by

$$g(\alpha) = (\operatorname{cl} f)(y + \alpha(x - y)).$$

For $\alpha \in (0,1]$, by the Line Segment Principle (Prop. 1.3.1), we have

$$y + \alpha(x - y) \in \operatorname{ri}(\operatorname{dom}(\operatorname{cl} f)),$$

so by part (a), $y + \alpha(x - y) \in ri(dom(f))$, and

$$g(\alpha) = (\operatorname{cl} f)(y + \alpha(x - y)) = f(y + \alpha(x - y)). \tag{1.15}$$

If $(\operatorname{cl} f)(y) = -\infty$, then $\operatorname{cl} f$ is improper and $(\operatorname{cl} f)(z) = -\infty$ for all $z \in \operatorname{dom}(\operatorname{cl} f)$, since an improper closed convex function cannot take a finite value at any point (cf. the discussion at the end of Section 1.1.2). Hence

$$f(y + \alpha(x - y)) = -\infty, \quad \forall \alpha \in (0, 1],$$

and the desired equation follows. If $(\operatorname{cl} f)(y) > -\infty$, then $(\operatorname{cl} f)(y)$ is finite, so $\operatorname{cl} f$ is proper and by part (a), f is also proper. It follows that the function g is real-valued, convex, and closed, and hence also continuous over [0,1] (Prop. 1.3.12). By taking the limit in Eq. (1.15),

$$(\operatorname{cl} f)(y) = g(0) = \lim_{\alpha \downarrow 0} g(\alpha) = \lim_{\alpha \downarrow 0} f(y + \alpha(x - y)).$$

Q.E.D.

Note a corollary of part (a) of the preceding proposition: an improper convex function f takes the value $-\infty$ at all $x \in \text{ri}(\text{dom}(f))$, since its closure does (cf. the discussion at the end of Section 1.1.2).

Calculus of Closure Operations

We now characterize the closure of functions obtained by linear composition and summation of convex functions.

Proposition 1.3.16: Let $f: \Re^m \mapsto [-\infty, \infty]$ be a convex function and A be an $m \times n$ matrix such that the range of A contains a point in ri(dom(f)). The function F defined by

$$F(x) = f(Ax),$$

is convex and

$$(\operatorname{cl} F)(x) = (\operatorname{cl} f)(Ax), \quad \forall \ x \in \Re^n.$$

Proof: Let z be a point in the range of A that belongs to $\operatorname{ri}(\operatorname{dom}(f))$, and let y be such that Ay = z. Then, since $\operatorname{dom}(F) = A^{-1}\operatorname{dom}(f)$ and by Prop. 1.3.9, $\operatorname{ri}(\operatorname{dom}(F)) = A^{-1}\operatorname{ri}(\operatorname{dom}(f))$, we see that $y \in \operatorname{ri}(\operatorname{dom}(F))$. By using Prop. 1.3.15(b), we have for every $x \in \Re^n$,

$$(\operatorname{cl} F)(x) = \lim_{\alpha \downarrow 0} F(x + \alpha(y - x)) = \lim_{\alpha \downarrow 0} f(Ax + \alpha(Ay - Ax)) = (\operatorname{cl} f)(Ax).$$

Q.E.D.

The following proposition is essentially a special case of the preceding one (cf. the discussion in Section 1.1.3).

Proposition 1.3.17: Let $f_i: \Re^n \mapsto [-\infty, \infty], i = 1, \dots, m$, be convex functions such that

$$\bigcap_{i=1}^{m} \operatorname{ri}(\operatorname{dom}(f_i)) \neq \emptyset. \tag{1.16}$$

The function F defined by

$$F(x) = f_1(x) + \dots + f_m(x),$$

is convex and

$$(\operatorname{cl} F)(x) = (\operatorname{cl} f_1)(x) + \dots + (\operatorname{cl} f_m)(x), \quad \forall \ x \in \Re^n.$$

Proof: We write F in the form F(x) = f(Ax), where A is the matrix defined by Ax = (x, ..., x), and $f: \Re^{mn} \mapsto (-\infty, \infty]$ is the function

$$f(x_1, \ldots, x_m) = f_1(x_1) + \cdots + f_m(x_m).$$

Since $dom(F) = \bigcap_{i=1}^{n} dom(f_i)$, Eq. (1.16) implies that

$$\bigcap_{i=1}^{n} \operatorname{ri}(\operatorname{dom}(f_i)) = \operatorname{ri}(\operatorname{dom}(F)) = \operatorname{ri}(A^{-1} \cdot \operatorname{dom}(f)) = A^{-1} \cdot \operatorname{ri}(\operatorname{dom}(f))$$

(cf. Props. 1.3.8 and 1.3.9). Thus Eq. (1.16) is equivalent to the range of A containing a point in $\operatorname{ri}(\operatorname{dom}(f))$, so that $(\operatorname{cl} F)(x) = (\operatorname{cl} f)(x, \ldots, x)$ (cf. Prop. 1.3.16). Let $y \in \cap_{i=1}^n \operatorname{ri}(\operatorname{dom}(f_i))$, so that $(y, \ldots, y) \in \operatorname{ri}(\operatorname{dom}(f))$. Then, from Prop. 1.3.15(b), $(\operatorname{cl} F)(x) = \lim_{\alpha \downarrow 0} f_1(x + \alpha(y - x)) + \cdots + \lim_{\alpha \downarrow 0} f_m(x + \alpha(y - x)) = (\operatorname{cl} f_1)(x) + \cdots + (\operatorname{cl} f_m)(x)$. Q.E.D.

Note that the relative interior assumption (1.16) is essential. To see this, let f_1 and f_2 be the indicator functions of two convex sets C_1 and C_2 such that $\operatorname{cl}(C_1 \cap C_2) \neq \operatorname{cl}(C_1) \cap \operatorname{cl}(C_2)$ (cf. the example following Prop. 1.3.8).

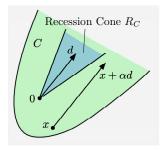


Figure 1.4.1. Illustration of the recession cone R_C of a convex set C. A direction of recession d has the property that $x + \alpha d \in C$ for all $x \in C$ and $\alpha \geq 0$.

1.4 RECESSION CONES

We will now develop some methodology to characterize the asymptotic behavior of convex sets and functions. This methodology is fundamental in several convex optimization contexts, including the issue of existence of optimal solutions, which will be discussed in Chapter 3.

Given a nonempty convex set C, we say that a vector d is a direction of recession of C if $x + \alpha d \in C$ for all $x \in C$ and $\alpha \geq 0$. Thus, d is a direction of recession of C if starting at any x in C and going indefinitely along d, we never cross the relative boundary of C to points outside C.

The set of all directions of recession is a cone containing the origin. It is called the *recession cone* of C and it is denoted by R_C (see Fig. 1.4.1). Thus $d \in R_C$ if $x + \alpha d \in C$ for all $x \in C$ and $\alpha \geq 0$. An important property of a *closed* convex set is that to test whether $d \in R_C$ it is enough to verify the property $x + \alpha d \in C$ for a *single* $x \in C$. This is part (b) of the following proposition.

Proposition 1.4.1: (Recession Cone Theorem) Let C be a nonempty closed convex set.

- (a) The recession cone R_C is closed and convex.
- (b) A vector d belongs to R_C if and only if there exists a vector $x \in C$ such that $x + \alpha d \in C$ for all $\alpha \geq 0$.

Proof: (a) If d_1, d_2 belong to R_C and γ_1, γ_2 are positive scalars such that $\gamma_1 + \gamma_2 = 1$, we have for any $x \in C$ and $\alpha \geq 0$

$$x + \alpha(\gamma_1 d_1 + \gamma_2 d_2) = \gamma_1(x + \alpha d_1) + \gamma_2(x + \alpha d_2) \in C,$$

where the last inclusion holds because C is convex, and $x + \alpha d_1$ and $x + \alpha d_2$ belong to C by the definition of R_C . Hence $\gamma_1 d_1 + \gamma_2 d_2 \in R_C$, implying that R_C is convex.

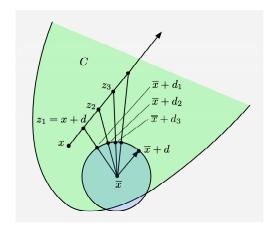


Figure 1.4.2. Construction for the proof of Prop. 1.4.1(b).

Let d be in the closure of R_C , and let $\{d_k\} \subset R_C$ be a sequence converging to d. For any $x \in C$ and $\alpha \geq 0$ we have $x + \alpha d_k \in C$ for all k, and because C is closed, $x + \alpha d \in C$. Hence $d \in R_C$, so R_C is closed.

(b) If $d \in R_C$, every vector $x \in C$ has the required property by the definition of R_C . Conversely, let d be such that there exists a vector $x \in C$ with $x + \alpha d \in C$ for all $\alpha \geq 0$. With no loss of generality, we assume that $d \neq 0$. We choose arbitrary $\overline{x} \in C$ and $\alpha > 0$, and we will show that $\overline{x} + \alpha d \in C$. In fact, it is sufficient to show that $\overline{x} + d \in C$, i.e., to assume that $\alpha = 1$, since the general case where $\alpha > 0$ can be reduced to the case where $\alpha = 1$ by replacing d with αd .

Let

$$z_k = x + kd,$$
 $k = 1, 2, \dots$

and note that $z_k \in C$ for all k, by our choice of x and d. If $\overline{x} = z_k$ for some k, then $\overline{x} + d = x + (k+1)d$, which belongs to C and we are done. We thus assume that $\overline{x} \neq z_k$ for all k, and we define

$$d_k = \frac{z_k - \overline{x}}{\|z_k - \overline{x}\|} \|d\|, \qquad k = 1, 2, \dots$$
 (1.17)

so that $\overline{x} + d_k$ is the intersection of the surface of the sphere centered at \overline{x} of radius ||d||, and the halfline that starts at \overline{x} and passes through z_k (see the construction of Fig. 1.4.2). We will now argue that $d_k \to d$, and that for large enough k, $\overline{x} + d_k \in C$, so using the closure of C, it follows that $\overline{x} + d \in C$.

Indeed, using the definition (1.17) of d_k , we have

$$\frac{d_k}{\|d\|} = \frac{\|z_k - x\|}{\|z_k - \overline{x}\|} \cdot \frac{z_k - x}{\|z_k - x\|} + \frac{x - \overline{x}}{\|z_k - \overline{x}\|} = \frac{\|z_k - x\|}{\|z_k - \overline{x}\|} \cdot \frac{d}{\|d\|} + \frac{x - \overline{x}}{\|z_k - \overline{x}\|}.$$

Because $\{z_k\}$ is an unbounded sequence,

$$\frac{\|z_k - x\|}{\|z_k - \overline{x}\|} \to 1, \qquad \frac{x - \overline{x}}{\|z_k - \overline{x}\|} \to 0,$$

so by combining the preceding relations, we have $d_k \to d$. The vector $\overline{x} + d_k$ lies between \overline{x} and z_k in the line segment connecting \overline{x} and z_k for all k such that $||z_k - \overline{x}|| \ge ||d||$, so by the convexity of C, we have $\overline{x} + d_k \in C$ for all sufficiently large k. Since $\overline{x} + d_k \to \overline{x} + d$ and C is closed, it follows that $\overline{x} + d \in C$. Q.E.D.

It is essential to assume that the set C is closed in the preceding proposition. For an example where part (a) fails without this assumption, consider the set

$$C = \{(x_1, x_2) \mid 0 < x_1, 0 < x_2\} \cup \{(0, 0)\}.$$

Its recession cone is equal to C, which is not closed. Part (b) also fails in this example, since for the direction d = (1,0) we have $x + \alpha d \in C$ for all $\alpha \geq 0$ and all $x \in C$, except for x = (0,0).

The following proposition gives some additional properties of recession cones.

Proposition 1.4.2: (Properties of Recession Cones) Let C be a nonempty closed convex set.

- (a) R_C contains a nonzero direction if and only if C is unbounded.
- (b) $R_C = R_{ri(C)}$.
- (c) For any collection of closed convex sets C_i , $i \in I$, where I is an arbitrary index set and $\bigcap_{i \in I} C_i \neq \emptyset$, we have

$$R_{\cap_{i\in I}C_i} = \cap_{i\in I}R_{C_i}.$$

(d) Let W be a compact and convex subset of \Re^m , and let A be an $m \times n$ matrix. The recession cone of the set

$$V = \{ x \in C \mid Ax \in W \}$$

(assuming this set is nonempty) is $R_C \cap N(A)$, where N(A) is the nullspace of A.

Proof: (a) Assuming that C is unbounded, we will show that R_C contains a nonzero direction (the reverse implication is clear). Choose any $x \in C$

and any unbounded sequence $\{z_k\} \subset C$. Consider the sequence $\{d_k\}$, where

$$d_k = \frac{z_k - x}{\|z_k - x\|},$$

and let d be a limit point of $\{d_k\}$ (compare with the construction of Fig. 1.4.2). Without loss of generality, assume that $||z_k - x||$ is monotonically increasing with k. For any fixed $\alpha \geq 0$, the vector $x + \alpha d_k$ lies between x and z_k in the line segment connecting x and z_k for all k such that $||z_k - x|| \geq \alpha$. Hence by the convexity of C, we have $x + \alpha d_k \in C$ for all sufficiently large k. Since $x + \alpha d$ is a limit point of $\{x + \alpha d_k\}$ and C is closed, we have $x + \alpha d \in C$. Hence, using also Prop. 1.4.1(b), it follows that the nonzero vector d is a direction of recession.

- (b) If $d \in R_{\mathrm{ri}(C)}$, then for a fixed $x \in \mathrm{ri}(C)$ and all $\alpha \geq 0$, we have $x + \alpha d \in \mathrm{ri}(C) \subset C$. Hence, by Prop. 1.4.1(b), we have $d \in R_C$. Conversely, if $d \in R_C$, then for any $x \in \mathrm{ri}(C)$, we have $x + \alpha d \in C$ for all $\alpha \geq 0$. It follows from the Line Segment Principle (Prop. 1.3.1) that $x + \alpha d \in \mathrm{ri}(C)$ for all $\alpha \geq 0$, so that d belongs to $R_{\mathrm{ri}(C)}$.
- (c) By the definition of direction of recession, $d \in R_{\cap_{i \in I} C_i}$ implies that $x + \alpha d \in \cap_{i \in I} C_i$ for all $x \in \cap_{i \in I} C_i$ and all $\alpha \geq 0$. By Prop. 1.4.1(b), this in turn implies that $d \in R_{C_i}$ for all i, so that $R_{\cap_{i \in I} C_i} \subset \cap_{i \in I} R_{C_i}$. Conversely, by the definition of direction of recession, if $d \in \cap_{i \in I} R_{C_i}$ and $x \in \cap_{i \in I} C_i$, we have $x + \alpha d \in \cap_{i \in I} C_i$ for all $\alpha \geq 0$, so $d \in R_{\cap_{i \in I} C_i}$. Thus, $\cap_{i \in I} R_{C_i} \subset R_{\cap_{i \in I} C_i}$.
- (d) Consider the closed convex set $\overline{V} = \{x \mid Ax \in W\}$, and choose some $x \in \overline{V}$. Then, by Prop. 1.4.1(b), $d \in R_{\overline{V}}$ if and only if $x + \alpha d \in \overline{V}$ for all $\alpha \geq 0$, or equivalently if and only if $A(x + \alpha d) \in W$ for all $\alpha \geq 0$. Since $Ax \in W$, the last statement is equivalent to $Ad \in R_W$. Thus, $d \in R_{\overline{V}}$ if and only if $Ad \in R_W$. Since W is compact, from part (a) we have $R_W = \{0\}$, so $R_{\overline{V}}$ is equal to $\{d \mid Ad = 0\}$, which is N(A). Since $V = C \cap \overline{V}$, using part (c), we have $R_V = R_C \cap N(A)$. Q.E.D.

For an example where part (a) of the preceding proposition fails, consider the unbounded convex set

$$C = \{(x_1, x_2) \mid 0 \le x_1 < 1, \ 0 \le x_2\} \cup \{(1, 0)\}.$$

By using the definition, it can be verified that C has no nonzero directions of recession. It can also be verified that (0,1) is a direction of recession of ri(C), so part (b) also fails. Finally, by letting

$$D = \{(x_1, x_2) \mid -1 \le x_1 \le 0, \ 0 \le x_2\},\$$

it can be seen that $(0,1) \in R_D$, so $R_{C \cap D} \neq R_C \cap R_D$ and part (c) fails as well.

Note that part (c) of the preceding proposition implies that if C and D are nonempty closed and convex sets such that $C \subset D$, then $R_C \subset R_D$. This can be seen by using part (c) to write $R_C = R_{C \cap D} = R_C \cap R_D$, from which we obtain $R_C \subset R_D$. This property can fail if the sets C and D are not closed; for example, if

$$C = \{(x_1, x_2) \mid 0 \le x_1 < 1, 0 \le x_2\}, \qquad D = C \cup \{(1, 0)\},\$$

then the vector (0,1) is a direction of recession of C but not of D.

Lineality Space

A subset of the recession cone of a convex set C that plays an important role in a number of interesting contexts is its *lineality space*, denoted by L_C . It is defined as the set of directions of recession d whose opposite, -d, are also directions of recession:

$$L_C = R_C \cap (-R_C).$$

Thus $d \in L_C$ if and only if the entire line $\{x + \alpha d \mid \alpha \in \Re\}$ is contained in C for every $x \in C$.

The lineality space inherits several of the properties of the recession cone that we have shown (Props. 1.4.1 and 1.4.2). We collect these properties in the following proposition.

Proposition 1.4.3: (Properties of Lineality Space) Let C be a nonempty closed convex subset of \Re^n .

- (a) L_C is a subspace of \Re^n .
- (b) $L_C = L_{ri(C)}$.
- (c) For any collection of closed convex sets C_i , $i \in I$, where I is an arbitrary index set and $\bigcap_{i \in I} C_i \neq \emptyset$, we have

$$L_{\cap_{i\in I}C_i} = \cap_{i\in I}L_{C_i}.$$

(d) Let W be a compact and convex subset of \Re^m , and let A be an $m \times n$ matrix. The lineality space of the set

$$V = \{ x \in C \mid Ax \in W \}$$

(assuming it is nonempty) is $L_C \cap N(A)$, where N(A) is the nullspace of A.

Proof: (a) Let d_1 and d_2 belong to L_C , and let α_1 and α_2 be nonzero scalars. We will show that $\alpha_1 d_1 + \alpha_2 d_2$ belongs to L_C . Indeed, we have

$$\alpha_1 d_1 + \alpha_2 d_2 = |\alpha_1| \left(\operatorname{sgn}(\alpha_1) d_1 \right) + |\alpha_2| \left(\operatorname{sgn}(\alpha_2) d_2 \right)$$

$$= \left(|\alpha_1| + |\alpha_2| \right) \left(\alpha \overline{d}_1 + (1 - \alpha) \overline{d}_2 \right), \tag{1.18}$$

where

$$\alpha = \frac{|\alpha_1|}{|\alpha_1| + |\alpha_2|}, \quad \overline{d}_1 = \operatorname{sgn}(\alpha_1)d_1, \quad \overline{d}_2 = \operatorname{sgn}(\alpha_2)d_2,$$

and for a nonzero scalar s, we use the notation sgn(s) = 1 or sgn(s) = -1depending on whether s is positive or negative, respectively. We now note that L_C is a convex cone, being the intersection of the convex cones R_C and $-R_C$. Hence, since \overline{d}_1 and \overline{d}_2 belong to L_C , any positive multiple of a convex combination of \overline{d}_1 and \overline{d}_2 belongs to L_C . It follows from Eq. (1.18) that $\alpha_1 d_1 + \alpha_2 d_2 \in L_C$.

(b) We have

$$L_{ri(C)} = R_{ri(C)} \cap \left(-R_{ri(C)}\right) = R_C \cap \left(-R_C\right) = L_C,$$

where the second equality follows from Prop. 1.4.2(b).

(c) We have

$$\begin{split} L_{\cap_{i \in I} C_i} &= \left(R_{\cap_{i \in I} C_i} \right) \cap \left(-R_{\cap_{i \in I} C_i} \right) \\ &= \left(\cap_{i \in I} R_{C_i} \right) \cap \left(- \cap_{i \in I} R_{C_i} \right) \\ &= \cap_{i \in I} \left(R_{C_i} \cap \left(-R_{C_i} \right) \right) \\ &= \cap_{i \in I} L_{C_i}, \end{split}$$

where the second equality follows from Prop. 1.4.2(c).

(d) We have

$$L_V = R_V \cap (-R_V)$$

$$= (R_C \cap N(A)) \cap ((-R_C) \cap N(A))$$

$$= (R_C \cap (-R_C)) \cap N(A)$$

$$= L_C \cap N(A),$$

where the second equality follows from Prop. 1.4.2(d). Q.E.D.

Example 1.4.1: (Sets Specified by Linear and Convex Quadratic Inequalities)

Consider a nonempty set of the form

$$C = \{x \mid x'Qx + c'x + b \le 0\},\$$

where Q is a symmetric positive semidefinite $n \times n$ matrix, c is a vector in \Re^n , and b is a scalar. A vector d is a direction of recession if and only if

$$(x + \alpha d)'Q(x + \alpha d) + c'(x + \alpha d) + b \le 0, \quad \forall \alpha > 0, \ x \in C,$$

or

$$x'Qx + c'x + b + \alpha(c + 2Qx)'d + \alpha^2 d'Qd \le 0, \quad \forall \ \alpha > 0, \ x \in C.$$
 (1.19)

Clearly, we cannot have d'Qd > 0, since then the left-hand side above would become arbitrarily large for a suitably large choice of α , so d'Qd = 0. Since Q is positive semidefinite, it can be written as Q = M'M for some matrix M, so that we have Md = 0, implying that Qd = 0. It follows that Eq. (1.19) is equivalent to

$$x'Qx + c'x + b + \alpha c'd \le 0, \quad \forall \alpha > 0, x \in C,$$

which is true if and only if $c'd \leq 0$. Thus,

$$R_C = \{d \mid Qd = 0, c'd \le 0\}.$$

Also, $L_C = R_C \cap (-R_C)$, so

$$L_C = \{d \mid Qd = 0, c'd = 0\}.$$

Consider now the case where C is nonempty and specified by any (possibly infinite) number of convex quadratic inequalities:

$$C = \{x \mid x'Q_{j}x + c'_{j}x + b_{j} \le 0, j \in J\},\$$

where J is some index set. Then using Props. 1.4.2(c) and 1.4.3(c), we have

$$R_C = \{d \mid Q_j d = 0, c'_j d \le 0, \forall j \in J\},\$$

$$L_C = \{d \mid Q_i d = 0, c'_i d = 0, \forall j \in J\}.$$

In particular, if C is a polyhedral set of the form

$$C = \{x \mid c_i'x + b_i \le 0, j = 1, \dots, r\},\$$

we have

$$R_C = \{d \mid c_i'd \le 0, j = 1, \dots, r\}, \qquad L_C = \{d \mid c_j'd = 0, j = 1, \dots, r\}.$$

Finally, let us prove a useful result that allows the decomposition of a convex set along a subspace of its lineality space (possibly the entire lineality space) and its orthogonal complement (see Fig. 1.4.3).

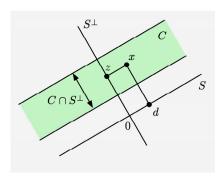


Figure 1.4.3. Illustration of the decomposition of a convex set C as

$$C = S + (C \cap S^{\perp}),$$

where S is a subspace contained in the lineality space L_C . A vector $x \in C$ is expressed as x = d + z with $d \in S$ and $z \in C \cap S^{\perp}$, as shown.

Proposition 1.4.4: (Decomposition of a Convex Set) Let C be a nonempty convex subset of \Re^n . Then, for every subspace S that is contained in the lineality space L_C , we have

$$C = S + (C \cap S^{\perp}).$$

Proof: We can decompose \Re^n as $S + S^{\perp}$, so for $x \in C$, let x = d + z for some $d \in S$ and $z \in S^{\perp}$. Because $-d \in S \subset L_C$, the vector -d is a direction of recession of C, so the vector x - d, which is equal to z, belongs to C, implying that $z \in C \cap S^{\perp}$. Thus, we have x = d + z with $d \in S$ and $z \in C \cap S^{\perp}$ showing that $C \subset S + (C \cap S^{\perp})$.

Conversely, if $x \in S + (C \cap S^{\perp})$, then x = d + z with $d \in S$ and $z \in C \cap S^{\perp}$. Thus, we have $z \in C$. Furthermore, because $S \subset L_C$, the vector d is a direction of recession of C, implying that $d + z \in C$. Hence $x \in C$, showing that $S + (C \cap S^{\perp}) \subset C$. Q.E.D.

In the special case where $S = L_C$ in Prop. 1.4.4, we obtain

$$C = L_C + (C \cap L_C^{\perp}). \tag{1.20}$$

Thus, C is the vector sum of two sets:

- (1) The set L_C , which consists of the lines contained in C, translated to pass through the origin.
- (2) The set $C \cap L_C^{\perp}$, which contains no lines; to see this, note that for any line $\{x + \alpha d \mid \alpha \in \Re\} \subset C \cap L_C^{\perp}$, we have $d \in L_C$ (since $x + \alpha d \in C$ for all $\alpha \in \Re$), so $d \perp (x + \alpha d)$ for all $\alpha \in \Re$, implying that d = 0.

Note that if $R_C = L_C$ and C is closed, the set $C \cap L_C^{\perp}$ contains no nonzero directions of recession, so it is compact [cf. Prop. 1.4.2(a)], and C can be decomposed into the sum of L_C and a compact set, as per Eq. (1.20).

1.4.1 Directions of Recession of a Convex Function

We will now develop a notion of direction of recession of a convex function. This notion is important in several contexts, including the existence of solutions of convex optimization problems, which will be discussed in Chapter 3. A key fact is that a convex function f can be described in terms of its epigraph, which is a convex set. The recession cone of $\operatorname{epi}(f)$ can be used to obtain the directions along which f does not increase monotonically. In particular, the "horizontal directions" in the recession cone of $\operatorname{epi}(f)$ correspond to the directions along which the level sets $\{x \mid f(x) \leq \gamma\}$ are unbounded. Along these directions, f is monotonically nonincreasing. This is the idea underlying the following proposition.

Proposition 1.4.5: Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a closed proper convex function and consider the level sets

$$V_{\gamma} = \{x \mid f(x) \le \gamma\}, \qquad \gamma \in \Re.$$

Then:

(a) All the nonempty level sets V_{γ} have the same recession cone, denoted $R_f,$ and given by

$$R_f = \{d \mid (d, 0) \in R_{epi(f)}\},\$$

where $R_{\text{epi}(f)}$ is the recession cone of the epigraph of f.

(b) If one nonempty level set V_{γ} is compact, then all of these level sets are compact.

Proof: (a) Fix a γ such that V_{γ} is nonempty. Let S be the " γ -slice" of epi(f),

$$S = \{(x, \gamma) \mid f(x) \le \gamma\},\$$

and note that

$$S = \operatorname{epi}(f) \cap \{(x, \gamma) \mid x \in \Re^n\}.$$

Using Prop. 1.4.2(c) [which applies since epi(f) is closed in view of the closedness of f], we have

$$R_S = R_{\text{epi}(f)} \cap \{(d,0) \mid d \in \Re^n\} = \{(d,0) \mid (d,0) \in R_{\text{epi}(f)}\}.$$

From this equation and the fact $S = \{(x, \gamma) \mid x \in V_{\gamma}\}$, the desired formula for $R_{V_{\gamma}}$ follows.

(b) From Prop. 1.4.2(a), a nonempty level set V_{γ} is compact if and only if the recession cone $R_{V_{\gamma}}$ does not contain a nonzero direction. By part (a), all nonempty level sets V_{γ} have the same recession cone, so if one of them is compact, all of them are compact. **Q.E.D.**

Note that closedness of f is essential for the level sets V_{γ} to have a common recession cone, as per Prop. 1.4.5(a). The reader may verify this by using as an example the convex but not closed function $f: \mathbb{R}^2 \mapsto (-\infty, \infty]$ given by

$$f(x_1, x_2) = \begin{cases} -x_1 & \text{if } x_1 > 0, \ x_2 \ge 0, \\ x_2 & \text{if } x_1 = 0, \ x_2 \ge 0, \\ \infty & \text{if } x_1 < 0 \text{ or } x_2 < 0. \end{cases}$$

Here, for $\gamma < 0$, we have $V_{\gamma} = \{(x_1, x_2) \mid x_1 \geq -\gamma, x_2 \geq 0\}$, so that $(0, 1) \in R_{V_{\gamma}}$, but $V_0 = \{(x_1, x_2) \mid x_1 > 0, x_2 \geq 0\} \cup \{(0, 0)\}$, so that $(0, 1) \notin R_{V_0}$.

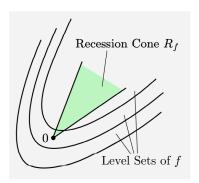


Figure 1.4.4. Illustration of the recession cone R_f of a closed proper convex function f. It is the (common) recession cone of the nonempty level sets of f.

For a closed proper convex function $f: \Re^n \mapsto (-\infty, \infty]$, the (common) recession cone R_f of the nonempty level sets is called the *recession* cone of f (cf. Fig. 1.4.4). A vector $d \in R_f$ is called a direction of recession of f.

The most intuitive way to look at directions of recession of f is from a descent viewpoint: if we start at any $x \in \text{dom}(f)$ and move indefinitely along a direction of recession, we must stay within each level set that contains x, or equivalently we must encounter exclusively points z with $f(z) \leq f(x)$. In words, a direction of recession of f is a direction of continuous nonascent for f. Conversely, if we start at some $x \in \text{dom}(f)$ and while moving along a direction d, we encounter a point z with f(z) > f(x), then d cannot be a direction of recession. By the convexity of the level sets of f, once we cross the relative boundary of a level set, we never cross it back again, and with a little thought, it can be seen that a direction that is not a direction of recession of f is a direction of eventual continuous ascent of f [see Figs. 1.4.5(e),(f)].

Constancy Space of a Convex Function

The lineality space of the recession cone R_f of a closed proper convex function f is denoted by L_f , and is the subspace of all $d \in \mathbb{R}^n$ such that both d and -d are directions of recession of f, i.e.,

$$L_f = R_f \cap (-R_f).$$

Equivalently, $d \in L_f$ if and only if both d and -d are directions of recession of each of the nonempty level sets $\{x \mid f(x) \leq \gamma\}$ [cf. Prop. 1.4.5(a)]. In view of the convexity of f, which implies that f is monotonically non-increasing along a direction of recession, we see that $d \in L_f$ if and only if

$$f(x + \alpha d) = f(x), \quad \forall x \in \text{dom}(f), \ \forall \ \alpha \in \Re.$$

Consequently, any $d \in L_f$ is called a direction in which f is constant, and L_f is called the constancy space of f.

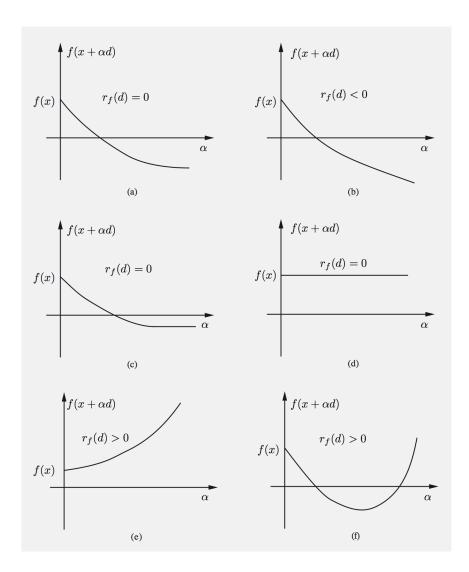


Figure 1.4.5. Ascent/descent behavior of a closed proper convex function starting at some $x \in \text{dom}(f)$ and moving along a direction d. If d is a direction of recession of f, there are two possibilities: either f decreases monotonically to a finite value or $-\infty$ [figures (a) and (b), respectively], or f reaches a value that is less or equal to f(x) and stays at that value [figures (c) and (d)]. If d is not a direction of recession of f, then eventually f increases monotonically to ∞ [figures (e) and (f)], i.e., for some $\overline{\alpha} \geq 0$ and all $\alpha_1, \alpha_2 \geq \overline{\alpha}$ with $\alpha_1 < \alpha_2$, we have

$$f(x + \alpha_1 d) < f(x + \alpha_2 d).$$

This behavior is determined only by d, and is independent of the choice of x within $\operatorname{dom}(f)$.

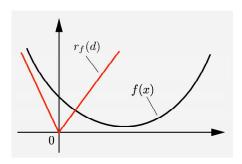


Figure 1.4.6. Illustration of the recession function r_f of a closed proper convex function f. Its epigraph is the recession cone of the epigraph of f.

As an example, if f is a quadratic function given by

$$f(x) = x'Qx + c'x + b,$$

where Q is a symmetric positive semidefinite $n \times n$ matrix, c is a vector in \Re^n , and b is a scalar, then its recession cone and constancy space are

$$R_f = \{d \mid Qd = 0, \ c'd \le 0\}, \qquad L_f = \{d \mid Qd = 0, \ c'd = 0\}$$

(cf. Example 1.4.1).

Recession Function

We saw that if d is a direction of recession of f, then f is asymptotically nonincreasing along each halfline $x + \alpha d$, but in fact a stronger property holds: it turns out that the asymptotic slope of f along d is independent of the starting point x. The "asymptotic slope" of a closed proper convex function along a direction is expressed by a function that we now introduce.

We first note that the recession cone $R_{\mathrm{epi}(f)}$ of the epigraph of a closed proper convex function $f: \Re^n \mapsto (-\infty, \infty]$ is itself the epigraph of another closed proper convex function. The reason is that for a given d, the set of scalars w such that $(d,w) \in R_{\mathrm{epi}(f)}$ is either empty or it is a closed interval that is unbounded above and is bounded below (since f is proper and hence its epigraph does not contain a vertical line). Thus $R_{\mathrm{epi}(f)}$ is the epigraph of a proper function, which must be closed and convex [since f, $\mathrm{epi}(f)$, and $R_{\mathrm{epi}(f)}$ are all closed and convex]. This function is called the recession function of f and is denoted r_f , i.e.,

$$\operatorname{epi}(r_f) = R_{\operatorname{epi}(f)};$$

see Fig. 1.4.6.

The recession function can be used to characterize the recession cone and constancy space of the function, as in the following proposition (cf. Fig. 1.4.5).

Proposition 1.4.6: Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a closed proper convex function. Then the recession cone and constancy space of f are given in terms of its recession function by

$$R_f = \{d \mid r_f(d) \le 0\}, \qquad L_f = \{d \mid r_f(d) = r_f(-d) = 0\}.$$

Proof: By Prop. 1.4.5(a) and the definition $epi(r_f) = R_{epi(f)}$ of r_f , the recession cone of f is

$$R_f = \{d \mid (d, 0) \in R_{\text{epi}(f)}\} = \{d \mid r_f(d) \le 0\}.$$

Since $L_f = R_f \cap (-R_f)$, it follows that $d \in L_f$ if and only if $r_f(d) \le 0$ and $r_f(-d) \le 0$. On the other hand, by the convexity of r_f , we have

$$r_f(d) + r_f(-d) \ge 2r_f(0) = 0, \quad \forall \ d \in \Re^n,$$

so it follows that $d \in L_f$ if and only if $r_f(d) = r_f(-d) = 0$. **Q.E.D.**

The following proposition provides an explicit formula for the recession function.

Proposition 1.4.7: Let $f: \Re^n \mapsto (-\infty, \infty]$ be a closed proper convex function. Then, for all $x \in \text{dom}(f)$ and $d \in \Re^n$,

$$r_f(d) = \sup_{\alpha > 0} \frac{f(x + \alpha d) - f(x)}{\alpha} = \lim_{\alpha \to \infty} \frac{f(x + \alpha d) - f(x)}{\alpha}.$$
 (1.21)

Proof: By definition, we have that $(d, \nu) \in R_{\text{epi}(f)}$ if and only if for all $(x, w) \in \text{epi}(f)$,

$$(x + \alpha d, w + \alpha \nu) \in \operatorname{epi}(f), \quad \forall \alpha > 0,$$

or equivalently, $f(x + \alpha d) \le f(x) + \alpha \nu$ for all $\alpha > 0$, which can be written as

$$\frac{f(x+\alpha d)-f(x)}{\alpha} \le \nu, \qquad \forall \ \alpha > 0.$$

Hence

$$(d, \nu) \in R_{\operatorname{epi}(f)}$$
 if and only if $\sup_{\alpha > 0} \frac{f(x + \alpha d) - f(x)}{\alpha} \le \nu$,

for all $x \in \text{dom}(f)$. Since $R_{\text{epi}(f)}$ is the epigraph of r_f , this implies the first equality in Eq. (1.21).

From the convexity of f, we see that the ratio

$$\frac{f(x+\alpha d)-f(x)}{\alpha}$$

is monotonically nondecreasing as a function of α over the range $(0, \infty)$. This implies the second equality in Eq. (1.21). Q.E.D.

The last expression in Eq. (1.21) leads to the interpretation of $r_f(d)$ as the "asymptotic slope" of f along the direction d. In fact, for differentiable convex functions $f: \Re^n \mapsto \Re$, this interpretation can be made more precise: we have

$$r_f(d) = \lim_{\alpha \to \infty} \nabla f(x + \alpha d)' d, \qquad \forall \ x \in \Re^n, \ d \in \Re^n.$$
 (1.22)

Indeed, for all x, d, and $\alpha > 0$, we have using Prop. 1.1.7(a),

$$\nabla f(x)'d \le \frac{f(x+\alpha d)-f(x)}{\alpha} \le \nabla f(x+\alpha d)'d,$$

so by taking the limit as $\alpha \to \infty$ and using Eq. (1.21), it follows that

$$\nabla f(x)'d \le r_f(d) \le \lim_{\alpha \to \infty} \nabla f(x + \alpha d)'d. \tag{1.23}$$

The left-hand side above holds for all x, so replacing x with $x + \alpha d$,

$$\nabla f(x + \alpha d)'d \le r_f(d), \quad \forall \alpha > 0.$$

By taking the limit as $\alpha \to \infty$, we obtain

$$\lim_{\alpha \to \infty} \nabla f(x + \alpha d)' d \le r_f(d), \tag{1.24}$$

and by combining Eqs. (1.23) and (1.24), we obtain Eq. (1.22).

The calculation of recession functions can be facilitated by nice formulas for the sum and the supremum of closed convex functions. The following proposition deals with the case of a sum.

Proposition 1.4.8: (Recession Function of a Sum) Let f_i : $\Re^n \mapsto (-\infty, \infty], i = 1, \dots, m$, be closed proper convex functions such that the function $f = f_1 + \cdots + f_m$ is proper. Then

$$r_f(d) = r_{f_1}(d) + \dots + r_{f_m}(d), \quad \forall \ d \in \Re^n.$$
 (1.25)

Proof: Without loss of generality, assume that m=2, and note that f_1+f_2 is closed proper convex (cf. Prop. 1.1.5). By using Eq. (1.21), we have for all $x \in \text{dom}(f_1+f_2)$ and $d \in \Re^n$,

$$r_{f_1+f_2}(d) = \lim_{\alpha \to \infty} \left\{ \frac{f_1(x+\alpha d) - f_1(x)}{\alpha} + \frac{f_2(x+\alpha d) - f_2(x)}{\alpha} \right\}$$
$$= \lim_{\alpha \to \infty} \left\{ \frac{f_1(x+\alpha d) - f_1(x)}{\alpha} \right\} + \lim_{\alpha \to \infty} \left\{ \frac{f_2(x+\alpha d) - f_2(x)}{\alpha} \right\}$$
$$= r_{f_1}(d) + r_{f_2}(d),$$

where the second equality holds because the limits involved exist. Q.E.D.

Note that for the formula (1.25) to hold, it is essential that f is proper, for otherwise its recession function is undefined. There is a similar result regarding the function

$$f(x) = \sup_{i \in I} f_i(x),$$

where I is an arbitrary index set, and $f_i: \Re^n \mapsto (-\infty, \infty], i \in I$, are closed proper convex functions such that f is proper. In particular, we have

$$r_f(d) = \sup_{i \in I} r_{f_i}(d), \qquad d \in \Re^n.$$
(1.26)

To show this, we simply note that the epigraph of r_f is the recession cone of the epigraph of f, the intersection of the epigraphs of f_i . Thus, the epigraph of r_f is the intersection of the recession cones of the epigraphs of f_i by Prop. 1.4.2(c), which yields the formula (1.26).

1.4.2 Nonemptiness of Intersections of Closed Sets

The notions of recession cone and lineality space can be used to generalize some of the fundamental properties of compact sets to closed convex sets. One such property is that a sequence $\{C_k\}$ of nonempty and compact sets with $C_{k+1} \subset C_k$ for all k has nonempty and compact intersection [cf. Prop. A.2.4(h)]. Another property is that the image of a compact set under a linear transformation is compact [cf. Prop. A.2.6(d)]. These properties may not hold when the sets involved are closed but unbounded (cf. Fig. 1.3.4), and some additional conditions are needed for their validity. In this section we develop such conditions, using directions of recession and related notions. We focus on the case where the sets involved are convex, but the analysis generalizes to the nonconvex case (see [BeT07]).

To understand the significance of set intersection results, consider a sequence of nonempty closed sets $\{C_k\}$ in \Re^n with $C_{k+1} \subset C_k$ for all k (such a sequence is said to be *nested*), and the question whether $\bigcap_{k=0}^{\infty} C_k$ is nonempty. Here are some of the contexts where this question arises:

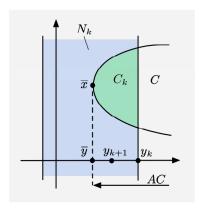


Figure 1.4.7. Set intersection argument to prove that the set AC closed when C is closed. Here A is the projection on the horizontal axis of points in the plane. For a sequence $\{y_k\} \subset AC$ that converges to some \overline{y} , in order to prove that $\overline{y} \in AC$, it is sufficient to prove that the intersection $\cap_{k=0}^{\infty} C_k$ is nonempty, where

$$C_k = C \cap N_k,$$

and

$$N_k = \{x \mid ||Ax - \overline{y}|| \le ||y_k - \overline{y}||\}.$$

(a) Does a function $f: \Re^n \mapsto \Re$ attain a minimum over a set X? This is true if and only if the intersection

$$\bigcap_{k=0}^{\infty} \left\{ x \in X \mid f(x) \le \gamma_k \right\}$$

is nonempty, where $\{\gamma_k\}$ is a scalar sequence with $\gamma_k \downarrow \inf_{x \in X} f(x)$.

(b) If C is a closed set and A is a matrix, is AC closed? To prove this, we may let $\{y_k\}$ be a sequence in AC that converges to some $\overline{y} \in \Re^n$, and then show that $\overline{y} \in AC$. If we introduce the sets $C_k = C \cap N_k$, where

$$N_k = \{x \mid ||Ax - \overline{y}|| \le ||y_k - \overline{y}||\},\$$

it is sufficient to show that $\bigcap_{k=0}^{\infty} C_k$ is nonempty (see Fig. 1.4.7).

We will next consider a nested sequence $\{C_k\}$ of nonempty closed convex sets, and in the subsequent propositions, we will derive several alternative conditions under which the intersection $\bigcap_{k=0}^{\infty} C_k$ is nonempty. These conditions involve a variety of assumptions about the recession cones, the lineality spaces, and the structure of the sets C_k .

Asymptotic Sequences of Convex Sets

The following line of analysis actually extends to nonconvex closed sets (see [BeT07]). However, in this book we will restrict ourselves to set intersections involving only convex sets.

Our analysis revolves around sequences $\{x_k\}$ such that $x_k \in C_k$ for each k. An important fact is that $\bigcap_{k=0}^{\infty} C_k$ is empty if and only if every sequence of this type is unbounded. Thus the idea is to introduce assumptions that guarantee that not all such sequences are unbounded. In fact it will be sufficient to restrict attention to unbounded sequences that escape to ∞ along common directions of recession of the sets C_k , as in the following definition.

Definition 1.4.1: Let $\{C_k\}$ be a nested sequence of nonempty closed convex sets. We say that $\{x_k\}$ is an asymptotic sequence of $\{C_k\}$ if $x_k \neq 0$, $x_k \in C_k$ for all k, and

$$||x_k|| \to \infty, \qquad \frac{x_k}{||x_k||} \to \frac{d}{||d||},$$

where d is some nonzero common direction of recession of the sets C_k ,

$$d \neq 0, \qquad d \in \cap_{k=0}^{\infty} R_{C_k}.$$

A special case is when all the sets C_k are equal. In particular, for a nonempty closed convex C, we say that $\{x_k\} \subset C$ is an asymptotic sequence of C if $\{x_k\}$ is asymptotic for the sequence $\{C_k\}$, where $C_k \equiv C$.

Note that given any unbounded sequence $\{x_k\}$ such that $x_k \in C_k$ for each k, there exists a subsequence $\{x_k\}_{k\in\mathcal{K}}$ that is asymptotic for the corresponding subsequence $\{C_k\}_{k\in\mathcal{K}}$. In fact, any limit point of $\{x_k/\|x_k\|\}$ is a common direction of recession of the sets C_k ; this can be seen by using the proof argument of Prop. 1.4.1(b). Thus, asymptotic sequences are in a sense representative of unbounded sequences with $x_k \in C_k$ for each k.

We now introduce a special type of set sequences that have favorable properties for our purposes.

Definition 1.4.2: Let $\{C_k\}$ be a nested sequence of nonempty closed convex sets. We say that an asymptotic sequence $\{x_k\}$ is retractive if for the direction d corresponding to $\{x_k\}$ as per Definition 1.4.1, there exists an index \overline{k} such that

$$x_k - d \in C_k, \quad \forall \ k \ge \overline{k}.$$

We say that the sequence $\{C_k\}$ is retractive if all its asymptotic sequences are retractive. In the special case $C_k \equiv C$, we say that the set C is retractive if all its asymptotic sequences are retractive.

Retractive set sequences are those whose asymptotic sequences still belong to the corresponding sets C_k (for sufficiently large k) when shifted by -d, where d is any corresponding direction of recession. For an example, consider a nested sequence consisting of "cylindrical" sets in the plane, such as $C_k = \{(x^1, x^2) \mid |x^1| \leq 1/k\}$, whose asymptotic sequences $\{(x^1_k, x^2_k)\}$ are retractive: they satisfy $x^1_k \to 0$, and either $x^2_k \to \infty$ [d = (0, 1)] or $x^2_k \to -\infty$ [d = (0, -1)] (see also Fig. 1.4.8). Some important types of

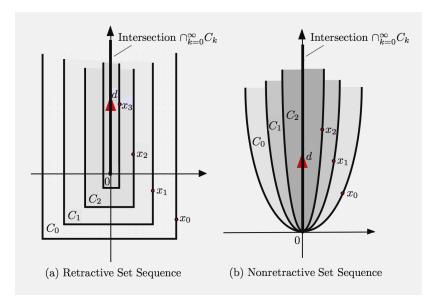


Figure 1.4.8. Illustration of retractive and nonretractive sequences in \Re^2 . For both set sequences, the intersection is the vertical half line $\{x \mid x_2 \geq 0\}$, and the common directions of recession are of the form $(0, d_2)$ with $d_2 \geq 0$. For the example on the right, any unbounded sequence $\{x_k\}$ such that x_k is on the boundary of the set C_k is asymptotic but not retractive.

set sequences can be shown to be retractive. As an aid in this regard, we note that intersections and Cartesian products (involving a finite number of sets) preserve retractiveness, as can be easily seen from the definition. In particular, if $\{C_k^1\},\ldots,\{C_k^r\}$ are retractive nested sequences of nonempty closed convex sets, the sequences $\{N_k\}$ and $\{T_k\}$ are retractive, where

$$N_k = C_k^1 \cap C_k^2 \cap \dots \cap C_k^r, \qquad T_k = C_k^1 \times C_k^2 \times \dots \times C_k^r, \qquad \forall \ k,$$

and we assume that all the sets N_k are nonempty.

The following proposition shows that a polyhedral set is retractive. Indeed, this is the most important type of retractive set for our purposes. Another retractive set of interest is the vector sum of a convex compact set and a polyhedral cone; we leave the proof of this for the reader. However, as a word of caution, we mention that a nonpolyhedral closed convex cone need not be retractive.

Proposition 1.4.9: A polyhedral set is retractive.

Proof: A closed halfspace is clearly retractive. A polyhedral set is the

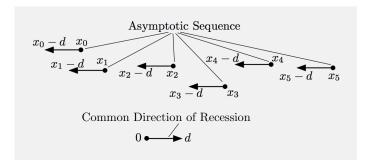


Figure 1.4.9. Geometric view of the proof idea of Prop. 1.4.10. An asymptotic sequence $\{x_k\}$ with corresponding direction of recession d eventually (for large k) gets closer to 0 when shifted by -d, so such a sequence cannot consist of the vectors of minimum norm from C_k without contradicting the retractiveness assumption.

nonempty intersection of a finite number of closed halfspaces, and set intersection preserves retractiveness. Q.E.D.

Set Intersection Theorems

The importance of retractive sequences is motivated by the following proposition.

Proposition 1.4.10: A retractive nested sequence of nonempty closed convex sets has nonempty intersection.

Proof: Let $\{C_k\}$ be the given sequence. For each k, let x_k be the vector of minimum norm in the closed set C_k (projection of the origin on C_k ; cf. Prop. 1.1.9). The proof involves two ideas:

- (a) The intersection $\bigcap_{k=0}^{\infty} C_k$ is empty if and only $\{x_k\}$ is unbounded, so there is a subsequence $\{x_k\}_{k\in\mathcal{K}}$ that is asymptotic.
- (b) If a subsequence $\{x_k\}_{k\in\mathcal{K}}$ of minimum norm vectors of C_k is asymptotic with corresponding direction of recession d, then $\{x_k\}_{k\in\mathcal{K}}$ cannot be retractive, because x_k would eventually (for large k) get closer to 0 when shifted by -d (see Fig. 1.4.9).

It will be sufficient to show that a subsequence $\{x_k\}_{k\in\mathcal{K}}$ is bounded. Then, since $\{C_k\}$ is nested, for each m, we have $x_k\in C_m$ for all $k\in\mathcal{K}$, $k\geq m$, and since C_m is closed, each of the limit points of $\{x_k\}_{k\in\mathcal{K}}$ will belong to each C_m and hence also to $\cap_{m=0}^{\infty} C_m$, thereby showing the result. Thus, we will prove the proposition by showing that there is no subsequence of $\{x_k\}$ that is unbounded.

Indeed, assume the contrary, let $\{x_k\}_{k\in\mathcal{K}}$ be a subsequence such that $\lim_{k\to\infty,\,k\in\mathcal{K}}\|x_k\|=\infty$, and let d be the limit of a subsequence $\{x_k/\|x_k\|\}_{k\in\overline{\mathcal{K}}}$, where $\overline{\mathcal{K}}\subset\mathcal{K}$. For each $k=0,1,\ldots$, define $z_k=x_m$, where m is the smallest index $m\in\overline{\mathcal{K}}$ with $k\le m$. Then, since $z_k\in C_k$ for all k and $\lim_{k\to\infty}\{z_k/\|z_k\|\}=d$, we see that d is a common direction of recession of C_k [cf. the proof of Prop. 1.4.1(b)] and $\{z_k\}$ is an asymptotic sequence corresponding to d. Using the retractiveness assumption, let \overline{k} be such that $z_k-d\in C_k$ for all $k\ge \overline{k}$. We have $d'z_k\to\infty$ since

$$\frac{d'z_k}{\|z_k\|} \to \|d\|^2 = 1,$$

so for all $k \geq \overline{k}$ with $2d'z_k > 1$, we obtain

$$||z_k - d||^2 = ||z_k||^2 - (2d'z_k - 1) < ||z_k||^2.$$

This is a contradiction, since for infinitely many k, z_k is the vector of minimum norm on C_k . Q.E.D.

For an example, consider the sequence $\{C_k\}$ of Fig. 1.4.8(a). Here the asymptotic sequences $\{(x_k^1, x_k^2)\}$ satisfy $x_k^1 \to 0$, $x_k^2 \to \infty$ and are retractive, and indeed the intersection $\cap_{k=0}^{\infty} C_k$ is nonempty. On the other hand, the condition for nonemptiness of $\cap_{k=0}^{\infty} C_k$ of the proposition is far from necessary, e.g., the sequence $\{C_k\}$ of Fig. 1.4.8(b) has nonempty intersection but is not retractive.

A simple example where the preceding proposition applies is a "cylindrical" set sequence, where $R_{C_k} \equiv L_{C_k} \equiv L$ for some subspace L. The following proposition gives an important extension.

Proposition 1.4.11: Let $\{C_k\}$ be a nested sequence of nonempty closed convex sets. Denote

$$R = \cap_{k=0}^{\infty} R_{C_k}, \qquad L = \cap_{k=0}^{\infty} L_{C_k}.$$

(a) If R = L, then $\{C_k\}$ is retractive, and $\bigcap_{k=0}^{\infty} C_k$ is nonempty. Furthermore,

$$\cap_{k=0}^{\infty} C_k = L + \tilde{C},$$

where \tilde{C} is some nonempty and compact set.

(b) Let X be a retractive closed convex set. Assume that all the sets $\overline{C}_k = X \cap C_k$ are nonempty, and that

$$R_X \cap R \subset L$$
.

Then, $\{\overline{C}_k\}$ is retractive, and $\bigcap_{k=0}^{\infty} \overline{C}_k$ is nonempty.

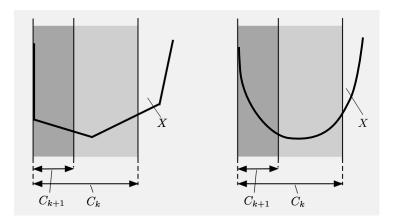


Figure 1.4.10. Illustration of the need to assume that X is retractive in Prop. 1.4.11(b). Here the intersection $\bigcap_{k=0}^{\infty} C_k$ is equal to the left vertical line. In the figure on the left, X is polyhedral, and the intersection $X \cap \left(\bigcap_{k=0}^{\infty} C_k\right)$ is nonempty. In the figure on the right, X is nonpolyhedral and nonretractive, and the intersection $X \cap \left(\bigcap_{k=0}^{\infty} C_k\right)$ is empty.

Proof: (a) The retractiveness of $\{C_k\}$ and consequent nonemptiness of $\bigcap_{k=0}^{\infty} C_k$ is the special case of part (b) where $X = \Re^n$. To show the given form of $\bigcap_{k=0}^{\infty} C_k$, we use the decomposition of Prop. 1.4.4, to obtain $\bigcap_{k=0}^{\infty} C_k = L + \tilde{C}$, where

$$\tilde{C} = (\cap_{k=0}^{\infty} C_k) \cap L^{\perp}.$$

The recession cone of \tilde{C} is $R \cap L^{\perp}$, and since R = L, it is equal to $\{0\}$. Hence by Prop. 1.4.2(a), \tilde{C} is compact.

(b) The common directions of recession of \overline{C}_k are those in $R_X \cap R$, so by the hypothesis they must belong to L. Thus, for any asymptotic sequence $\{x_k\}$ of $\{\overline{C}_k\}$, corresponding to $d \in R_X \cap R$, we have $d \in L$, and hence $x_k - d \in C_k$ for all k. Since X is retractive, we also have $x_k - d \in X$ and hence $x_k - d \in \overline{C}_k$, for sufficiently large k. Hence $\{x_k\}$ is retractive, so $\{\overline{C}_k\}$ is retractive, and by Prop. 1.4.10, $\bigcap_{k=0}^{\infty} \overline{C}_k$ is nonempty. Q.E.D.

Figure 1.4.10 illustrates the need to assume that X is retractive in Prop. 1.4.11(b). The following is an important application of the preceding set intersection result.

Proposition 1.4.12: (Existence of Solutions of Convex Quadratic Programs) Let Q be a symmetric positive semidefinite $n \times n$ matrix, let c and a_1, \ldots, a_r be vectors in \Re^n , and let b_1, \ldots, b_r be scalars. Assume that the optimal value of the problem

minimize
$$x'Qx + c'x$$

subject to $a'_j x \le b_j$, $j = 1, ..., r$,

is finite. Then the problem has at least one optimal solution.

Proof: Let f denote the cost function and let X be the polyhedral set of feasible solutions:

$$f(x) = x'Qx + c'x,$$
 $X = \{x \mid a'_j x \le b_j, \ j = 1, \dots, r\}.$

Let also f^* be the optimal value, let $\{\gamma_k\}$ be a scalar sequence with $\gamma_k \downarrow f^*$, and denote

$$\overline{C}_k = X \cap \{x \mid x'Qx + c'x \le \gamma_k\}.$$

We will use Prop. 1.4.11(b) to show that the set of optimal solutions, i.e., the intersection $\bigcap_{k=0}^{\infty} \overline{C}_k$, is nonempty. Indeed, let R_X be the recession cone of X, and let R and L be the common recession cone and lineality space of the sets $\{x \in \Re^n \mid x'Qx + c'x \leq \gamma_k\}$ (i.e., the recession cone and constancy space of f). By Example 1.4.1, we have

$$R = \{d \mid Qd = 0, c'd \le 0\},$$
 $L = \{d \mid Qd = 0, c'd = 0\},$ $R_X = \{d \mid a'_j d \le 0, j = 1, \dots, r\}.$

If d is such that $d \in R_X \cap R$ but $d \notin L$, then

$$Qd = 0, \qquad c'd < 0, \qquad a'_j d \le 0, \quad j = 1, \dots, r,$$

which implies that for any $x \in X$, we have $x + \alpha d \in X$ for all $\alpha \geq 0$, while $f(x + \alpha d) \to -\infty$ as $\alpha \to \infty$. This contradicts the finiteness of f^* , and shows that $R_X \cap R \subset L$. The nonemptiness of $\bigcap_{k=0}^{\infty} \overline{C}_k$ now follows from Prop. 1.4.11(b). **Q.E.D.**

1.4.3 Closedness Under Linear Transformations

The conditions just obtained regarding the nonemptiness of the intersection of a sequence of closed convex sets can be translated to conditions guaranteeing the closedness of the image, AC, of a closed convex set C under a linear transformation A. This is the subject of the following proposition.

Proposition 1.4.13: Let X and C be nonempty closed convex sets in \Re^n , and let A be an $m \times n$ matrix with nullspace denoted by N(A). If X is a retractive closed convex set and

$$R_X \cap R_C \cap N(A) \subset L_C$$

then $A(X \cap C)$ is a closed set.

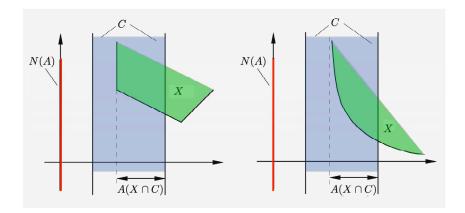


Figure 1.4.11. Illustration of the need to assume that the set X is retractive in Prop. 1.4.13. In both examples shown, the matrix A is the projection onto the horizontal axis, and its nullspace is the vertical axis. The condition $R_X \cap R_C \cap N(A) \subset L_C$ is satisfied. However, in the example on the right, X is not retractive, and the set $A(X \cap C)$ is not closed.

Proof: Let $\{y_k\}$ be a sequence in $A(X \cap C)$ converging to some \overline{y} . We will prove that $A(X \cap C)$ is closed by showing that $\overline{y} \in A(X \cap C)$. We introduce the sets

$$C_k = C \cap N_k$$

where

$$N_k = \left\{ x \mid ||Ax - \overline{y}|| \le ||y_k - \overline{y}|| \right\},\,$$

(see Fig. 1.4.7). The sets C_k are closed and convex, and their (common) recession cones and lineality spaces are $R_C \cap N(A)$ and $L_C \cap N(A)$, respectively [cf. Props. 1.4.2(d) and 1.4.3(d)]. Therefore, by Prop. 1.4.11(b), the intersection $X \cap \left(\bigcap_{k=0}^{\infty} C_k\right)$ is nonempty. Every point x in this intersection is such that $x \in X \cap C$ and $Ax = \overline{y}$, showing that $\overline{y} \in A(X \cap C)$. Q.E.D.

Figure 1.4.11 illustrates the need for the assumptions of Prop. 1.4.13. The proposition has some interesting special cases:

(a) Let $C = \Re^n$ and let X be a polyhedral set. Then, $L_C = \Re^n$ and the assumption of Prop. 1.4.13 is automatically satisfied, so it follows that AX is closed. Thus the image of a polyhedral set under a linear transformation is a closed set. Simple as this result may seem, it is especially important in optimization. For example, as a special case, it yields that the cone generated by vectors a_1, \ldots, a_r is a closed set, since it can be written as AC, where A is the matrix with columns a_1, \ldots, a_r and C is the polyhedral set of all $(\alpha_1, \ldots, \alpha_r)$ with $\alpha_j \geq 0$ for all j. This fact is central in the proof of Farkas' Lemma, an important result given in Section 2.3.1.

(b) Let $X = \Re^n$. Then, Prop. 1.4.13 yields that AC is closed if every direction of recession of C that belongs to N(A) belongs to the lineality space of C. This is true in particular if

$$R_C \cap N(A) = \{0\},$$

i.e., there is no nonzero direction of recession of C that lies in the nullspace of A. As a special case, this result can be used to obtain conditions that guarantee the closedness of the vector sum of closed convex sets. The idea is that the vector sum of a finite number of sets can be viewed as the image of their Cartesian product under a special type of linear transformation, as can be seen from the proof of the following proposition.

Proposition 1.4.14: Let C_1, \ldots, C_m be nonempty closed convex subsets of \Re^n such that the equality $d_1 + \cdots + d_m = 0$ for some vectors $d_i \in R_{C_i}$ implies that $d_i \in L_{C_i}$ for all $i = 1, \ldots, m$. Then $C_1 + \cdots + C_m$ is a closed set.

Proof: Let C be the Cartesian product $C_1 \times \cdots \times C_m$. Then, C is closed convex, and its recession cone and lineality space are given by

$$R_C = R_{C_1} \times \cdots \times R_{C_m}, \qquad L_C = L_{C_1} \times \cdots \times L_{C_m}.$$

Let A be the linear transformation that maps $(x_1, \ldots, x_m) \in \Re^{mn}$ into $x_1 + \cdots + x_m \in \Re^n$. The null space of A is the set of all (d_1, \ldots, d_m) such that $d_1 + \cdots + d_m = 0$. The intersection $R_C \cap N(A)$ consists of all (d_1, \ldots, d_m) such that $d_1 + \cdots + d_m = 0$ and $d_i \in R_{C_i}$ for all i. By the given condition, every $(d_1, \ldots, d_m) \in R_C \cap N(A)$ is such that $d_i \in L_{C_i}$ for all i, implying that $(d_1, \ldots, d_m) \in L_C$. Thus, $R_C \cap N(A) \subset L_C$, and by Prop. 1.4.13, the set AC is closed. Since

$$AC = C_1 + \dots + C_m,$$

the result follows. Q.E.D.

When specialized to just two sets, the above proposition implies that if C_1 and $-C_2$ are closed convex sets, then $C_1 - C_2$ is closed if there is no common nonzero direction of recession of C_1 and C_2 , i.e.

$$R_{C_1} \cap R_{C_2} = \{0\}.$$

This is true in particular if either C_1 or C_2 is bounded, in which case either $R_{C_1} = \{0\}$ or $R_{C_2} = \{0\}$, respectively.

Some other conditions asserting the closedness of vector sums can be derived from Prop. 1.4.13. For example, we can show that the vector sum of a finite number of polyhedral sets is closed, since it can be viewed as the image of their Cartesian product (clearly a polyhedral set) under a linear transformation (in fact this vector sum is polyhedral; see Section 2.3.2).

Another useful result is that if X is a polyhedral set, and C is a closed convex set, then X + C is closed if every direction of recession of X whose opposite is a direction of recession of C lies also in the lineality space of C (replace X and C by $X \times \mathbb{R}^n$ and $\mathbb{R}^n \times C$, respectively, in Prop. 1.4.13, and let A map Cartesian product to sum as in the proof of Prop. 1.4.14).

1.5 HYPERPLANES

Some of the most important principles in convexity and optimization, including duality, revolve around the use of hyperplanes, i.e., (n-1)-dimensional affine sets, which divide \Re^n into two halfspaces. For example, we will see that a closed convex set can be characterized in terms of hyperplanes: it is equal to the intersection of all the halfspaces that contain it. In the next section, we will apply this fundamental result to a convex function via its epigraph, and obtain an important dual description, encoded by another convex function, called the conjugate of the original.

A hyperplane in \Re^n is a set of the form

$$\{x \mid a'x = b\},\$$

where a is nonzero vector in \Re^n and b is a scalar. If \overline{x} is any vector in a hyperplane $H = \{x \mid a'x = b\}$, then we must have $a'\overline{x} = b$, so the hyperplane can be equivalently described as

$$H = \{x \mid a'x = a'\overline{x}\},\$$

or

$$H = \overline{x} + \{x \mid a'x = 0\}.$$

Thus, H is an affine set that is parallel to the subspace $\{x \mid a'x = 0\}$. The vector a is orthogonal to this subspace, and consequently, a is called the *normal* vector of H; see Fig. 1.5.1.

The sets

$$\{x \mid a'x \ge b\}, \qquad \{x \mid a'x \le b\},$$

are called the *closed halfspaces* associated with the hyperplane (also referred to as the *positive and negative halfspaces*, respectively). The sets

$$\{x \mid a'x > b\}, \qquad \{x \mid a'x < b\},$$

are called the *open halfspaces* associated with the hyperplane.

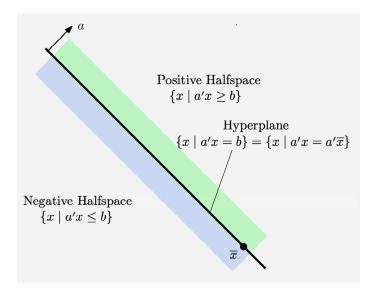


Figure 1.5.1. Illustration of the hyperplane $H = \{x \mid a'x = b\}$. If \overline{x} is any vector in the hyperplane, then the hyperplane can be equivalently described as

$$H = \{x \mid a'x = a'\overline{x}\} = \overline{x} + \{x \mid a'x = 0\}.$$

The hyperplane divides the space into two halfspaces as illustrated.

1.5.1 Hyperplane Separation

We say that two sets C_1 and C_2 are separated by a hyperplane $H = \{x \mid a'x = b\}$ if each lies in a different closed halfspace associated with H, i.e., if either

$$a'x_1 \le b \le a'x_2, \qquad \forall \ x_1 \in C_1, \ \forall \ x_2 \in C_2,$$

or

$$a'x_2 \le b \le a'x_1, \quad \forall x_1 \in C_1, \ \forall x_2 \in C_2.$$

We then also say that the hyperplane H separates C_1 and C_2 , or that H is a separating hyperplane of C_1 and C_2 . We use several different variants of this terminology. For example, the statement that two sets C_1 and C_2 can be separated by a hyperplane or that there exists a hyperplane separating C_1 and C_2 , means that there exists a vector $a \neq 0$ such that

$$\sup_{x \in C_1} a'x \le \inf_{x \in C_2} a'x;$$

[see Fig. 1.5.2(a)].

If a vector \overline{x} belongs to the closure of a set C, a hyperplane that separates C and the singleton set $\{\overline{x}\}$ is said to be supporting C at \overline{x} . Thus

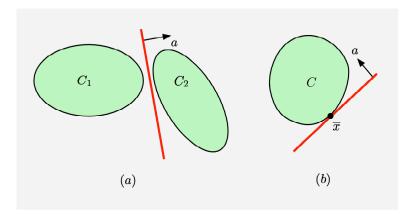


Figure 1.5.2. (a) Illustration of a hyperplane separating two sets C_1 and C_2 . (b) Illustration of a hyperplane supporting a set C at a point \overline{x} that belongs to the closure of C.

the statement that there exists a supporting hyperplane of C at \overline{x} means that there exists a vector $a \neq 0$ such that

$$a'\overline{x} \le a'x, \quad \forall \ x \in C,$$

or equivalently, since \overline{x} is a closure point of C,

$$a'\overline{x} = \inf_{x \in C} a'x.$$

As illustrated in Fig. 1.5.2(b), a supporting hyperplane of C is a hyperplane that "just touches" C.

We will prove several results regarding the existence of hyperplanes that separate two convex sets. Some of these results assert the existence of separating hyperplanes with special properties that will prove useful in various specialized contexts to be described later. The following proposition deals with the basic case where one of the two sets consists of a single vector. The proof is based on the Projection Theorem (Prop. 1.1.9) and is illustrated in Fig. 1.5.3.

Proposition 1.5.1: (Supporting Hyperplane Theorem) Let C be a nonempty convex subset of \mathbb{R}^n and let \overline{x} be a vector in \mathbb{R}^n . If \overline{x} is not an interior point of C, there exists a hyperplane that passes through \overline{x} and contains C in one of its closed halfspaces, i.e., there exists a vector $a \neq 0$ such that

$$a'\overline{x} \le a'x, \quad \forall \ x \in C.$$
 (1.27)

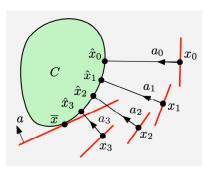


Figure 1.5.3. Illustration of the proof of the Supporting Hyperplane Theorem for the case where the vector \overline{x} belongs to $\operatorname{cl}(C)$, the closure of C. We choose a sequence $\{x_k\}$ of vectors that do not belong to $\operatorname{cl}(C)$, with $x_k \to \overline{x}$, and we project x_k on $\operatorname{cl}(C)$. We then consider, for each k, the hyperplane that is orthogonal to the line segment connecting x_k and its projection \hat{x}_k , and passes through x_k . These hyperplanes "converge" to a hyperplane that supports C at \overline{x} .

Proof: Consider cl(C), the closure of C, which is a convex set by Prop. 1.1.1(d). Let $\{x_k\}$ be a sequence of vectors such that $x_k \to \overline{x}$ and $x_k \notin cl(C)$ for all k; such a sequence exists because \overline{x} does not belong to the interior of C and hence does not belong to the interior of cl(C) [cf. Prop. 1.3.5(b)]. If \hat{x}_k is the projection of x_k on cl(C), we have by the optimality condition of the Projection Theorem (Prop. 1.1.9)

$$(\hat{x}_k - x_k)'(x - \hat{x}_k) \ge 0, \quad \forall \ x \in \operatorname{cl}(C).$$

Hence we obtain for all $x \in cl(C)$ and all k,

$$(\hat{x}_k - x_k)'x \ge (\hat{x}_k - x_k)'\hat{x}_k = (\hat{x}_k - x_k)'(\hat{x}_k - x_k) + (\hat{x}_k - x_k)'x_k \ge (\hat{x}_k - x_k)'x_k.$$

We can write this inequality as

$$a'_k x \ge a'_k x_k, \qquad \forall \ x \in \operatorname{cl}(C), \ \forall \ k,$$
 (1.28)

where

$$a_k = \frac{\hat{x}_k - x_k}{\|\hat{x}_k - x_k\|}.$$

We have $||a_k|| = 1$ for all k, so the sequence $\{a_k\}$ has a subsequence that converges to some $a \neq 0$. By considering Eq. (1.28) for all a_k belonging to this subsequence and by taking the limit as $k \to \infty$, we obtain Eq. (1.27). **Q.E.D.**

Note that if \overline{x} is a closure point of C, then the hyperplane of the preceding proposition supports C at \overline{x} . Note also that if C has empty interior, then any vector \overline{x} can be separated from C as in the proposition.

Proposition 1.5.2: (Separating Hyperplane Theorem) Let C_1 and C_2 be two nonempty convex subsets of \Re^n . If C_1 and C_2 are disjoint, there exists a hyperplane that separates them, i.e., there exists a vector $a \neq 0$ such that

$$a'x_1 \le a'x_2, \quad \forall \ x_1 \in C_1, \ \forall \ x_2 \in C_2.$$
 (1.29)

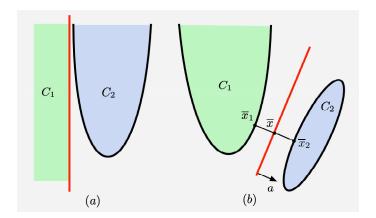


Figure 1.5.4. (a) An example of two disjoint convex sets that cannot be strictly separated. (b) Illustration of the construction of a strictly separating hyperplane.

Proof: Consider the convex set

$$C = C_2 - C_1 = \{x \mid x = x_2 - x_1, x_1 \in C_1, x_2 \in C_2\}.$$

Since C_1 and C_2 are disjoint, the origin does not belong to C, so by the Supporting Hyperplane Theorem (Prop. 1.5.1), there exists a vector $a \neq 0$ such that

$$0 \le a'x, \quad \forall \ x \in C,$$

which is equivalent to Eq. (1.29). Q.E.D.

We next consider a stronger form of separation of two sets C_1 and C_2 in \Re^n . We say that a hyperplane $\{x \mid a'x = b\}$ strictly separates C_1 and C_2 if it separates C_1 and C_2 while containing neither a point of C_1 nor a point of C_2 , i.e.,

$$a'x_1 < b < a'x_2, \quad \forall x_1 \in C_1, \ \forall x_2 \in C_2,$$

or

$$a'x_2 < b < a'x_1, \quad \forall x_1 \in C_1, \ \forall x_2 \in C_2.$$

Clearly, C_1 and C_2 must be disjoint in order that they can be strictly separated. However, this is not sufficient to guarantee strict separation (see Fig. 1.5.4). The following proposition provides conditions that guarantee the existence of a strictly separating hyperplane.

Proposition 1.5.3: (Strict Separation Theorem) Let C_1 and C_2 be two disjoint nonempty convex sets. There exists a hyperplane that strictly separates C_1 and C_2 under any one of the following five conditions:

- (1) $C_2 C_1$ is closed.
- (2) C_1 is closed and C_2 is compact.
- (3) C_1 and C_2 are polyhedral.
- (4) C_1 and C_2 are closed, and

$$R_{C_1} \cap R_{C_2} = L_{C_1} \cap L_{C_2},$$

where R_{C_i} and L_{C_i} denote the recession cone and the lineality space of C_i , i = 1, 2.

(5) C_1 is closed, C_2 is polyhedral, and $R_{C_1} \cap R_{C_2} \subset L_{C_1}$.

Proof: We will show the result under condition (1). The result will then follow under conditions (2)-(5), because these conditions imply condition (1) (see Prop. 1.4.14, and the discussion following its proof).

Assume that $C_2 - C_1$ is closed, and consider the vector of minimum norm (projection of the origin, cf. Prop. 1.1.9) in $C_2 - C_1$. This vector is of the form $\overline{x}_2 - \overline{x}_1$, where $\overline{x}_1 \in C_1$ and $\overline{x}_2 \in C_2$. Let

$$a = \frac{\overline{x}_2 - \overline{x}_1}{2}, \qquad \overline{x} = \frac{\overline{x}_1 + \overline{x}_2}{2}, \qquad b = a'\overline{x};$$

[cf. Fig. 1.5.4(b)]. Then, $a \neq 0$, since $\overline{x}_1 \in C_1$, $\overline{x}_2 \in C_2$, and C_1 and C_2 are disjoint. We will show that the hyperplane

$$\{x \mid a'x = b\}$$

strictly separates C_1 and C_2 , i.e., that

$$a'x_1 < b < a'x_2, \quad \forall x_1 \in C_1, \ \forall x_2 \in C_2.$$
 (1.30)

To this end we note that \overline{x}_1 is the projection of \overline{x}_2 on $\operatorname{cl}(C_1)$ (otherwise there would exist a vector $x_1 \in C_1$ with $\|\overline{x}_2 - x_1\| < \|\overline{x}_2 - \overline{x}_1\|$ - a contradiction of the minimum norm property of $\overline{x}_2 - \overline{x}_1$). Thus, we have

$$(\overline{x}_2 - \overline{x}_1)'(x_1 - \overline{x}_1) \le 0, \quad \forall x_1 \in C_1,$$

or equivalently, since $\overline{x} - \overline{x}_1 = a$,

$$a'x_1 \le a'\overline{x}_1 = a'\overline{x} + a'(\overline{x}_1 - \overline{x}) = b - ||a||^2 < b, \quad \forall x_1 \in C_1.$$

Thus, the left-hand side of Eq. (1.30) is proved. The right-hand side is proved similarly. **Q.E.D.**

Note that as a corollary of the preceding proposition, a closed set C can be strictly separated from a vector $\overline{x} \notin C$, i.e., from the singleton set $\{\overline{x}\}$. We will use this fact to provide the following important characterization of closed convex sets.

Proposition 1.5.4: The closure of the convex hull of a set C is the intersection of the closed halfspaces that contain C. In particular, a closed convex set is the intersection of the closed halfspaces that contain it.

Proof: Let H denote the intersection of all closed halfspaces that contain C. Since every closed halfspace containing C must also contain $\operatorname{cl}(\operatorname{conv}(C))$, it follows that $H \supset \operatorname{cl}(\operatorname{conv}(C))$.

To show the reverse inclusion, consider a vector $x \notin \operatorname{cl}(\operatorname{conv}(C))$ and a hyperplane strictly separating x and $\operatorname{cl}(\operatorname{conv}(C))$. The corresponding closed halfspace that contains $\operatorname{cl}(\operatorname{conv}(C))$ does not contain x, so $x \notin H$. Hence $H \subset \operatorname{cl}(\operatorname{conv}(C))$. Q.E.D.

1.5.2 Proper Hyperplane Separation

We now discuss a form of hyperplane separation, called *proper*, which turns out to be useful in some important optimization contexts, such as the duality theorems of Chapter 4 (Props. 4.4.1 and 4.5.1).

Let C_1 and C_2 be two subsets of \Re^n . We say that a hyperplane properly separates C_1 and C_2 if it separates C_1 and C_2 , and does not fully contain both C_1 and C_2 . Thus there exists a hyperplane that properly separates C_1 and C_2 if and only if there is a vector a such that

$$\sup_{x_1 \in C_1} a' x_1 \le \inf_{x_2 \in C_2} a' x_2, \qquad \inf_{x_1 \in C_1} a' x_1 < \sup_{x_2 \in C_2} a' x_2;$$

(see Fig. 1.5.5). If C is a subset of \Re^n and \overline{x} is a vector in \Re^n , we say that a hyperplane properly separates C and \overline{x} if it properly separates C and the singleton set $\{\overline{x}\}$.

Note that a convex set in \Re^n that has nonempty interior (and hence has dimension n) cannot be fully contained in a hyperplane (which has dimension n-1). Thus, in view of the Separating Hyperplane Theorem (Prop. 1.5.2), two disjoint convex sets one of which has nonempty interior can be properly separated. Similarly and more generally, two disjoint convex sets such that the affine hull of their union has dimension n can be properly separated. Figure 1.5.5(c) provides an example of two convex sets that cannot be properly separated.

The existence of a hyperplane that properly separates two convex sets is intimately tied to conditions involving the relative interiors of the sets.

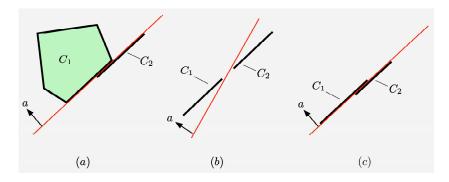


Figure 1.5.5. (a) and (b) Illustration of a properly separating hyperplanes. (c) Illustration of two convex sets that cannot be properly separated.

An important fact in this connection is that given a nonempty convex set C and a hyperplane H that contains C in one of its closed halfspaces, we have

$$C \subset H$$
 if and only if $ri(C) \cap H \neq \emptyset$. (1.31)

To see this, let H be of the form $\{x \mid a'x = b\}$ with $a'x \geq b$ for all $x \in C$. Then for a vector $\overline{x} \in \text{ri}(C)$, we have $\overline{x} \in H$ if and only if $a'\overline{x} = b$, i.e., a'x attains its minimum over C at \overline{x} . By Prop. 1.3.4, this is so if and only if a'x = b for all $x \in C$, i.e., $C \subset H$.

The following propositions provide relative interior assumptions that guarantee the existence of properly separating hyperplanes.

Proposition 1.5.5: (Proper Separation Theorem) Let C be a nonempty convex subset of \Re^n and let \overline{x} be a vector in \Re^n . There exists a hyperplane that properly separates C and \overline{x} if and only if $\overline{x} \notin ri(C)$.

Proof: Suppose that there exists a hyperplane H that properly separates C and \overline{x} . Then either $\overline{x} \notin H$, in which case $\overline{x} \notin C$ and $\overline{x} \notin ri(C)$, or else $\overline{x} \in H$ and C is not contained in H, in which case by Eq. (1.31), $ri(C) \cap H = \emptyset$, in which case again $\overline{x} \notin ri(C)$.

Conversely, assume that $\overline{x} \notin ri(C)$. To show the existence of a properly separating hyperplane, we consider two cases (see Fig. 1.5.6):

(a) $\overline{x} \notin \text{aff}(C)$. In this case, since aff(C) is closed and convex, by the Strict Separation Theorem [Prop. 1.5.3 under condition (2)] there exists a hyperplane that separates $\{\overline{x}\}$ and aff(C) strictly, and hence also properly separates C and \overline{x} .

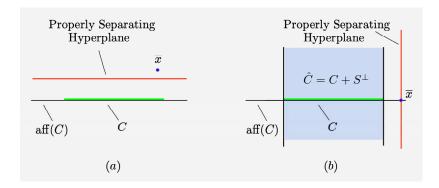


Figure 1.5.6. Illustration of the construction of a hyperplane that properly separates a convex set C and a point $\overline{x} \notin \operatorname{ri}(C)$ (cf. the proof of Prop. 1.5.5). In case (a), where $\overline{x} \notin \operatorname{aff}(C)$, the hyperplane is constructed as shown. In case (b), where $\overline{x} \in \operatorname{aff}(C)$, we consider the subspace S that is parallel to $\operatorname{aff}(C)$, we set $\hat{C} = C + S^{\perp}$, and we use the Supporting Hyperplane Theorem to separate \overline{x} from \hat{C} (Prop. 1.5.1).

(b) $\overline{x} \in \operatorname{aff}(C)$. In this case, let S be the subspace that is parallel to $\operatorname{aff}(C)$, and consider the set $\hat{C} = C + S^{\perp}$. From Prop. 1.3.7, we have $\operatorname{ri}(\hat{C}) = \operatorname{ri}(C) + S^{\perp}$, so that \overline{x} is not an interior point of \hat{C} [otherwise there must exist a vector $x \in \operatorname{ri}(C)$ such that $x - \overline{x} \in S^{\perp}$, which, since $x \in \operatorname{aff}(C)$, $\overline{x} \in \operatorname{aff}(C)$, and $x - \overline{x} \in S$, implies that $x - \overline{x} = 0$, thereby contradicting the hypothesis $\overline{x} \notin \operatorname{ri}(C)$]. By the Supporting Hyperplane Theorem (Prop. 1.5.1), it follows that there exists a vector $a \neq 0$ such that $a'x \geq a'\overline{x}$ for all $x \in \hat{C}$. Since \hat{C} has nonempty interior, a'x cannot be constant over \hat{C} , and

$$a'\overline{x} < \sup_{x \in \hat{C}} a'x = \sup_{x \in C, z \in S^{\perp}} a'(x+z) = \sup_{x \in C} a'x + \sup_{z \in S^{\perp}} a'z. \quad (1.32)$$

If we had $a'\overline{z} \neq 0$ for some $\overline{z} \in S^{\perp}$, we would also have

$$\inf_{\alpha \in \Re} a'(x + \alpha \overline{z}) = -\infty,$$

which contradicts the fact $a'(x+z) \ge a'\overline{x}$ for all $x \in C$ and $z \in S^{\perp}$. It follows that

$$a'z = 0, \quad \forall z \in S^{\perp},$$

which when combined with Eq. (1.32), yields

$$a'\overline{x} < \sup_{x \in C} a'x.$$

Thus the hyperplane $\{x \mid a'x = a'\overline{x}\}$ properly separates C and \overline{x} . Q.E.D.

Proposition 1.5.6: (Proper Separation of Two Convex Sets) Let C_1 and C_2 be two nonempty convex subsets of \Re^n . There exists a hyperplane that properly separates C_1 and C_2 if and only if

$$ri(C_1) \cap ri(C_2) = \emptyset.$$

Proof: Consider the convex set $C = C_2 - C_1$. By Prop. 1.3.7, we have

$$ri(C) = ri(C_2) - ri(C_1),$$

so the assumption $ri(C_1) \cap ri(C_2) = \emptyset$ is equivalent to $0 \notin ri(C)$. By using Prop. 1.5.5, it follows that there exists a hyperplane properly separating C and the origin, so we have

$$0 \le \inf_{x_1 \in C_1, x_2 \in C_2} a'(x_2 - x_1), \qquad 0 < \sup_{x_1 \in C_1, x_2 \in C_2} a'(x_2 - x_1),$$

if and only if $ri(C_1) \cap ri(C_2) = \emptyset$. This is equivalent to the desired assertion. **Q.E.D.**

The following proposition is a variant of Prop. 1.5.6. It shows that if C_2 is polyhedral and the slightly stronger condition $\operatorname{ri}(C_1) \cap C_2 = \emptyset$ holds, then there exists a properly separating hyperplane satisfying the extra restriction that it does not contain the nonpolyhedral set C_1 (rather than just the milder requirement that it does not contain either C_1 or C_2); see Fig. 1.5.7.

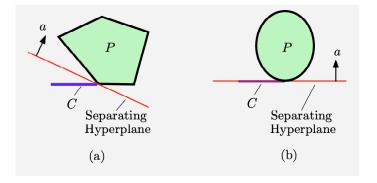


Figure 1.5.7. Illustration of the special proper separation property of a convex set C and a polyhedral set P, under the condition $\operatorname{ri}(C) \cap P = \emptyset$. In figure (a), the separating hyperplane can be chosen so that it does not contain C. If P is not polyhedral, as in figure (b), this may not be possible.

does not contain C if and only if

Proposition 1.5.7: (Polyhedral Proper Separation Theorem) Let C and P be two nonempty convex subsets of \Re^n such that P is polyhedral. There exists a hyperplane that separates C and P, and

$$ri(C) \cap P = \emptyset$$
.

Proof: First, as a general observation, we recall from our discussion of proper separation that given a convex set X and a hyperplane H that contains X in one of its closed halfspaces, we have

$$X \subset H$$
 if and only if $ri(X) \cap H \neq \emptyset$; (1.33)

cf. Eq. (1.31). We will use repeatedly this relation in the subsequent proof. Assume that there exists a hyperplane H that separates C and P, and does not contain C. Then, by Eq. (1.33), H cannot contain a point in $\mathrm{ri}(C)$, and since H separates C and P, we must have $\mathrm{ri}(C) \cap P = \emptyset$.

Conversely, assume that $ri(C) \cap P = \emptyset$. We will show that there exists a separating hyperplane that does not contain C. Denote

$$D = P \cap \operatorname{aff}(C)$$
.

If $D = \emptyset$, then since $\operatorname{aff}(C)$ and P are polyhedral, the Strict Separation Theorem [cf. Prop. 1.5.3 under condition (3)] applies and shows that there exists a hyperplane H that separates $\operatorname{aff}(C)$ and P strictly, and hence does not contain C.

We may thus assume that $D \neq \emptyset$. The idea now is to first construct a hyperplane that properly separates C and D, and then extend this hyperplane so that it suitably separates C and P. [If C had nonempty interior, the proof would be much simpler, since then $\operatorname{aff}(C) = \Re^n$ and D = P.]

By assumption, we have $ri(C) \cap P = \emptyset$ implying that

$$\operatorname{ri}(C) \cap \operatorname{ri}(D) \subset \operatorname{ri}(C) \cap (P \cap \operatorname{aff}(C)) = (\operatorname{ri}(C) \cap P) \cap \operatorname{aff}(C) = \emptyset.$$

Hence, by Prop. 1.5.6, there exists a hyperplane H that properly separates C and D. Furthermore, H does not contain C, since if it did, H would also contain $\mathrm{aff}(C)$ and hence also D, contradicting the proper separation property. Thus, C is contained in one of the closed halfspaces of H, but not in both. Let \overline{C} be the intersection of $\mathrm{aff}(C)$ and the closed halfspace of H that contains C; see Fig. 1.5.8. Note that H does not contain \overline{C} (since H does not contain C), and by Eq. (1.33), we have $H \cap \mathrm{ri}(\overline{C}) = \emptyset$, implying that

$$P\cap \mathrm{ri}(\overline{C})=\varnothing,$$

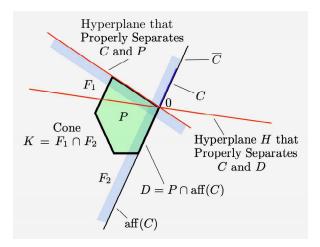


Figure 1.5.8. Illustration of the proof of Prop. 1.5.7 in the case where $D = P \cap \operatorname{aff}(C) \neq \emptyset$. The figure shows the construction of a hyperplane that properly separates C and P, and does not contain C, starting from the hyperplane H that properly separates C and D. In this two-dimensional example we have $M = \{0\}$, so $K = \operatorname{cone}(P) + M = \operatorname{cone}(P)$.

[if $\overline{x} \in P \cap ri(\overline{C})$ then $\overline{x} \in D \cap ri(\overline{C})$, a contradiction since D and $ri(\overline{C})$ lie in the opposite closed halfspaces of H and $H \cap ri(\overline{C}) = \emptyset$].

If $P \cap \overline{C} = \emptyset$, then by using again the Strict Separation Theorem [cf. Prop. 1.5.3 under condition (3)], we can construct a hyperplane that strictly separates P and \overline{C} . This hyperplane also strictly separates P and C, and we are done. We thus assume that $P \cap \overline{C} \neq \emptyset$, and by using a translation argument if necessary, we assume that

$$0 \in P \cap \overline{C}$$
,

as indicated in Fig. 1.5.8. The polyhedral set P can be represented as the intersection of halfspaces $\{x \mid a'_j x \leq b_j\}$ with $b_j \geq 0$ (since $0 \in P$) and with $b_j = 0$ for at least one j [since otherwise 0 would be in the interior of P; then, by the Line Segment Principle, for any $\overline{x} \in \text{ri}(\overline{C})$ the line segment connecting 0 and \overline{x} contains points in $\text{ri}(D) \cap \text{ri}(\overline{C})$, a contradiction of the fact that H properly separates D and \overline{C}]. Thus, we have

$$P = \{x \mid a'_{j}x \le 0, j = 1, \dots, m\} \cap \{x \mid a'_{j}x \le b_{j}, j = m + 1, \dots, \overline{m}\},\$$

for some integers $m \ge 1$ and $\overline{m} \ge m$, vectors a_j , and scalars $b_j > 0$. Let M be the relative boundary of \overline{C} , i.e.,

$$M = H \cap \operatorname{aff}(C),$$

and consider the cone

$$K = \{x \mid a'_i x \le 0, j = 1, \dots, m\} + M.$$

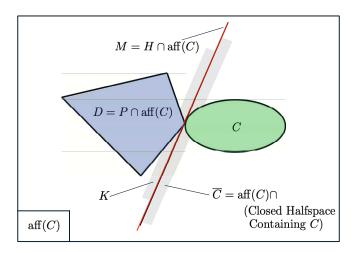


Figure 1.5.9. View of D and C within $\operatorname{aff}(C)$, and illustration of the construction of the cone $K = \operatorname{cone}(P) + M$ in the proof of Prop. 1.5.7.

Note that K = cone(P) + M (see Figs. 1.5.8 and 1.5.9).

We claim that $K \cap \operatorname{ri}(\overline{C}) = \emptyset$. The proof is by contradiction. If there exists $\overline{x} \in K \cap \operatorname{ri}(\overline{C})$, then \overline{x} can be expressed as $\overline{x} = \alpha w + v$ for some $\alpha > 0$, $w \in P$, and $v \in M$ [since $K = \operatorname{cone}(P) + M$ and $0 \in P$], so that $w = (\overline{x}/\alpha) - (v/\alpha) \in P$. On the other hand, since $\overline{x} \in \operatorname{ri}(\overline{C})$, $0 \in \overline{C} \cap M$, and M is a subset of the lineality space of \overline{C} [and hence also of the lineality space of $\operatorname{ri}(\overline{C})$], all vectors of the form $\overline{\alpha} \, \overline{x} + \overline{v}$, with $\overline{\alpha} > 0$ and $\overline{v} \in M$, belong to $\operatorname{ri}(\overline{C})$. In particular the vector $w = (\overline{x}/\alpha) - (v/\alpha)$ belongs to $\operatorname{ri}(\overline{C})$, so $w \in P \cap \operatorname{ri}(\overline{C})$. This is a contradiction since $P \cap \operatorname{ri}(\overline{C}) = \emptyset$, and it follows that $K \cap \operatorname{ri}(\overline{C}) = \emptyset$.

The cone K is polyhedral (since it is the vector sum of two polyhedral sets), so it is the intersection of some closed halfspaces F_1, \ldots, F_r that pass through 0 (cf. Fig. 1.5.8). Since $K = \operatorname{cone}(P) + M$, each of these closed halfspaces contains M, the relative boundary of the set \overline{C} , and furthermore \overline{C} is the closed half of a subspace. It follows that if any of the closed halfspaces F_1, \ldots, F_r contains a vector in $\operatorname{ri}(\overline{C})$, then that closed halfspace entirely contains \overline{C} . Hence, since K does not contain any point in $\operatorname{ri}(\overline{C})$, at least one of F_1, \ldots, F_r , say F_1 , does not contain any point in $\operatorname{ri}(\overline{C})$ (cf. Fig. 1.5.8). Therefore, the hyperplane corresponding to F_1 contains no points of $\operatorname{ri}(\overline{C})$, and hence also no points of $\operatorname{ri}(C)$. Thus, this hyperplane does not contain C, while separating K and C. Since K contains P, this hyperplane also separates P and C. Q.E.D.

Note that in the preceding proof, it is essential to introduce M, the relative boundary of the set \overline{C} , and to define K = cone(P) + M. If instead we define K = cone(P), then the corresponding halfspaces F_1, \ldots, F_r may all intersect $\text{ri}(\overline{C})$, and the proof argument fails (see Fig. 1.5.9).

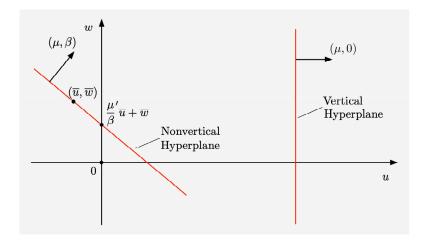


Figure 1.5.10. Illustration of vertical and nonvertical hyperplanes in \Re^{n+1} . A hyperplane with normal (μ, β) is nonvertical if $\beta \neq 0$, or, equivalently, if it intersects the (n+1)st axis at the unique point $\xi = (\mu/\beta)'\overline{u} + \overline{w}$, where $(\overline{u}, \overline{w})$ is any vector on the hyperplane.

1.5.3 Nonvertical Hyperplane Separation

In the context of optimization, supporting hyperplanes are often used in conjunction with epigraphs of functions defined on \Re^n . Since the epigraph is a subset of \Re^{n+1} , we consider hyperplanes in \Re^{n+1} and associate them with nonzero vectors of the form (μ, β) , where $\mu \in \Re^n$ and $\beta \in \Re$. We say that such a hyperplane is vertical if $\beta = 0$.

Note that if a hyperplane with normal (μ, β) is nonvertical, then it crosses the (n+1)st axis (the axis associated with w) at a unique point. In particular, if $(\overline{u}, \overline{w})$ is any vector on the hyperplane, the crossing point has the form $(0, \xi)$, where

$$\xi = \frac{\mu'}{\beta} \overline{u} + \overline{w},$$

since from the hyperplane equation, we have $(0,\xi)'(\mu,\beta) = (\overline{u},\overline{w})'(\mu,\beta)$. If the hyperplane is vertical, it either contains the entire (n+1)st axis, or else it does not cross it at all; see Fig. 1.5.10. Furthermore, a hyperplane H is vertical if and only if the recession cone of H, as well as the recession cones of the closed halfspaces associated with H, contain the (n+1)st axis.

Vertical lines in \Re^{n+1} are sets of the form $\{(\overline{u},w) \mid w \in \Re\}$, where \overline{u} is a fixed vector in \Re^n . If $f: \Re^n \mapsto (-\infty, \infty]$ is a proper convex function, then $\operatorname{epi}(f)$ cannot contain a vertical line, and it appears plausible that $\operatorname{epi}(f)$ is contained in a closed halfspace corresponding to some nonvertical hyperplane. We prove this fact in greater generality in the following proposition, which will be important for the development of duality.

Proposition 1.5.8: (Nonvertical Hyperplane Theorem) Let C be a nonempty convex subset of \Re^{n+1} that contains no vertical lines. Let the vectors in \Re^{n+1} be denoted by (u, w), where $u \in \Re^n$ and $w \in \Re$. Then:

(a) C is contained in a closed halfspace corresponding to a nonvertical hyperplane, i.e., there exist a vector $\mu \in \Re^n$, a scalar $\beta \neq 0$, and a scalar γ such that

$$\mu'u + \beta w \ge \gamma, \quad \forall (u, w) \in C.$$

(b) If $(\overline{u}, \overline{w})$ does not belong to cl(C), there exists a nonvertical hyperplane strictly separating $(\overline{u}, \overline{w})$ and C.

Proof: (a) Assume, to arrive at a contradiction, that every hyperplane containing C in one of its closed halfspaces is vertical. Then every hyperplane containing $\operatorname{cl}(C)$ in one of its closed halfspaces must also be vertical and its recession cone must contain the (n+1)st axis. By Prop. 1.5.4, $\operatorname{cl}(C)$ is the intersection of all closed halfspaces that contain it, so its recession cone contains the (n+1)st axis. Since the recession cones of $\operatorname{cl}(C)$ and $\operatorname{ri}(C)$ coincide [cf. Prop. 1.4.2(b)], for every $(\overline{u}, \overline{w}) \in \operatorname{ri}(C)$, the vertical line $\{(\overline{u}, w) \mid w \in \Re\}$ belongs to $\operatorname{ri}(C)$ and hence to C. This contradicts the assumption that C does not contain a vertical line.

(b) If $(\overline{u}, \overline{w}) \notin cl(C)$, then there exists a hyperplane strictly separating $(\overline{u}, \overline{w})$ and cl(C) [cf. Prop. 1.5.3 under condition (2)]. If this hyperplane is nonvertical, since $C \subset cl(C)$, we are done, so assume otherwise. Then, we have a nonzero vector $\overline{\mu}$ and a scalar $\overline{\gamma}$ such that

$$\overline{\mu}'u > \overline{\gamma} > \overline{\mu}'\overline{u}, \qquad \forall \ (u, w) \in \mathrm{cl}(C).$$
 (1.34)

The idea now is to combine this vertical hyperplane with a suitable nonvertical hyperplane in order to construct a nonvertical hyperplane that strictly separates $(\overline{u}, \overline{w})$ from cl(C) (see Fig. 1.5.11).

Since, by assumption, C does not contain a vertical line, $\operatorname{ri}(C)$ also does not contain a vertical line. Since the recession cones of $\operatorname{cl}(C)$ and $\operatorname{ri}(C)$ coincide [cf. Prop. 1.4.2(b)], it follows that $\operatorname{cl}(C)$ does not contain a vertical line. Hence, by part (a), there exists a nonvertical hyperplane containing $\operatorname{cl}(C)$ in one of its closed halfspaces, so that for some (μ, β) and γ , with $\beta \neq 0$, we have

$$\mu'u + \beta w \ge \gamma, \quad \forall (u, w) \in cl(C).$$

By multiplying this relation with an $\epsilon > 0$ and combining it with Eq. (1.34), we obtain

$$(\overline{\mu} + \epsilon \mu)' u + \epsilon \beta w > \overline{\gamma} + \epsilon \gamma, \quad \forall (u, w) \in cl(C), \ \forall \ \epsilon > 0.$$

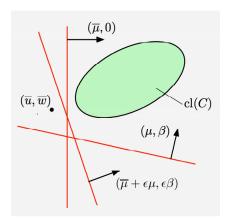


Figure 1.5.11. Construction of a strictly separating nonvertical hyperplane in the proof of Prop. 1.5.8(b).

Since $\overline{\gamma} > \overline{\mu}' \overline{u}$, there is a small enough ϵ such that

$$\overline{\gamma} + \epsilon \gamma > (\overline{\mu} + \epsilon \mu)' \overline{u} + \epsilon \beta \overline{w}.$$

From the above two relations, we obtain

$$(\overline{\mu} + \epsilon \mu)' u + \epsilon \beta w > (\overline{\mu} + \epsilon \mu)' \overline{u} + \epsilon \beta \overline{w}, \quad \forall (u, w) \in cl(C),$$

implying that there is a nonvertical hyperplane with normal

$$(\overline{\mu} + \epsilon \mu, \epsilon \beta)$$

that strictly separates $(\overline{u}, \overline{w})$ and $\operatorname{cl}(C)$. Since $C \subset \operatorname{cl}(C)$, this hyperplane also strictly separates $(\overline{u}, \overline{w})$ and C. Q.E.D.

1.6 CONJUGATE FUNCTIONS

We will now develop a concept that is fundamental in convex optimization. This is the conjugacy transformation, which associates with any function f, a convex function, called the conjugate of f. The idea here is to describe f in terms of the affine functions that are majorized by f. When f is closed proper convex, we will show that the description is accurate and the transformation is symmetric, i.e., f can be recovered by taking the conjugate of the conjugate of f. The conjugacy transformation thus provides an alternative view of a convex function, which often reveals interesting properties, and is useful for analysis and computation.

Consider an extended real-valued function $f: \Re^n \mapsto [-\infty, \infty]$. The conjugate function of f is the function $f^*: \Re^n \mapsto [-\infty, \infty]$ defined by

$$f^{\star}(y) = \sup_{x \in \mathbb{R}^n} \left\{ x'y - f(x) \right\}, \qquad y \in \mathbb{R}^n.$$
 (1.35)

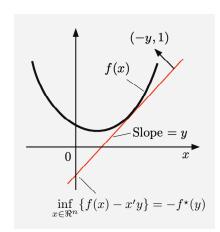


Figure 1.6.1. Visualization of the conjugate function

$$f^{\star}(y) = \sup_{x \in \Re^n} \left\{ x'y - f(x) \right\}$$

of a function f. The crossing point of the vertical axis with the hyperplane that has normal (-y, 1) and passes through a point $(\overline{x}, f(\overline{x}))$ on the graph of f is

$$f(\overline{x}) - \overline{x}'y$$
.

Thus, the crossing point corresponding to the hyperplane that supports the epigraph of f is

$$\inf_{x \in \Re^n} \{ f(x) - x'y \},\,$$

which by definition is equal to $-f^*(y)$.

Figure 1.6.1 provides a geometrical interpretation of the definition.

Note that regardless of the structure of f, the conjugate f^* is a closed convex function, since it is the pointwise supremum of the collection of affine functions

$$x'y - f(x)$$
, $\forall x \text{ such that } f(x) \text{ is finite,}$

(Prop. 1.1.6). Note also that f^* need not be proper, even if f is. We will show, however, that in the case where f is convex, f^* is proper if and only if f is.

Figure 1.6.2 shows some examples of conjugate functions. In this figure, all the functions are closed proper convex, and it can be verified that the conjugate of the conjugate yields the original function. This is a manifestation of a result that we will show shortly.

For a function $f: \mathbb{R}^n \mapsto [-\infty, \infty]$, consider the conjugate of the conjugate function f^* (or *double conjugate*). It is denoted by f^{**} , it is given by

$$f^{\star\star}(x) = \sup_{y \in \Re^n} \{y'x - f^\star(y)\}, \qquad x \in \Re^n,$$

and it can be constructed as shown in Fig. 1.6.3. As this figure suggests, and part (d) of the following proposition shows, by constructing f^{**} , we typically obtain the convex closure of f [the function that has as epigraph the closure of the convex hull of $\operatorname{epi}(f)$; cf. Section 1.3.3]. In particular, part (c) of the proposition shows that if f is closed proper convex, then $f^{**} = f$.

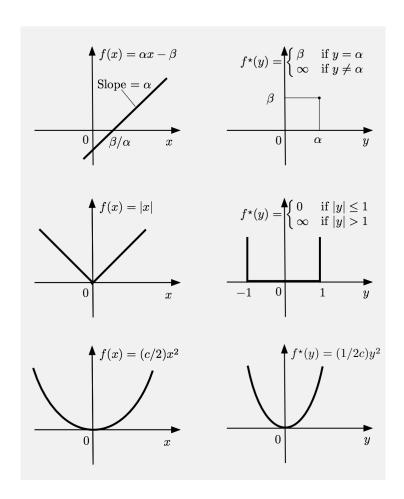


Figure 1.6.2. Some examples of conjugate functions. It can be verified that in each case, the conjugate of the conjugate is the original, i.e., the conjugates of the functions on the right are the corresponding functions on the left.

Proposition 1.6.1: (Conjugacy Theorem) Let $f: \Re^n \mapsto [-\infty, \infty]$ be a function, let f^* be its conjugate, and consider the double conjugate f^{**} . Then:

(a) We have

$$f(x) \ge f^{\star\star}(x), \quad \forall \ x \in \Re^n.$$

(b) If f is convex, then properness of any one of the functions f, f^* , and f^{**} implies properness of the other two.

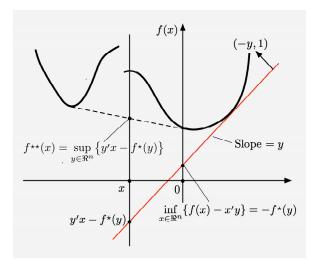


Figure 1.6.3. Visualization of the double conjugate (conjugate of the conjugate)

$$f^{\star\star}(x) = \sup_{y \in \Re^n} \left\{ y'x - f^{\star}(y) \right\}$$

of a function f, where f^* is the conjugate of f,

$$f^{\star}(y) = \sup_{x \in \Re^n} \left\{ x'y - f(x) \right\}.$$

For each $x \in \Re^n$, we consider the vertical line in \Re^{n+1} that passes through the point (x,0), and for each y in the effective domain of f^\star , we consider the crossing point of this line with the hyperplane with normal (-y,1) that supports the graph of f [and therefore passes through the point $(0,-f^\star(y))$]. This crossing point is $y'x-f^\star(y)$, so $f^{\star\star}(x)$ is equal to the highest crossing level. As the figure indicates (and Prop. 1.6.1 shows), the double conjugate $f^{\star\star}$ is the convex closure of f (barring the exceptional case where the convex closure takes the value $-\infty$ at some point, in which case the figure above is not valid).

(c) If f is closed proper convex, then

$$f(x) = f^{\star\star}(x), \quad \forall \ x \in \Re^n.$$

(d) The conjugates of f and its convex closure $\operatorname{\check{cl}} f$ are equal. Furthermore, if $\operatorname{\check{cl}} f$ is proper, then

$$(\check{\operatorname{cl}} f)(x) = f^{\star\star}(x), \quad \forall \ x \in \Re^n.$$

Proof: (a) For all x and y, we have

$$f^{\star}(y) \ge x'y - f(x),$$

so that

$$f(x) \ge x'y - f^*(y), \quad \forall x, y \in \Re^n.$$

Hence

$$f(x) \geq \sup_{y \in \Re^n} \left\{ x'y - f^{\star}(y) \right\} = f^{\star \star}(x), \qquad \forall \ x \in \Re^n.$$

(b) Assume that f is proper, in addition to being convex. Then its epigraph is nonempty and convex, and contains no vertical line. By applying the Nonvertical Hyperplane Theorem [Prop. 1.5.8(a)], with C being the set $\operatorname{epi}(f)$, it follows that there exists a nonvertical hyperplane with normal (y,1) that contains $\operatorname{epi}(f)$ in its positive halfspace. In particular, this implies the existence of a vector y and a scalar c such that

$$y'x + f(x) \ge c, \quad \forall x \in \Re^n.$$

We thus obtain

$$f^{\star}(-y) = \sup_{x \in \Re^n} \left\{ -y'x - f(x) \right\} \le -c,$$

so that f^* is not identically equal to ∞ . Also, by the properness of f, there exists a vector \overline{x} such that $f(\overline{x})$ is finite. For every $y \in \Re^n$, we have

$$f^{\star}(y) > y'\overline{x} - f(\overline{x}),$$

so $f^*(y) > -\infty$ for all $y \in \Re^n$. Thus, f^* is proper.

Conversely, assume that f^* is proper. The preceding argument shows that properness of f^* implies properness of its conjugate, f^{**} , so that $f^{**}(x) > -\infty$ for all $x \in \Re^n$. By part (a), $f(x) \geq f^{**}(x)$, so $f(x) > -\infty$ for all $x \in \Re^n$. Also, f cannot be identically equal to ∞ , since then by its definition, f^* would be identically equal to $-\infty$. Thus f is proper.

We have thus shown that a convex function is proper if and only if its conjugate is proper, and the result follows in view of the conjugacy relations between f, f^* , and f^{**} .

(c) We will apply the Nonvertical Hyperplane Theorem (Prop. 1.5.8), with C being the closed and convex set $\operatorname{epi}(f)$, which contains no vertical line since f is proper. Let (x,γ) belong to $\operatorname{epi}(f^{\star\star})$, i.e., $x\in\operatorname{dom}(f^{\star\star})$, $\gamma\geq f^{\star\star}(x)$, and suppose, to arrive at a contradiction, that (x,γ) does not belong to $\operatorname{epi}(f)$. Then by Prop. 1.5.8(b), there exists a nonvertical hyperplane with normal (y,ζ) , where $\zeta\neq 0$, and a scalar c such that

$$y'z + \zeta w < c < y'x + \zeta \gamma, \quad \forall (z, w) \in epi(f).$$

Since w can be made arbitrarily large, we have $\zeta < 0$, and without loss of generality, we can take $\zeta = -1$, so that

$$y'z - w < c < y'x - \gamma$$
, $\forall (z, w) \in epi(f)$.

Since $\gamma \geq f^{**}(x)$ and $(z, f(z)) \in \operatorname{epi}(f)$ for all $z \in \operatorname{dom}(f)$, we obtain

$$y'z - f(z) < c < y'x - f^{\star\star}(x), \quad \forall z \in \text{dom}(f).$$

Hence

$$\sup_{z \in \Re^n} \left\{ y'z - f(z) \right\} \le c < y'x - f^{\star\star}(x),$$

or

$$f^{\star}(y) < y'x - f^{\star\star}(x),$$

which contradicts the definition $f^{\star\star}(x) = \sup_{y \in \mathbb{R}^n} \{y'x - f^{\star}(y)\}$. Thus, we have $\operatorname{epi}(f^{\star\star}) \subset \operatorname{epi}(f)$, which implies that $f(x) \leq f^{\star\star}(x)$ for all $x \in \mathbb{R}^n$. This, together with part (a), shows that $f^{\star\star}(x) = f(x)$ for all x.

(d) Let \check{f}^* be the conjugate of $\check{\operatorname{cl}} f$. For any $y, -f^*(y)$ and $-\check{f}^*(y)$ are the supremum crossing levels of the vertical axis with the hyperplanes with normal (-y,1) that contain the sets $\operatorname{epi}(f)$ and $\operatorname{cl}(\operatorname{conv}(\operatorname{epi}(f)))$, respectively, in their positive closed halfspaces (cf. Fig. 1.6.1). Since the hyperplanes of this type are the same for the two sets, we have $f^*(y) = \check{f}^*(y)$ for all y. Thus, f^{**} is equal to the conjugate of \check{f}^* , which is $\check{\operatorname{cl}} f$ by part (c) when $\check{\operatorname{cl}} f$ is proper. Q.E.D.

The properness assumptions on f and $\operatorname{cl} f$ are essential for the validity of parts (c) and (d), respectively, of the preceding proposition. For an illustrative example, consider the closed convex (but improper) function

$$f(x) = \begin{cases} \infty & \text{if } x > 0, \\ -\infty & \text{if } x \le 0. \end{cases}$$
 (1.36)

We have $f = \operatorname{cl} f$, and it can be verified that $f^*(y) = \infty$ for all y and $f^{**}(x) = -\infty$ for all x, so that $f \neq f^{**}$ and $\operatorname{cl} f \neq f^{**}$.

For an example where f is proper (but not closed convex), while cl f is improper and we have cl $f \neq f^{\star\star}$, let

$$f(x) = \begin{cases} \log(-x) & \text{if } x < 0, \\ \infty & \text{if } x \ge 0. \end{cases}$$

Then $\operatorname{cl} f$ is equal to the function (1.36), and $\operatorname{cl} f \neq f^{\star\star}$.

The exceptional behavior in the preceding example can be attributed to a subtle difference in the constructions of the conjugate function and the convex closure: while the conjugate functions f^* and f^{**} are defined exclusively in terms of nonvertical hyperplanes via the construction of Fig.

1.6.3, the epigraph of the convex closure cl f is defined in terms of nonvertical and vertical hyperplanes. This difference is inconsequential when there exists at least one nonvertical hyperplane containing the epigraph of f in one of its closed halfspaces [this is true in particular if cl f is proper; see Props. 1.5.8(a) and 1.6.1(d)]. The reason is that, in this case, the epigraph of cl f can equivalently be defined by using just nonvertical hyperplanes [this can be seen using Prop. 1.5.8(b)].

Example 1.6.1: (Indicator/Support Function Conjugacy)

Given a nonempty set X, consider the *indicator function* of X, defined by

$$\delta_X(x) = \begin{cases} 0 & \text{if } x \in X, \\ \infty & \text{if } x \notin X. \end{cases}$$

The conjugate of δ_X is given by

$$\sigma_X(y) = \sup_{x \in X} y'x$$

and is called the *support function of* X (see Fig. 1.6.4). By the generic closedness and convexity properties of conjugate functions, σ_X is closed and convex. It is also proper since X is nonempty, so that $\sigma_X(0) = 0$ (an improper closed convex function cannot take finite values; cf. the discussion at the end of Section 1.1.2). Furthermore, the sets X, cl(X), conv(X), and cl(conv(X)) all have the same support function [cf. Prop. 1.6.1(d)].

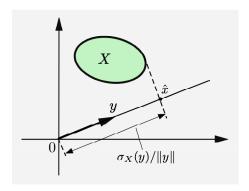


Figure 1.6.4. Visualization of the support function

$$\sigma_X(y) = \sup_{x \in X} y'x$$

of a set X. To determine the value $\sigma_X(y)$ for a given vector y, we project the set X on the line determined by y, and we find \hat{x} , the extreme point of projection in the direction y. Then

$$\sigma_X(y) = \|\hat{x}\| \cdot \|y\|.$$

Example 1.6.2: (Support Function of a Cone - Polar Cones)

Let C be a convex cone. By the preceding example, the conjugate of its indicator function δ_C is its support function,

$$\sigma_C(y) = \sup_{x \in C} y'x.$$

Since C is a cone, we see that

$$\sigma_C(y) = \begin{cases} 0 & \text{if } y'x \le 0, \ \forall \ x \in C, \\ \infty & \text{otherwise.} \end{cases}$$

Thus the support/conjugate function σ_C is the indicator function δ_{C^*} of the cone

$$C^* = \{ y \mid y'x \le 0, \ \forall \ x \in C \}, \tag{1.37}$$

called the polar cone of C. It follows that the conjugate of σ_C is the indicator function of the polar cone of C^* , and by the Conjugacy Theorem [Prop. 1.6.1(c)] it is also $\operatorname{cl} \delta_C$. Thus the polar cone of C^* is $\operatorname{cl}(C)$. In particular, if C is closed, the polar of its polar is equal to the original. This is a special case of the *Polar Cone Theorem*, which will be discussed in more detail in Section 2.2.

A special case of particular interest in when $C = \text{cone}(\{a_1, \ldots, a_r\})$, the cone generated by a finite set of vectors a_1, \ldots, a_r . Then it can be seen that

$$C^* = \{x \mid a_j' x \le 0, j = 1, \dots, r\}.$$

From this it follows that $(C^*)^* = C$ because C is a closed set [we have essentially shown this: $\operatorname{cone}\left(\{a_1,\ldots,a_r\}\right)$ is the image of the positive orthant $\{\alpha \mid \alpha \geq 0\}$ under the linear transformation that maps α to $\sum_{j=1}^r \alpha_j a_j$, and the image of any polyhedral set under a linear transformation is a closed set (see the discussion following the proof of Prop. 1.4.13)]. The assertion $(C^*)^* = C$ for the case $C = \operatorname{cone}\left(\{a_1,\ldots,a_r\}\right)$ is known as Farkas' Lemma and will be discussed further in what follows (Sections 2.3 and 5.1).

Let us finally note that we may define the polar cone of any set C via Eq. (1.37). However, unless C is a cone, the support function of C will not be an indicator function, nor will the relation $(C^*)^* = C$ hold. Instead, we will show in Section 2.2 that in general we have $(C^*)^* = \text{cl}(\text{cone}(C))$.

1.7 SUMMARY

In this section, we discuss how the material of this chapter is used in subsequent chapters. First, let us note that we have aimed to develop in this chapter the part of convexity theory that is needed for the optimization and duality analysis of Chapters 3-5, and no more. We have developed the basic principles of polyhedral convexity in Chapter 2 for completeness and a broader perspective, but this material is not needed for the mathematical development of Chapters 3-5. Despite the fact that we focus only on the essential, we still cover considerable ground, and it may help the reader to know how the various topics of this chapter connect with each other, and how they are used later. We thus provide a summary and guide on a section-by-section basis:

Section 1.1: The definitions and results on convexity and closure (Sections 1.1.1-1.1.3) are very basic and should be read carefully. The material of

Section 1.1.4 on differentiable convex functions, optimality conditions, and the projection theorem, is also basic. The reader may skip Prop. 1.1.10 on twice differentiable functions, which is used in nonessential ways later.

Section 1.2: The definitions of convex and affine hulls, and generated cones, as well as Caratheodory's theorem should also be read carefully. Proposition 1.2.2 on the compactness of the convex hull of a compact set is used only in the proof of Prop. 4.3.2, which is in turn used only in the specialized MC/MC framework of Section 5.7 on estimates of duality gap.

Section 1.3: The definitions and results on relative interior and closure up to and including Section 1.3.2 are pervasive in what follows. However, the somewhat tedious proofs of Props. 1.3.5-1.3.10 may be skipped at first reading. Similarly, the proof of continuity of a real-valued convex function (Prop. 1.3.11) is specialized and may be skipped. Regarding Section 1.3.3, it is important to understand the definitions, and gain intuition on closures and convex closures of functions, as they arise in the context of conjugacy. However, Props. 1.3.13-1.3.17 are used substantively later only for the development of minimax theory (Sections 3.4, 4.2.5, and 5.5) and the theory of directional derivatives (Section 5.4.4), which are themselves "terminal" and do not affect other sections.

Section 1.4: The material on directions of recession, up to and including Section 1.4.1, is very important for later developments, although the use of Props. 1.4.5-1.4.6 on recession functions is somewhat focused. In particular, Prop. 1.4.7 is used only to prove Prop. 1.4.8, and Props. 1.4.6-1.4.8 are used only for the development of existence of solutions criteria in Section 3.2. The set intersection analysis of Section 1.4.2, which may challenge some readers at first, is used for the development of some important theory: the existence of optimal solutions for linear and quadratic programs (Prop. 1.4.12), the criteria for preservations of closedness under linear transformations and vector sum (Props. 1.4.13 and 1.4.14), and the existence of solutions analysis of Chapter 3.

Section 1.5: The material on hyperplane separation is used extensively and should be read in its entirety. However, the long proof of the polyhedral proper separation theorem (Prop. 1.5.7, due to Rockafellar [Roc70], Th. 20.2) may be postponed for later. This theorem is important for our purposes. For example it is used (through Props. 4.5.1 and 4.5.2) to establish part of the Nonlinear Farkas' Lemma (Prop. 5.1.1), on which the constrained optimization and Fenchel duality theories rest.

Section 1.6: The material on conjugate functions and the Conjugacy Theorem (Prop. 1.6.1) is very basic. In our duality development, we use the theorem somewhat sparingly because we argue mostly in terms of the MC/MC framework of Chapter 4 and the Strong Duality Theorem (Prop. 4.3.1), which serves as an effective substitute.

Basic Concepts of Polyhedral Convexity

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In this chapter, we discuss polyhedral sets, i.e., nonempty sets specified by systems of a finite number of affine inequalities

$$a_j'x \le b_j, \qquad j = 1, \dots, r,$$

where a_1, \ldots, a_r are vectors in \Re^n , and b_1, \ldots, b_r are scalars. We start with the notion of an extreme point of a convex set. We then introduce a general duality relation between cones, and we specialize this relation to polyhedral cones and polyhedral sets. Then we apply the results to linear programming problems.†

2.1 EXTREME POINTS

A geometrically apparent property of a bounded polyhedral set in the plane is that it can be represented as the convex hull of a finite number of points, its "corner" points. This property has great analytical and algorithmic significance, and can be formalized through the notion of an extreme point, which we now introduce.

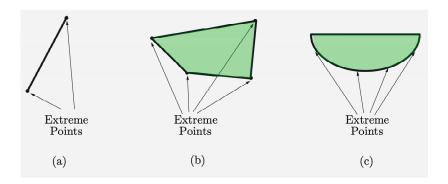


Figure 2.1.1. Illustration of extreme points of various convex sets in \Re^2 . For the set in (c), the extreme points are the ones that lie on the circular arc.

Given a nonempty convex set C, a vector $x \in C$ is said to be an extreme point of C if it does not lie strictly between the endpoints of any line segment contained in the set, i.e., if there do not exist vectors $y \in C$ and $z \in C$, with $y \neq x$ and $z \neq x$, and a scalar $\alpha \in (0,1)$ such that $x = \alpha y + (1 - \alpha)z$. It can be seen that an equivalent definition is that x cannot be expressed as a convex combination of some vectors of C, all of which are different from x.

[†] The material in this chapter is fundamental in itself, and is related to several of the topics treated later. However, the results of the chapter are not used directly in Chapters 3-5, and may be skipped by the reader.

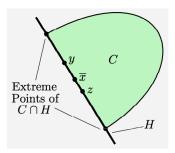


Figure 2.1.2. Construction to prove that the extreme points of $C \cap H$ are also extreme points of C (Prop. 2.1.1).

For some perspective, note that an interior point of a set cannot be an extreme point, so an open set has no extreme points. Also from the definition, we see that a convex cone may have at most one extreme point, the origin. We will show later in this section that a polyhedral set has at most a finite number of extreme points (possibly none). Moreover a concave (e.g., linear) function that attains a minimum over a polyhedral set with at least one extreme point, attains the minimum at some extreme point. Figure 2.1.1 illustrates the extreme points of various types of sets.

The following is a useful result for the analysis of extreme pointrelated issues. It is often used in proofs to obtain extreme points of convex sets by using the extreme points of lower dimensional subsets.

Proposition 2.1.1: Let C be a convex subset of \Re^n , and let H be a hyperplane that contains C in one of its closed halfspaces. Then the extreme points of $C \cap H$ are precisely the extreme points of C that belong to H.

Proof: Let \overline{x} be an extreme point of $C \cap H$. To show that \overline{x} is an extreme point of C, let $y \in C$ and $z \in C$ be such that $\overline{x} = \alpha y + (1 - \alpha)z$ for some $\alpha \in (0,1)$ (see Fig. 2.1.2). We will show that $\overline{x} = y = z$. Indeed, since $\overline{x} \in H$, the closed halfspace containing C is of the form $\{x \mid a'x \geq a'\overline{x}\}$, where $a \neq 0$, and H is of the form $\{x \mid a'x = a'\overline{x}\}$. Thus, we have $a'y \geq a'\overline{x}$ and $a'z \geq a'\overline{x}$, which in view of $\overline{x} = \alpha y + (1 - \alpha)z$, implies that $a'y = a'\overline{x}$ and $a'z = a'\overline{x}$. Therefore,

$$y \in C \cap H$$
, $z \in C \cap H$.

Since \overline{x} is an extreme point of $C \cap H$, it follows that $\overline{x} = y = z$, so that \overline{x} is an extreme point of C.

Conversely, if $\overline{x} \in H$ and \overline{x} is an extreme point of C, then \overline{x} cannot lie strictly between two distinct points of C, and hence also strictly between two distinct points of $C \cap H$. It follows that \overline{x} is an extreme point of $C \cap H$. Q.E.D.

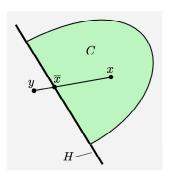


Figure 2.1.3. Construction used in the induction proof of Prop. 2.1.2 to show that if a closed convex set C does not contain a line, it must have an extreme point. Using the line segment connecting two points $x \in C$ and $y \notin C$, we obtain a relative boundary point \overline{x} of C. The proof then reduces to asserting the existence of an extreme point of the lower-dimensional set $C \cap H$, where H is a supporting hyperplane passing through \overline{x} , which is also an extreme point of C by Prop. 2.1.1.

The following proposition characterizes the existence of an extreme point.

Proposition 2.1.2: A nonempty closed convex subset of \Re^n has at least one extreme point if and only if it does not contain a line, i.e., a set of the form $\{x + \alpha d \mid \alpha \in \Re\}$, where x and d are vectors in \Re^n with $d \neq 0$.

Proof: Let C be a convex set that has an extreme point x. Assume, to arrive at a contradiction, that C contains a line $\{\overline{x} + \alpha d \mid \alpha \in \Re\}$, where $\overline{x} \in C$ and $d \neq 0$. Then, by the Recession Cone Theorem [Prop. 1.4.1(b)] and the closedness of C, both d and -d are directions of recession of C, so the line $\{x + \alpha d \mid \alpha \in \Re\}$ belongs to C. This contradicts the fact that x is an extreme point.

Conversely, we use induction on the dimension of the space to show that if C does not contain a line, it must have an extreme point. This is true for \Re , so assume it is true for \Re^{n-1} , where $n \geq 2$. We will show that any nonempty closed convex subset C of \Re^n , which does not contain a line, must have an extreme point. Since C does not contain a line, there must exist points $x \in C$ and $y \notin C$. The line segment connecting x and y intersects the relative boundary of C at some point \overline{x} , which belongs to C, since C is closed. Consider a hyperplane H passing through \overline{x} and containing C in one of its closed halfspaces (such a hyperplane exists by Prop. 1.5.1; see Fig. 2.1.3). The set $C \cap H$ lies in an (n-1)-dimensional space and does not contain a line, since C does not. Hence, by the induction hypothesis, it must have an extreme point. By Prop. 2.1.1, this extreme point must also be an extreme point of C. Q.E.D.

As an example of application of the preceding proposition, consider a nonempty set of the form

$${x \mid x \in C, x \ge 0},$$

where C is a closed convex set. Such sets arise commonly in optimization, and in view of the constraint $x \ge 0$, they do not contain a line. Hence, by Prop. 2.1.2, they always have at least one extreme point. The following is a useful generalization.

Proposition 2.1.3: Let C be a nonempty closed convex subset of \Re^n . Assume that for some $m \times n$ matrix A of rank n and some $b \in \Re^m$, we have

$$Ax \ge b, \quad \forall \ x \in C.$$

Then C has at least one extreme point.

Proof: The set $\{x \mid Ax \geq b\}$ cannot contain a line $\{\overline{x} + \alpha d \mid \alpha \in \Re\}$, with $d \neq 0$; if it did, we would have

$$A\overline{x} + \alpha Ad > 0, \quad \forall \ \alpha \in \Re,$$

a contradiction because $Ad \neq 0$, since A has rank n. Hence C does not contain a line, and the result follows from Prop. 2.1.2. Q.E.D.

Extreme Points of Polyhedral Sets

We recall the definition of a polyhedral set in \Re^n : it is a set that is nonempty and has the form

$$\{x \mid a_j' x \le b_j, j = 1, \dots, r\},\$$

where a_1, \ldots, a_r are vectors in \Re^n and b_1, \ldots, b_r are scalars.

The following proposition gives a characterization of extreme points of polyhedral sets that is central in the theory of linear programming (see Fig. 2.1.4). Parts (b) and (c) of the proposition are special cases of part (a), but we state them explicitly for completeness.

Proposition 2.1.4: Let P be a polyhedral subset of \Re^n .

(a) If P has the form

$$P = \{x \mid a'_{j}x \le b_{j}, j = 1, \dots, r\},\$$

where $a_j \in \mathbb{R}^n$, $b_j \in \mathbb{R}$, j = 1, ..., r, then a vector $v \in P$ is an extreme point of P if and only if the set

$$A_v = \{a_j \mid a_j'v = b_j, j \in \{1, \dots, r\}\}$$

contains n linearly independent vectors.

(b) If P has the form

$$P = \{x \mid Ax = b, x > 0\},\$$

where A is an $m \times n$ matrix and b is a vector in \Re^m , then a vector $v \in P$ is an extreme point of P if and only if the columns of A corresponding to the nonzero coordinates of v are linearly independent.

(c) If P has the form

$$P = \{x \mid Ax = b, c \le x \le d\},\$$

where A is an $m \times n$ matrix, b is a vector in \Re^m , and c, d are vectors in \Re^n , then a vector $v \in P$ is an extreme point of P if and only if the columns of A corresponding to the coordinates of v that lie strictly between the corresponding coordinates of v and v are linearly independent.

Proof: (a) If the set A_v contains fewer than n linearly independent vectors, then the system of equations

$$a_j'w = 0, \quad \forall \ a_j \in A_v,$$

has a nonzero solution, call it \overline{w} . For sufficiently small $\gamma > 0$, we have $v + \gamma \overline{w} \in P$ and $v - \gamma \overline{w} \in P$, thus showing that v is not an extreme point. Thus, if v is an extreme point, A_v must contain n linearly independent vectors.

Conversely, assume that A_v contains a subset \bar{A}_v consisting of n linearly independent vectors. Suppose that for some $y \in P$, $z \in P$, and $\alpha \in (0,1)$, we have

$$v = \alpha y + (1 - \alpha)z$$
.

Then, for all $a_i \in \bar{A}_v$, we have

$$b_j = a'_j v = \alpha a'_j y + (1 - \alpha) a'_j z \le \alpha b_j + (1 - \alpha) b_j = b_j.$$

Thus, v, y, and z are all solutions of the system of n linearly independent equations

$$a_j'w = b_j, \quad \forall \ a_j \in \bar{A}_v.$$

It follows that v = y = z, implying that v is an extreme point of P.

(b) We write P in terms of inequalities in the equivalent form

$$P = \{x \mid Ax \le b, -Ax \le -b, -x \le 0\}.$$

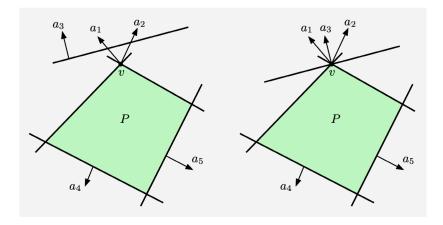


Figure 2.1.4. Illustration of extreme points of a 2-dimensional polyhedral set

$$P = \{x \mid a'_{j}x \le b_{j}, j = 1, \dots, r\},\$$

[see Prop. 2.1.4(a)]. For a vector $v \in P$ to be an extreme point it is necessary and sufficient that the set of vectors

$$A_v = \{a_j \mid a_j'v = b_j, j \in \{1, \dots, r\}\}$$

contains n=2 linearly independent vectors. In the case on the left-hand side, the set A_v consists of the two linearly independent vectors a_1 and a_2 . In the case on the right-hand side, the set A_v consists of the three vectors a_1, a_2, a_3 , any pair of which are linearly independent.

Let S_v be the set of vectors of the form $(0, \ldots, 0, -1, 0, \ldots 0)$, where -1 is in a position corresponding to a zero coordinate of v. By applying the result of part (a) to the above inequality representation of P, we see that v is an extreme point if and only if the rows of A together with the vectors in S_v contain n linearly independent vectors. Let \bar{A} be the matrix, which is the same as A except that all the columns corresponding to the zero coordinates of v are set to zero. It follows that v is an extreme point if and only if \bar{A} contains n-k linearly independent rows, where k is the number of vectors in S_v . The only (possibly) nonzero columns of \bar{A} are the n-k corresponding to nonzero coordinates of v, so these columns must be linearly independent. Since these columns are shared by \bar{A} and A, the result follows.

(c) The proof is essentially the same as the proof of part (b). Q.E.D.

By combining Props. 2.1.3 and 2.1.4(a), we obtain the following characterization of the existence of an extreme point of a polyhedral set.

Proposition 2.1.5: A polyhedral set in \Re^n of the form

$$\left\{x \mid a_j' x \le b_j, \ j = 1, \dots, r\right\}$$

has an extreme point if and only if the set $\{a_j \mid j=1,\ldots,r\}$ contains n linearly independent vectors.

Example 2.1.1: (Birkhoff-von Neumann Theorem)

In the space of $n \times n$ matrices $X = \{x_{ij} \mid i, j = 1, ..., n\}$, consider the polyhedral set

$$P = \left\{ X \mid x_{ij} \ge 0, \sum_{j=1}^{n} x_{ij} = 1, \sum_{i=1}^{n} x_{ij} = 1, i, j = 1, \dots, n \right\}.$$

Matrices in this set are called *doubly stochastic*, in view of the fact that both their rows and columns are probability distributions.

An important optimization problem that involves P is the assignment problem. Here, the elements of some finite set (say, n persons) are to be matched on a one-to-one basis, to the elements of another set (say, n objects), and x_{ij} is a variable taking value 0 or 1, and representing whether person i is or is not matched to object j. Thus a feasible solution to the assignment problem is a matrix that has a single 1 in each row and a single 1 in each column, with 0s in all other positions. Such a matrix is called a permutation matrix, since when multiplying a vector $x \in \Re^n$ it produces a vector whose coordinates are a permutation of the coordinates of x. The assignment problem is to minimize a linear cost of the form

$$\sum_{i=1}^{n} \sum_{j=1}^{n} a_{ij} x_{ij}$$

subject to X being a permutation matrix.

The Birkhoff-von Neumann Theorem states that the extreme points of the polyhedral set of doubly stochastic matrices are the permutation matrices. Thus if we minimize $\sum_{i=1}^n \sum_{j=1}^n a_{ij} x_{ij}$ subject to $X \in P$, using an algorithm that is guaranteed to find an optimal solution that is an extreme point (the simplex method, among others, offers this guarantee), we also solve the assignment problem. To prove the Birkhoff-von Neumann Theorem, we note that if X is a permutation matrix, and $Y, Z \in P$ are such that

$$x_{ij} = \alpha y_{ij} + (1 - \alpha)z_{ij}, \quad \forall i, j = 1, \dots, n,$$

for some $\alpha \in (0, 1)$, then since $y_{ij}, z_{ij} \in [0, 1]$, we have $y_{ij} = z_{ij} = 0$ if $x_{ij} = 0$, and $y_{ij} = z_{ij} = 1$ if $x_{ij} = 1$. Thus, if

$$X = \alpha Y + (1 - \alpha)Z,$$

we have X = Y = Z, implying that X is an extreme point of P.

We now argue that by Prop. 2.1.1, X is an extreme point of the polyhedral set $P \cap H$, where H is the hyperplane of matrices whose component at row \overline{i} and column \overline{j} is equal to 1. Repeating the preceding argument with P replaced by $P \cap H$, we see that there must exist a second row of X with a single nonzero component. Continuing in this manner, by finding at each step a new row and column with a single nonzero component, we establish that all rows and columns of X have a single nonzero component, showing that X is a permutation matrix.

2.2 POLAR CONES

We will now consider an important type of cone that is associated with a nonempty set C, and provides a dual and equivalent description of C, when C is itself a closed cone. In particular, the *polar cone* of C, denoted C^* , is given by

$$C^* = \{ y \mid y'x < 0, \ \forall \ x \in C \},\$$

(see Fig. 2.2.1). We have already discussed briefly polar cones in Example 1.6.2.

Clearly, C^* is a cone, and being the intersection of a collection of closed halfspaces, it is closed and convex (regardless of whether C is closed and/or convex). If C is a subspace, it can be seen that C^* is equal to the orthogonal subspace C^{\perp} . The following proposition generalizes the equality $C = (C^{\perp})^{\perp}$, which holds in the case where C is a subspace. Part (b) of the proposition was essentially proved in Example 1.6.2 using the Conjugacy Theorem (Prop. 1.6.1), and the conjugacy of indicator and support functions, but we will give here an alternative, more elementary proof that uses the Projection Theorem.

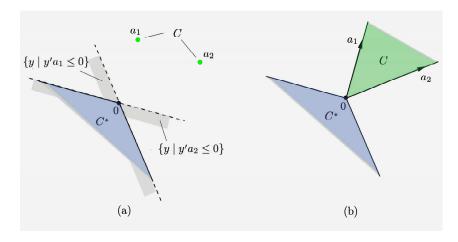


Figure 2.2.1. Illustration of the polar cone C^* of a subset C of \Re^2 . In (a), C consists of just two points, a_1 and a_2 , and C^* is the intersection of the two closed halfspaces $\{y \mid y'a_1 \leq 0\}$ and $\{y \mid y'a_2 \leq 0\}$. In (b), C is the convex cone

$${x \mid x = \mu_1 a_1 + \mu_2 a_2, \, \mu_1 \ge 0, \, \mu_2 \ge 0},$$

and the polar cone C^* is the same as in case (a).

Proposition 2.2.1:

(a) For any nonempty set C, we have

$$C^* = (\text{cl}(C))^* = (\text{conv}(C))^* = (\text{cone}(C))^*.$$

(b) (Polar Cone Theorem) For any nonempty cone C, we have

$$(C^*)^* = \operatorname{cl}(\operatorname{conv}(C)).$$

In particular, if C is closed and convex, we have $(C^*)^* = C$.

Proof: (a) For any two sets X and Y with $X \supset Y$, we have $X^* \subset Y^*$, from which it follows that $(\operatorname{cl}(C))^* \subset C^*$. Conversely, if $y \in C^*$, then $y'x_k \leq 0$ for all k and all sequences $\{x_k\} \subset C$, so that $y'x \leq 0$ for all $x \in \operatorname{cl}(C)$. Hence $y \in (\operatorname{cl}(C))^*$, implying that $C^* \subset (\operatorname{cl}(C))^*$.

Also, since $\operatorname{conv}(C) \supset C$, we have $\left(\operatorname{conv}(C)\right)^* \subset C^*$. Conversely, if $y \in C^*$, then $y'x \leq 0$ for all $x \in C$, so that $y'z \leq 0$ for all z that are convex combinations of vectors $x \in C$. Hence, $y \in \left(\operatorname{conv}(C)\right)^*$ and $C^* \subset \left(\operatorname{conv}(C)\right)^*$. A nearly identical argument also shows that $C^* = \left(\operatorname{cone}(C)\right)^*$.

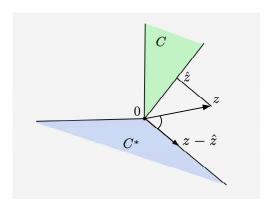


Figure 2.2.2. Illustration of the proof of the Polar Cone Theorem for the case where C is a closed convex cone. It is easily seen that $C \subset (C^*)^*$. To prove that $(C^*)^* \subset C$, we show that for any $z \in \Re^n$ and its projection on C, call it \hat{z} , we have $z - \hat{z} \in C^*$, so if $z \in (C^*)^*$, the geometry shown in the figure [(angle between z and $z - \hat{z}$) $< \pi/2$] is impossible, and we must have $z - \hat{z} = 0$, i.e., $z \in C$.

(b) We first show the result for the case where C is closed and convex. Indeed, if this is so, then for any $x \in C$, we have $x'y \leq 0$ for all $y \in C^*$, which implies that $x \in (C^*)^*$. Hence, $C \subset (C^*)^*$.

To prove that $(C^*)^* \subset C$, consider any $z \in \Re^n$, and its projection on C, denoted \hat{z} (see Fig. 2.2.2). By the Projection Theorem (Prop. 1.1.9),

$$(z - \hat{z})'(x - \hat{z}) \le 0, \quad \forall \ x \in C.$$

By taking in this relation x=0 and $x=2\hat{z}$ (which belong to C since C is a closed cone), it is seen that

$$(z - \hat{z})'\hat{z} = 0.$$

Combining the last two relations, we obtain $(z - \hat{z})'x \leq 0$ for all $x \in C$. Therefore, $(z - \hat{z}) \in C^*$. Hence, if $z \in (C^*)^*$, we have $(z - \hat{z})'z \leq 0$, which when added to $-(z - \hat{z})'\hat{z} = 0$ yields $||z - \hat{z}||^2 \leq 0$. It follows that $z = \hat{z}$ and $z \in C$.

Using the result just shown for the case where ${\cal C}$ is closed and convex, it follows that

$$\left(\left(\operatorname{cl}(\operatorname{conv}(C))\right)^*\right)^* = \operatorname{cl}(\operatorname{conv}(C)).$$

By using part (a), we have

$$C^* = (\operatorname{conv}(C))^* = (\operatorname{cl}(\operatorname{conv}(C)))^*.$$

By combining the above two relations, we obtain $(C^*)^* = \text{cl}(\text{conv}(C))$. **Q.E.D.**

Note that for any nonempty set C, from part (a) of Prop. 2.2.1, we have $C^* = (\operatorname{cl}(\operatorname{cone}(C)))^*$, so from part (b), $(C^*)^* = \operatorname{cl}(\operatorname{cone}(C))$.

2.3 POLYHEDRAL SETS AND FUNCTIONS

One of the distinctive characteristics of polyhedral sets is that they can be described by a finite number of vectors and scalars. In this section, we focus on several alternative finite characterizations of polyhedral sets, starting with polyhedral cones and their polar cones.

2.3.1 Polyhedral Cones and Farkas' Lemma

We will introduce two different ways to view polyhedral cones, and use polarity to show that these two views are essentially equivalent. We recall that a polyhedral cone $C \subset \Re^n$ has the form

$$C = \{x \mid a_j' x \le 0, j = 1, \dots, r\},\$$

where a_1, \ldots, a_r are some vectors in \Re^n , and r is a positive integer; see Fig. 2.3.1(a).

We say that a cone $C \subset \mathbb{R}^n$ is finitely generated, if it is generated by a finite set of vectors, i.e., if it has the form

$$C = \operatorname{cone}(\{a_1, \dots, a_r\}) = \left\{ x \mid x = \sum_{j=1}^r \mu_j a_j, \, \mu_j \ge 0, \, j = 1, \dots, r \right\},$$

where a_1, \ldots, a_r are some vectors in \Re^n , and r is a positive integer; see Fig. 2.3.1(b).

It turns out that polyhedral and finitely generated cones are connected by a polarity relation, as shown in the following proposition.

Proposition 2.3.1: (Farkas' Lemma) Let a_1, \ldots, a_r be vectors in \Re^n . Then, $\{x \mid a'_j x \leq 0, j = 1, \ldots, r\}$ and cone $(\{a_1, \ldots, a_r\})$ are closed cones that are polar to each other.

Proof: The cone $\{x \mid a_j'x \leq 0, j=1,\ldots,r\}$ is closed, being the intersection of closed halfspaces. To see that cone $(\{a_1,\ldots,a_r\})$ is closed, note that it is the image of the positive orthant $\{\mu \mid \mu \geq 0\}$ under the linear transformation that maps μ to $\sum_{j=1}^r \mu_j a_j$, and that the image of any polyhedral set under a linear transformation is a closed set (see the discussion following the proof of Prop. 1.4.13).

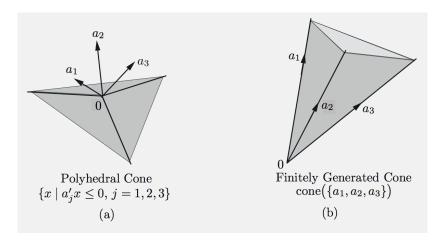


Figure 2.3.1. (a) Polyhedral cone defined by the inequality constraints $a'_j x \leq 0$, j = 1, 2, 3. (b) Cone generated by the vectors a_1, a_2, a_3 .

To show that the two cones are polar to each other, we use Prop. 2.2.1(a) to write

$$(\{a_1,\ldots,a_r\})^* = (\operatorname{cone}(\{a_1,\ldots,a_r\}))^*,$$

and we note that by the definition of polarity,

$$(\{a_1,\ldots,a_r\})^* = \{x \mid a'_j x \le 0, j = 1,\ldots,r\}.$$

Q.E.D.

Farkas' Lemma is often stated in an alternative (and equivalent) form that involves polyhedra specified by equality constraints. In particular, if $x, e_1, \ldots, e_m, a_1, \ldots, a_r$ are vectors in \Re^n , we have $x'y \leq 0$ for all $y \in \Re^n$ such that

$$y'e_i = 0, \quad \forall \ i = 1, ..., m, \qquad y'a_j \le 0, \quad \forall \ j = 1, ..., r,$$

if and only if x can be expressed as

$$x = \sum_{i=1}^{m} \lambda_i e_i + \sum_{j=1}^{r} \mu_j a_j,$$

where λ_i and μ_j are some scalars with $\mu_j \geq 0$ for all j.

For the proof, define $a_{r+i} = e_i$ and $a_{r+m+i} = -e_i$, i = 1, ..., m. The result can be stated as $P^* = C$, where

$$P = \{y \mid y'a_j \le 0, j = 1, \dots, r + 2m\}, \qquad C = \operatorname{cone}(\{a_1, \dots, a_{r+2m}\}).$$

Since by Prop. 2.3.1, $P = C^*$, and C is closed and convex, we have $P^* = (C^*)^* = C$ by the Polar Cone Theorem [Prop. 2.2.1(b)].

2.3.2 Structure of Polyhedral Sets

We will now prove a major theorem in polyhedral convexity, which will lead to a useful representation of polyhedral sets. The proof is visually intuitive and it is original to our knowledge. Alternative proofs can be found in [BNO03] and [Roc70].

Proposition 2.3.2: (Minkowski-Weyl Theorem) A cone is polyhedral if and only if it is finitely generated.

Proof: We first show that a finitely generated cone is polyhedral. Consider the cone generated by vectors a_1, \ldots, a_r in \Re^n , and assume that these vectors span \Re^n , i.e., that n of these vectors are linearly independent. We will show that cone $(\{a_1, \ldots, a_r\})$ is the intersection of halfspaces corresponding to hyperplanes/subspaces spanned by (n-1) linearly independent vectors from $\{a_1, \ldots, a_r\}$. Since there is a finite number of such halfspaces, it will follow that cone $(\{a_1, \ldots, a_r\})$ is polyhedral.

Indeed, if cone $(\{a_1, \ldots, a_r\}) = \Re^n$, it is polyhedral, and we are done, so assume otherwise. A vector b that does not belong to cone $(\{a_1, \ldots, a_r\})$ can be strictly separated from it by a hyperplane [cf. Prop. 1.5.3 under condition (2)], so there exists $\xi \in \Re^n$ and $\gamma \in \Re$ such that $b'\xi > \gamma > x'\xi$ for all $x \in \text{cone}(\{a_1, \ldots, a_r\})$. Since $\text{cone}(\{a_1, \ldots, a_r\})$ is a cone, each of its vectors x must satisfy $0 \ge x'\xi$ (otherwise $x'\xi$ could be increased without bound), and since $0 \in \text{cone}(\{a_1, \ldots, a_r\})$, we must have $\gamma > 0$. Thus,

$$b'\xi > \gamma > 0 \ge x'\xi$$
, $\forall x \in \text{cone}(\{a_1, \dots, a_r\})$.

It follows that the set

$$P_b = \{ y \mid b'y \ge 1, a'_i y \le 0, j = 1, \dots, r \}$$

is nonempty (it contains ξ/γ – see Fig. 2.3.2). Since $\{a_1, \ldots, a_r\}$ contains n linearly independent vectors, by Prop. 2.1.5, P_b has at least one extreme point, denoted \overline{y} . By Prop. 2.1.4(a), there are two possibilities:

- (1) We have $b'\overline{y} = 1$ and the set $\{a_j \mid a'_j \overline{y} = 0\}$ contains exactly n 1 linearly independent vectors.
- (2) The set $\{a_j \mid a'_j \overline{y} = 0\}$ contains exactly n linearly independent vectors.

The second alternative is impossible since \overline{y} is nonzero in view of $b'\overline{y} \ge 1$, while the first alternative implies that the hyperplane $\{x \mid \overline{y}'x = 0\}$, which is an (n-1)-dimensional subspace, is spanned by n-1 linearly independent vectors from $\{a_1, \ldots, a_r\}$, separates b and cone $(\{a_1, \ldots, a_r\})$,

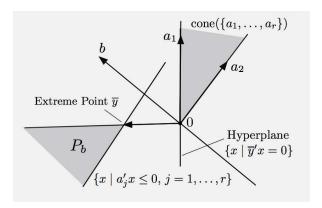


Figure 2.3.2. The polyhedral set

$$P_b = \{ y \mid b'y \ge 1, \ a'_i y \le 0, \ j = 1, \dots, r \},\$$

corresponding to a vector $b \notin \text{cone}(\{a_1, \ldots, a_r\})$, in the proof of the Minkowski-Weyl Theorem (cf. Prop. 2.3.2). It has a nonzero extreme point \overline{y} , which defines a halfspace $\{x \mid \overline{y}'x \leq 0\}$ that contains $\text{cone}(\{a_1, \ldots, a_r\})$. These halfspaces are finite in number and their intersection defines $\text{cone}(\{a_1, \ldots, a_r\})$.

and does not contain b. By letting b range over all vectors that are not in $\operatorname{cone}(\{a_1,\ldots,a_r\})$, it follows that $\operatorname{cone}(\{a_1,\ldots,a_r\})$ is the intersection of some halfspaces corresponding to hyperplanes/subspaces spanned by (n-1) linearly independent vectors from $\{a_1,\ldots,a_r\}$. Since there are only a finite number of such subspaces, $\operatorname{cone}(\{a_1,\ldots,a_r\})$ is polyhedral.

In the case where $\{a_1, \ldots, a_r\}$ does not contain n linearly independent vectors, let S be the subspace spanned by a_1, \ldots, a_r . The vectors defining the finitely generated cone

$$S^{\perp} + \operatorname{cone}(\{a_1, \ldots, a_r\})$$

contain a linearly independent set, so by the result proved so far, this cone is a polyhedral set. By writing

$$\operatorname{cone}(\{a_1,\ldots,a_r\}) = S \cap (S^{\perp} + \operatorname{cone}(\{a_1,\ldots,a_r\})),$$

we see that $cone(\{a_1,\ldots,a_r\})$ is the intersection of two polyhedral sets, so it is polyhedral.

Conversely, let $C = \{x \mid a_j'x \leq 0, j = 1, \ldots, r\}$ be a polyhedral cone. We will show that C is finitely generated. Indeed, by Farkas' Lemma, C^* is equal to cone $(\{a_1, \ldots, a_r\})$, which is polyhedral, as we have just shown. Thus, $C^* = \{x \mid c_j'x \leq 0, j = 1, \ldots, \bar{r}\}$ for some vectors c_j . Applying Farkas' Lemma again, it follows that $C = \text{cone}(\{c_1, \ldots, c_{\bar{r}}\})$. Q.E.D.

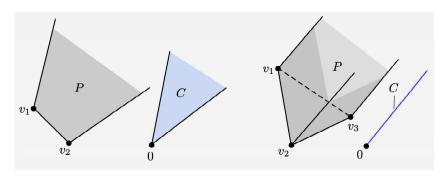


Figure 2.3.3. Examples of a Minkowski-Weyl representation of a two-dimensional and a three-dimensional polyhedral set P. It has the form

$$P = \operatorname{conv}(\{v_1, \dots, v_m\}) + C,$$

where v_1, \ldots, v_m are vectors and C is a finitely generated cone.

We now establish a fundamental result, showing that a polyhedral set can be represented as the sum of a finitely generated cone and the convex hull of a finite set of points (see Fig. 2.3.3).

Proposition 2.3.3: (Minkowski-Weyl Representation) A set P is polyhedral if and only if there is a nonempty finite set $\{v_1, \ldots, v_m\}$ and a finitely generated cone C such that $P = \text{conv}(\{v_1, \ldots, v_m\}) + C$, i.e.,

$$P = \left\{ x \mid x = \sum_{j=1}^{m} \mu_j v_j + y, \sum_{j=1}^{m} \mu_j = 1, \, \mu_j \ge 0, \, j = 1, \dots, m, \, y \in C \right\}.$$

Proof: Assume that P is polyhedral. Then, it has the form

$$P = \{x \mid a'_j x \le b_j, j = 1, \dots, r\},\$$

for some $a_j \in \mathbb{R}^n$, $b_j \in \mathbb{R}$, j = 1, ..., r. Consider the polyhedral cone in \mathbb{R}^{n+1} given by

$$\hat{P} = \{(x, w) \mid 0 \le w, a'_j x \le b_j w, j = 1, \dots, r\}$$

(see Fig. 2.3.4), and note that

$$P = \{x \mid (x, 1) \in \hat{P}\}.$$

By the Minkowski-Weyl Theorem (Prop. 2.3.2), \hat{P} is finitely generated, so it has the form

$$\hat{P} = \left\{ (x, w) \mid x = \sum_{j=1}^{m} \nu_j \tilde{v}_j, \ w = \sum_{j=1}^{m} \nu_j \tilde{d}_j, \ \nu_j \ge 0, \ j = 1, \dots, m \right\},\,$$

for some $\tilde{v}_j \in \mathbb{R}^n$, $\tilde{d}_j \in \mathbb{R}$, j = 1, ..., r. Since $w \geq 0$ for all vectors $(x, w) \in \hat{P}$, we see that $\tilde{d}_j \geq 0$ for all j. Let

$$J^+ = \{j \mid \tilde{d}_j > 0\}, \qquad J^0 = \{j \mid \tilde{d}_j = 0\}.$$

By replacing, for all $j \in J^+$, $\nu_j \tilde{d}_j$ with μ_j and $\tilde{v}_j / \tilde{d}_j$ with v_j , and for all $j \in J^0$, ν_j with μ_j and \tilde{v}_j with v_j , we obtain the equivalent description

$$\hat{P} = \left\{ (x, w) \mid x = \sum_{j \in J^+ \cup J^0} \mu_j v_j, \ w = \sum_{j \in J^+} \mu_j, \ \mu_j \ge 0, \ j \in J^+ \cup J^0 \right\}. \tag{2.1}$$

Since $P = \{x \mid (x, 1) \in \hat{P}\}$, we obtain

$$P = \left\{ x \mid x = \sum_{j \in J^+ \cup J^0} \mu_j v_j, \sum_{j \in J^+} \mu_j = 1, \, \mu_j \ge 0, \, j \in J^+ \cup J^0 \right\}. \quad (2.2)$$

Thus, P is the vector sum of $\operatorname{conv}(\{v_j \mid j \in J^+\})$ and the finitely generated cone

$$\left\{ \sum_{j \in J^0} \mu_j v_j \mid \mu_j \ge 0, j \in J^0 \right\}.$$

To prove that the vector sum of $\operatorname{conv}(\{v_1,\ldots,v_m\})$ and a finitely generated cone is a polyhedral set, we reverse the preceding argument: starting from Eq. (2.2), we express P as $\{x \mid (x,1) \in \hat{P}\}$, where \hat{P} is the finitely generated cone of Eq. (2.1). We then use the Minkowski-Weyl Theorem to assert that this cone is polyhedral, and we finally construct a polyhedral set description using the equation $P = \{x \mid (x,1) \in \hat{P}\}$. Q.E.D.

As indicated by the examples of Fig. 2.3.3, the finitely generated cone C in a Minkowski-Weyl representation of a polyhedral set P is just the recession cone of P. Let us now use the Minkowski-Weyl representation to show that the major algebraic operations on polyhedral sets preserve their polyhedral character.

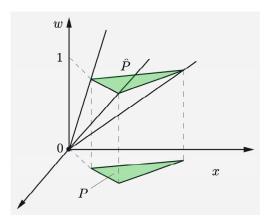


Figure 2.3.4. Illustration of the proof of the Minkowski-Weyl representation. The cone of \Re^{n+1}

$$\hat{P} = \{(x, w) \mid 0 \le w, a'_j x \le b_j w, j = 1, \dots, r\},\$$

is derived from the polyhedral set

$$P = \{x \mid a'_{j}x \le b_{j}, j = 1, \dots, r\}.$$

Proposition 2.3.4: (Algebraic Operations on Polyhedral Sets)

- (a) The intersection of polyhedral sets is polyhedral, if it is nonempty.
- (b) The Cartesian product of polyhedral sets is polyhedral.
- (c) The image of a polyhedral set under a linear transformation is a polyhedral set.
- (d) The vector sum of two polyhedral sets is polyhedral.
- (e) The inverse image of a polyhedral set under a linear transformation is polyhedral.

Proof: Parts (a) and (b) are evident from the definition of a polyhedral set. To show part (c), let the polyhedral set P be represented as

$$P = \text{conv}(\{v_1, \dots, v_m\}) + \text{cone}(\{a_1, \dots, a_r\}),$$

and let A be a matrix. We have

$$AP = \operatorname{conv}(\{Av_1, \dots, Av_m\}) + \operatorname{cone}(\{Aa_1, \dots, Aa_r\}).$$

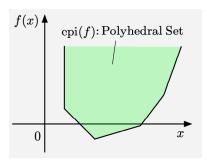


Figure 2.3.5. Illustration of a polyhedral function. By definition, the function must be proper, and its epigraph must be a polyhedral set.

It follows that AP has a Minkowski-Weyl representation, and by Prop. 2.3.3, it is polyhedral. Part (d) follows from part (c) since $P_1 + P_2$ can be viewed as the image of the polyhedral set $P_1 \times P_2$ under the linear transformation $(x_1, x_2) \mapsto (x_1 + x_2)$.

To prove part (e) note that the inverse image of a polyhedral set

$$\{y \mid a_i'y \le b_j, j = 1, \dots, r\}$$

under the linear transformation A is the set

$$\{x \mid a_j' Ax \leq b_j, j = 1, \dots, r\},\$$

which is clearly polyhedral. Q.E.D.

2.3.3 Polyhedral Functions

Polyhedral sets can also be used to define functions with polyhedral structure. In particular, we say that a function $f: \Re^n \mapsto (-\infty, \infty]$ is polyhedral if its epigraph is a polyhedral set in \Re^{n+1} ; see Fig. 2.3.5. Note that a polyhedral function f is, by definition, closed, convex, and also proper [since f cannot take the value $-\infty$, and $\operatorname{epi}(f)$ is closed, convex, and nonempty (based on our convention that only nonempty sets can be polyhedral)]. The following proposition provides a useful representation of polyhedral functions.

Proposition 2.3.5: Let $f: \Re^n \mapsto (-\infty, \infty]$ be a convex function. Then f is polyhedral if and only if dom(f) is a polyhedral set and

$$f(x) = \max_{j=1,\dots,m} \{a'_j x + b_j\}, \qquad \forall \ x \in \text{dom}(f),$$

where a_i are vectors in \Re^n , b_i are scalars, and m is a positive integer.

Proof: If f has the representation given above, then epi(f) is written as

$$epi(f) = \{(x, w) \mid x \in dom(f)\} \cap \{(x, w) \mid a'_j x + b_j \le w, j = 1, \dots, m\}.$$

Since the two sets in the right-hand side are polyhedral, their intersection, epi(f), [which is nonempty since dom(f) is polyhedral and hence nonempty] is also polyhedral. Hence f is polyhedral.

Conversely, if f is polyhedral, its epigraph is a polyhedral set, and can be represented as

$$\{(x,w) \mid a'_j x + b_j \le c_j w, j = 1, \dots, r\},\$$

where a_j are some vectors in \Re^n , and b_j and c_j are some scalars. Since for any $(x, w) \in \operatorname{epi}(f)$, we have $(x, w + \gamma) \in \operatorname{epi}(f)$ for all $\gamma \geq 0$, it follows that $c_j \geq 0$, so by normalizing if necessary, we may assume without loss of generality that either $c_j = 0$ or $c_j = 1$. If $c_j = 0$ for all j, then f would not be proper, contradicting the fact that a polyhedral function is proper. Hence we must have, for some m with $1 \leq m \leq r$, $c_j = 1$ for $j = 1, \ldots, m$, and $c_j = 0$ for $j = m + 1, \ldots, r$, i.e.,

$$epi(f) = \{(x, w) \mid a_i'x + b_j \le w, \ j = 1, \dots, m, \ a_i'x + b_j \le 0, \ j = m+1, \dots, r\}.$$

Thus the effective domain of f is the polyhedral set

$$dom(f) = \{x \mid a'_j x + b_j \le 0, j = m + 1, \dots, r\},\$$

and we have

$$f(x) = \max_{j=1,\dots,m} \{a'_j x + b_j\}, \qquad \forall \ x \in \text{dom}(f).$$

Q.E.D.

Some common operations on polyhedral functions, such as sum and linear composition preserve their polyhedral character as shown by the following two propositions.

Proposition 2.3.6: The sum of two polyhedral functions f_1 and f_2 , such that $dom(f_1) \cap dom(f_2) \neq \emptyset$, is a polyhedral function.

Proof: By Prop. 2.3.5, $dom(f_1)$ and $dom(f_2)$ are polyhedral sets in \Re^n , and

$$f_1(x) = \max\{a'_1 x + b_1, \dots, a'_m x + b_m\}, \quad \forall x \in \text{dom}(f_1),$$

$$f_2(x) = \max\{\overline{a}_1'x + \overline{b}_1, \dots, \overline{a}_{\overline{m}}'x + \overline{b}_{\overline{m}}\}, \quad \forall \ x \in \text{dom}(f_2),$$

where a_i and \overline{a}_i are vectors in \Re^n , and b_i and \overline{b}_i are scalars. The domain of $f_1 + f_2$ is $\text{dom}(f_1) \cap \text{dom}(f_2)$, which is polyhedral since $\text{dom}(f_1)$ and $\text{dom}(f_2)$ are polyhedral. Furthermore, for all $x \in \text{dom}(f_1 + f_2)$,

$$f_1(x) + f_2(x) = \max \left\{ a'_1 x + b_1, \dots, a'_m x + b_m \right\}$$

$$+ \max \left\{ \overline{a}'_1 x + \overline{b}_1, \dots, \overline{a}'_m x + \overline{b}_{\overline{m}} \right\}$$

$$= \max_{1 \le i \le m, \ 1 \le j \le \overline{m}} \left\{ a'_i x + b_i + \overline{a}'_j x + \overline{b}_j \right\}$$

$$= \max_{1 \le i \le m, \ 1 \le j \le \overline{m}} \left\{ (a_i + \overline{a}_j)' x + (b_i + \overline{b}_j) \right\}.$$

Therefore, by Prop. 2.3.5, $f_1 + f_2$ is polyhedral. **Q.E.D.**

Proposition 2.3.7: If A is a matrix and g is a polyhedral function such that dom(g) contains a point in the range of A, the function f given by f(x) = g(Ax) is polyhedral.

Proof: Since $g: \Re^m \mapsto (-\infty, \infty]$ is polyhedral, by Prop. 2.3.5, dom(g) is a polyhedral set in \Re^m and g is given by

$$g(y) = \max\{a'_1 y + b_1, \dots, a'_m y + b_m\}, \quad \forall y \in \text{dom}(g),$$

for some vectors a_i in \Re^m and scalars b_i . We have

$$dom(f) = \{x \mid f(x) < \infty\} = \{x \mid g(Ax) < \infty\} = \{x \mid Ax \in dom(g)\}.$$

Thus, dom(f) is the inverse image of the polyhedral set dom(g) under the linear transformation A. By the assumption that dom(g) contains a point in the range of A, it follows that dom(f) is nonempty, while dom(f) is polyhedral. Furthermore, for all $x \in dom(f)$, we have

$$f(x) = g(Ax)$$

$$= \max\{a'_1 Ax + b_1, \dots, a'_m Ax + b_m\}$$

$$= \max\{(A'a_1)'x + b_1, \dots, (A'a_m)'x + b_m\}.$$

Thus, by Prop. 2.3.5, the function f is polyhedral. **Q.E.D.**

2.4 POLYHEDRAL ASPECTS OF OPTIMIZATION

Polyhedral convexity plays a very important role in optimization. One reason is that many practical problems can be readily formulated in terms of

polyhedral sets and functions. Another reason is that for polyhedral constraint sets and/or linear cost functions, it is often possible to show stronger optimization results than those available for general convex constraint sets and/or general cost functions. We have seen a few such instances so far. In particular:

- (1) A linear (or, more generally, concave) function that is not constant over a convex constraint set C, can only attain its minimum at a relative boundary point of C (Prop. 1.3.4).
- (2) A linear (or, more generally, convex quadratic) function that is bounded below over a nonempty polyhedral set C, attains a minimum over C (Prop. 1.4.12).

In this section, we explore some further consequences of polyhedral convexity in optimization, with an emphasis on linear programming (the minimization of a linear function over a polyhedral set). In particular, we prove one of the fundamental linear programming results: if a linear function f attains a minimum over a polyhedral set C that has at least one extreme point, then f attains a minimum at some extreme point of C (as well as possibly at some other nonextreme points). This is a special case of the following more general result, which holds when f is concave, and C is closed and convex.

Proposition 2.4.1: Let C be a closed convex subset of \mathbb{R}^n that has at least one extreme point. A concave function $f: C \mapsto \mathbb{R}$ that attains a minimum over C attains the minimum at some extreme point of C.

Proof: Let x^* minimize f over C. If $x^* \in \text{ri}(C)$ [see Fig. 2.4.1(a)], by Prop. 1.3.4, f must be constant over C, so it attains a minimum at an extreme point of C (since C has at least one extreme point by assumption). If $x^* \notin \text{ri}(C)$, then by Prop. 1.5.5, there exists a hyperplane H_1 properly separating x^* and C. Since $x^* \in C$, H_1 must contain x^* , so by the proper separation property, H_1 cannot contain C, and it follows that the intersection $C \cap H_1$ has dimension smaller than the dimension of C.

If $x^* \in \operatorname{ri}(C \cap H_1)$ [see Fig. 2.4.1(b)], then f must be constant over $C \cap H_1$, so it attains a minimum at an extreme point of $C \cap H_1$ (since C contains an extreme point, it does not contain a line by Prop. 2.1.2, and hence $C \cap H_1$ does not contain a line, which implies that $C \cap H_1$ has an extreme point). By Prop. 2.1.1, this optimal extreme point is also an extreme point of C. If $x^* \notin \operatorname{ri}(C \cap H_1)$, there exists a hyperplane H_2 properly separating x^* and $C \cap H_1$. Again, since $x^* \in C \cap H_1$, H_2 contains x^* , so it cannot contain $C \cap H_1$, and it follows that the intersection $C \cap H_1 \cap H_2$ has dimension smaller than the dimension of $C \cap H_1$.

If $x^* \in ri(C \cap H_1 \cap H_2)$ [see Fig. 2.4.1(c)], then f must be constant

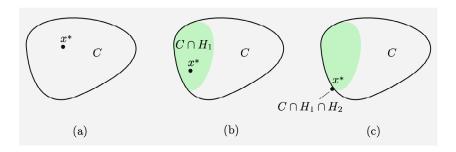


Figure 2.4.1. An outline of the argument of the proof of Prop. 2.4.1 for a 3-dimensional set C.

over $C \cap H_1 \cap H_2$, etc. Since with each new hyperplane, the dimension of the intersection of C with the generated hyperplanes is reduced, this process will be repeated at most n times, until x^* is a relative interior point of some set $C \cap H_1 \cap \cdots \cap H_k$, at which time an extreme point of $C \cap H_1 \cap \cdots \cap H_k$ will be obtained. Through a reverse argument, repeatedly applying Prop. 2.1.1, it follows that this extreme point is an extreme point of C. Q.E.D.

We now specialize the preceding result to the case of a linear program, where the cost function f is linear.

Proposition 2.4.2: (Fundamental Theorem of Linear Programming) Let P be a polyhedral set that has at least one extreme point. A linear function that is bounded below over P attains a minimum at some extreme point of P.

Proof: Since the cost function is bounded below over P, it attains a minimum (Prop. 1.4.12). The result now follows from Prop. 2.4.1. **Q.E.D.**

Figure 2.4.2 illustrates the possibilities for a linear programming problem. There are two cases:

- (a) The constraint set P contains an extreme point (equivalently, by Prop. 2.1.2, P does not contain a line). In this case, the linear cost function is either unbounded below over P or else it attains a minimum at an extreme point of P. For an example, let $P = [0, \infty)$. Then, the minimum of $1 \cdot x$ over P is attained at the extreme point 0, the minimum of $0 \cdot x$ over P is attained at every point of P, including the extreme point 0, and the minimum of $-1 \cdot x$ over P is not attained while the cost is unbounded below.
- (b) The constraint set P does not contain an extreme point (equivalently,

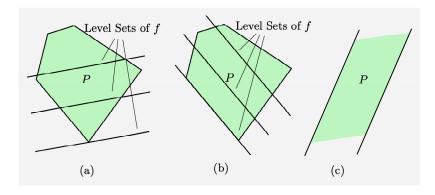


Figure 2.4.2. Illustration of the fundamental theorem of linear programming. In (a) and (b), the constraint set P has at least one extreme point and the linear cost function f is bounded over P. Then f either attains a minimum at a unique extreme point as in (a), or it attains a minimum at one or more extreme points as well as at an infinite number of nonextreme points as in (b). In (c), the constraint set has no extreme points because it contains a line (cf. Prop. 2.1.2), and the linear cost function is either unbounded below over P or attains a minimum at an infinite set of (nonextreme) points whose lineality space is equal to the lineality space of P.

by Prop. 2.1.2, P contains a line). In this case, L_P , the lineality space of P, is a subspace of dimension greater or equal to one. If the linear cost function f is bounded below over P, it attains a minimum by Prop. 1.4.12. However, because every direction in L_P must be a direction along which f is constant (otherwise f would be unbounded below over P), the set of minima is a polyhedral set whose lineality space is equal to L_P . Therefore, the set of minima must be unbounded. For an example, let $P = \Re$. Then, the minimum of $1 \cdot x$ over P is not attained and the cost is unbounded below, while the set of minima of $0 \cdot x$ over P is P.

The theory of polyhedral sets and their extreme points can be used for the development of linear programming algorithms, such as the simplex method and other related methods. These algorithms are outside our scope, so we refer to standard textbooks, such as Dantzig [Dan63], Chvatal [Chv83], and Bertsimas and Tsitsiklis [BeT97].

Basic Concepts of Convex Optimization

Contents				
 3.1. Constrained Optimization 3.2. Existence of Optimal Solutions 3.3. Partial Minimization of Convex Functions 3.4. Saddle Point and Minimax Theory 				p. 118 p. 122

In this chapter, we introduce some basic concepts of convex optimization and minimax theory, with a special focus on the question of existence of optimal solutions.

3.1 CONSTRAINED OPTIMIZATION

Consider the problem

minimize
$$f(x)$$

subject to $x \in X$,

where $f: \Re^n \mapsto (-\infty, \infty]$ is a function and X is a nonempty subset of \Re^n . Any vector $x \in X \cap \text{dom}(f)$ is said to be a *feasible solution* of the problem (we also use the terms *feasible vector* or *feasible point*). If there is at least one feasible solution, i.e., $X \cap \text{dom}(f) \neq \emptyset$, we say that the problem is *feasible*; otherwise we say that the problem is *infeasible*. Thus, when f is extended real-valued, we view only the points in $X \cap \text{dom}(f)$ as candidates for optimality, and we view dom(f) as an implicit constraint set. Furthermore, feasibility of the problem is equivalent to $\inf_{x \in X} f(x) < \infty$.

We say that a vector x^* is a minimum of f over X if

$$x^* \in X \cap \text{dom}(f)$$
, and $f(x^*) = \inf_{x \in X} f(x)$.

We also call x^* a minimizing point or minimizer or global minimum of f over X. Alternatively, we say that f attains a minimum over X at x^* , and we indicate this by writing

$$x^* \in \arg\min_{x \in X} f(x).$$

If x^* is known to be the unique minimizer of f over X, with slight abuse of notation, we also write

$$x^* = \arg\min_{x \in X} f(x).$$

We use similar terminology for maxima, i.e., a vector $x^* \in X$ such that $f(x^*) = \sup_{x \in X} f(x)$ is said to be a maximum of f over X if x^* is a minimum of (-f) over X, and we indicate this by writing

$$x^* \in \arg\max_{x \in X} f(x).$$

If $X = \Re^n$ or if the domain of f is the set X (instead of \Re^n), we also call x^* a (global) minimum or (global) maximum of f (without the qualifier "over X").

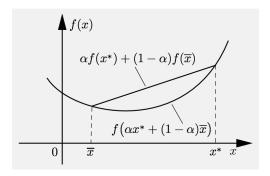


Figure 3.1.1. Illustration of why local minima of convex functions are also global (cf. Prop. 3.1.1). Given x^* and \overline{x} with $f(\overline{x}) < f(x^*)$, every point of the form

$$x_{\alpha} = \alpha x^* + (1 - \alpha)\overline{x}, \qquad \alpha \in (0, 1),$$

satisfies $f(x_{\alpha}) < f(x^*)$. Thus x^* cannot be a local minimum that is not global.

Local Minima

Often in optimization problems we have to deal with a weaker form of minimum, one that is optimum only when compared with points that are "nearby." In particular, given a subset X of \Re^n and a function $f: \Re^n \mapsto (-\infty, \infty]$, we say that a vector x^* is a local minimum of f over X if $x^* \in X \cap \text{dom}(f)$ and there exists some $\epsilon > 0$ such that

$$f(x^*) \le f(x), \quad \forall x \in X \text{ with } ||x - x^*|| < \epsilon.$$

If $X = \Re^n$ or the domain of f is the set X (instead of \Re^n), we also call x^* a local minimum of f (without the qualifier "over X"). A local minimum x^* is said to be *strict* if there is no other local minimum within some open sphere centered at x^* . Local maxima are defined similarly.

In practical applications we are typically interested in global minima, yet we have to contend with local minima because of the inability of many optimality conditions and algorithms to distinguish between the two types of minima. This can be a major practical difficulty, but an important implication of convexity of f and X is that all local minima are also global, as shown in the following proposition and in Fig. 3.1.1.

Proposition 3.1.1: If X is a convex subset of \Re^n and $f: \Re^n \mapsto (-\infty, \infty]$ is a convex function, then a local minimum of f over X is also a global minimum. If in addition f is strictly convex, then there exists at most one global minimum of f over X.

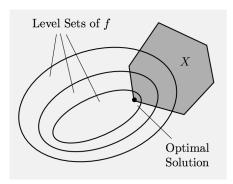


Figure 3.2.1. View of the set of optimal solutions of the problem

minimize
$$f(x)$$

subject to $x \in X$,

as the intersection of all the nonempty level sets of the form

$$\{x \in X \mid f(x) \le \gamma\}, \quad \gamma \in \Re.$$

Proof: Let f be convex, and assume to arrive at a contradiction, that x^* is a local minimum of f over X that is not global (see Fig. 3.1.1). Then, there must exist an $\overline{x} \in X$ such that $f(\overline{x}) < f(x^*)$. By convexity, for all $\alpha \in (0,1)$,

$$f(\alpha x^* + (1 - \alpha)\overline{x}) \le \alpha f(x^*) + (1 - \alpha)f(\overline{x}) < f(x^*).$$

Thus, f has strictly lower value than $f(x^*)$ at every point on the line segment connecting x^* with \overline{x} , except at x^* . Since X is convex, the line segment belongs to X, thereby contradicting the local minimality of x^* .

Let f be strictly convex, let x^* be a global minimum of f over X, and let x be a point in X with $x \neq x^*$. Then the midpoint $y = (x+x^*)/2$ belongs to X since X is convex, and by strict convexity, $f(y) < 1/2(f(x) + f(x^*))$, while by the optimality of x^* , we have $f(x^*) \leq f(y)$. These two relations imply that $f(x^*) < f(x)$, so x^* is the unique global minimum. **Q.E.D.**

3.2 EXISTENCE OF OPTIMAL SOLUTIONS

A basic question in optimization problems is whether an optimal solution exists. It can be seen that the set of minima of a real-valued function f over a nonempty set X, call it X^* , is equal to the intersection of X and the level sets of f that have a common point with X:

$$X^* = \cap_{k=0}^{\infty} \{ x \in X \mid f(x) \le \gamma_k \},\,$$

where $\{\gamma_k\}$ is any scalar sequence with $\gamma_k \downarrow \inf_{x \in X} f(x)$ (see Fig. 3.2.1).

From this characterization of X^* , it follows that the set of minima is nonempty and compact if the sets

$$\{x \in X \mid f(x) \le \gamma\},\$$

are compact (since the intersection of a nested sequence of nonempty and compact sets is nonempty and compact). This is the essence of the classical theorem of Weierstrass (Prop. A.2.7), which states that a continuous function attains a minimum over a compact set. We will provide a more general version of this theorem, and to this end, we introduce some terminology.

We say that a function $f: \Re^n \mapsto (-\infty, \infty]$ is coercive if for every sequence $\{x_k\}$ such that $\|x_k\| \to \infty$, we have $\lim_{k\to\infty} f(x_k) = \infty$. Note that as a consequence of the definition, if $\operatorname{dom}(f)$ is bounded, then f is coercive. Furthermore, all the nonempty level sets of a coercive function are bounded.

Proposition 3.2.1: (Weierstrass' Theorem) Consider a closed proper function $f: \mathbb{R}^n \mapsto (-\infty, \infty]$, and assume that any one of the following three conditions holds:

- (1) dom(f) is bounded.
- (2) There exists a scalar $\overline{\gamma}$ such that the level set

$$\left\{x\mid f(x)\leq\overline{\gamma}\right\}$$

is nonempty and bounded.

(3) f is coercive.

Then the set of minima of f over \Re^n is nonempty and compact.

Proof: It is sufficient to show that each of the three conditions implies that the nonempty level sets $V_{\gamma} = \{x \mid f(x) \leq \gamma\}$ are compact for all $\gamma \leq \overline{\gamma}$, where $\overline{\gamma}$ is such that $V_{\overline{\gamma}}$ is nonempty and compact, and then use the fact that the set of minima of f is the intersection of its nonempty level sets. (Note that f is assumed proper, so it has some nonempty level sets.) Since f is closed, its level sets are closed (cf. Prop. 1.1.2). It is evident that under each of the three conditions the level sets are also bounded for γ less or equal to some $\overline{\gamma}$, so they are compact. **Q.E.D.**

The most common application of Weierstrass' Theorem is when we want to minimize a real-valued function $f: \mathbb{R}^n \to \mathbb{R}$ over a nonempty set X. Then, by applying the proposition to the extended real-valued function

$$\tilde{f}(x) = \begin{cases} f(x) & \text{if } x \in X, \\ \infty & \text{otherwise,} \end{cases}$$

we see that the set of minima of f over X is nonempty and compact if X is closed, f is lower semicontinuous at each $x \in X$ (implying, by Prop. 1.1.3, that \tilde{f} is closed), and one of the following conditions holds:

- (1) X is bounded.
- (2) Some set $\{x \in X \mid f(x) \leq \overline{\gamma}\}$ is nonempty and bounded.
- (3) \tilde{f} is coercive, or equivalently, for every sequence $\{x_k\} \subset X$ such that $||x_k|| \to \infty$, we have $\lim_{k \to \infty} f(x_k) = \infty$.

The following is essentially Weierstrass' Theorem specialized to convex functions.

Proposition 3.2.2: Let X be a closed convex subset of \Re^n , and let $f: \Re^n \mapsto (-\infty, \infty]$ be a closed convex function with $X \cap \text{dom}(f) \neq \emptyset$. The set of minima of f over X is nonempty and compact if and only if X and f have no common nonzero direction of recession.

Proof: Let $f^* = \inf_{x \in X} f(x)$ and note that $f^* < \infty$ since $X \cap \text{dom}(f) \neq \emptyset$. Let $\{\gamma_k\}$ be a scalar sequence with $\gamma_k \downarrow f^*$, and consider the sets

$$V_k = \{ x \mid f(x) \le \gamma_k \}.$$

Then the set of minima of f over X is

$$X^* = \bigcap_{k=1}^{\infty} (X \cap V_k).$$

The sets $X \cap V_k$ are nonempty and have $R_X \cap R_f$ as their common recession cone, which is also the recession cone of X^* , when $X^* \neq \emptyset$ [cf. Props. 1.4.5, 1.4.2(c)]. It follows using Prop. 1.4.2(a) that X^* is nonempty and compact if and only if $R_X \cap R_f = \{0\}$. Q.E.D.

If X and f of the above proposition have a common direction of recession, then either the optimal solution set is empty [take for example, $X = \Re$ and $f(x) = e^x$] or else it is nonempty and unbounded [take for example, $X = \Re$ and $f(x) = \max\{0, x\}$]. Here is another result that addresses an important special case where the set of minima is compact.

Proposition 3.2.3: (Existence of Solution, Sum of Functions) Let $f_i: \Re^n \mapsto (-\infty, \infty], i = 1, \dots, m$, be closed proper convex functions such that the function $f = f_1 + \dots + f_m$ is proper. Assume that the recession function of a single function f_i satisfies $r_{f_i}(d) = \infty$ for all $d \neq 0$. Then the set of minima of f is nonempty and compact.

Proof: By Prop. 3.2.2, the set of minima of f is nonempty and compact if and only if $R_f = \{0\}$, which by Prop. 1.4.6, is true if and only if $r_f(d) > 0$ for all $d \neq 0$. The result now follows from Prop. 1.4.8. **Q.E.D.**

As an example of application of the preceding proposition, if one of the functions f_i is a positive definite quadratic function, the set of minima of the sum f is nonempty and compact. In fact in this case f has a unique minimum because the positive definite quadratic is strictly convex, which makes f strictly convex.

The next proposition addresses cases where the optimal solution set is not compact.

Proposition 3.2.4: (Existence of Solution, Noncompact Level Sets) Let X be a closed convex subset of \mathbb{R}^n , and let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a closed convex function with $X \cap \text{dom}(f) \neq \emptyset$. The set of minima of f over X, denoted X^* , is nonempty under any one of the following two conditions:

- (1) $R_X \cap R_f = L_X \cap L_f$.
- (2) $R_X \cap R_f \subset L_f$ and X is a polyhedral set.

Furthermore, under condition (1),

$$X^* = \tilde{X} + (L_X \cap L_f),$$

where \tilde{X} is some nonempty and compact set.

Proof: Let $f^* = \inf_{x \in X} f(x)$ and note that $f^* < \infty$ since $X \cap \text{dom}(f) \neq \emptyset$. Let $\{\gamma_k\}$ be a scalar sequence with $\gamma_k \downarrow f^*$, consider the level sets

$$V_k = \{ x \mid f(x) \le \gamma_k \},\,$$

and note that

$$X^* = \bigcap_{k=1}^{\infty} (X \cap V_k).$$

Let condition (1) hold. The sets $X \cap V_k$ are nonempty, closed, convex, and nested. Furthermore, they have the same recession cone, $R_X \cap R_f$, and the same lineality space $L_X \cap L_f$, while by assumption, $R_X \cap R_f = L_X \cap L_f$. By Prop. 1.4.11(a), it follows that X^* is nonempty and has the form

$$X^* = \tilde{X} + (L_X \cap L_f),$$

where \tilde{X} is some nonempty compact set.

Let condition (2) hold. The sets V_k are nested and $X \cap V_k$ is nonempty for all k. Furthermore, all the sets V_k have the same recession cone, R_f , and the same lineality space, L_f , while by assumption, $R_X \cap R_f \subset L_f$, and X is polyhedral and hence retractive (cf. Prop. 1.4.9). By Prop. 1.4.11(b), it follows that X^* is nonempty. **Q.E.D.** Note that in the special case $X = \Re^n$, conditions (1) and (2) of Prop. 3.2.4 coincide. Figure 3.2.3(b) provides a counterexample showing that if X is nonpolyhedral, the condition

$$R_X \cap R_f \subset L_f$$

is not sufficient to guarantee the existence of optimal solutions or even the finiteness of f^* . This counterexample also shows that the cost function may be bounded below and attain a minimum over any closed halfline contained in the constraint set, and yet it may not attain a minimum over the entire set. Recall, however, that in the special cases of linear and quadratic programming problems, boundedness from below of the cost function over the constraint set guarantees the existence of an optimal solution (cf. Prop. 1.4.12).

3.3 PARTIAL MINIMIZATION OF CONVEX FUNCTIONS

In our development of duality and minimax theory we will often encounter functions obtained by minimizing other functions partially, i.e., with respect to some of their variables. It is then useful to be able to deduce properties of the function obtained, such as convexity and closedness, from corresponding properties of the original.

There is an important geometric relation between the epigraph of a given function and the epigraph of its partially minimized version: except for some boundary points, the latter is obtained by projection from the former [see part (b) of the following proposition, and Fig. 3.3.1]. This is the key to understanding the properties of partially minimized functions.

Proposition 3.3.1: Consider a function $F: \Re^{n+m} \mapsto (-\infty, \infty]$ and the function $f: \Re^n \mapsto [-\infty, \infty]$ defined by

$$f(x) = \inf_{z \in \Re^m} F(x, z).$$

Then:

- (a) If F is convex, then f is also convex.
- (b) We have

$$P(\operatorname{epi}(F)) \subset \operatorname{epi}(f) \subset \operatorname{cl}(P(\operatorname{epi}(F))),$$
 (3.1)

where $P(\cdot)$ denotes projection on the space of (x, w), i.e., for any subset S of \Re^{n+m+1} , $P(S) = \{(x, w) \mid (x, z, w) \in S\}$.

Proof: (a) If $\operatorname{epi}(f) = \emptyset$, i.e., $f(x) = \infty$ for all $x \in \Re^n$, then $\operatorname{epi}(f)$ is convex, so f is convex. Assume that $\operatorname{epi}(f) \neq \emptyset$, and let $(\overline{x}, \overline{w})$ and (\tilde{x}, \tilde{w})

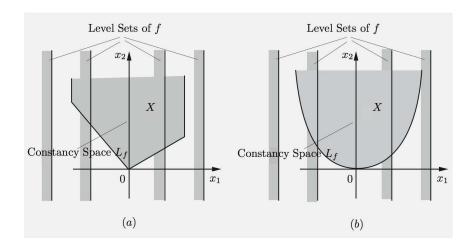


Figure 3.2.3. Illustration of the issues regarding existence of an optimal solution assuming $R_X \cap R_f \subset L_f$, i.e., that every common direction of recession of X and f is a direction in which f is constant [cf. Prop. 3.2.4 under condition (2)].

In both problems illustrated in (a) and (b) the cost function is

$$f(x_1, x_2) = e^{x_1}$$
.

In the problem of (a), the constraint set X is the polyhedral set shown in the figure, while in the problem of (b), X is specified by a quadratic inequality:

$$X = \{(x_1, x_2) \mid x_1^2 \le x_2\},\$$

as shown in the figure. In both cases we have

$$R_X = \{ (d_1, d_2) \mid d_1 = 0, \ d_2 \ge 0 \},\$$

$$R_f = \Big\{ (d_1, d_2) \mid d_1 \leq 0, \ d_2 \in \Re \Big\}, \qquad L_f = \Big\{ (d_1, d_2) \mid d_1 = 0, \ d_2 \in \Re \Big\},$$

so that $R_X \cap R_f \subset L_f$.

In the problem of (a) it can be seen that an optimal solution exists. In the problem of (b), however, we have $f(x_1,x_2)>0$ for all (x_1,x_2) , while for $x_1=-\sqrt{x_2}$ where $x_2\geq 0$, we have $(x_1,x_2)\in X$ with

$$\lim_{x_2 \to \infty} f(-\sqrt{x_2}, x_2) = \lim_{x_2 \to \infty} e^{-\sqrt{x_2}} = 0,$$

implying that $f^* = 0$. Thus f cannot attain the minimum value f^* over X. Note that f attains a minimum over the intersection of any line with X.

If in the problem of (b) the cost function were instead $f(x_1,x_2)=x_1$, we would still have $R_X\cap R_f\subset L_f$ and f would still attain a minimum over the intersection of any line with X, but it can be seen that $f^*=-\infty$. If the constraint set were instead $X=\left\{(x_1,x_2)\mid |x_1|\leq x_2\right\}$, which is polyhedral, we would still have $f^*=-\infty$, but then the condition $R_X\cap R_f\subset L_f$ would be violated.

be two points in $\operatorname{epi}(f)$. Then $f(\overline{x}) < \infty$, $f(\tilde{x}) < \infty$, and there exist sequences $\{\overline{z}_k\}$ and $\{\tilde{z}_k\}$ such that

$$F(\overline{x}, \overline{z}_k) \to f(\overline{x}), \qquad F(\tilde{x}, \tilde{z}_k) \to f(\tilde{x}).$$

Using the definition of f and the convexity of F, we have for all $\alpha \in [0,1]$ and k,

$$f(\alpha \overline{x} + (1 - \alpha)\tilde{x}) \le F(\alpha \overline{x} + (1 - \alpha)\tilde{x}, \alpha \overline{z}_k + (1 - \alpha)\tilde{z}_k)$$

$$< \alpha F(\overline{x}, \overline{z}_k) + (1 - \alpha)F(\tilde{x}, \tilde{z}_k).$$

By taking the limit as $k \to \infty$, we obtain

$$f(\alpha \overline{x} + (1 - \alpha)\tilde{x}) \le \alpha f(\overline{x}) + (1 - \alpha)f(\tilde{x}) \le \alpha \overline{w} + (1 - \alpha)\tilde{w}.$$

It follows that the point $\alpha(\overline{x}, \overline{w}) + (1 - \alpha)(\tilde{x}, \tilde{w})$ belongs to epi(f). Thus epi(f) is convex, implying that f is convex.

(b) To show the left-hand side of Eq. (3.1), let $(x, w) \in P(\operatorname{epi}(F))$, so that there exists \overline{z} such that $(x, \overline{z}, w) \in \operatorname{epi}(F)$, or equivalently $F(x, \overline{z}) \leq w$. Then

$$f(x) = \inf_{z \in \Re^m} F(x, z) \le w,$$

implying that $(x, w) \in epi(f)$.

To show the right-hand side, note that for any $(x, w) \in \operatorname{epi}(f)$ and every k, there exists a z_k such that

$$(x, z_k, w + 1/k) \in \operatorname{epi}(F),$$

so that
$$(x, w + 1/k) \in P(\operatorname{epi}(F))$$
 and $(x, w) \in \operatorname{cl}(P(\operatorname{epi}(F)))$. Q.E.D.

Among other things, part (b) of the preceding proposition asserts that if F is closed, and if the projection operation preserves closedness of its epigraph, then partial minimization of F yields a closed function. Note also a connection between closedness of $P(\operatorname{epi}(F))$ and the attainment of the infimum of F(x,z) over z. As illustrated in Fig. 3.3.1, for a fixed x, F(x,z) attains a minimum over z if and only if (x,f(x)) belongs to $P(\operatorname{epi}(F))$. Thus if $P(\operatorname{epi}(F))$ is closed, F(x,z) attains a minimum over z for all x such that f(x) is finite.

We now provide criteria guaranteeing that closedness is preserved under partial minimization, while simultaneously the partial minimum is attained.

Proposition 3.3.2: Let $F: \Re^{n+m} \mapsto (-\infty, \infty]$ be a closed proper convex function, and consider the function f given by

$$f(x) = \inf_{z \in \mathbb{R}^m} F(x, z), \qquad x \in \mathbb{R}^n.$$

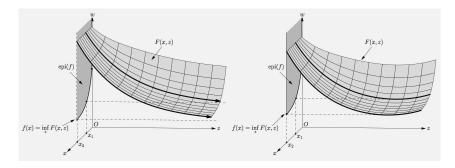


Figure 3.3.1. Illustration of partial minimization and the relation

$$P\Big(\mathrm{epi}(F)\Big)\subset\mathrm{epi}(f)\subset\mathrm{cl}\Big(P\Big(\mathrm{epi}(F)\Big)\Big).$$

In the figure on the right, where equality holds throughout in the above relation, the minimum over z is attained for all x such that f(x) is finite. In the figure on the left the minimum over z is not attained and $P\Big(\operatorname{epi}(F)\Big)$ is not equal to $\operatorname{epi}(f)$ because it is missing the points $\Big(x,f(x)\Big)$.

Assume that for some $\overline{x} \in \Re^n$ and $\overline{\gamma} \in \Re$ the set

$$\{z \mid F(\overline{x}, z) \le \overline{\gamma}\}$$

is nonempty and compact. Then f is closed proper convex. Furthermore, for each $x \in \text{dom}(f)$, the set of minima in the definition of f(x) is nonempty and compact.

Proof: We first note that by Prop. 3.3.1(a), f is convex. By Prop. 1.4.5(a), the recession cone of F has the form

$$R_F = \{(d_x, d_z) \mid (d_x, d_z, 0) \in R_{\text{epi}(F)}\}.$$

The (common) recession cone of the nonempty level sets of $F(\overline{x},\cdot)$ has the form

$$\{d_z \mid (0, d_z) \in R_F\},\$$

and by the compactness hypothesis, using also Props. 1.4.2(a) and 1.4.5(b), it consists of just the origin. Thus there exists no vector $d_z \neq 0$ such that $(0, d_z, 0) \in R_{\operatorname{epi}(F)}$. Equivalently, there is no nonzero vector in $R_{\operatorname{epi}(F)}$ that belongs to the nullspace of $P(\cdot)$, where $P(\cdot)$ denotes projection on the space of (x, w). Therefore by the closedness of $\operatorname{epi}(F)$ and Prop. 1.4.13, $P(\operatorname{epi}(F))$ is a closed set, and from the relation (3.1) it follows that f is closed.

For any $x \in \text{dom}(f)$ the (common) recession cone of the nonempty level sets of $F(x,\cdot)$ consists of just the origin. Therefore the set of minima of F(x,z) over $z \in \Re^m$ is nonempty and compact (cf. Prop. 3.2.2). Furthermore, for all $x \in \text{dom}(f)$, we have $f(x) > -\infty$, and since dom(f) is nonempty (it contains \overline{x}), it follows that f is proper. **Q.E.D.**

Note that as the preceding proof indicates, the assumption that there exists a vector $\overline{x} \in \Re^n$ and a scalar $\overline{\gamma}$ such that the level set

$$\{z \mid F(\overline{x}, z) \leq \overline{\gamma}\}$$

is nonempty and compact is equivalent to assuming that all the nonempty level sets of the form $\{z \mid F(x,z) \leq \gamma\}$ are compact. This is so because all these sets have the same recession cone, namely $\{d_z \mid (0,d_z) \in R_F\}$.

A simple but useful special case of the preceding proposition is the following.

Proposition 3.3.3: Let X and Z be nonempty convex sets of \mathbb{R}^n and \mathbb{R}^m , respectively, let $F: X \times Z \mapsto \mathbb{R}$ be a closed convex function, and assume that Z is compact. Then the function f given by

$$f(x) = \inf_{z \in Z} F(x, z), \qquad x \in X,$$

is a real-valued convex function over X.

Proof: Apply Prop. 3.3.2 with the partially minimized function being equal to F(x, z) for $x \in X$ and $z \in Z$, and ∞ otherwise. **Q.E.D.**

The following is a slight generalization of Prop. 3.3.2.

Proposition 3.3.4: Let $F: \Re^{n+m} \mapsto (-\infty, \infty]$ be a closed proper convex function, and consider the function f given by

$$f(x) = \inf_{z \in \Re^m} F(x, z), \qquad x \in \Re^n.$$

Assume that for some $\overline{x} \in \Re^n$ and $\overline{\gamma} \in \Re$ the set

$$\{z \mid F(\overline{x}, z) \leq \overline{\gamma}\}$$

is nonempty and its recession cone is equal to its lineality space. Then f is closed proper convex. Furthermore, for each $x \in \text{dom}(f)$, the set of minima in the definition of f(x) is nonempty.

Proof: The proof is nearly identical to the one of Prop. 3.3.2, using Props. 1.4.13 and 3.2.4. **Q.E.D.**

3.4 SADDLE POINT AND MINIMAX THEORY

Let us consider a function $\phi: X \times Z \mapsto \Re$, where X and Z are nonempty subsets of \Re^n and \Re^m , respectively. We wish to either

minimize
$$\sup_{z \in Z} \phi(x, z)$$

subject to $x \in X$

or

Our main interest will be to derive conditions guaranteeing that

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z), \tag{3.2}$$

and that the infima and the suprema above are attained. Minimax problems are encountered in several important contexts.

One major context is zero sum games. In the simplest such game there are two players: the first may choose one out of n moves and the second may choose one out of m moves. If moves i and j are selected by the first and the second player, respectively, the first player gives a specified amount a_{ij} to the second. The objective of the first player is to minimize the amount given to the other player, and the objective of the second player is to maximize this amount. The players use mixed strategies, whereby the first player selects a probability distribution $x = (x_1, \ldots, x_n)$ over his n possible moves and the second player selects a probability distribution z = (z_1,\ldots,z_m) over his m possible moves. Since the probability of selecting i and j is $x_i z_i$, the expected amount to be paid by the first player to the second is $\sum_{i,j} a_{ij} x_i z_j$ or x'Az, where A is the $n \times m$ matrix with components a_{ij} . If each player adopts a worst case viewpoint, whereby he optimizes his choice against the worst possible selection by the other player, the first player must minimize $\max_z x'Az$ and the second player must maximize $\min_x x'Az$. The main result, a special case of a theorem we will prove in Chapter 5, is that these two optimal values are equal, implying that there is an amount that can be meaningfully viewed as the value of the game for its participants.

Another major context is duality theory for inequality-constrained problems of the form

minimize
$$f(x)$$

subject to $x \in X$, $g(x) \le 0$, (3.3)

where $f: \mathbb{R}^n \to \mathbb{R}$, $g_j: \mathbb{R}^n \to \mathbb{R}$ are given functions, and X is a nonempty subset of \mathbb{R}^n . A common approach here is to introduce a vector $\mu = (\mu_1, \ldots, \mu_r) \in \mathbb{R}^r$, consisting of multipliers for the inequality constraints, and to form the Lagrangian function

$$L(x, \mu) = f(x) + \sum_{j=1}^{r} \mu_j g_j(x).$$

We may then form a dual problem of optimization over μ of a cost function defined in terms of the Lagrangian. In particular, we will introduce later the dual problem

$$\text{maximize} \quad \inf_{x \in X} L(x, \mu)$$

subject to
$$\mu \geq 0$$
.

Note that the original problem (3.3) can also be written as

$$\mbox{minimize} \quad \sup_{\mu \geq 0} L(x,\mu)$$

subject to
$$x \in X$$

[if x violates any of the constraints $g_j(x) \leq 0$, we have $\sup_{\mu \geq 0} L(x,\mu) = \infty$, and if it does not, we have $\sup_{\mu \geq 0} L(x,\mu) = f(x)$]. A major question is whether there is no duality gap, i.e., whether the optimal primal and dual values are equal. This is so if and only if

$$\sup_{\mu \ge 0} \inf_{x \in X} L(x, \mu) = \inf_{x \in X} \sup_{\mu \ge 0} L(x, \mu), \tag{3.4}$$

where L is the Lagrangian function.

In Section 5.5 we will prove a classical result, von Neumann's Saddle Point Theorem, which guarantees the minimax equality (3.2), as well as the attainment of the infima and suprema, assuming convexity/concavity assumptions on ϕ , and compactness assumptions on X and Z. Unfortunately, von Neumann's Theorem is not fully adequate for the development of constrained optimization duality theory, because compactness of Z and to some extent compactness of X turn out to be restrictive assumptions [for example Z corresponds to the set $\{\mu \mid \mu \geq 0\}$ in Eq. (3.4), which is not compact].

A first observation regarding the potential validity of the minimax equality (3.2) is that we always have

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) \le \inf_{x \in X} \sup_{z \in Z} \phi(x, z), \tag{3.5}$$

[for every $\overline{z} \in Z$, write

$$\inf_{x \in X} \phi(x, \overline{z}) \le \inf_{x \in X} \sup_{z \in Z} \phi(x, z)$$

and take the supremum over $\overline{z} \in Z$ of the left-hand side]. We refer to this relation as the *minimax inequality*. It is sufficient to show the reverse inequality in order for the minimax equality (3.2) to hold. However, special conditions are required for the reverse inequality to be true.

Saddle Points

The following definition formalizes pairs of vectors that attain the infimum and the supremum in the minimax equality (3.2).

Definition 3.4.1: A pair of vectors $x^* \in X$ and $z^* \in Z$ is called a saddle point of ϕ if

$$\phi(x^*, z) \le \phi(x^*, z^*) \le \phi(x, z^*), \quad \forall x \in X, \forall z \in Z.$$

Note that (x^*, z^*) is a saddle point if and only if $x^* \in X$, $z^* \in Z$, and

$$\sup_{z \in Z} \phi(x^*, z) = \phi(x^*, z^*) = \inf_{x \in X} \phi(x, z^*), \tag{3.6}$$

i.e., if " x^* minimizes against z^* " and " z^* maximizes against x^* ". We have the following characterization of a saddle point.

Proposition 3.4.1: A pair (x^*, z^*) is a saddle point of ϕ if and only if the minimax equality (3.2) holds, and x^* is an optimal solution of the problem

minimize
$$\sup_{z \in Z} \phi(x, z)$$
 (3.7)

subject to $x \in X$,

while z^* is an optimal solution of the problem

$$\begin{array}{ll} \text{maximize} & \inf_{x \in X} \phi(x,z) \\ \text{subject to} & z \in Z. \end{array}$$
 (3.8)

Proof: Suppose that x^* is an optimal solution of problem (3.7) and z^* is an optimal solution of problem (3.8). Then we have

$$\sup_{z\in Z}\inf_{x\in X}\phi(x,z)=\inf_{x\in X}\phi(x,z^*)\leq \phi(x^*,z^*)\leq \sup_{z\in Z}\phi(x^*,z)=\inf_{x\in X}\sup_{z\in Z}\phi(x,z).$$

If the minimax equality [cf. Eq. (3.2)] holds, then equality holds throughout above, so that

$$\sup_{z \in Z} \phi(x^*, z) = \phi(x^*, z^*) = \inf_{x \in X} \phi(x, z^*),$$

i.e., (x^*, z^*) is a saddle point of ϕ [cf. Eq. (3.6)].

Conversely, if (x^*, z^*) is a saddle point, then using Eq. (3.6), we have

$$\inf_{x \in X} \sup_{z \in Z} \phi(x, z) \le \sup_{z \in Z} \phi(x^*, z) = \phi(x^*, z^*) = \inf_{x \in X} \phi(x, z^*) \le \sup_{z \in Z} \inf_{x \in X} \phi(x, z).$$
(3.9)

Combined with the minimax inequality (3.5), this relation shows that the minimax equality (3.2) holds. Therefore, equality holds throughout in Eq. (3.9), which implies that x^* and z^* are optimal solutions of problems (3.7) and (3.8), respectively. **Q.E.D.**

Note a simple consequence of Prop. 3.4.1: assuming the minimax equality (3.2) holds, the set of saddle points is the Cartesian product $X^* \times Z^*$, where X^* and Z^* are the sets of optimal solutions of problems (3.7) and (3.8), respectively. In other words x^* and z^* can be independently chosen from the sets X^* and Z^* to form a saddle point. Note also that if the minimax equality (3.2) does not hold, there is no saddle point, even if the sets X^* and Z^* are nonempty.

To obtain saddle points using Prop. 3.4.1, we may calculate the "sup" and "inf" functions appearing in the proposition, then minimize and maximize them, respectively, and obtain the corresponding sets of minima X^* and maxima Z^* [cf. Eqs. (3.7) and (3.8)]. If the optimal values are equal (i.e., $\inf_x \sup_z \phi = \sup_z \inf_x \phi$), the set of saddle points is $X^* \times Z^*$. Otherwise, there are no saddle points. The following example illustrates this process.

Example 3.4.1:

Consider the case where

$$\phi(x,z) = \frac{1}{2}x'Qx + x'z - \frac{1}{2}z'Rz, \qquad X = Z = \Re^n,$$

and Q and R are symmetric invertible matrices. If Q and R are positive definite, a straightforward calculation shows that

$$\sup_{z \in \Re^n} \phi(x, z) = \frac{1}{2} x'(Q + R^{-1})x, \qquad \inf_{x \in \Re^n} \phi(x, z) = -\frac{1}{2} z'(Q^{-1} + R)z.$$

Therefore, $\inf_x \sup_z \phi(x,z) = \sup_z \inf_x \phi(x,z) = 0$, we have

$$X^* = Z^* = \{0\},\,$$

and it follows that (0,0) is the unique saddle point.

Assume now that Q is not positive semidefinite, but R is positive definite and such that $Q + R^{-1}$ is positive definite. Then it can be seen that

$$\sup\nolimits_{z\in\Re^n}\phi(x,z)=\tfrac{1}{2}x'(Q+R^{-1})x,\qquad\inf\nolimits_{x\in\Re^n}\phi(x,z)=-\infty,\ \ \forall\ z\in\Re^n.$$

Here, the sets X^* and Z^* are nonempty $(X^* = \{0\})$ and $Z^* = \Re^n$. However, we have $0 = \inf_x \sup_z \phi(x, z) > \sup_z \inf_x \phi(x, z) = -\infty$, and there are no saddle points.

$Geometric\ Duality\ Framework$

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Duality in optimization is often viewed as a manifestation of the fundamental description of a closed convex set as the intersection of all closed halfspaces containing the set (cf. Prop. 1.5.4). When specialized to the epigraph of a function $f: \mathbb{R}^n \to [-\infty, \infty]$, this description leads to the formalism of the conjugate function of f:

$$f^{\star}(y) = \sup_{x \in \Re^n} \left\{ x'y - f(x) \right\}. \tag{4.1}$$

Indeed, if f is closed proper convex, by the Conjugacy Theorem (Prop. 1.6.1), we have $f(x) = f^{**}(x) = \sup_{y \in \Re^n} \{x'y - f^*(y)\}$, so that

$$\begin{aligned} \text{epi}(f) &= \big\{ (x, w) \mid f(x) \le w \big\} \\ &= \big\{ (x, w) \mid \sup_{y \in \Re^n} \big\{ x'y - f^*(y) \big\} \le w \big\} \\ &= \cap_{y \in \Re^n} \big\{ (x, w) \mid x'y - w \le f^*(y) \big\}. \end{aligned}$$

Thus the conjugate f^* defines epi(f) as the intersection of closed halfspaces.

In this chapter, we focus on a geometric framework for duality analysis, referred to as *min common/max crossing* (MC/MC for short). It is inspired by the above description of epigraphs, and it is related to conjugacy, but it does not involve an algebraic definition such as Eq. (4.1). For this reason it is simpler and better suited for geometric visualization and analysis in many important convex optimization contexts.

4.1 MIN COMMON/MAX CROSSING DUALITY

Our framework aims to capture the most essential characteristics of duality in two simple geometrical problems, defined by a nonempty subset M of \Re^{n+1} .

- (a) Min Common Point Problem: Consider all vectors that are common to M and the (n + 1)st axis. We want to find one whose (n + 1)st component is minimum.
- (b) Max Crossing Point Problem: Consider nonvertical hyperplanes that contain M in their corresponding "upper" closed halfspace, i.e., the closed halfspace whose recession cone contains the vertical halfline $\{(0,w) \mid w \geq 0\}$ (see Fig. 4.1.1). We want to find the maximum crossing point of the (n+1)st axis with such a hyperplane.

We refer to the two problems as the min common/max crossing (MC/MC) framework, and we will show that it can be used to develop much of the core theory of convex optimization in a unified way.

Figure 4.1.1 suggests that the optimal value of the max crossing problem is no larger than the optimal value of the min common problem, and

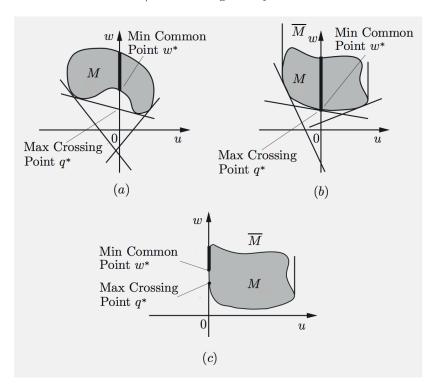


Figure 4.1.1. Illustration of the optimal values of the min common and max crossing problems. In (a), the two optimal values are not equal. In (b), when M is "extended upwards" along the (n+1)st axis, it yields the set

$$\begin{split} \overline{M} &= M + \Big\{ (0,w) \mid w \geq 0 \Big\} \\ &= \Big\{ (u,w) \mid \text{there exists } \overline{w} \text{ with } \overline{w} \leq w \text{ and } (u,\overline{w}) \in M \Big\}, \end{split}$$

which is convex and admits a nonvertical supporting hyperplane passing through $(0,w^*)$. As a result, the two optimal values are equal. In (c), the set \overline{M} is convex but not closed, and there are points $(0,\overline{w})$ on the vertical axis with $\overline{w} < w^*$ that lie in the closure of \overline{M} . Here q^* is equal to the minimum such value of \overline{w} , and we have $q^* < w^*$.

that under favorable circumstances the two optimal values are equal. In this chapter, we will provide conditions that guarantee equality of the optimal values and the existence of optimal solutions. In the next chapter we will use these conditions to analyze a variety of duality-related issues.

Mathematically, the min common problem is

Its optimal value is denoted by w^* , i.e.,

$$w^* = \inf_{(0,w) \in M} w.$$

To describe mathematically the max crossing problem, we recall that a nonvertical hyperplane in \Re^{n+1} is specified by its normal vector $(\mu, 1) \in \Re^{n+1}$, and a scalar ξ as

$$H_{\mu,\xi} = \{(u, w) \mid w + \mu' u = \xi\}.$$

Such a hyperplane crosses the (n+1)st axis at $(0,\xi)$. For M to be contained in the "upper" closed halfspace that corresponds to $H_{\mu,\xi}$ [the one that contains the vertical halfline $\{(0,w) \mid w \geq 0\}$ in its recession cone], it is necessary and sufficient that

$$\xi \le w + \mu' u, \qquad \forall \ (u, w) \in M,$$

or equivalently

$$\xi \le \inf_{(u,w)\in M} \{w + \mu'u\}.$$

For a fixed normal $(\mu, 1)$, the maximum crossing level ξ over all hyperplanes $H_{\mu,\xi}$ is denoted by $q(\mu)$ and is given by

$$q(\mu) = \inf_{(u,w) \in M} \{ w + \mu' u \}; \tag{4.2}$$

(see Fig. 4.1.2). The max crossing problem is to maximize over all $\mu \in \Re^n$ the maximum crossing level corresponding to μ , i.e.,

maximize
$$q(\mu)$$
 subject to $\mu \in \Re^n$.

We also refer to this as the dual problem, we denote by q^* its optimal value,

$$q^* = \sup_{\mu \in \Re^n} q(\mu),$$

and we refer to $q(\mu)$ as the *crossing* or *dual* function.

We will show shortly that we always have $q^* \leq w^*$; we refer to this as weak duality. When $q^* = w^*$, we say that strong duality holds or that there is no duality gap.

Note that both w^* and q^* remain unaffected if M is replaced by its "upwards extension"

$$\overline{M} = M + \{(0, w) \mid w \ge 0\}$$

$$= \{(u, w) \mid \text{there exists } \overline{w} \text{ with } \overline{w} \le w \text{ and } (u, \overline{w}) \in M\}$$

$$(4.3)$$

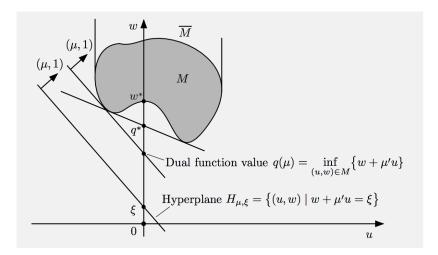


Figure 4.1.2. Mathematical specification of the max crossing problem. When considering crossing points by nonvertical hyperplanes, it is sufficient to restrict attention to hyperplanes of the form

$$H_{\mu,\xi} = \{(u, w) \mid w + \mu' u = \xi\},\$$

that have normal $(\mu, 1)$ and cross the vertical axis at a point ξ with $\xi \leq w + \mu' u$ for all $(u, w) \in M$, or equivalently $\xi \leq \inf_{(u, w) \in M} \{w + \mu' u\}$. For a given $\mu \in \Re^n$, the highest crossing level is

$$q(\mu) = \inf_{(u,w) \in M} \{w + \mu' u\}.$$

To find the max crossing point q^* we must maximize $q(\mu)$ over $\mu \in \Re^n$.

(cf. Fig. 4.1.1). It is often more convenient to work with \overline{M} because in many cases of interest \overline{M} is convex while M is not. However, on occasion M has some interesting properties (such as compactness) that are masked when passing to \overline{M} , in which case it may be preferable to work with M.

Note also that we do not exclude the possibility that either w^* or q^* (or both) are infinite. In particular, we have $w^* = \infty$ if the min common problem has no feasible solution $[M \cap \{(0,w) \mid w \in \Re\} = \varnothing]$. Similarly, we have $q^* = -\infty$ if the max crossing problem has no feasible solution, which occurs in particular if \overline{M} contains a vertical line, i.e., a set of the form $\{(x,w) \mid w \in \Re\}$ for some $x \in \Re^n$.

We now establish some generic properties of the MC/MC framework.

Proposition 4.1.1: The dual function q is concave and upper semi-continuous.

Proof: By definition [cf. Eq. (4.2)], q is the infimum of a collection of affine functions, so -q is closed (cf. Prop. 1.1.6). The result follows from Prop. 1.1.2. **Q.E.D.**

We next establish the weak duality property, which is intuitively apparent from Fig. 4.1.1.

Proposition 4.1.2: (Weak Duality Theorem) We have $q^* \leq w^*$.

Proof: For every $(u, w) \in M$ and $\mu \in \mathbb{R}^n$, we have

$$q(\mu)=\inf_{(u,w)\in M}\{w+\mu'u\}\leq \inf_{(0,w)\in M}w=w^*,$$

so by taking the supremum of the left-hand side over $\mu \in \Re^n$, we obtain $q^* \leq w^*$. Q.E.D.

As Fig. 4.1.2 indicates, the feasible solutions of the max crossing problem are restricted by the horizontal directions of recession of \overline{M} . This is the essence of the following proposition.

Proposition 4.1.3: Assume that the set

$$\overline{M} = M + \{(0, w) \mid w \ge 0\}$$

is convex. Then the set of feasible solutions of the max crossing problem, $\big\{\mu\mid q(\mu)>-\infty\big\}$, is contained in the cone

$$\{\mu \mid \mu'd \ge 0 \text{ for all } d \text{ with } (d,0) \in R_{\overline{M}}\},\$$

where $R_{\overline{M}}$ is the recession cone of \overline{M} .

Proof: Let $(\overline{u}, \overline{w}) \in \overline{M}$. If $(d, 0) \in R_{\overline{M}}$, then $(\overline{u} + \alpha d, \overline{w}) \in \overline{M}$ for all $\alpha \geq 0$, so that for all $\mu \in \Re^n$,

$$q(\mu) = \inf_{(u,w) \in M} \{w + \mu'u\} = \inf_{(u,w) \in \overline{M}} \{w + \mu'u\} \leq \overline{w} + \mu'\overline{u} + \alpha\mu'd, \quad \forall \; \alpha \geq 0.$$

Thus, if $\mu'd < 0$, we must have $q(\mu) = -\infty$, implying that $\mu'd \ge 0$ for all μ with $q(\mu) > -\infty$. **Q.E.D.**

As an example, consider the case where \overline{M} is the vector sum of a convex set and the nonnegative orthant of \Re^{n+1} . Then it can be seen that

the set $\{d \mid (d,0) \in R_{\overline{M}}\}$ contains the nonnegative orthant of \Re^n , so the preceding proposition implies that $\mu \geq 0$ for all μ such that $q(\mu) > -\infty$. This case arises in optimization problems with inequality constraints (see Section 5.3).

4.2 SOME SPECIAL CASES

We will now consider some examples illustrating the use of the MC/MC framework. In all of these examples the set M is the epigraph of some function.

4.2.1 Connection to Conjugate Convex Functions

Consider the case where the set M is the epigraph of a function $p: \Re^n \mapsto [-\infty, \infty]$. Then M coincides with its upwards extension \overline{M} of Eq. (4.3) (cf. Fig. 4.1.1), and the min common value is

$$w^* = p(0).$$

The dual function q of Eq. (4.2) is given by

$$q(\mu) = \inf_{(u,w) \in \operatorname{epi}(p)} \{ w + \mu' u \} = \inf_{\{(u,w) \mid p(u) \leq w\}} \{ w + \mu' u \},$$

and finally

$$q(\mu) = \inf_{u \in \Re^m} \{ p(u) + \mu' u \}. \tag{4.4}$$

Thus, we have $q(\mu) = -p^{\star}(-\mu)$, where p^{\star} is the conjugate of p:

$$p^{\star}(\mu) = \sup_{u \in \Re^n} \bigl\{ \mu' u - p(u) \bigr\}.$$

Furthermore,

$$q^* = \sup_{\mu \in \Re^n} q(\mu) = \sup_{\mu \in \Re^n} \left\{ 0 \cdot (-\mu) - p^*(-\mu) \right\} = p^{**}(0),$$

where p^{**} is the conjugate of p^{*} (double conjugate of p); see Fig. 4.2.1. In particular, if $p = p^{**}$ (e.g., if p is closed proper convex), then $w^{*} = q^{*}$.

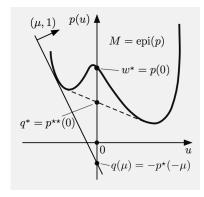


Figure 4.2.1. Illustration of the MC/MC framework for the case where the set M is the epigraph of a function $p: \Re^n \mapsto [-\infty, \infty]$. We have

$$q(\mu) = -p^*(-\mu),$$

$$w^* = p(0), \qquad q^* = p^{**}(0),$$

where $p^{\star\star}$ is the double conjugate of p. If p is closed proper convex, by the Conjugacy Theorem (cf. Prop. 1.6.1), we have $p=p^{\star\star}$, so $w^*=q^*$.

Note that there is a symmetry between the min common and max crossing problems: if we consider the min common problem defined by the epigraph of -q, the corresponding dual function is $-p^{**}$ [or -p when p is closed proper and convex, by the Conjugacy Theorem (cf. Prop. 1.6.1)]. Moreover, by Prop. 1.6.1(d), the min common problems corresponding to $M = \operatorname{epi}(p)$ and $M = \operatorname{epi}(\operatorname{cl} p)$ have the same dual function, namely $q(\mu) = -p^*(-\mu)$. Furthermore, if $\operatorname{cl} p$ is proper, then $\operatorname{cl} p = p^{**}$, so the min common problem corresponding to $M = \operatorname{epi}(p^{**})$ also has the same dual function.

4.2.2 General Optimization Duality

Consider the problem of minimizing a function $f: \Re^n \mapsto [-\infty, \infty]$. We introduce a function $F: \Re^{n+r} \mapsto [-\infty, \infty]$ of the pair (x, u), which satisfies

$$f(x) = F(x,0), \qquad \forall \ x \in \Re^n. \tag{4.5}$$

Let the function $p: \Re^r \mapsto [-\infty, \infty]$ be defined by

$$p(u) = \inf_{x \in \Re^n} F(x, u). \tag{4.6}$$

For a helpful insight, we may view u as a perturbation, and p(u) as the optimal value of an optimization problem whose cost function is perturbed by u. When u=0, the perturbed problem coincides with the original problem of minimizing f. Note that p is defined by a partial minimization, so its closedness may be checked using the theory of Section 3.3.

Consider the MC/MC framework with

$$M = epi(p).$$

The min common value w^* is the minimal value of f, since

$$w^* = p(0) = \inf_{x \in \mathbb{R}^n} F(x, 0) = \inf_{x \in \mathbb{R}^n} f(x).$$

By Eq. (4.4), the dual function is

$$q(\mu) = \inf_{u \in \Re^r} \big\{ p(u) + \mu' u \big\} = \inf_{(x,u) \in \Re^{n+r}} \big\{ F(x,u) + \mu' u \big\}, \tag{4.7}$$

and the max crossing problem is

maximize
$$q(\mu)$$

subject to $\mu \in \Re^r$.

Note that from Eq. (4.7), an alternative expression for q is

$$q(\mu) = -\sup_{(x,u) \in \Re^{n+r}} \left\{ -\mu' u - F(x,u) \right\} = -F^*(0,-\mu),$$

where F^* is the conjugate of F, viewed as a function of (x, u). Since

$$q^* = \sup_{\mu \in \Re^r} q(\mu) = -\inf_{\mu \in \Re^r} F^{\star}(0, -\mu) = -\inf_{\mu \in \Re^r} F^{\star}(0, \mu),$$

the strong duality relation $w^* = q^*$ can be written as

$$\inf_{x \in \Re^n} F(x,0) = -\inf_{\mu \in \Re^r} F^{\star}(0,\mu).$$

4.2.3 Optimization with Inequality Constraints

Different choices of perturbation structure and function F, as in Eqs. (4.5) and (4.6), yield corresponding MC/MC frameworks and dual problems. An example of this type, to be considered in detail later, is minimization with inequality constraints:

minimize
$$f(x)$$

subject to $x \in X$, $g(x) \le 0$, (4.8)

where X is a nonempty subset of \Re^n , $f: X \mapsto \Re$ is a given function, and $g(x) = (g_1(x), \ldots, g_r(x))$ with $g_j: X \mapsto \Re$ being given functions. We introduce a "perturbed constraint set" of the form

$$C_u = \{ x \in X \mid g(x) \le u \}, \qquad u \in \Re^r, \tag{4.9}$$

and the function

$$F(x, u) = \begin{cases} f(x) & \text{if } x \in C_u, \\ \infty & \text{otherwise,} \end{cases}$$

which satisfies the condition F(x,0) = f(x) for all $x \in C_0$ [cf. Eq. (4.5)]. The function p of Eq. (4.6) is given by

$$p(u) = \inf_{x \in \mathbb{R}^n} F(x, u) = \inf_{x \in X, g(x) \le u} f(x), \tag{4.10}$$

and is known as the *primal function* or *perturbation function* (see Fig. 4.2.2). It captures the essential structure of the constrained minimization problem, relating to duality and other properties, such as sensitivity (the change in optimal cost as a result of changes in constraint level).

Consider now the MC/MC framework corresponding to M = epi(p). From Eq. (4.4) [or Eq. (4.7)],

$$q(\mu) = \inf_{u \in \Re^r} \left\{ p(u) + \mu' u \right\}$$

$$= \inf_{x \in X, g(x) \le u} \left\{ f(x) + \mu' u \right\}$$

$$= \begin{cases} \inf_{x \in X} \left\{ f(x) + \mu' g(x) \right\} & \text{if } \mu \ge 0, \\ -\infty & \text{otherwise.} \end{cases}$$

$$(4.11)$$

Later, in Section 5.3, we will view q as the standard dual function, obtained by minimizing over $x \in X$ the function $f(x) + \mu' g(x)$, which is known as the Lagrangian function (compare also with the discussion of duality and its connection with minimax theory in Section 3.4).

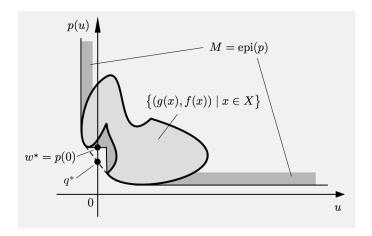


Figure 4.2.2. Illustration of the perturbation function p of Eq. (4.10). Note that epi(p) is the sum of the positive orthant with the set $\{(g(x), f(x)) \mid x \in X\}$.

Example 4.2.1: (Linear Programming Duality)

Consider the linear program

minimize
$$c'x$$

subject to $a'_j x \ge b_j$, $j = 1, ..., r$,

where $c \in \mathbb{R}^n$, $a_j \in \mathbb{R}^n$, and $b_j \in \mathbb{R}$, $j = 1, \ldots, r$. For $\mu \geq 0$, the corresponding dual function has the form

$$q(\mu) = \inf_{x \in \mathbb{R}^n} \left\{ c'x + \sum_{j=1}^r \mu_j (b_j - a'_j x) \right\} = \begin{cases} b'\mu & \text{if } \sum_{j=1}^r a_j \mu_j = c, \\ -\infty & \text{otherwise,} \end{cases}$$

while for all other $\mu \in \Re^r$, $q(\mu) = -\infty$ [cf. Eq. (4.11)]. Thus the dual problem is

maximize
$$b'\mu$$
 subject to $\sum_{j=1}^{r} a_j \mu_j = c, \quad \mu \ge 0.$

4.2.4 Augmented Lagrangian Duality

In connection with the inequality-constrained problem (4.8), let us consider the function

$$F_c(x, u) = \begin{cases} f(x) + \frac{c}{2} ||u||^2 & \text{if } x \in C_u, \\ \infty & \text{otherwise,} \end{cases}$$

where C_u is the perturbed constraint set (4.9), and c is a positive scalar. Then instead of the perturbation function p of Eq. (4.10), we have

$$p_c(u) = \inf_{x \in \mathbb{R}^n} F_c(x, u) = \inf_{x \in X, \ q(x) < u} \left\{ f(x) + \frac{c}{2} ||u||^2 \right\}$$

[yet $p(0) = p_c(0)$, so the min common value is left unchanged]. The corresponding dual function is

$$q_c(\mu) = \inf_{u \in \Re^r} \left\{ p_c(u) + \mu' u \right\} = \inf_{x \in X, \, g(x) \le u} \left\{ f(x) + \mu' u + \frac{c}{2} \|u\|^2 \right\}.$$

For a fixed $x \in X$, we find the infimum above by separately minimizing, over each component u_j of u, the one-dimensional quadratic function $\mu_j u_j + \frac{c}{2}(u_j)^2$ subject to the constraint $g_j(x) \leq u_j$. The minimum is attained at

$$g_j^+(x, \mu, c) = \max\left\{-\frac{\mu_j}{c}, g_j(x)\right\}, \quad j = 1, \dots, r,$$

and substituting into the expression for q_c , we obtain

$$q_c(\mu) = \inf_{x \in X} \left\{ f(x) + \mu' g^+(x, \mu, c) + \frac{c}{2} \|g^+(x, \mu, c)\|^2 \right\}, \qquad \mu \in \Re^r,$$

where $g^+(x, \mu, c)$ is the vector with components $g_i^+(x, \mu, c)$, $j = 1, \dots, r$.

The expression within braces in the right-hand side above is known as the augmented Lagrangian function. It plays a central role in important practical algorithms that aim to maximize the "penalized" dual function $q_c(\mu)$ over μ (we refer to the monograph [Ber82] and standard nonlinear programming textbooks). Note that [in contrast with the corresponding expression for q in Eq. (4.11)], q_c is often real-valued (for example, when f and g_j are continuous functions, and X is a compact set), and in some important contexts it turns out to be differentiable. Based on this property, augmented Lagrangian functions provide a form of regularization of the dual problem that is significant in algorithmic practice.

4.2.5 Minimax Problems

Consider a function $\phi: X \times Z \mapsto \Re$, where X and Z are nonempty subsets of \Re^n and \Re^m , respectively. As in Section 3.4, we wish to either

minimize
$$\sup_{z \in Z} \phi(x, z)$$

subject to $x \in X$

or

maximize
$$\inf_{x \in X} \phi(x, z)$$

subject to $z \in Z$.

An important question is whether the minimax equality

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z), \tag{4.12}$$

holds, and whether the infimum and the supremum above are attained. This is significant in a zero sum game context, as well as in optimization duality theory; cf. Section 3.4.

We introduce the function $p: \Re^m \mapsto [-\infty, \infty]$ given by

$$p(u) = \inf_{x \in X} \sup_{z \in Z} \bigl\{ \phi(x,z) - u'z \bigr\}, \qquad u \in \Re^m, \tag{4.13} \label{eq:4.13}$$

which can be viewed as a perturbation function. It characterizes how the "infsup" of the function ϕ changes when the linear perturbation term u'z is subtracted from ϕ . We consider the MC/MC framework with

$$M = epi(p),$$

so that the min common value is

$$w^* = p(0) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z). \tag{4.14}$$

We will show that the max crossing value q^* is equal to the "supinf" value of a function derived from ϕ via a convexification operation.

We recall that given a set $X \subset \Re^n$ and a function $f: X \mapsto [-\infty, \infty]$, the convex closure of f, denoted by $\operatorname{cl} f$, is the function whose epigraph is the closure of the convex hull of the epigraph of f (cf. Section 1.3.3). The concave closure of f, denoted by $\operatorname{cl} f: \Re^n \mapsto [-\infty, \infty]$, is the opposite of the convex closure of -f, i.e.,

$$\hat{\operatorname{cl}} f = -\check{\operatorname{cl}} (-f).$$

It is the smallest concave and upper semicontinuous function that majorizes f, i.e., $\hat{\operatorname{cl}} f \leq g$ for any $g: X \mapsto [-\infty, \infty]$ that is concave and upper semicontinuous with $g \geq f$ (Prop. 1.3.14). Note that by Prop. 1.3.13, we have

$$\sup_{x \in X} f(x) = \sup_{x \in X} (\widehat{\operatorname{cl}} f)(x) = \sup_{x \in \operatorname{conv}(X)} (\widehat{\operatorname{cl}} f)(x) = \sup_{x \in \Re^n} (\widehat{\operatorname{cl}} f)(x). \tag{4.15}$$

The following proposition derives the form of the dual function, under a mild assumption on $(\hat{\operatorname{cl}}\phi)(x,\cdot)$, the concave closure of the function $\phi(x,\cdot)$.

Proposition 4.2.1: Let X and Z be nonempty subsets of \Re^n and \Re^m , respectively, and let $\phi: X \times Z \mapsto \Re$ be a function. Assume that $(-\operatorname{cl}\phi)(x,\cdot)$ is proper for all $x \in X$, and consider the MC/MC framework corresponding to $M = \operatorname{epi}(p)$, where p is given by Eq. (4.13). Then the dual function is given by

$$q(\mu) = \inf_{x \in X} (\hat{\operatorname{cl}} \phi)(x, \mu), \qquad \forall \ \mu \in \Re^m.$$
 (4.16)

Proof: Let us write

$$p(u) = \inf_{x \in X} p_x(u),$$

where

$$p_x(u) = \sup_{z \in Z} \{\phi(x, z) - u'z\}, \qquad x \in X,$$
 (4.17)

and note that

$$\inf_{u \in \Re^m} \left\{ p_x(u) + u'\mu \right\} = -\sup_{u \in \Re^m} \left\{ u'(-\mu) - p_x(u) \right\} = -p_x^{\star}(-\mu), \quad (4.18)$$

where p_x^* is the conjugate of p_x . From Eq. (4.17), we see that (except for a sign change of its argument) p_x is the conjugate of $(-\phi)(x,\cdot)$ (restricted to Z), so from the Conjugacy Theorem [Prop. 1.6.1(d)], using also the properness assumption on $(-\hat{\operatorname{cl}}\phi)(x,\cdot)$, we have

$$p_x^{\star}(-\mu) = -(\hat{\text{cl}}\,\phi)(x,\mu).$$
 (4.19)

By using Eq. (4.4) for q, and Eqs. (4.18), (4.19), we obtain for all $\mu \in \Re^m$,

$$q(\mu) = \inf_{u \in \mathbb{R}^m} \left\{ p(u) + u'\mu \right\}$$

$$= \inf_{u \in \mathbb{R}^m} \inf_{x \in X} \left\{ p_x(u) + u'\mu \right\}$$

$$= \inf_{x \in X} \inf_{u \in \mathbb{R}^m} \left\{ p_x(u) + u'\mu \right\}$$

$$= \inf_{x \in X} \left\{ -p_x^*(-\mu) \right\}$$

$$= \inf_{x \in X} \left(\hat{\operatorname{cl}} \phi \right)(x, \mu),$$

$$(4.20)$$

which is the desired relation. Q.E.D.

Note that without the properness assumption on $(-\hat{\operatorname{cl}}\phi)(x,\cdot)$ for all $x \in X$ the conclusion of Prop. 4.2.1 may fail, since then Eq. (4.19) may not hold. The key conclusion from the proposition is that, while w^* is equal to the "infsup" value of ϕ , the dual value q^* is equal to the "supinf" value of the concave closure $\hat{\operatorname{cl}}\phi$. In particular, we may observe the following:

(a) In general, we have

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) \le q^* \le w^* = \inf_{x \in X} \sup_{z \in Z} \phi(x, z), \tag{4.21}$$

where the first inequality follows using the calculation

$$\begin{split} q(\mu) &= \inf_{u \in \Re^m} \left\{ p(u) + u'\mu \right\} \\ &= \inf_{u \in \Re^m} \inf_{x \in X} \sup_{z \in Z} \left\{ \phi(x,z) + u'(\mu - z) \right\} \\ &\geq \inf_{x \in X} \phi(x,\mu), \end{split}$$

[cf. Eq. (4.4)], and taking supremum over $\mu \in Z$ (the inequality above is obtained by setting $z = \mu$), the second inequality is the weak duality relation, and the last equality is Eq. (4.14). Therefore the minimax equality (4.12) always implies the strong duality relation $q^* = w^*$.

(b) If

$$\phi(x,z) = (\hat{\operatorname{cl}}\phi)(x,z), \quad \forall \ x \in X, z \in Z,$$

as in the case where $-\phi(x,\cdot)$ is closed and convex for all $x\in X$, then using Prop. 4.2.1 and the properness of $(-\hat{\operatorname{cl}}\phi)(x,\cdot)$ for all $x\in X$, we have

$$q^* = \sup_{z \in \Re^m} q(z) = \sup_{z \in \Re^m} \inf_{x \in X} (\widehat{\operatorname{cl}} \phi)(x, z) = \sup_{z \in Z} \inf_{x \in X} \phi(x, z).$$

In view of Eq. (4.21), it follows that if $\phi = \hat{cl} \phi$, the minimax equality (4.12) is equivalent to the strong duality relation $q^* = w^*$.

(c) From Eq. (4.15), we have $\sup_{z\in Z}\phi(x,z)=\sup_{z\in\Re^m}(\hat{\operatorname{cl}}\,\phi)(x,z),$ so

$$w^* = \inf_{x \in X} \sup_{z \in \Re^m} (\hat{\operatorname{cl}} \phi)(x, z).$$

Furthermore, assuming that $(-\hat{\operatorname{cl}}\phi)(x,\cdot)$ is proper for all $x\in X$ so that Prop. 4.2.1 applies, we have

$$q^* = \sup_{z \in \Re^m} \inf_{x \in X} (\widehat{\operatorname{cl}} \phi)(x, z).$$

Thus, w^* and q^* are the "infsup" and "supinf" values of $\hat{\operatorname{cl}}\,\phi$, respectively.

(d) If $(\hat{\operatorname{cl}}\phi)(\cdot,z)$ is closed and convex for each z, then the "infsup" and "supinf" values of the convex/concave function $\hat{\operatorname{cl}}\phi$ will ordinarily be expected to be equal (see Section 5.5). In this case we will have strong duality $(q^*=w^*)$, but this will not necessarily imply that the minimax equality (4.12) holds, because strict inequality may hold in the first relation of Eq. (4.21). If $(\hat{\operatorname{cl}}\phi)(\cdot,z)$ is not closed and convex for some z, then the strong duality relation $q^*=w^*$ should not ordinarily be expected. The size of the duality gap w^*-q^* will depend on the difference between $(\hat{\operatorname{cl}}\phi)(\cdot,z)$ and its convex closure, and may be investigated in some cases by using methods to be presented in Section 5.7.

Example 4.2.2: (Finite Set Z)

Let

$$\phi(x,z) = z'f(x),$$

where $f: X \mapsto \mathbb{R}^m$, X is a subset of \mathbb{R}^n , and f(x) is viewed as a column vector whose components are functions $f_j: X \mapsto \mathbb{R}, j = 1, ..., m$. Suppose that Z is the finite set $Z = \{e_1, ..., e_m\}$, where e_j is the jth unit vector (the jth column of the $m \times m$ identity matrix). Then we have

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \max \left\{ \inf_{x \in X} f_1(x), \dots, \inf_{x \in X} f_m(x) \right\},\,$$

and

$$w^* = \inf_{x \in X} \sup_{z \in Z} \phi(x, z) = \inf_{x \in X} \max \{f_1(x), \dots, f_m(x)\}.$$

Let \overline{Z} denote the unit simplex in \Re^m (the convex hull of Z). We have

$$(\widehat{\operatorname{cl}}\,\phi)(x,z) = \begin{cases} z'f(x) & \text{if } z \in \overline{Z}, \\ -\infty & \text{if } z \notin \overline{Z}, \end{cases}$$

and by Prop. 4.2.1,

$$q^* = \sup_{z \in Z} \inf_{x \in X} (\widehat{\operatorname{cl}} \phi)(x, z) = \sup_{z \in \overline{Z}} \inf_{x \in X} z' f(x).$$

Note that it can easily happen that $\sup_{z \in Z} \inf_{x \in X} \phi(x, z) < q^*$. For example, if X = [0, 1], $f_1(x) = x$, and $f_2(x) = 1 - x$, it is straightforward to verify that

$$\sup_{z\in Z}\inf_{x\in X}\phi(x,z)=0<\frac{1}{2}=q^*.$$

If X is compact and f_1, \ldots, f_m are continuous convex functions (as in the preceding example), we may show that $q^* = w^*$ by using results to be presented in Section 5.5. On the other hand, if f_1, \ldots, f_m are not convex, we may have $q^* < w^*$.

Under certain conditions, q^* can be associated with a special minimax problem derived from the original by introducing mixed (randomized) strategies. This is so in the following classical game context.

Example 4.2.3: (Finite Zero Sum Games)

Consider a minimax problem where the sets X and Z are finite:

$$X = \{d_1 \dots, d_n\}, \qquad Z = \{e_1, \dots, e_m\},\$$

where d_i is the *i*th column of the $n \times n$ identity matrix, and e_j is the *j*th column of the $m \times m$ identity matrix. Let

$$\phi(x,z) = x'Az,$$

where A is an $n \times m$ matrix. This corresponds to the classical game context, discussed at the beginning of Section 3.4, where upon selection of $x=d_i$ and $z=e_j$, the payoff is the ijth component of A. Let \overline{X} and \overline{Z} be the unit simplexes in \Re^n and \Re^m , respectively. Then it can be seen, similar to the preceding example, that $q^* = \max_{z \in \overline{Z}} \min_{x \in X} x' Az$, and using theory to be developed in Section 5.5, it can be shown that

$$q^* = \max_{z \in \overline{Z}} \min_{x \in \overline{X}} x'Az = \min_{x \in \overline{X}} \max_{z \in \overline{Z}} x'Az.$$

Note that \overline{X} and \overline{Z} can be interpreted as sets of mixed strategies, so q^* is the value of the corresponding mixed strategy game.

4.3 STRONG DUALITY THEOREM

We will now establish conditions for strong duality to hold in the MC/MC framework, i.e., $q^* = w^*$. To avoid degenerate cases, we will often exclude the case $w^* = \infty$, which corresponds to an infeasible min common problem.

An important point, around which much of our analysis revolves, is that when w^* is finite, the vector $(0, w^*)$ is a closure point of the set \overline{M} of Eq. (4.3), so if we assume that \overline{M} is closed convex and admits a nonvertical supporting hyperplane at $(0, w^*)$, then we have $q^* = w^*$ while the optimal values q^* and w^* are attained. Between the "unfavorable" case where $q^* < w^*$, and the "most favorable" case where $q^* = w^*$ while the optimal values q^* and w^* are attained, there are several intermediate cases.

The following proposition provides a necessary and sufficient condition for $q^* = w^*$, but does not address the attainment of the optimal values. Aside from convexity of \overline{M} , it requires that points on the vertical axis that lie below w^* cannot be approached through a sequence from M; this has the flavor of a "lower semicontinuity" property for M at $(0, w^*)$.

Proposition 4.3.1: (MC/MC Strong Duality) Consider the min common and max crossing problems, and assume the following:

- (1) Either $w^* < \infty$, or else $w^* = \infty$ and M contains no vertical lines.
- (2) The set

$$\overline{M} = M + \{(0, w) \mid w \ge 0\}$$

is convex.

Then, we have $q^* = w^*$ if and only if for every sequence $\{(u_k, w_k)\} \subset M$ with $u_k \to 0$, there holds $w^* \le \liminf_{k \to \infty} w_k$.

Proof: Assume that $w^* \leq \liminf_{k \to \infty} w_k$ for all $\{(u_k, w_k)\} \subset M$ with $u_k \to 0$. We will show that $q^* = w^*$. The proof first deals with the easier cases where $w^* = -\infty$ and $w^* = \infty$, and then focuses on the case where w^* is finite. In the latter case, the key idea is that the assumption $w^* \leq \liminf_{k \to \infty} w_k$ implies that $(0, w^* - \epsilon)$ is not a closure point of \overline{M} for any $\epsilon > 0$, so [by Prop. 1.5.8(b)] it can be strictly separated from \overline{M} by a nonvertical hyperplane, which must cross the vertical axis between $w^* - \epsilon$ and w^* . This implies that $w^* - \epsilon < q^* \leq w^*$, and as $\epsilon \downarrow 0$, $q^* = w^*$.

Consider the case $w^* = -\infty$. Then, by the Weak Duality Theorem (Prop. 4.1.2), we also have $q^* = -\infty$, so the conclusion trivially follows. Consider next the case where $w^* = \infty$ and M contains no vertical lines. Since $\liminf_{k\to\infty} w_k = w^* = \infty$ for all $\{(u_k, w_k)\} \subset M$ with $u_k \to 0$, it follows that the vertical axis of \Re^{n+1} contains no closure points of \overline{M} . Hence by the Nonvertical Hyperplane Theorem [Prop. 1.5.8(b)], for any vector

 $(0, w) \in \Re^{n+1}$, there exists a nonvertical hyperplane strictly separating (0, w) and \overline{M} . The crossing point of this hyperplane with the vertical axis lies between w and q^* , so $w < q^*$ for all $w \in \Re$, which implies that $q^* = w^* = \infty$.

Consider finally the case where w^* is a real number. We first show by contradiction that for any $\epsilon > 0$, $(0, w^* - \epsilon) \notin \operatorname{cl}(\overline{M})$. If this were not so, i.e., $(0, w^* - \epsilon)$ is a closure point of \overline{M} for some $\epsilon > 0$, there would exist a sequence $\{(u_k, \overline{w}_k)\} \subset \overline{M}$ converging to $(0, w^* - \epsilon)$. In view of the definition of \overline{M} , this implies the existence of another sequence $\{(u_k, w_k)\} \subset M$ with $u_k \to 0$ and $w_k \leq \overline{w}_k$ for all k, so that

$$\liminf_{k \to \infty} w_k \le w^* - \epsilon,$$

which contradicts the assumption $w^* \leq \liminf_{k \to \infty} w_k$.

We next argue by contradiction that \overline{M} does not contain any vertical lines. If this were not so, by convexity of \overline{M} , the direction (0,-1) would be a direction of recession of $\operatorname{cl}(\overline{M})$. Because $(0,w^*)\in\operatorname{cl}(\overline{M})$ [since $(0,w^*)$ is the "infimum" common point of \overline{M} and the vertical axis], it follows that the entire halfline $\{(0,w^*-\epsilon)\mid\epsilon\geq 0\}$ belongs to $\operatorname{cl}(\overline{M})$ [cf. Prop. 1.4.1(a)], contradicting what was shown earlier.

Since \overline{M} does not contain any vertical lines and the vector $(0, w^* - \epsilon)$ does not belong to $\operatorname{cl}(\overline{M})$ for any $\epsilon > 0$, by Prop. 1.5.8(b), it follows that there exists a nonvertical hyperplane strictly separating $(0, w^* - \epsilon)$ and \overline{M} . This hyperplane crosses the (n+1)st axis at a unique vector $(0,\xi)$, which must lie between $(0,w^*-\epsilon)$ and $(0,w^*)$, i.e., $w^*-\epsilon < \xi \leq w^*$. Furthermore, ξ cannot exceed the optimal value q^* of the max crossing problem, which together with weak duality $(q^* \leq w^*)$, implies that

$$w^* - \epsilon < q^* \le w^*.$$

Since ϵ can be arbitrarily small, it follows that $q^* = w^*$.

Conversely, assume $q^* = w^*$ and consider a sequence $\{(u_k, w_k)\} \subset M$ such that $u_k \to 0$. We will show that $w^* \leq \liminf_{k \to \infty} w_k$. Indeed we have

$$q(\mu) = \inf_{(u,w) \in M} \{w + \mu' u\} \leq w_k + \mu' u_k, \qquad \forall \ k, \quad \forall \ \mu \in \Re^n,$$

and by taking the limit as $k \to \infty$, we obtain $q(\mu) \le \liminf_{k \to \infty} w_k$. Hence

$$w^* = q^* = \sup_{\mu \in \mathbb{R}^n} q(\mu) \le \liminf_{k \to \infty} w_k.$$

Q.E.D.

For an example where assumption (1) of the preceding proposition is violated and $q^* < w^*$, let M consist of a vertical line that does not pass through the origin. Then we have $q^* = -\infty$, $w^* = \infty$, and yet the

condition $w^* \leq \liminf_{k \to \infty} w_k$, for every sequence $\{(u_k, w_k)\} \subset M$ with $u_k \to 0$, trivially holds.

An important corollary of Prop. 4.3.1 is that if $M = \operatorname{epi}(p)$ where $p: \Re^n \mapsto [-\infty, \infty]$ is a convex function with $p(0) = w^* < \infty$, then we have $q^* = w^*$ if and only if p is lower semicontinuous at 0.

We will now consider cases where M (rather than \overline{M}) is convex, and where M has some compactness structure but may not be convex.

Proposition 4.3.2: Consider the MC/MC framework, assuming that $w^* < \infty$.

(a) Let M be closed and convex. Then $q^* = w^*$. Furthermore, the function

$$p(u) = \inf\{w \mid (u, w) \in M\}, \qquad u \in \Re^n,$$

is convex and its epigraph is the set

$$\overline{M} = M + \{(0, w) \mid w \ge 0\}.$$

If in addition $-\infty < w^*$, then p is closed and proper.

- (b) q^* is equal to the optimal value of the min common problem corresponding to cl(conv(M)).
- (c) If M is of the form

$$M = \tilde{M} + \{(u,0) \mid u \in C\},\$$

where \tilde{M} is a compact set and C is a closed convex set, then q^* is equal to the optimal value of the min common problem corresponding to $\operatorname{conv}(M)$.

Proof: (a) Since M is closed, for each $u \in \text{dom}(p)$, the infimum defining p(u) is either $-\infty$ or else it is attained. In view of the definition of \overline{M} , this implies that \overline{M} is the epigraph of p. Furthermore \overline{M} is convex, being the vector sum of two convex sets, so p is convex.

If $w^* = -\infty$, then $q^* = w^*$ by weak duality. It will suffice to show that if $w^* > -\infty$, then p is closed and proper, since by Prop. 4.3.1, this will also imply that $q^* = w^*$. We note that (0, -1) is not a direction of recession of M (since $w^* > -\infty$). Since \overline{M} is the vector sum of M and $\{(0, w) \mid w \geq 0\}$, this implies, using Prop. 1.4.14, that \overline{M} is closed as well as convex. Since $\overline{M} = \operatorname{epi}(p)$, it follows that p is closed. Since w^* is finite, p is also proper (an improper closed convex function cannot take

finite values; cf. the discussion at the end of Section 1.1.2).

- (b) The max crossing value q^* is the same for M and $\operatorname{cl}(\operatorname{conv}(M))$, since the closed halfspaces containing M are the ones that $\operatorname{contain} \operatorname{cl}(\operatorname{conv}(M))$ (cf. Prop. 1.5.4). Since $\operatorname{cl}(\operatorname{conv}(M))$ is closed and convex, the result follows from part (a).
- (c) The max crossing value is the same for M and $\operatorname{conv}(M)$, since the closed halfspaces containing M are the ones that $\operatorname{contain} \operatorname{conv}(M)$. It can be seen that

$$\operatorname{conv}(M) = \operatorname{conv}(\tilde{M}) + \{(u, 0) \mid u \in C\}.$$

Since \tilde{M} is compact, $\operatorname{conv}(\tilde{M})$ is also compact (cf. Prop. 1.2.2), so $\operatorname{conv}(M)$ is the vector sum of two closed convex sets one of which is compact. Hence, by Prop. 1.4.14, $\operatorname{conv}(M)$ is closed, and the result follows from part (a). Q.E.D.

The preceding proposition involves a few subtleties. First note that in part (a), if M is closed and convex, but $w^* = -\infty$, then p is convex, but it need not be closed. For an example in \Re^2 , consider the closed and convex set

$$M = \{(u, w) \mid w \le -1/(1 - |u|), |u| < 1\}.$$

Then.

$$\overline{M} = \{(u, w) \mid |u| < 1\},\$$

so \overline{M} is convex but not closed, implying that p (which is $-\infty$ for |u| < 1 and ∞ otherwise) is not closed.

Also note that if M is closed but violates the assumptions of Prop. 4.3.2(c), the min common values corresponding to conv(M) and cl(conv(M)) may differ. For an example in \Re^2 , consider the set

$$M = \{(0,0)\} \cup \{(u,w) \mid u > 0, w \le -1/u\}.$$

Then

$$\mathrm{conv}(M) = \big\{ (0,0) \} \cup \big\{ (u,w) \mid u > 0, \, w < 0 \big\}.$$

We have $q^* = -\infty$ but the min common value corresponding to $\operatorname{conv}(M)$ is $w^* = 0$, while the min common value corresponding to $\operatorname{cl}(\operatorname{conv}(M))$ is equal to $-\infty$ [consistent with part (b)].

4.4 EXISTENCE OF DUAL OPTIMAL SOLUTIONS

We now discuss the nonemptiness and the structure of the optimal solution set of the max crossing problem. The following proposition, in addition to the equality $q^* = w^*$, guarantees the attainment of the maximum crossing point under appropriate assumptions. Key among these assumptions is

that 0 is a relative interior point of the projection of M (or \overline{M}) on the horizontal plane [cf. Fig. 4.1.1(b)]. A subsequent proposition will show that for compactness (as well as nonemptiness) of the set of max crossing hyperplanes, it is necessary and sufficient that 0 be an interior point of the projection of M on the horizontal plane.

Proposition 4.4.1: (MC/MC Existence of Max Crossing Solutions) Consider the MC/MC framework and assume the following:

- $(1) -\infty < w^*$.
- (2) The set

$$\overline{M} = M + \{(0, w) \mid w \ge 0\}$$

is convex.

(3) The origin is a relative interior point of the set

$$D = \{u \mid \text{there exists } w \in \Re \text{ with } (u, w) \in \overline{M}\}.$$

Then $q^* = w^*$, and there exists at least one optimal solution of the max crossing problem.

Proof: Condition (3) implies that the vertical axis contains a point of M, so that $w^* < \infty$. Hence in view of condition (1), w^* is finite.

We note that $(0, w^*) \notin ri(\overline{M})$, since by Prop. 1.3.10,

$$\operatorname{ri}(\overline{M}) = \{(u, w) \mid u \in \operatorname{ri}(D), \overline{w} < w \text{ for some } (u, \overline{w}) \in \overline{M}\},\$$

and

$$w^* = \inf_{(0,\overline{w})\in\overline{M}} \overline{w}.$$

Therefore, by the Proper Separation Theorem (Prop. 1.5.5), there exists a hyperplane that passes through $(0, \underline{w}^*)$, contains \overline{M} in one of its closed halfspaces, but does not fully contain \overline{M} , i.e., there exists (μ, β) such that

$$\beta w^* \le \mu' u + \beta w, \qquad \forall \ (u, w) \in \overline{M}, \tag{4.22}$$

$$\beta w^* < \sup_{(u,w)\in \overline{M}} \{\mu' u + \beta w\}. \tag{4.23}$$

Since for any $(\overline{u}, \overline{w}) \in M$, the set \overline{M} contains the halfline $\{(\overline{u}, w) \mid \overline{w} \leq w\}$, it follows from Eq. (4.22) that $\beta \geq 0$. If $\beta = 0$, then from Eq. (4.22), we have

$$0 \le \mu' u, \quad \forall \ u \in D.$$

Thus, the linear function $\mu'u$ attains its minimum over the set D at 0, which is a relative interior point of D by condition (3). Since D is convex, being the projection on the space of u of the set \overline{M} , which is convex by condition (2), it follows from Prop. 1.3.4 that $\mu'u$ is constant over D, i.e.,

$$\mu'u = 0, \quad \forall \ u \in D.$$

This, however, contradicts Eq. (4.23). Therefore, we must have $\beta > 0$, and by appropriate normalization if necessary, we can assume that $\beta = 1$. From Eq. (4.22), we then obtain

$$w^* \le \inf_{(u,w) \in \overline{M}} \{\mu' u + w\} \le \inf_{(u,w) \in M} \{\mu' u + w\} = q(\mu) \le q^*.$$

Since $q^* \leq w^*$, by the Weak Duality Theorem (Prop. 4.1.2), equality holds throughout in the above relation, and we must have $q(\mu) = q^* = w^*$. Thus μ is an optimal solution of the max crossing problem. **Q.E.D.**

Note that if $w^* = -\infty$, by weak duality we have $q^* \le w^*$, so that $q^* = w^* = -\infty$. This means that $q(\mu) = -\infty$ for all $\mu \in \Re^n$, and that the dual problem is infeasible. The following proposition supplements the preceding one, and characterizes the optimal solution set of the max crossing problem.

Proposition 4.4.2: Let the assumptions of Prop. 4.4.1 hold. Then Q^* , the set of optimal solutions of the max crossing problem, has the form

$$Q^* = \left(\operatorname{aff}(D)\right)^{\perp} + \tilde{Q},$$

where \tilde{Q} is a nonempty, convex, and compact set. In particular, Q^* is compact if and only if the origin is an interior point of D.

Proof: By Prop. 4.4.1, q^* is finite and Q^* is nonempty. Since q is concave and upper semicontinuous (cf. Prop. 4.1.1), and $Q^* = \{\mu \mid q(\mu) \geq q^*\}$, it follows that Q^* is convex and closed. We will first show that the recession cone R_{Q^*} and the lineality space L_{Q^*} of Q^* are both equal to $\left(\operatorname{aff}(D)\right)^{\perp}$ [note here that $\operatorname{aff}(D)$ is a subspace since it contains the origin]. The proof of this is based on the generic relation $L_{Q^*} \subset R_{Q^*}$ and the following two relations

$$\left(\operatorname{aff}(D)\right)^{\perp} \subset L_{Q^*}, \qquad R_{Q^*} \subset \left(\operatorname{aff}(D)\right)^{\perp},$$

which we show next.

Let d be a vector in $\left(\operatorname{aff}(D)\right)^{\perp}$, so that d'u = 0 for all $u \in D$. For any vector $\mu \in Q^*$ and any scalar α , we then have

$$q(\mu+\alpha d)=\inf_{(u,w)\in\overline{M}}\bigl\{(\mu+\alpha d)'u+w\bigr\}=\inf_{(u,w)\in\overline{M}}\{\mu'u+w\}=q(\mu),$$

so that $\mu + \alpha d \in Q^*$. Hence $d \in L_{Q^*}$, and it follows that $\left(\operatorname{aff}(D)\right)^{\perp} \subset L_{Q^*}$. Let d be a vector in R_{Q^*} , so that for any $\mu \in Q^*$ and $\alpha \geq 0$,

$$q(\mu + \alpha d) = \inf_{(u,w) \in \overline{M}} \{ (\mu + \alpha d)' u + w \} = q^*.$$

Since $0 \in ri(D)$, for any $u \in aff(D)$, there exists a positive scalar γ such that the vectors γu and $-\gamma u$ are in D. By the definition of D, there exist scalars w^+ and w^- such that the pairs $(\gamma u, w^+)$ and $(-\gamma u, w^-)$ are in \overline{M} . Using the preceding equation, it follows that for any $\mu \in Q^*$, we have

$$(\mu + \alpha d)'(\gamma u) + w^{+} \ge q^{*}, \qquad \forall \ \alpha \ge 0,$$
$$(\mu + \alpha d)'(-\gamma u) + w^{-} > q^{*}, \qquad \forall \ \alpha > 0.$$

If $d'u \neq 0$, then for sufficiently large $\alpha \geq 0$, one of the preceding two relations will be violated. Thus we must have d'u = 0, showing that $d \in (\operatorname{aff}(D))^{\perp}$ and implying that

$$R_{Q^*} \subset \left(\operatorname{aff}(D)\right)^{\perp}$$
.

This relation, together with the generic relation $L_{Q^*} \subset R_{Q^*}$ and the relation $(\operatorname{aff}(D))^{\perp} \subset L_{Q^*}$ proved earlier, shows that

$$\left(\operatorname{aff}(D)\right)^{\perp} \subset L_{Q^*} \subset R_{Q^*} \subset \left(\operatorname{aff}(D)\right)^{\perp}.$$

Therefore

$$L_{Q^*} = R_{Q^*} = \left(\operatorname{aff}(D)\right)^{\perp}.$$

We now use the decomposition result of Prop. 1.4.4, to assert that

$$Q^* = L_{Q^*} + (Q^* \cap L_{Q^*}^{\perp}).$$

Since $L_{Q^*} = (\operatorname{aff}(D))^{\perp}$, we obtain

$$Q^* = \left(\operatorname{aff}(D)\right)^{\perp} + \tilde{Q},$$

where $\tilde{Q} = Q^* \cap \operatorname{aff}(D)$. Furthermore, by Prop. 1.4.2(c), we have

$$R_{\tilde{Q}} = R_{Q^*} \cap R_{\mathrm{aff}(D)}.$$

Since $R_{Q^*} = (\operatorname{aff}(D))^{\perp}$, as shown earlier, and $R_{\operatorname{aff}(D)} = \operatorname{aff}(D)$, the recession cone $R_{\tilde{Q}}$ consists of the zero vector only, implying that the set \tilde{Q} is compact.

From the formula $Q^* = (\operatorname{aff}(D))^{\perp} + \tilde{Q}$, it follows that Q^* is compact if and only if $(\operatorname{aff}(D))^{\perp} = \{0\}$, or equivalently $\operatorname{aff}(D) = \Re^n$. Since 0 is a relative interior point of D by assumption, this is equivalent to 0 being an interior point of D. **Q.E.D.**

4.5 DUALITY AND POLYHEDRAL CONVEXITY

We now strengthen the result on the existence of max crossing solutions (Prop. 4.4.1) for a special case where the upwards extension

$$\overline{M} = M + \{(0, w) \mid w \ge 0\}$$

of the set M has partially polyhedral structure. In particular, we will consider the case where \overline{M} is a vector sum of the form

$$\overline{M} = \tilde{M} - \{(u,0) \mid u \in P\},$$
 (4.24)

where \tilde{M} is convex and P is polyhedral. Then the corresponding set

$$D = \{u \mid \text{there exists } w \in \Re \text{ with } (u, w) \in \overline{M}\},$$

can be written as

$$D = \tilde{D} - P$$

where

$$\tilde{D} = \{ u \mid \text{there exists } w \in \Re \text{ with } (u, w) \in \tilde{M} \}.$$
 (4.25)

To understand the nature of the following proposition, we note that from Props. 4.4.1 and 4.4.2, assuming that $-\infty < w^*$, we have:

(a) $q^* = w^*$ and Q^* is nonempty, provided that $0 \in \text{ri}(D)$, which is equivalent to

$$\operatorname{ri}(\tilde{D}) \cap \operatorname{ri}(P) \neq \emptyset,$$
 (4.26)

since, by Prop. 1.3.7, $ri(D) = ri(\tilde{D}) - ri(P)$.

(b) $q^* = w^*$, and Q^* is nonempty and compact, provided that $0 \in \text{int}(D)$, which is true in particular if either

$$\operatorname{int}(\tilde{D}) \cap P \neq \emptyset$$
,

or

$$\tilde{D} \cap \operatorname{int}(P) \neq \emptyset$$
.

The following proposition shows in part that when P is polyhedral, these results can be strengthened, and in particular, the condition $\operatorname{ri}(\tilde{D}) \cap \operatorname{ri}(P) \neq \emptyset$ [Eq. (4.26)] can be replaced by the condition

$$ri(\tilde{D}) \cap P \neq \emptyset$$
.

The proof is similar to the proofs of Props. 4.4.1 and 4.4.2, but uses the Polyhedral Proper Separation Theorem (Prop. 1.5.7) to exploit the polyhedral structure of P.

Proposition 4.5.1: Consider the MC/MC framework, and assume the following:

- $(1) -\infty < w^*$.
- (2) The set \overline{M} has the form

$$\overline{M} = \tilde{M} - \{(u,0) \mid u \in P\},\$$

where \tilde{M} and P are convex sets.

(3) Either $\operatorname{ri}(\tilde{D}) \cap \operatorname{ri}(P) \neq \emptyset$, or P is polyhedral and $\operatorname{ri}(\tilde{D}) \cap P \neq \emptyset$, where \tilde{D} is the set given by Eq. (4.25).

Then $q^* = w^*$, and Q^* , the set of optimal solutions of the max crossing problem, is a nonempty subset of R_P^* , the polar cone of the recession cone of P. Furthermore, Q^* is compact if $\operatorname{int}(\tilde{D}) \cap P \neq \emptyset$.

Proof: We assume that P is polyhedral and $ri(\tilde{D}) \cap P \neq \emptyset$. The proof for the case where P is just convex and $ri(\tilde{D}) \cap ri(P) \neq \emptyset$ is similar: we just use the Proper Separation Theorem (Prop. 1.5.6) in place of the Polyhedral Proper Separation Theorem (Prop. 1.5.7) in the following argument. We consider the sets

$$C_1 = \{(u, v) \mid v > w \text{ for some } (u, w) \in \tilde{M}\},$$

$$C_2 = \{(u, w^*) \mid u \in P\},\$$

(cf. Fig. 4.5.1).

It can be seen that C_1 and C_2 are nonempty and convex, and C_2 is polyhedral. Furthermore, C_1 and C_2 are disjoint. To see this, note that if \overline{u} is such that $\overline{u} \in P$ and there exists (\overline{u}, w) with $w^* > w$, then $(0, w) \in M$, which is impossible since w^* is the min common value. Therefore, by Prop. 1.5.7, there exists a hyperplane that separates C_1 and C_2 , and does not contain C_1 , i.e., a vector $(\overline{\mu}, \beta)$ such that

$$\beta w^* + \overline{\mu}' z \le \beta v + \overline{\mu}' u, \qquad \forall (u, v) \in C_1, \ \forall \ z \in P,$$
 (4.27)

$$\inf_{(u,v)\in C_1} \left\{ \beta v + \overline{\mu}' u \right\} < \sup_{(u,v)\in C_1} \left\{ \beta v + \overline{\mu}' u \right\}. \tag{4.28}$$

From Eq. (4.27), since (0,1) is a direction of recession of C_1 , we see that $\beta \geq 0$. If $\beta = 0$, then for a vector $\overline{u} \in ri(D) \cap P$, Eq. (4.27) yields

$$\overline{\mu}'\overline{u} \le \inf_{u \in \tilde{D}} \overline{\mu}'u,$$

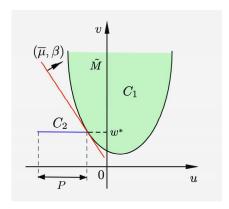


Figure 4.5.1. Illustration of the sets

$$C_1 = \{(u, v) \mid v > w \text{ for some } (u, w) \in \tilde{M}\}, \qquad C_2 = \{(u, w^*) \mid u \in P\},$$

and the hyperplane separating them in the proof of Prop. 4.5.1.

so \overline{u} attains the minimum of the linear function $\overline{\mu}'u$ over \tilde{D} , and from Prop. 1.3.4, it follows that $\overline{\mu}'u$ is constant over \tilde{D} . On the other hand, from Eq. (4.28) we have $\inf_{u\in \tilde{D}}\overline{\mu}'u < \sup_{u\in \tilde{D}}\overline{\mu}'u$, a contradiction. Hence $\beta>0$, and by normalizing $(\overline{\mu},\beta)$ if necessary, we may assume that $\beta=1$.

Thus, from Eq. (4.27), we have

$$w^* + \overline{\mu}'z \le \inf_{(u,v) \in C_1} \{v + \overline{\mu}'u\}, \qquad \forall \ z \in P, \tag{4.29}$$

which in particular implies that $\overline{\mu}'d \leq 0$ for all $d \in R_P$. Hence $\overline{\mu} \in R_P^*$. From Eq. (4.29), we also obtain

$$\begin{split} w^* &\leq \inf_{(u,v) \in C_1, \, z \in P} \left\{ v + \overline{\mu}'(u-z) \right\} \\ &= \inf_{(u,v) \in \tilde{M} - \left\{ (z,0) | z \in P \right\}} \left\{ v + \overline{\mu}'u \right\} \\ &= \inf_{(u,v) \in \overline{M}} \left\{ v + \overline{\mu}'u \right\} \\ &= q(\overline{\mu}). \end{split}$$

Using the weak duality relation $q^* \leq w^*$, we have $q(\overline{\mu}) = q^* = w^*$. To show that $Q^* \subset R_P^*$, note that for all μ , we have

$$q(\mu) = \inf_{(u,w) \in \overline{M}} \{w + \mu' u\} = \inf_{(u,w) \in \tilde{M}, \, z \in P} \{w + \mu' (u - z)\},$$

so $q(\mu) = -\infty$ if $\mu'd > 0$ for some $d \in R_P$. Hence if $\mu \in Q^*$, we must have $\mu'd \leq 0$ for all $d \in R_P$, or $\mu \in R_P^*$.

The proof that Q^* is compact if $\operatorname{int}(\tilde{D}) \cap P \neq \emptyset$ is similar to the one of Prop. 4.4.2 (see the discussion preceding the proposition). **Q.E.D.**

We now discuss an interesting special case of Prop. 4.5.1. It corresponds to \tilde{M} being a linearly transformed epigraph of a convex function f, and P being a polyhedral set such as the nonpositive orthant. This special case applies to constrained optimization duality, and will be used in the proof of the Nonlinear Farkas' Lemma given in Section 5.1.

Proposition 4.5.2: Consider the MC/MC framework, and assume that:

- $(1) -\infty < w^*.$
- (2) The set \overline{M} is defined in terms of a polyhedral set P, an $r \times n$ matrix A, a vector $b \in \Re^r$, and a convex function $f : \Re^n \mapsto (-\infty, \infty]$ as follows:

$$\overline{M} = \{(u, w) \mid Ax - b - u \in P \text{ for some } (x, w) \in \operatorname{epi}(f)\}.$$

(3) There is a vector $\overline{x} \in \text{ri}(\text{dom}(f))$ such that $A\overline{x} - b \in P$.

Then $q^* = w^*$ and Q^* , the set of optimal solutions of the max crossing problem, is a nonempty subset of R_P^* , the polar cone of the recession cone of P. Furthermore, Q^* is compact if the matrix A has rank r and there is a vector $\overline{x} \in \text{int}(\text{dom}(f))$ such that $A\overline{x} - b \in P$.

Proof: Let

$$\tilde{M} = \{ (Ax - b, w) \mid (x, w) \in \operatorname{epi}(f) \}.$$

The following calculation relates \tilde{M} and \overline{M} :

$$\begin{split} \tilde{M} - \left\{ (z,0) \mid z \in P \right\} \\ &= \left\{ (u,w) \mid u = Ax - b - z \text{ for some } (x,w) \in \operatorname{epi}(f) \text{ and } z \in P \right\} \\ &= \left\{ (u,w) \mid Ax - b - u \in P \text{ for some } (x,w) \in \operatorname{epi}(f) \right\} \\ &= \overline{M}. \end{split}$$

Thus the framework of Prop. 4.5.1 applies. Furthermore, the set \tilde{D} of Eq. (4.25) is given by

$$\tilde{D} = \big\{ u \mid \text{there exists } w \in \Re \text{ with } (u, w) \in \tilde{M} \big\} = \big\{ Ax - b \mid x \in \text{dom}(f) \big\}.$$

Hence, the relative interior assumption (3) implies that the corresponding relative interior assumption of Prop. 4.5.1 is satisfied. Furthermore, if A

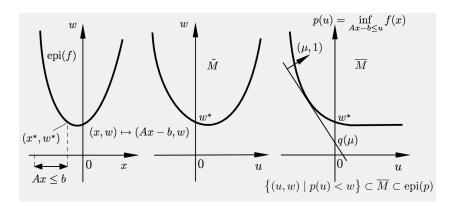


Figure 4.5.2. The MC/MC framework for minimizing f(x) subject to $Ax \leq b$. The set \overline{M} is

$$\overline{M} = \left\{ (u, w) \mid Ax - b \le u \text{ for some } (x, w) \in \operatorname{epi}(f) \right\}$$

(cf. Prop. 4.5.2) and can be written as a vector sum:

$$\overline{M} = \tilde{M} + \{(u,0) \mid u \ge 0\},\$$

where \tilde{M} is obtained by transformation of epi(f),

$$\tilde{M} = \left\{ (Ax - b, w) \mid (x, w) \in \operatorname{epi}(f) \right\}.$$

Also \overline{M} is related to the perturbation function $p(u) = \inf_{Ax-b < u} f(x)$ as follows:

$$\{(u,w) \mid p(u) < w\} \subset \overline{M} \subset \operatorname{epi}(p).$$

In particular, we have $w^* = p(0) = \inf_{Ax < b} f(x)$.

has rank r, we have $A\overline{x} - b \in \operatorname{int}(\tilde{D})$ for all $\overline{x} \in \operatorname{int}(\operatorname{dom}(f))$. The result follows from the conclusion of Prop. 4.5.1. Q.E.D.

The special case where P is the negative orthant $P=\{u\mid u\leq 0\}$, is noteworthy as it corresponds to a major convex optimization model. To see this, consider the minimization of a convex function $f:\Re^n\mapsto (-\infty,\infty]$ subject to $Ax\leq b$, and note that the set \overline{M} in Prop. 4.5.2 is related to the epigraph of the perturbation function

$$p(u) = \inf_{Ax - b \le u} f(x),$$

(cf. Fig. 4.5.2). The min common value is equal to p(0), the optimal value:

$$w^* = p(0) = \inf_{Ax \le b} f(x)$$

(cf. Section 4.2.3). The max crossing problem is to maximize over $\mu \in \mathbb{R}^r$ the function q given by

$$q(\mu) = \begin{cases} \inf_{x \in \Re^n} \left\{ f(x) + \mu'(Ax - b) \right\} & \text{if } \mu \ge 0, \\ -\infty & \text{otherwise,} \end{cases}$$

[cf. Eq. (4.11)], and is obtained by maximization of the Lagrangian function $f(x) + \mu'(Ax - b)$ when $\mu \geq 0$. Proposition 4.5.2 provides the principal duality result for this problem, namely that assuming the existence of $\overline{x} \in \text{ri}(\text{dom}(f))$ such that $A\overline{x} \leq b$, strong duality holds and there exists a dual optimal solution.

4.6 SUMMARY

Several propositions regarding the structure of the MC/MC framework were given in this chapter, and it may be helpful to discuss how they are used in specific duality contexts. These are:

- (1) Propositions 4.3.1 and 4.3.2, which assert the strong duality property $q^* = w^*$ under convexity of \overline{M} , and various conditions on M that involve a "lower semicontinuity structure," and/or closure and compactness. Note that these propositions say nothing about existence of optimal primal or dual solutions.
- (2) Propositions 4.4.1-4.5.2, which under various relative interior conditions, assert (in addition to $q^* = w^*$) that the max crossing solution set is nonempty and describe some of its properties, such as the structure of its recession cone and/or its compactness.

Proposition 4.3.1 is quite general and will be used to assert strong duality in Section 5.3.4 (convex programming without relative interior conditions - Prop. 5.3.7), and Section 5.5.1 (minimax problems - Prop. 5.5.1). Proposition 4.3.2 is more specialized, and assumes some structure for the set M (rather than \overline{M}). It is helpful in providing perspective, and it is used in Section 5.7 for analysis of the size of the duality gap.

Propositions 4.4.1 and 4.4.2 are also quite general, and are used throughout Chapter 5. For example, they are used to prove part (a) of the Nonlinear Farkas' Lemma (Prop. 5.1.1) on which much of convex programming duality rests (Section 5.3). Propositions 4.5.1 and 4.5.2 are more specialized, but still broadly applicable. They assume a polyhedral structure that is relevant to duality theory for problems involving linear constraints (cf. the problem discussed at the end of Section 4.5). In particular, they are used to prove part (b) of the Nonlinear Farkas' Lemma (and by extension the Fenchel and conic duality theorems of Section 5.3), as well as various results in subdifferential calculus (Section 5.4) and (implicitly) in theorems of the alternative (Section 5.6).

Duality and Optimization

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In this chapter we develop many of the basic analytical results of constrained optimization and minimax theory, including duality, and optimality conditions. We also develop the core of subdifferential theory, as well as theorems of the alternative. These lead to additional optimization-related results, such as subgradient-based optimality conditions, sensitivity interpretations of dual optimal solutions, and conditions for compactness of the solution set of a linear program. We finally consider nonconvex optimization and minimax problems, and discuss estimates of the duality gap. The MC/MC duality results of the preceding chapter are the principal tools.

5.1 NONLINEAR FARKAS' LEMMA

We will first prove a nonlinear version of Farkas' Lemma that captures the essence of convex programming duality. The lemma involves a nonempty convex set $X \subset \Re^n$, and functions $f: X \mapsto \Re$ and $g_j: X \mapsto \Re$, $j = 1, \ldots, r$. We denote $g(x) = (g_1(x), \ldots, g_r(x))'$, and assume the following.

Assumption 5.1.1: The functions f and g_j , j = 1, ..., r, are convex, and

$$f(x) \ge 0, \quad \forall x \in X \text{ with } g(x) \le 0.$$

The Nonlinear Farkas' Lemma asserts, under some conditions, that there exists a nonvertical hyperplane that passes through the origin and contains the set

$$\left\{ \left(g(x),f(x)\right)\mid x\in X\right\}$$

in its positive halfspace. Figure 5.1.1 provides a geometric interpretation and suggests the strong connection with the MC/MC framework.

Proposition 5.1.1: (Nonlinear Farkas' Lemma) Let Assumption 5.1.1 hold and let Q^* be the subset of \Re^r given by

$$Q^* = \{ \mu \mid \mu \ge 0, f(x) + \mu' g(x) \ge 0, \forall x \in X \}.$$

Assume that one of the following two conditions holds:

- (1) There exists $\overline{x} \in X$ such that $g_i(\overline{x}) < 0$ for all j = 1, ..., r.
- (2) The functions g_j , $j=1,\ldots,r$, are affine, and there exists $\overline{x} \in ri(X)$ such that $g(\overline{x}) \leq 0$.

Then Q^* is nonempty, and under condition (1) it is also compact.

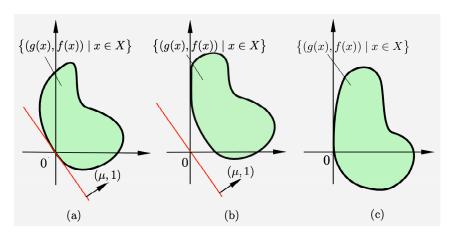


Figure 5.1.1. Geometrical interpretation of the Nonlinear Farkas' Lemma. Assuming that $f(x) \geq 0$ for all $x \in X$ with $g(x) \leq 0$, the lemma asserts the existence of a nonvertical hyperplane in \Re^{r+1} , with normal $(\mu, 1)$, that passes through the origin and contains the set

$$\left\{ \left(g(x), f(x) \right) \mid x \in X \right\}$$

in its positive halfspace. Figures (a) and (b) show examples where such a hyperplane exists, and figure (c) shows an example where it does not. In Fig. (a) there exists a point $\overline{x} \in X$ with $g(\overline{x}) < 0$.

Proof: Let condition (1) hold. We consider the MC/MC framework corresponding to the subset of \Re^{r+1} given by

$$M = \{(u, w) \mid \text{there exists } x \in X \text{ such that } g(x) \le u, \ f(x) \le w\}$$

(cf. Fig. 5.1.2). We will apply Props. 4.4.1 and 4.4.2 to assert that the set of max crossing hyperplanes is nonempty and compact. To this end, we will verify that the assumptions of these propositions are satisfied. In particular, we will show that:

(i) The optimal value w^* of the corresponding min common problem,

$$w^* = \inf\{w \mid (0, w) \in M\},\$$

satisfies $-\infty < w^*$.

(ii) The set

$$\overline{M} = M + \{(0, w) \mid w \ge 0\},\$$

is convex. (Note here that $\overline{M} = M$.)

(iii) The set

$$D = \{u \mid \text{there exists } w \in \Re \text{ such that } (u, w) \in \overline{M} \}$$

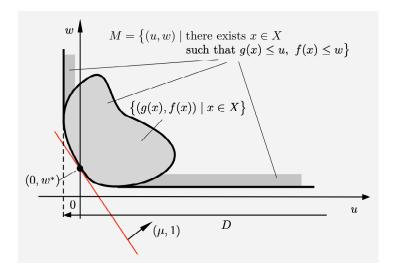


Figure 5.1.2. Illustration of the sets used in the MC/MC framework of the proof of Prop. 5.1.1 under condition (1). We have

$$M = \overline{M} = \{(u, w) \mid \text{there exists } x \in X \text{ such that } g(x) \le u, \ f(x) \le w \}$$

and

$$\begin{split} D &= \left\{ u \mid \text{ there exists } w \in \Re \text{ such that } (u,w) \in \overline{M} \right\} \\ &= \left\{ u \mid \text{ there exists } x \in X \text{ such that } g(x) \leq u \right\}. \end{split}$$

The existence of $\overline{x} \in X$ such that $g(\overline{x}) < 0$ is equivalent to $0 \in \operatorname{int}(D)$. Note that M is convex, even though the set $\left\{\left(g(x), f(x)\right) \mid x \in X\right\}$ need not be [take for example $X = \Re, \ f(x) = x, \ g(x) = x^2$].

contains the origin in its interior.

To show (i), note that since $f(x) \ge 0$ for all $x \in X$ with $g(x) \le 0$, we have $w \ge 0$ for all $(0, w) \in M$, so that $w^* \ge 0$.

To show (iii), note that D can also be written as

$$D = \big\{ u \mid \text{ there exists } x \in X \text{ such that } g(x) \leq u \big\}.$$

Since D contains the set $g(\overline{x}) + \{u \mid u \geq 0\}$, the condition $g(\overline{x}) < 0$ for some $\overline{x} \in X$, implies that $0 \in \operatorname{int}(D)$.

There remains to show (ii), i.e., that the set \overline{M} is convex. Since $\overline{M} = M$, we will prove that M is convex. To this end, we consider vectors $(u, w) \in M$ and $(\tilde{u}, \tilde{w}) \in M$, and we show that their convex combinations lie in M. By the definition of M, for some $x \in X$ and $\tilde{x} \in X$, we have

$$f(x) \le w$$
, $g_j(x) \le u_j$, $\forall j = 1, \dots, r$,

$$f(\tilde{x}) \leq \tilde{w}, \quad g_j(\tilde{x}) \leq \tilde{u}_j, \quad \forall j = 1, \dots, r.$$

For any $\alpha \in [0, 1]$, we multiply these relations with α and $1-\alpha$, respectively, and add them. By using the convexity of f and g_j for all j, we obtain

$$f(\alpha x + (1 - \alpha)\tilde{x}) < \alpha f(x) + (1 - \alpha)f(\tilde{x}) < \alpha w + (1 - \alpha)\tilde{w},$$

$$g_j(\alpha x + (1-\alpha)\tilde{x}) \le \alpha g_j(x) + (1-\alpha)g_j(\tilde{x}) \le \alpha u_j + (1-\alpha)\tilde{u}_j, \quad \forall j = 1, \dots, r.$$

By convexity of X, we have $\alpha x + (1 - \alpha)\tilde{x} \in X$ for all $\alpha \in [0, 1]$, so the preceding inequalities imply that the convex combination $\alpha(u, w) + (1 - \alpha)(\tilde{u}, \tilde{w})$ belongs to M, showing that M is convex.

Thus our assumptions imply that all the assumptions of Props. 4.4.1 and 4.4.2 hold. From Prop. 4.4.1, we obtain $w^* = q^* = \sup_{\mu} q(\mu)$, where q is the dual function,

$$q(\mu) = \inf_{(u,w) \in M} \{w + \mu'u\} = \begin{cases} \inf_{x \in X} \left\{ f(x) + \mu'g(x) \right\} & \text{if } \mu \geq 0, \\ -\infty & \text{otherwise.} \end{cases}$$

From Prop. 4.4.2, the optimal solution set $\tilde{Q} = \{ \mu \mid q(\mu) \geq w^* \}$ is nonempty and compact. Furthermore, according to the definition of Q^* , we have

$$Q^* = \big\{ \mu \mid \mu \geq 0, \, f(x) + \mu' g(x) \geq 0, \, \forall \, \, x \in X \big\} = \big\{ \mu \mid q(\mu) \geq 0 \big\}.$$

Thus Q^* and \tilde{Q} are level sets of the closed proper convex function -q, with $Q^* \supset \tilde{Q}$ (in view of $w^* \geq 0$). Since \tilde{Q} is nonempty and compact, so is Q^* [cf. Prop. 1.4.5(b)].

Let condition (2) hold, and let the constraint $g(x) \leq 0$ be written as

$$Ax - b \le 0$$
,

where A is an $r \times n$ matrix and b is a vector in \Re^n . We introduce a MC/MC framework with polyhedral structure and apply Prop. 4.5.2, with P being the nonpositive orthant and the set \overline{M} defined by

$$\overline{M} = \big\{(u,w) \mid Ax - b - u \leq 0, \text{ for some } (x,w) \in \operatorname{epi}(\tilde{f})\big\},\$$

where

$$\tilde{f}(x) = \begin{cases} f(x) & \text{if } x \in X, \\ \infty & \text{if } x \notin X. \end{cases}$$

Since by assumption, $f(x) \ge 0$ for all $x \in X$ with $Ax - b \le 0$, the min common value satisfies

$$w^* = \inf_{Ax - b \le 0} \tilde{f}(x) \ge 0.$$

By Prop. 4.5.2 and the discussion following its proof, there exists $\mu \geq 0$ such that

$$q^* = q(\mu) = \inf_{x \in \Re^n} \{ \tilde{f}(x) + \mu'(Ax - b) \}.$$

Since $q^* = w^* \ge 0$, it follows that $\tilde{f}(x) + \mu'(Ax - b) \ge 0$ for all $x \in \Re^n$, or $f(x) + \mu'(Ax - b) \ge 0$ for all $x \in X$, so $\mu \in Q^*$. Q.E.D.

By selecting f and g_j to be linear, and X to be the entire space in the Nonlinear Farkas' Lemma, we obtain a version of Farkas' Lemma (cf. Section 2.3.1) as a special case.

Proposition 5.1.2: (Linear Farkas' Lemma) Let A be an $m \times n$ matrix and c be a vector in \Re^m .

(a) The system Ay = c, $y \ge 0$ has a solution if and only if

$$A'x \le 0 \qquad \Rightarrow \qquad c'x \le 0. \tag{5.1}$$

(b) The system $Ay \geq c$ has a solution if and only if

$$A'x = 0, x \ge 0 \Rightarrow c'x \le 0.$$

Proof: (a) If $y \in \mathbb{R}^n$ is such that Ay = c, $y \geq 0$, then y'A'x = c'x for all $x \in \mathbb{R}^m$, which implies Eq. (5.1). Conversely, let us apply the Nonlinear Farkas' Lemma under condition (2) with f(x) = -c'x, g(x) = A'x, and $X = \mathbb{R}^m$. We see that the relation (5.1) implies the existence of $\mu \geq 0$ such that

$$-c'x + \mu'A'x \ge 0, \qquad \forall \ x \in \Re^m,$$

or equivalently $(A\mu - c)'x \ge 0$ for all $x \in \Re^m$, or $A\mu = c$.

(b) This part follows by writing the system $Ay \geq c$ in the equivalent form

$$Ay^+ - Ay^- - z = c,$$
 $y^+ \ge 0, y^- \ge 0, z \ge 0,$

and by applying part (a). Q.E.D.

5.2 LINEAR PROGRAMMING DUALITY

We will now derive one of the most important results in optimization: the linear programming duality theorem. Consider the problem

minimize
$$c'x$$

subject to $a'_j x \ge b_j$, $j = 1, ..., r$,

where $c \in \mathbb{R}^n$, $a_j \in \mathbb{R}^n$, and $b_j \in \mathbb{R}$, j = 1, ..., r. We refer to this as the primal problem. We consider the dual problem

maximize
$$b'\mu$$

subject to $\sum_{j=1}^{r} a_j \mu_j = c, \quad \mu \ge 0,$

which was derived from the MC/MC duality framework in Section 4.2.3. We denote the primal and dual optimal values by f^* and q^* , respectively.

We will show that $f^* = q^*$, assuming that the problem is feasible. While we may argue in terms of the MC/MC framework and the duality theory of Chapter 4, we will use instead the Linear Farkas' Lemma (Prop. 5.1.2), which captures the essential argument. The first step in the proof is to assert that weak duality holds, i.e., $q^* \leq f^*$ (cf. Prop. 4.1.2). [We can also prove this relation with a simple argument: if x and μ are feasible solutions of the primal and dual problems, we have $b'\mu \leq c'x$. Indeed,

$$b'\mu = \sum_{j=1}^{r} b_j \mu_j + \left(c - \sum_{j=1}^{r} a_j \mu_j\right)' x = c'x + \sum_{j=1}^{r} \mu_j (b_j - a_j' x) \le c'x, \quad (5.2)$$

where the inequality follows from the feasibility of x and μ . By taking supremum of the left-hand side over all feasible μ and infimum of the right-hand side over all feasible x, we obtain $q^* \leq f^*$.] The second step is to demonstrate the existence of feasible vectors x^* and μ^* such that $b'\mu^* = c'x^*$, assuming f^* or q^* is finite. This will follow from Farkas' Lemma, as geometrically illustrated in Fig. 5.2.1.

Proposition 5.2.1: (Linear Programming Duality Theorem)

- (a) If either f^* or q^* is finite, then $f^* = q^*$ and both the primal and the dual problem have optimal solutions.
- (b) If $f^* = -\infty$, then $q^* = -\infty$.
- (c) If $q^* = \infty$, then $f^* = \infty$.

Proof: (a) If f^* is finite, by Prop. 1.4.12, there exists a primal optimal solution x^* . Let

$$J = \{ j \in \{1, \dots, r\} \mid a_j' x^* = b_j \}.$$

We claim that

$$c'y \ge 0, \quad \forall y \text{ such that } a'_j y \ge 0 \text{ for all } j \in J.$$
 (5.3)

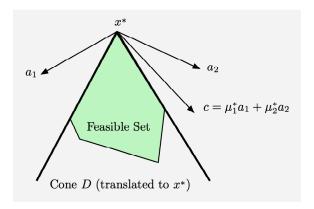


Figure 5.2.1. Illustration of the use of Farkas' Lemma to prove the linear programming duality theorem. Let x^* be a primal optimal solution, and let $J=\{j\mid a_j'x^*=b_j\}.$ Then, we have $c'y\geq 0$ for all y in the cone of "feasible directions"

$$D = \{ y \mid a_i' y \ge 0, \, \forall \ j \in J \}$$

(cf. Prop. 1.1.8). By Farkas' Lemma [Prop. 5.1.2(a)], this implies that c can be expressed in terms of some scalars $\mu_i^* \geq 0$ as

$$c = \sum_{j=1}^{r} \mu_{j}^{*} a_{j}, \qquad \mu_{j}^{*} \ge 0, \ \forall \ j \in J, \qquad \mu_{j}^{*} = 0, \ \forall \ j \notin J.$$

Taking inner product with x^* , we obtain $c'x^* = b'\mu^*$, which in view of $q^* \leq f^*$, shows that $q^* = f^*$ and that μ^* is optimal.

Indeed, each y such that $a'_i y \ge 0$ for all $j \in J$ is a feasible direction at x^* , in the sense that $x^* + \alpha y$ is a feasible point for all $\alpha > 0$ that are sufficiently small. It follows that the inequality c'y < 0 would violate the optimality of x^* , thus showing Eq. (5.3).

By Farkas' Lemma [Prop. 5.1.2(a)], Eq. (5.3) implies that

$$c = \sum_{j=1}^{r} \mu_j^* a_j,$$

for some $\mu^* \in \Re^r$ with

$$\mu_j^* \ge 0, \ \forall \ j \in J, \qquad \mu_j^* = 0, \ \forall \ j \notin J.$$

Thus, μ^* is a dual-feasible solution, and taking inner product with x^* , and using the fact $a_j'x^*=b_j$ for $j\in J$, and $\mu_j^*=0$ for $j\notin J$, we obtain

$$c'x^* = \sum_{j=1}^r \mu_j^* a_j' x^* = \sum_{j=1}^r \mu_j^* b_j = b'\mu^*.$$

This, together with Eq. (5.2), implies that $q^* = f^*$ and that μ^* is optimal. If q^* is finite, again by Prop. 1.4.12, there exists a dual optimal solution μ^* , and a similar argument proves the result.

- (b) If $f^* = -\infty$, the inequality $q^* < f^*$ implies that $q^* = -\infty$.
- (c) If $q^* = \infty$, the inequality $q^* \leq f^*$ implies that $f^* = \infty$. Q.E.D.

The one possibility left open by the proposition is to have $f^* = \infty$ and $q^* = -\infty$ (both primal and dual problems are infeasible). It turns out that this is possible: a trivial example is the infeasible scalar problem $\min_{0:x\geq 1} x$ whose dual is the infeasible scalar problem $\max_{0:\mu=1,\,\mu\geq 0} \mu$.

Another result, related to the duality theorem, is the following necessary and sufficient condition for primal and dual optimality.

Proposition 5.2.2: (Linear Programming Optimality Conditions) A pair of vectors (x^*, μ^*) form a primal and dual optimal solution pair if and only if x^* is primal-feasible, μ^* is dual-feasible, and

$$\mu_j^*(b_j - a_j'x^*) = 0, \quad \forall \ j = 1, \dots, r.$$
 (5.4)

Proof: If x^* is primal-feasible and μ^* is dual-feasible, then [cf. Eq. (5.2)]

$$b'\mu^* = \sum_{j=1}^r b_j \mu_j^* + \left(c - \sum_{j=1}^r a_j \mu_j^*\right)' x^* = c'x^* + \sum_{j=1}^r \mu_j^* (b_j - a_j' x^*). \quad (5.5)$$

Thus, if Eq. (5.4) holds, we have $b'\mu^* = c'x^*$, so Eq. (5.2) implies that x^* is a primal optimal solution and μ^* is a dual optimal solution.

Conversely, if (x^*, μ^*) form a primal and dual optimal solution pair, then x^* is primal-feasible, μ^* is dual-feasible, and by the duality theorem [Prop. 5.2.1(a)], we have $b'\mu^* = c'x^*$. From Eq. (5.5), we obtain Eq. (5.4). **Q.E.D.**

The condition (5.4) is known as *complementary slackness*. For a primal-feasible vector x^* and a dual-feasible vector μ^* , complementary slackness can be written in several equivalent ways, such as:

$$a'_j x^* > b_j \Rightarrow \mu_j^* = 0, \quad \forall j = 1, \dots, r,$$

 $\mu_j^* > 0 \Rightarrow a'_j x^* = b_j, \quad \forall j = 1, \dots, r.$

5.3 CONVEX PROGRAMMING DUALITY

We will now derive duality results and optimality conditions for convex programming, i.e., problems with convex cost function and constraints.

5.3.1 Strong Duality Theorem – Inequality Constraints

We first focus on the problem

minimize
$$f(x)$$

subject to $x \in X$, $q(x) < 0$, (5.6)

We refer to this as the *primal problem*. The duality results of this section assume that the problem is feasible, i.e., $f^* < \infty$, where f^* is the optimal value:

$$f^* = \inf_{x \in X, g(x) \le 0} f(x).$$

We will explore the connection with the dual problem derived from the MC/MC framework of Section 4.2.3, by using the Nonlinear Farkas' Lemma (Prop. 5.1.1) as our primary analytical tool.

Consider the Lagrangian function

$$L(x,\mu) = f(x) + \mu' g(x), \qquad x \in X, \ \mu \in \Re^r,$$

and the function q given by

$$q(\mu) = \begin{cases} \inf_{x \in X} L(x, \mu) & \text{if } \mu \ge 0, \\ -\infty & \text{otherwise.} \end{cases}$$

We refer to q as the *dual function* and to the problem

maximize
$$q(\mu)$$

subject to $\mu \in \Re^r$,

as the dual problem. The dual optimal value is

$$q^* = \sup_{\mu \in \Re^r} q(\mu).$$

Note the weak duality relation $q^* \leq f^*$ (cf. Prop. 4.1.2).

To investigate the question of existence of optimal solutions of the dual problem, we may use the MC/MC general results of Sections 4.3-4.5. However, for convex programming, the essential nature of these results is captured in the Nonlinear Farkas' Lemma (Prop. 5.1.1), which we will use instead in our analysis. To see the connection, assume that the optimal value f^* is finite. Then, we have

$$0 \le f(x) - f^*, \quad \forall \ x \in X \text{ with } g(x) \le 0,$$

so by replacing f(x) by $f(x) - f^*$ and by applying the Nonlinear Farkas' Lemma (assuming that one of the two conditions of the lemma holds), we see that the set

$$Q^* = \big\{ \mu \mid \mu^* \geq 0, \, 0 \leq f(x) - f^* + \mu^{*\prime} g(x), \, \forall \; x \in X \big\},$$

is nonempty [and also compact if condition (1) of the lemma is satisfied]. The vectors $\mu^* \in Q^*$ are precisely those for which $\mu^* \geq 0$ and

$$f^* \le \inf_{x \in X} \{ f(x) + \mu^{*'} g(x) \} = q(\mu^*) \le q^*.$$

Using the weak duality relation $q^* \leq f^*$, we see that $f^* = q(\mu^*) = q^*$ if and only if $\mu^* \in Q^*$, i.e., Q^* coincides with the set of optimal solutions of the dual problem. We state the conclusion as a proposition.

Proposition 5.3.1: (Convex Programming Duality - Existence of Dual Optimal Solutions) Consider the problem (5.6). Assume that f^* is finite, and that one of the following two conditions holds:

- (1) There exists $\overline{x} \in X$ such that $g_j(\overline{x}) < 0$ for all $j = 1, \dots, r$.
- (2) The functions g_j , j = 1, ..., r, are affine, and there exists $\overline{x} \in ri(X)$ such that $g(\overline{x}) \leq 0$.

Then $q^* = f^*$ and the set of optimal solutions of the dual problem is nonempty. Under condition (1) this set is also compact.

The interior condition (1) in the preceding proposition is known in the nonlinear programming literature as the *Slater condition*.

5.3.2 Optimality Conditions

The following proposition gives necessary and sufficient conditions for optimality, which generalize the linear programming conditions of Prop. 5.2.2.

Proposition 5.3.2: (Optimality Conditions) Consider the problem (5.6). There holds $q^* = f^*$, and (x^*, μ^*) are a primal and dual optimal solution pair if and only if x^* is feasible, $\mu^* \geq 0$, and

$$x^* \in \arg\min_{x \in X} L(x, \mu^*), \qquad \mu_j^* g_j(x^*) = 0, \quad j = 1, \dots, r.$$
 (5.7)

Proof: If $q^* = f^*$, and x^* and μ^* are primal and dual optimal solutions,

$$f^* = q^* = q(\mu^*) = \inf_{x \in X} L(x, \mu^*) \le L(x^*, \mu^*) = f(x^*) + \sum_{j=1}^r \mu_j^* g_j(x^*) \le f(x^*),$$

where the last inequality follows since $\mu_j^* \geq 0$ and $g_j(x^*) \leq 0$ for all j. Hence equality holds throughout above and Eq. (5.7) holds.

Conversely, if x^* is feasible, $\mu^* \geq 0$, and Eq. (5.7) is satisfied, then

$$q(\mu^*) = \inf_{x \in X} L(x, \mu^*) = L(x^*, \mu^*) = f(x^*) + \sum_{j=1}^r \mu_j^* g_j(x^*) = f(x^*).$$

In view of the weak duality relation $q^* \leq f^*$, it follows that $q^* = f^*$, x^* is primal optimal, and μ^* is dual optimal. **Q.E.D.**

Note that the preceding proof makes no use of the convexity assumptions on f, g, or X, and indeed Prop. 5.3.2 holds without these assumptions. On the other hand the proposition is meaningful only when $q^* = f^*$, which can rarely be guaranteed without convexity assumptions. The condition $\mu_j^* g_j(x^*) = 0$ is known as *complementary slackness* and generalizes the corresponding linear programming condition (5.4).

When using Prop. 5.3.2 as a sufficient optimality condition, it is important to realize that for $x^* \in X$ to be primal optimal, it is not sufficient that it attain the minimum of $L(x, \mu^*)$ over X for some dual optimal solution μ^* ; this minimum may be attained by some x^* that are infeasible [violate the constraints $g(x) \leq 0$] or by some x^* that are feasible, but are not optimal (these must violate complementary slackness).

Example 5.3.1: (Quadratic Programming Duality)

Consider the quadratic programming problem

minimize
$$\frac{1}{2}x'Qx + c'x$$

subject to $Ax < b$,

where Q is a symmetric positive definite $n \times n$ matrix, A is an $r \times n$ matrix, b is a vector in \Re^r , and c is a vector in \Re^n . This is problem (5.6) with

$$f(x) = \frac{1}{2}x'Qx + c'x, \qquad g(x) = Ax - b, \qquad X = \Re^n.$$

Assuming the problem is feasible, it has a unique optimal solution x^* , since the cost function is strictly convex and coercive.

The dual function is

$$q(\mu) = \inf_{x \in \Re^n} L(x,\mu) = \inf_{x \in \Re^n} \Big\{ \tfrac{1}{2} x' Q x + c' x + \mu' (Ax - b) \Big\}.$$

The infimum is attained for $x = -Q^{-1}(c + A'\mu)$, and after substituting this expression in the preceding relation, a straightforward calculation yields

$$q(\mu) = -\frac{1}{2}\mu'AQ^{-1}A'\mu - \mu'(b + AQ^{-1}c) - \frac{1}{2}c'Q^{-1}c.$$

The dual problem, after changing the minus sign to convert the maximization to a minimization and dropping the constant $\frac{1}{2}c'Q^{-1}c$, can be written as

minimize
$$\frac{1}{2}\mu'P\mu + t'\mu$$

subject to $\mu \ge 0$,

where

$$P = AQ^{-1}A', t = b + AQ^{-1}c.$$

Note that the dual problem has simpler constraints than the primal. Furthermore, if the row dimension r of A is smaller than its column dimension n, the dual problem is defined on a space of smaller dimension than the primal, and this can be algorithmically significant.

By applying Prop. 5.3.1 under condition (2), we see that $f^* = q^*$ and that the dual problem has an optimal dual solution (the relative interior assumption of the proposition is trivially satisfied since $X = \Re^n$).

According to the optimality conditions of Prop. 5.3.2, (x^*, μ^*) is a primal and dual optimal solution pair if and only if $Ax^* \leq b$, $\mu^* \geq 0$, and the two conditions of Eq. (5.7) hold. The first of these conditions $[x^*$ minimizes $L(x, \mu^*)$ over $x \in \Re^n$ yields

$$x^* = -Q^{-1}(c + A'\mu^*).$$

The second is the complementary slackness condition $(Ax^* - b)'\mu^* = 0$, encountered in linear programming (cf. Prop. 5.2.2). It can be written as

$$\mu_i^* > 0 \quad \Rightarrow \quad a_i' x^* = b_i, \qquad \forall \ j = 1, \dots, r,$$

where a'_{j} is the jth row of A, and b_{j} is the jth component of b.

5.3.3 Partially Polyhedral Constraints

The preceding analysis for the inequality-constrained problem (5.6) can be refined by making more specific assumptions regarding available polyhedral structure in the constraint functions and the abstract constraint set X. We first consider some of the simplest cases where there is a mixture of polyhedral and nonpolyhedral constraints. We then provide a general duality theorem that covers simultaneously a broad variety of structures.

Consider an extension of problem (5.6) where there are additional linear equality constraints:

minimize
$$f(x)$$

subject to $x \in X$, $g(x) \le 0$, $Ax = b$, (5.8)

where X is a convex set, $g(x) = (g_1(x), \ldots, g_r(x))'$, $f: X \mapsto \Re$ and $g_j: X \mapsto \Re$, $j = 1, \ldots, r$, are convex functions, A is an $m \times n$ matrix, and $b \in \Re^m$. We can deal with this problem by simply converting the constraint Ax = b to the equivalent set of linear inequality constraints

$$Ax \le b, \qquad -Ax \le -b, \tag{5.9}$$

with corresponding dual variables $\lambda^+ \geq 0$ and $\lambda^- \geq 0$. The Lagrangian function is

$$f(x) + \mu' g(x) + (\lambda^{+} - \lambda^{-})' (Ax - b),$$

and by introducing a dual variable

$$\lambda = \lambda^{+} - \lambda^{-} \tag{5.10}$$

with no sign restriction, it can be written as

$$L(x, \mu, \lambda) = f(x) + \mu' g(x) + \lambda' (Ax - b).$$

The dual problem is

$$\begin{aligned} & \text{maximize} & & q(\mu,\lambda) \equiv \inf_{x \in X} L(x,\mu,\lambda) \\ & \text{subject to} & & \mu \geq 0, \ \lambda \in \Re^m. \end{aligned}$$

In the special case of a problem with just linear equality constraints:

minimize
$$f(x)$$

subject to $x \in X$, $Ax = b$, (5.11)

the Lagrangian function is

$$L(x,\lambda) = f(x) + \lambda'(Ax - b),$$

and the dual problem is

$$\label{eq:problem} \begin{aligned} \text{maximize} & & q(\lambda) \equiv \inf_{x \in X} L(x, \lambda) \\ \text{subject to} & & \lambda \in \Re^m. \end{aligned}$$

The following two propositions are obtained by transcription of Prop. 5.3.1 [under condition (2) which deals with linear constraints] and Prop. 5.3.2, using the transformations (5.9), (5.10). The straightforward details are left for the reader.

Proposition 5.3.3: (Convex Programming - Linear Equality Constraints) Consider problem (5.11).

- (a) Assume that f^* is finite and that there exists $\overline{x} \in \text{ri}(X)$ such that $A\overline{x} = b$. Then $f^* = q^*$ and there exists at least one dual optimal solution.
- (b) There holds $f^* = q^*$, and (x^*, λ^*) are a primal and dual optimal solution pair if and only if x^* is feasible and

$$x^* \in \arg\min_{x \in X} L(x, \lambda^*). \tag{5.12}$$

Proposition 5.3.4: (Convex Programming - Linear Equality and Inequality Constraints) Consider problem (5.8).

- (a) Assume that f^* is finite, that the functions g_j are linear, and that there exists $\overline{x} \in \text{ri}(X)$ such that $A\overline{x} = b$ and $g(\overline{x}) \leq 0$. Then $g^* = f^*$ and there exists at least one dual optimal solution.
- (b) There holds $f^* = q^*$, and (x^*, μ^*, λ^*) are a primal and dual optimal solution pair if and only if x^* is feasible, $\mu^* \geq 0$, and

$$x^* \in \arg\min_{x \in X} L(x, \mu^*, \lambda^*), \qquad \mu_j^* g_j(x^*) = 0, \quad j = 1, \dots, r.$$

The following is an extension of part (a) of the preceding proposition to the case where the inequality constraints may be nonlinear.

Proposition 5.3.5: (Convex Programming - Linear Equality and Nonlinear Inequality Constraints) Consider problem (5.8). Assume that f^* is finite, that there exists $\overline{x} \in X$ such that $A\overline{x} = b$ and $g(\overline{x}) < 0$, and that there exists $\tilde{x} \in \text{ri}(X)$ such that $A\tilde{x} = b$. Then $q^* = f^*$ and there exists at least one dual optimal solution.

Proof: By applying Prop. 5.3.1 under condition (1), we see that there exists $\mu^* \geq 0$ such that

$$f^* = \inf_{x \in X, Ax = b} \left\{ f(x) + \mu^*'g(x) \right\}.$$

By applying Prop. 5.3.3 to the minimization problem above, we see that there exists λ^* such that

$$f^* = \inf_{x \in X} \left\{ f(x) + \mu^{*'}g(x) + \lambda^{*'}(Ax - b) \right\} = \inf_{x \in X} L(x, \mu^*, \lambda^*) = q(\lambda^*, \mu^*).$$

From weak duality we have $q(\mu^*, \lambda^*) \le q^* \le f^*$, so it follows that $q^* = f^*$ and that (μ^*, λ^*) is a dual optimal solution. **Q.E.D.**

The following example illustrates how we can have $q^* < f^*$ in convex programming.

Example 5.3.2: (Strong Duality Counterexample)

Consider the two-dimensional problem

minimize
$$f(x)$$

subject to $x_1 = 0$, $x \in X = \{x \mid x \ge 0\}$,

where

$$f(x) = e^{-\sqrt{x_1 x_2}}, \quad \forall \ x \in X.$$

Here it can be verified that f is convex (its Hessian is positive definite in the interior of X). Since for feasibility we must have $x_1 = 0$, we see that $f^* = 1$. On the other hand, the dual function is

$$q(\lambda) = \inf_{x \ge 0} \left\{ e^{-\sqrt{x_1 x_2}} + \lambda x_1 \right\} = \begin{cases} 0 & \text{if } \lambda \ge 0, \\ -\infty & \text{otherwise,} \end{cases}$$

since when $\lambda \geq 0$, the expression in braces above is nonnegative for $x \geq 0$ and can approach zero by taking $x_1 \to 0$ and $x_1x_2 \to \infty$. It follows that $q^* = 0$. Thus, there is a duality gap, $f^* - q^* = 1$. Here the relative interior assumption of Prop. 5.3.3(a) is violated.

For an example where $q^* = f^*$ but there exists no dual optimal solution, take

$$X = \Re, \qquad f(x) = x, \qquad g(x) = x^2.$$

Then $x^* = 0$ is the only feasible/optimal solution, and we have

$$q(\mu)=\inf_{x\in\Re}\{x+\mu x^2\}=-\frac{1}{4\mu},\qquad\forall\;\mu>0,$$

and $q(\mu) = -\infty$ for $\mu \le 0$, so that $q^* = f^* = 0$. However, there is no $\mu^* \ge 0$ such that $q(\mu^*) = q^* = 0$. This is a typical constrained optimization situation where there are no Lagrange multipliers (as usually defined in nonlinear programming; see e.g., [Ber99]), because some form of "regularity" condition fails to hold. We will revisit this example in Section 5.3.4.

Mixtures of Constraints

We finally consider a more refined polyhedral structure that allows an essentially arbitrary mixture of polyhedral and nonpolyhedral constraints. In particular, we consider the problem

minimize
$$f(x)$$

subject to $x \in X$, $g(x) \le 0$, $Ax = b$, (5.13)

where X is the intersection of a polyhedral set P and a convex set C,

$$X = P \cap C$$
,

 $g(x) = (g_1(x), \dots, g_r(x))'$, the functions $f : \mathbb{R}^n \to \mathbb{R}$ and $g_j : \mathbb{R}^n \to \mathbb{R}$, $j = 1, \dots, r$, are defined over \mathbb{R}^n , A is an $m \times n$ matrix, and $b \in \mathbb{R}^m$.

We will assume that some of the functions g_j are polyhedral, i.e., each is specified as the maximum of a finite number of linear functions (cf.

Section 2.3.3). We will also assume that f and g_j are convex over C (rather than just X). This stronger convexity assumption is significant for taking advantage of the (partially) polyhedral character of X (for example, it fails to hold in Example 5.3.2, with a duality gap resulting).

Proposition 5.3.6: (Convex Programming - Mixed Polyhedral and Nonpolyhedral Constraints) Consider problem (5.13). Assume that f^* is finite and that for some \overline{r} with $1 \leq \overline{r} \leq r$, the functions g_j , $j = 1, \ldots, \overline{r}$, are polyhedral, and the functions f and g_j , $j = \overline{r} + 1, \ldots, r$, are convex over C. Assume further that:

(1) There exists a vector $\tilde{x} \in ri(C)$ in the set

$$\tilde{P} = P \cap \{x \mid Ax = b, g_j(x) \le 0, j = 1, \dots, \overline{r}\}.$$

(2) There exists $\overline{x} \in \tilde{P} \cap C$ such that $g_j(\overline{x}) < 0$ for all $j = \overline{r} + 1, \dots, r$. Then $q^* = f^*$ and there exists at least one dual optimal solution.

Proof: Consider the problem

minimize f(x)subject to $x \in \tilde{P} \cap C$, $g_j(x) \leq 0$, $j = \overline{r} + 1, \dots, r$,

which is equivalent to problem (5.13). We use Prop. 5.3.1 and the assumption (2) to assert that there exist $\mu_i^* \geq 0$, $j = \overline{r} + 1, \dots, r$, such that

$$f^* = \inf_{x \in \tilde{P} \cap C} \left\{ f(x) + \sum_{j=\overline{r}+1}^r \mu_j^* g_j(x) \right\}.$$
 (5.14)

We next consider the minimization problem in the preceding equation, and we introduce explicit representations for the polyhedral set P and the polyhedral functions $g_j(x) \leq 0, j = 1, \ldots, \overline{r}$, in terms of linear inequalities and linear functions:

$$P = \{x \mid e'_{i0}x \le d_{i0}, i = 1, \dots, m_0\},$$

$$g_j(x) = \max_{i=1,\dots,m_j} \{e'_{ij}x - d_{ij}\}, \qquad j = 1,\dots,\overline{r},$$

where e_{ij} are some vectors in \Re^n and d_{ij} are corresponding scalars. We write Eq. (5.14) as

$$f^* = \inf_{x \in C, Ax = b, e'_{ij} x - d_{ij} \le 0, j = 0, \dots, \overline{r}, i = 1, \dots, m_j} \left\{ f(x) + \sum_{j = \overline{r} + 1}^r \mu_j^* g_j(x) \right\}.$$
(5.15)

We now use the assumption (1) and Prop. 5.3.4 in connection with the minimization problem in the right-hand side of the preceding equation, to assert that there exist a vector λ^* and scalars $\nu_{ii}^* \geq 0$ such that

$$f^* = \inf_{x \in C} \left\{ f(x) + \sum_{j=\overline{r}+1}^r \mu_j^* g_j(x) + \lambda^{*'} (Ax - b) + \sum_{j=1}^{\overline{r}} \sum_{i=1}^{m_j} \nu_{ij}^* (e'_{ij}x - d_{ij}) \right\}$$

$$\leq \inf_{x \in C \cap P} \left\{ f(x) + \sum_{j=\overline{r}+1}^r \mu_j^* g_j(x) + \lambda^{*'} (Ax - b) + \sum_{j=1}^{\overline{r}} \sum_{i=1}^{m_j} \nu_{ij}^* (e'_{ij}x - d_{ij}) \right\}$$

(the inequality follows since we are taking infimum over a subset of C). Since for all $x \in P$, we have $\nu_{i0}^*(e'_{i0}x - d_{i0}) \le 0$, it follows that

$$f^* \leq \inf_{x \in C \cap P} \left\{ f(x) + \sum_{j=\overline{r}+1}^r \mu_j^* g_j(x) + \lambda^{*'} (Ax - b) + \sum_{j=1}^{\overline{r}} \sum_{i=1}^{m_j} \nu_{ij}^* (e'_{ij} x - d_{ij}) \right\}$$

$$\leq \inf_{x \in C \cap P} \left\{ f(x) + \sum_{j=\overline{r}+1}^r \mu_j^* g_j(x) + \lambda^{*'} (Ax - b) + \sum_{j=1}^{\overline{r}} \left(\sum_{i=1}^{m_j} \nu_{ij}^* \right) g_j(x) \right\}$$

$$= \inf_{x \in C \cap P} \left\{ f(x) + \mu^{*'} g(x) + \lambda^{*'} (Ax - b) \right\}$$

$$= q(\mu^*, \lambda^*),$$

where $\mu^* = (\mu_1^*, ..., \mu_r^*)$ with

$$\mu_j^* = \sum_{i=1}^{m_j} \nu_{ij}^*, \qquad j = 1, \dots, \overline{r}.$$

In view of the weak duality relation $q(\mu^*, \lambda^*) \leq q^* \leq f^*$, it follows that $q^* = f^*$ and that (μ^*, λ^*) is a dual optimal solution. **Q.E.D.**

Note that the preceding proposition contains as special cases Props. 5.3.3(a), 5.3.4(a), and 5.3.5.

5.3.4 Duality and Existence of Optimal Primal Solutions

Our approach so far for establishing strong duality in constrained optimization has been based on the Nonlinear Farkas' Lemma. It provides assumptions guaranteeing that the dual problem has an optimal solution (even if there may be no primal optimal solution; cf. Prop. 5.3.1). We will now give an alternative approach, which under some compactness assumptions, guarantees strong duality and that there exists an optimal primal solution (even if there may be no dual optimal solution).

We focus on the convex programming problem

minimize
$$f(x)$$

subject to $x \in X$, $g(x) \le 0$, (5.16)

where X is a convex set, $g(x) = (g_1(x), \dots, g_r(x))'$, and $f: X \mapsto \Re$ and $g_j: X \mapsto \Re$, $j = 1, \dots, r$, are convex functions. We consider the MC/MC framework with $M = \operatorname{epi}(p)$, where p is the perturbation function

$$p(u) = \inf_{x \in X, g(x) \le u} f(x)$$

(cf. Section 4.2.3). From Prop. 4.3.1, we know that strong duality holds if p is closed proper convex (see also Fig. 4.2.1). With this in mind, let

$$F(x,u) = \begin{cases} f(x) & \text{if } x \in X, \ g(x) \le u, \\ \infty & \text{otherwise,} \end{cases}$$

and note that

$$p(u) = \inf_{x \in \Re^n} F(x, u),$$

so p is obtained from F by partial minimization. By using the corresponding analysis of Section 3.3, we obtain the following.

Proposition 5.3.7: (Convex Programming Duality - Existence of Primal Optimal Solutions) Assume that the problem (5.16) is feasible, that the convex functions f and g_j are closed, and that the function

$$F(x,0) = \begin{cases} f(x) & \text{if } g(x) \le 0, \ x \in X, \\ \infty & \text{otherwise,} \end{cases}$$

has compact level sets. Then $f^* = q^*$ and the set of optimal solutions of the primal problem is nonempty and compact.

Proof: From Prop. 3.3.2 applied to the function F, the partial minimum function p is convex and closed. From Prop. 4.3.1, it follows that $f^* = q^*$. Furthermore, since F(x,0) has compact level sets, the set of minima of F(x,0), which is equal to the set of optimal solutions of the primal problem, is nonempty and compact. **Q.E.D.**

The compactness assumption of the preceding proposition is satisfied, in particular, if either X is compact, or if X is closed and f has compact level sets. More generally, it is satisfied if X is closed, and X, f, and g_j , $j=1,\ldots,r$, have no common nonzero direction of recession. The proposition, however, does not guarantee the existence of a dual optimal solution, as illustrated by the following example.

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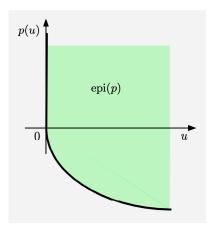


Figure 5.3.1. The perturbation function p for Example 5.3.3:

$$p(u) = \inf_{x^2 \le u} x = \begin{cases} -\sqrt{u} & \text{if } u \ge 0, \\ \infty & \text{if } u < 0. \end{cases}$$

Here p is lower semicontinuous at 0 and there is no duality gap. However, there is no dual optimal solution.

Example 5.3.3: (Nonexistence of Dual Optimal Solutions)

Consider the one-dimensional problem considered following Example 5.3.2, where

$$f(x) = x,$$
 $g(x) = x^2,$ $X = \Re.$

The perturbation function is shown in Fig. 5.3.1. It is convex and closed. Also there is a unique primal optimal solution $x^* = 0$ and there is no duality gap, consistent with the preceding proposition (the compactness assumption is satisfied). The dual function is

$$q(\mu) = \inf_{x \in \Re} \left\{ x + \mu x^2 \right\} = \left\{ \begin{array}{ll} -1/(4\mu) & \text{if } \mu > 0, \\ -\infty & \text{if } \mu \leq 0. \end{array} \right.$$

Thus there is no dual optimal solution, something that is also apparent from Fig. 5.3.1.

A result that is related to Prop. 5.3.7 is that if x^* is the unique minimum over X of the Lagrangian $L(\cdot, \mu^*)$ for some dual optimal solution μ^* , and if X, $F(\cdot,0)$, and $L(\cdot,\mu^*)$ are closed, then x^* is the unique primal optimal solution. The reason is that $L(x,\mu^*) \leq F(x,0)$ for all $x \in X$, which implies that $F(\cdot,0)$ has compact level sets, so there exists a primal optimal solution. Since by Prop. 5.3.2, all primal optimal solutions minimize $L(\cdot,\mu^*)$ over X, it follows that x^* is the unique primal optimal solution.

5.3.5 Fenchel Duality

We will now analyze another important optimization framework, which can be embedded within the convex programming framework discussed so far. Consider the problem

minimize
$$f_1(x) + f_2(Ax)$$

subject to $x \in \Re^n$, (5.17)

where A is an $m \times n$ matrix, $f_1 : \mathbb{R}^n \mapsto (-\infty, \infty]$ and $f_2 : \mathbb{R}^m \mapsto (-\infty, \infty]$ are closed convex functions, and we assume that there exists a feasible solution. We convert it to the following equivalent problem in the variables $x_1 \in \mathbb{R}^n$ and $x_2 \in \mathbb{R}^m$:

minimize
$$f_1(x_1) + f_2(x_2)$$

subject to $x_1 \in \text{dom}(f_1), x_2 \in \text{dom}(f_2), x_2 = Ax_1.$ (5.18)

We can view this as a convex programming problem with the linear equality constraint $x_2 = Ax_1$ [cf. problem (5.11)]. The dual function is

$$q(\lambda) = \inf_{x_1 \in \text{dom}(f_1), \ x_2 \in \text{dom}(f_2)} \left\{ f_1(x_1) + f_2(x_2) + \lambda'(x_2 - Ax_1) \right\}$$
$$= \inf_{x_1 \in \Re^n} \left\{ f_1(x_1) - \lambda' Ax_1 \right\} + \inf_{x_2 \in \Re^n} \left\{ f_2(x_2) + \lambda' x_2 \right\}.$$

The dual problem, after a sign change to convert it to a minimization problem, takes the form

minimize
$$f_1^{\star}(A'\lambda) + f_2^{\star}(-\lambda)$$

subject to $\lambda \in \Re^m$, (5.19)

where f_1^{\star} and f_2^{\star} are the conjugate functions of f_1 and f_2 , respectively:

$$f_1^{\star}(\lambda) = \sup_{x \in \mathbb{R}^n} \{ \lambda' x - f_1(x) \}, \quad f_2^{\star}(\lambda) = \sup_{x \in \mathbb{R}^n} \{ \lambda' x - f_2(x) \}, \quad \lambda \in \mathbb{R}^n.$$

Note that the primal and dual problems have a similar/symmetric form. Figure 5.3.2 illustrates the duality between problems (5.17) and (5.19).

We now apply Prop. 5.3.3 to the primal problem (5.18) and obtain the following proposition.

Proposition 5.3.8: (Fenchel Duality)

- (a) If f^* is finite and $(A \cdot \text{ri}(\text{dom}(f_1))) \cap \text{ri}(\text{dom}(f_2)) \neq \emptyset$, then $f^* = q^*$ and there exists at least one dual optimal solution.
- (b) There holds $f^* = q^*$, and (x^*, λ^*) is a primal and dual optimal solution pair if and only if

$$x^* \in \arg\min_{x \in \Re^n} \left\{ f_1(x) - x'A'\lambda^* \right\} \text{ and } Ax^* \in \arg\min_{z \in \Re^n} \left\{ f_2(z) + z'\lambda^* \right\}.$$

$$(5.20)$$

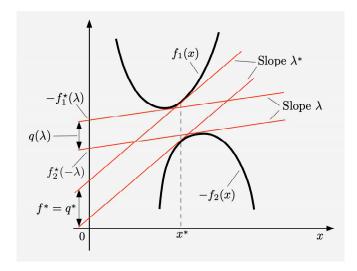


Figure 5.3.2. Illustration of Fenchel duality for the case where A is the identity matrix. The dual function value $q(\lambda)$ is constructed as in the figure. As λ changes, the dual value $q(\lambda) = -\left(f_1^*(\lambda) + f_2^*(-\lambda)\right)$ reaches its maximum at a vector λ^* such that corresponding hyperplanes support the epigraphs of f_1 and f_2 at a common point x^* , which is a primal optimal solution.

Proof: (a) Using the Cartesian product formula for relative interiors (cf. Prop. 1.3.10), it follows that the relative interior assumption of Prop. 5.3.3(a), applied to problem (5.18), is satisfied if for some $\overline{x}_1 \in \text{ri}(\text{dom}(f_1))$ and $\overline{x}_2 \in \text{ri}(\text{dom}(f_2))$ we have $\overline{x}_2 = A\overline{x}_1$. This is equivalent to our relative interior assumption, so Prop. 5.3.3(a) can be used, yielding the desired result.

(b) Similarly, we apply Prop. 5.3.3(b) to problem (5.18). Q.E.D.

Note that in the case where $dom(f_1) = \Re^n$, $dom(f_2) = \Re^m$, and f_1 and f_2 are differentiable, the optimality condition (5.20) is equivalent to

$$A'\lambda^* = \nabla f_1(x^*), \qquad \lambda^* = -\nabla f_2(Ax^*);$$

see Fig. 5.3.2. This condition will be generalized to the nondifferentiable case, by using subgradients (see Section 5.4.1, following the Conjugate Subgradient Theorem).

By reversing the roles of the (symmetric) primal and dual problems, and by applying the preceding conditions to the dual problem (5.19), we can obtain alternative criteria for strong duality. In particular, if q^* is finite and $\operatorname{ri}(\operatorname{dom}(f_1^*)) \cap (A' \cdot \operatorname{ri}(-\operatorname{dom}(f_2^*))) \neq \emptyset$, then $f^* = q^*$ and there exists at least one primal optimal solution.

Finally, we note that a more refined Fenchel duality theorem can be obtained if there is polyhedral structure in f_1 and f_2 . For example, if f_1 is

polyhedral, the relative interior condition of Prop. 5.3.8(a) can be replaced by the weaker condition $A \cdot \text{dom}(f_1) \cap \text{ri}(\text{dom}(f_2)) \neq \emptyset$. This follows by applying Prop. 5.3.6 instead of Prop. 5.3.3(a) in the preceding proof. More generally, the same can be shown if f_1 is of the form

$$f_1(x) = \begin{cases} \tilde{f}(x) & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$

where X is a polyhedral set and $\tilde{f}: \Re^n \mapsto \Re$ is a function that is convex over a convex set C with $X \subset \text{ri}(C)$. Similarly, if f_1 and f_2 are both polyhedral, the relative interior condition is unnecessary. In this case the problem (5.17) is equivalent to a linear program for which strong duality holds and a dual optimal solution exists under just the condition that f^* is finite (cf. Prop. 5.2.1).

5.3.6 Conic Duality

Consider the problem

minimize
$$f(x)$$

subject to $x \in C$, (5.21)

where $f: \Re^n \mapsto (-\infty, \infty]$ is a closed proper convex function and C is a closed convex cone in \Re^n . This is known as a *conic program*, and some of its special cases (semidefinite programming, second order cone programming) have many practical applications, for which we refer to the literature, e.g., [BeN01], [BoV04].

We apply Fenchel duality with A equal to the identity and the definitions

$$f_1(x) = f(x),$$
 $f_2(x) = \begin{cases} 0 & \text{if } x \in C, \\ \infty & \text{if } x \notin C. \end{cases}$

The corresponding conjugates are

$$f_1^{\star}(\lambda) = \sup_{x \in \Re^n} \left\{ \lambda' x - f(x) \right\}, \qquad f_2^{\star}(\lambda) = \sup_{x \in C} \lambda' x = \begin{cases} 0 & \text{if } \lambda \in C^*, \\ \infty & \text{if } \lambda \notin C^*, \end{cases}$$

where

$$C^* = \{ \lambda \mid \lambda' x \le 0, \, \forall \, x \in C \}$$

is the polar cone of C (note that f_2^{\star} is the support function of C; cf. Example 1.6.1). The dual problem [cf. Eq. (5.19)] is

minimize
$$f^*(\lambda)$$

subject to $\lambda \in \hat{C}$, (5.22)

where f^* is the conjugate of f and \hat{C} is the negative polar cone (also called the *dual cone* of C):

$$\hat{C} = -C^* = \{ \lambda \mid \lambda' x \ge 0, \, \forall \, x \in C \}.$$

Note the symmetry between primal and dual problems. The strong duality relation $f^* = q^*$ can be written as

$$\inf_{x \in C} f(x) = -\inf_{\lambda \in \hat{C}} f^{\star}(\lambda).$$

The following proposition translates the conditions of Prop. 5.3.8 for asserting that there is no duality gap and the dual problem has an optimal solution.

Proposition 5.3.9: (Conic Duality Theorem) Assume that the optimal value of the primal conic problem (5.21) is finite, and that $\operatorname{ri}(\operatorname{dom}(f)) \cap \operatorname{ri}(C) \neq \emptyset$. Then, there is no duality gap and that the dual problem (5.22) has an optimal solution.

Using the symmetry of the primal and dual problems, we also obtain that there is no duality gap and the primal problem (5.21) has an optimal solution if the optimal value of the dual conic problem (5.22) is finite and $\operatorname{ri}(\operatorname{dom}(f^*)) \cap \operatorname{ri}(\hat{C}) \neq \emptyset$. It is also possible to exploit polyhedral structure in f and/or C, using Prop. 5.3.6 (cf. the discussion at the end of the preceding section). Furthermore, we may derive primal and dual optimality conditions using Prop. 5.3.8(b).

5.4 SUBGRADIENTS AND OPTIMALITY CONDITIONS

In this section we introduce the notion of a subgradient of a convex function at a point. Subgradients serve as a substitute for gradients when the function is nondifferentiable: like gradients in differentiable cost minimization, they enter in optimality conditions and find wide use in algorithms.

We will develop a connection of subgradients with the MC/MC framework. In particular, we will show that the subdifferential can be identified with the set of max crossing hyperplanes in a suitable MC/MC framework. Through this connection we will obtain some of the basic results of subdifferential theory by using the MC/MC theory of Chapter 4, the Nonlinear Farkas' Lemma (cf. Prop. 5.1.1), and the constrained optimization duality theory (cf. Prop. 5.3.6).

Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a proper convex function. We say that a vector $g \in \mathbb{R}^n$ is a *subgradient* of f at a point $x \in \text{dom}(f)$ if

$$f(z) \ge f(x) + g'(z - x), \qquad \forall \ z \in \Re^n. \tag{5.23}$$

The set of all subgradients of f at x is called the *subdifferential of* f at x and is denoted by $\partial f(x)$. By convention, $\partial f(x)$ is considered empty for all $x \notin \text{dom}(f)$. Generally, $\partial f(x)$ is closed and convex, since based

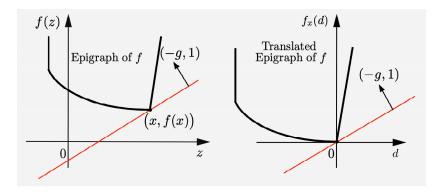


Figure 5.4.1. Illustration of the subdifferential $\partial f(x)$ of a convex function f and its connection with the MC/MC framework. The subgradient inequality (5.23) can be written as

$$f(z) - z'g \ge f(x) - x'g, \qquad \forall \ z \in \Re^n. \tag{5.24}$$

Thus, g is a subgradient of f at x if and only if the hyperplane in \Re^{n+1} that has normal (-g,1) and passes through (x,f(x)) supports the epigraph of f, as shown in the left-hand side figure.

The right-hand side figure shows that $\partial f(x)$ is the set of max crossing solutions in the MC/MC framework where M is the epigraph of f_x , the x-translation of f.

on the subgradient inequality (5.23), it is the intersection of a collection of closed halfspaces. Note that we restrict attention to proper functions (subgradients are not useful and make no sense for improper functions).

As Fig. 5.4.1 illustrates, g is a subgradient of f at x if and only if the hyperplane in \Re^{n+1} that has normal (-g,1) and passes through (x,f(x)) supports the epigraph of f. From this geometric view, it is evident that there is a strong connection with the MC/MC framework. In particular, for any $x \in \text{dom}(f)$, consider the x-translation of f, which is the function f_x , whose epigraph is the epigraph of f translated so that (x, f(x)) is moved to the origin of \Re^{n+1} :

$$f_x(d) = f(x+d) - f(x), \qquad d \in \Re^n.$$

Then the subdifferential $\partial f(x)$ is the set of all max crossing solutions for the MC/MC framework corresponding to the set

$$M = epi(f_x) = epi(f) - \{(x, f(x))\}$$
 (5.25)

(cf. the right-hand side of Fig. 5.4.1). Based on this fact, we can use the MC/MC theory to obtain results regarding the existence of subgradients, as in the following proposition.

Proposition 5.4.1: Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a proper convex function. For every $x \in \text{ri}(\text{dom}(f))$,

$$\partial f(x) = S^{\perp} + G,$$

where S is the subspace that is parallel to the affine hull of dom(f), and G is a nonempty convex and compact set. In particular, if $x \in int(dom(f))$, then $\partial f(x)$ is nonempty and compact.

Proof: The result follows by applying Props. 4.4.1 and 4.4.2 to the set M given by Eq. (5.25). **Q.E.D.**

It follows from the preceding proposition that if f is real-valued, then $\partial f(x)$ is nonempty and compact for all $x \in \Re^n$. If f is extended real-valued, $\partial f(x)$ can be unbounded, and it can be empty not only for $x \notin \text{dom}(f)$, but also for some x in the boundary of dom(f). As an example, consider the function

$$f(x) = \begin{cases} -\sqrt{x} & \text{if } 0 \le x \le 1, \\ \infty & \text{otherwise.} \end{cases}$$

Its subdifferential is

$$\partial f(x) = \begin{cases} -\frac{1}{2\sqrt{x}} & \text{if } 0 < x < 1, \\ \left[-\frac{1}{2}, \infty \right) & \text{if } x = 1, \\ \emptyset & \text{if } x \le 0 \text{ or } 1 < x, \end{cases}$$

so it is empty or unbounded at boundary points within dom(f) (the points 0 and 1, respectively).

An important property is that if f is differentiable at some $x \in \operatorname{int}(\operatorname{dom}(f))$, its gradient $\nabla f(x)$ is the unique subgradient at x. Indeed, by Prop. 1.1.7(a), $\nabla f(x)$ satisfies the subgradient inequality (5.23), so it is a subgradient at x. To show uniqueness, note that if g is a subgradient at x, we have

$$f(x) + \alpha g'd \leq f(x + \alpha d) = f(x) + \alpha \nabla f(x)'d + o(|\alpha|), \quad \forall \ \alpha \in \Re, \ d \in \Re^n.$$

By letting $d = \nabla f(x) - g$, we obtain

$$0 \le \alpha (\nabla f(x) - g)' d + o(|\alpha|) = \alpha ||\nabla f(x) - g||^2 + o(|\alpha|).$$

In particular, we have

$$\|\nabla f(x) - g\|^2 \le -\frac{o(|\alpha|)}{\alpha}, \quad \forall \alpha < 0,$$

Taking $\alpha \uparrow 0$, we obtain $\nabla f(x) - g = 0$.

Let us also consider another important special case, where f is the indicator function of a convex set.

Example 5.4.1: (Subdifferential of an Indicator Function)

Let us derive the subdifferential of the indicator function of a nonempty convex set C:

$$\delta_C(x) = \begin{cases} 0 & \text{if } x \in C, \\ \infty & \text{if } x \notin C. \end{cases}$$

For all $x \notin C$, we have $\partial \delta_C(x) = \emptyset$, by convention. For $x \in C$, we have $g \in \partial \delta_C(x)$ if and only if

$$\delta_C(z) \ge \delta_C(x) + q'(z - x), \quad \forall z \in C,$$

or equivalently $g'(z-x) \leq 0$ for all $z \in C$. For $x \in C$, the set of all g satisfying this relation is called the *normal cone of* C at x and is denoted by $N_C(x)$:

$$N_C(x) = \{g \mid g'(z - x) \le 0, \forall z \in C\}.$$

Thus the normal cone $N_C(x)$ is the polar cone of $C - \{x\}$, the set C translated so that x is moved to the origin (see Fig. 5.4.2).

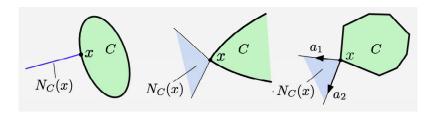


Figure 5.4.2. Illustration of the normal cone of a convex set C at a point $x \in C$, i.e., the subdifferential of the indicator function of C at x. For $x \in \text{int}(C)$ we have $N_C(x) = \{0\}$, while for $x \notin \text{int}(C)$, the normal cone contains at least one half line, as in the examples shown. In the case of the polyhedral set on the right, $N_C(x)$ is the cone generated by the normal vectors to the hyperplanes that correspond to the active inequalities at x.

Finally, let us show an important property of real-valued convex functions.

Proposition 5.4.2: (Subdifferential Boundedness and Lipschitz Continuity) Let $f: \mathbb{R}^n \to \mathbb{R}$ be a real-valued convex function, and let X be a nonempty compact subset of \mathbb{R}^n .

- (a) The set $\bigcup_{x \in X} \partial f(x)$ is nonempty and bounded.
- (b) The function f is Lipschitz continuous over X, i.e., there exists a scalar L such that

$$|f(x) - f(z)| \le L ||x - z||, \quad \forall x, z \in X.$$

Proof: (a) Nonemptiness follows from Prop. 5.4.1. To prove boundedness, assume the contrary, so that there exists a sequence $\{x_k\} \subset X$, and an unbounded sequence $\{g_k\}$ with

$$g_k \in \partial f(x_k), \qquad 0 < ||g_k|| < ||g_{k+1}||, \quad k = 0, 1, \dots$$

We denote $d_k = g_k/\|g_k\|$. Since $g_k \in \partial f(x_k)$, we have

$$f(x_k + d_k) - f(x_k) \ge g'_k d_k = ||g_k||.$$

Since both $\{x_k\}$ and $\{d_k\}$ are bounded, they contain convergent subsequences. We assume without loss of generality that $\{x_k\}$ and $\{d_k\}$ converge to some vectors. Therefore, by the continuity of f (cf. Prop. 1.3.11), the left-hand side of the preceding relation is bounded. Hence the right-hand side is also bounded, thereby contradicting the unboundedness of $\{g_k\}$.

(b) Let x and z be any two points in X. By the subgradient inequality (5.23), we have

$$f(x) + g'(z - x) \le f(z), \quad \forall g \in \partial f(x),$$

so that

$$f(x) - f(z) \le ||g|| \cdot ||x - z||, \quad \forall g \in \partial f(x).$$

By part (a), $\bigcup_{y \in X} \partial f(y)$ is bounded, so that for some constant L > 0, we have

$$||g|| \le L, \quad \forall g \in \partial f(y), \quad \forall y \in X,$$
 (5.26)

and therefore,

$$f(x) - f(z) \le L \|x - z\|.$$

By exchanging the roles of x and z, we similarly obtain

$$f(z) - f(x) \le L ||x - z||,$$

and by combining the preceding two relations, we see that

$$|f(x) - f(z)| \le L ||x - z||,$$

showing that f is Lipschitz continuous over X. Q.E.D.

Note that the proof of part (b) shows how to determine the Lipschitz constant L: it is the maximum subgradient norm, over all subgradients in $\bigcup_{x \in X} \partial f(x)$ [cf. Eq. (5.26)].

5.4.1 Subgradients of Conjugate Functions

We will now derive an important relation between the subdifferentials of a proper convex function $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ and its conjugate f^* . Using the definition of conjugacy, we have

$$x'y \le f(x) + f^*(y), \quad \forall \ x \in \Re^n, \ y \in \Re^n.$$

This is known as the *Fenchel inequality*. A pair (x, y) satisfies this inequality as an equation if and only if x attains the supremum in the definition

$$f^{\star}(y) = \sup_{z \in \mathbb{R}^n} \{ y'z - f(z) \}.$$

Pairs of this type are connected with the subdifferentials of f and f^* , as shown in the following proposition, and illustrated in Fig. 5.4.3.

Proposition 5.4.3: (Conjugate Subgradient Theorem) Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a proper convex function and let f^* be its conjugate. The following two relations are equivalent for a pair of vectors (x,y):

- (i) $x'y = f(x) + f^*(y)$.
- (ii) $y \in \partial f(x)$.

If in addition f is closed, the relations (i) and (ii) are equivalent to

(iii) $x \in \partial f^*(y)$.

Proof: A pair (x, y) satisfies (i) if and only if x attains the supremum in the definition $f^*(y) = \sup_{z \in \Re^n} \{y'z - f(z)\}$, and by Eq. (5.24), this is true if and only if $y \in \partial f(x)$. This shows that (i) and (ii) are equivalent. If in addition f is closed, then by the Conjugacy Theorem [Prop. 1.6.1(c)], f is equal to the conjugate of f^* , so by using the equivalence just shown with the roles of f and f^* reversed, we see that (i) is equivalent to (iii). **Q.E.D.**

Note that the closure assumption in condition (iii) of the Conjugate Subgradient Theorem is necessary, because by the Conjugacy Theorem [Prop. 1.6.1(d)], the conjugate of f^* is cl f, so the relation $x \in \partial f^*(y)$ implies that

$$x'y = (\operatorname{cl} f)(x) + f^{\star}(y)$$

[by the equivalence of conditions (i) and (ii)]. On the other hand, for $x \in \partial f^*(y)$, we may have $(\operatorname{cl} f)(x) < f(x)$ and hence $x'y < f(x) + f^*(y)$

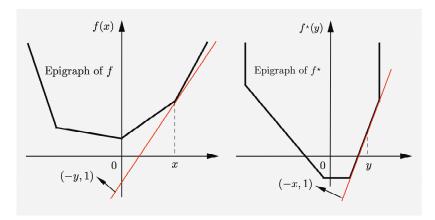


Figure 5.4.3. Illustration of a pair (x,y) satisfying Fenchel's inequality as an equation, i.e., $x'y = f(x) + f^*(y)$. If f is closed, this equality is equivalent to

$$y \in \partial f(x)$$
 and $x \in \partial f^*(y)$.

Since x minimizes a proper convex function f if and only if $0 \in \partial f(x)$ [a consequence of the subgradient inequality (5.23)], a corollary is that the set of minima of a closed proper convex f is $\partial f^*(0)$, while the set of minima of f^* is $\partial f(0)$.

[for example, take f to be the indicator function of the interval (-1,1), $f^*(y) = |y|, x = 1$, and y = 0].

For an application of the Conjugate Subgradient Theorem, note that the necessary and sufficient optimality condition (5.20) in the Fenchel Duality Theorem can be equivalently written as

$$A'\lambda^* \in \partial f_1(x^*), \qquad \lambda^* \in -\partial f_2(Ax^*)$$

(cf. Fig. 5.3.2).

The following proposition gives some useful corollaries of the Conjugate Subgradient Theorem:

Proposition 5.4.4: Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a closed proper convex function and let f^* be its conjugate.

- (a) f^* is differentiable at a vector $y \in \operatorname{int}(\operatorname{dom}(f^*))$ if and only if the supremum of x'y f(x) over $x \in \mathbb{R}^n$ is uniquely attained.
- (b) The set of minima of f is given by

$$\arg\min_{x\in\Re^n}f(x)=\partial f^\star(0),$$

Proof: Both parts follow from the fact

$$\arg\max_{x \in \Re^n} \{x'y - f(x)\} = \partial f^{\star}(y),$$

which is a consequence of Prop. 5.4.3. Q.E.D.

Proposition 5.4.4(a) can be used to characterize differentiability properties of f^* in terms of strict convexity properties of f, which guarantee the unique attainment of the supremum of x'y - f(x) over $x \in \Re^n$. This is the basis for the Legendre transformation, a precursor to the conjugacy transformation (see [Roc70], §26). Proposition 5.4.4(b) shows that the set of minima of f is nonempty and compact if and only if $0 \in \text{int}(\text{dom}(f^*))$ (cf. Prop. 5.4.1).

The following example applies the Conjugate Subgradient Theorem to establish an important interpretation of dual optimal solutions within the MC/MC framework.

Example 5.4.2: (Sensitivity Interpretation of Dual Optimal Solutions)

Consider the MC/MC framework for the case where the set M is the epigraph of a function $p: \mathbb{R}^n \mapsto [-\infty, \infty]$. Then the dual function is

$$q(\mu) = \inf_{u \in \Re^m} \{p(u) + \mu' u\} = -p^*(-\mu),$$

where p^* is the conjugate of p (cf. Section 4.2.1). Assume that p is proper convex, and that strong duality holds, i.e.,

$$p(0) = w^* = q^* = \sup_{\mu \in \Re^m} \{-p^*(-\mu)\}.$$

Let Q^* be the set of dual optimal solutions, i.e.,

$$Q^* = \left\{ \mu^* \mid p(0) + p^*(-\mu^*) = 0 \right\}.$$

Then it follows from Prop. 5.4.3 that $\mu^* \in Q^*$ if and only if $-\mu^* \in \partial p(0)$, i.e.,

$$Q^* = -\partial p(0).$$

This leads to various sensitivity interpretations of dual optimal solutions (assuming strong duality holds). The most interesting case is when p is convex and differentiable at 0, in which case $-\nabla p(0)$ is equal to the unique dual optimal solution μ^* . As an example, for the constrained optimization problem

$$\label{eq:force_function} \begin{aligned} & \text{minimize} & & f(x) \\ & \text{subject to} & & x \in X, \quad g(x) \leq 0, \end{aligned}$$

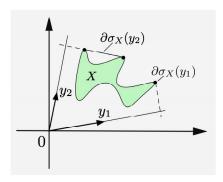


Figure 5.4.4. For a fixed y, the subdifferential $\partial \sigma_X(y)$ of the support function of X is the set of maxima of y'x over $x \in \text{cl}(\text{conv}(X))$ (cf. Example 5.4.3).

of Section 5.3, where

$$p(u) = \inf_{x \in X, g(x) \le u} f(x),$$

we have

$$\mu_j^* = -\frac{\partial p(0)}{\partial u_j}, \qquad j = 1, \dots, r,$$

so μ_j^* is the rate of improvement of the optimal primal cost as the jth constraint $g_j(x) \leq 0$ is violated.

Example 5.4.3: (Subdifferential of a Support Function)

Let us derive the subdifferential of the support function σ_X of a nonempty set X at a vector \overline{y} . Note that σ_X is closed proper convex, since it is the conjugate of the indicator function of X; cf. Example 1.6.1. To calculate $\partial \sigma_X(\overline{y})$, we introduce the closed proper convex function

$$r(y) = \sigma_X(y + \overline{y}), \qquad y \in \Re^n$$

and we note that $\partial \sigma_X(\overline{y}) = \partial r(0)$. The conjugate of r is

$$r^{\star}(x) = \sup_{y \in \Re^n} \{ y'x - \sigma_X(y + \overline{y}) \},$$

or

$$r^{\star}(x) = \sup_{y \in \Re^n} \left\{ (y + \overline{y})' x - \sigma_X(y + \overline{y}) \right\} - \overline{y}' x,$$

and finally

$$r^{\star}(x) = \delta(x) - \overline{y}'x,$$

where δ is the indicator function of $\operatorname{cl}(\operatorname{conv}(X))$ (cf. Example 1.6.1). Letting $r = f^*$ in Prop. 5.4.4(b), we see that $\partial r(0)$ is the set of minima of $\delta(x) - \overline{y}'x$, or equivalently $\partial \sigma_X(\overline{y})$ is the set of maxima of $\overline{y}'x$ over $x \in \operatorname{cl}(\operatorname{conv}(X))$ (see Fig. 5.4.4).

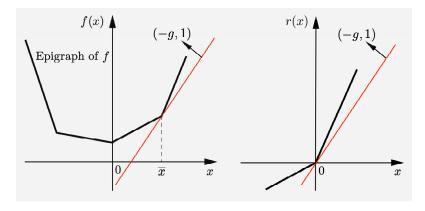


Figure 5.4.5. Construction used to derive the subdifferential at \overline{x} of the function

$$f(x) = \max\{a'_1x + b_1, \dots, a'_rx + b_r\},\$$

shown on the left (cf. Example 5.4.4). Consider the function on the right,

$$r(x) = \max \left\{ a_j' x \mid j \in A_{\overline{x}} \right\},\,$$

with $A_{\overline{x}} = \{j \mid a'_j \overline{x} + b_j = f(\overline{x})\}$. We claim that $\partial f(\overline{x}) = \partial r(0)$, from which we obtain $\partial f(\overline{x}) = \operatorname{conv}\left(\left\{a_j \mid j \in A_{\overline{x}}\right\}\right)$ [since r is the support function of the finite set $\{a_j \mid j \in A_{\overline{x}}\}$; cf. Example 5.4.3]. To verify that $\partial f(\overline{x}) = \partial r(0)$, note that for all $g \in \partial r(0)$, we have

$$f(x) - f(\overline{x}) \ge r(x - \overline{x}) \ge r(0) + g'(x - \overline{x}) = g'(x - \overline{x}), \quad \forall x \in \Re^n$$

It follows that $g \in \partial f(\overline{x})$ and $\partial r(0) \subset \partial f(\overline{x})$.

Conversely, let $g \in \partial f(\overline{x})$. Then for x sufficiently close to \overline{x} , we have $f(x) - f(\overline{x}) = r(x - \overline{x})$, so the inequality $f(x) - f(\overline{x}) \ge g'(x - \overline{x})$ implies that

$$r(x - \overline{x}) \ge r(0) + g'(x - \overline{x}).$$

From the definition of r, it follows that

$$r(x - \overline{x}) \ge r(0) + g'(x - \overline{x}), \quad \forall x \in \Re^n,$$

so $g \in \partial r(0)$ and $\partial f(\overline{x}) \subset \partial r(0)$.

Example 5.4.4: (Subdifferential of a Real-Valued Polyhedral Function)

Let us derive the subdifferential of a real-valued polyhedral function of the form

$$f(x) = \max\{a_1'x + b_1, \dots, a_r'x + b_r\},\$$

where $a_1, \ldots, a_r \in \mathbb{R}^n$ and $b_1, \ldots, b_r \in \mathbb{R}$. For a fixed $\overline{x} \in \mathbb{R}^n$, consider the set of "active" indices at x, that is, the ones that attain the maximum in the definition of $f(\overline{x})$:

$$A_{\overline{x}} = \{ j \mid a_j' \overline{x} + b_j = f(\overline{x}) \}.$$

Consider also the function

$$r(x) = \max \{ a_j' x \mid j \in A_{\overline{x}} \},\,$$

which is obtained by translating $\operatorname{epi}(f)$ so that $(\overline{x}, f(\overline{x}))$ is moved to the origin and the "inactive" functions $a'_j x + b_j$, $j \notin A_{\overline{x}}$, are "discarded" (see Fig. 5.4.5). It is evident from the figure, and can also be easily shown algebraically (see the caption of Fig. 5.4.5), that

$$\partial f(\overline{x}) = \partial r(0).$$

We now note that r is the support function of the finite set $\{a_j \mid j \in A_{\overline{x}}\}$, so from the result of Example 5.4.3, it follows that $\partial r(0)$ is the convex hull of this set. Thus,

$$\partial f(\overline{x}) = \operatorname{conv}(\{a_j \mid j \in A_{\overline{x}}\}).$$

5.4.2 Subdifferential Calculus

We will now generalize some of the basic theorems of ordinary differentiation by using the convex programming results of Section 5.3 (and by extension the Nonlinear Farkas' Lemma and the MC/MC framework). The following proposition generalizes the differentiation rule

$$\nabla F(x) = A' \nabla f(Ax)$$

for the function F(x) = f(Ax), where f is differentiable.

Proposition 5.4.5: (Chain Rule) Let $f: \Re^m \mapsto (-\infty, \infty]$ be a convex function, let A be an $m \times n$ matrix, and assume that the function F given by

$$F(x) = f(Ax)$$

is proper. Then

$$\partial F(x) \supset A' \partial f(Ax), \quad \forall \ x \in \Re^n.$$

Furthermore, if either f is polyhedral or else the range of A contains a point in the relative interior of dom(f), we have

$$\partial F(x) = A' \partial f(Ax), \quad \forall \ x \in \Re^n.$$

Proof: If $x \notin \text{dom}(F)$, then $\partial F(x) = A'\partial f(Ax) = \emptyset$. For any $x \in \text{dom}(F)$, if $d \in A'\partial f(Ax)$, there exists a $g \in \partial f(Ax)$ such that d = A'g. We have for all $z \in \Re^n$,

$$F(z) - F(x) - (z - x)'d = f(Az) - f(Ax) - (z - x)'A'g$$

$$= f(Az) - f(Ax) - (Az - Ax)'g$$

$$\geq 0,$$

where the inequality follows since $g \in \partial f(Ax)$. Hence $d \in \partial F(x)$, and we have $\partial F(x) \supset A'\partial f(Ax)$.

To prove the reverse inclusion under the given assumption, we let $d \in \partial F(x)$ and we show that $d \in A'\partial f(Ax)$ by viewing x as the solution of an optimization problem defined by d. Indeed, we have

$$F(z) \ge F(x) + (z - x)'d \ge 0, \quad \forall z \in \Re^n,$$

or

$$f(Az) - z'd \ge f(Ax) - x'd, \quad \forall z \in \Re^n.$$

Thus (Ax, x) solves the following optimization problem in the variables (y, z):

minimize
$$f(y) - z'd$$

subject to $y \in \text{dom}(f)$, $Az = y$. (5.27)

If f is polyhedral, dom(f) is polyhedral and f can be replaced in the above problem by a real-valued polyhedral function, so that we can use Prop. 5.3.6. If instead the range of A contains a point in ri(dom(f)), we use Prop. 5.3.3. In either case, we conclude that there is no duality gap, and that there exists a dual optimal solution λ , such that

$$(Ax, x) \in \arg\min_{y \in \Re^m} \inf_{z \in \Re^n} \{f(y) - z'd + \lambda'(Az - y)\}$$

[cf. Eq. (5.12)]. Since the minimization over z is unconstrained, we must have $d=A'\lambda$, thus obtaining

$$Ax \in \arg\min_{y \in \Re^m} \{f(y) - \lambda'y\},\$$

or

$$f(y) > f(Ax) + \lambda'(y - Ax), \quad \forall y \in \Re^m.$$

Hence $\lambda \in \partial f(Ax)$, so that $d = A'\lambda \in A'\partial f(Ax)$. It follows that $\partial F(x) \subset A'\partial f(Ax)$. Q.E.D.

From the preceding proof, it is seen that the assumption that f is polyhedral in Prop. 5.4.5 may be replaced by the weaker assumption that for some convex function $\tilde{f}: \Re^n \mapsto \Re$ and polyhedral set X, f is of form

$$f(x) = \begin{cases} \tilde{f}(x) & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$
 (5.28)

where \tilde{f} is convex over a convex set C with $X \subset \text{ri}(C)$ [replace f with \tilde{f} and dom(f) with $C \cap X$ in Eq. (5.27), and use Prop. 5.3.6].

As a special case of Prop. 5.4.5, we obtain the following.

Proposition 5.4.6: (Subdifferential of Sum of Functions) Let $f_i: \Re^n \mapsto (-\infty, \infty], i = 1, \dots, m$, be convex functions, and assume that the function $F = f_1 + \dots + f_m$ is proper. Then

$$\partial F(x) \supset \partial f_1(x) + \dots + \partial f_m(x), \quad \forall x \in \Re^n.$$

Furthermore, if $\bigcap_{i=1}^m \text{ri}(\text{dom}(f_i)) \neq \emptyset$, we have

$$\partial F(x) = \partial f_1(x) + \dots + \partial f_m(x), \quad \forall x \in \Re^n.$$

More generally, the same is true if for some \overline{m} with $1 \leq \overline{m} \leq m$, the functions f_i , $i = 1, ..., \overline{m}$, are polyhedral and

$$\left(\bigcap_{i=1}^{\overline{m}}\operatorname{dom}(f_i)\right)\cap\left(\bigcap_{i=\overline{m}+1}^{m}\operatorname{ri}\left(\operatorname{dom}(f_i)\right)\right)\neq\emptyset. \tag{5.29}$$

Proof: We can write F in the form F(x) = f(Ax), where A is the matrix defined by Ax = (x, ..., x), and $f : \Re^{mn} \mapsto (-\infty, \infty]$ is the function

$$f(x_1, \ldots, x_m) = f_1(x_1) + \cdots + f_m(x_m).$$

If $\bigcap_{i=1}^{m} \operatorname{ri}(\operatorname{dom}(f_i)) \neq \emptyset$, the range of A contains a point in $\operatorname{ri}(\operatorname{dom}(f))$ (cf. the proof of Prop. 1.3.17). The result then follows from Prop. 5.4.5. If the weaker assumption (5.29) holds, a refinement of the proof of Prop. 5.4.5 is needed. In particular, the optimization problem (5.27) takes the form

minimize
$$\sum_{i=1}^{m} (f_i(y_i) - z_i'd_i)$$
subject to $y_i \in \text{dom}(f_i), \quad z_i = y_i, \quad i = 1, \dots, m.$

By using Prop. 5.3.6 and the assumption (5.29) it follows that strong duality holds and there exists a dual optimal solution for this problem. The proof then proceeds similar to the one of Prop. 5.4.5. Q.E.D.

For an example illustrating the need for the assumption (5.29), let $f_1(x) = -\sqrt{x}$ for $x \ge 0$ and $f_1(x) = \infty$ for x < 0, and let $f_2(x) = f_1(-x)$. Then $\partial f_1(0) = \partial f_2(0) = \emptyset$, but $\partial (f_1 + f_2)(0) = \Re$. We finally note a slight extension of the preceding proposition: instead of assuming that the functions f_i , $i = 1, \ldots, \overline{m}$, are polyhedral, we may assume that they are of the form (5.28), and the same line of proof applies.

5.4.3 Optimality Conditions

It can be seen from the definition of subgradient that a vector x^* minimizes f over \Re^n if and only if $0 \in \partial f(x^*)$. We will now generalize this condition to constrained problems.

Proposition 5.4.7: Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a proper convex function, let X be a nonempty convex subset of \mathbb{R}^n , and assume that one of the following four conditions holds:

- (1) $\operatorname{ri}(\operatorname{dom}(f)) \cap \operatorname{ri}(X) \neq \emptyset$.
- (2) f is polyhedral and $dom(f) \cap ri(X) \neq \emptyset$.
- (3) X is polyhedral and $ri(dom(f)) \cap X \neq \emptyset$.
- (4) f and X are polyhedral, and $dom(f) \cap X \neq \emptyset$.

Then, a vector x^* minimizes f over X if and only if there exists $g \in \partial f(x^*)$ such that -g belongs to the normal cone $N_X(x^*)$, or equivalently,

$$g'(x - x^*) \ge 0, \qquad \forall \ x \in X. \tag{5.30}$$

Proof: The problem can be written as

minimize
$$f(x) + \delta_X(x)$$

subject to $x \in \Re^n$,

where δ_X is the indicator function of X. We have that x^* is an optimal solution if and only if

$$0 \in \partial (f + \delta_X)(x^*) = \partial f(x^*) + \partial \delta_X(x^*) = \partial f(x^*) + N_X(x^*),$$

where the first equality follows from the assumptions (1)-(4) and Prop. 5.4.6, and the second equality follows from the form of subdifferential of δ_X derived in Example 5.4.1. Thus x^* is an optimal solution if and only if there exists $g \in \partial f(x^*)$ such that $-g \in N_X(x^*)$. Q.E.D.

When f is real-valued, the relative interior condition (1) of the preceding proposition is automatically satisfied [we have $dom(f) = \Re^n$]. If in addition, f is differentiable, the optimality condition (5.30) reduces to the one of Prop. 1.1.8:

$$\nabla f(x^*)'(x - x^*) \ge 0, \quad \forall \ x \in X.$$

Also, conditions (2) and (4) can be weakened: f can be assumed to be of the form (5.28), instead of being polyhedral.

5.4.4 Directional Derivatives

Optimization algorithms are often based on iterative cost function improvement. This process is often guided by the directional derivative of the cost function. In the case of a proper convex function $f: \Re^n \mapsto (-\infty, \infty]$, the directional derivative at any $x \in \text{dom}(f)$ in a direction $d \in \Re^n$, is defined by

$$f'(x;d) = \lim_{\alpha \downarrow 0} \frac{f(x+\alpha d) - f(x)}{\alpha}.$$
 (5.31)

An important fact here is that the ratio in Eq. (5.31) is monotonically nonincreasing as $\alpha \downarrow 0$, so that the limit above is well-defined (see Fig. 5.4.6). To verify this, note that for any $\overline{\alpha} > 0$, the convexity of f implies that for all $\alpha \in (0, \overline{\alpha})$,

$$f(x + \alpha d) \le \frac{\alpha}{\overline{\alpha}} f(x + \overline{\alpha} d) + \left(1 - \frac{\alpha}{\overline{\alpha}}\right) f(x) = f(x) + \frac{\alpha}{\overline{\alpha}} \left(f(x + \overline{\alpha} d) - f(x)\right),$$

so that

$$\frac{f(x+\alpha d) - f(x)}{\alpha} \le \frac{f(x+\overline{\alpha}d) - f(x)}{\overline{\alpha}}, \quad \forall \ \alpha \in (0,\overline{\alpha}).$$
 (5.32)

Thus the limit in Eq. (5.31) is well-defined (as a real number, or ∞ , or $-\infty$) and an alternative definition of f'(x;d) is

$$f'(x;d) = \inf_{\alpha > 0} \frac{f(x + \alpha d) - f(x)}{\alpha}, \qquad d \in \Re^n.$$
 (5.33)

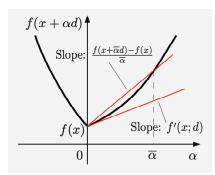


Figure 5.4.6. Illustration of the directional derivative of a convex function f. The ratio

$$\frac{f(x+\alpha d) - f(x)}{\alpha}$$

is monotonically nonincreasing as $\alpha \downarrow 0$ and converges to f'(x;d) [cf. Eqs. (5.32), (5.33)].

It can be shown that $f'(x;\cdot)$ is convex for all $x\in \mathrm{dom}(f)$. For this, it is sufficient to prove that the set of all (d,w) such that f'(x;d)< w is convex, and the verification is straightforward using the convexity of f and Eq. (5.33). If $x\in \mathrm{int}\big(\mathrm{dom}(f)\big)$, we have $f'(x;d)<\infty$ and $f'(x;-d)<\infty$ [cf. Eq. (5.33)], so the convexity of $f'(x;\cdot)$ implies that

$$0 = f'(x;0) \le \frac{1}{2}f'(x;-d) + \frac{1}{2}f'(x;d),$$

or

$$-f'(x; -d) \le f'(x; d), \quad \forall x \in \operatorname{int}(\operatorname{dom}(f)), \ d \in \Re^n.$$

This inequality, combined with $f'(x;d) < \infty$ and $f'(x;-d) < \infty$, shows that

$$-\infty < f'(x;d) < \infty, \quad \forall x \in \operatorname{int}(\operatorname{dom}(f)), d \in \Re^n,$$

i.e., $f'(x; \cdot)$ is real-valued. More generally, the same argument shows that f'(x; d) is a real number for all $x \in ri(dom(f))$ and all d in the subspace that is parallel to aff(dom(f)).

The directional derivative is related to the support function of the subdifferential $\partial f(x)$, as indicated in the following proposition.

Proposition 5.4.8: (Support Function of the Subdifferential) Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let $(\operatorname{cl} f')(x; \cdot)$ be the closure of the directional derivative $f'(x; \cdot)$.

- (a) For all $x \in \text{dom}(f)$ such that $\partial f(x)$ is nonempty, $(\operatorname{cl} f')(x;\cdot)$ is the support function of $\partial f(x)$.
- (b) For all $x \in \text{ri}(\text{dom}(f))$, $f'(x;\cdot)$ is closed and it is the support function of $\partial f(x)$.

Proof: (a) Let us fix $x \in \text{dom}(f)$. From the subgradient inequality, we have

$$g \in \partial f(x)$$
 if and only if $\frac{f(x + \alpha d) - f(x)}{\alpha} \ge g'd$, $\forall d \in \mathbb{R}^n$, $\alpha > 0$.

Therefore, using the equivalent definition (5.33) of directional derivative, we obtain

$$g \in \partial f(x)$$
 if and only if $f'(x;d) \ge g'd, \ \forall \ d \in \Re^n$. (5.34)

Let δ be the conjugate function of $f'(x;\cdot)$:

$$\delta(y) = \sup_{d \in \Re^n} \{ d'y - f'(x; d) \}.$$

Since for any $\gamma > 0$, we have [using Eq. (5.33)]

$$f'(x; \gamma d) = \gamma f'(x; d),$$

it follows that

$$\gamma \delta(y) = \sup_{d \in \Re^n} \left\{ \gamma d' y - \gamma f'(x; d) \right\} = \sup_{d \in \Re^n} \left\{ (\gamma d)' y - f'(x; \gamma d) \right\},$$

and finally

$$\delta(y) = \gamma \delta(y), \quad \forall \gamma > 0.$$

This implies that δ takes only the values 0 and ∞ , so it is the indicator function of some closed convex set, call it Y. In fact, we have

$$Y = \left\{ y \mid \delta(y) \le 0 \right\}$$

$$= \left\{ y \mid \sup_{d \in \Re^n} \left\{ d'y - f'(x; d) \right\} \le 0 \right\}$$

$$= \left\{ y \mid d'y \le f'(x; d), \ \forall \ d \in \Re^n \right\}.$$

Combining this with Eq. (5.34), we obtain $Y = \partial f(x)$.

In conclusion, the conjugate of the convex function $f'(x;\cdot)$ is the indicator function of $\partial f(x)$, which is a proper function. By the Conjugacy Theorem (Prop. 1.6.1), $f'(x;\cdot)$ is also proper, and using the conjugacy between indicator and support function of a closed convex set (cf. Example 1.6.1), it follows that $(\operatorname{cl} f')(x;\cdot)$ is the support function of $\partial f(x)$.

(b) For $x \in \text{ri}(\text{dom}(f))$, by the definition (5.31), f'(x;d) is finite if and only if d belongs to the subspace S that is parallel to aff (dom(f)). Thus, S is equal to both the domain of $f'(x;\cdot)$ and its relative interior. Hence, by Prop. 1.3.15, $f'(x;\cdot)$ coincides with $(\text{cl }f')(x;\cdot)$, and the result follows from part (a), since $\partial f(x) \neq \emptyset$. Q.E.D.

Generally, for all $x \in \text{dom}(f)$ for which $\partial f(x) \neq \emptyset$, we have

$$f'(x;d) \geq \sup_{g \in \partial f(x)} d'g, \qquad \forall \ d \in \Re^n;$$

this follows from the argument of the preceding proof [cf. Eq. (5.34)]. However, strict inequality may hold for x in the relative boundary of dom(f), even if $\partial f(x) \neq \emptyset$. As a result, it is essential to use $(\operatorname{cl} f')(x;\cdot)$ rather than $f'(x;\cdot)$ in Prop. 5.4.8(a). For example, consider the function of two variables

$$f(x_1, x_2) = \begin{cases} 0 & \text{if } x_1^2 + (x_2 - 1)^2 \le 1, \\ \infty & \text{otherwise.} \end{cases}$$

Then the subdifferential at the origin is

$$\partial f(0) = \{x \mid x_1 = 0, x_2 \le 0\},\$$

while for d = (1, 0), we have

$$f'(0;d) = \infty > 0 = \sup_{g \in \partial f(0)} d'g.$$

What is happening here is that $f'(0;\cdot)$ is not closed and does not coincide with $(\operatorname{cl} f')(0;\cdot)$ [the support function of $\partial f(0)$ by the preceding proposition]. Thus, we have

$$\infty=f'(0;d)>(\operatorname{cl} f')(0;d)=\sup_{g\in\partial f(0)}d'g=0.$$

Example 5.4.5: (Directional Derivative and Subdifferential of the Max Function)

Let us derive the directional derivative of the function

$$f(x) = \max\{f_1(x), \dots, f_r(x)\},\$$

where $f_j: \Re^n \to \Re$, j = 1, ..., r, are convex functions. For $x \in \Re^n$, let

$$A_x = \{ j \mid f_j(x) = f(x) \}.$$

For any $x, d \in \Re^n$, and $\alpha > 0$, we have

$$\frac{f(x+\alpha d)-f(x)}{\alpha} \ge \frac{f_j(x+\alpha d)-f_j(x)}{\alpha}, \qquad \forall \ j \in A_x,$$

so by taking the limit as $\alpha \downarrow 0$, we obtain

$$f'(x;d) \ge f_j'(x;d), \qquad \forall \ j \in A_x. \tag{5.35}$$

Consider a sequence $\{\alpha_k\}$ with $\alpha_k \downarrow 0$, and let $x_k = x + \alpha_k d$. For each k, let \overline{j} be an index such that $\overline{j} \in A_{x_k}$ for infinitely many k, and by restricting attention to the corresponding subsequence, assume without loss of generality that $\overline{j} \in A_{x_k}$ for all k. Then,

$$f_{\overline{j}}(x_k) \ge f_j(x_k), \quad \forall k, j,$$

and by taking the limit as $k \to \infty$, and using the continuity of f_j , we have

$$f_{\overline{i}}(x) \ge f_i(x), \quad \forall j.$$

It follows that $\overline{j} \in A_x$, so that

$$f'(x;d) = \lim_{k \to \infty} \frac{f(x + \alpha_k d) - f(x)}{\alpha_k} = \lim_{k \to \infty} \frac{f_{\overline{j}}(x + \alpha_k d) - f_{\overline{j}}(x)}{\alpha_k} = f'_{\overline{j}}(x;d).$$

Combining this relation with Eq. (5.35), we obtain

$$f'(x;d) = \max\{f'_j(x;d) \mid j \in A_x\}, \qquad \forall \ x, d \in \Re^n.$$

Since by Prop. 5.4.8(b), $f'_j(x;\cdot)$ are the support functions of $\partial f_j(x)$, the preceding equation shows that $f'(x;\cdot)$ is the support function of $\bigcup_{j\in A_x}\partial f_j(x)$, and hence also of the closure of $\operatorname{conv}\left(\bigcup_{j\in A_x}\partial f_j(x)\right)$. By Prop. 5.4.1, the sets $\partial f_j(x)$ are compact, so that $\bigcup_{j\in A_x}\partial f_j(x)$ is compact, and hence also by Prop. 1.2.2, $\operatorname{conv}\left(\bigcup_{j\in A_x}\partial f_j(x)\right)$ is compact. On the other hand, by Prop. 5.4.8(b), $f'(x;\cdot)$ is also the support function of $\partial f(x)$. We thus conclude that

$$\partial f(x) = \operatorname{conv} (\cup_{j \in A_x} \partial f_j(x)).$$

Let us finally note that in convex optimization algorithms, directional derivatives and subgradients arise typically in contexts where f is real-valued. In this case there are no anomalies: the directional derivatives are real-valued, and they are the support functions of the corresponding nonempty and compact subdifferentials.

5.5 MINIMAX THEORY

We will now prove theorems regarding the validity of the minimax equality and the existence of saddle points by specializing the MC/MC theorems of Chapter 4. We will assume throughout this section the following.

Assumption 5.5.1: (Convexity/Concavity and Closedness) X and Z are nonempty convex subsets of \Re^n and \Re^m , respectively, and $\phi: X \times Z \mapsto \Re$ is a function such that $\phi(\cdot, z): X \mapsto \Re$ is convex and closed for each $z \in Z$, and $-\phi(x, \cdot): Z \mapsto \Re$ is convex and closed for each $x \in X$.

The analysis will revolve around the function

$$p(u) = \inf_{x \in X} \sup_{z \in Z} \{\phi(x, z) - u'z\}, \qquad u \in \Re^m,$$
 (5.36)

whose epigraph defines the set M used in the MC/MC framework (cf. Section 4.2.5). Note that under Assumption 5.5.1, p is convex. The reason is that p is obtained by partial minimization,

$$p(u) = \inf_{x \in \Re^n} F(x, u),$$

where

$$F(x,u) = \begin{cases} \sup_{z \in Z} \{ \phi(x,z) - u'z \} & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$

is a convex function, since it is the supremum of a collection of convex functions (cf. Props. 1.1.6 and 3.3.1).

5.5.1 Minimax Duality Theorems

We will use the MC/MC Strong Duality Theorem (Prop. 4.3.1) to prove the following proposition.

Proposition 5.5.1: Assume that the function p of Eq. (5.36) satisfies either $p(0) < \infty$, or else $p(0) = \infty$ and $p(u) > -\infty$ for all $u \in \mathbb{R}^m$. Then

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z)$$

if and only if p is lower semicontinuous at u = 0.

Proof: The proof consists of showing that with an appropriate selection of the set M, the assumptions of the proposition are essentially equivalent

to the corresponding assumptions of the MC/MC strong duality theorem (Prop. 4.3.1).

We choose the set M in the MC/MC strong duality theorem to be the epigraph of p,

$$M = \overline{M} = \{(u, w) \mid u \in \Re^m, \ p(u) \le w\},\$$

which is convex in view of the convexity of p noted earlier. Thus, condition (2) of the theorem is satisfied.

From the definition of p, we have

$$w^* = p(0) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z).$$

It follows that the assumption $p(0) < \infty$, or else $p(0) = \infty$ and $p(u) > -\infty$ for all u, is equivalent to condition (1) of the MC/MC strong duality theorem.

Finally, with the preceding definition of M, the condition of lower semicontinuity at u = 0, i.e.,

$$p(0) \le \liminf_{k \to \infty} p(u_k)$$

for all $\{u_k\}$ with $u_k \to 0$, is equivalent to $w^* \leq \liminf_{k \to \infty} w_k$ for every sequence $\{(u_k, w_k)\} \subset M$. Thus, by the conclusion of the MC/MC strong duality theorem, the lower semicontinuity at u = 0 holds if and only if $q^* = w^*$, which is in turn equivalent to the minimax equality (see the discussion following Prop. 4.2.1). Q.E.D.

The finiteness assumptions on p in the preceding proposition are essential. As an example, consider the case where x and z are scalars and

$$\Phi(x,z) = x + z, \qquad X = \{x \mid x \le 0\}, \qquad Z = \{z \mid z \ge 0\}.$$

We have

$$p(u) = \inf_{x \le 0} \sup_{z \ge 0} \{x + z - uz\} = \begin{cases} \infty & \text{if } u < 1, \\ -\infty & \text{if } u \ge 1, \end{cases}$$

so p is closed convex, but the finiteness assumptions of the proposition are violated, and we also have

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = -\infty < \infty = \inf_{x \in X} \sup_{z \in Z} \phi(x, z).$$

We will now use the MC/MC existence of solutions results of Section 4.4 (cf. Props. 4.4.1 and 4.4.2) to derive conditions for the attainment of the supremum in the minimax equality. What we need is an assumption that 0 lies in the relative interior or the interior of dom(p) and that $p(0) > -\infty$.

We then obtain the following result, which also asserts that the supremum in the minimax equality is attained (this follows from the corresponding attainment assertions of Props. 4.4.1 and 4.4.2).

Proposition 5.5.2: Assume that $0 \in ri(dom(p))$ and $p(0) > -\infty$. Then

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z),$$

and the supremum over Z in the left-hand side is finite and is attained. Furthermore, the set of $z \in Z$ attaining this supremum is compact if and only if 0 lies in the interior of dom(p).

The preceding two minimax theorems indicate that the properties of the function p around u=0 are critical to guarantee the minimax equality. Here is an illustrative example:

Example 5.5.1:

Let

$$X = \left\{ x \in \Re^2 \mid x \ge 0 \right\}, \quad Z = \left\{ z \in \Re \mid z \ge 0 \right\}, \quad \phi(x, z) = e^{-\sqrt{x_1 x_2}} + z x_1,$$

which can be shown to satisfy the convexity/concavity and closedness Assumption 5.5.1. (This is essentially the constrained minimization Example 5.3.2, recast as a minimax problem.) For all $z \ge 0$, we have

$$\inf_{x>0} \phi(x,z) = \inf_{x\geq 0} \left\{ e^{-\sqrt{x_1 x_2}} + zx_1 \right\} = 0,$$

since the expression in braces is nonnegative for $x \ge 0$ and can approach zero by taking $x_1 \to 0$ and $x_1x_2 \to \infty$. Hence,

$$\sup_{z \ge 0} \inf_{x > 0} \phi(x, z) = 0.$$

We also have for all $x \geq 0$,

$$\sup_{z \ge 0} \phi(x, z) = \sup_{z \ge 0} \left\{ e^{-\sqrt{x_1 x_2}} + z x_1 \right\} = \begin{cases} 1 & \text{if } x_1 = 0, \\ \infty & \text{if } x_1 > 0. \end{cases}$$

Hence,

$$\inf_{x>0} \sup_{z>0} \phi(x,z) = 1,$$

so

$$\inf_{x \ge 0} \sup_{z \ge 0} \phi(x, z) > \sup_{z \ge 0} \inf_{x \ge 0} \phi(x, z).$$

Here, the function p is given by

$$p(u) = \inf_{x \ge 0} \sup_{z \ge 0} \left\{ e^{-\sqrt{x_1 x_2}} + z(x_1 - u) \right\} = \begin{cases} \infty & \text{if } u < 0, \\ 1 & \text{if } u = 0, \\ 0 & \text{if } u > 0. \end{cases}$$

Thus, p is not lower semicontinuous at 0, the assumptions of Props. 5.5.1 are violated, and the minimax equality does not hold.

The following example illustrates how the minimax equality may hold, while the supremum over $z \in Z$ is not attained because the relative interior assumption of Prop. 5.5.2 is not satisfied.

Example 5.5.2:

Let

$$X = \Re, \qquad Z = \{z \in \Re \mid z \ge 0\}, \qquad \phi(x, z) = x + zx^2,$$

which satisfy Assumption 5.5.1. (This is essentially the constrained minimization Example 5.3.3, recast as a minimax problem.) For all $z \geq 0$, we have

$$\inf_{x \in \Re} \phi(x, z) = \inf_{x \in \Re} \left\{ x + zx^2 \right\} = \left\{ \begin{array}{ll} -1/(4z) & \text{if } z > 0, \\ -\infty & \text{if } z = 0. \end{array} \right.$$

Hence,

$$\sup_{z>0}\inf_{x\in\Re}\phi(x,z)=0.$$

We also have for all $x \in \Re$,

$$\sup_{z \ge 0} \phi(x, z) = \sup_{z \ge 0} \left\{ x + zx^2 \right\} = \begin{cases} 0 & \text{if } x = 0, \\ \infty & \text{if } x \ne 0. \end{cases}$$

Hence,

$$\inf_{x \in \Re} \sup_{z \ge 0} \phi(x, z) = 0,$$

and the minimax equality holds. However, the problem

maximize
$$\inf_{x \in \Re} \phi(x, z)$$

subject to $z \in Z$

does not have an optimal solution. Here we have

$$F(x,u) = \sup_{z \ge 0} \{x + zx^2 - uz\} = \begin{cases} x & \text{if } x^2 \le u, \\ \infty & \text{if } x^2 > u, \end{cases}$$

and

$$p(u) = \inf_{x \in \Re} \sup_{z \ge 0} F(x, u) = \begin{cases} -\sqrt{u} & \text{if } u \ge 0, \\ \infty & \text{if } u < 0. \end{cases}$$

It can be seen that $0 \notin ri(dom(p))$, thus violating the assumption of Prop. 5.5.2.

5.5.2 Saddle Point Theorems

We will now use the two minimax theorems just derived (Props. 5.5.1 and 5.5.2) to obtain more specific conditions for the validity of the minimax equality and the existence of saddle points. Again, Assumption 5.5.1 is assumed throughout. The preceding analysis has underscored the importance of the function

$$p(u) = \inf_{x \in \Re^n} F(x, u), \tag{5.37}$$

where

$$F(x,u) = \begin{cases} \sup_{z \in Z} \{\phi(x,z) - u'z\} & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$
 (5.38)

and suggests a two-step process to ascertain the validity of the minimax equality and the existence of a saddle point:

- (1) Show that p is closed and convex, thereby showing that the minimax equality holds by using Prop. 5.5.1.
- (2) Verify that the infimum of $\sup_{z\in Z} \phi(x,z)$ over $x\in X$, and the supremum of $\inf_{x\in X} \phi(x,z)$ over $z\in Z$ are attained, thereby showing that the set of saddle points is nonempty (cf. Prop. 3.4.1).

Step (1) requires two types of assumptions:

- (a) The convexity/concavity and closedness Assumption 5.5.1. This guarantees that F is convex and closed (being the pointwise supremum over $z \in Z$ of closed convex functions), and also guarantees that p is convex.
- (b) Conditions that guarantee that the partial minimization in the definition of p preserves closedness, so that p is also closed.

Step (2) requires that either Weierstrass' Theorem can be applied, or else that some other suitable condition for existence of an optimal solution is satisfied (cf. Section 3.3). Fortunately, conditions that guarantee that the partial minimization in the definition of F preserves closedness as in (b), also guarantee the existence of corresponding optimal solutions.

As an example of this line of analysis, we obtain the following classical result.

Proposition 5.5.3: (Classical Saddle Point Theorem) Let the sets X and Z be compact. Then the set of saddle points of ϕ is nonempty and compact.

Proof: We note that F is convex and closed. Using the convexity of F and Prop. 3.3.1, we see that the function p of Eq. (5.37) is convex. Using the compactness of Z, F is real-valued over $X \times \Re^m$, and from the compactness

of X and Prop. 3.3.3, it follows that p is also real-valued and therefore continuous. Hence, the minimax equality holds by Prop. 5.5.1.

Finally, the function $\sup_{z\in Z}\phi(x,z)$ is equal to F(x,0), so it is closed, and the set of its minima over $x\in X$ is nonempty and compact by Weierstrass' Theorem. Similarly the set of maxima of the function $\inf_{x\in X}\phi(x,z)$ over $z\in Z$ is nonempty and compact. From Prop. 3.4.1 it follows that the set of saddle points is nonempty and compact. **Q.E.D.**

We will now derive alternative and more general saddle point theorems, using a similar line of analysis. To formulate these theorems, we consider the functions $t: \mathbb{R}^n \mapsto (-\infty, \infty]$ and $r: \mathbb{R}^m \mapsto (-\infty, \infty]$ given by

$$t(x) = \begin{cases} \sup_{z \in Z} \phi(x, z) & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$

and

$$r(z) = \begin{cases} -\inf_{x \in X} \phi(x, z) & \text{if } z \in Z, \\ \infty & \text{if } z \notin Z. \end{cases}$$

Note that by Assumption 5.5.1, t is closed and convex, being the supremum of closed and convex functions. Furthermore, since $t(x) > -\infty$ for all x, we have

t is proper if and only if
$$\sup_{x \in X} \phi(x, z) < \infty$$
.

Similarly, the function r is closed and convex, and

$$r$$
 is proper if and only if $-\infty < \sup_{z \in Z} \inf_{x \in X} \phi(x, z)$.

The next two propositions provide conditions for the minimax equality to hold. These propositions are subsequently used to prove results about nonemptiness and compactness of the set of saddle points.

Proposition 5.5.4: Assume that t is proper and that the level sets $\{x \mid t(x) \leq \gamma\}, \gamma \in \Re$, are compact. Then

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z)$$

and the infimum over X in the right-hand side above is attained at a set of points that is nonempty and compact.

Proof: The function p is defined by partial minimization of the function F of Eq. (5.38), i.e.,

$$p(u) = \inf_{x \in \Re^n} F(x, u)$$

[cf. Eq. (5.37)]. We note that

$$t(x) = F(x, 0),$$

so F is proper, since t is proper and F is closed [by the Recession Cone Theorem (Prop. 1.4.1), $\operatorname{epi}(F)$ contains a vertical line if and only if $\operatorname{epi}(t)$ contains a vertical line]. Furthermore, the compactness assumption on the level sets of t can be translated to the compactness assumption of Prop. 3.3.2 (with 0 playing the role of the vector \overline{x}). It follows from the result of that proposition that p is closed and proper, and that p(0) is finite. By Prop. 5.5.1, it follows that the minimax equality holds. Finally, the infimum over X in the right-hand side of the minimax equality is attained at the set of minima of t, which is nonempty and compact since t is proper and has compact level sets. **Q.E.D.**

Example 5.5.3:

Let us show that

$$\min_{\|x\| \le 1} \max_{z \in S+C} x'z = \max_{z \in S+C} \min_{\|x\| \le 1} x'z,$$

where S is a subspace, and C is a nonempty, convex, and compact subset of \Re^n . By defining

$$X = \left\{x \mid \|x\| \le 1\right\}, \qquad Z = S + C,$$

$$\phi(x, z) = x'z, \qquad \forall \ (x, z) \in X \times Z,$$

we see that Assumption 5.5.1 is satisfied, so we can apply Prop. 5.5.4. We have

$$\begin{split} t(x) &= \begin{cases} \sup_{z \in S+C} x'z & \text{if } \|x\| \leq 1, \\ \infty & \text{otherwise,} \end{cases} \\ &= \begin{cases} \sup_{z \in C} x'z & \text{if } x \in S^{\perp}, \, \|x\| \leq 1, \\ \infty & \text{otherwise.} \end{cases} \end{split}$$

Since $\sup_{z \in C} x'z$, viewed as a function of x over \Re^n , is continuous, the level sets of t are compact. It follows from Prop. 5.5.4 that the minimax equality holds. It also turns out that a saddle point exists. We will show this shortly, after we develop some additional machinery.

Proposition 5.5.5: Assume that t is proper, and that the recession cone and the constancy space of t are equal. Then

$$\sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z)$$

and the infimum over X in the right-hand side above is attained.

Proof: The proof is similar to the one of Prop. 5.5.4. We use Prop. 3.3.4 in place of Prop. 3.3.2. **Q.E.D.**

By combining the preceding two propositions, we obtain conditions for existence of a saddle point.

Proposition 5.5.6: Assume that either t is proper or r is proper.

- (a) If the level sets $\{x \mid t(x) \leq \gamma\}$ and $\{z \mid r(z) \leq \gamma\}$, $\gamma \in \Re$, of t and r are compact, the set of saddle points of ϕ is nonempty and compact.
- (b) If the recession cones of t and r are equal to the constancy spaces of t and r, respectively, the set of saddle points of ϕ is nonempty.

Proof: We assume that t is proper. If instead r is proper, we reverse the roles of x and z.

(a) From Prop. 5.5.4, it follows that the minimax equality holds, and that the infimum over X of $\sup_{z\in Z}\phi(x,z)$ is finite and is attained at a nonempty and compact set. Therefore,

$$-\infty < \sup_{z \in Z} \inf_{x \in X} \phi(x, z) = \inf_{x \in X} \sup_{z \in Z} \phi(x, z) < \infty,$$

and we can reverse the roles of x and z, and apply Prop. 5.5.4 again to show that the supremum over Z of $\inf_{x \in X} \phi(x, z)$ is attained at a nonempty and compact set.

(b) The proof is similar to the one of part (a), except that we use Prop. 5.5.5 instead of Prop. 5.5.4. Q.E.D.

Example 5.5.2: (continued)

To illustrate the difference between Props. 5.5.4 and 5.5.6, let

$$X = \Re, \qquad Z = \{z \in \Re \mid z > 0\}, \qquad \phi(x, z) = x + zx^2.$$

A straightforward calculation yields

$$t(x) = \begin{cases} 0 & \text{if } x = 0, \\ \infty & \text{if } x \neq 0, \end{cases} \qquad r(z) = \begin{cases} \frac{1}{4z} & \text{if } z > 0, \\ \infty & \text{if } z \leq 0. \end{cases}$$

Thus, t satisfies the assumptions of Prop. 5.5.4 and the minimax equality holds, but r violates the assumptions of Prop. 5.5.6 and there is no saddle point, since the supremum over z of $-r(z) = \inf_{x \in \Re} \Phi(x, z)$ is not attained.

Example 5.5.3: (continued)

Let

$$X = \{x \mid ||x|| \le 1\}, \qquad Z = S + C, \qquad \phi(x, z) = x'z,$$

where S is a subspace, and C is a nonempty, convex, and compact subset of $\Re^n.$ We have

$$\begin{split} t(x) &= \begin{cases} \sup_{z \in S+C} x'z & \text{if } \|x\| \leq 1, \\ \infty & \text{otherwise,} \end{cases} \\ &= \begin{cases} \sup_{z \in C} x'z & \text{if } x \in S^{\perp}, \, \|x\| \leq 1, \\ \infty & \text{otherwise,} \end{cases} \end{split}$$

$$\begin{split} r(z) &= \begin{cases} \sup_{\|x\| \leq 1} -x'z & \text{if } z \in S+C, \\ \infty & \text{otherwise,} \end{cases} \\ &= \begin{cases} \|z\| & \text{if } z \in S+C, \\ \infty & \text{otherwise.} \end{cases} \end{split}$$

Since the level sets of both t and r are compact, it follows from Prop. 5.5.6(a) that the minimax equality holds, and that the set of saddle points is nonempty and compact.

The compactness of the level sets of t can be guaranteed by simpler sufficient conditions. In particular, the level sets $\{x \mid t(x) \leq \gamma\}$ are compact if any one of the following two conditions holds:

- (1) The set X is compact [since $\{x \mid t(x) \leq \gamma\}$ is closed, by the closedness of t, and is contained in X].
- (2) For some $\overline{z} \in Z$, $\overline{\gamma} \in \Re$, the set

$$\left\{ x \in X \mid \phi(x, \overline{z}) \le \overline{\gamma} \right\}$$

is nonempty and compact [since then all sets $\{x \in X \mid \phi(x,\overline{z}) \leq \gamma\}$ are compact, and a nonempty level set $\{x \mid t(x) \leq \gamma\}$ is contained in $\{x \in X \mid \phi(x,\overline{z}) \leq \gamma\}$].

Furthermore, any one of above two conditions also guarantees that r is proper; for example under condition (2), the infimum over $x \in X$ in the relation

$$r(\overline{z}) = -\inf_{x \in X} \phi(x, \overline{z})$$

is attained by Weierstrass' Theorem, so that $r(\overline{z}) < \infty$.

By a symmetric argument, we also see that the level sets of r are compact under any one of the following two conditions:

- (1) The set Z is compact.
- (2) For some $\overline{x} \in X$, $\overline{\gamma} \in \Re$, the set

$$\{z \in Z \mid \phi(\overline{x}, z) \geq \overline{\gamma}\}$$

is nonempty and compact.

Furthermore, any one of these conditions guarantees that t is proper.

Thus, by combining the preceding discussion and Prop. 5.5.6(a), we obtain the following result, which generalizes the classical saddle point theorem (Prop. 5.5.3), and provides sufficient conditions for the existence of a saddle point.

Proposition 5.5.7: (Saddle Point Theorem) The set of saddle points of ϕ is nonempty and compact under any one of the following conditions:

- (1) X and Z are compact.
- (2) Z is compact, and for some $\overline{z} \in Z$, $\overline{\gamma} \in \Re$, the level set

$$\{x \in X \mid \phi(x, \overline{z}) \le \overline{\gamma}\}\$$

is nonempty and compact.

(3) X is compact, and for some $\overline{x} \in X$, $\overline{\gamma} \in \Re$, the level set

$$\{z \in Z \mid \phi(\overline{x}, z) \ge \overline{\gamma}\}$$

is nonempty and compact.

(4) For some $\overline{x} \in X$, $\overline{z} \in Z$, $\overline{\gamma} \in \Re$, the level sets

$$\big\{x\in X\mid \phi(x,\overline{z})\leq \overline{\gamma}\big\}, \qquad \big\{z\in Z\mid \phi(\overline{x},z)\geq \overline{\gamma}\big\},$$

are nonempty and compact.

Proof: From the discussion preceding the proposition, it is seen that under Assumption 5.5.1, t and r are proper, and the level sets of t and r are compact. The result follows from Prop. 5.5.6(a). Q.E.D.

5.6 THEOREMS OF THE ALTERNATIVE

Theorems of the alternative are important tools in optimization, which address the feasibility (possibly strict) of affine inequalities. We will show that these theorems can be viewed as special cases of MC/MC duality (cf. Props. 4.4.1, 4.4.2). This connection is demonstrated in the proofs of some classical theorems.

One of them is Gordan's Theorem, which dates to 1873. In its classical version it states that there exists a vector $x \in \mathbb{R}^n$ such that

 $a_1'x < 0, \ldots, a_r'x < 0$, if and only if $\operatorname{cone}(\{a_1, \ldots, a_r\})$ does not contain a line (both vectors d and -d for some $d \neq 0$); see alternatives (i) and (ii) in the following proposition, for b = 0. We provide a slight extension of this theorem.

Proposition 5.6.1: (Gordan's Theorem) Let A be an $m \times n$ matrix and b be a vector in \Re^m . The following are equivalent:

(i) There exists a vector $x \in \Re^n$ such that

$$Ax < b$$
.

(ii) For every vector $\mu \in \Re^m$,

$$A'\mu = 0, \quad b'\mu \le 0, \quad \mu \ge 0 \qquad \Rightarrow \qquad \mu = 0.$$

(iii) Any polyhedral set of the form

$$\{\mu \mid A'\mu = c, b'\mu \le d, \mu \ge 0\},$$
 (5.39)

where $c \in \mathbb{R}^n$ and $d \in \mathbb{R}$, is compact.

Proof: We will show that (i) and (ii) are equivalent, and then that (ii) and (iii) are equivalent. The equivalence of (i) and (ii) is geometrically evident, once the proper MC/MC framework is considered (see Fig. 5.6.1). For completeness we provide the (somewhat lengthy) details. Consider the set

$$M = \{(u, w) \mid w \ge 0, Ax - b \le u \text{ for some } x \in \Re^n\},\$$

its projection on the u axis

$$D = \{u \mid Ax - b \le u \text{ for some } x \in \Re^n\},\$$

and the corresponding MC/MC framework. Let w^* and q^* be the min common and max crossing values, respectively. Clearly, the system $Ax \leq b$ has a solution if and only if $w^* = 0$. Also, if (ii) holds, we have

$$A'\mu = 0, \quad \mu \ge 0 \qquad \Rightarrow \qquad b'\mu \ge 0,$$

which by the linear version of Farkas' Lemma (Prop. 5.1.2), implies that the system $Ax \leq b$ has a solution. In conclusion, both (i) and (ii) imply that the system $Ax \leq b$ has a solution, which is in turn equivalent to $w^* = q^* = 0$. Thus in proving the equivalence of (i) and (ii), we may assume that the system $Ax \leq b$ has a solution and $w^* = q^* = 0$.

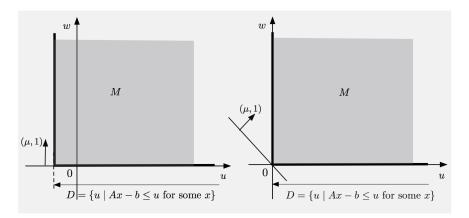


Figure 5.6.1. MC/MC framework for showing the equivalence of conditions (i) and (ii) of Gordan's Theorem. We consider the set

$$M = \left\{ (u, w) \mid w \ge 0, \, Ax - b \le u \text{ for some } x \in \Re^n \right\},\,$$

its projection on the u axis

$$D = \{ u \mid Ax - b \le u \text{ for some } x \in \mathbb{R}^n \},$$

and the corresponding MC/MC framework. As shown by the figure on the left, condition (i) of Gordan's Theorem is equivalent to $0 \in \operatorname{int}(D)$, and is also equivalent to 0 (the horizontal hyperplane) being the unique optimal solution of the max crossing problem. It is seen that the latter condition can be written as

$$\mu \ge 0$$
, $0 \le \mu'(Ax - b) + w$, $\forall x \in \Re^n, w \ge 0$ $\Rightarrow \mu = 0$

which is equivalent to condition (ii). In the figure on the right both conditions (i) and (ii) are violated.

The cost function of the max crossing problem is

$$q(\mu) = \inf_{(u,w) \in M} \{ w + \mu' u \}.$$

Since M contains vectors (u, w) with arbitrarily large components [for each $(u, w) \in M$, we have $(\overline{u}, \overline{w}) \in M$ for all $\overline{u} \geq u$ and $\overline{w} \geq w$], it follows that $q(\mu) = -\infty$ for all μ that are not in the nonnegative orthant, and we have

$$q(\mu) = \begin{cases} \inf_{Ax - b \le u} \mu' u & \text{if } \mu \ge 0, \\ -\infty & \text{otherwise,} \end{cases}$$

or equivalently,

$$q(\mu) = \begin{cases} \inf_{x \in \Re^n} \mu'(Ax - b) & \text{if } \mu \ge 0, \\ -\infty & \text{otherwise.} \end{cases}$$
 (5.40)

We now note that condition (i) is equivalent to 0 being an interior point of D (to see this note that if \overline{x} satisfies $A\overline{x} - b < 0$, then the set $\{u \mid A\overline{x} - b \leq u\}$ contains 0 in its interior and is contained in the set D). By Prop. 4.4.2, it follows that (i) is equivalent to $w^* = q^* = 0$ and the set of optimal solutions of the max crossing problem being nonempty and compact. Thus, using the form (5.40) of q, we see that (i) is equivalent to $\mu = 0$ being the only $\mu \geq 0$ satisfying $q(\mu) \geq 0$ or equivalently, satisfying $A'\mu = 0$ and $\mu'b \leq 0$. It follows that (i) is equivalent to (ii).

To show that (ii) is equivalent to (iii), we note that the recession cone of the set (5.39) is

$$\{\mu \mid A'\mu = 0, b'\mu \le 0, \mu \ge 0\}.$$

Thus (ii) states that the recession cone of the set (5.39) consists of just the origin, which is equivalent to (iii); cf. Prop. 1.4.2(a). Q.E.D.

There are several theorems of the alternative involving strict inequalities. Among the ones involving linear constraints, the following is the most general.

Proposition 5.6.2: (Motzkin's Transposition Theorem) Let A and B be $p \times n$ and $q \times n$ matrices, and let $b \in \Re^p$ and $c \in \Re^q$ be vectors. The system

$$Ax < b$$
, $Bx < c$

has a solution if and only if for all $\mu \in \Re^p$ and $\nu \in \Re^q$, with $\mu \geq 0$, $\nu \geq 0$, the following two conditions hold:

$$A'\mu + B'\nu = 0 \qquad \Rightarrow \qquad b'\mu + c'\nu \ge 0, \tag{5.41}$$

$$A'\mu + B'\nu = 0, \ \mu \neq 0 \qquad \Rightarrow \qquad b'\mu + c'\nu > 0.$$
 (5.42)

Proof: Consider the set

$$M = \{(u, w) \mid w \ge 0, Ax - b \le u \text{ for some } x \text{ with } Bx \le c\},\$$

its projection on the u axis

$$D = \{u \mid Ax - b \le u \text{ for some } x \text{ with } Bx \le c\},\$$

and the corresponding MC/MC framework. The cost function of the \max crossing problem is

$$q(\mu) = \begin{cases} \inf_{(u,w) \in M} \{w + \mu'u\} & \text{if } \mu \geq 0, \\ -\infty & \text{otherwise,} \end{cases}$$

or

$$q(\mu) = \begin{cases} \inf_{Ax - b \le u, Bx \le c} \mu' u & \text{if } \mu \ge 0, \\ -\infty & \text{otherwise.} \end{cases}$$
 (5.43)

Similar to the proof of Gordan's Theorem (Prop. 5.6.1), and using the linear version of Farkas' Lemma (Prop. 5.1.2), we may assume that the system $Ax \leq b$, $Bx \leq c$ has a solution.

Now the system Ax < b, $Bx \le c$ has a solution if and only if 0 is an interior point of D, so by Prop. 4.4.2, it follows that the system Ax < b, $Bx \le c$ has a solution if and only if the set of optimal solutions of the max crossing problem is nonempty and compact. Since q(0) = 0 and $\sup_{\mu \in \Re^m} q(\mu) = 0$, using also Eq. (5.43), we see that the set of dual optimal solutions is nonempty and compact if and only if $q(\mu) < 0$ for all $\mu \ge 0$ with $\mu \ne 0$ [if $q(\mu) = 0$ for some nonzero $\mu \ge 0$, it can be seen from Eq. (5.43) that we must have $q(\gamma\mu) = 0$ for all $\gamma > 0$]. We will complete the proof by showing that these conditions hold if and only if conditions (5.41) and (5.42) hold for all $\mu \ge 0$ and $\nu \ge 0$.

We calculate the infimum of the linear program in (x, u) in the right-hand side of Eq. (5.43) by using duality, and simplifying the dual problem. In particular, we have for all $\mu \geq 0$, after a straightforward calculation,

$$q(\mu) = \inf_{Ax - b \le u, \, Bx \le c} \mu' u = \sup_{A'\mu + B'\nu = 0, \, \nu \ge 0} (-b'\mu - c'\nu).$$

(In these relations, μ is fixed; the infimum is over (x, u), and the supremum is over ν .) Using this equation, we see that q(0) = 0 and $q(\mu) < 0$ for all $\mu \geq 0$ with $\mu \neq 0$ if and only if conditions (5.41) and (5.42) hold for all $\mu \geq 0$ and $\nu \geq 0$. **Q.E.D.**

Let us derive another interesting theorem of the alternative. Like the Gordan and Motzkin Theorems, it can be proved by means of a suitable MC/MC formulation. We will give an alternative proof that uses Motzkin's theorem. (Actually, Motzkin's Theorem, being the most general theorem of the alternative that involves linear equations and inequalities, can also be similarly used to prove Gordan's Theorem. However, the proof we gave, based on the connection with the MC/MC framework, is more intuitive.)

Proposition 5.6.3: (Stiemke's Transposition Theorem) Let A be an $m \times n$ matrix, and let c be a vector in \Re^m . The system

$$Ax = c, \qquad x > 0$$

has a solution if and only if

$$A'\mu \ge 0$$
 and $c'\mu \le 0$ \Rightarrow $A'\mu = 0$ and $c'\mu = 0$.

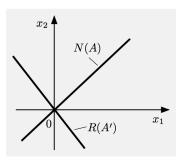


Figure 5.6.2. Geometric interpretation of Stiemke's Transposition Theorem for the case c = 0. The nullspace N(A) of A meets the interior of the nonnegative orthant if and only if the range R(A') of A' meets the nonnegative orthant at just the origin.

Proof: We first consider the case where c=0, i.e., we show that the system Ax=0, x>0 has a solution if and only if

$$A'\mu \ge 0 \qquad \Rightarrow \qquad A'\mu = 0.$$

Indeed, we can equivalently write the system Ax = 0, x > 0 in the form

$$-x < 0$$
, $Ax \le 0$, $-Ax \le 0$,

so that Motzkin's Theorem can be applied with the identifications $A \sim -I$, $B \sim [A - A], b \sim 0, c \sim 0$. We obtain that the system has a solution if and only if the system

$$-y + A'z^{+} - A'z^{=0}, y \ge 0, z^{+} \ge 0, z^{-} \ge 0,$$

has no solution with $y \neq 0$, or equivalently, by introducing the transformation $\mu = z^+ - z^-$, if and only if the system

$$A'\mu \ge 0, \qquad A'\mu \ne 0$$

has no solution.

We now assume that $c \neq 0$, and we note that the system Ax = c, x > 0 has a solution if and only if the system

$$(A -c)\begin{pmatrix} x \\ z \end{pmatrix} = 0, \qquad x > 0, \ z > 0,$$

has a solution. By applying the result already shown for the case c=0, we see that the above system has a solution if and only if

$$\begin{pmatrix} A' \\ -c' \end{pmatrix} \mu \ge 0 \qquad \Rightarrow \qquad \begin{pmatrix} A' \\ -c' \end{pmatrix} \mu = 0,$$

or equivalently if and only if

$$A'\mu \ge 0$$
 and $c'\mu \le 0$ \Rightarrow $A'\mu = 0$ and $c'\mu = 0$.

Q.E.D.

A geometric interpretation of Stiemke's Theorem for the case c=0 is given in Fig. 5.6.2.

Compactness of the Solution Set of a Linear Program

The theorems of Gordan and Stiemke (Props. 5.6.1 and 5.6.3) can be used to provide necessary and sufficient conditions for the compactness of the primal and the dual optimal solution sets of linear programs (cf. Example 4.2.1). We say that the primal linear program

minimize
$$c'x$$

subject to $a'_j x \ge b_j$, $j = 1, ..., r$, (5.44)

is strictly feasible if there exists a primal-feasible vector $x \in \Re^n$ with $a'_j x > b_j$ for all $j = 1, \ldots, r$. Similarly, we say that the dual linear program

maximize
$$b'\mu$$

subject to $\sum_{j=1}^{r} a_j \mu_j = c$, $\mu \ge 0$, (5.45)

is strictly feasible if there exists a dual-feasible vector μ with $\mu > 0$. We have the following proposition. Part (b) requires that the matrix defining the constraints has rank n, the dimension of the primal vector x.

Proposition 5.6.4: Consider the primal and dual linear programs [cf. Eqs. (5.44), (5.45)], and assume that their common optimal value is finite. Then:

- (a) The dual optimal solution set is compact if and only if the primal problem is strictly feasible.
- (b) Assuming that the set $\{a_1, \ldots, a_r\}$ contains n linearly independent vectors, the primal optimal solution set is compact if and only if the dual problem is strictly feasible.

Proof: We first note that, by the duality theorem (Prop. 5.2.1), the finiteness of the optimal value of the primal problem implies that the dual optimal solution set, denoted D^* , is nonempty.

(a) The dual optimal solution set is

$$D^* = \{ \mu \mid \mu_1 a_1 + \dots + \mu_r a_r = c, \, \mu_1 b_1 + \dots + \mu_r b_r \ge q^*, \, \mu \ge 0 \},$$

The result now follows from the equivalence of conditions (i) and (iii) in Gordan's Theorem (Prop. 5.6.1).

(b) The primal optimal solution set is

$$P^* = \{x \mid c'x \le f^*, \ a'_i x \ge b_j, \ j = 1, \dots, r\}.$$

The dual problem is strictly feasible if there exists $\mu > 0$ such that $A\mu = c$, where A is the matrix with columns a_j , j = 1, ..., r. From Stiemke's Transposition Theorem (Prop. 5.6.3), we see that this is true if and only if every vector d in the recession cone of P^* , which is

$$R_{P^*} = \{d \mid c'd \leq 0, \ a'_i d \geq 0, \ j = 1, \dots, r\},\$$

satisfies c'd = 0 and $a'_j d = 0$ for all j, or just $a'_j d = 0$ for all j (since c is a linear combination of a_1, \ldots, a_r , by the feasibility of the dual problem). Since the set $\{a_1, \ldots, a_r\}$ contains n linearly independent vectors, it follows that R_{P^*} consists of just the origin. Hence, by Prop. 1.4.2(a), it follows that P^* is compact. **Q.E.D.**

For an example where the dual problem is strictly feasible, but the primal optimal solution set is not compact, consider the problem

minimize
$$x_1 + x_2$$

subject to $x_1 + x_2 \ge 1$.

The dual problem,

maximize
$$\mu$$
 subject to $\begin{pmatrix} 1 \\ 1 \end{pmatrix} \mu = \begin{pmatrix} 1 \\ 1 \end{pmatrix}, \quad \mu \geq 0,$

is strictly feasible, but the primal problem has an unbounded optimal solution set. The difficulty here is that the linear independence assumption of Prop. 5.6.4(b) is violated (since n=2 but r=1).

5.7 NONCONVEX PROBLEMS

In this section we focus on the MC/MC framework in the absence of the convex structure that guarantees strong duality. We aim to estimate the duality gap. Consider first the case where the set M is the epigraph of a function $p: \Re^n \mapsto [-\infty, \infty]$. Then $w^* = p(0)$, and as discussed in Section 4.2.1, $q(\mu) = -p^*(-\mu)$, where p^* is the conjugate of p. Thus

$$q^* = \sup_{\mu \in \mathbb{R}^n} q(\mu) = \sup_{\mu \in \mathbb{R}^n} \left\{ 0 \cdot (-\mu) - p^*(-\mu) \right\} = p^{**}(0), \tag{5.46}$$

where $p^{\star\star}$ is the conjugate of p^{\star} (double conjugate of p). The duality gap is

$$w^* - q^* = p(0) - p^{**}(0),$$

and can be interpreted as a "measure of nonconvexity" of p at 0.

Consider next the case where M has the form

$$M = \tilde{M} + \{(u, 0) \mid u \in C\},\$$

where \tilde{M} is a compact set and C is a closed convex set, or more generally, where $\operatorname{conv}(M)$ is closed and $q^* > -\infty$. Then, by Prop. 4.3.2(c), the duality gap $w^* - q^*$ is equal to $w^* - \overline{w}^*$, where w^* and \overline{w}^* are the min common values corresponding to the sets M and $\operatorname{conv}(M)$.

In either case, we see that estimates of the duality gap should depend on how much M differs from its convex hull along the vertical axis. A special case where interesting estimates of this type can be obtained arises in separable optimization problems, as we now discuss.

5.7.1 Duality Gap in Separable Problems

Suppose that x consists of m components x_1, \ldots, x_m of corresponding dimensions n_1, \ldots, n_m , and the problem has the form

minimize
$$\sum_{i=1}^{m} f_i(x_i)$$
subject to
$$\sum_{i=1}^{m} g_i(x_i) \le 0, \quad x_i \in X_i, \quad i = 1, \dots, m,$$

$$(5.47)$$

where $f_i: \Re^{n_i} \mapsto \Re$ and $g_i: \Re^{n_i} \mapsto \Re^r$ are given functions, and X_i are given subsets of \Re^{n_i} . An important nonconvex instance is when the functions f_i and g_i are linear, and the sets X_i are finite, in which case we obtain a discrete/integer programming problem.

We call a problem of the form (5.47) *separable*. Its salient feature is that the minimization involved in the calculation of the dual function

$$q(\mu) = \inf_{\substack{x_i \in X_i \\ i=1,\dots,m}} \left\{ \sum_{i=1}^m \left(f_i(x_i) + \mu' g_i(x_i) \right) \right\} = \sum_{i=1}^m q_i(\mu),$$

is decomposed into the m simpler minimizations

$$q_i(\mu) = \inf_{x_i \in X_i} \{ f_i(x_i) + \mu' g_i(x_i) \}, \quad i = 1, \dots, m.$$

These minimizations are often conveniently done either analytically or computationally, in which case the dual function can be easily evaluated.

When the cost and/or the constraints are not convex, the separable structure is helpful in another, somewhat unexpected way. In particular, in this case the duality gap turns out to be relatively small and can often be shown to diminish to zero relative to the optimal primal value as the number

m of separable terms increases. As a result, one can often obtain a near-optimal primal solution, starting from a dual optimal solution. In integer programming problems, this may obviate the need for computationally intensive procedures such as branch-and-bound.

The small duality gap is a consequence of the structure of the set of achievable constraint-cost pairs

$$S = \{ (g(x), f(x)) \mid x \in X \},\$$

where

$$g(x) = \sum_{i=1}^{m} g_i(x_i), \qquad f(x) = \sum_{i=1}^{m} f_i(x_i).$$

Note that this set plays a central role in the proof of the Nonlinear Farkas' Lemma (see Fig. 5.1.2). In the case of a separable problem, it can be written as a vector sum of m sets, one for each separable term, i.e.,

$$S = S_1 + \dots + S_m,$$

where

$$S_i = \{ (g_i(x_i), f_i(x_i)) \mid x_i \in X_i \}.$$

A key fact is that generally, a set that is the vector sum of a large number of possibly nonconvex but roughly similar sets "tends to be convex" in the sense that any vector in its convex hull can be closely approximated by a vector in the set. As a result, the duality gap tends to be relatively small. The analytical substantiation is based on the following theorem, which roughly states that at most r+1 out of the m convex sets "contribute to the nonconvexity" of their vector sum, regardless of the value of m.

Proposition 5.7.1: (Shapley-Folkman Theorem) Let S_i , i = 1, ..., m, be nonempty subsets of \Re^{r+1} , with m > r+1, and let $S = S_1 + \cdots + S_m$. Then every vector $s \in \text{conv}(S)$ can be represented as $s = s_1 + \cdots + s_m$, where $s_i \in \text{conv}(S_i)$ for all i = 1, ..., m, and $s_i \notin S_i$ for at most r+1 indices i.

Proof: We clearly have $\operatorname{conv}(S) = \operatorname{conv}(S_1) + \cdots + \operatorname{conv}(S_m)$ (since convex combinations commute with linear transformations). Thus any $s \in \operatorname{conv}(S)$ can be written as $s = \sum_{i=1}^m y_i$ with $y_i \in \operatorname{conv}(S_i)$, so that $y_i = \sum_{j=1}^{t_i} a_{ij} y_{ij}$ for some $a_{ij} > 0$, $\sum_{j=1}^{t_i} a_{ij} = 1$, and $y_{ij} \in S_i$. Consider the following vectors of \Re^{r+1+m} :

$$z = \begin{pmatrix} s \\ 1 \\ 1 \\ \vdots \\ 1 \end{pmatrix}, \qquad z_{1j} = \begin{pmatrix} y_{1j} \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, \qquad \cdots \qquad z_{mj} = \begin{pmatrix} y_{mj} \\ 0 \\ 0 \\ \vdots \\ 1 \end{pmatrix},$$

so that $z = \sum_{i=1}^{m} \sum_{j=1}^{t_i} a_{ij} z_{ij}$. We now view z as a vector in the cone generated by the vectors z_{ij} , and use Caratheodory's Theorem [Prop. 1.2.1(a)] to write $z = \sum_{i=1}^{m} \sum_{j=1}^{t_i} b_{ij} z_{ij}$ for some nonnegative scalars b_{ij} , at most r+1+m of which are strictly positive. Focusing on the first component of z, we have

$$s = \sum_{i=1}^{m} \sum_{j=1}^{t_i} b_{ij} y_{ij}, \qquad \sum_{j=1}^{t_i} b_{ij} = 1, \quad \forall \ i = 1, \dots, m.$$

Let $s_i = \sum_{j=1}^{t_i} b_{ij} y_{ij}$, so that $s = s_1 + \dots + s_m$ with $s_i \in \text{conv}(S_i)$ for all i. For each $i = 1, \dots, m$, at least one of b_{i1}, \dots, b_{it_i} must be positive, so there are at most only r+1 additional coefficients b_{ij} that can be positive (since at most r+1+m of the b_{ij} are positive). It follows that for at least m-r-1 indices i, we have $b_{ik}=1$ for some k and $b_{ij}=0$ for all $j \neq k$. For these indices, we have $s_i \in S_i$. Q.E.D.

Let us use the Shapley-Folkman Theorem to estimate the duality gap in the case of a linear program with integer 0-1 constraints:

minimize
$$\sum_{i=1}^{m} c_i x_i$$
subject to
$$\sum_{i=1}^{m} a_{ji} x_i \le b_j, \quad j = 1, \dots, r,$$

$$x_i = 0 \text{ or } 1, \quad i = 1, \dots, m.$$

$$(5.48)$$

Let f^* and q^* denote the optimal primal and dual values, respectively. Note that the "relaxed" linear program, where the integer constraints are replaced by $x_i \in [0,1], i=1,\ldots,m$, has the same dual, so its optimal value is q^* . The set S can be written as

$$S = S_1 + \cdots + S_m - (b, 0),$$

where $b = (b_1, ..., b_r)$ and each S_i consists of just two elements corresponding to $x_i = 0, 1$, i.e.,

$$S_i = \{(0, \dots, 0, 0), (a_{1i}, \dots, a_{ri}, c_i)\}.$$

Thus, S consists of a total of 2^m points. A natural idea is to solve the "relaxed" program, and then try to suitably "round" the fractional (i.e., non-integer) components of the relaxed optimal solution to an integer, thereby obtaining a suboptimal solution of the original integer program (5.48).

Let us denote

$$\gamma = \max_{i=1,\dots,m} |c_i|, \qquad \delta = \max_{i=1,\dots,m} \delta_i,$$

where

$$\delta_{i} = \begin{cases} 0 & \text{if } a_{1i}, \dots, a_{ri} \ge 0, \\ 0 & \text{if } a_{1i}, \dots, a_{ri} \le 0, \\ \max_{j=1,\dots,r} |a_{ji}| & \text{otherwise.} \end{cases}$$
 (5.49)

Note that δ_i is an upper bound to the maximum amount of constraint violation that can result when a fractional variable $x_i \in (0,1)$ is rounded suitably (up or down). The following proposition shows what can be achieved with such a rounding procedure.†

Proposition 5.7.2: Assume that the relaxed version of the linear/integer problem (5.48) is feasible. Then there exists $\overline{x} = (\overline{x}_1, \dots, \overline{x}_m)$, with $\overline{x}_i \in \{0,1\}$ for all i, which violates each of the inequality constraints of problem (5.48) by at most $(r+1)\delta$, and has cost that is at most $q^* + (r+1)\gamma$.

Proof: We note that the relaxed problem, being a feasible linear program with compact constraint set, has an optimal solution with optimal value q^* . Since

$$\operatorname{conv}(S) = \operatorname{conv}(S_1) + \dots + \operatorname{conv}(S_m) - (b, 0),$$

we see that $\operatorname{conv}(S)$ is the set of constraint-cost pairs of the relaxed problem, so it contains a vector of the form (u^*, q^*) with $u^* \leq 0$. By the Shapley-Folkman Theorem, there is an index subset I with at most r+1elements such that (u^*, q^*) can be written as

$$u^* = \sum_{i \in I} \overline{u}_i + \sum_{i \notin I} u_i, \qquad q^* = \sum_{i \in I} \overline{w}_i + \sum_{i \notin I} w_i,$$

where $(\overline{u}_i, \overline{w}_i) \in \text{conv}(S_i)$ for $i \in I$ and $(u_i, w_i) \in S_i$ for $i \notin I$. Each pair (u_i, w_i) , $i \notin I$ corresponds to an integer component $\overline{x}_i \in \{0, 1\}$. Each pair $(\overline{u}_i, \overline{w}_i)$, $i \in I$ may be replaced/rounded by one of the two elements (u_i, w_i) of S_i , yielding again an integer component $\overline{x}_i \in \{0, 1\}$, with an increase in cost of at most γ and an increase of the level of each inequality constraint of at most δ_i . We thus obtain an integer vector \overline{x} that violates each inequality constraint by at most $(r+1)\delta$, and has cost that is at most $q^* + (r+1)\gamma$. Q.E.D.

The preceding proof also indicates the rounding mechanism for obtaining the vector \overline{x} in the proposition. In practice, the simplex method

[†] A slightly stronger bound $[r\gamma]$ and $r\delta$ in place of $(r+1)\gamma$ and $(r+1)\delta$, respectively] can be shown with an alternative proof that uses the theory of the simplex method. However, the proof given here generalizes to the case where f_i and g_i are nonlinear (see the subsequent Prop. 5.7.4).

provides a relaxed problem solution with no more than r noninteger components, which are then rounded as indicated in the proof to obtain \overline{x} . Note that \overline{x} may not be feasible, and indeed it is possible that the relaxed problem is feasible, while the original integer problem is not. For example the constraints $x_1 - x_2 \le -1/2$, $x_1 + x_2 \le 1/2$ are feasible for the relaxed problem but not for the original integer problem.

On the other hand if for each j, each of the constraint coefficients a_{j1}, \ldots, a_{jm} is either 0 or has the same sign, we have $\delta = 0$ and a feasible solution that is $(r+1)\gamma$ -optimal can be found. Assuming this condition, let us consider now a sequence of similar problems where the number of inequality constraints r is kept constant, but the dimension m grows to infinity. Assuming that

$$\beta_1 m \le |f^*| \le \beta_2 m$$

for some $\beta_1, \beta_2 > 0$, we see that the rounding error $(r+1)\gamma$ (which bounds the duality gap $f^* - q^*$) diminishes in relation to f^* , i.e., its ratio to f^* tends to 0 as $m \to \infty$. In particular, we have

$$\lim_{m\to\infty}\frac{f^*-q^*}{f^*}\to 0.$$

We will now generalize the line of analysis of the preceding proposition to the nonlinear separable problem (5.47) and obtain similar results. In particular, under assumptions that parallel the condition $\delta = 0$ discussed earlier, we will prove that the duality gap satisfies

$$f^* - q^* \le (r+1) \max_{i=1,\dots,m} \gamma_i,$$

where for each i, γ_i is a nonnegative scalar that depends on the functions f_i , g_i , and the set X_i . This suggests that as $m \to \infty$, the duality gap diminishes relative to f^* as $m \to \infty$. We first derive a general estimate for the case where the set M in the MC/MC framework is a vector sum, and then we specialize to the case of the separable problem.

Proposition 5.7.3: Consider the MC/MC framework corresponding to a set $M \subset \Re^{n+1}$ given by

$$M = M_1 + \cdots + M_m$$

where the sets M_i have the form

$$M_i = \tilde{M}_i + \{(u,0) \mid u \in C_i\}, \qquad i = 1, \dots, m,$$

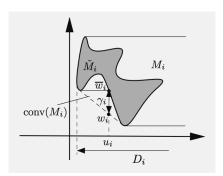


Figure 5.7.1. Geometric interpretation of the bound coefficient γ_i in Prop. 5.7.3 for the case where

$$M_i = \tilde{M}_i + \left\{ (u, 0) \mid u \in C_i \right\}$$

and

$$C_i = \{u \mid u \ge 0\}.$$

and for each i, \tilde{M}_i is compact and C_i is convex, and such that $C_1 + \cdots + C_m$ is closed. For $i = 1, \ldots, m$, let

$$D_i = \{u_i \mid \text{there exists } w_i \text{ such that } (u_i, w_i) \in M_i\},\$$

and assume that for each vector $(u_i, w_i) \in \text{conv}(M_i)$, we have $u_i \in D_i$. Let

$$\gamma_i = \sup_{u_i \in D_i} \inf \{ \overline{w}_i - w_i \mid (u_i, \overline{w}_i) \in M_i, (u_i, w_i) \in \text{conv}(M_i) \}.$$

Then assuming that w^* and q^* are finite, we have

$$w^* - q^* \le (n+1) \max_{i=1,...,m} \gamma_i.$$

Proof: Note that the meaning of γ_i is that we can approximate any vector $(u_i, w_i) \in \text{conv}(M_i)$ with a vector $(u_i, \overline{w}_i) \in M_i$, while incurring an increase $\overline{w}_i - w_i$ of the vertical component w_i that is at most $\gamma_i + \epsilon$, where ϵ is arbitrarily small (cf. Fig. 5.7.1).

The idea of the proof is to use the Shapley-Folkman Theorem to approximate a vector $(u_1 + \cdots + u_m, w_1 + \cdots + w_m) \in \text{conv}(M)$ with a vector $(u_1 + \cdots + u_m, \overline{w}_1 + \cdots + \overline{w}_m) \in M$ such that $\overline{w}_i = w_i$ for all except at most n+1 indexes i for which $\overline{w}_i - w_i \leq \gamma_i + \epsilon$.

By the result of Prop. 4.3.2(c), q^* is finite and is equal to the optimal value of the MC/MC problem corresponding to

$$\operatorname{conv}(M) = \operatorname{conv}(\tilde{M}_1) + \dots + \operatorname{conv}(\tilde{M}_m) + \{(u,0) \mid u \in C_1 + \dots + C_m\}.$$

Since \tilde{M}_i is compact, so is $\operatorname{conv}(\tilde{M}_i)$ (cf. Prop. 1.2.2). Since $C_1 + \cdots + C_m$ is closed, $\operatorname{conv}(M)$ is closed (cf. Prop. 1.4.14), and it follows that the vector $(0, q^*)$ belongs to $\operatorname{conv}(M)$. By the Shapley-Folkman Theorem (Prop.

5.7.1), there is an index subset I with at most n+1 elements such that

$$0 = \sum_{i=1}^{m} u_i, \qquad q^* = \sum_{i \in I} w_i + \sum_{i \notin I} \overline{w}_i,$$

where $(u_i, w_i) \in \text{conv}(M_i)$ for $i \in I$ and $(u_i, \overline{w_i}) \in M_i$ for $i \notin I$. By assumption, for any $\epsilon > 0$ and $i \in I$, there exist vectors of the form $(u_i, \overline{w_i}) \in M_i$ with $\overline{w_i} \leq w_i + \gamma_i + \epsilon$. Hence, $(0, \overline{w}) \in M$, where

$$\overline{w} = \sum_{i \in I} \overline{w}_i + \sum_{i \notin I} \overline{w}_i.$$

Thus, $w^* \leq \overline{w} = \sum_{i \in I} \overline{w}_i + \sum_{i \notin I} \overline{w}_i$, and it follows that

$$w^* \le \sum_{i \in I} \overline{w}_i + \sum_{i \notin I} \overline{w}_i \le \sum_{i \in I} (w_i + \gamma_i + \epsilon) + \sum_{i \notin I} \overline{w}_i \le q^* + (n+1) \max_{i=1,\dots,m} (\gamma_i + \epsilon).$$

By taking $\epsilon \downarrow 0$, the result follows. **Q.E.D.**

We now apply the preceding proposition to the general separable problem of this section.

Proposition 5.7.4: Consider the separable problem (5.47). Assume that the sets

$$\tilde{M}_i = \{ (g_i(x_i), f_i(x_i)) \mid x_i \in X_i \}, \quad i = 1, \dots, m,$$

are nonempty and compact, and that for each i, given any $x_i \in \text{conv}(X_i)$, there exists $\tilde{x}_i \in X_i$ such that

$$g_i(\tilde{x}_i) \le (\operatorname{cl} g_i)(x_i), \tag{5.50}$$

where $\operatorname{cl} g_i$ is the function whose components are the convex closures of the corresponding components of g_i . Then

$$f^* - q^* \le (r+1) \max_{i=1,...,m} \gamma_i,$$

where

$$\gamma_i = \sup \{ \tilde{f}_i(x_i) - (\operatorname{cl} f_i)(x_i) \mid x_i \in \operatorname{conv}(X_i) \},$$
 (5.51)

čl f_i is the convex closure of f_i , and \tilde{f}_i is given by

$$\tilde{f}_i(x_i) = \inf\{f_i(\tilde{x}_i) \mid g_i(\tilde{x}_i) \le (\tilde{\operatorname{cl}} g_i)(x_i), \, \tilde{x}_i \in X_i\}, \quad \forall \, x_i \in \operatorname{conv}(X_i).$$

Note that Eq. (5.50) is a vector (rather than scalar) inequality, so it is not automatically satisfied. In the context of the integer programming problem (5.48), this inequality is satisfied if for each i, all the coefficients a_{1i}, \ldots, a_{ri} have the same sign [cf. Eq. (5.49)]. It is also satisfied if X_i is convex and the components of g_i are convex over X_i . We will prove Prop. 5.7.4 by first showing the following lemma, which can be visualized with the aid of Fig. 5.7.1:

Lemma 5.7.1: Let $h: \Re^n \mapsto \Re^r$ and $\ell: \Re^n \mapsto \Re$ be given functions, and let X be a nonempty set. Denote

$$\tilde{M} = \{ (h(x), \ell(x)) \mid x \in X \}, \qquad M = \tilde{M} + \{ (u, 0) \mid u \ge 0 \}.$$

Then every $(u, w) \in \text{conv}(M)$ satisfies

$$(\check{\operatorname{cl}} h)(x) \le u, \qquad (\check{\operatorname{cl}} \ell)(x) \le w,$$

for some $x \in \text{conv}(X)$, where $\text{cl}\,\ell$ is the convex closure of ℓ , and $\text{cl}\,h$ is the function whose components are the convex closures of the corresponding components of h.

Proof: Any $(u, w) \in \text{conv}(M)$ can be written as

$$(u, w) = \sum_{j=1}^{s} \alpha_j \cdot (u_j, w_j),$$

where s is some positive integer and

$$\sum_{j=1}^{s} \alpha_{j} = 1, \quad \alpha_{j} \ge 0, \quad (u_{j}, w_{j}) \in M, \quad j = 1, \dots, r.$$

By the definition of M, there exists $x_j \in X$ such that $h(x_j) \leq u_j, \ell(x_j) = w_j$. Let $x = \sum_{j=1}^s \alpha_j x_j$, so that $x \in \text{conv}(X)$. Then

$$(\check{\operatorname{cl}} h)(x) \le \sum_{j=1}^{s} \alpha_j(\check{\operatorname{cl}} h)(x_j) \le \sum_{j=1}^{s} \alpha_j h(x_j) \le u,$$

$$(\check{\operatorname{cl}} \ell)(x) \le \sum_{j=1}^{s} \alpha_j(\check{\operatorname{cl}} \ell)(x_j) \le \sum_{j=1}^{s} \alpha_j \ell(x_j) = w,$$

where we used the definition of convex closure and Prop. 1.3.14. Q.E.D.

Proof of Prop. 5.7.4: Denote

$$M_i = \tilde{M}_i + \{(u,0) \mid u \ge 0\},\$$

and consider the MC/MC framework corresponding to

$$M = M_1 + \dots + M_m.$$

By applying Prop. 4.3.2(c) with $C = \{u \mid u \ge 0\}$, we have

$$w^* = \inf_{(0,w) \in M} w, \qquad q^* = \inf_{(0,w) \in \text{conv}(M)} w.$$

We will now apply the result of Prop. 5.7.3 to derive the desired bound.

To this end, we use Lemma 5.7.1 to assert that for each i = 1, ..., m and vector $(u_i, w_i) \in \text{conv}(M_i)$, there exists $x_i \in \text{conv}(X_i)$ satisfying

$$(\operatorname{cl} g_i)(x_i) \le u_i, \qquad (\operatorname{cl} f_i)(x_i) \le w_i.$$

Since from the definition (5.51) of γ_i , we have $\tilde{f}_i(x_i) \leq (\operatorname{cl} f_i)(x_i) + \gamma_i$, we obtain

$$\tilde{f}_i(x_i) \le w_i + \gamma_i.$$

It follows from the definition of \tilde{f}_i that there exists $\tilde{x}_i \in X_i$ such that

$$g_i(\tilde{x}_i) \le (\operatorname{cl} g_i)(x_i) \le u_i, \qquad f_i(\tilde{x}_i) \le w_i + \gamma_i.$$

Hence $(u_i, f_i(\tilde{x}_i)) \in M_i$, and by letting $\overline{w}_i = f_i(\tilde{x}_i)$, we have that for each vector $(u_i, w_i) \in \text{conv}(M_i)$, there exists a vector of the form $(u_i, \overline{w}_i) \in M_i$ with $\overline{w}_i - w_i \leq \gamma_i$. By applying Prop. 5.7.3, we obtain

$$w^* - q^* \le (r+1) \max_{i=1,...,m} \gamma_i.$$

Q.E.D.

As an example, consider the special case of the integer programming problem (5.48), where f_i and g_i are linear, and $X_i = \{0, 1\}$. Then the assumption (5.50) is equivalent to $\delta_i = 0$ [cf. Eq. (5.49)], and Props. 5.7.2 and 5.7.4 give the same estimate. For further discussion and visualization of the duality gap estimate of Prop. 5.7.4, we refer to [Ber82], Section 5.6.1. The estimate suggests that many nonconvex/integer separable problems with a fixed number of constraints become easier to solve as their dimension increases, because the corresponding duality gap diminishes.

5.7.2 Duality Gap in Minimax Problems

Consider a minimax problem involving a function $\phi: X \times Z \mapsto \Re$ defined over nonempty sets X and Z.† We saw in Sections 4.2.5 and 5.5 that minimax theory can be developed within the context of the MC/MC framework involving the set

$$M = \operatorname{epi}(p),$$

where $p: \Re^m \mapsto [-\infty, \infty]$ is the function given by

$$p(u) = \inf_{x \in X} \sup_{z \in Z} \big\{ \phi(x, z) - u'z \big\}, \qquad u \in \Re^m.$$

Let us assume that $(-\hat{\operatorname{cl}}\phi)(x,\cdot)$ is proper for all $x\in X$, so that the corresponding dual function is

$$q(\mu) = \inf_{x \in X} (\hat{\operatorname{cl}} \phi)(x, \mu)$$

(cf. Prop. 4.2.1). Then we have

$$\sup_{z \in Z} \inf_{x \in X} \phi(x,z) \leq \sup_{z \in Z} \inf_{x \in X} (\widehat{\operatorname{cl}} \phi)(x,z) = \sup_{\mu \in \Re^m} q(\mu) = q^*.$$

Furthermore, $q^* = p^{\star\star}(0)$ [cf. Eq. (5.46)], so that

$$\sup_{z \in Z} \inf_{x \in X} \phi(x,z) \leq q^* = p^{\star\star}(0) \leq p(0) = w^* = \inf_{x \in X} \sup_{z \in Z} \phi(x,z).$$

Thus the gap between "infsup" and "supinf" can be decomposed into the sum of two terms:

$$\inf_{x \in X} \sup_{z \in Z} \phi(x,z) - \sup_{z \in Z} \inf_{x \in X} \phi(x,z) = \overline{G} + \underline{G},$$

where

$$\overline{G} = \inf_{x \in X} \sup_{z \in Z} \phi(x, z) - p^{\star\star}(0) = p(0) - p^{\star\star}(0),$$

$$\underline{G} = q^* - \sup_{z \in Z} \inf_{x \in X} \phi(x,z) = \sup_{z \in \Re^m} \inf_{x \in X} (\widehat{\operatorname{cl}} \, \phi)(x,z) - \sup_{z \in Z} \inf_{x \in X} \phi(x,z).$$

The term \overline{G} is the duality gap $w^* - q^*$ of the MC/MC framework. It can be attributed to the lack of convexity and/or closure of p, and in view of the definition of p, it can also be attributed to lack of convexity and/or closure of ϕ with respect to x. The term \underline{G} can be attributed to the lack of concavity and/or upper semicontinuity of ϕ with respect to z.

In cases where ϕ has a separable form in x, and ϕ is concave and upper semicontinuous with respect to z, we have $\underline{G}=0$, while the Shapley-Folkman Theorem can be used to estimate \overline{G} , similar to Prop. 5.7.4. Similarly, we can estimate the duality gap if ϕ is separable in z, and ϕ is convex and lower semicontinuous with respect to x, or if ϕ is separable in both x and z.

[†] The ideas of this subsection are due to Mengdi Wang.

APPENDIX A:

Mathematical Background

In this appendix, we list some basic definitions, notational conventions, and results from linear algebra and real analysis. We assume that the reader is familiar with these subjects, so no proofs are given. For additional related material, we refer to textbooks such as Hoffman and Kunze [HoK71], Lancaster and Tismenetsky [LaT85], and Strang [Str76] (linear algebra), and Ash [Ash72], Ortega and Rheinboldt [OrR70], and Rudin [Rud76] (real analysis).

Set Notation

If X is a set and x is an element of X, we write $x \in X$. A set can be specified in the form $X = \{x \mid x \text{ satisfies } P\}$, as the set of all elements satisfying property P. The union of two sets X_1 and X_2 is denoted by $X_1 \cup X_2$, and their intersection by $X_1 \cap X_2$. The symbols \exists and \forall have the meanings "there exists" and "for all," respectively. The empty set is denoted by \emptyset .

The set of real numbers (also referred to as scalars) is denoted by \Re . The set \Re augmented with $+\infty$ and $-\infty$ is called the set of extended real numbers. We write $-\infty < x < \infty$ for all real numbers x, and $-\infty \le x \le \infty$ for all extended real numbers x. We denote by [a,b] the set of (possibly extended) real numbers x satisfying $a \le x \le b$. A rounded, instead of square, bracket denotes strict inequality in the definition. Thus (a,b], [a,b), and (a,b) denote the set of all x satisfying $a < x \le b$, $a \le x < b$, and a < x < b, respectively. Furthermore, we use the natural extensions of the rules of arithmetic: $x \cdot 0 = 0$ for every extended real number x, $x \cdot \infty = \infty$ if x > 0, $x \cdot \infty = -\infty$ if x < 0, and $x + \infty = \infty$ and $x - \infty = -\infty$ for every scalar x. The expression $\infty - \infty$ is meaningless and is never allowed to occur.

Inf and Sup Notation

The supremum of a nonempty set X of scalars, denoted by $\sup X$, is defined as the smallest scalar y such that $y \geq x$ for all $x \in X$. If no such scalar exists, we say that the supremum of X is ∞ . Similarly, the infimum of X, denoted by $\inf X$, is defined as the largest scalar y such that $y \leq x$ for all $x \in X$, and is equal to $-\infty$ if no such scalar exists. For the empty set, we use the convention

$$\sup \emptyset = -\infty, \quad \inf \emptyset = \infty.$$

If $\sup X$ is equal to a scalar \overline{x} that belongs to the set X, we say that \overline{x} is the maximum point of X and we write $\overline{x} = \max X$. Similarly, if $\inf X$ is equal to a scalar \overline{x} that belongs to the set X, we say that \overline{x} is the minimum point of X and we write $\overline{x} = \min X$. Thus, when we write $\max X$ (or $\min X$) in place of $\sup X$ (or $\inf X$, respectively), we do so just for emphasis: we indicate that it is either evident, or it is known through earlier analysis, or it is about to be shown that the maximum (or minimum, respectively) of the set X is attained at one of its points.

Vector Notation

We denote by \Re^n the set of *n*-dimensional real vectors. For any $x \in \Re^n$, we use x_i to indicate its *i*th *coordinate*, also called its *i*th *component*. Vectors in \Re^n will be viewed as column vectors, unless the contrary is explicitly stated. For any $x \in \Re^n$, x' denotes the transpose of x, which is an *n*-dimensional row vector. The *inner product* of two vectors $x, y \in \Re^n$ is defined by $x'y = \sum_{i=1}^n x_i y_i$. Two vectors $x, y \in \Re^n$ satisfying x'y = 0 are called *orthogonal*.

If x is a vector in \Re^n , the notations x>0 and $x\geq 0$ indicate that all components of x are positive and nonnegative, respectively. For any two vectors x and y, the notation x>y means that x-y>0. The notations $x\geq y, x< y$, etc., are to be interpreted accordingly.

Function Notation

If f is a function, we use the notation $f: X \mapsto Y$ to indicate the fact that f is defined on a nonempty set X (its domain) and takes values in a set Y (its range). Thus when using the notation $f: X \mapsto Y$, we implicitly assume that X is nonempty. If $f: X \mapsto Y$ is a function, and U and V are subsets of X and Y, respectively, the set $\{f(x) \mid x \in U\}$ is called the image or $forward\ image\ of\ U\ under\ f$, and the set $\{x \in X \mid f(x) \in V\}$ is called the $inverse\ image\ of\ V\ under\ f$.

A.1 LINEAR ALGEBRA

If X is a set and λ is a scalar, we denote by λX the set $\{\lambda x \mid x \in X\}$. If X_1 and X_2 are two subsets of \Re^n , we denote by $X_1 + X_2$ the set

$${x_1 + x_2 \mid x_1 \in X_1, x_2 \in X_2},$$

which is referred to as the vector sum of X_1 and X_2 . We use a similar notation for the sum of any finite number of subsets. In the case where one of the subsets consists of a single vector \overline{x} , we simplify this notation as follows:

$$\overline{x} + X = \{ \overline{x} + x \mid x \in X \}.$$

We also denote by $X_1 - X_2$ the set

$${x_1 - x_2 \mid x_1 \in X_1, x_2 \in X_2}.$$

Given sets $X_i \subset \Re^{n_i}$, i = 1, ..., m, the Cartesian product of the X_i , denoted by $X_1 \times \cdots \times X_m$, is the set

$$\{(x_1,\ldots,x_m) \mid x_i \in X_i, i = 1,\ldots,m\},\$$

which is viewed as a subset of $\Re^{n_1+\cdots+n_m}$.

Subspaces and Linear Independence

A nonempty subset S of \Re^n is called a subspace if $ax + by \in S$ for every $x,y \in S$ and every $a,b \in \Re$. An affine set in \Re^n is a translated subspace, i.e., a set X of the form $X = \overline{x} + S = \{\overline{x} + x \mid x \in S\}$, where \overline{x} is a vector in \Re^n and S is a subspace of \Re^n , called the subspace parallel to X. Note that there can be only one subspace S associated with an affine set in this manner. [To see this, let X = x + S and $X = \overline{x} + \overline{S}$ be two representations of the affine set X. Then, we must have $x = \overline{x} + \overline{s}$ for some $\overline{s} \in \overline{S}$ (since $x \in X$), so that $X = \overline{x} + \overline{s} + S$. Since we also have $X = \overline{x} + \overline{S}$, it follows that $S = \overline{S} - \overline{s} = \overline{S}$.] A nonempty set X is a subspace if and only if it contains the origin, and every line that passes through any pair of its points that are distinct, i.e., it contains 0 and all points $\alpha x + (1 - \alpha)y$, where $\alpha \in \Re$ and $x,y \in X$ with $x \neq y$. Similarly X is affine if and only if it contains every line that passes through any pair of its points that are distinct. The span of a finite collection $\{x_1, \ldots, x_m\}$ of elements of \Re^n , denoted by $span(x_1, \ldots, x_m)$, is the subspace consisting of all vectors y of the form $y = \sum_{k=1}^m \alpha_k x_k$, where each α_k is a scalar.

The vectors $x_1, \ldots, x_m \in \mathbb{R}^n$ are called *linearly independent* if there exists no set of scalars $\alpha_1, \ldots, \alpha_m$, at least one of which is nonzero, such that $\sum_{k=1}^m \alpha_k x_k = 0$. An equivalent definition is that $x_1 \neq 0$, and for every k > 1, the vector x_k does not belong to the span of x_1, \ldots, x_{k-1} .

If S is a subspace of \Re^n containing at least one nonzero vector, a basis for S is a collection of vectors that are linearly independent and whose span is equal to S. Every basis of a given subspace has the same number of vectors. This number is called the dimension of S. By convention, the subspace $\{0\}$ is said to have dimension zero. Every subspace of nonzero dimension has a basis that is orthogonal (i.e., any pair of distinct vectors from the basis is orthogonal). The dimension of an affine set $\overline{x} + S$ is the dimension of the corresponding subspace S. An (n-1)-dimensional affine set is called a hyperplane. It is a set specified by a single linear equation, i.e., a set of the form $\{x \mid a'x = b\}$, where $a \neq 0$ and $b \in \Re$.

Given any set X, the set of vectors that are orthogonal to all elements of X is a subspace denoted by X^{\perp} :

$$X^{\perp} = \{ y \mid y'x = 0, \, \forall \, x \in X \}.$$

If S is a subspace, S^{\perp} is called the *orthogonal complement* of S. Any vector x can be uniquely decomposed as the sum of a vector from S and a vector from S^{\perp} . Furthermore, we have $(S^{\perp})^{\perp} = S$.

Matrices

For any matrix A, we use A_{ij} , $[A]_{ij}$, or a_{ij} to denote its ijth component. The transpose of A, denoted by A', is defined by $[A']_{ij} = a_{ji}$. For any two matrices A and B of compatible dimensions, the transpose of the product matrix AB satisfies (AB)' = B'A'. The inverse of a square and invertible A is denoted A^{-1} .

If X is a subset of \Re^n and A is an $m \times n$ matrix, then the *image of* X under A is denoted by AX (or $A \cdot X$ if this enhances notational clarity):

$$AX = \{Ax \mid x \in X\}.$$

If Y is a subset of \Re^m , the inverse image of Y under A is denoted by $A^{-1}Y$:

$$A^{-1}Y = \{x \mid Ax \in Y\}.$$

Let A be a square matrix. We say that A is symmetric if A' = A. A symmetric matrix has real eigenvalues and an orthogonal set of corresponding real eigenvectors. We say that A is diagonal if $[A]_{ij} = 0$ whenever $i \neq j$. We use I to denote the identity matrix (this is the diagonal matrix whose diagonal components are equal to 1).

A symmetric $n \times n$ matrix A is called positive definite if x'Ax > 0 for all $x \in \Re^n$, $x \neq 0$. It is called positive semidefinite if $x'Ax \geq 0$ for all $x \in \Re^n$. Throughout this book, the notion of positive definiteness applies exclusively to symmetric matrices. Thus whenever we say that a matrix is positive (semi)definite, we implicitly assume that the matrix is

symmetric, although we usually add the term "symmetric" for clarity. A positive semidefinite matrix A can be written as A = M'M for some matrix M. A symmetric matrix is positive definite (or semidefinite) if and only if its eigenvalues are positive (nonnegative, respectively).

Let A be an $m \times n$ matrix. The range space of A, denoted by R(A), is the set of all vectors $y \in \Re^m$ such that y = Ax for some $x \in \Re^n$. The nullspace of A, denoted by N(A), is the set of all vectors $x \in \Re^n$ such that Ax = 0. It is seen that the range space and the null space of A are subspaces. The rank of A is the dimension of the range space of A. The rank of A is equal to the maximal number of linearly independent columns of A, and is also equal to the maximal number of linearly independent rows of A. The matrix A and its transpose A' have the same rank. We say that A has full rank, if its rank is equal to $\min\{m,n\}$. This is true if and only if either all the rows of A are linearly independent, or all the columns of A are linearly independent. A symmetric matrix is positive (semi)definite if and only if its eignevalues are positive (nonnegative, respectively)

The range space of an $m \times n$ matrix A is equal to the orthogonal complement of the nullspace of its transpose, i.e., $R(A) = N(A')^{\perp}$. Another way to state this result is that given vectors $a_1, \ldots, a_n \in \Re^m$ (the columns of A) and a vector $x \in \Re^m$, we have x'y = 0 for all y such that $a'_i y = 0$ for all i if and only if $x = \lambda_1 a_1 + \cdots + \lambda_n a_n$ for some scalars $\lambda_1, \ldots, \lambda_n$. This is a special case of Farkas' Lemma, an important result for constrained optimization, which will be discussed in Section 2.3.1.

A function $f: \Re^n \to \Re$ is said to be affine if it has the form f(x) = a'x + b for some $a \in \Re^n$ and $b \in \Re$. Similarly, a function $f: \Re^n \to \Re^m$ is said to be affine if it has the form f(x) = Ax + b for some $m \times n$ matrix A and some $b \in \Re^m$. If b = 0, f is said to be a linear function or linear transformation. Sometimes, with slight abuse of terminology, an equation or inequality involving a linear function, such as a'x = b or $a'x \leq b$, is referred to as a linear equation or inequality, respectively.

A.2 TOPOLOGICAL PROPERTIES

Definition A.2.1: A norm $\|\cdot\|$ on \Re^n is a function that assigns a scalar $\|x\|$ to every $x \in \Re^n$ and that has the following properties:

- (a) $||x|| \ge 0$ for all $x \in \Re^n$.
- (b) $\|\alpha x\| = |\alpha| \cdot \|x\|$ for every scalar α and every $x \in \Re^n$.
- (c) ||x|| = 0 if and only if x = 0.
- (d) $||x+y|| \le ||x|| + ||y||$ for all $x, y \in \Re^n$ (this is referred to as the triangle inequality).

The Euclidean norm of a vector $x = (x_1, \ldots, x_n)$ is defined by

$$||x|| = (x'x)^{1/2} = \left(\sum_{i=1}^{n} |x_i|^2\right)^{1/2}.$$

We will use the Euclidean norm almost exclusively in this book. In particular, in the absence of a clear indication to the contrary, $\|\cdot\|$ will denote the Euclidean norm. The Schwarz inequality states that for any two vectors x and y, we have

$$|x'y| \le ||x|| \cdot ||y||,$$

with equality holding if and only if $x = \alpha y$ for some scalar α . The Pythagorean Theorem states that for any two vectors x and y that are orthogonal, we have

$$||x + y||^2 = ||x||^2 + ||y||^2.$$

Sequences

A sequence $\{x_k \mid k=1,2,\ldots\}$ (or $\{x_k\}$ for short) of scalars is said to converge if there exists a scalar x such that for every $\epsilon>0$ we have $|x_k-x|<\epsilon$ for every k greater than some integer K (that depends on ϵ). The scalar x is said to be the limit of $\{x_k\}$, and the sequence $\{x_k\}$ is said to converge to x; symbolically, $x_k \to x$ or $\lim_{k\to\infty} x_k = x$. If for every scalar k there exists some k (that depends on k) such that k0 for all k1 for every scalar k2 for all k2 for all k3 some integer k5 such that k6 for all k7 for every scalar k8 scalar for all k8 converges" or "the limit of k8 exists" or "k8 has a limit" is that the limit of k8 is a scalar.

A scalar sequence $\{x_k\}$ is said to be bounded above (respectively, below) if there exists some scalar b such that $x_k \leq b$ (respectively, $x_k \geq b$) for all k. It is said to be bounded if it is bounded above and bounded below. The sequence $\{x_k\}$ is said to be monotonically nonincreasing (respectively, nondecreasing) if $x_{k+1} \leq x_k$ (respectively, $x_{k+1} \geq x_k$) for all k. If $x_k \to x$ and $\{x_k\}$ is monotonically nonincreasing (nondecreasing), we also use the notation $x_k \downarrow x$ ($x_k \uparrow x$, respectively).

Proposition A.2.1: Every bounded and monotonically nonincreasing or nondecreasing scalar sequence converges.

Note that a monotonically nondecreasing sequence $\{x_k\}$ is either bounded, in which case it converges to some scalar x by the above proposition, or else it is unbounded, in which case $x_k \to \infty$. Similarly, a mono-

tonically nonincreasing sequence $\{x_k\}$ is either bounded and converges, or it is unbounded, in which case $x_k \to -\infty$.

Given a scalar sequence $\{x_k\}$, let

$$y_m = \sup\{x_k \mid k \ge m\}, \qquad z_m = \inf\{x_k \mid k \ge m\}.$$

The sequences $\{y_m\}$ and $\{z_m\}$ are nonincreasing and nondecreasing, respectively, and therefore have a limit whenever $\{x_k\}$ is bounded above or is bounded below, respectively (Prop. A.2.1). The limit of y_m is denoted by $\limsup_{k\to\infty} x_k$, and is referred to as the *upper limit* of $\{x_k\}$. The limit of z_m is denoted by $\liminf_{k\to\infty} x_k$, and is referred to as the *lower limit* of $\{x_k\}$. If $\{x_k\}$ is unbounded above, we write $\limsup_{k\to\infty} x_k = \infty$, and if it is unbounded below, we write $\liminf_{k\to\infty} x_k = -\infty$.

Proposition A.2.2: Let $\{x_k\}$ and $\{y_k\}$ be scalar sequences.

(a) We have

$$\inf\{x_k\mid k\geq 0\}\leq \liminf_{k\to\infty}x_k\leq \limsup_{k\to\infty}x_k\leq \sup\{x_k\mid k\geq 0\}.$$

(b) $\{x_k\}$ converges if and only if

$$-\infty < \liminf_{k \to \infty} x_k = \limsup_{k \to \infty} x_k < \infty.$$

Furthermore, if $\{x_k\}$ converges, its limit is equal to the common scalar value of $\liminf_{k\to\infty} x_k$ and $\limsup_{k\to\infty} x_k$.

(c) If $x_k \leq y_k$ for all k, then

$$\liminf_{k \to \infty} x_k \le \liminf_{k \to \infty} y_k, \qquad \limsup_{k \to \infty} x_k \le \limsup_{k \to \infty} y_k.$$

(d) We have

$$\liminf_{k \to \infty} x_k + \liminf_{k \to \infty} y_k \le \liminf_{k \to \infty} (x_k + y_k),$$

$$\limsup_{k \to \infty} x_k + \limsup_{k \to \infty} y_k \ge \limsup_{k \to \infty} (x_k + y_k).$$

A sequence $\{x_k\}$ of vectors in \Re^n is said to converge to some $x \in \Re^n$ if the *i*th component of x_k converges to the *i*th component of x for every i. We use the notations $x_k \to x$ and $\lim_{k \to \infty} x_k = x$ to indicate convergence for vector sequences as well. The sequence $\{x_k\}$ is called bounded if each

of its corresponding component sequences is bounded. It can be seen that $\{x_k\}$ is bounded if and only if there exists a scalar c such that $\|x_k\| \leq c$ for all k. An infinite subset of a sequence $\{x_k\}$ is called a *subsequence* of $\{x_k\}$. Thus a subsequence can itself be viewed as a sequence, and can be represented as a set $\{x_k \mid k \in \mathcal{K}\}$, where \mathcal{K} is an infinite subset of positive integers (the notation $\{x_k\}_{\mathcal{K}}$ will also be used).

A vector $x \in \Re^n$ is said to be a *limit point* of a sequence $\{x_k\}$ if there exists a subsequence of $\{x_k\}$ that converges to x.† The following is a classical result that will be used often.

Proposition A.2.3: (Bolzano-Weierstrass Theorem) A bounded sequence in \Re^n has at least one limit point.

$o(\cdot)$ Notation

For a function $h: \Re^n \mapsto \Re^m$ we write $h(x) = o(||x||^p)$, where p is a positive integer, if

$$\lim_{k \to \infty} \frac{h(x_k)}{\|x_k\|^p} = 0,$$

for all sequences $\{x_k\}$ such that $x_k \to 0$ and $x_k \neq 0$ for all k.

Closed and Open Sets

We say that x is a closure point of a subset X of \Re^n if there exists a sequence $\{x_k\} \subset X$ that converges to x. The closure of X, denoted $\operatorname{cl}(X)$, is the set of all closure points of X.

Definition A.2.2: A subset X of \Re^n is called *closed* if it is equal to its closure. It is called *open* if its complement, $\{x \mid x \notin X\}$, is closed. It is called *bounded* if there exists a scalar c such that $||x|| \le c$ for all $x \in X$. It is called *compact* if it is closed and bounded.

Given $x^* \in \Re^n$ and $\epsilon > 0$, the sets $\{x \mid \|x - x^*\| < \epsilon\}$ and $\{x \mid \|x - x^*\| \le \epsilon\}$ are called an *open sphere* and a *closed sphere* centered at

[†] Some authors prefer the alternative term "cluster point" of a sequence, and use the term "limit point of a set S" to indicate a point \overline{x} such that $\overline{x} \notin S$ and there exists a sequence $\{x_k\} \subset S$ that converges to \overline{x} . With this terminology, \overline{x} is a cluster point of a sequence $\{x_k \mid k=1,2,\ldots\}$ if and only if $(\overline{x},0)$ is a limit point of the set $\{(x_k,1/k)\mid k=1,2,\ldots\}$. Our use of the term "limit point" of a sequence is quite popular in optimization and should not lead to any confusion.

 x^* . Sometimes the terms open ball and closed ball are used, respectively. A consequence of the definitions, is that a subset X of \Re^n is open if and only if for every $x \in X$ there is an open sphere that is centered at x and is contained in X. A neighborhood of a vector x is an open set containing x.

Definition A.2.3: We say that x is an *interior point* of a subset X of \Re^n if there exists a neighborhood of x that is contained in X. The set of all interior points of X is called the *interior* of X, and is denoted by $\operatorname{int}(X)$. A vector $x \in \operatorname{cl}(X)$ which is not an interior point of X is said to be a *boundary point* of X. The set of all boundary points of X is called the *boundary* of X.

Proposition A.2.4:

- (a) The union of a finite collection of closed sets is closed.
- (b) The intersection of any collection of closed sets is closed.
- (c) The union of any collection of open sets is open.
- (d) The intersection of a finite collection of open sets is open.
- (e) A set is open if and only if all of its elements are interior points.
- (f) Every subspace of \Re^n is closed.
- (g) A set X is compact if and only if every sequence of elements of X has a subsequence that converges to an element of X.
- (h) If $\{X_k\}$ is a sequence of nonempty and compact sets such that $X_{k+1} \subset X_k$ for all k, then the intersection $\bigcap_{k=0}^{\infty} X_k$ is nonempty and compact.

The topological properties of sets in \Re^n , such as being open, closed, or compact, do not depend on the norm being used. This is a consequence of the following proposition.

Proposition A.2.5: (Norm Equivalence Property)

(a) For any two norms $\|\cdot\|$ and $\|\cdot\|'$ on $\Re^n,$ there exists a scalar c such that

$$||x|| \le c||x||', \qquad \forall \ x \in \Re^n.$$

(b) If a subset of \Re^n is open (respectively, closed, bounded, or compact) with respect to some norm, it is open (respectively, closed, bounded, or compact) with respect to all other norms.

Continuity

Let $f: X \mapsto \Re^m$ be a function, where X is a subset of \Re^n , and let x be a vector in X. If there exists a vector $y \in \Re^m$ such that the sequence $\{f(x_k)\}$ converges to y for every sequence $\{x_k\} \subset X$ such that $\lim_{k \to \infty} x_k = x$, we write $\lim_{z \to x} f(z) = y$. If there exists a vector $y \in \Re^m$ such that the sequence $\{f(x_k)\}$ converges to y for every sequence $\{x_k\} \subset X$ such that $\lim_{k \to \infty} x_k = x$ and $x_k \leq x$ (respectively, $x_k \geq x$) for all k, we write $\lim_{z \uparrow x} f(z) = y$ [respectively, $\lim_{z \downarrow x} f(z)$].

Definition A.2.4: Let X be a subset of \Re^n .

- (a) A function $f: X \mapsto \Re^m$ is called *continuous* at a vector $x \in X$ if $\lim_{z \to x} f(z) = f(x)$.
- (b) A function $f: X \mapsto \Re^m$ is called *right-continuous* (respectively, left-continuous) at a vector $x \in X$ if $\lim_{z \downarrow x} f(z) = f(x)$ [respectively, $\lim_{z \uparrow x} f(z) = f(x)$].
- (c) A real-valued function $f: X \mapsto \Re$ is called upper semicontinuous (respectively, lower semicontinuous) at a vector $x \in X$ if $f(x) \ge \limsup_{k \to \infty} f(x_k)$ [respectively, $f(x) \le \liminf_{k \to \infty} f(x_k)$] for every sequence $\{x_k\} \subset X$ that converges to x.

If $f: X \mapsto \Re^m$ is continuous at every vector in a subset of its domain X, we say that f is continuous over that subset. If $f: X \mapsto \Re^m$ is continuous at every vector in its domain X, we say that f is continuous (without qualification). We use similar terminology for right-continuous, left-continuous, upper semicontinuous, and lower semicontinuous functions.

Proposition A.2.6:

- (a) Any vector norm on \Re^n is a continuous function.
- (b) Let $f: \Re^m \mapsto \Re^p$ and $g: \Re^n \mapsto \Re^m$ be continuous functions. The composition $f \cdot g: \Re^n \mapsto \Re^p$, defined by $(f \cdot g)(x) = f(g(x))$, is a continuous function.

- (c) Let $f: \Re^n \mapsto \Re^m$ be continuous, and let Y be an open (respectively, closed) subset of \Re^m . Then the inverse image of Y, $\left\{x \in \Re^n \mid f(x) \in Y\right\}$, is open (respectively, closed).
- (d) Let $f: \mathbb{R}^n \to \mathbb{R}^m$ be continuous, and let X be a compact subset of \mathbb{R}^n . Then the image of X, $\{f(x) \mid x \in X\}$, is compact.

If $f: \Re^n \to \Re$ is a continuous function and $X \subset \Re^n$ is compact, by Prop. A.2.6(c), the sets

$$V_{\gamma} = \{ x \in X \mid f(x) \le \gamma \}$$

are nonempty and compact for all $\gamma \in \Re$ with $\gamma > f^*$, where

$$f^* = \inf_{x \in X} f(x).$$

Since the set of minima of f is the intersection of the nonempty and compact sets V_{γ_k} for any sequence $\{\gamma_k\}$ with $\gamma_k \downarrow f^*$ and $\gamma_k > f^*$ for all k, it follows from Prop. A.2.4(h) that the set of minima is nonempty. This proves the classical theorem of Weierstrass, which is discussed further and generalized considerably in Section 3.2.

Proposition A.2.7: (Weierstrass' Theorem for Continuous Functions) A continuous function $f: \mathbb{R}^n \to \mathbb{R}$ attains a minimum over any compact subset of \mathbb{R}^n .

A.3 DERIVATIVES

Let $f:\Re^n\mapsto\Re$ be some function, fix some $x\in\Re^n,$ and consider the expression

$$\lim_{\alpha \to 0} \frac{f(x + \alpha e_i) - f(x)}{\alpha},$$

where e_i is the *i*th unit vector (all components are 0 except for the *i*th component which is 1). If the above limit exists, it is called the *i*th partial derivative of f at the vector x and it is denoted by $(\partial f/\partial x_i)(x)$ or $\partial f(x)/\partial x_i$ (x_i in this section will denote the *i*th component of the vector x). Assuming all of these partial derivatives exist, the gradient of f at x is defined as the column vector

$$\nabla f(x) = \begin{bmatrix} \frac{\partial f(x)}{\partial x_1} \\ \vdots \\ \frac{\partial f(x)}{\partial x_n} \end{bmatrix}.$$

For any $d \in \mathbb{R}^n$, we define the one-sided directional derivative of f at a vector x in the direction d by

$$f'(x;d) = \lim_{\alpha \downarrow 0} \frac{f(x + \alpha d) - f(x)}{\alpha},$$

provided that the limit exists.

If the directional derivative of f at a vector x exists in all directions and f'(x;d) is a linear function of d, we say that f is differentiable at x. It can be seen that f is differentiable at x if and only if the gradient $\nabla f(x)$ exists and satisfies $\nabla f(x)'d = f'(x;d)$ for all $d \in \Re^n$, or equivalently

$$f(x + \alpha d) = f(x) + \alpha \nabla f(x)'d + o(|\alpha|), \quad \forall \alpha \in \Re.$$

The function f is called differentiable over a subset S of \mathbb{R}^n if it is differentiable at every $x \in S$. The function f is called differentiable (without qualification) if it is differentiable at all $x \in \mathbb{R}^n$.

If f is differentiable over an open set S and $\nabla f(\cdot)$ is continuous at all $x \in S$, f is said to be *continuously differentiable over* S. It can then be shown that for any $x \in S$ and norm $\|\cdot\|$,

$$f(x+d) = f(x) + \nabla f(x)'d + o(||d||), \quad \forall d \in \Re^n.$$

If each one of the partial derivatives of a function $f: \Re^n \mapsto \Re$ is a continuously differentiable function of x over an open set S, we say that f is twice continuously differentiable over S. We then denote by

$$\frac{\partial^2 f(x)}{\partial x_i \partial x_j}$$

the *i*th partial derivative of $\partial f/\partial x_j$ at a vector $x \in \Re^n$. The *Hessian* of f at x, denoted by $\nabla^2 f(x)$, is the matrix whose components are the above second derivatives. The matrix $\nabla^2 f(x)$ is symmetric. The following theorem will be useful to us.

Proposition A.3.1: (Mean Value Theorem) Let $f: \mathbb{R}^n \to \mathbb{R}$ be differentiable over an open sphere S, let x be a vector in S, and let d be such that $x + d \in S$. Then there exists an $\alpha \in [0, 1]$ such that

$$f(x+d) = f(x) + \nabla f(x + \alpha d)'d.$$

If in addition f is twice continuously differentiable over S, there exists an $\alpha \in [0,1]$ such that

$$f(x+d) = f(x) + \nabla f(x)'d + \frac{1}{2}d'\nabla^2 f(x+\alpha d)d.$$

Subdifferential of the Maximum of a Convex Function

A case of great interest in optimization involves functions of the form

$$f(x) = \max_{z \in Z} \phi(x, z).$$

The directional derivative and the subdifferential of f can be described in terms of the directional derivative and the subdifferential of ϕ , evaluated at points \overline{z} where the maximum is attained, as shown by the following proposition, which extends Example 5.4.5. The proposition was stated and proved in the form given here in the books [BeT89] (Prop. A.43) and [BNO03] (Prop. 4.5.1).

Proposition A.3.2: (Danskin's Theorem) Let $Z \subset \mathbb{R}^m$ be a compact set, and let $\phi : \mathbb{R}^n \times Z \mapsto \mathbb{R}$ be continuous and such that $\phi(\cdot, z) : \mathbb{R}^n \mapsto \mathbb{R}$ is convex for each $z \in Z$.

(a) The function $f: \Re^n \mapsto \Re$ given by

$$f(x) = \max_{z \in Z} \phi(x, z) \tag{A.1}$$

is convex and has directional derivative given by

$$f'(x;y) = \max_{z \in Z(x)} \phi'(x,z;y),$$

where $\phi'(x, z; y)$ is the directional derivative of the function $\phi(\cdot, z)$ at x in the direction y, and Z(x) is the set of maximizing points in Eq. (A.1)

$$Z(x) = \left\{ \overline{z} \; \middle| \; \phi(x, \overline{z}) = \max_{z \in Z} \phi(x, z) \right\}.$$

In particular, if Z(x) consists of a unique point \overline{z} and $\phi(\cdot, \overline{z})$ is differentiable at x, then f is differentiable at x, and $\nabla f(x) = \nabla_x \phi(x, \overline{z})$, where $\nabla_x \phi(x, \overline{z})$ is the vector with coordinates

$$\frac{\partial \phi(x,\overline{z})}{\partial x_i}, \qquad i = 1, \dots, n.$$

(b) If $\phi(\cdot, z)$ is differentiable for all $z \in Z$ and $\nabla_x \phi(x, \cdot)$ is continuous on Z for each x, then

$$\partial f(x) = \operatorname{conv} \{ \nabla_x \phi(x, z) \mid z \in Z(x) \}, \quad \forall \ x \in \Re^n. \quad (A.2)$$

In particular, if ϕ is linear in x for all $z \in Z$, i.e.,

$$\phi(x,z) = a_z'x + b_z, \qquad \forall \ z \in Z,$$

then

$$\partial f(x) = \operatorname{conv}\{a_z \mid z \in Z(x)\}.$$

The preceding proposition derives its origin from a theorem by Danskin [Dan67] that provides a formula for the directional derivative of the maximum of a (not necessarily convex) directionally differentiable function. When adapted to a convex function f, this formula yields the expression (A.2) for $\partial f(x)$.

Danskin's Theorem for Extended Real-Valued Convex Functions

Let us also note an extension of Danskin's Theorem [Prop. A.3.2(b)], which provides a more general formula for the subdifferential $\partial f(x)$ of the function

$$f(x) = \sup_{z \in Z} \phi(x, z), \tag{A.3}$$

where Z is a compact set. This version of the theorem does not require that $\phi(\cdot, z)$ is differentiable. Instead it assumes that $\phi(\cdot, z)$ is an extended real-valued closed proper convex function for each $z \in Z$, that $\operatorname{int}(\operatorname{dom}(f))$ [the interior of the set $\operatorname{dom}(f) = \{x \mid f(x) < \infty\}$] is nonempty, and that ϕ is continuous on the set $\operatorname{int}(\operatorname{dom}(f)) \times Z$. Then for all $x \in \operatorname{int}(\operatorname{dom}(f))$, we have

$$\partial f(x) = \operatorname{conv} \{ \partial \phi(x, z) \mid z \in Z(x) \}, \tag{A.4}$$

where $\partial \phi(x, z)$ is the subdifferential of $\phi(\cdot, z)$ at x for any $z \in Z$, and Z(x) is the set of maximizing points in Eq. (A.3). This result was first obtained by the author as part of his Ph.D. thesis (see [Ber71], Prop. A.22, for a formal statement and proof).

Note that the nonemptiness of $\operatorname{int}(\operatorname{dom}(f))$ is an essential assumption for the formula (A.4) to hold. In particular, the formula may not hold if instead we just assume that the relative interior of $\operatorname{dom}(f)$ is nonempty. For an example, consider the two spheres in \Re^2

$$S_1 = \{(x_1, x_2) \mid (x_1 - 1)^2 + x_2^2 \le 1\}, \qquad S_2 = \{(x_1, x_2) \mid (x_1 + 1)^2 + x_2^2 \le 1\},$$
 let f_1 and f_2 be the indicator functions of S_1 and S_2 , respectively,

$$f_1(x) = \begin{cases} 0 & \text{if } x \in S_1, \\ \infty & \text{if } x \notin S_1, \end{cases}, \qquad f_2(x) = \begin{cases} 0 & \text{if } x \in S_2, \\ \infty & \text{if } x \notin S_2, \end{cases}$$

and let

$$f(x) = \max \{f_1(x), f_2(x)\} = \begin{cases} 0 & \text{if } x = 0, \\ \infty & \text{if } x \neq 0. \end{cases}$$

Then it can be seen that the formula (A.4) does not hold at x = 0.

Subdifferential of the Expected Value of a Convex Function

Another important subdifferential formula relates to the subgradients of an expected value function

$$f(x) = E\{F(x,\omega)\},\$$

where ω is a random variable taking values in a set Ω , and $F(\cdot, \omega) : \Re^n \mapsto \Re$ is a real-valued convex function such that f is real-valued (note that f is easily verified to be convex). If ω takes a finite number of values with probabilities $p(\omega)$, then the formulas

$$f'(x;d) = E\{F'(x,\omega;d)\}, \qquad \partial f(x) = E\{\partial F(x,\omega)\},$$
 (A.5)

hold because they can be written in terms of finite sums as

$$f'(x;d) = \sum_{\omega \in \Omega} p(\omega) F'(x,\omega;d), \qquad \partial f(x) = \sum_{\omega \in \Omega} p(\omega) \partial F(x,\omega),$$

so Prop. A.3.2(e) applies. However, the formulas (A.5) hold even in the case where Ω is uncountably infinite, with appropriate mathematical interpretation of the integral of set-valued functions $E\{\partial F(x,\omega)\}$ as the set of integrals

$$\int_{\omega \in \Omega} g(x, \omega) \, dP(\omega),\tag{A.6}$$

where $g(x,\omega) \in \partial F(x,\omega)$, $\omega \in \Omega$ (measurability issues must be addressed in this context). For a formal proof and analysis, see the author's papers [Ber72], [Ber73], which also provide a necessary and sufficient condition for f to be differentiable, even when $F(\cdot,\omega)$ is not. In this connection, it is important to note that the integration over ω in Eq. (A.6) may smooth out the nondifferentiabilities of $F(\cdot,\omega)$ if ω is a "continuous" random variable. This property can be used in turn in algorithms, including schemes that bring to bear the methodology of differentiable optimization.

There is a very extensive literature on convex analysis and optimization and it is beyond our scope to give a complete bibliography. We are providing instead a brief historical account and list some of the main textbooks in the field.

Among early classical works on convexity, we mention Caratheodory [Car11], Minkowski [Min11], and Steinitz [Ste13], [Ste14], [Ste16]. In particular, Caratheodory gave the theorem on convex hulls that carries his name, while Steinitz developed the theory of relative interiors and recession cones. Minkowski is credited with initiating the theory of hyperplane separation of convex sets and the theory of support functions (a precursor to conjugate convex functions). Furthermore, Minkowski and Farkas (whose work, published in Hungarian, spans a 30-year period starting around 1894), are credited with laying the foundations of polyhedral convexity.

The work of Fenchel was instrumental in launching the modern era of convex analysis, when the subject came to a sharp focus thanks to its rich applications in optimization and game theory. In his 1951 lecture notes [Fen51], Fenchel laid the foundations of convex duality theory, and together with related works by von Neumann on saddle points and game theory, and Kuhn and Tucker on nonlinear programming [KuT51], inspired much subsequent work on convexity and its connections with optimization. Furthermore, Fenchel developed most of the topics that are fundamental in our exposition, such as the theory of conjugate convex functions (introduced earlier in a more limited form by Legendre), and the theory of subgradients.

There are several books that relate to both convex analysis and optimization. The book by Rockafellar [Roc70] has been an important influence to subsequent convex optimization books, including the present work, but unfortunately it is written in a style that does not facilitate intuition through visualization (it does not contain a single figure). The book by Rockafellar and Wets [RoW98] is an extensive treatment of "variational analysis," a broad spectrum of topics that integrate classical analysis, convexity, and optimization of both convex and nonconvex (possibly nonsmooth) functions. Stoer and Witzgall [StW70] discuss similar topics as Rockafellar [Roc70] but less comprehensively. Ekeland and Temam

[EkT76], and Zalinescu [Zal02] develop the subject in infinite dimensional spaces. Hiriart-Urruty and Lemarechal [HiL93] emphasize algorithms for dual and nondifferentiable optimization. Rockafellar [Roc84] focuses on convexity and duality in network optimization, and an important generalization, called monotropic programming. Bertsekas [Ber98] also gives a detailed coverage of this material, which owes much to the early work of Minty [Min60] on network optimization. Schrijver [Sch86] provides an extensive account of polyhedral convexity with applications to integer programming and combinatorial optimization, and gives many historical references. Bonnans and Shapiro [BoS00] emphasize sensitivity analysis and discuss infinite dimensional problems as well. Borwein and Lewis BoL00 develop many of the concepts in Rockafellar and Wets [RoW98], but more succinctly. The author's earlier book with Nedić and Ozdaglar [BNO03] also straddles the boundary between convex and variational analysis. Ben-Tal and Nemirovski [BeN01] focus on conic and semidefinite programming [see also the 2005 class notes by Nemirovski (on line)]. Auslender and Teboulle [AuT03] emphasize the question of existence of solutions for convex as well as nonconvex optimization problems, and related issues in duality theory and variational inequalities (a more recent extensive treatment of the existence of solutions question is given in the paper by Bertsekas and Tseng [BeT07]). Boyd and Vanderbergue [BoV04] discuss many applications of convex optimization.

We also note a few books that focus on the geometry and other properties of convex sets, but have limited connection with duality and optimization: Bonnesen and Fenchel [BoF34], Eggleston [Egg58], Valentine [Val64], Grunbaum [Gru67], Webster [Web94], and Barvinok [Bar02].

The MC/MC framework differentiates the present book from alternative treatments of convex optimization. It was initially developed by the author in joint research with A. Nedić and A. Ozdaglar, which is described in the book [BNO03]. The present account is improved and more comprehensive. In particular, it contains more streamlined proofs and some new results, particularly in connection with minimax problems (Sections 4.2.5 and 5.7.2), and nonconvex problems (Section 5.7, which generalizes the work on duality gap estimates in [Ber82], Section 5.6.1).

In a nutshell, the MC/MC framework aims to reduce duality theory to its bare essentials. The starting point is the two dual descriptions of a closed convex set: as the union of its points and as the intersection of the halfspaces that contain it. Thus a closed convex function f admits two dual descriptions: as the union of the points of its epigraph epi(f) and as the intersection of the halfspaces that contain epi(f) (the latter description defines the conjugate f^*).

Fenchel duality provides dual descriptions of optimization problems through the use of conjugate functions, but in the author's view, this is an indirect and far too complicated starting point for optimization duality. A simpler and more fundamental approach is to start from a single set and

the two dual problems that it defines: the min common point problem and max crossing point problem. Developing duality theory around these two problems leads to a simpler, unified, and visually transparent analysis of the central frameworks of convex optimization: Lagrange duality, Fenchel duality, min-max theory, and theorems of the alternative.

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Convex Optimization Theory

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Exercises and Solutions

Many of the exercises and solutions given in this manual were developed as part of my 2003 convex optimization book (coauthored with Angelia Nedić and Asuman Ozdaglar), and are posted on the internet at that book's web site. The contribution of my coauthors in the development of these exercises and their solutions is gratefully acknowledged. Since some of the exercises and/or their solutions have been modified and also new exercises have been added, all errors are my sole responsibility.

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CHAPTER 1: EXERCISES AND SOLUTIONS

SECTION 1.1: Convex Sets and Functions

1.1

Let C be a nonempty subset of \Re^n , and let λ_1 and λ_2 be positive scalars. Show that if C is convex, then $(\lambda_1 + \lambda_2)C = \lambda_1C + \lambda_2C$ [cf. Prop. 1.1.1(c)]. Show by example that this need not be true when C is not convex.

Solution: We always have $(\lambda_1 + \lambda_2)C \subset \lambda_1C + \lambda_2C$, even if C is not convex. To show the reverse inclusion assuming C is convex, note that a vector x in $\lambda_1C + \lambda_2C$ is of the form $x = \lambda_1x_1 + \lambda_2x_2$, where $x_1, x_2 \in C$. By convexity of C, we have

$$\frac{\lambda_1}{\lambda_1 + \lambda_2} x_1 + \frac{\lambda_2}{\lambda_1 + \lambda_2} x_2 \in C,$$

and it follows that

$$x = \lambda_1 x_1 + \lambda_2 x_2 \in (\lambda_1 + \lambda_2)C,$$

so $\lambda_1 C + \lambda_2 C \subset (\lambda_1 + \lambda_2) C$. For a counterexample when C is not convex, let C be a set in \Re^n consisting of two vectors, 0 and $x \neq 0$, and let $\lambda_1 = \lambda_2 = 1$. Then C is not convex, and $(\lambda_1 + \lambda_2)C = 2C = \{0, 2x\}$, while $\lambda_1 C + \lambda_2 C = C + C = \{0, x, 2x\}$, showing that $(\lambda_1 + \lambda_2)C \neq \lambda_1 C + \lambda_2 C$.

1.2 (Properties of Cones)

Show that:

- (a) The intersection $\cap_{i \in I} C_i$ of a collection $\{C_i \mid i \in I\}$ of cones is a cone.
- (b) The Cartesian product $C_1 \times C_2$ of two cones C_1 and C_2 is a cone.
- (c) The vector sum $C_1 + C_2$ of two cones C_1 and C_2 is a cone.
- (d) The image and the inverse image of a cone under a linear transformation is a cone.
- (e) A subset C is a convex cone if and only if it is closed under addition and positive scalar multiplication, i.e., $C+C\subset C$, and $\gamma C\subset C$ for all $\gamma>0$.

Solution: (a) Let $x \in \cap_{i \in I} C_i$ and let α be a positive scalar. Since $x \in C_i$ for all $i \in I$ and each C_i is a cone, the vector αx belongs to C_i for all $i \in I$. Hence, $\alpha x \in \cap_{i \in I} C_i$, showing that $\cap_{i \in I} C_i$ is a cone.

- (b) Let $x \in C_1 \times C_2$ and let α be a positive scalar. Then $x = (x_1, x_2)$ for some $x_1 \in C_1$ and $x_2 \in C_2$, and since C_1 and C_2 are cones, it follows that $\alpha x_1 \in C_1$ and $\alpha x_2 \in C_2$. Hence, $\alpha x = (\alpha x_1, \alpha x_2) \in C_1 \times C_2$, showing that $C_1 \times C_2$ is a cone.
- (c) Let $x \in C_1 + C_2$ and let α be a positive scalar. Then, $x = x_1 + x_2$ for some $x_1 \in C_1$ and $x_2 \in C_2$, and since C_1 and C_2 are cones, $\alpha x_1 \in C_1$ and $\alpha x_2 \in C_2$. Hence, $\alpha x = \alpha x_1 + \alpha x_2 \in C_1 + C_2$, showing that $C_1 + C_2$ is a cone.
- (d) First we prove that $A\cdot C$ is a cone, where A is a linear transformation and $A\cdot C$ is the image of C under A. Let $z\in A\cdot C$ and let α be a positive scalar. Then, Ax=z for some $x\in C$, and since C is a cone, $\alpha x\in C$. Because $A(\alpha x)=\alpha z$, the vector αz is in $A\cdot C$, showing that $A\cdot C$ is a cone. Next we prove that the inverse image $A^{-1}\cdot C$ of C under A is a cone. Let $x\in A^{-1}\cdot C$ and let α be a positive scalar. Then $Ax\in C$, and since C is a cone, $\alpha Ax\in C$. Thus, the vector $A(\alpha x)\in C$, implying that $\alpha x\in A^{-1}\cdot C$, and showing that $A^{-1}\cdot C$ is a cone.
- (e) Let C be a convex cone. Then $\gamma C \subset C$, for all $\gamma > 0$, by the definition of cone. Furthermore, by convexity of C, for all $x,y \in C$, we have $z \in C$, where

$$z = \frac{1}{2}(x+y).$$

Hence $(x+y)=2z\in C$, since C is a cone, and it follows that $C+C\subset C$. Conversely, assume that $C+C\subset C$, and $\gamma C\subset C$. Then C is a cone. Furthermore, if $x,y\in C$ and $\alpha\in(0,1)$, we have $\alpha x\in C$ and $(1-\alpha)y\in C$, and $\alpha x+(1-\alpha)y\in C$ (since $C+C\subset C$). Hence C is convex.

1.3 (Convexity under Composition)

Let C be a nonempty convex subset of \Re^n . Let also $f=(f_1,\ldots,f_m)$, where $f_i:C\mapsto\Re$, $i=1,\ldots,m$, are convex functions, and let $g:\Re^m\mapsto\Re$ be a function that is convex and monotonically nondecreasing over a convex set that contains the set $\{f(x)\mid x\in C\}$, in the sense that for all u,\overline{u} in this set such that $u\leq\overline{u}$, we have $g(u)\leq g(\overline{u})$. Show that the function h defined by h(x)=g(f(x)) is convex over C. If in addition, m=1,g is monotonically increasing and f is strictly convex, then h is strictly convex.

Solution: Let $x, y \in \Re^n$ and let $\alpha \in [0, 1]$. By the definitions of h and f, we have

$$h(\alpha x + (1 - \alpha)y) = g(f(\alpha x + (1 - \alpha)y))$$

$$= g(f_1(\alpha x + (1 - \alpha)y), \dots, f_m(\alpha x + (1 - \alpha)y))$$

$$\leq g(\alpha f_1(x) + (1 - \alpha)f_1(y), \dots, \alpha f_m(x) + (1 - \alpha)f_m(y))$$

$$= g\Big(\alpha\big(f_1(x), \dots, f_m(x)\big) + (1 - \alpha)\big(f_1(y), \dots, f_m(y)\big)\Big)$$

$$\leq \alpha g\big(f_1(x), \dots, f_m(x)\big) + (1 - \alpha)g\big(f_1(y), \dots, f_m(y)\big)$$

$$= \alpha g\big(f(x)\big) + (1 - \alpha)g\big(f(y)\big)$$

$$= \alpha h(x) + (1 - \alpha)h(y),$$

where the first inequality follows by convexity of each f_i and monotonicity of g, while the second inequality follows by convexity of g.

If m=1, g is monotonically increasing, and f is strictly convex, then the first inequality is strict whenever $x \neq y$ and $\alpha \in (0,1)$, showing that h is strictly convex.

1.4 (Examples of Convex Functions)

Show that the following functions from \Re^n to $(-\infty, \infty]$ are convex:

(a) $f_1(x_1,\ldots,x_n) = \begin{cases} -(x_1x_2\cdots x_n)^{\frac{1}{n}} & \text{if } x_1 > 0,\ldots,x_n > 0, \\ \infty & \text{otherwise.} \end{cases}$

- (b) $f_2(x) = \ln(e^{x_1} + \dots + e^{x_n}).$
- (c) $f_3(x) = ||x||^p$ with $p \ge 1$.
- (d) $f_4(x) = \frac{1}{f(x)}$, where f is concave and $0 < f(x) < \infty$ for all x.
- (e) $f_5(x) = \alpha f(x) + \beta$, where $f: \Re^n \to \Re$ is a convex function, and α and β are scalars, with $\alpha > 0$.
- (f) $f_6(x) = e^{\beta x'Ax}$, where A is a positive semidefinite symmetric $n \times n$ matrix and β is a positive scalar.
- (g) $f_7(x) = f(Ax + b)$, where $f: \Re^m \to \Re$ is a convex function, A is an $m \times n$ matrix, and b is a vector in \Re^m .

Solution: (a) Denote $X = \text{dom}(f_1)$. It can be seen that f_1 is twice continuously differentiable over X and its Hessian matrix is given by

$$\nabla^2 f_1(x) = \frac{f_1(x)}{n^2} \begin{bmatrix} \frac{1-n}{x_1^2} & \frac{1}{x_1 x_2} & \cdots & \frac{1}{x_1 x_n} \\ \frac{1}{x_2 x_1} & \frac{1-n}{x_2^2} & \cdots & \frac{1}{x_2 x_n} \\ & & \vdots & \\ \frac{1}{x_n x_1} & \frac{1}{x_1 x_2} & \cdots & \frac{1-n}{x_n^2} \end{bmatrix}$$

for all $x = (x_1, \ldots, x_n) \in X$. From this, direct computation shows that for all $z = (z_1, \ldots, z_n) \in \Re^n$ and $x = (x_1, \ldots, x_n) \in X$, we have

$$z'\nabla^2 f_1(x)z = \frac{f_1(x)}{n^2} \left(\left(\sum_{i=1}^n \frac{z_i}{x_i} \right)^2 - n \sum_{i=1}^n \left(\frac{z_i}{x_i} \right)^2 \right).$$

Note that this quadratic form is nonnegative for all $z \in \mathbb{R}^n$ and $x \in X$, since $f_1(x) < 0$, and for any real numbers $\alpha_1, \ldots, \alpha_n$, we have

$$(\alpha_1 + \dots + \alpha_n)^2 \le n(\alpha_1^2 + \dots + \alpha_n^2),$$

in view of the fact that $2\alpha_j\alpha_k \leq \alpha_j^2 + \alpha_k^2$. Hence, $\nabla^2 f_1(x)$ is positive semidefinite for all $x \in X$, and it follows from Prop. 1.1.10(a) that f_1 is convex.

(b) We show that the Hessian of f_2 is positive semidefinite at all $x \in \Re^n$. Let $\beta(x) = e^{x_1} + \cdots + e^{x_n}$. Then a straightforward calculation yields

$$z'\nabla^2 f_2(x)z = \frac{1}{\beta(x)^2} \sum_{i=1}^n \sum_{j=1}^n e^{(x_i + x_j)} (z_i - z_j)^2 \ge 0, \quad \forall \ z \in \Re^n.$$

Hence by Prop. 1.1.10(a), f_2 is convex.

- (c) The function $f_3(x) = ||x||^p$ can be viewed as a composition g(f(x)) of the scalar function $g(t) = t^p$ with $p \ge 1$ and the function f(x) = ||x||. In this case, g is convex and monotonically increasing over the nonnegative axis, the set of values that f can take, while f is convex over \Re^n (since any vector norm is convex). Using Exercise 1.3, it follows that the function $f_3(x) = ||x||^p$ is convex over \Re^n .
- (d) The function $f_4(x) = \frac{1}{f(x)}$ can be viewed as a composition g(h(x)) of the function $g(t) = -\frac{1}{t}$ for t < 0 and the function h(x) = -f(x) for $x \in \Re^n$. In this case, the g is convex and monotonically increasing in the set $\{t \mid t < 0\}$, while h is convex over \Re^n . Using Exercise 1.3, it follows that the function $f_4(x) = \frac{1}{f(x)}$ is convex over \Re^n .
- (e) The function $f_5(x) = \alpha f(x) + \beta$ can be viewed as a composition g(f(x)) of the function $g(t) = \alpha t + \beta$, where $t \in \Re$, and the function f(x) for $x \in \Re^n$. In this case, g is convex and monotonically increasing over \Re (since $\alpha \geq 0$), while f is convex over \Re^n . Using Exercise 1.3, it follows that f_5 is convex over \Re^n .
- (f) The function $f_6(x) = e^{\beta x'Ax}$ can be viewed as a composition g(f(x)) of the function $g(t) = e^{\beta t}$ for $t \in \Re$ and the function f(x) = x'Ax for $x \in \Re^n$. In this case, g is convex and monotonically increasing over \Re , while f is convex over \Re^n (since A is positive semidefinite). Using Exercise 1.3, it follows that f_6 is convex over \Re^n .
- (g) This part is straightforward using the definition of a convex function.

1.5 (Ascent/Descent Behavior of a Convex Function)

Let $f: \Re \mapsto \Re$ be a convex function.

(a) (Monotropic Property) Use the definition of convexity to show that f is "turning upwards" in the sense that if x_1, x_2, x_3 are three scalars such that $x_1 < x_2 < x_3$, then

$$\frac{f(x_2) - f(x_1)}{x_2 - x_1} \le \frac{f(x_3) - f(x_2)}{x_3 - x_2}.$$

(b) Use part (a) to show that there are four possibilities as x increases to ∞ : (1) f(x) decreases monotonically to $-\infty$, (2) f(x) decreases monotonically to a finite value, (3) f(x) reaches some value and stays at that value, (4) f(x) increases monotonically to ∞ when $x \geq \overline{x}$ for some $\overline{x} \in \Re$.

Solution: (a) Let x_1, x_2, x_3 be three scalars such that $x_1 < x_2 < x_3$. Then we can write x_2 as a convex combination of x_1 and x_3 as follows

$$x_2 = \frac{x_3 - x_2}{x_3 - x_1} x_1 + \frac{x_2 - x_1}{x_3 - x_1} x_3,$$

so that by convexity of f, we obtain

$$f(x_2) \le \frac{x_3 - x_2}{x_3 - x_1} f(x_1) + \frac{x_2 - x_1}{x_3 - x_1} f(x_3).$$

This relation and the fact

$$f(x_2) = \frac{x_3 - x_2}{x_3 - x_1} f(x_2) + \frac{x_2 - x_1}{x_3 - x_1} f(x_2),$$

imply that

$$\frac{x_3 - x_2}{x_3 - x_1} (f(x_2) - f(x_1)) \le \frac{x_2 - x_1}{x_3 - x_1} (f(x_3) - f(x_2)).$$

By multiplying the preceding relation with $x_3 - x_1$ and by dividing it with $(x_3 - x_2)(x_2 - x_1)$, we obtain

$$\frac{f(x_2) - f(x_1)}{x_2 - x_1} \le \frac{f(x_3) - f(x_2)}{x_3 - x_2}.$$

(b) Let $\{x_k\}$ be an increasing scalar sequence, i.e., $x_1 < x_2 < x_3 < \cdots$. Then according to part (a), we have for all k

$$\frac{f(x_2) - f(x_1)}{x_2 - x_1} \le \frac{f(x_3) - f(x_2)}{x_3 - x_2} \le \dots \le \frac{f(x_{k+1}) - f(x_k)}{x_{k+1} - x_k}.$$
 (1.1)

Since $(f(x_k) - f(x_{k-1}))/(x_k - x_{k-1})$ is monotonically nondecreasing, we have

$$\frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}} \to \gamma, \tag{1.2}$$

where γ is either a real number or ∞ . Furthermore,

$$\frac{f(x_{k+1}) - f(x_k)}{x_{k+1} - x_k} \le \gamma, \qquad \forall \ k. \tag{1.3}$$

We now show that γ is independent of the sequence $\{x_k\}$. Let $\{y_j\}$ be any increasing scalar sequence. For each j, choose x_{k_j} such that $y_j < x_{k_j}$ and $x_{k_1} < x_{k_2} < \cdots < x_{k_j}$, so that we have $y_j < y_{j+1} < x_{k_{j+1}} < x_{k_{j+2}}$. By part (a), it follows that

$$\frac{f(y_{j+1}) - f(y_j)}{y_{j+1} - y_j} \le \frac{f(x_{k_{j+2}}) - f(x_{k_{j+1}})}{x_{k_{j+2}} - x_{k_{j+1}}},$$

and letting $j \to \infty$ yields

$$\lim_{j \to \infty} \frac{f(y_{j+1}) - f(y_j)}{y_{j+1} - y_j} \le \gamma.$$

Similarly, by exchanging the roles of $\{x_k\}$ and $\{y_j\}$, we can show that

$$\lim_{j \to \infty} \frac{f(y_{j+1}) - f(y_j)}{y_{j+1} - y_j} \ge \gamma.$$

Thus the limit in Eq. (1.2) is independent of the choice for $\{x_k\}$, and Eqs. (1.1) and (1.3) hold for any increasing scalar sequence $\{x_k\}$.

We consider separately each of the three possibilities $\gamma < 0, \gamma = 0$, and $\gamma > 0$. First, suppose that $\gamma < 0$, and let $\{x_k\}$ be any increasing sequence. By using Eq. (1.3), we obtain

$$f(x_k) = \sum_{j=1}^{k-1} \frac{f(x_{j+1}) - f(x_j)}{x_{j+1} - x_j} (x_{j+1} - x_j) + f(x_1)$$

$$\leq \sum_{j=1}^{k-1} \gamma(x_{j+1} - x_j) + f(x_1)$$

$$= \gamma(x_k - x_1) + f(x_1),$$

and since $\gamma < 0$ and $x_k \to \infty$, it follows that $f(x_k) \to -\infty$. To show that f decreases monotonically, pick any x and y with x < y, and consider the sequence $x_1 = x$, $x_2 = y$, and $x_k = y + k$ for all $k \ge 3$. By using Eq. (1.3) with k = 1, we have

$$\frac{f(y) - f(x)}{y - x} \le \gamma < 0,$$

so that f(y) - f(x) < 0. Hence f decreases monotonically to $-\infty$, corresponding to case (1).

Suppose now that $\gamma = 0$, and let $\{x_k\}$ be any increasing sequence. Then, by Eq. (1.3), we have $f(x_{k+1}) - f(x_k) \leq 0$ for all k. If $f(x_{k+1}) - f(x_k) < 0$ for all k, then f decreases monotonically. To show this, pick any x and y with x < y, and consider a new sequence given by $y_1 = x$, $y_2 = y$, and $y_k = x_{K+k-3}$ for all $k \geq 3$, where K is large enough so that $y < x_K$. By using Eqs. (1.1) and (1.3) with $\{y_k\}$, we have

$$\frac{f(y) - f(x)}{y - x} \le \frac{f(x_{K+1}) - f(x_K)}{x_{K+1} - x_K} < 0,$$

implying that f(y) - f(x) < 0. Hence f decreases monotonically, and it may decrease to $-\infty$ or to a finite value, corresponding to cases (1) or (2), respectively.

If for some K we have $f(x_{K+1}) - f(x_K) = 0$, then by Eqs. (1.1) and (1.3) where $\gamma = 0$, we obtain $f(x_k) = f(x_K)$ for all $k \geq K$. To show that f stays at the value $f(x_K)$ for all $x \geq x_K$, choose any x such that $x > x_K$, and define $\{y_k\}$ as $y_1 = x_K$, $y_2 = x$, and $y_k = x_{N+k-3}$ for all $k \geq 3$, where N is large enough so that $x < x_N$. By using Eqs. (1.1) and (1.3) with $\{y_k\}$, we have

$$\frac{f(x) - f(x_K)}{x - x_K} \le \frac{f(x_N) - f(x)}{x_N - x} \le 0,$$

so that $f(x) \leq f(x_K)$ and $f(x_N) \leq f(x)$. Since $f(x_K) = f(x_N)$, we have $f(x) = f(x_K)$. Hence $f(x) = f(x_K)$ for all $x \geq x_K$, corresponding to case (3).

Finally, suppose that $\gamma > 0$, and let $\{x_k\}$ be any increasing sequence. Since $(f(x_k) - f(x_{k-1}))/(x_k - x_{k-1})$ is nondecreasing and tends to γ [cf. Eqs. (1.2) and (1.3)], there is a positive integer K and a positive scalar ϵ with $\epsilon < \gamma$ such that

$$\epsilon \le \frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}}, \qquad \forall \ k \ge K. \tag{1.4}$$

Therefore, for all k > K

$$f(x_k) = \sum_{j=K}^{k-1} \frac{f(x_{j+1}) - f(x_j)}{x_{j+1} - x_j} (x_{j+1} - x_j) + f(x_K) \ge \epsilon(x_k - x_K) + f(x_K),$$

implying that $f(x_k) \to \infty$. To show that f(x) increases monotonically to ∞ for all $x \geq x_K$, pick any x < y satisfying $x_K < x < y$, and consider a sequence given by $y_1 = x_K$, $y_2 = x$, $y_3 = y$, and $y_k = x_{N+k-4}$ for $k \geq 4$, where N is large enough so that $y < x_N$. By using Eq. (1.4) with $\{y_k\}$, we have

$$\epsilon \le \frac{f(y) - f(x)}{y - x}.$$

Thus f(x) increases monotonically to ∞ for all $x \geq x_K$, corresponding to case (4) with $\overline{x} = x_K$.

1.6 (Posynomials)

A posynomial is a function of positive scalar variables y_1, \ldots, y_n of the form

$$g(y_1, \dots, y_n) = \sum_{i=1}^m \beta_i y_1^{a_{i1}} \cdots y_n^{a_{in}},$$

where a_{ij} and β_i are scalars, such that $\beta_i > 0$ for all i. Show the following:

- (a) A posynomial need not be convex.
- (b) By a logarithmic change of variables, where we set

$$f(x) = \ln(g(y_1, \dots, y_n)),$$
 $b_i = \ln \beta_i, \ \forall i, \quad x_j = \ln y_j, \ \forall j,$

we obtain a convex function

$$f(x) = \ln \exp(Ax + b), \quad \forall x \in \Re^n,$$

where $\exp(z) = e^{z_1} + \dots + e^{z_m}$ for all $z \in \mathbb{R}^m$, A is an $m \times n$ matrix with components a_{ij} , and $b \in \mathbb{R}^m$ is a vector with components b_i .

(c) Every function $g: \Re^n \mapsto \Re$ of the form

$$g(y) = g_1(y)^{\gamma_1} \cdots g_r(y)^{\gamma_r},$$

where g_k is a posynomial and $\gamma_k > 0$ for all k, can be transformed by a logarithmic change of variables into a convex function f given by

$$f(x) = \sum_{k=1}^{r} \gamma_k \ln \exp(A_k x + b_k),$$

with the matrix A_k and the vector b_k being associated with the posynomial g_k for each k.

Solution: (a) Consider the following posynomial for which we have n = m = 1 and $\beta = \frac{1}{2}$,

$$g(y) = y^{\frac{1}{2}}, \qquad \forall \ y > 0.$$

This function is not convex.

(b) Consider the following change of variables, where we set

$$f(x) = \ln(g(y_1, \dots, y_n)), \qquad b_i = \ln \beta_i, \ \forall i, \qquad x_j = \ln y_j, \ \forall j.$$

With this change of variables, f(x) can be written as

$$f(x) = \ln \left(\sum_{i=1}^{m} e^{b_i + a_{i1}x_1 + \dots + a_{in}x_n} \right).$$

Note that f(x) can also be represented as

$$f(x) = \ln \exp(Ax + b), \quad \forall \ x \in \Re^n,$$

where $\ln \exp(z) = \ln(e^{z_1} + \cdots + e^{z_m})$ for all $z \in \Re^m$, A is an $m \times n$ matrix with entries a_{ij} , and $b \in \Re^m$ is a vector with components b_i . Let $f_2(z) = \ln(e^{z_1} + \cdots + e^{z_m})$. This function is convex by Exercise 1.4(b). With this identification, f(x) can be viewed as the composition $f(x) = f_2(Ax + b)$, which is convex by Exercise 1.4(g).

(c) Consider a function $g: \Re^n \mapsto \Re$ of the form

$$g(y) = g_1(y)^{\gamma_1} \cdots g_r(y)^{\gamma_r},$$

where g_k is a posynomial and $\gamma_k > 0$ for all k. Using a change of variables similar to part (b), we see that we can represent the function $f(x) = \ln g(y)$ as

$$f(x) = \sum_{k=1}^{r} \gamma_k \ln \exp(A_k x + b_k),$$

with the matrix A_k and the vector b_k being associated with the posynomial g_k for each k. Since f(x) is the weighted sum of convex functions with nonnegative coefficients [part (b)], it follows that f(x) is convex.

1.7 (Arithmetic-Geometric Mean Inequality)

Show that if $\alpha_1, \ldots, \alpha_n$ are positive scalars with $\sum_{i=1}^n \alpha_i = 1$, then for every set of positive scalars x_1, \ldots, x_n , we have

$$x_1^{\alpha_1} x_2^{\alpha_2} \cdots x_n^{\alpha_n} \le \alpha_1 x_1 + a_2 x_2 + \cdots + \alpha_n x_n$$

with equality if and only if $x_1 = x_2 = \cdots = x_n$. Hint: Show that $(-\ln x)$ is a strictly convex function on $(0, \infty)$.

Solution: Consider the function $f(x) = -\ln(x)$. Since $\nabla^2 f(x) = 1/x^2 > 0$ for all x > 0, the function $-\ln(x)$ is strictly convex over $(0, \infty)$. Therefore, for all positive scalars x_1, \ldots, x_n and $\alpha_1, \ldots, \alpha_n$ with $\sum_{i=1}^n \alpha_i = 1$, we have

$$-\ln(\alpha_1 x_1 + \dots + \alpha_n x_n) \le -\alpha_1 \ln(x_1) - \dots - \alpha_n \ln(x_n),$$

which is equivalent to

$$e^{\ln(\alpha_1 x_1 + \dots + \alpha_n x_n)} \ge e^{\alpha_1 \ln(x_1) + \dots + \alpha_n \ln(x_n)} = e^{\alpha_1 \ln(x_1)} \dots e^{\alpha_n \ln(x_n)},$$

or

$$\alpha_1 x_1 + \dots + \alpha_n x_n \ge x_1^{\alpha_1} \cdots x_n^{\alpha_n},$$

as desired. Since $-\ln(x)$ is strictly convex, the above inequality is satisfied with equality if and only if the scalars x_1, \ldots, x_n are all equal.

1.8 (Young and Holder Inequalities)

Use the result of Exercise 1.7 to verify Young's inequality

$$xy \le \frac{x^p}{p} + \frac{y^q}{q}, \quad \forall \ x \ge 0, \ \forall \ y \ge 0,$$

where p > 0, q > 0, and

$$1/p + 1/q = 1$$
.

Then, use Young's inequality to verify Holder's inequality

$$\sum_{i=1}^{n} |x_i y_i| \le \left(\sum_{i=1}^{n} |x_i|^p\right)^{1/p} \left(\sum_{i=1}^{n} |y_i|^q\right)^{1/q}.$$

Solution: According to Exercise 1.7, we have

$$u^{\frac{1}{p}}v^{\frac{1}{q}} \le \frac{u}{p} + \frac{v}{q}, \qquad \forall \ u > 0, \quad \forall \ v > 0,$$

where 1/p + 1/q = 1, p > 0, and q > 0. The above relation also holds if u = 0 or v = 0. By setting $u = x^p$ and $v = y^q$, we obtain Young's inequality

$$xy \le \frac{x^p}{p} + \frac{y^q}{q}, \quad \forall \ x \ge 0, \quad \forall \ y \ge 0.$$

To show Holder's inequality, note that it holds if $x_1 = \cdots = x_n = 0$ or $y_1 = \cdots = y_n = 0$. If x_1, \ldots, x_n and y_1, \ldots, y_n are such that $(x_1, \ldots, x_n) \neq 0$ and $(y_1, \ldots, y_n) \neq 0$, then by using

$$x = \frac{|x_i|}{\left(\sum_{j=1}^n |x_j|^p\right)^{1/p}}$$
 and $y = \frac{|y_i|}{\left(\sum_{j=1}^n |y_j|^q\right)^{1/q}}$

in Young's inequality, we have for all i = 1, ..., n,

$$\frac{|x_i|}{\left(\sum_{j=1}^n |x_j|^p\right)^{1/p}} \frac{|y_i|}{\left(\sum_{j=1}^n |y_j|^q\right)^{1/q}} \le \frac{|x_i|^p}{p\left(\sum_{j=1}^n |x_j|^p\right)} + \frac{|y_i|^q}{q\left(\sum_{j=1}^n |y_j|^q\right)}.$$

By adding these inequalities over i = 1, ..., n, we obtain

$$\frac{\sum_{i=1}^{n} |x_i| \cdot |y_i|}{\left(\sum_{j=1}^{n} |x_j|^p\right)^{1/p} \left(\sum_{j=1}^{n} |y_j|^q\right)^{1/q}} \le \frac{1}{p} + \frac{1}{q} = 1,$$

which implies Holder's inequality.

1.9 (Characterization of Differentiable Convex Functions)

Let $f: \Re^n \to \Re$ be a differentiable function. Show that f is convex over a nonempty convex set C if and only if

$$(\nabla f(x) - \nabla f(y))'(x - y) \ge 0, \quad \forall x, y \in C.$$

Note: The condition above says that the function f, restricted to the line segment connecting x and y, has monotonically nondecreasing gradient.

Solution: If f is convex, then by Prop. 1.1.7(a), we have

$$f(y) \ge f(x) + \nabla f(x)'(y - x), \quad \forall x, y \in C.$$

By exchanging the roles of x and y in this relation, we obtain

$$f(x) \ge f(y) + \nabla f(y)'(x - y), \quad \forall x, y \in C,$$

and by adding the preceding two inequalities, it follows that

$$\left(\nabla f(y) - \nabla f(x)\right)'(x - y) \ge 0. \tag{1.5}$$

Conversely, let Eq. (1.5) hold, and let x and y be two points in C. Define the function $h: \Re \mapsto \Re$ by

$$h(t) = f(x + t(y - x)).$$

Consider some $t, t' \in [0, 1]$ such that t < t'. By convexity of C, we have that x + t(y - x) and x + t'(y - x) belong to C. Using the chain rule and Eq. (1.5), we have

$$\left(\frac{dh(t')}{dt} - \frac{dh(t)}{dt}\right)(t' - t)
= \left(\nabla f(x + t'(y - x)) - \nabla f(x + t(y - x))\right)'(y - x)(t' - t)
\ge 0.$$

Thus, dh/dt is nondecreasing on [0,1] and for any $t \in (0,1)$, we have

$$\frac{h(t) - h(0)}{t} = \frac{1}{t} \int_0^t \frac{dh(\tau)}{d\tau} d\tau \le h(t) \le \frac{1}{1 - t} \int_t^1 \frac{dh(\tau)}{d\tau} d\tau = \frac{h(1) - h(t)}{1 - t}.$$

Equivalently,

$$th(1) + (1-t)h(0) \ge h(t),$$

and from the definition of h, we obtain

$$tf(y) + (1-t)f(x) \ge f(ty + (1-t)x).$$

Since this inequality has been proved for arbitrary $t \in [0,1]$ and $x,y \in C$, we conclude that f is convex.

1.10 (Strong Convexity)

Let $f: \Re^n \to \Re$ be a function that is continuous over a closed convex set $C \subset \text{dom}(f)$, and let $\sigma > 0$. We say that f is strongly convex over C with coefficient σ if for all $x, y \in C$ and all $\alpha \in [0, 1]$, we have

$$f(\alpha x + (1 - \alpha)y) + \frac{\sigma}{2}\alpha(1 - \alpha)\|x - y\|^2 \le \alpha f(x) + (1 - \alpha)f(y).$$

(a) Show that if f is strongly convex over C with coefficient σ , then f is strictly convex over C. Furthermore, there exists a unique $x^* \in C$ that minimizes f over C, and we have

$$f(x) \ge f(x^*) + \frac{\sigma}{2} ||x - x^*||^2, \quad \forall \ x \in C.$$

- (b) Assume that int(C), the interior of C, is nonempty, and that f is continuously differentiable over int(C). Show that the following are equivalent:
 - (i) f is strongly convex with coefficient σ over C.
 - (ii) We have

$$(\nabla f(x) - \nabla f(y))'(x - y) \ge \sigma ||x - y||^2, \quad \forall x, y \in \text{int}(C).$$

(iii) We have

$$f(y) \ge f(x) + \nabla f(x)'(y-x) + \frac{\sigma}{2} ||x-y||^2, \quad \forall x, y \in \text{int}(C).$$

Furthermore, if f is twice continuously differentiable over int(C), the above three properties are equivalent to:

(iv) The matrix $\nabla^2 f(x) - \sigma I$ is positive semidefinite for every $x \in \text{int}(C)$, where I is the identity matrix.

Solution: (a) The strict convexity of f over C is evident from the definition of strong convexity and the hypothesis. Strict convexity also implies that there can be at most one minimum of f over C.

To show existence of a vector x^* that minimizes f over C, we show that every level set $\{x \in C \mid f(x) \leq \gamma\}$ is bounded and hence compact (since C is closed and f is continuous over C), and then use Weierstrass' Theorem. Assume to arrive at a contradiction that a level set $L = \{x \in C \mid f(x) \leq \gamma\}$ is unbounded, and let $\{x_k\} \subset L$ be an unbounded sequence. We assume with no loss of generality that $||x_k - x_0|| \geq 1$ for all k. Let $\alpha_k = 1/||x_k - x_0||$, and note that $\alpha_k \to 0$. Define

$$y_k = \alpha_k x_k + (1 - \alpha_k) x_0 = \frac{x_k - x_0}{\|x_k - x_0\|} + x_0,$$

and note that $||y_k - x_0|| = 1$ for all $k \ge 1$. By strong convexity of f, we have

$$f(y_k) \le \alpha_k f(x_k) + (1 - \alpha_k) f(x_0) - \frac{1}{2} \sigma \alpha_k (1 - \alpha_k) \|x_k - x_0\|^2$$

$$= \alpha_k f(x_k) + (1 - \alpha_k) f(x_0) - \frac{1}{2} \sigma (1 - \alpha_k) \|x_k - x_0\|$$

$$\le \gamma - \frac{1}{2} \sigma (1 - \alpha_k) \|x_k - x_0\|.$$

Hence $f(y_k) \to -\infty$, which contradicts the boundedness of $\{y_k\}$ and the continuity of f.

To show the inequality $f(x) \ge f(x^*) + (\sigma/2)||x - x^*||^2$, we write for any $x \in C$ and $\alpha \in (0,1)$,

$$\alpha f(x) + (1 - \alpha)f(x^*) \ge f(\alpha x + (1 - \alpha)x^*) + \frac{1}{2}\sigma\alpha(1 - \alpha)\|x - x^*\|^2$$
$$\ge f(x^*) + \frac{1}{2}\sigma\alpha(1 - \alpha)\|x - x^*\|^2.$$

It follows that $f(x) \ge f(x^*) + (\sigma/2)(1-\alpha)||x-x^*||^2$, and by taking the limit as $\alpha \to 0$, we obtain the desired inequality.

(b) We first show that (i) implies (ii). We have, using the definition of strong convexity,

$$f(y) + \alpha \nabla f(y)'(x-y) \leq f\left(y + \alpha(x-y)\right) \leq \alpha f(x) + (1-\alpha)f(y) - \frac{\sigma}{2}\alpha(1-\alpha)\|x-y\|^2,$$

for all $x, y \in \text{int}(C)$ and $\alpha \in (0, 1)$, from which

$$f(y) + \nabla f(y)'(x - y) \le f(x) - \frac{\sigma}{2}(1 - \alpha)||x - y||^2.$$

Similarly,

$$f(x) + \nabla f(x)'(y - x) \le f(y) - \frac{\sigma}{2}(1 - \alpha)||x - y||^2,$$
 (1.6)

and adding these two inequalities:

$$(\nabla f(y) - \nabla f(x))'(x - y) \le -\sigma(1 - \alpha)||x - y||^2,$$

or

$$(\nabla f(y) - \nabla f(x))'(y - x) \ge \sigma(1 - \alpha) ||x - y||^2.$$

Taking the limit as $\alpha \to 0$, we obtain

$$(\nabla f(y) - \nabla f(x))'(y - x) \ge \sigma ||x - y||^2.$$

Next we show that (ii) implies (i). For any $\alpha \in (0,1)$ and $x_1, x_2 \in \text{int}(C)$ with $x_1 \neq x_2$, let

$$x_{\alpha} = \alpha x_1 + (1 - \alpha)x_2.$$

We have

$$f(x_a) = f(x_1) + \int_0^1 \nabla f(x_1 + t(x_\alpha - x_1))'(x_a - x_1) dt,$$

$$f(x_a) = f(x_2) + \int_0^1 \nabla f(x_2 + t(x_a - x_2))'(x_a - x_2) dt.$$

Multiplying these relations with α and $1 - \alpha$, respectively, adding, and collecting terms using the relations $x_{\alpha} - x_1 = (1 - \alpha)(x_2 - x_1)$, $x_{\alpha} - x_2 = \alpha(x_1 - x_2)$, and

$$(x_1 + t(x_a - x_1)) - (x_2 + t(x_a - x_2)) = (1 - t)(x_1 - x_2),$$

we obtain

$$\alpha f(x_1) + (1 - \alpha)f(x_2) - f(x_a)$$

$$= \alpha (1 - \alpha) \int_0^1 \left(\nabla f(x_1 + t(x_a - x_1)) - \nabla f(x_2 + t(x_a - x_2)) \right)'(x_1 - x_2) dt$$

$$\geq \sigma \alpha (1 - \alpha) ||x_1 - x_2||^2 \int_0^1 (1 - t) dt$$

$$= \frac{1}{2} \sigma \alpha (1 - \alpha) ||x_1 - x_2||^2,$$

verifying the strong convexity inequality for x_1, x_2 in the interior of C [and using the continuity of f, for x_1, x_2 in the boundary of C as well].

Next we show that (iii) is equivalent to (i) and (ii). Indeed, by taking the limit in Eqs. (1.6) as $\alpha \to 0$, we see that (i) implies (iii). Conversely if (iii) holds, we have

$$f(y) \ge f(x) + \nabla f(x)'(y - x) + \frac{\sigma}{2} ||x - y||^2, \quad \forall x, y \in \text{int}(C),$$

and

$$f(x) \geq f(y) + \nabla f(y)'(x-y) + \frac{\sigma}{2} \|x-y\|^2, \qquad \forall \; x,y \in \mathrm{int}(C).$$

By adding these two relations, we obtain (ii).

Assume now that f is twice continuously differentiable over $\operatorname{int}(C)$. First we show that (iv) implies (ii). Let $x,y\in\operatorname{int}(C)$ and consider the function $g:\Re\mapsto\Re$ defined by

$$g(t) = \nabla f(tx + (1-t)y)'(x-y).$$

Using the Mean Value Theorem, we have

$$\left(\nabla f(x) - \nabla f(y)\right)'(x - y) = g(1) - g(0) = \frac{dg(t)}{dt}$$

for some $t \in [0,1]$. On the other hand,

$$\frac{dg(t)}{dt} = (x-y)'\nabla^2 f(tx + (1-t)y)(x-y) \ge \sigma ||x-y||^2,$$

where the last inequality holds because the matrix

$$\nabla^2 f(tx + (1-t)y) - \sigma I$$

is positive semidefinite. Combining the last two relations, we obtain the desired inequality.

We finally show that (i) implies (iv). For any $\alpha \in (0,1)$ and $x_1, x_2 \in \text{int}(C)$ with $x_1 \neq x_2$, let

$$x_{\alpha} = \alpha x_1 + (1 - \alpha)x_2.$$

Using the 2nd order Mean Value Theorem, we have

$$f(x_1) = f(x_\alpha) + \nabla f(x_\alpha)'(x_1 - x_\alpha) + \frac{1}{2}(x_1 - x_\alpha)'\nabla^2 f(\tilde{x}_\alpha)(x_1 - x_\alpha),$$

$$f(x_2) = f(x_\alpha) + \nabla f(x_\alpha)'(x_2 - x_\alpha) + \frac{1}{2}(x_2 - x_\alpha)'\nabla^2 f(\hat{x}_\alpha)(x_2 - x_\alpha),$$

where \tilde{x}_{α} and \hat{x}_{α} are vectors that lie in the intervals connecting x_{α} with x_1 and x_2 , respectively. Multiplying these relations with α and $1-\alpha$, respectively, adding, canceling the terms involving $\nabla f(x_{\alpha})$, and using the relations $x_{\alpha} - x_1 = (1-\alpha)(x_2-x_1)$ and $x_{\alpha} - x_2 = \alpha(x_1-x_2)$ and the definition of strong convexity, we obtain

$$f(x_{\alpha}) + \frac{1}{2}\sigma\alpha(1-\alpha)\|x_1 - x_2\|^2 \le \alpha f(x_1) + (1-\alpha)f(x_2)$$

$$= f(x_{\alpha}) + \frac{1}{2}\alpha(1-\alpha)^2(x_1 - x_2)'\nabla^2 f(\tilde{x}_{\alpha})(x_1 - x_2)$$

$$+ \frac{1}{2}\alpha^2(1-\alpha)(x_1 - x_2)'\nabla^2 f(\hat{x}_{\alpha})(x_1 - x_2)$$

and finally,

$$\sigma \|x_1 - x_2\|^2 \le (x_1 - x_2)' \big((1 - \alpha) \nabla^2 f(\tilde{x}_\alpha) + \alpha \nabla^2 f(\hat{x}_\alpha) \big) (x_1 - x_2).$$

Dividing by $||x_1 - x_2||^2$ and letting x_2 approach x_1 , we obtain

$$\sigma \leq d' \nabla^2 f(x_1) d$$

where

$$d = \frac{(x_1 - x_2)}{\|x_1 - x_2\|}.$$

Since x_1 and x_2 were chosen arbitrarily within $\operatorname{int}(C)$, it follows that the matrix $\nabla^2 f(x) - \sigma I$ is positive semidefinite for every $x \in \operatorname{int}(C)$. Since by the convexity of C, every point in the boundary of C can be approached from the interior and $\nabla^2 f$ is continuous, $\nabla^2 f(x) - \sigma I$ is also positive semidefinite for every x in the boundary of C.

SECTION 1.2: Convex and Affine Hulls

1.11 (Characterization of Convex Hulls and Affine Hulls)

Let X be a nonempty subset of \Re^n .

(a) Show that the convex hull of X coincides with the set of all convex combinations of its elements, i.e.,

$$\operatorname{conv}(X) = \left\{ \sum_{i \in I} \alpha_i x_i \mid I \text{: a finite set, } \sum_{i \in I} \alpha_i = 1, \, a_i \ge 0, \, x_i \in X, \, \forall \, i \in I \right\}.$$

(b) Show that the affine hull of X coincides with the set of all linear combinations of its elements with coefficients adding to 1, i.e.,

$$\operatorname{aff}(X) = \left\{ \sum_{i \in I} \alpha_i x_i \mid I: \text{ a finite set, } \sum_{i \in I} \alpha_i = 1, \, x_i \in X, \, \forall \, i \in I \right\}.$$

Show also that if m is the dimension of $\operatorname{aff}(X)$, there exist vectors $\overline{x}, \overline{x}_1, \ldots, \overline{x}_m$ from X such that $\overline{x}_1 - \overline{x}, \ldots, \overline{x}_m - \overline{x}$ form a basis for the subspace that is parallel to $\operatorname{aff}(X)$.

Solution: (a) The elements of X belong to $\operatorname{conv}(X)$, so all their convex combinations belong to $\operatorname{conv}(X)$ since $\operatorname{conv}(X)$ is a convex set. On the other hand, consider any two convex combinations of elements of X, $x = \lambda_1 x_1 + \cdots + \lambda_m x_m$ and $y = \mu_1 y_1 + \cdots + \mu_r y_r$, where $x_i \in X$ and $y_j \in X$. The vector

$$(1-\alpha)x + \alpha y = (1-\alpha)(\lambda_1 x_1 + \dots + \lambda_m x_m) + \alpha(\mu_1 y_1 + \dots + \mu_r y_r),$$

where $0 \le \alpha \le 1$, is another convex combination of elements of X.

Thus, the set of convex combinations of elements of X is itself a convex set, which contains X, and is contained in conv(X). Hence it must coincide with conv(X), which by definition is the intersection of all convex sets containing X.

(b) The set

$$A = \left\{ \sum_{i \in I} \alpha_i x_i \mid I \text{ is a finite set, } \sum_{i \in I} \alpha_i = 1, \, x_i \in X, \, \forall \, i \in I \right\}$$

contains every line that passes through any pair of its points, so it is affine. Since it also contains X, it must contain aff(X).

To show the reverse inclusion, we note that

$$A = \overline{x} + S$$
,

where \overline{x} is some vector in X and S is a subspace that must have the form

$$S = \left\{ \sum_{i \in I} \alpha_i(x_i - \overline{x}) \mid I \text{ is a finite set, } \sum_{i \in I} \alpha_i = 1, x_i \in X, \forall i \in I \right\}.$$

By taking one of the vectors x_i to be \overline{x} , we see that

$$S = \left\{ \sum_{i \in I} \beta_i(x_i - \overline{x}) \mid I \text{ is a finite set, } \beta_i \in \Re, \ x_i \in X, \ \forall \ i \in I \right\}.$$

It follows that S has a basis of the form $\overline{x}_1 - \overline{x}, \dots, \overline{x}_m - \overline{x}$, where $\overline{x}_1, \dots, \overline{x}_m \in X$, and m is the dimension of S. The subspace that is parallel to an affine set that contains X must contain the basis $\overline{x}_1 - \overline{x}, \dots, \overline{x}_m - \overline{x}$. Hence any affine set that contains X, including aff(X), must contain $A = \overline{x} + S$. The proof of aff(X) = A is complete.

1.12

Let $\{C_i \mid i \in I\}$ be an arbitrary collection of convex sets in \Re^n , and let C be the convex hull of the union of the collection. Show that

$$C = \bigcup_{\overline{I} \subset I, \ \overline{I}: \text{ finite set }} \left\{ \sum_{i \in \overline{I}} \alpha_i C_i \ \middle| \ \sum_{i \in \overline{I}} \alpha_i = 1, \ \alpha_i \geq 0, \ \forall \ i \in \overline{I} \right\},$$

i.e., the convex hull of the union of the sets C_i is equal to the set of all convex combinations of vectors that come from different sets C_i .

Solution: By Exercise 1.11, C is the set of all convex combinations $x = \alpha_1 y_1 + \cdots + \alpha_m y_m$, where m is a positive integer, and the vectors y_1, \ldots, y_m belong to the union of the sets C_i . Actually, we can get C just by taking those combinations in which the vectors are taken from different sets C_i .

Indeed, if two of the vectors, y_1 and y_2 belong to the same C_i , then the term $\alpha_1 y_1 + \alpha_2 y_2$ can be replaced by αy , where $\alpha = \alpha_1 + \alpha_2$ and

$$y = (\alpha_1/\alpha)y_1 + (\alpha_2/\alpha)y_2 \in C_i.$$

Thus, C is the union of the vector sums of the form

$$\alpha_1 C_{i_1} + \cdots + \alpha_m C_{i_m}$$

with

$$\alpha_i \ge 0, \ \forall \ i = 1, \dots, m, \quad \sum_{i=1}^{m} \alpha_i = 1,$$

and the indices i_1, \ldots, i_m are all different, proving our claim.

1.13 (Generated Cones and Convex Hulls I)

Show that:

(a) For a nonempty convex subset C of \Re^n , we have

$$cone(C) = \bigcup_{x \in C} \{ \gamma x \mid \gamma \ge 0 \}.$$

- (b) A cone C is convex if and only if $C + C \subset C$.
- (c) For any two convex cones C_1 and C_2 containing the origin, we have

$$C_1 + C_2 = \text{conv}(C_1 \cup C_2), \qquad C_1 \cap C_2 = \bigcup_{\alpha \in [0,1]} (\alpha C_1 \cap (1-\alpha)C_2).$$

Solution: (a) Let $y \in \text{cone}(C)$. If y = 0, then $y \in \bigcup_{x \in C} \{\gamma x \mid \gamma \geq 0\}$. If $y \neq 0$, then by definition of cone(C), we have

$$y = \sum_{i=1}^{m} \lambda_i x_i,$$

for some positive integer m, nonnegative scalars λ_i , and vectors $x_i \in C$. Since $y \neq 0$, we cannot have all λ_i equal to zero, implying that $\sum_{i=1}^m \lambda_i > 0$. Because $x_i \in C$ for all i and C is convex, the vector

$$x = \sum_{i=1}^{m} \frac{\lambda_i}{\sum_{j=1}^{m} \lambda_j} x_i$$

belongs to C. For this vector, we have

$$y = \left(\sum_{i=1}^{m} \lambda_i\right) x,$$

with $\sum_{i=1}^{m} \lambda_i > 0$, implying that $y \in \bigcup_{x \in C} \{ \gamma x \mid \gamma \geq 0 \}$ and showing that

$$cone(C) \subset \bigcup_{x \in C} \{ \gamma x \mid \gamma \ge 0 \}.$$

The reverse inclusion follows from the definition of cone(C).

- (b) Let C be a cone such that $C+C\subset C$, and let $x,y\in C$ and $\alpha\in[0,1]$. Then since C is a cone, $\alpha x\in C$ and $(1-\alpha)y\in C$, so that $\alpha x+(1-\alpha)y\in C+C\subset C$, showing that C is convex. Conversely, let C be a convex cone and let $x,y\in C$. Then, since C is a cone, $2x\in C$ and $2y\in C$, so that by the convexity of C, $x+y=\frac{1}{2}(2x+2y)\in C$, showing that $C+C\subset C$.
- (c) First we prove that $C_1 + C_2 \subset \text{conv}(C_1 \cup C_2)$. Choose any $x \in C_1 + C_2$. Since $C_1 + C_2$ is a cone [see Exercise 1.2(c)], the vector 2x is in $C_1 + C_2$, so that $2x = x_1 + x_2$ for some $x_1 \in C_1$ and $x_2 \in C_2$. Therefore,

$$x = \frac{1}{2}x_1 + \frac{1}{2}x_2,$$

showing that $x \in \text{conv}(C_1 \cup C_2)$.

Next, we show that $\operatorname{conv}(C_1 \cup C_2) \subset C_1 + C_2$. Since $0 \in C_1$ and $0 \in C_2$, it follows that

$$C_i = C_i + 0 \subset C_1 + C_2, \qquad i = 1, 2,$$

implying that

$$C_1 \cup C_2 \subset C_1 + C_2$$
.

By taking the convex hull of both sides in the above inclusion and by using the convexity of $C_1 + C_2$, we obtain

$$conv(C_1 \cup C_2) \subset conv(C_1 + C_2) = C_1 + C_2.$$

We finally show that

$$C_1 \cap C_2 = \bigcup_{\alpha \in [0,1]} (\alpha C_1 \cap (1-\alpha)C_2).$$

We claim that for all α with $0 < \alpha < 1$, we have

$$\alpha C_1 \cap (1 - \alpha)C_2 = C_1 \cap C_2.$$

Indeed, if $x \in C_1 \cap C_2$, it follows that $x \in C_1$ and $x \in C_2$. Since C_1 and C_2 are cones and $0 < \alpha < 1$, we have $x \in \alpha C_1$ and $x \in (1 - \alpha)C_2$. Conversely, if $x \in \alpha C_1 \cap (1 - \alpha)C_2$, we have

$$\frac{x}{\alpha} \in C_1$$
,

and

$$\frac{x}{(1-\alpha)} \in C_2.$$

Since C_1 and C_2 are cones, it follows that $x \in C_1$ and $x \in C_2$, so that $x \in C_1 \cap C_2$.

If $\alpha = 0$ or $\alpha = 1$, we obtain

$$\alpha C_1 \cap (1 - \alpha)C_2 = \{0\} \subset C_1 \cap C_2,$$

since C_1 and C_2 contain the origin. Thus, the result follows.

1.14 (Generated Cones and Convex Hulls II)

Let X be a nonempty set. Show that:

- (a) X, conv(X), and cl(X) have the same affine hull.
- (b) cone(X) = cone(conv(X)).
- (c) $\operatorname{aff}(\operatorname{conv}(X)) \subset \operatorname{aff}(\operatorname{cone}(X))$. Give an example where the inclusion is strict.
- (d) If the origin belongs to conv(X), then aff(conv(X)) = aff(cone(X)).
- (e) If A is a matrix, $A \operatorname{conv}(X) = \operatorname{conv}(AX)$.

Solution: (a) We first show that X and cl(X) have the same affine hull. Since $X \subset cl(X)$, there holds

$$\operatorname{aff}(X) \subset \operatorname{aff}(\operatorname{cl}(X)).$$

Conversely, because $X \subset \operatorname{aff}(X)$ and $\operatorname{aff}(X)$ is closed, we have $\operatorname{cl}(X) \subset \operatorname{aff}(X)$, implying that

$$\operatorname{aff}(\operatorname{cl}(X)) \subset \operatorname{aff}(X).$$

We now show that X and $\operatorname{conv}(X)$ have the same affine hull. By using a translation argument if necessary, we assume without loss of generality that X contains the origin, so that both $\operatorname{aff}(X)$ and $\operatorname{aff}(\operatorname{conv}(X))$ are subspaces. Since $X \subset \operatorname{conv}(X)$, evidently $\operatorname{aff}(X) \subset \operatorname{aff}(\operatorname{conv}(X))$. To show the reverse inclusion, let the dimension of $\operatorname{aff}(\operatorname{conv}(X))$ be m, and let x_1, \ldots, x_m be linearly independent vectors in $\operatorname{conv}(X)$ that span $\operatorname{aff}(\operatorname{conv}(X))$. Then every $x \in \operatorname{aff}(\operatorname{conv}(X))$ is a linear combination of the vectors x_1, \ldots, x_m , i.e., there exist scalars β_1, \ldots, β_m such that

$$x = \sum_{i=1}^{m} \beta_i x_i.$$

By the definition of convex hull, each x_i is a convex combination of vectors in X, so that x is a linear combination of vectors in X, implying that $x \in \text{aff}(X)$. Hence, $\text{aff}(\text{conv}(X)) \subset \text{aff}(X)$.

(b) Since $X \subset \text{conv}(X)$, clearly $\text{cone}(X) \subset \text{cone}(\text{conv}(X))$. Conversely, let $x \in \text{cone}(\text{conv}(X))$. Then x is a nonnegative combination of some vectors in conv(X), i.e., for some positive integer p, vectors $x_1, \ldots, x_p \in \text{conv}(X)$, and nonnegative scalars $\alpha_1, \ldots, \alpha_p$, we have

$$x = \sum_{i=1}^{p} \alpha_i x_i.$$

Each x_i is a convex combination of some vectors in X, so that x is a nonnegative combination of some vectors in X, implying that $x \in \text{cone}(X)$. Hence $\text{cone}(\text{conv}(X)) \subset \text{cone}(X)$.

(c) Since $\operatorname{conv}(X)$ is the set of all convex combinations of vectors in X, and $\operatorname{cone}(X)$ is the set of all nonnegative combinations of vectors in X, it follows that $\operatorname{conv}(X) \subset \operatorname{cone}(X)$. Therefore

$$\operatorname{aff}(\operatorname{conv}(X)) \subset \operatorname{aff}(\operatorname{cone}(X)).$$

For an example showing that the above inclusion can be strict, consider the set $X = \{(1,1)\}$ in \Re^2 . Then $\operatorname{conv}(X) = X$, so that

$$aff(conv(X)) = X = \{(1,1)\},\$$

and the dimension of conv(X) is zero. On the other hand, $cone(X) = \{(\alpha, \alpha) \mid \alpha \geq 0\}$, so that

$$aff(cone(X)) = \{(x_1, x_2) \mid x_1 = x_2\},\$$

and the dimension of cone(X) is one.

(d) In view of parts (a) and (c), it suffices to show that

$$\operatorname{aff}(\operatorname{cone}(X)) \subset \operatorname{aff}(X).$$

It is always true that $0 \in \operatorname{cone}(X)$, so $\operatorname{aff}(\operatorname{cone}(X))$ is a subspace. Let the dimension of $\operatorname{aff}(\operatorname{cone}(X))$ be m, and let x_1,\ldots,x_m be linearly independent vectors in $\operatorname{cone}(X)$ that span $\operatorname{aff}(\operatorname{cone}(X))$ [cf. Exercise 1.11(b)]. Since every vector in $\operatorname{aff}(\operatorname{cone}(X))$ is a linear combination of x_1,\ldots,x_m , and since each x_i is a nonnegative combination of some vectors in X, it follows that every vector in $\operatorname{aff}(\operatorname{cone}(X))$ is a linear combination of some vectors in X. In view of the assumption that $0 \in \operatorname{conv}(X)$, the affine hull of $\operatorname{conv}(X)$ is a subspace, which implies by part (a) that the affine hull of X is a subspace. Hence, $\operatorname{aff}(X)$ is the set of linear combinations of vectors from X. It follows that every vector in $\operatorname{aff}(\operatorname{cone}(X))$ belongs to $\operatorname{aff}(X)$, showing that $\operatorname{aff}(\operatorname{cone}(X)) \subset \operatorname{aff}(X)$.

(e) If $y \in \text{conv}(AX)$, then for some $x_1, x_2 \in X$,

$$y = \alpha A x_1 + (1 - \alpha) A x_2 = A(\alpha x_1 + (1 - \alpha) x_2) \in A \operatorname{conv}(X).$$

Hence $\operatorname{conv}(AX) \subset A\operatorname{conv}(X)$.

Conversely, if $y \in A \operatorname{conv}(X)$, then for some $x_1, x_2 \in X$,

$$y = A(\alpha x_1 + (1 - \alpha)x_2) = A(\alpha x_1 + (1 - \alpha)x_2) \in \text{conv}(AX).$$

Hence $A \operatorname{conv}(X) \subset \operatorname{conv}(AX)$.

1.15

Let $\{f_i \mid i \in I\}$ be an arbitrary collection of proper convex functions $f_i : \mathbb{R}^n \mapsto (-\infty, \infty]$. Define

$$f(x) = \inf \{ w \mid (x, w) \in \operatorname{conv}(\bigcup_{i \in I} \operatorname{epi}(f_i)) \}, \quad x \in \Re^n$$

Show that f(x) is given by

$$f(x) = \inf \left\{ \sum_{i \in \overline{I}} \alpha_i f_i(x_i) \mid \sum_{i \in \overline{I}} \alpha_i x_i = x, \ x_i \in \Re^n, \ \sum_{i \in \overline{I}} \alpha_i = 1, \ \alpha_i \ge 0, \ \forall \ i \in \overline{I}, \right.$$
$$\overline{I} \subset I, \ \overline{I} : \ \text{finite} \right\}.$$

Solution: By definition, f(x) is the infimum of the values of w such that $(x, w) \in C$, where C is the convex hull of the union of nonempty convex sets $\operatorname{epi}(f_i)$. By Exercise 1.12, $(x, w) \in C$ if and only if (x, w) can be expressed as a convex combination of the form

$$(x, w) = \sum_{i \in \overline{I}} \alpha_i(x_i, w_i) = \left(\sum_{i \in \overline{I}} \alpha_i x_i, \sum_{i \in \overline{I}} \alpha_i w_i\right),$$

where $\overline{I} \subset I$ is a finite set and $(x_i, w_i) \in \operatorname{epi}(f_i)$ for all $i \in \overline{I}$. Thus, f(x) can be expressed as

$$f(x) = \inf \left\{ \sum_{i \in \overline{I}} \alpha_i w_i \mid (x, w) = \sum_{i \in \overline{I}} \alpha_i (x_i, w_i), \right.$$
$$(x_i, w_i) \in \operatorname{epi}(f_i), \ \alpha_i \ge 0, \ \forall \ i \in \overline{I}, \ \sum_{i \in \overline{I}} \alpha_i = 1 \right\}.$$

Since the set $\{(x_i, f_i(x_i)) \mid x_i \in \Re^n\}$ is contained in $\operatorname{epi}(f_i)$, we obtain

$$f(x) \le \inf \left\{ \sum_{i \in \overline{I}} \alpha_i f_i(x_i) \mid x = \sum_{i \in \overline{I}} \alpha_i x_i, \ x_i \in \Re^n, \ \alpha_i \ge 0, \ \forall \ i \in \overline{I}, \ \sum_{i \in \overline{I}} \alpha_i = 1 \right\}.$$

On the other hand, by the definition of $\operatorname{epi}(f_i)$, for each $(x_i, w_i) \in \operatorname{epi}(f_i)$ we have $w_i \geq f_i(x_i)$, implying that

$$f(x) \ge \inf \left\{ \sum_{i \in \overline{I}} \alpha_i f_i(x_i) \mid x = \sum_{i \in \overline{I}} \alpha_i x_i, \ x_i \in \Re^n, \ \alpha_i \ge 0, \ \forall \ i \in \overline{I}, \ \sum_{i \in \overline{I}} \alpha_i = 1 \right\}.$$

By combining the last two relations, we obtain

$$f(x) = \inf \left\{ \sum_{i \in \overline{I}} \alpha_i f_i(x_i) \mid x = \sum_{i \in \overline{I}} \alpha_i x_i, \ x_i \in \Re^n, \ \alpha_i \ge 0, \ \forall \ i \in \overline{I}, \ \sum_{i \in \overline{I}} \alpha_i = 1 \right\},$$

where the infimum is taken over all representations of x as a convex combination of elements x_i such that only finitely many coefficients α_i are nonzero.

1.16 (Minimization of Linear Functions)

Show that minimization of a linear function over a set is equivalent to minimization over its convex hull, i.e.,

$$\inf_{x \in \text{conv}(X)} c'x = \inf_{x \in X} c'x,$$

if $X \subset \Re^n$ and $c \in \Re^n$. Furthermore, the infimum in the left-hand side above is attained if and only if the infimum in the right-hand side is attained.

Solution: Since $X \subset \text{conv}(X)$, we have

$$\inf_{x \in \text{conv}(X)} c'x \le \inf_{x \in X} c'x. \tag{1.7}$$

Also, any $\overline{x} \in \text{conv}(X)$ can be written as $\overline{x} = \sum_{i=1}^{m} \alpha_i x_i$, for some $x_1, \ldots, x_m \in X$ and some scalars $\alpha_1, \ldots, \alpha_m \geq 0$ with $\sum_{i=1}^{m} \alpha_i = 1$. Hence, since $c'x_i \geq \inf_{x \in X} c'x$, we have

$$c'\overline{x} = \sum_{i=1}^{m} \alpha_i c' x_i \ge \left(\sum_{i=1}^{m} \alpha_i\right) \inf_{x \in X} c' x = \inf_{x \in X} c' x, \qquad \forall \ \overline{x} \in \text{conv}(X).$$

Taking the infimum of the left-hand side over $\overline{x} \in \text{conv}(X)$,

$$\inf_{\overline{x} \in \text{conv}(X)} c' \overline{x} \ge \inf_{x \in X} c' x. \tag{1.8}$$

Combining Eqs. (1.7) and (1.8), we obtain

$$\inf_{x \in \text{conv}(X)} c'x = \inf_{x \in X} c'x.$$

Since $X \subset \operatorname{conv}(X)$ and $\inf_{x \in \operatorname{conv}(X)} c'x = \inf_{x \in X} c'x$, every point that attains the infimum of c'x over X, attains the infimum of c'x over $\operatorname{conv}(X)$. For the converse, assume that the infimum of c'x over $\operatorname{conv}(X)$ is attained at some $\overline{x} \in \operatorname{conv}(X)$. Then, $\overline{x} = \sum_{i=1}^m \alpha_i x_i$, for some $x_1, \ldots, x_m \in X$ and some scalars $\alpha_1, \ldots, \alpha_m \geq 0$ with $\sum_{i=1}^m \alpha_i = 1$, and we have

$$\inf_{x \in X} c'x = \left(\sum_{i=1}^m \alpha_i\right) \inf_{x \in X} c'x \le \sum_{i=1}^m \alpha_i c'x_i = c'\overline{x} = \inf_{x \in \operatorname{conv}(X)} c'x = \inf_{x \in X} c'x.$$

Since the left-hand and right-hand sides are equal, it follows that equality holds throughout above, which can happen only if $c'x_i = \inf_{x \in X} c'x$ for all i with $\alpha_i > 0$. Thus the infimum of c'x over X is attained.

1.17 (Extension of Caratheodory's Theorem)

Let X_1 and X_2 be nonempty subsets of \Re^n , and let $X = \operatorname{conv}(X_1) + \operatorname{cone}(X_2)$. Show that every vector x in X can be represented in the form

$$x = \sum_{i=1}^{k} \alpha_i x_i + \sum_{i=k+1}^{m} \alpha_i y_i,$$

where m is a positive integer with $m \leq n+1$, the vectors x_1, \ldots, x_k belong to X_1 , the vectors y_{k+1}, \ldots, y_m belong to X_2 , and the scalars $\alpha_1, \ldots, \alpha_m$ are nonnegative with $\alpha_1 + \cdots + \alpha_k = 1$. Furthermore, the vectors $x_2 - x_1, \ldots, x_k - x_1, y_{k+1}, \ldots, y_m$ are linearly independent.

Solution: The proof will be an application of Caratheodory's Theorem [Prop. 1.2.1(a)] to the subset of \Re^{n+1} given by

$$Y = \{(x,1) \mid x \in X_1\} \cup \{(y,0) \mid y \in X_2\}.$$

If $x \in X$, then

$$x = \sum_{i=1}^{k} \gamma_i x_i + \sum_{i=k+1}^{m} \gamma_i y_i,$$

where the vectors x_1, \ldots, x_k belong to X_1 , the vectors y_{k+1}, \ldots, y_m belong to X_2 , and the scalars $\gamma_1, \ldots, \gamma_m$ are nonnegative with $\gamma_1 + \cdots + \gamma_k = 1$. Equivalently, $(x,1) \in \text{cone}(Y)$. By Caratheodory's Theorem part (a), we have that

$$(x,1) = \sum_{i=1}^{k} \alpha_i(x_i,1) + \sum_{i=k+1}^{m} \alpha_i(y_i,0),$$

for some positive scalars $\alpha_1, \ldots, \alpha_m$ and vectors

$$(x_1,1),\ldots(x_k,1), (y_{k+1},0),\ldots, (y_m,0),$$

which are linearly independent (implying that $m \leq n+1$) or equivalently,

$$x = \sum_{i=1}^{k} \alpha_i x_i + \sum_{i=k+1}^{m} \alpha_i y_i, \qquad 1 = \sum_{i=1}^{k} \alpha_i.$$

Finally, to show that the vectors $x_2 - x_1, \ldots, x_k - x_1, y_{k+1}, \ldots, y_m$ are linearly independent, assume to arrive at a contradiction, that there exist $\lambda_2, \ldots, \lambda_m$, not all 0, such that

$$\sum_{i=2}^{k} \lambda_i (x_i - x_1) + \sum_{i=k+1}^{m} \lambda_i y_i = 0.$$

Equivalently, defining $\lambda_1 = -(\lambda_2 + \cdots + \lambda_m)$, we have

$$\sum_{i=1}^{k} \lambda_i(x_i, 1) + \sum_{i=k+1}^{m} \lambda_i(y_i, 0) = 0,$$

which contradicts the linear independence of the vectors

$$(x_1,1),\ldots,(x_k,1),(y_{k+1},0),\ldots,(y_m,0)$$

1.18

Let X be a nonempty bounded subset of \Re^n . Show that

$$\operatorname{cl}(\operatorname{conv}(X)) = \operatorname{conv}(\operatorname{cl}(X)).$$

In particular, if X is compact, then conv(X) is compact (cf. Prop. 1.2.2).

Solution: The set cl(X) is compact since X is bounded by assumption. Hence, by Prop. 1.2.2, its convex hull, conv(cl(X)), is compact, and it follows that

$$\operatorname{cl}(\operatorname{conv}(X)) \subset \operatorname{cl}(\operatorname{conv}(\operatorname{cl}(X))) = \operatorname{conv}(\operatorname{cl}(X)).$$

It is also true that

$$\operatorname{conv}(\operatorname{cl}(X)) \subset \operatorname{conv}(\operatorname{cl}(\operatorname{conv}(X))) = \operatorname{cl}(\operatorname{conv}(X)),$$

since by Prop. 1.1.1(d), the closure of a convex set is convex. Hence, the result follows.

1.19 (Convex Hulls and Generated Cones of Cartesian Products)

Given nonempty sets $X_i \subset \Re^{n_i}$, i = 1, ..., m, let $X = X_1 \times \cdots \times X_m$ be their Cartesian product. Show that:

- (a) The convex hull (closure, affine hull) of X is equal to the Cartesian product of the convex hulls (closures, affine hulls, respectively) of the X_i .
- (b) If all the sets X_1, \ldots, X_m contain the origin, then

$$cone(X) = cone(X_1) \times \cdots \times cone(X_m).$$

Furthermore, the result fails if one of the sets does not contain the origin.

Solution: (a) We first show that the convex hull of X is equal to the Cartesian product of the convex hulls of the sets X_i , i = 1, ..., m. Let y be a vector that belongs to conv(X). Then, by definition, for some k, we have

$$y = \sum_{i=1}^{k} \alpha_i y_i$$
, with $\alpha_i \ge 0$, $i = 1, ..., m$, $\sum_{i=1}^{k} \alpha_i = 1$,

where $y_i \in X$ for all i. Since $y_i \in X$, we have that $y_i = (x_1^i, \dots, x_m^i)$ for all i, with $x_1^i \in X_1, \dots, x_m^i \in X_m$. It follows that

$$y = \sum_{i=1}^{k} \alpha_i(x_1^i, \dots, x_m^i) = \left(\sum_{i=1}^{k} \alpha_i x_1^i, \dots, \sum_{i=1}^{k} \alpha_i x_m^i\right),$$

thereby implying that $y \in \text{conv}(X_1) \times \cdots \times \text{conv}(X_m)$.

To prove the reverse inclusion, assume that y is a vector in $conv(X_1) \times \cdots \times conv(X_m)$. Then, we can represent y as $y = (y_1, \dots, y_m)$ with $y_i \in conv(X_i)$, i.e., for all $i = 1, \dots, m$, we have

$$y_i = \sum_{j=1}^{k_i} \alpha_j^i x_j^i, \qquad x_j^i \in X_i, \ \forall j, \qquad \alpha_j^i \ge 0, \ \forall j, \quad \sum_{j=1}^{k_i} \alpha_j^i = 1.$$

First, consider the vectors

$$(x_1^1, x_{r_1}^2, \dots, x_{r_{m-1}}^m), (x_2^1, x_{r_1}^2, \dots, x_{r_{m-1}}^m), \dots, (x_{k_i}^1, x_{r_1}^2, \dots, x_{r_{m-1}}^m),$$

for all possible values of r_1, \ldots, r_{m-1} , i.e., we fix all components except the first one, and vary the first component over all possible x_i^1 's used in the

convex combination that yields y_1 . Since all these vectors belong to X, their convex combination given by

$$\left(\left(\sum_{j=1}^{k_1} \alpha_j^1 x_j^1\right), x_{r_1}^2, \dots, x_{r_{m-1}}^m\right)$$

belongs to the convex hull of X for all possible values of r_1, \ldots, r_{m-1} . Now, consider the vectors

$$\left(\left(\sum_{j=1}^{k_1} \alpha_j^1 x_j^1 \right), x_1^2, \dots, x_{r_{m-1}}^m \right), \dots, \left(\left(\sum_{j=1}^{k_1} \alpha_j^1 x_j^1 \right), x_{k_2}^2, \dots, x_{r_{m-1}}^m \right),$$

i.e., fix all components except the second one, and vary the second component over all possible x_j^2 's used in the convex combination that yields y_2 . Since all these vectors belong to conv(X), their convex combination given by

$$\left(\left(\sum_{j=1}^{k_1} \alpha_j^1 x_j^1 \right), \left(\sum_{j=1}^{k_2} \alpha_j^2 x_j^2 \right), \dots, x_{r_{m-1}}^m \right)$$

belongs to the convex hull of X for all possible values of r_2, \ldots, r_{m-1} . Proceeding in this way, we see that the vector given by

$$\left(\left(\sum_{j=1}^{k_1} \alpha_j^1 x_j^1\right), \left(\sum_{j=1}^{k_2} \alpha_j^2 x_j^2\right), \dots, \left(\sum_{j=1}^{k_m} \alpha_j^m x_j^m\right)\right)$$

belongs to conv(X), thus proving our claim.

Next, we show the corresponding result for the closure of X. Assume that $y=(x_1,\ldots,x_m)\in\operatorname{cl}(X)$. This implies that there exists some sequence $\{y^k\}\subset X$ such that $y^k\to y$. Since $y^k\in X$, we have that $y^k=(x_1^k,\ldots,x_m^k)$ with $x_i^k\in X_i$ for each i and k. Since $y^k\to y$, it follows that $x_i\in\operatorname{cl}(X_i)$ for each i, and hence $y\in\operatorname{cl}(X_1)\times\cdots\times\operatorname{cl}(X_m)$. Conversely, suppose that $y=(x_1,\ldots,x_m)\in\operatorname{cl}(X_1)\times\cdots\times\operatorname{cl}(X_m)$. This implies that there exist sequences $\{x_i^k\}\subset X_i$ such that $x_i^k\to x_i$ for each $i=1,\ldots,m$. Since $x_i^k\in X_i$ for each i and k, we have that $y^k=(x_1^k,\ldots,x_m^k)\in X$ and $\{y^k\}$ converges to $y=(x_1,\ldots,x_m)$, implying that $y\in\operatorname{cl}(X)$.

Finally, we show the corresponding result for the affine hull of X. Let's assume, by using a translation argument if necessary, that all the X_i 's contain the origin, so that $\operatorname{aff}(X_1), \ldots, \operatorname{aff}(X_m)$ as well as $\operatorname{aff}(X)$ are all subspaces.

Assume that $y \in \text{aff}(X)$. Let the dimension of aff(X) be r, and let y^1, \ldots, y^r be linearly independent vectors in X that span aff(X). Thus, we can represent y as

$$y = \sum_{i=1}^{r} \beta^i y^i,$$

where β^1, \ldots, β^r are scalars. Since $y^i \in X$, we have that $y^i = (x_1^i, \ldots, x_m^i)$ with $x_j^i \in X_j$. Thus,

$$y = \sum_{i=1}^{r} \beta^{i}(x_{1}^{i}, \dots, x_{m}^{i}) = \left(\sum_{i=1}^{r} \beta^{i} x_{1}^{i}, \dots, \sum_{i=1}^{r} \beta^{i} x_{m}^{i}\right),$$

implying that $y \in \operatorname{aff}(X_1) \times \cdots \times \operatorname{aff}(X_m)$. Now, assume that $y \in \operatorname{aff}(X_1) \times \cdots \times \operatorname{aff}(X_m)$. Let the dimension of $\operatorname{aff}(X_i)$ be r_i , and let $x_i^1, \dots, x_i^{r_i}$ be linearly independent vectors in X_i that span $\operatorname{aff}(X_i)$. Thus, we can represent y as

$$y = \left(\sum_{j=1}^{r_1} \beta_1^j x_1^j, \dots, \sum_{j=1}^{r_m} \beta_m^j x_m^j\right).$$

Since each X_i contains the origin, we have that the vectors

$$\left(\sum_{j=1}^{r_1} \beta_1^j x_1^j, 0, \dots, 0\right), \left(0, \sum_{j=1}^{r_2} \beta_2^j x_2^j, 0, \dots, 0\right), \dots, \left(0, \dots, \sum_{j=1}^{r_m} \beta_m^j x_m^j\right),$$

belong to $\operatorname{aff}(X)$, and so does their sum, which is the vector y. Thus, $y \in \operatorname{aff}(X)$, concluding the proof.

(b) Assume that $y \in \text{cone}(X)$. We can represent y as

$$y = \sum_{i=1}^{r} \alpha^{i} y^{i},$$

for some r, where $\alpha^1, \ldots, \alpha^r$ are nonnegative scalars and $y_i \in X$ for all i. Since $y^i \in X$, we have that $y^i = (x_1^i, \ldots, x_m^i)$ with $x_j^i \in X_j$. Thus,

$$y = \sum_{i=1}^{r} \alpha^{i}(x_{1}^{i}, \dots, x_{m}^{i}) = \left(\sum_{i=1}^{r} \alpha^{i} x_{1}^{i}, \dots, \sum_{i=1}^{r} \alpha^{i} x_{m}^{i}\right),$$

implying that $y \in \text{cone}(X_1) \times \cdots \times \text{cone}(X_m)$.

Conversely, assume that $y \in \text{cone}(X_1) \times \cdots \times \text{cone}(X_m)$. Then, we can represent y as

$$y = \left(\sum_{j=1}^{r_1} \alpha_1^j x_1^j, \dots, \sum_{j=1}^{r_m} \alpha_m^j x_m^j\right),$$

where $x_i^j \in X_i$ and $\alpha_i^j \geq 0$ for each i and j. Since each X_i contains the origin, we have that the vectors

$$\left(\sum_{j=1}^{r_1} \alpha_1^j x_1^j, 0, \dots, 0\right), \left(0, \sum_{j=1}^{r_2} \alpha_2^j x_2^j, 0, \dots, 0\right) \dots, \left(0, \dots, \sum_{j=1}^{r_m} \alpha_m^j x_m^j\right),$$

belong to the cone(X), and so does their sum, which is the vector y. Thus, $y \in cone(X)$, concluding the proof.

Finally, consider the example where

$$X_1 = \{0, 1\} \subset \Re, \qquad X_2 = \{1\} \subset \Re.$$

For this example, $cone(X_1) \times cone(X_2)$ is given by the nonnegative quadrant, whereas cone(X) is given by the two halflines $\alpha(0,1)$ and $\alpha(1,1)$ for $\alpha \geq 0$ and the region that lies between them.

SECTION 1.3: Relative Interior and Closure

1.20 (Characterization of Twice Continuously Differentiable Convex Functions)

Let C be a nonempty convex subset of \Re^n and let $f: \Re^n \to \Re$ be twice continuously differentiable over \Re^n . Let S be the subspace that is parallel to the affine hull of C. Show that f is convex over C if and only if $y'\nabla^2 f(x)y \ge 0$ for all $x \in C$ and $y \in S$. [In particular, when C has nonempty interior, f is convex over C if and only if $\nabla^2 f(x)$ is positive semidefinite for all $x \in C$.]

Solution: Suppose that $f: \Re^n \to \Re$ is convex over C. We first show that for all $x \in \mathrm{ri}(C)$ and $y \in S$, we have $y' \nabla^2 f(x) y \geq 0$. Assume to arrive at a contradiction, that there exists some $\overline{x} \in \mathrm{ri}(C)$ such that for some $y \in S$, we have

$$y'\nabla^2 f(\overline{x})y < 0.$$

Without loss of generality, we may assume that ||y|| = 1. Using the continuity of $\nabla^2 f$, we see that there is an open ball $B(\overline{x}, \epsilon)$ centered at \overline{x} with radius ϵ such that $B(\overline{x}, \epsilon) \cap \operatorname{aff}(C) \subset C$ [since $\overline{x} \in \operatorname{ri}(C)$], and

$$y'\nabla^2 f(x)y < 0, \qquad \forall \ x \in B(\overline{x}, \epsilon).$$
 (1.9)

For all positive scalars α with $\alpha < \epsilon$, we have

$$f(\bar{x} + \alpha y) = f(\bar{x}) + \alpha \nabla f(\bar{x})' y + \frac{1}{2} y' \nabla^2 f(\bar{x} + \bar{\alpha} y) y,$$

for some $\bar{\alpha} \in [0, \alpha]$. Furthermore, $\|(\overline{x} + \overline{\alpha}y) - \overline{x}\| \le \epsilon$ [since $\|y\| = 1$ and $\bar{\alpha} < \epsilon$]. Hence, from Eq. (1.9), it follows that

$$f(\bar{x} + \alpha y) < f(\bar{x}) + \alpha \nabla f(\bar{x})'y, \quad \forall \alpha \in [0, \epsilon).$$

On the other hand, by the choice of ϵ and the assumption that $y \in S$, the vectors $\bar{x} + \alpha y$ are in C for all α with $\alpha \in [0, \epsilon)$, which is a contradiction in

view of the convexity of f over C. Hence, we have $y'\nabla^2 f(x)y \geq 0$ for all $y \in S$ and all $x \in \text{ri}(C)$.

Next, let \overline{x} be a point in C that is not in the relative interior of C. Then, by the Line Segment Principle, there is a sequence $\{x_k\} \subset \mathrm{ri}(C)$ such that $x_k \to \overline{x}$. As seen above, $y' \nabla^2 f(x_k) y \geq 0$ for all $y \in S$ and all k, which together with the continuity of $\nabla^2 f$ implies that

$$y'\nabla^2 f(\overline{x})y = \lim_{k \to \infty} y'\nabla^2 f(x_k)y \ge 0, \quad \forall y \in S.$$

It follows that $y'\nabla^2 f(x)y \geq 0$ for all $x \in C$ and $y \in S$.

Conversely, assume that $y'\nabla^2 f(x)y\geq 0$ for all $x\in C$ and $y\in S.$ For all $x,z\in C$ we have

$$f(z) = f(x) + (z - x)'\nabla f(x) + \frac{1}{2}(z - x)'\nabla^2 f(x + \alpha(z - x))(z - x)$$

for some $\alpha \in [0,1]$. Since $x, z \in C$, we have that $(z-x) \in S$, and using the convexity of C and our assumption, it follows that

$$f(z) \ge f(x) + (z - x)' \nabla f(x), \quad \forall x, z \in C.$$

From Prop. 1.1.7(a), we conclude that f is convex over C.

1.21

Construct an example of a point in a nonconvex set X that has the prolongation property of Prop. 1.3.3 but is not a relative interior point of X.

Solution: Take two intersecting lines in the plane, and consider the point of intersection.

For another example, take the union of two circular disks in the plane, which have a single common point, and consider the common point.

1.22

Let C be a nonempty convex subset of \Re^n . Show that

$$ri(C) = int(C + S^{\perp}) \cap C$$
,

where S is the subspace that is parallel to the affine hull of C.

Solution: For any vector $a \in \Re^n$, we have $\operatorname{ri}(C+a) = \operatorname{ri}(C) + a$ (cf. Prop. 1.3.7). Therefore, we can assume without loss of generality that $0 \in C$, and $\operatorname{aff}(C)$ coincides with S.

Let $x \in ri(C)$. Then there exists some open ball $B(x,\epsilon)$ centered at x with radius $\epsilon > 0$ such that

$$B(x,\epsilon) \cap S \subset C.$$
 (1.10)

We now show that $B(x,\epsilon)\subset C+S^\perp$. Let z be a vector in $B(x,\epsilon)$. Then, we can express z as $z=x+\alpha y$ for some vector $y\in\Re^n$ with $\|y\|=1$, and some $\alpha\in[0,\epsilon)$. Since S and S^\perp are orthogonal subspaces, y can be uniquely decomposed as $y=y_S+y_{S^\perp}$, where $y_S\in S$ and $y_{S^\perp}\in S^\perp$. Since $\|y\|=1$, this implies that $\|y_S\|\leq 1$ (Pythagorean Theorem), and using Eq. (1.10), we obtain

$$x + \alpha y_S \in B(x, \epsilon) \cap S \subset C$$

from which it follows that the vector $z = x + \alpha y$ belongs to $C + S^{\perp}$, implying that $B(x, \epsilon) \subset C + S^{\perp}$. This shows that $x \in \text{int}(C + S^{\perp}) \cap C$.

Conversely, let $x \in \operatorname{int}(C+S^{\perp}) \cap C$. We have that $x \in C$ and there exists some open ball $B(x,\epsilon)$ centered at x with radius $\epsilon > 0$ such that $B(x,\epsilon) \subset C+S^{\perp}$. Since C is a subset of S, it can be seen that $(C+S^{\perp}) \cap S = C$. Therefore,

$$B(x, \epsilon) \cap S \subset C$$
,

implying that $x \in ri(C)$.

1.23

Let x_0, \ldots, x_m be vectors in \Re^n such that $x_1 - x_0, \ldots, x_m - x_0$ are linearly independent. The convex hull of x_0, \ldots, x_m is called an *m-dimensional simplex*, and x_0, \ldots, x_m are called the *vertices* of the simplex.

- (a) Show that the dimension of a convex set is the maximum of the dimensions of all the simplices contained in the set.
- (b) Use part (a) to show that a nonempty convex set has a nonempty relative interior.

Solution: (a) Let C be the given convex set. The convex hull of any subset of C is contained in C. Therefore, the maximum dimension of the various simplices contained in C is the largest m for which C contains m+1 vectors x_0, \ldots, x_m such that $x_1 - x_0, \ldots, x_m - x_0$ are linearly independent.

Let $K = \{x_0, \ldots, x_m\}$ be such a set with m maximal, and let $\operatorname{aff}(K)$ denote the affine hull of set K. Then, we have $\dim(\operatorname{aff}(K)) = m$, and since $K \subset C$, it follows that $\operatorname{aff}(K) \subset \operatorname{aff}(C)$.

We claim that $C \subset \operatorname{aff}(K)$. To see this, assume that there exists some $x \in C$, which does not belong to $\operatorname{aff}(K)$. This implies that the set $\{x, x_0, \ldots, x_m\}$ is a set of m + 2 vectors in C such that $x - x_0, x_1 - x_0 = 1$

 $x_0, \ldots, x_m - x_0$ are linearly independent, contradicting the maximality of m. Hence, we have $C \subset \text{aff}(K)$, and it follows that

$$\operatorname{aff}(K) = \operatorname{aff}(C),$$

thereby implying that $\dim(C) = m$.

(b) We first consider the case where C is n-dimensional with n > 0 and show that the interior of C is not empty. By part (a), an n-dimensional convex set contains an n-dimensional simplex. We claim that such a simplex S has a nonempty interior. Indeed, applying an affine transformation if necessary, we can assume that the vertices of S are the vectors $(0,0,\ldots,0),(1,0,\ldots,0),\ldots,(0,0,\ldots,1)$, i.e.,

$$S = \left\{ (x_1, \dots, x_n) \mid x_i \ge 0, \ \forall \ i = 1, \dots, n, \ \sum_{i=1}^n x_i \le 1 \right\}.$$

The interior of the simplex S,

$$int(S) = \left\{ (x_1, \dots, x_n) \mid x_i > 0, \ \forall \ i = 1, \dots, n, \ \sum_{i=1}^n x_i < 1 \right\},$$

is nonempty, which in turn implies that int(C) is nonempty.

For the case where $\dim(C) < n$, consider the n-dimensional set $C+S^{\perp}$, where S^{\perp} is the orthogonal complement of the subspace parallel to $\operatorname{aff}(C)$. Since $C+S^{\perp}$ is a convex set, it follows from the above argument that $\operatorname{int}(C+S^{\perp})$ is nonempty. Let $x\in\operatorname{int}(C+S^{\perp})$. We can represent x as $x=x_C+x_{S^{\perp}}$, where $x_C\in C$ and $x_{S^{\perp}}\in S^{\perp}$. It can be seen that $x_C\in\operatorname{int}(C+S^{\perp})$. Since

$$ri(C) = int(C + S^{\perp}) \cap C$$

[cf. Exercise 1.22(a)], it follows that $x_c \in ri(C)$, so ri(C) is nonempty.

1.24 (Characterizations of Relative Interior)

Let C be a nonempty convex set.

- (a) Show the following refinement of the Prolongation Lemma (Prop. 1.3.3): $x \in ri(C)$ if and only if for every $\overline{x} \in aff(C)$, there exists a $\gamma > 0$ such that $x + \gamma(x \overline{x}) \in C$.
- (b) Show that cone(C) = aff(C) if and only if $0 \in ri(C)$.

Solution: (a) Let $x \in \text{ri}(C)$. We will show that for every $\overline{x} \in \text{aff}(C)$, there exists a $\gamma > 1$ such that $x + (\gamma - 1)(x - \overline{x}) \in C$. This is true if $\overline{x} = x$, so assume that $\overline{x} \neq x$. Since $x \in \text{ri}(C)$, there exists $\epsilon > 0$ such that

$$\{z \mid ||z - x|| < \epsilon\} \cap \operatorname{aff}(C) \subset C.$$

Choose a point $\overline{x}_{\epsilon} \in C$ in the intersection of the ray $\{x + \alpha(\overline{x} - x) \mid \alpha \geq 0\}$ and the set $\{z \mid ||z - x|| < \epsilon\} \cap \text{aff}(C)$. Then, for some positive scalar α_{ϵ} ,

$$x - \overline{x}_{\epsilon} = \alpha_{\epsilon}(x - \overline{x}).$$

Since $x \in ri(C)$ and $\overline{x}_{\epsilon} \in C$, by Prop. 1.3.1(c), there is $\gamma_{\epsilon} > 1$ such that

$$x + (\gamma_{\epsilon} - 1)(x - \overline{x}_{\epsilon}) \in C$$
,

which in view of the preceding relation implies that

$$x + (\gamma_{\epsilon} - 1)\alpha_{\epsilon}(x - \overline{x}) \in C.$$

The result follows by letting $\gamma = 1 + (\gamma_{\epsilon} - 1)\alpha_{\epsilon}$ and noting that $\gamma > 1$, since $(\gamma_{\epsilon} - 1)\alpha_{\epsilon} > 0$.

The converse assertion follows from the fact $C \subset \operatorname{aff}(C)$ and Prop. 1.3.1(c).

(b) Assume that $0 \in \text{ri}(C)$. Then, the inclusion $\text{cone}(C) \subset \text{aff}(C)$ is evident. For the reverse inclusion, note that if $\overline{x} \in \text{aff}(C)$, then $-\overline{x} \in \text{aff}(C)$, so applying part (a) with x = 0, we have that $\gamma \overline{x} \in C$ for some $\gamma > 0$. Hence $\overline{x} \in \text{cone}(C)$ and $\text{aff}(C) \subset \text{cone}(C)$.

Conversely, assume that $\operatorname{aff}(C)=\operatorname{cone}(C)$. We will show that $0\in\operatorname{ri}(C)$. Indeed if this is not so, by applying part (a) with x=0, it follows that there exists $\overline{x}\in\operatorname{aff}(C)$ such that $\gamma(-\overline{x})\notin C$ for all $\gamma>0$. Hence $-\overline{x}\notin\operatorname{cone}(C)$, a contradiction.

1.25

Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a convex function, let γ be a scalar, and let C be a nonempty convex subset of \mathbb{R}^n .

- (a) Show that if $f(x) < \gamma$ for some x, then $f(x) < \gamma$ for some $x \in \text{ri}(\text{dom}(f))$.
- (b) Show that if $C \subset \text{ri}(\text{dom}(f))$ and $f(x) < \gamma$ for some $x \in \text{cl}(C)$, then $f(x) < \gamma$ for some $x \in \text{ri}(C)$.
- (c) Show that if $C \subset \text{dom}(f)$ and $f(x) \geq \gamma$ for all $x \in C$, then $f(x) \geq \gamma$ for all $x \in \text{cl}(C)$.

Solution: (a) Assume the contrary, i.e., that $f(x) \geq \gamma$ for all $x \in \text{ri}(\text{dom}(f))$. Let \overline{x} be such that $f(\overline{x}) < \gamma$ and let \tilde{x} be any vector in ri(dom(f)). By the Line Segment Principle, all the points on the line segment connecting \overline{x} and \tilde{x} , except possibly \overline{x} , belong to ri(dom(f)) and therefore,

$$f(\alpha \tilde{x} + (1 - \alpha)\overline{x}) \ge \gamma, \quad \forall \alpha \in (0, 1].$$

Thus, we have

$$\alpha f(\tilde{x}) + (1 - \alpha)f(\overline{x}) \ge f(\alpha \tilde{x} + (1 - \alpha)\overline{x}) \ge \gamma, \quad \forall \alpha \in (0, 1].$$

By letting $\alpha \to 0$, it follows that $f(\overline{x}) \ge \gamma$, a contradiction.

(b) Define

$$g(x) = \begin{cases} f(x) & \text{if } x \in \text{cl}(C), \\ \infty & \text{otherwise.} \end{cases}$$

Then

$$ri(C) \subset dom(g) \subset cl(C),$$

so that $\operatorname{ri}(\operatorname{dom}(g)) = \operatorname{ri}(C)$, by Prop. 1.3.5. By hypothesis, there is an \overline{x} with $g(\overline{x}) < \gamma$, so by part (a), there exists an $\tilde{x} \in \operatorname{ri}(\operatorname{dom}(g))$ with $g(\tilde{x}) < \gamma$. This vector belongs to $\operatorname{ri}(C)$ and satisfies $f(\tilde{x}) < \alpha$.

(c) Assume the contrary, i.e., that $f(x) < \gamma$ for some $x \in \text{cl}(C)$. Then, by part (b), we have $f(x) < \gamma$ for some $x \in \text{ri}(C)$, which contradicts the hypothesis.

1.26

Let C_1 and C_2 be two nonempty convex sets such that $C_1 \subset C_2$.

- (a) Give an example showing that $ri(C_1)$ need not be a subset of $ri(C_2)$.
- (b) Assuming that the sets C_1 and C_2 have the same affine hull, show that $ri(C_1) \subset ri(C_2)$.
- (c) Assuming that the set $ri(C_1) \cap ri(C_2)$ is nonempty, show that $ri(C_1) \subset ri(C_2)$.
- (d) Assuming that the set $C_1 \cap ri(C_2)$ is nonempty, show that the set $ri(C_1) \cap ri(C_2)$ is nonempty.

Solution: (a) Let C_1 be the segment $\{(x_1, x_2) \mid 0 \le x_1 \le 1, x_2 = 0\}$ and let C_2 be the box $\{(x_1, x_2) \mid 0 \le x_1 \le 1, 0 \le x_2 \le 1\}$. We have

$$ri(C_1) = \{(x_1, x_2) \mid 0 < x_1 < 1, \ x_2 = 0\},\$$

$$ri(C_2) = \{(x_1, x_2) \mid 0 < x_1 < 1, \ 0 < x_2 < 1\}.$$

Thus $C_1 \subset C_2$, while $ri(C_1) \cap ri(C_2) = \emptyset$.

- (b) Let $x \in \text{ri}(C_1)$, and consider a open ball B centered at x such that $B \cap \text{aff}(C_1) \subset C_1$. Since $\text{aff}(C_1) = \text{aff}(C_2)$ and $C_1 \subset C_2$, it follows that $B \cap \text{aff}(C_2) \subset C_2$, so $x \in \text{ri}(C_2)$. Hence $\text{ri}(C_1) \subset \text{ri}(C_2)$.
- (c) Because $C_1 \subset C_2$, we have

$$ri(C_1) = ri(C_1 \cap C_2).$$

Since $ri(C_1) \cap ri(C_2) \neq \emptyset$, there holds

$$ri(C_1 \cap C_2) = ri(C_1) \cap ri(C_2)$$

(Prop. 1.3.8). Combining the preceding two relations, we obtain $ri(C_1) \subset ri(C_2)$.

(d) Let x_2 be in the intersection of C_1 and $\mathrm{ri}(C_2)$, and let x_1 be in the relative interior of C_1 [ri (C_1) is nonempty by Prop. 1.3.2]. If $x_1 = x_2$, then we are done, so assume that $x_1 \neq x_2$. By the Line Segment Principle, all the points on the line segment connecting x_1 and x_2 , except possibly x_2 , belong to the relative interior of C_1 . Since $C_1 \subset C_2$, the vector x_1 is in C_2 , so that by the Line Segment Principle, all the points on the line segment connecting x_1 and x_2 , except possibly x_1 , belong to the relative interior of C_2 . Hence, all the points on the line segment connecting x_1 and x_2 , except possibly x_1 and x_2 , belong to the intersection $\mathrm{ri}(C_1) \cap \mathrm{ri}(C_2)$, showing that $\mathrm{ri}(C_1) \cap \mathrm{ri}(C_2)$ is nonempty.

1.27

Let C be a nonempty set.

- (a) If C is convex and compact, and the origin is not in the relative boundary of C, then cone(C) is closed.
- (b) Give examples showing that the assertion of part (a) fails if C is unbounded or the origin is in the relative boundary of C.
- (c) If C is compact and the origin is not in the relative boundary of conv(C), then cone(C) is closed. *Hint*: Use part (a) and Exercise 1.14(b).

Solution: (a) If $0 \in C$, then $0 \in \text{ri}(C)$ since 0 is not on the relative boundary of C. By Exercise 1.24(b), it follows that cone(C) coincides with aff(C), which is a closed set. If $0 \notin C$, let y be in the closure of cone(C) and let $\{y_k\} \subset \text{cone}(C)$ be a sequence converging to y. By Exercise 1.13, for every y_k , there exists a nonnegative scalar α_k and a vector $x_k \in C$ such that $y_k = \alpha_k x_k$. Since $\{y_k\} \to y$, the sequence $\{y_k\}$ is bounded, implying that

$$\alpha_k ||x_k|| \le \sup_{m>0} ||y_m|| < \infty, \quad \forall k.$$

We have $\inf_{m\geq 0} ||x_m|| > 0$, since $\{x_k\} \subset C$ and C is a compact set not containing the origin, so that

$$0 \le \alpha_k \le \frac{\sup_{m \ge 0} \|y_m\|}{\inf_{m > 0} \|x_m\|} < \infty, \quad \forall k.$$

Thus, the sequence $\{(\alpha_k, x_k)\}$ is bounded and has a limit point (α, x) such that $\alpha \geq 0$ and $x \in C$. By taking a subsequence of $\{(\alpha_k, x_k)\}$ that converges to (α, x) , and by using the facts $y_k = \alpha_k x_k$ for all k and $\{y_k\} \to y$, we see that $y = \alpha x$ with $\alpha \geq 0$ and $x \in C$. Hence, $y \in \text{cone}(C)$, showing that cone(C) is closed.

(b) To see that the assertion in part (a) fails when C is unbounded, let C be the line $\{(x_1, x_2) \mid x_1 = 1, x_2 \in \Re\}$ in \Re^2 not passing through the origin. Then, $\operatorname{cone}(C)$ is the nonclosed set $\{(x_1, x_2) \mid x_1 > 0, x_2 \in \Re\} \cup \{(0, 0)\}$.

To see that the assertion in part (a) fails when C contains the origin on its relative boundary, let C be the closed ball $\{(x_1, x_2) \mid (x_1 - 1)^2 + x_2^2 \le 1\}$ in \Re^2 . Then, cone(C) is the nonclosed set $\{(x_1, x_2) \mid x_1 > 0, x_2 \in \Re\} \cup \{(0, 0)\}$ (see Fig. 1.3.2).

(c) Since C is compact, the convex hull of C is compact (cf. Prop. 1.2.2). Because conv(C) does not contain the origin on its relative boundary, by part (a), the cone generated by conv(C) is closed. By Exercise 1.14(b), cone(conv(C)) coincides with cone(C) implying that cone(C) is closed.

1.28 (Closure and Relative Interior of Cones)

- (a) Let C be a nonempty convex cone. Show that $\operatorname{cl}(C)$ and $\operatorname{ri}(C)$ is also a convex cone.
- (b) Let $C = \operatorname{cone}(\{x_1, \dots, x_m\})$. Show that

$$ri(C) = \left\{ \sum_{i=1}^{m} \alpha_i x_i \mid \alpha_i > 0, \ i = 1, \dots, m \right\}.$$

Solution: (a) Let $x \in \operatorname{cl}(C)$ and let α be a positive scalar. Then, there exists a sequence $\{x_k\} \subset C$ such that $x_k \to x$, and since C is a cone, $\alpha x_k \in C$ for all k. Furthermore, $\alpha x_k \to \alpha x$, implying that $\alpha x \in \operatorname{cl}(C)$. Hence, $\operatorname{cl}(C)$ is a cone, and it also convex since the closure of a convex set if convex.

By Prop. 1.3.2(a), the relative interior of a convex set is convex. To show that $\operatorname{ri}(C)$ is a cone, let $x \in \operatorname{ri}(C)$. Then, $x \in C$ and since C is a cone, $\alpha x \in C$ for all $\alpha > 0$. By the Line Segment Principle, all the points on the line segment connecting x and αx , except possibly αx , belong to $\operatorname{ri}(C)$. Since this is true for every $\alpha > 0$, it follows that $\alpha x \in \operatorname{ri}(C)$ for all $\alpha > 0$, showing that $\operatorname{ri}(C)$ is a cone.

(b) Consider the linear transformation A that maps $(\alpha_1, \ldots, \alpha_m) \in \mathbb{R}^m$ into $\sum_{i=1}^m \alpha_i x_i \in \mathbb{R}^n$. Note that C is the image of the nonempty convex set

$$\{(\alpha_1,\ldots,\alpha_m)\mid \alpha_1\geq 0,\ldots,\alpha_m\geq 0\}$$

under A. Therefore, by using Prop. 1.3.6, we have

$$\operatorname{ri}(C) = \operatorname{ri}\left(A \cdot \left\{ (\alpha_1, \dots, \alpha_m) \mid \alpha_1 \ge 0, \dots, \alpha_m \ge 0 \right\} \right)$$

$$= A \cdot \operatorname{ri}\left(\left\{ (\alpha_1, \dots, \alpha_m) \mid \alpha_1 \ge 0, \dots, \alpha_m \ge 0 \right\} \right)$$

$$= A \cdot \left\{ (\alpha_1, \dots, \alpha_m) \mid \alpha_1 > 0, \dots, \alpha_m > 0 \right\}$$

$$= \left\{ \sum_{i=1}^m \alpha_i x_i \mid \alpha_1 > 0, \dots, \alpha_m > 0 \right\}.$$

1.29 (Closure and Relative Interior of Level Sets)

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let $\gamma > \inf_{x \in \Re^n} f(x)$.

(a) Show that

$$\operatorname{ri}\left(\left\{x\mid f(x)\leq\gamma\right\}\right)=\operatorname{ri}\left(\left\{x\mid f(x)<\gamma\right\}\right)=\left\{x\in\operatorname{ri}\left(\operatorname{dom}(f)\right)\mid f(x)<\gamma\right\},$$

$$\operatorname{cl}\left(\left\{x\mid f(x)\leq\gamma\right\}\right)=\operatorname{cl}\left(\left\{x\mid f(x)<\gamma\right\}\right)=\left\{x\mid (\operatorname{cl}f)(x)\leq\gamma\right\}.$$

- (b) The sets $\{x \mid f(x) \leq \gamma\}$ and $\{x \mid f(x) < \gamma\}$ have the same dimension as dom(f).
- (c) If f is real-valued, $\{x \mid f(x) \leq \gamma\}$ has nonempty interior. Furthermore, for all $\gamma, \overline{\gamma}$ with $\inf_{x \in \Re^n} f(x) < \gamma < \overline{\gamma}$, the interior of $\{x \mid f(x) \leq \gamma\}$ is contained in the interior of $\{x \mid f(x) \leq \overline{\gamma}\}$.

Solution: We have for every $\gamma \in \Re$

$$\{(x,\gamma) \mid f(x) \le \gamma\} = \operatorname{epi}(f) \cap M, \tag{1.11}$$

where M is the set

$$M = \big\{ (x,\gamma) \mid x \in \Re^n \big\}.$$

By Prop. 1.3.10,

$$ri(epi(f)) = \{(x, w) \mid x \in ri(dom(f)), f(x) < w\},$$
 (1.12)

so for $\gamma > \inf_{x \in \Re^n} f(x)$, using also Exercise 1.25(a), we have $\operatorname{ri}(\operatorname{epi}(f)) \cap M \neq \emptyset$. It follows from Eq. (1.11) and Prop. 1.3.8 that

$$\operatorname{ri}\left(\left\{(x,\gamma)\mid f(x)\leq\gamma\right\}\right)=\operatorname{ri}\left(\operatorname{epi}(f)\right)\cap M,$$

$$\operatorname{cl}(\{(x,\gamma) \mid f(x) \leq \gamma\}) = \operatorname{cl}(\operatorname{epi}(f)) \cap M.$$

The last two equations together with Eq. (1.12) show that for every $\gamma > \inf_{x \in \mathbb{R}^n} f(x)$, we have

$$\operatorname{ri}\left(\left\{x\mid f(x)\leq\gamma\right\}\right)=\left\{x\in\operatorname{ri}(\operatorname{dom}(f))\mid f(x)<\gamma\right\},$$
$$\operatorname{cl}\left(\left\{x\mid f(x)\leq\gamma\right\}\right)=\left\{x\mid (\operatorname{cl}f)(x)\leq\gamma\right\}.$$

Next, we show that

$$\operatorname{cl}(\left\{x \mid f(x) \le \gamma\right\}) = \operatorname{cl}(\left\{x \mid f(x) < \gamma\right\}). \tag{1.13}$$

Clearly,

$$\operatorname{cl}(\{x \mid f(x) \le \gamma\}) \supset \operatorname{cl}(\{x \mid f(x) < \gamma\}).$$

To show the reverse inclusion, let $\overline{x} \in \text{cl}(\{x \mid f(x) \leq \gamma\})$, or equivalently, $(\text{cl}\, f)(\overline{x}) \leq \gamma$. Also, choose \tilde{x} such that $\tilde{x} \in \text{ri}(\text{dom}(f))$ and $f(\tilde{x}) < \gamma$ [such a vector exists by Exercise 1.25(a), in view of the assumption $\gamma > \inf_{x \in \mathbb{R}^n} f(x)$]. Then, by Prop. 1.3.15, along the line segment connecting \tilde{x} and \overline{x} , there is a sequence $\{x_k\} \subset \text{ri}(\text{dom}(f))$ that converges to \overline{x} and satisfies $f(x_k) < \gamma$ for all k. It follows that $\overline{x} \in \text{cl}(\{x \mid f(x) < \gamma\})$, showing that

$$\mathrm{cl}\big(\big\{x\mid f(x)\leq\gamma\big\}\big)\subset\mathrm{cl}\big(\big\{x\mid f(x)<\gamma\big\}\big),$$

and thereby proving Eq. (1.13).

Next note that since the sets $\{x \mid f(x) \leq \gamma\}$ and $\{x \mid f(x) < \gamma\}$ have the same closure, by Prop. 1.3.5(c), they have the same relative interior, i.e.,

$$\mathrm{ri}\big(\big\{x\mid f(x)\leq\gamma\big\}\big)=\mathrm{ri}\big(\big\{x\mid f(x)<\gamma\big\}\big).$$

Finally, since the sets $\{x \mid f(x) \leq \gamma\}$ and $\{x \mid f(x) < \gamma\}$ have the same closure and relative interior, they also have the same affine hull, and hence the same dimension.

1.30 (Relative Interior Intersection Lemma)

Let C_1 and C_2 be convex sets. Show that

$$C_1 \cap \operatorname{ri}(C_2) \neq \emptyset$$
 if and only if $\operatorname{ri}(C_1 \cap \operatorname{aff}(C_2)) \cap \operatorname{ri}(C_2) \neq \emptyset$

Hint: Choose $\overline{x} \in \text{ri}(C_1 \cap \text{aff}(C_2))$ and $x \in C_1 \cap \text{ri}(C_2)$ [which belongs to $C_1 \cap \text{aff}(C_2)$], consider the line segment connecting x and \overline{x} , and use the Line Segment Principle to conclude that points close to x belong to $\text{ri}(C_1 \cap \text{aff}(C_2)) \cap \text{ri}(C_2)$.

Solution: Let $x \in C_1 \cap \text{ri}(C_2)$ and $\overline{x} \in \text{ri}(C_1 \cap \text{aff}(C_2))$. Let L be the line segment connecting x and \overline{x} . Then L belongs to $C_1 \cap \text{aff}(C_2)$ since

both of its endpoints belong to $C_1 \cap \operatorname{aff}(C_2)$. Hence, by the Line Segment Principle, all points of L except possibly x, belong to $\operatorname{ri}(C_1 \cap \operatorname{aff}(C_2))$. On the other hand, by the definition of relative interior, all points of L that are sufficiently close to x belong to $\operatorname{ri}(C_2)$, and these points, except possibly for x belong to $\operatorname{ri}(C_1 \cap \operatorname{aff}(C_2)) \cap \operatorname{ri}(C_2)$.

1.31 (Closedness of Finitely Generated Cones)

Let a_1, \ldots, a_r be vectors of \Re^n . Then the generated cone

$$C = \operatorname{cone}(\{a_1, \dots, a_r\}) = \left\{ x \mid x = \sum_{j=1}^r \mu_j a_j, \, \mu_j \ge 0, \, j = 1, \dots, r \right\}$$

is closed. Note and Hint: One way to show this is by noting that C can be written as AX where A is the matrix with columns a_1, \ldots, a_r and X is the polyhedral set of all (μ_1, \ldots, μ_r) with $\mu_j \geq 0$ for all j. The result then follows from Prop. 1.4.13. The purpose of this exercise is to explore an alternative and more elementary method of proof. To this end, use induction on the number of vectors r. When r = 1, C is either $\{0\}$ (if $a_1 = 0$) or a halfline, and is therefore closed. Suppose, for some $r \geq 1$, all cones of the form

$$C_r = \left\{ x \mid x = \sum_{j=1}^r \mu_j a_j, \, \mu_j \ge 0 \right\},$$

are closed. Then, show that a cone of the form

$$C_{r+1} = \left\{ x \mid x = \sum_{j=1}^{r+1} \mu_j a_j, \, \mu_j \ge 0 \right\}$$

is also closed.

Solution: Without loss of generality, assume that $||a_j|| = 1$ for all j. There are two cases: (i) The vectors $-a_1, \ldots, -a_{r+1}$ belong to C_{r+1} , in which case C_{r+1} is the subspace spanned by a_1, \ldots, a_{r+1} and is therefore closed, and (ii) The negative of one of the vectors, say $-a_{r+1}$, does not belong to C_{r+1} . In this case, consider the cone C_r , which is closed by the induction hypothesis. Let

$$m = \min_{x \in C_r, ||x|| = 1} a'_{r+1} x.$$

Since, the set $\{x \in C_r \mid ||x|| = 1\}$ is nonempty and compact, the minimum above is attained at some x^* by Weierstrass' theorem. We have, using the Schwartz inequality,

$$m = a'_{r+1}x^* \ge -||a_{r+1}|| \cdot ||x^*|| = -1,$$

with equality if and only if $x^* = -a_{r+1}$. It follows that

$$m > -1$$
,

since otherwise we would have $x^* = -a_{r+1}$, which violates the hypothesis $(-a_{r+1}) \notin C_r$. Let $\{x_k\}$ be a convergent sequence in C_{r+1} . We will prove that its limit belongs to C_{r+1} , thereby showing that C_{r+1} is closed. Indeed, for all k, we have $x_k = \xi_k a_{r+1} + y_k$, where $\xi_k \geq 0$ and $y_k \in C_r$. Using the fact $||a_{r+1}|| = 1$, we obtain

$$||x_k||^2 = \xi_k^2 + ||y_k||^2 + 2\xi_k a'_{r+1} y_k$$

$$\geq \xi_k^2 + ||y_k||^2 + 2m\xi_k ||y_k||$$

$$= (\xi_k - ||y_k||)^2 + 2(1+m)\xi_k ||y_k||.$$

Since $\{x_k\}$ converges, $\xi_k \ge 0$, and 1 + m > 0, it follows that the sequences $\{\xi_k\}$ and $\{y_k\}$ are bounded and hence, they have limit points denoted by ξ and y, respectively. The limit of $\{x_k\}$ is

$$\lim_{k \to \infty} (\xi_k a_{r+1} + y_k) = \xi a_{r+1} + y,$$

which belongs to C_{r+1} , since $\xi \geq 0$ and $y \in C_r$ (by the closure hypothesis on C_r). We conclude that C_{r+1} is closed, completing the proof.

1.32 (Improper Convex Functions)

Let $f: \Re^n \mapsto [-\infty, \infty]$ be a convex function with $dom(f) \neq \emptyset$.

(a) Show that if f is improper, then

$$f(x) = -\infty, \quad \forall x \in ri(dom(f)).$$

Furthermore,

$$(\operatorname{cl} f)(x) = \begin{cases} -\infty & \text{if } x \in \operatorname{cl}(\operatorname{dom}(f)), \\ \infty & \text{otherwise.} \end{cases}$$

(b) Show that if $f(x) < \infty$ for all $x \in \Re^n$, then either $f(x) = -\infty$ for all $x \in \Re^n$ or $f(x) > -\infty$ for all $x \in \Re^n$.

Solution: (a) Since f is improper, there exists some $\overline{x} \in \text{dom}(f)$ such that $f(\overline{x}) = -\infty$. Let $x \in \text{ri}(\text{dom}(f))$. Then by the Prolongation Principle [Prop. 1.3.1(c)], there is a vector $y \in \text{ri}(\text{dom}(f))$ such that $y \neq x$ and x lies in the line segment connecting y and \overline{x} . Thus, for some $\alpha \in (0,1)$, we have $x = \alpha y + (1 - \alpha)\overline{x}$, so by convexity of f,

$$f(x) \le \alpha f(y) + (1 - \alpha) f(\overline{x}).$$

Since $f(y) < \infty$ and $f(\overline{x}) = -\infty$, it follows that $f(x) = -\infty$.

We next note that if $x \notin \text{cl}(\text{dom}(f))$, then $(x, w) \notin \text{cl}(\text{epi}(f))$ for all $w \in \Re$, so that $(\text{cl} f)(x) = \infty$.

We finally show that $(\operatorname{cl} f)(x) = -\infty$ for all $x \in \operatorname{cl}(\operatorname{dom}(f))$. Assume, to arrive at a contradiction, that for some $x \in \operatorname{cl}(\operatorname{dom}(f))$, we have $(\operatorname{cl} f)(x) > -\infty$. By the Line Segment Principle, there exists a sequence $\{x_k\} \subset \operatorname{ri}(\operatorname{dom}(f))$ that converges to x. Since by part (a), we have $f(x_k) = -\infty$, we have that $(x_k, w) \in \operatorname{epi}(f)$ for every k and $w \in \Re$. It follows that

$$(x, w) \in \operatorname{cl}(\operatorname{epi}(f)) = \operatorname{epi}(\operatorname{cl} f), \quad \forall w \in \Re.$$

This implies that $(\operatorname{cl} f)(x) = -\infty$ for all $x \in \operatorname{cl}(\operatorname{dom}(f))$.

- (b) We have $dom(f) = \Re^n$, so either f is improper, in which case by part (a) we have $f(x) = -\infty$ for all $x \in \Re^n$, or f is proper, in which case we
- have $f(x) > -\infty$ for all $x \in \Re^n$.

1.33 (Lipschitz Continuity of Convex Functions)

Let $f: \mathbb{R}^n \to \mathbb{R}$ be a convex function and X be a bounded set in \mathbb{R}^n . Show that f is Lipschitz continuous over X, i.e., there exists a positive scalar L such that

$$|f(x) - f(y)| \le L||x - y||, \qquad \forall \ x, y \in X.$$

Note: This result is also shown with a different proof in Section 5.4, using the theory of subgradients.

Solution: Let ϵ be a positive scalar and let C_{ϵ} be the set given by

$$C_{\epsilon} = \{z \mid ||z - x|| \le \epsilon, \text{ for some } x \in \operatorname{cl}(X)\}.$$

We claim that the set C_{ϵ} is compact. Indeed, since X is bounded, so is its closure, which implies that $||z|| \leq \max_{x \in \operatorname{cl}(X)} ||x|| + \epsilon$ for all $z \in C_{\epsilon}$, showing that C_{ϵ} is bounded. To show the closedness of C_{ϵ} , let $\{z_k\}$ be a sequence in C_{ϵ} converging to some z. By the definition of C_{ϵ} , there is a corresponding sequence $\{x_k\}$ in $\operatorname{cl}(X)$ such that

$$||z_k - x_k|| \le \epsilon, \qquad \forall \ k. \tag{1.14}$$

Because $\operatorname{cl}(X)$ is compact, $\{x_k\}$ has a subsequence converging to some $x \in \operatorname{cl}(X)$. Without loss of generality, we may assume that $\{x_k\}$ converges to $x \in \operatorname{cl}(X)$. By taking the limit in Eq. (1.14) as $k \to \infty$, we obtain $||z - x|| \le \epsilon$ with $x \in \operatorname{cl}(X)$, showing that $z \in C_{\epsilon}$. Hence, C_{ϵ} is closed.

We now show that f has the Lipschitz property over X. Let x and y be two distinct points in X. Then, by the definition of C_{ϵ} , the point

$$z = y + \frac{\epsilon}{\|y - x\|}(y - x)$$

is in C_{ϵ} . Thus

$$y = \frac{\|y - x\|}{\|y - x\| + \epsilon} z + \frac{\epsilon}{\|y - x\| + \epsilon} x,$$

showing that y is a convex combination of $z \in C_{\epsilon}$ and $x \in C_{\epsilon}$. By convexity of f, we have

$$f(y) \le \frac{\|y - x\|}{\|y - x\| + \epsilon} f(z) + \frac{\epsilon}{\|y - x\| + \epsilon} f(x),$$

implying that

$$f(y) - f(x) \le \frac{\|y - x\|}{\|y - x\| + \epsilon} \left(f(z) - f(x) \right) \le \frac{\|y - x\|}{\epsilon} \left(\max_{u \in C_{\epsilon}} f(u) - \min_{v \in C_{\epsilon}} f(v) \right),$$

where in the last inequality we use Weierstrass' theorem (f is continuous over \Re^n and C_{ϵ} is compact). By switching the roles of x and y, we similarly obtain

$$f(x) - f(y) \le \frac{\|x - y\|}{\epsilon} \left(\max_{u \in C_{\epsilon}} f(u) - \min_{v \in C_{\epsilon}} f(v) \right),$$

which combined with the preceding relation yields $|f(x)-f(y)| \le L||x-y||$, where

$$L = \left(\max_{u \in C_{\epsilon}} f(u) - \min_{v \in C_{\epsilon}} f(v)\right) / \epsilon.$$

1.34 (Uniform Approximation Lemma)

Let C be a convex and compact set, and let $\{f_i \mid i \in I\}$ be a family of convex functions $f_i : C \mapsto \Re$ such that

$$\sup_{i \in I} f_i(x) = 0, \qquad \forall \ x \in C.$$

Then for every $\epsilon > 0$ there exists an index $\overline{i} \in I$ such that

$$-\epsilon \le f_{\overline{i}}(x) \le 0, \quad \forall \ x \in C.$$

Hint: Let \hat{x} be a point in the relative interior of C. For any $x \in C$ with $x \neq \hat{x}$, consider the line that starts at x and passes through \hat{x} , and let r(x) be the point at which it meets D, the relative boundary of C. Choose an index \bar{i} such that

$$f_{\overline{i}}(\hat{x}) \ge -\left(\frac{\max_{x \in C} \|x - \hat{x}\|}{\min_{x \in D} \|x - \hat{x}\|} + 1\right)^{-1} \epsilon.$$

From the paper: Yu, H., and Bertsekas, D. P., "On Near-Optimality of the Set of Finite-State Controllers for Average Cost POMDP," Mathematics of Operations Research, Vol. 33, pp. 1-11, 2008.

Solution: Let \hat{x} be a point in the relative interior of C. For any $x \in C$ with $x \neq \hat{x}$, consider the line that starts at x and passes through \hat{x} , and let r(x) be the point at which it meets D, the relative boundary of C (D is the set of points in C that are not relative interior points of C). Choose an index \bar{i} such that

$$0 \ge f_{\bar{i}}(\hat{x}) \ge -\left(\frac{\max_{x \in C} \|x - \hat{x}\|}{\min_{x \in D} \|\hat{x} - x\|} + 1\right)^{-1} \epsilon.$$
 (1.15)

Using the convexity and nonpositivity of $f_{\overline{i}}$, we have

$$f_{\overline{i}}(\hat{x}) \leq \frac{\|\hat{x} - r(x)\|}{\|x - \hat{x}\| + \|\hat{x} - r(x)\|} f_{\overline{i}}(x) + \frac{\|x - \hat{x}\|}{\|x - \hat{x}\| + \|\hat{x} - r(x)\|} f_{\overline{i}}(r(x))$$

$$\leq \frac{\|\hat{x} - r(x)\|}{\|x - \hat{x}\| + \|\hat{x} - r(x)\|} f_{\overline{i}}(x).$$

From this relation and Eq. (1.15), we obtain for all $x \in C$ with $x \neq \hat{x}$

$$\begin{split} f_{\overline{i}}(x) &\geq \frac{\|x - \hat{x}\| + \|\hat{x} - r(x)\|}{\|\hat{x} - r(x)\|} f_{\overline{i}}(\hat{x}) \\ &= \left(\frac{\|x - \hat{x}\|}{\|\hat{x} - r(x)\|} + 1\right) f_{\overline{i}}(\hat{x}) \\ &\geq \left(\frac{\max_{x \in C} \|x - \hat{x}\|}{\min_{x \in D} \|\hat{x} - x\|} + 1\right) f_{\overline{i}}(\hat{x}) \\ &\geq -\epsilon. \end{split}$$

SECTION 1.4: Recession Cones

1.35 (Recession Cones of Nonclosed Sets)

Let C be a nonempty convex set.

(a) Show that

$$R_C \subset R_{\mathrm{cl}(C)}, \qquad \mathrm{cl}(R_C) \subset R_{\mathrm{cl}(C)}.$$

Give an example where the inclusion $\operatorname{cl}(R_C) \subset R_{\operatorname{cl}(C)}$ is strict.

(b) Let \overline{C} be a closed convex set such that $C \subset \overline{C}$. Show that $R_C \subset R_{\overline{C}}$. Give an example showing that the inclusion can fail if \overline{C} is not closed.

Solution: (a) Let $y \in R_C$. Then, by the definition of R_C , $x + \alpha y \in C$ for every $x \in C$ and every $\alpha \geq 0$. Since $C \subset \operatorname{cl}(C)$, it follows that $x + \alpha y \in \operatorname{cl}(C)$ for some $x \in \operatorname{cl}(C)$ and every $\alpha \geq 0$, which, in view of part (b) of the Recession Cone Theorem (cf. Prop. 1.4.1), implies that $y \in R_{\operatorname{cl}(C)}$. Hence

$$R_C \subset R_{\operatorname{cl}(C)}$$
.

By taking closures in this relation and by using the fact that $R_{\text{cl}(C)}$ is closed [part (a) of the Recession Cone Theorem], we obtain $\text{cl}(R_C) \subset R_{\text{cl}(C)}$.

To see that the inclusion $\operatorname{cl}(R_C) \subset R_{\operatorname{cl}(C)}$ can be strict, consider the set

$$C = \{(x_1, x_2) \mid 0 \le x_1, \ 0 \le x_2 < 1\} \cup \{(0, 1)\},\$$

whose closure is

$$cl(C) = \{(x_1, x_2) \mid 0 \le x_1, \ 0 \le x_2 \le 1\}.$$

The recession cones of C and its closure are

$$R_C = \{(0,0)\}, \qquad R_{\mathrm{cl}(C)} = \{(x_1, x_2) \mid 0 \le x_1, \ x_2 = 0\}.$$

Thus, $\operatorname{cl}(R_C) = \{(0,0)\}$, and $\operatorname{cl}(R_C)$ is a strict subset of $R_{\operatorname{cl}(C)}$.

(b) Let $y \in R_C$ and let x be a vector in C. Then we have $x + \alpha y \in C$ for all $\alpha \geq 0$. Thus for the vector x, which belongs to \overline{C} , we have $x + \alpha y \in \overline{C}$ for all $\alpha \geq 0$, and it follows from part (b) of the Recession Cone Theorem (cf. Prop. 1.4.1) that $y \in R_{\overline{C}}$. Hence, $R_C \subset R_{\overline{C}}$.

To see that the inclusion $R_C \subset R_{\overline{C}}$ can fail when \overline{C} is not closed, consider the sets

$$C = \{(x_1, x_2) \mid x_1 \ge 0, \ x_2 = 0\}, \qquad \overline{C} = \{(x_1, x_2) \mid x_1 \ge 0, \ 0 \le x_2 < 1\}.$$

Their recession cones are

$$R_C = C = \{(x_1, x_2) \mid x_1 \ge 0, \ x_2 = 0\}, \qquad R_{\overline{C}} = \{(0, 0)\},\$$

showing that R_C is not a subset of $R_{\overline{C}}$.

1.36 (Recession Cones of Relative Interiors)

Let C be a nonempty convex set.

- (a) Show that $R_{ri(C)} = R_{cl(C)}$.
- (b) Show that a vector y belongs to $R_{ri(C)}$ if and only if there exists a vector $x \in ri(C)$ such that $x + \alpha y \in ri(C)$ for every $\alpha \geq 0$.
- (c) Let \overline{C} be a convex set such that $\overline{C} = \operatorname{ri}(\overline{C})$ and $C \subset \overline{C}$. Show that $R_C \subset R_{\overline{C}}$. Give an example showing that the inclusion can fail if $\overline{C} \neq \operatorname{ri}(\overline{C})$.

Solution: (a) The inclusion $R_{\mathrm{ri}(C)} \subset R_{\mathrm{cl}(C)}$ follows from Exercise 1.35(b). Conversely, let $y \in R_{\mathrm{cl}(C)}$, so that by the definition of $R_{\mathrm{cl}(C)}$, $x + \alpha y \in \mathrm{cl}(C)$ for every $x \in \mathrm{cl}(C)$ and every $\alpha \geq 0$. In particular, $x + \alpha y \in \mathrm{cl}(C)$ for every $x \in \mathrm{ri}(C)$ and every $\alpha \geq 0$. By the Line Segment Principle, all points on

the line segment connecting x and $x + \alpha y$, except possibly $x + \alpha y$, belong to $\mathrm{ri}(C)$, implying that $x + \alpha y \in \mathrm{ri}(C)$ for every $x \in \mathrm{ri}(C)$ and every $\alpha \geq 0$. Hence, $y \in R_{\mathrm{ri}(C)}$, showing that $R_{\mathrm{cl}(C)} \subset R_{\mathrm{ri}(C)}$.

(b) If $y \in R_{ri(C)}$, then by the definition of $R_{ri(C)}$ for every vector $x \in ri(C)$ and $\alpha \geq 0$, the vector $x + \alpha y$ is in ri(C), which holds in particular for some $x \in ri(C)$ [note that ri(C) is nonempty by Prop. 1.3.1(b)].

Conversely, let y be such that there exists a vector $x \in \operatorname{ri}(C)$ with $x + \alpha y \in \operatorname{ri}(C)$ for all $\alpha \geq 0$. Hence, there exists a vector $x \in \operatorname{cl}(C)$ with $x + \alpha y \in \operatorname{cl}(C)$ for all $\alpha \geq 0$, which, by part (b) of the Recession Cone Theorem (cf. Prop. 1.4.1), implies that $y \in R_{\operatorname{cl}(C)}$. Using part (a), it follows that $y \in R_{\operatorname{ri}(C)}$, completing the proof.

(c) Using Exercise 1.35(c) and the assumption that $C \subset \overline{C}$ [which implies that $C \subset \overline{\text{cl}}(C)$], we have

$$R_C \subset R_{\operatorname{cl}(\overline{C})} = R_{\operatorname{ri}(\overline{C})} = R_{\overline{C}},$$

where the equalities follow from part (a) and the assumption that $\overline{C} = ri(\overline{C})$.

To see that the inclusion $R_C \subset R_{\overline{C}}$ can fail when $\overline{C} \neq \text{ri}(\overline{C})$, consider the sets

$$C = \{(x_1, x_2) \mid x_1 \ge 0, \ 0 < x_2 < 1\}, \qquad \overline{C} = \{(x_1, x_2) \mid x_1 \ge 0, \ 0 \le x_2 < 1\},$$

for which we have $C \subset \overline{C}$ and

$$R_C = \{(x_1, x_2) \mid x_1 \ge 0, \ x_2 = 0\}, \qquad R_{\overline{C}} = \{(0, 0)\},\$$

showing that R_C is not a subset of $R_{\overline{C}}$.

1.37 (Closure Under Linear Transformations)

Let C be a nonempty convex subset of \Re^n and let A be an $m \times n$ matrix. Show that if $R_{\operatorname{cl}(C)} \cap N(A) = \{0\}$, then

$$\operatorname{cl}(A \cdot C) = A \cdot \operatorname{cl}(C), \qquad A \cdot R_{\operatorname{cl}(C)} = R_{A \cdot \operatorname{cl}(C)}.$$

Give an example showing that $A \cdot R_{\operatorname{cl}(C)}$ and $R_{A \cdot \operatorname{cl}(C)}$ may differ when $R_{\operatorname{cl}(C)} \cap N(A) \neq \{0\}$.

Solution: Let y be in the closure of $A \cdot C$. We will show that y = Ax for some $x \in cl(C)$. For every $\epsilon > 0$, the set

$$C_{\epsilon} = \operatorname{cl}(C) \cap \left\{ x \mid ||y - Ax|| \le \epsilon \right\}$$

is closed. Since $A \cdot C \subset A \cdot \operatorname{cl}(C)$ and $y \in \operatorname{cl}(A \cdot C)$, it follows that y is in the closure of $A \cdot \operatorname{cl}(C)$, so that C_{ϵ} is nonempty for every $\epsilon > 0$. Furthermore, the recession cone of the set $\{x \mid \|Ax - y\| \leq \epsilon\}$ coincides with the null space N(A), so that $R_{C_{\epsilon}} = R_{\operatorname{cl}(C)} \cap N(A)$. By assumption we have $R_{\operatorname{cl}(C)} \cap N(A) = \{0\}$, and by part (c) of the Recession Cone Theorem (cf. Prop. 1.4.1), it follows that C_{ϵ} is bounded for every $\epsilon > 0$. Now, since the sets C_{ϵ} are nested nonempty compact sets, their intersection $\bigcap_{\epsilon > 0} C_{\epsilon}$ is nonempty. For any x in this intersection, we have $x \in \operatorname{cl}(C)$ and Ax - y = 0, showing that $y \in A \cdot \operatorname{cl}(C)$. Hence, $\operatorname{cl}(A \cdot C) \subset A \cdot \operatorname{cl}(C)$. The converse $A \cdot \operatorname{cl}(C) \subset \operatorname{cl}(A \cdot C)$ is clear, since for any $x \in \operatorname{cl}(C)$ and sequence $\{x_k\} \subset C$ converging to x, we have $Ax_k \to Ax$, showing that $Ax \in \operatorname{cl}(A \cdot C)$. Therefore,

$$cl(A \cdot C) = A \cdot cl(C). \tag{1.16}$$

We now show that $A \cdot R_{\operatorname{cl}(C)} = R_{A \cdot \operatorname{cl}(C)}$. Let $y \in A \cdot R_{\operatorname{cl}(C)}$. Then, there exists a vector $u \in R_{\operatorname{cl}(C)}$ such that Au = y, and by the definition of $R_{\operatorname{cl}(C)}$, there is a vector $x \in \operatorname{cl}(C)$ such that $x + \alpha u \in \operatorname{cl}(C)$ for every $\alpha \geq 0$. Therefore, $Ax + \alpha Au \in A \cdot \operatorname{cl}(C)$ for every $\alpha \geq 0$, which, together with $Ax \in A \cdot \operatorname{cl}(C)$ and Au = y, implies that y is a direction of recession of the closed set $A \cdot \operatorname{cl}(C)$ [cf. Eq. (1.16)]. Hence, $A \cdot R_{\operatorname{cl}(C)} \subset R_{A \cdot \operatorname{cl}(C)}$.

Conversely, let $y \in R_{A \cdot \operatorname{cl}(C)}$. We will show that $y \in A \cdot R_{\operatorname{cl}(c)}$. This is true if y = 0, so assume that $y \neq 0$. By definition of direction of recession, there is a vector $z \in A \cdot \operatorname{cl}(C)$ such that $z + \alpha y \in A \cdot \operatorname{cl}(C)$ for every $\alpha \geq 0$. Let $x \in \operatorname{cl}(C)$ be such that Ax = z, and for every positive integer k, let $x_k \in \operatorname{cl}(C)$ be such that $Ax_k = z + ky$. Since $y \neq 0$, the sequence $\{Ax_k\}$ is unbounded, implying that $\{x_k\}$ is also unbounded (if $\{x_k\}$ were bounded, then $\{Ax_k\}$ would be bounded, a contradiction). Because $x_k \neq x$ for all k, we can define

$$u_k = \frac{x_k - x}{\|x_k - x\|}, \qquad \forall \ k.$$

Let u be a limit point of $\{u_k\}$, and note that $u \neq 0$. It can be seen that u is a direction of recession of $\operatorname{cl}(C)$ [this can be done similar to the proof of part (c) of the Recession Cone Theorem (cf. Prop. 1.4.1)]. By taking an appropriate subsequence if necessary, we may assume without loss of generality that $\lim_{k\to\infty} u_k = u$. Then, by the choices of u_k and x_k , we have

$$Au = \lim_{k \to \infty} Au_k = \lim_{k \to \infty} \frac{Ax_k - Ax}{\|x_k - x\|} = \lim_{k \to \infty} \frac{k}{\|x_k - x\|} y,$$

implying that $\lim_{k\to\infty}\frac{k}{\|x_k-x\|}$ exists. Denote this limit by λ . If $\lambda=0$, then u is in the null space N(A), implying that $u\in R_{\operatorname{cl}(C)}\cap N(A)$. By the given condition $R_{\operatorname{cl}(C)}\cap N(A)=\{0\}$, we have u=0 contradicting the fact $u\neq 0$. Thus, λ is positive and $Au=\lambda y$, so that $A(u/\lambda)=y$. Since $R_{\operatorname{cl}(C)}$ is a cone [part (a) of the Recession Cone Theorem] and $u\in R_{\operatorname{cl}(C)}$, the vector u/λ

is in $R_{\operatorname{cl}(C)}$, so that y belongs to $A \cdot R_{\operatorname{cl}(C)}$. Hence, $R_{A \cdot \operatorname{cl}(C)} \subset A \cdot R_{\operatorname{cl}(C)}$, completing the proof.

As an example showing that $A \cdot R_{\operatorname{cl}(C)}$ and $R_{A \cdot \operatorname{cl}(C)}$ may differ when $R_{\operatorname{cl}(C)} \cap N(A) \neq \{0\}$, consider the set

$$C = \{(x_1, x_2) \mid x_1 \in \Re, \ x_2 \ge x_1^2\},\$$

and the linear transformation A that maps $(x_1, x_2) \in \mathbb{R}^2$ into $x_1 \in \mathbb{R}$. Then, C is closed and its recession cone is

$$R_C = \{(x_1, x_2) \mid x_1 = 0, \ x_2 \ge 0\},\$$

so that $A \cdot R_C = \{0\}$, where 0 is scalar. On the other hand, $A \cdot C$ coincides with \Re , so that $R_{A \cdot C} = \Re \neq A \cdot R_C$.

1.38

Let C be a nonempty convex subset of \Re^n , and A be an $m \times n$ matrix. Show that if $R_{\operatorname{cl}(C)} \cap N(A)$ is a subspace of the lineality space of $\operatorname{cl}(C)$, then

$$\operatorname{cl}(A \cdot C) = A \cdot \operatorname{cl}(C), \qquad A \cdot R_{\operatorname{cl}(C)} = R_{A \cdot \operatorname{cl}(C)}.$$

Note: This is a refinement of Exercise 1.37.

Solution: Let S be defined by

$$S = R_{\operatorname{cl}(C)} \cap N(A),$$

and note that S is a subspace of $L_{cl(C)}$ by the given assumption. Then, by Lemma 1.4.4, we have

$$\operatorname{cl}(C) = (\operatorname{cl}(C) \cap S^{\perp}) + S,$$

so that the images of $\operatorname{cl}(C)$ and $\operatorname{cl}(C) \cap S^{\perp}$ under A coincide [since $S \subset N(A)$], i.e.,

$$A \cdot \operatorname{cl}(C) = A \cdot \left(\operatorname{cl}(C) \cap S^{\perp}\right). \tag{1.17}$$

Because $A \cdot C \subset A \cdot \operatorname{cl}(C)$, we have

$$\operatorname{cl}(A \cdot C) \subset \operatorname{cl}(A \cdot \operatorname{cl}(C)),$$

which in view of Eq. (1.17) gives

$$\mathrm{cl}(A\cdot C)\subset\mathrm{cl}\Big(A\cdot \left(\mathrm{cl}(C)\cap S^\perp\right)\Big).$$

Define

$$\overline{C} = \operatorname{cl}(C) \cap S^{\perp}$$

so that the preceding relation becomes

$$\operatorname{cl}(A \cdot C) \subset \operatorname{cl}(A \cdot \overline{C}).$$
 (1.18)

The recession cone of \overline{C} is given by

$$R_{\overline{C}} = R_{\operatorname{cl}(C)} \cap S^{\perp}, \tag{1.19}$$

[cf. part (e) of the Recession Cone Theorem, Prop. 1.4.1], for which, since $S = R_{cl(C)} \cap N(A)$, we have

$$R_{\overline{G}} \cap N(A) = S \cap S^{\perp} = \{0\}.$$

Therefore, by Prop. 1.4.13, the set $A \cdot \overline{C}$ is closed, implying that $\operatorname{cl}(A \cdot \overline{C}) = A \cdot \overline{C}$. By the definition of \overline{C} , we have $A \cdot \overline{C} \subset A \cdot \operatorname{cl}(C)$, implying that $\operatorname{cl}(A \cdot \overline{C}) \subset A \cdot \operatorname{cl}(C)$ which together with Eq. (1.18) yields $\operatorname{cl}(A \cdot C) \subset A \cdot \operatorname{cl}(C)$. The converse $A \cdot \operatorname{cl}(C) \subset \operatorname{cl}(A \cdot C)$ is clear, since for any $x \in \operatorname{cl}(C)$ and sequence $\{x_k\} \subset C$ converging to x, we have $Ax_k \to Ax$, showing that $Ax \in \operatorname{cl}(A \cdot C)$. Therefore,

$$cl(A \cdot C) = A \cdot cl(C). \tag{1.20}$$

We next show that $A \cdot R_{\operatorname{cl}(C)} = R_{A \cdot \operatorname{cl}(C)}$. Let $y \in A \cdot R_{\operatorname{cl}(C)}$. Then, there exists a vector $u \in R_{\operatorname{cl}(C)}$ such that Au = y, and by the definition of $R_{\operatorname{cl}(C)}$, there is a vector $x \in \operatorname{cl}(C)$ such that $x + \alpha u \in \operatorname{cl}(C)$ for every $\alpha \geq 0$. Therefore, $Ax + \alpha Au \in \operatorname{Acl}(C)$ for some $x \in \operatorname{cl}(C)$ and for every $\alpha \geq 0$, which together with $Ax \in A \cdot \operatorname{cl}(C)$ and Au = y implies that y is a recession direction of the closed set $A \cdot \operatorname{cl}(C)$ [Eq. (1.20)]. Hence, $A \cdot R_{\operatorname{cl}(C)} \subset R_{A \cdot \operatorname{cl}(C)}$.

Conversely, in view of Eq. (1.17) and the definition of \overline{C} , we have

$$R_{A \cdot \operatorname{cl}(C)} = R_{A \cdot \overline{C}}.$$

Since $R_{\overline{C}} \cap N(A) = \{0\}$ and \overline{C} is closed, by Exercise 1.37, it follows that

$$R_{A \cdot \overline{C}} = A \cdot R_{\overline{C}},$$

which combined with Eq. (1.19) implies that

$$A \cdot R_{\overline{C}} \subset A \cdot R_{\operatorname{cl}(C)}$$
.

The preceding three relations yield $R_{A \cdot \text{cl}(C)} \subset A \cdot R_{\text{cl}(C)}$, completing the proof.

1.39 (Recession Cones of Vector Sums)

This exercise is an extension of Prop. 1.4.14 to nonclosed sets. Let C_1, \ldots, C_m be nonempty convex subsets of \Re^n such that the equality $d_1 + \cdots + d_m = 0$ with $d_i \in R_{\operatorname{cl}(C_i)}$ implies that $d_i \in L_{\operatorname{cl}(C_i)}$ for all i. Then

$$\operatorname{cl}(C_1 + \dots + C_m) = \operatorname{cl}(C_1) + \dots + \operatorname{cl}(C_m),$$

$$R_{\operatorname{cl}(C_1 + \dots + C_m)} = R_{\operatorname{cl}(C_1)} + \dots + R_{\operatorname{cl}(C_m)}.$$

Solution: Let C be the Cartesian product $C_1 \times \cdots \times C_m$. Then,

$$\operatorname{cl}(C) = \operatorname{cl}(C_1) \times \cdots \times \operatorname{cl}(C_m),$$
 (1.21)

and using Exercise 1.38, its recession cone and lineality space are given by

$$R_{\operatorname{cl}(C)} = R_{\operatorname{cl}(C_1)} \times \cdots \times R_{\operatorname{cl}(C_m)}, \tag{1.22}$$

$$L_{\mathrm{cl}(C)} = L_{\mathrm{cl}(C_1)} \times \cdots \times L_{\mathrm{cl}(C_m)}.$$

Let A be a linear transformation that maps $(x_1,\ldots,x_m)\in\Re^{mn}$ into $x_1+\cdots+x_m\in\Re^n$. Then, the intersection $R_{\operatorname{cl}}(C)\cap N(A)$ consists of points (y_1,\ldots,y_m) such that $y_1+\cdots+y_m=0$ with $y_i\in R_{\operatorname{cl}(C_i)}$ for all i. By the given condition, every vector (y_1,\ldots,y_m) in the intersection $R_{\operatorname{cl}(C)}\cap N(A)$ is such that $y_i\in L_{\operatorname{cl}(C_i)}$ for all i, implying that (y_1,\ldots,y_m) belongs to the lineality space $L_{\operatorname{cl}(C)}$. Thus, $R_{\operatorname{cl}(C)}\cap N(A)\subset L_{\operatorname{cl}(C)}\cap N(A)$. On the other hand by definition of the lineality space, we have $L_{\operatorname{cl}(C)}\subset R_{\operatorname{cl}(C)}$, so that $L_{\operatorname{cl}(C)}\cap N(A)\subset R_{\operatorname{cl}(C)}\cap N(A)$. Hence, $R_{\operatorname{cl}(C)}\cap N(A)=L_{\operatorname{cl}(C)}\cap N(A)$, implying that $R_{\operatorname{cl}(C)}\cap N(A)$ is a subspace of $L_{\operatorname{cl}(C)}$. By Exercise 1.38, we have $\operatorname{cl}(A\cdot C)=A\cdot\operatorname{cl}(C)$ and $R_{A\cdot\operatorname{cl}(C)}=A\cdot R_{\operatorname{cl}(C)}$, from which by using the relation $A\cdot C=C_1+\cdots+C_m$, and Eqs. (1.21) and (1.22), we obtain

$$\operatorname{cl}(C_1 + \cdots + C_m) = \operatorname{cl}(C_1) + \cdots + \operatorname{cl}(C_m),$$

$$R_{\operatorname{cl}(C_1+\cdots+C_m)} = R_{\operatorname{cl}(C_1)} + \cdots + R_{\operatorname{cl}(C_m)}.$$

1.40 (Retractiveness of Convex Cones)

(a) Show that a nonpolyhedral closed convex cone need not be retractive, by using as an example the cone $C = \{(u, v, w) \mid ||(u, v)|| \leq w\}$, the recession direction d = (1, 0, 1), and the corresponding asymptotic sequence $\{(k, \sqrt{k}, \sqrt{k^2 + k})\}$. (This is the, so-called, second order cone, which plays an important role in conic programming; see Chapter 5.)

(b) Verify that the cone C of part (a) can be written as the intersection of an infinite number of closed halfspaces, thereby showing that a nested set sequence obtained by intersection of an infinite number of retractive nested set sequences need not be retractive.

Solution: (a) Clearly, d=(1,0,1) is the recession direction associated with the asymptotic sequence $\{x_k\}$, where $x_k=(k,\sqrt{k},\sqrt{k^2+k})$. On the other hand, it can be verified by straightforward calculation that the vector

$$x_k - d = (k - 1, \sqrt{k}, \sqrt{k^2 + k} - 1)$$

does not belong to C. Indeed, denoting

$$u_k = k - 1,$$
 $v_k = \sqrt{k},$ $w_k = \sqrt{k^2 + k} - 1,$

we have

$$||(u_k, v_k)||^2 = (k-1)^2 + k = k^2 - k + 1,$$

while

$$w_k^2 = (\sqrt{k^2 + k} - 1)^2 = k^2 + k + 1 - 2\sqrt{k^2 + k},$$

and it can be seen that

$$||(u_k, v_k)||^2 > w_k^2, \quad \forall k \ge 1.$$

(b) Since by the Schwarz inequality, we have

$$\max_{\|(w,y)\|=1} (uw + vy) = \|(u,v)\|,$$

it follows that the cone

$$C = \{(u, v, w) \mid ||(u, v)|| \le w\}$$

can be written as

$$C = \cap_{\|(w,y)\|=1} \big\{ (u,v,w) \mid uw + vy \le w \big\}.$$

Hence C is the intersection of an infinite number of closed halfspaces.

1.41 (Radon's Theorem)

Let x_1, \ldots, x_m be vectors in \Re^n , where $m \ge n+2$. Show that there exists a partition of the index set $\{1, \ldots, m\}$ into two disjoint sets I and J such that

$$\operatorname{conv}(\{x_i \mid i \in I\}) \cap \operatorname{conv}(\{x_j \mid j \in J\}) \neq \emptyset.$$

As an illustration, show that given four points in the plane, either the (possibly degenerate) triangle formed by three of the points contains the fourth, or else the four points define a (possibly degenerate) quadrilateral. Hint: The system of n+1 equations in the m unknowns $\lambda_1, \ldots, \lambda_m$,

$$\sum_{i=1}^{m} \lambda_i x_i = 0, \qquad \sum_{i=1}^{m} \lambda_i = 0,$$

has a nonzero solution λ^* . Let $I = \{i \mid \lambda_i^* \ge 0\}$ and $J = \{j \mid \lambda_j^* < 0\}$.

Solution: Consider the system of n+1 equations in the m unknowns $\lambda_1, \ldots, \lambda_m$

$$\sum_{i=1}^{m} \lambda_i x_i = 0, \qquad \sum_{i=1}^{m} \lambda_i = 0.$$

Since m > n + 1, there exists a nonzero solution, call it λ^* . Let

$$I=\{i\mid \lambda_i^*\geq 0\}, \qquad J=\{j\mid \lambda_j^*< 0\},$$

and note that I and J are nonempty, and that

$$\sum_{k \in I} \lambda_k^* = \sum_{k \in J} (-\lambda_k^*) > 0.$$

Consider the vector

$$x^* = \sum_{i \in I} \alpha_i x_i,$$

where

$$\alpha_i = \frac{\lambda_i^*}{\sum_{k \in I} \lambda_k^*}, \quad i \in I.$$

In view of the equations $\sum_{i=1}^{m} \lambda_i^* x_i = 0$ and $\sum_{i=1}^{m} \lambda_i^* = 0$, we also have

$$x^* = \sum_{j \in J} \alpha_j x_j,$$

where

$$\alpha_j = \frac{-\lambda_j^*}{\sum_{k \in J} (-\lambda_k^*)}, \quad j \in J.$$

It is seen that the scalars α_i and α_j are nonnegative, and that

$$\sum_{i \in I} \alpha_i = \sum_{j \in J} \alpha_j = 1,$$

so x^* belongs to the intersection

$$\operatorname{conv}(\{x_i \mid i \in I\}) \cap \operatorname{conv}(\{x_j \mid j \in J\}).$$

Given four distinct points in the plane (i.e., m=4 and n=2), Radon's Theorem guarantees the existence of a partition into two subsets, the convex hulls of which intersect. Assuming, there is no subset of three points lying on the same line, there are two possibilities:

- (1) Each set in the partition consists of two points, in which case the convex hulls intesect and define the diagonals of a quadrilateral.
- (2) One set in the partition consists of three points and the other consists of one point, in which case the triangle formed by the three points must contain the fourth.

In the case where three of the points define a line segment on which they lie, and the fourth does not, the triangle formed by the two ends of the line segment and the point outside the line segment form a triangle that contains the fourth point. In the case where all four of the points lie on a line segment, the degenerate triangle formed by three of the points, including the two ends of the line segment, contains the fourth point.

1.42 (Helly's Theorem [Hel21])

Consider a collection S of convex sets in \Re^n , with at least n+1 members, and assume that the intersection of every subcollection of n+1 sets has nonempty intersection.

(a) Assuming that S is a finite collection, show that the entire collection has nonempty intersection. Hint: Use induction. Assume that the conclusion holds for every collection of M sets, where $M \geq n+1$, and show that the conclusion holds for every collection of M+1 sets. In particular, let C_1, \ldots, C_{M+1} be a collection of M+1 convex sets, and consider the collection of M+1 sets B_1, \ldots, B_{M+1} , where

$$B_j = \bigcap_{\substack{i=1,\dots,M+1,\ i\neq j}} C_i, \qquad j = 1,\dots,M+1.$$

Note that, by the induction hypothesis, each set B_j is the intersection of a collection of M sets that have the property that every subcollection of n+1 (or fewer) sets has nonempty intersection. Hence each set B_j is nonempty. Let x_j be a vector in B_j . Apply Radon's Theorem

(Exercise 1.41) to the vectors x_1, \ldots, x_{M+1} . Show that any vector in the intersection of the corresponding convex hulls belongs to the intersection of C_1, \ldots, C_{M+1} .

- (b) Assuming that the members of S are compact, show that the entire collection has nonempty intersection. *Hint*: Use part (a) and the fact that if a collection of compact sets has empty intersection, so does one of its finite subcollections [cf. Prop. A.2.4(i)].
- (c) Use part (a) to show that given a finite family of vertical intervals on the plane every three of which can be intersected by a line, the entire family can be intersected by the same line.

Solution: (a) Consider the induction argument of the hint, let B_j be defined as in the hint, and for each j, let x_j be a vector in B_j . Since $M+1 \geq n+2$, we can apply Radon's Theorem to the vectors x_1, \ldots, x_{M+1} . Thus, there exist nonempty and disjoint index subsets I and J such that $I \cup J = \{1, \ldots, M+1\}$, nonnegative scalars $\alpha_1, \ldots, \alpha_{M+1}$, and a vector x^* such that

$$x^* = \sum_{i \in I} \alpha_i x_i = \sum_{j \in J} \alpha_j x_j, \qquad \sum_{i \in I} \alpha_i = \sum_{j \in J} \alpha_j = 1.$$

It can be seen that for every $i \in I$, a vector in B_i belongs to the intersection $\cap_{j \in J} C_j$. Therefore, since x^* is a convex combination of vectors in B_i , $i \in I$, x^* also belongs to the intersection $\cap_{j \in J} C_j$. Similarly, by reversing the role of I and J, we see that x^* belongs to the intersection $\cap_{i \in I} C_I$. Thus, x^* belongs to the intersection of the entire collection C_1, \ldots, C_{M+1} .

- (b) Evident from the hint.
- (c) Consider a finite family $\{S_i \mid i=1,\ldots,m\}$ of vertical line segments on the plane:

$$S_i = \{(x, y) \mid x = x_i, \underline{y}_i \le y \le \overline{y}_i\}, \qquad i = 1, \dots, m,$$

where $m \geq 3$, and $x_i, \underline{y}_i, \overline{y}_i$, i = 1, ..., m, are given scalars. For each i, consider the set of lines that intersect S_i :

$$C_i = \{(a, b) \mid \underline{y}_i \le ax_i + b \le \overline{y}_i\}.$$

The sets C_i are convex, and every three of them have a common point. By Helly's Theorem, it follows that all the sets C_i have a common point, which is a line that intersects all the intervals S_i .

1.43 (Minimization of Max Functions)

Consider the minimization over \Re^n of the function

$$\max\{f_1(x),\ldots,f_M(x)\},\,$$

where $f_i: \mathbb{R}^n \mapsto (-\infty, \infty]$, $i = 1, \dots, M$, are convex functions, and assume that the optimal value, denoted f^* , is finite. Show that there exists a subset I of $\{1, \dots, M\}$, containing no more than n + 1 indices, such that

$$\inf_{x \in \Re^n} \left\{ \max_{i \in I} f_i(x) \right\} = f^*.$$

Hint: Consider the convex sets $X_i = \{x \mid f_i(x) < f^*\}$, argue by contradiction, and apply Helly's Theorem (Exercise 1.42). Note: The result of this exercise relates to the following question: what is the minimal number of functions f_i that we need to include in the cost function $\max_i f_i(x)$ in order to attain the optimal value f^* ? According to the result, the number is no more than n+1. For applications of this result in structural design and Chebyshev approximation, see Ben Tal and Nemirovski [BeN01].

Solution: Assume the contrary, i.e., that for every index set $I \subset \{1, ..., M\}$, which contains no more than n+1 indices, we have

$$\inf_{x \in \mathbb{R}^n} \left\{ \max_{i \in I} f_i(x) \right\} < f^*.$$

This means that for every such I, the intersection $\cap_{i \in I} X_i$ is nonempty, where

$$X_i = \{ x \mid f_i(x) < f^* \}.$$

From Helly's Theorem, it follows that the entire collection $\{X_i \mid i = 1, ..., M\}$ has nonempty intersection, thereby implying that

$$\inf_{x \in \Re^n} \left\{ \max_{i=1,\dots,M} f_i(x) \right\} < f^*.$$

This contradicts the definition of f^* .

1.44 (Set Intersections and Helly's Theorem)

Show that the conclusions of Prop. 1.4.11(b) hold if the assumption that the sets C_k are nonempty and nested is replaced by the weaker assumption that any subcollection of n+1 (or fewer) sets from the sequence $\{C_k\}$ has nonempty intersection. *Hint*: Consider the sets \overline{C}_k given by

$$\overline{C}_k = \bigcap_{i=0}^k C_i, \qquad k = 1, 2, \dots,$$

and use Helly's Theorem [Exercise 1.42(a)] to show that they are nonempty.

Solution: Helly's Theorem implies that the sets \overline{C}_k defined in the hint are nonempty. These sets are also nested and satisfy the assumptions of Prop. 1.4.11(b). Therefore, the intersection $\bigcap_{i=1}^{\infty} \overline{C}_i$ is nonempty. Since

$$\cap_{i=1}^{\infty} \overline{C}_i \subset \cap_{i=1}^{\infty} C_i,$$

the result follows.

1.45 (Kirchberger's Theorem [Kir1903])

Let S be a finite subset of \Re^n with at least n+2 points, and let $S=B\cup R$ be a partition of S in two disjoint subsets B (the "blue" points) and R (the "red" points). Suppose that every subset \overline{S} of n+2 points of S can be linearly separated, in the sense that there is a vector \overline{a} and a scalar \overline{c} such that $\overline{a}'b+\overline{c}<0$ for all $b\in \overline{S}\cap B$ and $\overline{a}'r+\overline{c}>0$ for all $r\in \overline{S}\cap R$. Use Helly's Theorem (Exercise 1.42) to show that the entire set S can be linearly separated, i.e., that there is a vector S and a scalar S such that S and S are S of or all S and S are S of or all S and S are S can be consider the set S of or vectors S and S are S are S and S are S

$$\sum_{i=1}^{n} x_i b_i + x_{n+1} < 0,$$

and for each $r \in R$, consider the set H(r) of vectors (x_1, \ldots, x_{n+1}) such that

$$\sum_{i=1}^{n} x_i r_i + x_{n+1} > 0.$$

Let C be the collection of the sets G(b) and H(r) as b and r ranges over B and R, respectively. Use Helly's Theorem (Exercise 1.42) to show that C has nonempty intersection.

Solution: For each $b \in B$ consider the set G(b) of vectors (x_1, \ldots, x_{n+1}) such that

$$\sum_{i=1}^{n} x_i b_i + x_{n+1} < 0,$$

and for each $r \in R$, consider the set H(r) of vectors (x_1, \ldots, x_{n+1}) such that

$$\sum_{i=1}^{n} x_i r_i + x_{n+1} > 0.$$

Let \mathcal{C} be the collection of the convex sets G(b) and H(r) as b and r ranges over B and R, respectively. By assumption, for any subset $C \subset B \cup R$, consisting of n+2 points, the sets $B \cap C$ and $R \cap C$ can be linearly separated, so there exist $\overline{a} \in \mathbb{R}^n$ and $\overline{c} \in \mathbb{R}$ such that

$$\overline{a}'b + \overline{c} < 0, \quad \forall b \in B \cap C,$$

$$\overline{a}'r + \overline{c} > 0, \quad \forall \ r \in R \cap C.$$

Thus, $(\overline{a}, \overline{c}) \in L(b)$ for all $b \in B \cap C$, and $(\overline{a}, \overline{c}) \in G(r)$ for all $r \in R \cap C$. It follows that C is a finite family of convex sets in \Re^{n+1} , which contains at least n+2 members and every collection of n+2 of these members has

nonempty intersection. By Helly's Theorem, there is a vector (a,c) that belongs to all members of \mathcal{C} , and for which we have a'x + c < 0 for all $x \in B$ and a'x + c > 0 for all $x \in B$. (Proof given in Webster [Web94], and credited to H. Rademacher and I. J. Shoenberg, "Helly's Theorem on Convex Domains and Tchebycheff's Approximation Problem," Canadian J. of Math., Vol. 2, 1950, pp. 245-256.)

1.46 (Krasnosselsky's Theorem [Kra46])

Let S be a nonempty compact subset of \Re^n . For any two points x and y of S, we say that x is visible from y if the line segment connecting x and y belongs to S. Assume that S has the property that for any subset of n+1 points of S, there is a point of S from which all n+1 points are visible. Show that there is a point in S from which all points of S are visible. Hint: For each $y \in S$, let S_y be the set of points of S that are visible from Y. Show that the set $C_y = \operatorname{conv}(S_y)$ is compact, and consider the family of sets $\{C_y \mid y \in S\}$. Use Helly's Theorem (Exercise 1.42) to show that there is a vector $a \in S$ that belongs to $\cap_{y \in S} \operatorname{conv}(S_y)$. Show that $a \in \cap_{y \in S} S_y$ (this last part is not simple).

Solution: For each $y \in S$, let S_y be the set of points of S that are visible from y. The set S_y is easily seen to be closed, and hence its convex hull, $C_y = \operatorname{conv}(S_y)$, is compact by Prop. 1.2.2. Consider the family of sets $\{C_y \mid y \in S\}$. Let y_0, \ldots, y_n be points in S. By the hypothesis, there is a vector $x \in S$ from which y_0, \ldots, y_n are visible. Thus $x \in S_{y_0} \cap \cdots \cap S_{y_n}$, and hence also $x \in C_{y_0} \cap \cdots \cap C_{y_n}$. It follows that any subcollection of n+1 sets from the family $\{C_y \mid y \in S\}$ is nonempty. By Helly's Theorem [Exercise 1.42(b)], the entire family is nonempty. Thus, there exists a vector a such that

$$a \in \operatorname{conv}(S_y), \quad \forall \ y \in S.$$

We claim now that every $y \in S$ is visible from a. Assume the contrary, so there exists a vector $b \in S$ and a vector c in the line segment connecting a and b such that $c \notin S$. Let C be a closed ball of nonzero radius, which is centered at c and does not intersect S. Let

$$\alpha = \inf\{\lambda \ge 0 \mid S \cap (C + \lambda(b - c)) \ne \emptyset\},\$$

denote the closed ball $C + \alpha(b - c)$ by D and denote its center by d. Then by construction, S meets the boundary of D but not its interior. Let e a vector in $S \cap D$. We will show that $a \notin \text{conv}(S_e)$, thus arriving at a contradiction.

Indeed, consider the halfspaces

$$H^{-} = \{ z \mid (z - e)'(e - d) < 0 \}, \qquad H^{+} = \{ z \mid (z - e)'(e - d) \ge 0 \}.$$

Then, by elementary geometry, it follows that $a \in H^-$, while $S_e \subset H^+$ and hence also $\operatorname{conv}(S_e) \subset H^+$. Since $H^- \cap H^+ = \emptyset$, it follows that $a \notin \operatorname{conv}(S_e)$, a contradiction. (Proof given in Webster [Web94].)

SECTION 1.5: Hyperplanes

1.47

- (a) Let C_1 be a convex set with nonempty interior and C_2 be a nonempty convex set that does not intersect the interior of C_1 . Show that there exists a hyperplane such that one of the associated closed halfspaces contains C_2 , and does not intersect the interior of C_1 .
- (b) Show by an example that we cannot replace interior with relative interior in the statement of part (a).

Solution: (a) In view of the assumption that $int(C_1)$ and C_2 are disjoint and convex [cf Prop. 1.1.1(d)], it follows from the Separating Hyperplane Theorem that there exists a vector $a \neq 0$ such that

$$a'x_1 \le a'x_2$$
, $\forall x_1 \in \text{int}(C_1)$, $\forall x_2 \in C_2$.

Let $b = \inf_{x_2 \in C_2} a'x_2$. Then, from the preceding relation, we have

$$a'x \le b, \quad \forall \ x \in \text{int}(C_1).$$
 (1.23)

We claim that the closed halfspace $\{x \mid a'x \geq b\}$, which contains C_2 , does not intersect int (C_1) .

Assume to arrive at a contradiction that there exists some $\overline{x}_1 \in \operatorname{int}(C_1)$ such that $a'\overline{x}_1 \geq b$. Since $\overline{x}_1 \in \operatorname{int}(C_1)$, we have that there exists some $\epsilon > 0$ such that $\overline{x}_1 + \epsilon a \in \operatorname{int}(C_1)$, and

$$a'(\overline{x}_1 + \epsilon a) \ge b + \epsilon ||a||^2 > b.$$

This contradicts Eq. (1.23). Hence, we have

$$\operatorname{int}(C_1) \subset \{x \mid a'x < b\}.$$

(b) Consider the sets

$$C_1 = \{(x_1, x_2) \mid x_1 = 0\},\$$

$$C_2 = \{(x_1, x_2) \mid x_1 > 0, \ x_2 x_1 \ge 1\}.$$

These two sets are convex and C_2 is disjoint from $ri(C_1)$, which is equal to C_1 . The only separating hyperplane is the x_2 axis, which corresponds to having a = (0, 1), as defined in part (a). For this example, there does not exist a closed halfspace that contains C_2 but is disjoint from $ri(C_1)$.

Let C be a nonempty convex set in \Re^n , and let M be a nonempty affine set in \Re^n . Show that $M \cap \operatorname{ri}(C) = \emptyset$ is a necessary and sufficient condition for the existence of a hyperplane H containing M, and such that $\operatorname{ri}(C)$ is contained in one of the open halfspaces associated with H.

Solution: If there exists a hyperplane H with the properties stated, then $M \cap C = \emptyset$, so the condition $M \cap \operatorname{ri}(C) = \emptyset$ holds. Conversely, if $M \cap \operatorname{ri}(C) = \emptyset$, then M and C can be properly separated by Prop. 1.5.6. This hyperplane can be chosen to contain M since M is affine. If this hyperplane contains a point in $\operatorname{ri}(C)$, then it must contain all of C by Prop. 1.3.4. This contradicts the proper separation property, thus showing that $\operatorname{ri}(C)$ is contained in one of the open halfspaces.

1.49

Let C_1 and C_2 be nonempty convex subsets of \Re^n such that C_2 is a cone.

- (a) Suppose that there exists a hyperplane that separates C_1 and C_2 properly. Show that there exists a hyperplane which separates C_1 and C_2 properly and passes through the origin.
- (b) Suppose that there exists a hyperplane that separates C_1 and C_2 strictly. Show that there exists a hyperplane that passes through the origin such that one of the associated closed halfspaces contains the cone C_2 and does not intersect C_1 .

Solution: (a) If C_1 and C_2 can be separated properly, we have from the Proper Separation Theorem that there exists a vector $a \neq 0$ such that

$$\inf_{x \in C_1} a'x \ge \sup_{x \in C_2} a'x,\tag{1.24}$$

$$\sup_{x \in C_1} a'x > \inf_{x \in C_2} a'x. \tag{1.25}$$

Let

$$b = \sup_{x \in C_2} a'x. \tag{1.26}$$

and consider the hyperplane

$$H = \{x \mid a'x = b\}.$$

Since C_2 is a cone, we have

$$\lambda a'x = a'(\lambda x) \le b < \infty, \quad \forall x \in C_2, \ \forall \ \lambda > 0.$$

This relation implies that $a'x \leq 0$, for all $x \in C_2$, since otherwise it is possible to choose λ large enough and violate the above inequality for some $x \in C_2$. Hence, it follows from Eq. (1.26) that $b \leq 0$. Also, by letting $\lambda \to 0$ in the preceding relation, we see that $b \geq 0$. Therefore, we have that b = 0 and the hyperplane H contains the origin.

(b) If C_1 and C_2 can be separated strictly, we have by definition that there exists a vector $a \neq 0$ and a scalar β such that

$$a'x_2 < \beta < a'x_1, \quad \forall x_1 \in C_1, \quad \forall x_2 \in C_2.$$
 (1.27)

We choose b to be

$$b = \sup_{x \in C_2} a'x,\tag{1.28}$$

and consider the closed halfspace

$$K = \{x \mid a'x \le b\},\$$

which contains C_2 . By Eq. (1.27), we have

$$b \le \beta < a'x$$
, $\forall x \in C_1$,

so the closed halfspace K does not intersect C_1 .

Since C_2 is a cone, an argument similar to the one in part (a) shows that b = 0, and hence the hyperplane associated with the closed halfspace K passes through the origin, and has the desired properties.

1.50 (Separation Properties of Cones)

Define a homogeneous halfspace to be a closed halfspace associated with a hyperplane that passes through the origin. Show that:

- (a) A nonempty closed convex cone is the intersection of the homogeneous halfspaces that contain it.
- (b) The closure of the convex cone generated by a nonempty set X is the intersection of all the homogeneous halfspaces containing X.

Solution: (a) C is contained in the intersection of the homogeneous closed halfspaces that contain C, so we focus on proving the reverse inclusion. Let $x \notin C$. Since C is closed and convex by assumption, by using the Strict Separation Theorem, we see that the sets C and $\{x\}$ can be separated strictly. From Exercise 1.49(c), this implies that there exists a hyperplane that passes through the origin such that one of the associated closed halfspaces contains C, but is disjoint from x. Hence, if $x \notin C$, then x cannot belong to the intersection of the homogeneous closed halfspaces containing C, proving that C contains that intersection.

(b) A homogeneous halfspace is in particular a closed convex cone containing the origin, and such a cone includes X if and only if it includes $\operatorname{cl}(\operatorname{cone}(X))$. Hence, the intersection of all closed homogeneous halfspaces containing X and the intersection of all closed homogeneous halfspaces containing $\operatorname{cl}(\operatorname{cone}(X))$ coincide. From what has been proved in $\operatorname{part}(a)$, the latter intersection is equal to $\operatorname{cl}(\operatorname{cone}(X))$.

1.51 (Strong Separation)

Let C_1 and C_2 be nonempty convex subsets of \Re^n , and let B denote the unit ball in \Re^n , $B = \{x \mid ||x|| \le 1\}$. A hyperplane H is said to separate strongly C_1 and C_2 if there exists an $\epsilon > 0$ such that $C_1 + \epsilon B$ is contained in one of the open halfspaces associated with H and $C_2 + \epsilon B$ is contained in the other. Show that:

- (a) The following three conditions are equivalent.
 - (i) There exists a hyperplane separating strongly C_1 and C_2 .
 - (ii) There exists a vector $a \in \Re^n$ such that $\inf_{x \in C_1} a'x > \sup_{x \in C_2} a'x$.
 - (iii) $\inf_{x_1 \in C_1, x_2 \in C_2} ||x_1 x_2|| > 0$, i.e., $0 \notin \text{cl}(C_2 C_1)$.
- (b) If C_1 and C_2 are disjoint, any one of the five conditions for strict separation, given in Prop. 1.5.3, implies that C_1 and C_2 can be strongly separated.

Solution: (a) We first show that (i) implies (ii). Suppose that C_1 and C_2 can be separated strongly. By definition, this implies that for some nonzero vector $a \in \mathbb{R}^n$, $b \in \mathbb{R}$, and $\epsilon > 0$, we have

$$C_1 + \epsilon B \subset \{x \mid a'x > b\},\$$

$$C_2 + \epsilon B \subset \{x \mid a'x < b\},\$$

where B denotes the closed unit ball. Since $a \neq 0$, we also have

$$\inf\{a'y \mid y \in B\} < 0, \quad \sup\{a'y \mid y \in B\} > 0.$$

Therefore, it follows from the preceding relations that

$$b \le \inf\{a'x + \epsilon a'y \mid x \in C_1, y \in B\} < \inf\{a'x \mid x \in C_1\},\$$

$$b \ge \sup\{a'x + \epsilon a'y \mid x \in C_2, y \in B\} > \sup\{a'x \mid x \in C_2\}.$$

Thus, there exists a vector $a \in \mathbb{R}^n$ such that

$$\inf_{x \in C_1} a'x > \sup_{x \in C_2} a'x,$$

proving (ii).

Next, we show that (ii) implies (iii). Suppose that (ii) holds, i.e., there exists some vector $a \in \Re^n$ such that

$$\inf_{x \in C_1} a'x > \sup_{x \in C_2} a'x,\tag{1.29}$$

Using the Schwartz inequality, we see that

$$0 < \inf_{x \in C_1} a'x - \sup_{x \in C_2} a'x$$

$$= \inf_{x_1 \in C_1, \ x_2 \in C_2} a'(x_1 - x_2),$$

$$\leq \inf_{x_1 \in C_1, \ x_2 \in C_2} ||a|| ||x_1 - x_2||.$$

It follows that

$$\inf_{x_1 \in C_1, x_2 \in C_2} ||x_1 - x_2|| > 0,$$

thus proving (iii).

Finally, we show that (iii) implies (i). If (iii) holds, we have for some $\epsilon>0,$

$$\inf_{x_1 \in C_1, x_2 \in C_2} ||x_1 - x_2|| > 2\epsilon > 0.$$

From this we obtain for all $x_1 \in C_1$, all $x_2 \in C_2$, and for all y_1, y_2 with $||y_1|| \le \epsilon$, $||y_2|| \le \epsilon$,

$$||(x_1+y_1)-(x_2+y_2)|| \ge ||x_1-x_2||-||y_1||-||y_2|| > 0,$$

which implies that $0 \notin (C_1 + \epsilon B) - (C_2 + \epsilon B)$. Therefore, the convex sets $C_1 + \epsilon B$ and $C_2 + \epsilon B$ are disjoint. By the Separating Hyperplane Theorem, we see that $C_1 + \epsilon B$ and $C_2 + \epsilon B$ can be separated, i.e., $C_1 + \epsilon B$ and $C_2 + \epsilon B$ lie in opposite closed halfspaces associated with the hyperplane that separates them. Then, the sets $C_1 + (\epsilon/2)B$ and $C_2 + (\epsilon/2)B$ lie in opposite open halfspaces, which by definition implies that C_1 and C_2 can be separated strongly.

(b) Since C_1 and C_2 are disjoint, we have $0 \notin (C_1 - C_2)$. Any one of conditions (2)-(5) of Prop. 1.5.3 imply condition (1) of that proposition (see the discussion in the proof of Prop. 1.5.3), which states that the set $C_1 - C_2$ is closed, i.e.,

$$cl(C_1 - C_2) = C_1 - C_2.$$

Hence, we have $0 \notin \operatorname{cl}(C_1 - C_2)$, which implies that

$$\inf_{x_1 \in C_1, x_2 \in C_2} ||x_1 - x_2|| > 0.$$

From part (a), it follows that there exists a hyperplane separating C_1 and C_2 strongly.

1.52 (Characterization of Closed Convex Sets)

This exercise generalizes Prop. 1.5.4. Let C be a nonempty closed convex subset of \Re^{n+1} . Show that if C contains no vertical lines, then C is the intersection of the closed halfspaces that contain it and correspond to nonvertical hyperplanes.

Solution: The set C is contained in the intersection of the closed halfspaces that contain C and correspond to nonvertical hyperplanes, so we focus on proving the reverse inclusion. Let $x \notin C$. Since by assumption C does not contain any vertical lines, by Prop. 1.5.8, there exists a closed halfspace that corresponds to a nonvertical hyperplane, contains C, but does not contain x. Hence, if $x \notin C$, then x cannot belong to the intersection of the closed halfspaces containing C and corresponding to nonvertical hyperplanes, proving that C contains that intersection.

SECTION 1.6: Conjugate Functions

1.53 (Logarithmic/Exponential Conjugacy)

Let $f: \Re \mapsto \Re$ be the exponential function

$$f(x) = e^x$$
.

Show that the conjugate is

$$f^{\star}(y) = \begin{cases} y \ln y - y & \text{if } y > 0, \\ 0 & \text{if } y = 0, \\ \infty & \text{if } y < 0. \end{cases}$$

Solution: The conjugate is

$$f^{\star}(y) = \sup_{x \in \Re} \{xy - e^x\}.$$

For y < 0, by taking $x \to -\infty$, we see that $xy - e^x$ can be made arbitrarily large, so $f^*(y) = \infty$. For y = 0, we have

$$f^*(0) = \sup_{x \in \Re} \{-e^x\} = -\inf_{x \in \Re} e^x = 0.$$

Finally, for y > 0, by setting the derivative of $xy - e^x$ to zero, we see that the supremum of $xy - e^x$ is obtained for $x = \ln y$, and by substitution, we obtain $f^*(y) = y \ln y - y$.

1.54 (Conjugates of p-Norms)

Let $f: \Re^n \mapsto \Re$ be the function

$$f(x) = \frac{1}{p} \sum_{i=1}^{n} |x_i|^p$$

where 1 < p. Show that the conjugate is

$$f^{\star}(y) = \frac{1}{q} \sum_{i=1}^{n} |y_i|^q,$$

where q is defined by the relation

$$\frac{1}{p} + \frac{1}{q} = 1.$$

Solution: Consider first the case n = 1. Let x and y be scalars. By setting the derivative of $xy - (1/p)|x|^p$ to zero, and we see that the supremum over x is attained when $\operatorname{sgn}(x)|x|^{p-1} = y$, which implies that $xy = |x|^p$ and $|x|^{p-1} = |y|$. By substitution in the formula for the conjugate, we obtain

$$f^{\star}(y) = |x|^p - \frac{1}{p}|x|^p = \left(1 - \frac{1}{p}\right)|x|^p = \frac{1}{q}|y|^{\frac{p}{p-1}} = \frac{1}{q}|y|^q. \tag{1.30}$$

We now note that for any function $f:\Re^n\mapsto (-\infty,\infty]$ that has the form

$$f(x) = f_1(x_1) + \cdots + f_n(x_n),$$

where $x = (x_1, \dots, x_n)$ and $f_i : \Re \mapsto (-\infty, \infty], i = 1, \dots, n$, the conjugate is given by

$$f^*(y) = f_1^*(y_1) + \dots + f_n^*(y_n),$$

where $f_i^*: \Re \mapsto (-\infty, \infty]$ is the conjugate of f_i , i = 1, ..., n. By combining this fact with the formula (1.30), we obtain the desired result.

1.55 (Conjugate of a Quadratic)

Let

$$f(x) = \frac{1}{2}x'Qx + a'x + b,$$

where Q is a symmetric positive semidefinite $n \times n$ matrix, a is a vector in \Re^n , and b is a scalar. Derive the conjugate of f.

Solution: Let us assume first that Q is nonsingular. Then the maximum of x'y - f(x) over x is attained when $Qx + a = \nabla f(x) = y$. By substitution, we obtain

$$f^*(y) = x'y - f(x) = x'(Qx + a) - \frac{1}{2}x'Qx - a'x - b = \frac{1}{2}x'Qx - b,$$

and finally, using $x = Q^{-1}(y - a)$,

$$f^*(y) = \frac{1}{2}(y-a)'Q^{-1}(y-a) - b.$$

Consider now the general case where Q may be singular. Then if the equation y=Qx+a has no solution, i.e., y-a does not belong to the range R(Q) of Q, we have $f^{\star}(y)=\infty$. Otherwise, let x be the solution of the equation y=Qx+a that has minimum Euclidean norm. Then x is linearly related to y-a and can be written as $x=Q^{\dagger}(y-a)$ where Q^{\dagger} is a symmetric positive semidefinite matrix (this is the *pseudoinverse* of Q; see e.g., Luenberger, Optimization by Vector Space Methods, 1969, p. 165). Similar to the case where Q is invertible, we have $f^{\star}(y)=(1/2)x'Qx-b$, and it follows, using $x=Q^{\dagger}(y-a)$, that

$$f^{\star}(y) = \begin{cases} \frac{1}{2}(y-a)'Q^{\dagger}(y-a) - b & \text{if } y-a \in R(Q), \\ \infty & \text{otherwise.} \end{cases}$$

1.56

(a) Show that if $f_1: \mathbb{R}^n \mapsto (-\infty, \infty]$ and $f_2: \mathbb{R}^n \mapsto (-\infty, \infty]$ are closed proper convex functions, with conjugates denoted by f_1^* and f_2^* , respectively, we have

$$f_1(x) \le f_2(x), \quad \forall \ x \in \Re^n,$$

if and only if

$$f_1^{\star}(y) \ge f_2^{\star}(y), \qquad \forall \ y \in \Re^n.$$

(b) Show that if C_1 and C_2 are nonempty closed convex sets, we have

$$C_1 \subset C_2$$
,

if and only if

$$\sigma_{C_1}(y) \le \sigma_{C_2}(y), \quad \forall y \in \Re^n.$$

Construct an example showing that closedness of C_1 and C_2 is a necessary assumption.

Solution: (a) If $f_1(x) \leq f_2(x)$ for all x, we have for all $y \in \Re^n$,

$$f_1^{\star}(y) = \sup_{x \in \Re^n} \{ x'y - f_1(x) \} \ge \sup_{x \in \Re^n} \{ x'y - f_1(x) \} = f_2^{\star}(y).$$

The reverse implication follows from the fact that f_1 and f_2 are the conjugates of f_1^* and f_2^* , respectively.

(b) Consider the indicator functions δ_{C_1} and δ_{C_2} of C_1 and C_2 . We have

$$C_1 \subset C_2$$
 if and only if $\delta_{C_1}(x) \geq \delta_{C_2}(x)$, $\forall x \in \Re^n$.

Since σ_{C_1} and σ_{C_2} are the conjugates of δ_{C_1} and δ_{C_2} , respectively, the result follows from part (a).

To see that the assumption of closedness of C_1 and C_2 is needed, consider two convex sets that have the same closure, but none of the two is contained in the other, such as for example (0,1] and [0,1).

1.57 (Essentially One-Dimensional Functions)

We say that a closed proper convex function $f: \Re^n \mapsto (-\infty, \infty]$ is essentially one-dimensional if it has the form

$$f(x) = \overline{f}(a'x),$$

where a is a vector in \Re^n and $\overline{f}: \Re \mapsto (-\infty, \infty]$ is a scalar closed proper convex function. We say that a closed proper convex function $f: \Re^n \mapsto (-\infty, \infty]$ is domain one-dimensional if the affine hull of dom(f) is either a single point or a line, i.e.,

$$\operatorname{aff}(\operatorname{dom}(f)) = \{\gamma a + b \mid \gamma \in \Re\},\$$

where a and b are some vectors in \Re^n .

- (a) The conjugate of an essentially one-dimensional function is a domain one-dimensional function such that the affine hull of its domain is a subspace.
- (b) The conjugate of a domain one-dimensional function is the sum of an essentially one-dimensional function and a linear function.

Solution: (a) Let $f: \Re^n \mapsto (-\infty, \infty]$ be essentially one-dimensional, so that

$$f(x)=\overline{f}(a'x),$$

where a is a vector in \Re^n and $\overline{f}: \Re \mapsto (-\infty, \infty]$ is a scalar closed proper convex function. If a = 0, then f is a constant function, so its conjugate

is domain one-dimensional, since its domain is $\{0\}$. We may thus assume that $a \neq 0$. We claim that the conjugate

$$f^{\star}(y) = \sup_{x \in \Re^n} \left\{ y'x - \overline{f}(a'x) \right\},\tag{1.31}$$

takes infinite values if y is outside the one-dimensional subspace spanned by a, implying that f^* is domain one-dimensional with the desired property. Indeed, let y be of the form $y = \xi a + v$, where ξ is a scalar, and v is a nonzero vector with $v \perp a$. If we take $x = \gamma a + \delta v$ in Eq. (1.31), where γ is such that $\gamma ||a||^2 \in \text{dom}(\overline{f})$, we obtain

$$\begin{split} f^{\star}(y) &= \sup_{x \in \Re^n} \left\{ y'x - \overline{f}(a'x) \right\} \\ &\geq \sup_{\delta \in \Re} \left\{ (\xi a + v)'(\gamma a + \delta v) - \overline{f}\left(\gamma \|a\|^2\right) \right\} \\ &= \xi \gamma \|a\|^2 - \overline{f}\left(\gamma \|a\|^2\right) + \sup_{\delta \in \Re} \left\{ \delta \|v\|^2 \right\}, \end{split}$$

so it follows that $f^*(y) = \infty$.

(b) Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be domain one-dimensional, so that

$$\operatorname{aff}(\operatorname{dom}(f)) = \{\gamma a + b \mid \gamma \in \Re\},\$$

for some vectors a and b. If a = b = 0, the domain of f is $\{0\}$, so its conjugate is the function taking the constant value -f(0) and is essentially one-dimensional. If b = 0 and $a \neq 0$, then the conjugate is

$$f^{\star}(y) = \sup_{x \in \Re^n} \left\{ y'x - f(x) \right\} = \sup_{\gamma \in \Re} \left\{ \gamma a'y - f(\gamma a) \right\},$$

so $f^{\star}(y) = \overline{f}^{\star}(a'y)$ where \overline{f}^{\star} is the conjugate of the scalar function $\overline{f}(\gamma) = f(\gamma a)$. Since \overline{f} is closed, convex, and proper, the same is true for \overline{g} , and it follows that f^{\star} is essentially one-dimensional. Finally, consider the case where $b \neq 0$. Then we use a translation argument and write $f(x) = \hat{f}(x-b)$, where \hat{f} is a function such that the affine hull of its domain is the subspace spanned by a. The conjugate of \hat{f} is essentially one-dimensional (by the preceding argument), and the conjugate of f is obtained by adding b'y to it.

1.58

Calculate the support functions of the following sets:

- $(1) C = \{x \mid ||x|| \le 1\}.$
- (2) $C = \{a\} + \gamma \{x \mid ||x|| \le 1\}$, where $a \in \Re^n$ and $\gamma > 0$.

- (3) $C = \{x \mid x_1 + \dots + x_n = 1, x_i \ge 0, i = 1, \dots, n\}.$
- (4) $C = \{x \mid |x_i| \le 1, i = 1, \dots, n\}.$
- (5) $C = \{x \mid \sum_{i=1}^{n} |x_i| \le 1\}.$

Solution: Answers:

- (1) $\sigma_C(y) = ||y||$.
- (2) $\sigma_C(y) = a'y + \gamma ||y||.$
- (3) $\sigma_C(y) = \max_{i=1,...,n} y_i$.
- (4) $\sigma_C(y) = \sum_{i=1}^n |y_i|$.
- (5) $\sigma_C(y) = \max_{i=1,...,n} |y_i|$.

1.59

Show that a subset of \Re^n is bounded if and only if its support function is everywhere finite.

Solution:

1.60

Show that the support function of the unit ball of \Re^n with respect to a norm $\|\cdot\|_a$ is another norm $\|\cdot\|_b$. Verify that if $\|\cdot\|_a$ is the l_1 -norm then $\|\cdot\|_b$ is the l_{∞} -norm, and that if $\|\cdot\|_a$ is the l_{∞} -norm then $\|\cdot\|_b$ is the l_1 -norm.

Solution:

1.61 (Support Function of a Bounded Ellipsoid)

Let X be an ellipsoid of the form

$$X = \{x \mid (x - \overline{x})'Q(x - \overline{x}) \le b\},\$$

where Q is a symmetric positive definite matrix, \overline{x} is a vector, and b is a positive scalar. Show that the support function of X is

$$\sigma_X(y) = y'\overline{x} + (by'Q^{-1}y)^{1/2}, \quad \forall y \in \Re^n.$$

Solution: To calculate $\sigma_X(y)$, we write

$$X = \{\overline{x}\} + \overline{X},$$

where

$$\overline{X} = \{ x \mid x'Qx \le b \},\$$

we calculate the support function $\sigma_{\overline{X}}(y)$, and we use the equation

$$\sigma_X(y) = y'\overline{x} + \sigma_{\overline{X}}(y). \tag{1.32}$$

To calculate

$$\sigma_{\overline{X}}(y) = \sup_{x'Qx \le b} y'x,$$

we introduce the transformation $z = Q^{1/2}x$ and we write

$$\sigma_{\overline{X}}(y) = \sup_{\|z\| \leq b^{1/2}} y'Q^{-1/2}z.$$

It can be seen that for $y \neq 0$, the supremum over z above is attained at

$$z(y) = b^{1/2} \frac{Q^{-1/2}y}{\|Q^{-1/2}y\|},$$

and by substitution in the expression for $\sigma_{\overline{X}}(y)$, we have

$$\sigma_{\overline{X}}(y) = (b \, y' Q^{-1} y)^{1/2}$$
.

Thus, using Eq. (1.32), we finally obtain

$$\sigma_X(y) = y'\overline{x} + (by'Q^{-1}y)^{1/2}, \quad \forall y \in \Re^n.$$

1.62 (Support Function of Sum and Union)

Let X_1, \ldots, X_r , be nonempty subsets of \Re^n . Derive formulas for $X_1 + \cdots + X_r$, $\operatorname{conv}(X_1) + \cdots + \operatorname{conv}(X_r)$, $\bigcup_{i=1}^r X_i$, and $\operatorname{conv}\left(\bigcup_{i=1}^r X_i\right)$.

Solution: Let $X = X_1 + \cdots + X_r$. We have for all $y \in \Re^n$,

$$\sigma_X(y) = \sup_{x \in X_1 + \dots + X_r} x'y$$

$$= \sup_{x_1 \in X_1, \dots, x_r \in X_r} (x_1 + \dots + x_r)'y$$

$$= \sup_{x_1 \in X_1} x'_1 y + \dots + \sup_{x_r \in X_r} x'_r y$$

$$= \sigma_{X_1}(y) + \dots + \sigma_{X_r}(y).$$

Since X_j and $conv(X_j)$ have the same support function, it follows that

$$\sigma_{X_1}(y) + \cdots + \sigma_{X_r}(y)$$

is also the support function of

$$\operatorname{conv}(X_1) + \cdots + \operatorname{conv}(X_r).$$

Let also $X = \bigcup_{j=1}^r X_j$. We have

$$\sigma_X(y) = \sup_{x \in X} y'x = \max_{j=1,\dots,r} \sup_{x \in X_j} y'x = \max_{j=1,\dots,r} \sigma_{X_j}(y).$$

This is also the support function of conv $(\bigcup_{j=1}^r X_j)$.

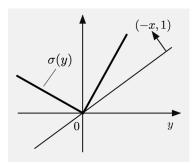


Figure 1.1. Visualization of the conjugate of a positively homogeneous function σ (cf. Exercise 1.63). The values of its conjugate correspond to crossing levels of the vertical axis by hyperplanes supporting the epigraph of σ . There can be only one such level, namely 0, so the conjugate of σ takes only the values 0 and ∞ , i.e., it is the indicator function of the set whose elements x correspond to normals (-x, 1) of hyperplanes supporting epi(f) at 0, as shown in the figure.

1.63 (Positively Homogeneous Functions and Support Functions)

A function $f: \Re^n \mapsto [-\infty, \infty]$ is called *positively homogeneous* if its epigraph is a cone in \Re^{n+1} . Equivalently, f is positively homogeneous if and only if

$$f(\gamma x) = \gamma f(x), \quad \forall \ \gamma > 0, \ \forall \ x \in \Re^n.$$

(a) Show that if f is a proper convex positively homogeneous function, then for all $x_1, \ldots, x_m \in \mathbb{R}^n$ and $\gamma_1, \ldots, \gamma_m > 0$, we have

$$f(\gamma_1 x_1 + \dots + \gamma_m x_m) \le \gamma_1 f(x_1) + \dots + \gamma_m f(x_m).$$

(b) Clearly, the support function σ_X of a nonempty set X is closed proper convex and positively homogeneous. Show that the reverse is also true, namely that the closure of any proper convex positively homogeneous function is the support function of some set. In particular, show that if $\sigma: \Re^n \mapsto (-\infty, \infty]$ is a proper convex positively homogeneous function, then the conjugate of σ is the indicator function of the closed convex set

$$X = \big\{ x \mid y'x \le \sigma(y), \, \forall \, \, y \in \Re^n \big\},\,$$

and $\operatorname{cl}\sigma$ is the support function of X (see Fig. 1.1 for a geometric interpretation of this result).

Solution: (a) Let

$$\overline{\gamma} = \gamma_1 + \dots + \gamma_m,$$

and

$$\overline{\gamma}_i = \frac{\gamma_i}{\overline{\gamma}}, \qquad i = 1, \dots, m.$$

By convexity of f, we have

$$f(\overline{\gamma}_1 x_1 + \dots + \overline{\gamma}_m x_m) \leq \overline{\gamma}_1 f(x_1) + \dots + \overline{\gamma}_m f(x_m).$$

Since f is positively homogeneous, this inequality can be written as

$$\frac{1}{\overline{\gamma}}f(\gamma_1x_1+\cdots+\gamma_mx_m)\leq \frac{1}{\overline{\gamma}}(\gamma_1f(x_1)+\cdots+\gamma_mf(x_m)),$$

and the result follows.

(b) Let δ be the conjugate of σ :

$$\delta(x) = \sup_{y \in \Re^n} \{ y'x - \sigma(y) \}.$$

Since σ is positively homogeneous, we have for any $\gamma > 0$,

$$\gamma\,\delta(x) = \sup_{y\in\Re^n} \bigl\{\gamma y' x - \gamma\,\sigma(y)\bigr\} = \sup_{y\in\Re^n} \bigl\{(\gamma y)' x - \sigma(\gamma y)\bigr\}.$$

The right-hand sides of the preceding two relations are equal, so we obtain

$$\delta(x) = \gamma \, \delta(x), \quad \forall \, \gamma > 0,$$

which implies that δ takes only the values 0 and ∞ (since σ and hence also its conjugate δ is proper). Thus, δ is the indicator function of a set, call it X, and we have

$$\begin{split} X &= \left\{ x \mid \delta(x) \leq 0 \right\} \\ &= \left\{ x \mid \sup_{y \in \Re^n} \left\{ y'x - \sigma(y) \right\} \leq 0 \right\} \\ &= \left\{ x \mid y'x \leq \sigma(y), \, \forall \, \, y \in \Re^n \right\}. \end{split}$$

Finally, since δ is the conjugate of σ , we see that $\operatorname{cl} \sigma$ is the conjugate of δ ; cf. the Conjugacy Theorem [Prop. 1.6.1(c)]. Since δ is the indicator function of X, it follows that $\operatorname{cl} \sigma$ is the support function of X.

1.64 (Support Function of Domain)

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let f^* be its conjugate.

- (a) The support function of dom(f) is the recession function of f^* .
- (b) If f is closed, the support function of $dom(f^*)$ is the recession function of f.

Solution: (a) From the definition

$$f^{\star}(y) = \sup_{x \in \Re^n} \{x'y - f(x)\},\,$$

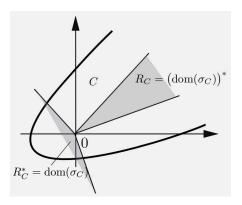


Figure 1.2. Illustration of the polarity of the recession cone R_C of a closed convex set C and the domain $dom(\sigma_C)$ of its support function (cf. Exercise 1.65). For any $y \notin R_C^*$ we have

$$\sup_{x \in C} y'x = \infty,$$

while for any $y \in R_C^*$, we must have $\sup_{x \in C} y'x < \infty$.

we see that f^* is the pointwise supremum of the affine functions

$$h_{(x,w)}(y) = x'y - w,$$

as (x, w) ranges over $\operatorname{epi}(f)$. Therefore, $\operatorname{epi}(f^*)$ is the intersection of the epigraphs of $h_{(x,w)}$ as (x,w) ranges over $\operatorname{epi}(f)$. Hence, by the Recession Cone Theorem [Prop. 1.4.1(e)], the recession cone of $\operatorname{epi}(f^*)$ is the intersection of the recession cones of the epigraphs of $h_{(x,w)}$ as (x,w) ranges over $\operatorname{epi}(f)$. Since the epigraph of $h_{(x,w)}$ is $\{(y,u) \mid x'y-w \leq u\}$, its recession cone is $\{(y,u) \mid x'y \leq u\}$, and we have

$$R_{\operatorname{epi}(f^{\star})} = \bigcap_{x \in \operatorname{dom}(f)} \{(y, u) \mid x'y \leq u\}.$$

Since $R_{\text{epi}(f^*)}$ is the epigraph of the recession function r_f^* of f^* , it follows that

$$r_f^{\star}(y) = \sup_{x \in \text{dom}(f)} x'y,$$

so r_f^{\star} is the support function of dom(f).

(b) If f is closed, then by the Conjugacy Theorem [Prop. 1.6.1(c)], it is the conjugate of f^* , and the result follows by applying part (a) with the roles of f and f^* interchanged.

1.65 (Polarity of Recession Cone and Domain of Support Function)

Let C be a nonempty closed convex set in \Re^n . Then, the recession cone of C is equal to the polar cone of the domain of σ_C :

$$R_C = (\operatorname{dom}(\sigma_C))^*.$$

Solution: We apply the result of Exercise 1.64(b) with f and f^* equal to the indicator function δ_C and support function σ_C of C, respectively. We obtain

$$r_{\delta_C} = \sigma_{\text{dom}(\sigma_C)},$$

from which, by using also the formula of Example 1.6.2,

$$R_C = \left\{ y \mid r_{\delta_C}(y) \le 0 \right\} = \left\{ y \mid \sigma_{\operatorname{dom}(\sigma_C)}(y) \le 0 \right\} = \left(\operatorname{dom}(\sigma_C)\right)^*.$$

1.66 (Generated Functions - Support Function of 0-Level Set)

Given a proper convex function $f: \Re^n \mapsto (-\infty, \infty]$, the closure of the cone generated by $\operatorname{epi}(f)$, is the epigraph of a closed convex positively homogeneous function, called the *closed function generated by* f, and denoted by $\operatorname{gen} f$. The epigraph of $\operatorname{gen} f$ is the intersection of all the halfspaces that contain $\operatorname{epi}(f)$ and contain 0 in their boundary. Alternatively, the epigraph of $\operatorname{gen} f$ is the intersection of all the closed cones that contain $\operatorname{epi}(f)$.

- (a) Show that if the level set $\{y \mid h(y) \leq 0\}$ is nonempty, the generated function gen f is proper.
- (b) Let $f: \Re^n \mapsto (-\infty, \infty]$ be a closed proper convex function and let f^* be its conjugate. Show that if the level set $\{y \mid f^*(y) \leq 0\}$ [respectively $\{x \mid f(x) \leq 0\}$] is nonempty, its support function is the closed function generated by f (respectively f^*).

Solution: (a) If $(\operatorname{cl} f)(0) > 0$, by the Nonvertical Hyperplane Theorem (Prop. 1.5.8), there exists a nonvertical hyperplane passing through the origin and containing $\operatorname{epi}(f)$ in one of its closed halfspaces, implying that the epigraph of $\operatorname{gen} f$ does not contain a line, so $\operatorname{gen} f$ is proper. Any y such that $f^*(y) \leq 0$, or equivalently $y'x \leq f(x)$ for all x, defines a nonvertical hyperplane that separates the origin from $\operatorname{epi}(f)$.

(b) Let σ be the closed function generated by f. Then, since $\{y \mid f^*(y) \leq 0\} \neq \emptyset$, by part (a), σ is proper, and by Exercise 1.63, σ is the support function of the set

$$Y = \{ y \mid y'x \le \sigma(x), \, \forall \, x \in \Re^n \}.$$

Since $epi(\sigma)$ is the intersection of all the halfspaces that contain epi(f) and contain 0 in their boundary, the set Y can be written as

$$Y = \{ y \mid y'x \le f(x), \, \forall \, x \in \Re^n \},\$$

from which we obtain

$$Y = \left\{ y \ \middle| \ \sup_{x \in \Re^n} \left\{ y'x - f(x) \right\} \le 0 \right\} = \left\{ y \ \middle| \ f^\star(y) \le 0 \right\}.$$

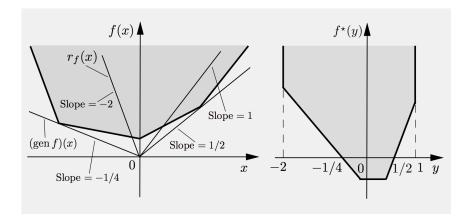


Figure 1.3. Illustration of Exercise 1.66. The recession function r_f and the closed generated function gen f are the support functions of the sets $\mathrm{dom}(f^\star)$ and $\left\{y\mid f^\star(y)\leq 0\right\}$, respectively.

Note that the method used to characterize the 0-level set of f can be applied to any level set. In particular, a nonempty level set $L_{\gamma} = \{x \mid f(x) \leq \gamma\}$ is the 0-level set of the function f_{γ} defined by $f_{\gamma}(x) = f(x) - \gamma$, and its support function is the closed function generated by h_{γ} , the conjugate of f_{γ} , which is given by $h_{\gamma}(y) = h(y) + \gamma$. The preceding analysis is illustrated in Fig. 1.3.

1.67 (Conjugates Involving Invertible Linear Transformations)

Let $p: \Re^n \mapsto [-\infty, \infty]$ be a convex function, let A be an invertible $n \times n$ matrix, let a and c be vectors in \Re^n , and let b be a scalar. Calculate the conjugate of the convex function

$$f(x) = p(A(x-c)) + a'x + b.$$

Solution: Using the transformation z = A(x - c), we can write the conjugate as

$$\begin{split} f^{\star}(y) &= \sup_{x \in \Re^n} \left\{ x'y - p\big(A(x-c)\big) - a'x - b \right\} \\ &= \sup_{z \in \Re^n} \left\{ (A^{-1}z + c)'y - p(z) - (A^{-1}z + c)'a - b \right\} \\ &= \sup_{z \in \Re^n} \left\{ (A^{-1}z)'(y-a) - p(z) \right\} + c'(y-a) - b \end{split}$$

and finally

$$f^*(y) = p^*((A')^{-1}(y-a)) + c'y + d,$$

where p^* is the conjugate of p and

$$d = -(c'a + b).$$

Note the symmetry between f and f^* .

1.68 (Conjugate of a Partially Affine Function)

A partially affine function f is defined as a function such that dom(f) is an affine set, and f is affine on dom(f). Show that the conjugate of a partially affine function is another partially affine function. Hint: Write f as

$$f(x) = \delta_{\text{aff}(f)}(x) + a'x + b$$

and apply the result of Exercise 1.67.

Solution: See p. 107 of [Roc70]. Write f as

$$f(x) = \delta_{\text{aff}(f)}(x) + a'x + b$$

and apply the result of Exercise 1.67.

1.69

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function and let f^* be its conjugate. Show that:

- (a) The constancy space of f^* is the orthogonal complement of the subspace parallel to aff(dom(f)). *Hint*: Use Exercise 1.64.
- (b) If f is closed, the subspace parallel to aff(dom(h)) is the orthogonal complement of the constancy space of f.

Solution: *Hint*: Use Exercise 1.64. See Theorem 13.3 of [Roc70].

1.70

Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a closed proper convex function, and let f^* be its conjugate. Show that dom(f) is an affine set if and only if the recession function of f^* satisfies $r_h(d) = \infty$ for all d that are not in the constancy space of f. Hint: Use Exercise 1.64.

Solution: *Hint*: Use Exercise 1.64. See Theorem 13.3 of [Roc70].

1.71 (Co-finite Functions)

A closed proper convex function $f: \Re^n \mapsto (-\infty, \infty]$ is said to be co-finite if its recession function is

$$r_f(d) = \begin{cases} 0 & \text{if } d = 0, \\ \infty & \text{if } d \neq 0. \end{cases}$$

Show that f is co-finite if and only if its conjugate is real-valued. *Hint*: Use Exercise 1.64.

Solution: See Theorem 13.3 of [Roc70].

1.72 (Lipschitz Continuity and Domain Boundedness)

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let f^* be its conjugate. Show that $\text{dom}(f^*)$ is bounded if and only if f is real-valued and there exists a scalar $L \geq 0$ such that

$$|f(x) - f(y)| \le L||x - y||, \quad \forall x \in \Re^n, y \in \Re^n.$$

Furthermore, the minimal value of L for which this Lipschitz condition holds is

$$\sup_{y\in \mathrm{dom}(f^\star)}\|y\|.$$

Hint: Use Exercise 1.64.

Solution: See Theorem 13.3, Corollary 13.3.3 of [Roc70].

1.73

Let C be a nonempty convex subset of \Re^n and let σ_C be its support function.

- (a) $x \in cl(C)$ if and only if $y'x \leq \sigma_C(y)$ for all $y \in \Re^n$.
- (b) $x \in ri(C)$ if and only if $y'x \le \sigma_C(y)$ for all $y \in \mathbb{R}^n$ and $y'x < \sigma_C(y)$ for all $y \in \mathbb{R}^n$ with $-\sigma_C(-y) \ne \sigma_C(y)$.
- (c) x is an interior point of C if and only if $y'x < \sigma_C(y)$ for all $y \neq 0$.

Solution: See Theorem 13.1 of [Roc70].

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Convex Optimization Theory

Chapter 2 Exercises and Solutions

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CHAPTER 2: EXERCISES AND SOLUTIONS

SECTION 2.1: Extreme Points

2.1

Show by example that the set of extreme points of a nonempty compact set need not be closed. *Hint*: Consider a line segment $C_1 = \{(x_1, x_2, x_3) \mid x_1 = 0, x_2 = 0, -1 \leq x_3 \leq 1\}$ and a circular disk $C_2 = \{(x_1, x_2, x_3) \mid (x_1 - 1)^2 + x_2^2 \leq 1, x_3 = 0\}$, and verify that the set $\operatorname{conv}(C_1 \cup C_2)$ is compact, while its set of extreme points is not closed.

Solution: For the sets C_1 and C_2 as given in the hint, the set $C_1 \cup C_2$ is compact, and its convex hull is also compact by Prop. 1.2.2. The set of extreme points of $\operatorname{conv}(C_1 \cup C_2)$ is not closed, since it consists of the two end points of the line segment C_1 , namely (0,0,-1) and (0,0,1), and all the points $x = (x_1, x_2, x_3)$ such that

$$x \neq 0$$
, $(x_1 - 1)^2 + x_2^2 = 1$, $x_3 = 0$.

2.2 (Krein-Milman Theorem)

Show that a convex and compact subset of \Re^n is equal to the convex hull of its extreme points.

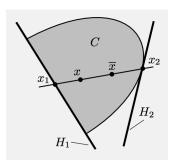


Figure 2.1. Construction used in the induction proof of the Krein-Milman Theorem (Exercise 2.2): any vector x of a convex and compact set C can be represented as a convex combination of extreme points of C. If \overline{x} is another point in C, the points x_1 and x_2 shown can be represented as convex combinations of extreme points of the lower dimensional convex and compact sets $C \cap H_1$ and $C \cap H_2$, which are also extreme points of C by Prop. 2.1.1.

Solution: By convexity, C contains the convex hull of its extreme points. To show the reverse inclusion, we use induction on the dimension of the space. On the real line, a compact convex set C is a line segment whose

endpoints are the extreme points of C, so every point in C is a convex combination of the two endpoints. Suppose now that every vector in a compact and convex subset of \Re^{n-1} can be represented as a convex combination of extreme points of the set. We will show that the same is true for compact and convex subsets of \Re^n .

Let C be a compact and convex subset of \Re^n , and choose any $x \in C$. If x is the only point in C, it is an extreme point and we are done, so assume that \overline{x} is another point in C, and consider the line that passes through x and \overline{x} . Since C is compact, the intersection of this line and C is a compact line segment whose endpoints, say x_1 and x_2 , belong to the relative boundary of C. Let H_1 be a hyperplane that passes through x_1 and contains C in one of its closed halfspaces. Similarly, let H_2 be a hyperplane that passes through x_2 and contains C in one of its closed halfspaces (see Fig. 2.1). The intersections $C \cap H_1$ and $C \cap H_2$ are compact convex sets that lie in the hyperplanes H_1 and H_2 , respectively. By viewing H_1 and H_2 as (n-1)-dimensional spaces, and by using the induction hypothesis, we see that each of the sets $C \cap H_1$ and $C \cap H_2$ is the convex hull of its extreme points. Hence, x_1 is a convex combination of some extreme points of $C \cap H_1$, and x_2 is a convex combination of some extreme points of $C \cap H_2$. By Prop. 2.1.1, all the extreme points of $C \cap H_1$ and all the extreme points of $C \cap H_2$ are also extreme points of C, so both x_1 and x_2 are convex combinations of some extreme points of C. Since x lies in the line segment connecting x_1 and x_2 , it follows that x is a convex combination of some extreme points of C, showing that C is contained in the convex hull of the extreme points of C.

2.3

Let C be a nonempty convex subset of \Re^n , and let A be an $m \times n$ matrix with linearly independent columns. Show that a vector $x \in C$ is an extreme point of C if and only if Ax is an extreme point of the image AC. Show by example that if the columns of A are linearly dependent, then Ax can be an extreme point of AC, for some non-extreme point x of C.

Solution: Suppose that x is not an extreme point of C. Then $x = \alpha x_1 + (1 - \alpha)x_2$ for some $x_1, x_2 \in C$ with $x_1 \neq x$ and $x_2 \neq x$, and a scalar $\alpha \in (0,1)$, so that $Ax = \alpha Ax_1 + (1 - \alpha)Ax_2$. Since the columns of A are linearly independent, we have $Ay_1 = Ay_2$ if and only if $y_1 = y_2$. Therefore, $Ax_1 \neq Ax$ and $Ax_2 \neq Ax$, implying that Ax is a convex combination of two distinct points in AC, i.e., Ax is not an extreme point of AC.

Suppose now that Ax is not an extreme point of AC, so that $Ax = \alpha Ax_1 + (1 - \alpha)Ax_2$ for some $x_1, x_2 \in C$ with $Ax_1 \neq Ax$ and $Ax_2 \neq Ax$, and a scalar $\alpha \in (0,1)$. Then, $A(x - \alpha x_1 - (1 - \alpha)x_2) = 0$ and since the columns of A are linearly independent, it follows that $x = \alpha x_1 - (1 - \alpha)x_2$.

Furthermore, because $Ax_1 \neq Ax$ and $Ax_2 \neq Ax$, we must have $x_1 \neq x$ and $x_2 \neq x$, implying that x is not an extreme point of C.

As an example showing that if the columns of A are linearly dependent, then Ax can be an extreme point of AC, for some non-extreme point x of C, consider the 1×2 matrix $A = [1 \ 0]$, whose columns are linearly dependent. The polyhedral set C given by

$$C = \{(x_1, x_2) \mid x_1 \ge 0, \ 0 \le x_2 \le 1\}$$

has two extreme points, (0,0) and (0,1). Its image $AC \subset \Re$ is given by

$$AC = \{x_1 \mid x_1 \ge 0\},\$$

whose unique extreme point is $x_1 = 0$. The point $x = (0, 1/2) \in C$ is not an extreme point of C, while its image Ax = 0 is an extreme point of AC. Actually, all the points in C on the line segment connecting (0,0) and (0,1), except for (0,0) and (0,1), are non-extreme points of C that are mapped under A into the extreme point 0 of AC.

2.4

Let C be a nonempty closed convex subset of \Re^n . Show that the following are equivalent.

- (i) All boundary points of C are extreme points of C.
- (ii) Every hyperplane that supports C at some point intersects C only at that point.
- (iii) Every line intersects the boundary of C at no more than two points.

Solution: The result is clearly true if C consists of a single point, so assume that C consists of more than one point.

We first show that (i) implies (ii). Assume that all boundary points of C are extreme points. If there is a hyperplane that supports C and intersects C at two distinct points, the entire line segment connecting the two points would lie on the boundary of C, but the midpoint of this line segment would not be an extreme point - a contradiction.

Next we show that (ii) implies (iii). Assume that every hyperplane that supports C at some point intersects C only at that point. Suppose that there is a line that intersects the boundary of C at three distinct boundary points x_1, x_2, x_3 , with x_2 being the midpoint. Consider a hyperplane H that supports C at x_2 , i.e., a vector $a \neq 0$ such that

$$a'x \ge a'x_2, \quad \forall \ x \in C.$$

Then since by the hypothesis, H intersects C only at x_2 , we must have $a'x_1 > a'x_2$ and $a'x_3 > a'x_2$, which is a contradiction since x_2 lies strictly between x_1 and x_3 .

Finally, we show that (iii) implies (i). Assume that every line intersects the boundary of C at no more than two points. If there is a boundary point x_2 that is not extreme and therefore lies strictly between two points $x_1, x_3 \in C$, then either x_1 or x_3 must be an interior point, for otherwise the line that passes through x_1, x_2, x_3 would contain more than two boundary points. Thus, by the Line Segment Principle (Prop. 1.3.1), every point that lies strictly between x_1 and x_3 , including x_2 , is an interior point of C. This contradicts the hypothesis that x_2 is a boundary point of C.

2.5 (Matrix Inequalities)

Let A be a symmetric $n \times n$ matrix with components denoted a_{ij} and eigenvalues denoted $\lambda_1, \ldots, \lambda_n$, and let Λ_A be the set of all vectors of \Re^n obtained by permutations of these eigenvalues.

(a) Let C be a convex set that contains Λ_A , and let $f: C \mapsto \Re^n$ be a convex function. Show that for any orthonormal set of vectors v_1, \ldots, v_n in \Re^n , we have

$$f(v_1'Av_1,\ldots,v_n'Av_n) \le \max_{(\xi_1,\ldots,\xi_n)\in\Lambda_A} f(\xi_1,\ldots,\xi_n).$$

Hint: Let S be the doubly stochastic matrix with components $s_{ij} = (v'_i u_j)^2$, where u_1, \ldots, u_n are orthonormal eigenvectors corresponding to the eigenvalues $\lambda_1, \ldots, \lambda_n$. Show that $v = S\lambda$, where

$$v = (v_1' A v_1, \dots, v_n' A v_n), \qquad \lambda = (\lambda_1, \dots, \lambda_n),$$

and use the Birkhoff-von Neumann Theorem.

(b) Let A be positive semidefinite. Show that for any orthonormal set of vectors v_1, \ldots, v_n in \Re^n , we have

$$\det A = \lambda_1 \cdots \lambda_n \le v_1' A v_1 \cdots v_n' A v_n.$$

Furthermore, the inequality is sharp in the sense that it is satisfied as an equality for some orthonormal set of vectors. *Hint*: Use part (a) with $f(x_1, \ldots, x_n) = -(x_1 \cdots x_n)^{1/2}$, and C equal to the nonnegative orthant.

(c) (Hadamard's Determinant Inequality) We have

$$(\det A)^2 \le (a_{11}^2 + \dots + a_{n1}^2) \dots (a_{1n}^2 + \dots + a_{nn}^2).$$

Furthermore, if in addition A is positive semidefinite, we have

$$\det A \leq a_{11} \cdots a_{nn}$$
.

Solution: (a) Let u_1, \ldots, u_n be orthonormal eigenvectors corresponding to the eigenvalues $\lambda_1, \ldots, \lambda_n$. The orthogonality of u_1, \ldots, u_n implies that

$$v_i = (v_i'u_1)u_1 + \dots + (v_i'u_n)u_n, \qquad i = 1, \dots, n.$$

Using this relation, it is straightforward to verify that

$$v = S\lambda$$
,

where

$$v = (v_1' A v_1, \dots, v_n' A v_n), \qquad \lambda = (\lambda_1, \dots, \lambda_n),$$

and S is the $n \times n$ matrix with components $s_{ij} = (v'_i u_j)^2$ for all i and j. We now note that S is a doubly stochastic matrix. The reason is that we have for each i,

$$||v_i||^2 = ||(v_i'u_1)u_1 + \dots + (v_i'u_n)u_n||^2,$$

so that by using the orthonormality of u_1, \ldots, u_n , we have

$$||v_i||^2 = (v_i'u_1)^2 + \dots + (v_i'u_n)^2.$$

This implies that S is doubly stochastic, since $||v_i|| = 1$ by assumption, and the *i*th row of the matrix S is $((v_i'u_1)^2, \ldots, (v_i'u_n)^2)$.

The Birkhoff-von Neumann Theorem asserts that S can be expressed as a convex combination of permutation matrices, i.e., there exist $\mu_j \geq 0$, $j = 1, \ldots, m$, with $\sum_{j=1}^{m} \mu_j = 1$, and such that

$$S = \mu_1 P_1 + \dots + \mu_m P_m,$$

where P_1, \ldots, P_m are permutation matrices. Hence,

$$v = S\lambda = \mu_1(P_1\lambda) + \cdots + \mu_m(P_m\lambda).$$

Since the vectors $P_j\lambda$, $j=1,\ldots,m$, belong to Λ_A , they also belong to C. Since v is a convex combination of $P_j\lambda$, $j=1,\ldots,m$, it follows that $v \in C$. Thus, using the convexity of f, we have

$$f(v) \le \mu_1 f(P_1 \lambda) + \dots + \mu_m f(P_m \lambda) \le \max_{(\xi_1, \dots, \xi_n) \in \Lambda_A} f(\xi_1, \dots, \xi_n).$$

- (b) The inequality follows from part (a) and the hint. The inequality is satisfied as an equality if the vectors v_1, \ldots, v_n are normalized eigenvectors corresponding to $\lambda_1, \ldots, \lambda_n$.
- (c) Let B = A'A. We apply part (b) to B with the orthonormal vectors being the unit vectors e_1, \ldots, e_n of \Re^n . We obtain

$$\det B \le e_1' B e_1 \cdots e_n' B e_n = (a_{11}^2 + \cdots + a_{n1}^2) \cdots (a_{1n}^2 + \cdots + a_{nn}^2),$$

where the last equality can be verified by straightforward calculation. Since $\det B = (\det A)^2$, the desired inequality follows.

If A positive semidefinite, we apply part (b) to A with the orthonormal vectors being the unit vectors e_1, \ldots, e_n of \Re^n , to obtain

$$\det A \leq a_{11} \cdots a_{nn}$$
.

2.6 (Faces)

Let P be a polyhedral set. For any hyperplane H that passes through a boundary point of P and contains P in one of its halfspaces, we say that the set $F = P \cap H$ is a *face* of P. Show the following:

- (a) Each face is a polyhedral set.
- (b) Each extreme point of P, viewed as a singleton set, is a face.
- (c) If P is not an affine set, there is a face of P whose dimension is $\dim(P) 1$.
- (d) The number of distinct faces of P is finite.

Solution: (a) Let P be a polyhedral set in \Re^n , and let $F = P \cap H$ be a face of P, where H is a hyperplane passing through some boundary point \overline{x} of P and containing P in one of its halfspaces. Then H is given by $H = \{x \mid a'x = a'\overline{x}\}$ for some nonzero vector $a \in \Re^n$. By replacing $a'x = a'\overline{x}$ with two inequalities $a'x \leq a'\overline{x}$ and $-a'x \leq -a'\overline{x}$, we see that H is a polyhedral set in \Re^n . Since the intersection of two nondisjoint polyhedral sets is a polyhedral set, the set $F = P \cap H$ is polyhedral.

(b) Let P be given by

$$P = \{x \mid a'_j x \le b_j, \ j = 1, \dots, r\},\$$

for some vectors $a_j \in \mathbb{R}^n$ and scalars b_j . Let v be an extreme point of P, and without loss of generality assume that the first n inequalities define v, i.e., the first n of the vectors a_j are linearly independent and such that

$$a_j'v = b_j, \quad \forall \ j = 1, \dots, n$$

[cf. Prop. 2.1.4(a)]. Define the vector $a \in \Re^n$, the scalar b, and the hyperplane H as follows

$$a = \frac{1}{n} \sum_{j=1}^{n} a_j, \qquad b = \frac{1}{n} \sum_{j=1}^{n} b_j, \qquad H = \{x \mid a'x = b\}.$$

Then, we have

$$a'v = b$$

so that H passes through v. Moreover, for every $x \in P$, we have $a'_j x \leq b_j$ for all j, implying that $a'x \leq b$ for all $x \in P$. Thus, H contains P in one of its halfspaces.

We will next prove that $P \cap H = \{v\}$. We start by showing that for every $\overline{v} \in P \cap H$, we must have

$$a_j'\overline{v} = b_j, \qquad \forall \ j = 1, \dots, n.$$
 (2.1)

To arrive at a contradiction, assume that $a'_j \overline{v} < b_j$ for some $\overline{v} \in P \cap H$ and $j \in \{1, ..., n\}$. Without loss of generality, we can assume that the strict inequality holds for j = 1, so that

$$a_1'\overline{v} < b_1, \qquad a_j'\overline{v} \le b_j, \qquad \forall \ j = 2, \dots, n.$$

By multiplying each of the above inequalities with 1/n and by summing the obtained inequalities, we obtain

$$\frac{1}{n}\sum_{j=1}^n a_j'\overline{v} < \frac{1}{n}\sum_{j=1}^n b_j,$$

implying that $a'\overline{v} < b$, which contradicts the fact that $\overline{v} \in H$. Hence, Eq. (2.1) holds, and since the vectors a_1, \ldots, a_n are linearly independent, it follows that $v = \overline{v}$, showing that $P \cap H = \{v\}$.

As discussed in Section 2.1, every extreme point of P is a relative boundary point of P. Since every relative boundary point of P is also a boundary point of P, it follows that every extreme point of P is a boundary point of P. Thus, v is a boundary point of P, and as shown earlier, H passes through v and contains P in one of its halfspaces. By definition, it follows that $P \cap H = \{v\}$ is a face of P.

(c) Since P is not an affine set, it cannot consist of a single point, so we must have $\dim(P) > 0$. Let P be given by

$$P = \{x \mid a'_{j}x \le b_{j}, \ j = 1, \dots, r\},\$$

for some vectors $a_j \in \mathbb{R}^n$ and scalars b_j . Also, let A be the matrix with rows a'_j and b be the vector with components b_j , so that

$$P = \{x \mid Ax \le b\}.$$

An inequality $a'_j x \leq b_j$ of the system $Ax \leq b$ is redundant if it is implied by the remaining inequalities in the system. If the system $Ax \leq b$ has no redundant inequalities, we say that the system is nonredundant. An inequality $a'_j x \leq b_j$ of the system $Ax \leq b$ is an implicit equality if $a'_j x = b_j$ for all x satisfying $Ax \leq b$.

By removing the redundant inequalities if necessary, we may assume that the system $Ax \leq b$ defining P is nonredundant. Since P is not an affine set, there exists an inequality $a'_{j_0}x \leq b_{j_0}$ that is not an implicit equality of the system $Ax \leq b$. Consider the set

$$F = \{ x \in P \mid a'_{j_0} x = b_{j_0} \}.$$

Note that $F \neq \emptyset$, since otherwise $a'_{j_0}x \leq b_{j_0}$ would be a redundant inequality of the system $Ax \leq b$, contradicting our earlier assumption that the system is nonredundant. Note also that every point of F is a boundary point

of P. Thus, F is the intersection of P and the hyperplane $\{x \mid a'_{j_0}x = b_{j_0}\}$ that passes through a boundary point of P and contains P in one of its halfspaces, i.e., F is a face of P. Since $a'_{j_0}x \leq b_{j_0}$ is not an implicit equality of the system $Ax \leq b$, the dimension of F is $\dim(P) - 1$.

(d) Let P be a polyhedral set given by

$$P = \{x \mid a_i' x \le b_j, \ j = 1, \dots, r\},\$$

with $a_j \in \mathbb{R}^n$ and $b_j \in \mathbb{R}$, or equivalently

$$P = \{x \mid Ax \le b\},\$$

where A is an $r \times n$ matrix and $b \in \Re^r$. We will show that F is a face of P if and only if F is nonempty and

$$F = \{ x \in P \mid a'_{j}x = b_{j}, \ j \in J \},\$$

where $J \subset \{1, ..., r\}$. From this it will follow that the number of distinct faces of P is finite.

By removing the redundant inequalities if necessary, we may assume that the system $Ax \leq b$ defining P is nonredundant. Let F be a face of P, so that $F = P \cap H$, where H is a hyperplane that passes through a boundary point of P and contains P in one of its halfspaces. Let $H = \{x \mid c'x = c\overline{x}\}$ for a nonzero vector $c \in \Re^n$ and a boundary point \overline{x} of P, so that

$$F = \{ x \in P \mid c'x = c\overline{x} \}$$

and

$$c'x < c\overline{x}, \quad \forall \ x \in P.$$

These relations imply that the set of points x such that $Ax \leq b$ and $c'x \leq c\overline{x}$ coincides with P, and since the system $Ax \leq b$ is nonredundant, it follows that $c'x \leq c\overline{x}$ is a redundant inequality of the system $Ax \leq b$ and $c'x \leq c\overline{x}$. Therefore, the inequality $c'x \leq c\overline{x}$ is implied by the inequalities of $Ax \leq b$, so that there exists some $\mu \in \Re^r$ with $\mu \geq 0$ such that

$$\sum_{j=1}^{r} \mu_j a_j = c, \qquad \sum_{j=1}^{r} \mu_j b_j = c' \overline{x}.$$

Let $J = \{j \mid \mu_j > 0\}$. Then, for every $x \in P$, we have

$$c'x = c\overline{x} \iff \sum_{j \in J} \mu_j a'_j x = \sum_{j \in J} \mu_j b_j \iff a'_j x = b_j, \ j \in J, \quad (2.2)$$

implying that

$$F = \left\{ x \in P \mid a'_j x = b_j, j \in J \right\}.$$

Conversely, let F be a nonempty set given by

$$F = \{x \in P \mid a_j' x = b_j, j \in J\},\$$

for some $J \subset \{1, \ldots, r\}$. Define

$$c = \sum_{j \in J} a_j, \qquad \beta = \sum_{j \in J} b_j.$$

Then, we have

$$\{x \in P \mid a_j' x = b_j, j \in J\} = \{x \in P \mid c' x = \beta\},\$$

[cf. Eq. (2.2) where $\mu_j = 1$ for all $j \in J$]. Let $H = \{x \mid c'x = \beta\}$, so that in view of the preceding relation, we have that $F = P \cap H$. Since every point of F is a boundary point of P, it follows that H passes through a boundary point of P. Furthermore, for every $x \in P$, we have $a'_j x \leq b_j$ for all $j \in J$, implying that $c'x \leq \beta$ for every $x \in P$. Thus, H contains P in one of its halfspaces. Hence, F is a face.

2.7 (Isomorphic Polyhedral Sets)

Let P and Q be polyhedral sets in \Re^n and \Re^m , respectively. We say that P and Q are *isomorphic* if there exist affine functions $f: P \mapsto Q$ and $g: Q \mapsto P$ such that

$$x = g(f(x)), \quad \forall \ x \in P, \qquad y = f(g(y)), \quad \forall \ y \in Q.$$

- (a) Show that if P and Q are isomorphic, then their extreme points are in one-to-one correspondence.
- (b) Let A be an $r \times n$ matrix and b be a vector in \Re^r , and let

$$P = \{x \in \Re^n \mid Ax < b, \ x > 0\},\$$

$$Q = \{(x, z) \in \Re^{n+r} \mid Ax + z = b, \ x \ge 0, \ z \ge 0\}.$$

Show that P and Q are isomorphic.

Solution: (a) Let P and Q be isomorphic polyhedral sets, and let $f: P \mapsto Q$ and $g: Q \mapsto P$ be affine functions such that

$$x = g(f(x)), \quad \forall \ x \in P, \qquad y = f(g(y)), \quad \forall \ y \in Q.$$

Assume that x^* is an extreme point of P and let $y^* = f(x^*)$. We will show that y^* is an extreme point of Q. Since x^* is an extreme point of P, by

Exercise 2.6(b), it is also a face of P, and therefore, there exists a vector $c \in \Re^n$ such that

$$c'x < c'x^*, \quad \forall \ x \in P, \ x \neq x^*.$$

For any $y \in Q$ with $y \neq y^*$, we have

$$f(g(y)) = y \neq y^* = f(x^*),$$

implying that

$$g(y) \neq g(y^*) = x^*,$$
 with $g(y) \in P$.

Hence,

$$c'g(y) < c'g(y^*), \qquad \forall \ y \in Q, \ y \neq y^*.$$

Let the affine function g be given by g(y) = By + d for some $n \times m$ matrix B and vector $d \in \Re^n$. Then, we have

$$c'(By+d) < c'(By^*+d), \qquad \forall \ y \in Q, \ y \neq y^*,$$

implying that

$$(B'c)'y < (B'c)'y^*, \qquad \forall \ y \in Q, \ y \neq y^*.$$

If y^* were not an extreme point of Q, then we would have $y^* = \alpha y_1 + (1 - \alpha)y_2$ for some distinct points $y_1, y_2 \in Q$, $y_1 \neq y^*$, $y_2 \neq y^*$, and $\alpha \in (0, 1)$, so that

$$(B'c)'y^* = \alpha(B'c)'y_1 + (1-\alpha)(B'c)'y_2 < (B'c)'y^*,$$

which is a contradiction. Hence, y^* is an extreme point of Q.

Conversely, if y^* is an extreme point of Q, then by using a symmetrical argument, we can show that x^* is an extreme point of P.

(b) For the sets

$$\begin{split} P &= \{x \in \Re^n \mid Ax \le b, \ x \ge 0\}, \\ Q &= \big\{(x,z) \in \Re^{n+r} \mid Ax + z = b, \ x \ge 0, \ z \ge 0\big\}, \end{split}$$

let f and g be given by

$$f(x) = (x, b - Ax), \quad \forall x \in P,$$

 $g(x, z) = x, \quad \forall (x, z) \in Q.$

Evidently, f and g are affine functions. Furthermore, clearly

$$f(x) \in Q,$$
 $g(f(x)) = x,$ $\forall x \in P,$ $g(x,z) \in P,$ $f(g(x,z)) = x,$ $\forall (x,z) \in Q.$

Hence, P and Q are isomorphic.

SECTION 2.2: Polar Cones

2.8 (Cone Decomposition Theorem)

Let C be a nonempty closed convex cone in \Re^n and let x be a vector in \Re^n . Show that:

(a) \hat{x} is the projection of x on C if and only if

$$\hat{x} \in C$$
, $(x - \hat{x})'\hat{x} = 0$, $x - \hat{x} \in C^*$.

- (b) The following two statements are equivalent:
 - (i) x_1 and x_2 are the projections of x on C and C^* , respectively.
 - (ii) $x = x_1 + x_2$ with $x_1 \in C$, $x_2 \in C^*$, and $x_1'x_2 = 0$.

Solution: (a) Let \hat{x} be the projection of x on C, which exists and is unique since C is closed and convex. By the Projection Theorem (Prop. 1.1.9), we have

$$(x - \hat{x})'(y - \hat{x}) \le 0, \quad \forall y \in C.$$

Since C is a cone, we have $(1/2)\hat{x} \in C$ and $2\hat{x} \in C$, and by taking $y = (1/2)\hat{x}$ and $y = 2\hat{x}$ in the preceding relation, it follows that

$$(x - \hat{x})'\hat{x} = 0.$$

By combining the preceding two relations, we obtain

$$(x - \hat{x})'y \le 0, \quad \forall y \in C,$$

implying that $x - \hat{x} \in C^*$.

Conversely, if $\hat{x} \in C$, $(x - \hat{x})'\hat{x} = 0$, and $x - \hat{x} \in C^*$, then it follows that

$$(x - \hat{x})'(y - \hat{x}) < 0, \qquad \forall \ y \in C,$$

and by the Projection Theorem, \hat{x} is the projection of x on C.

(b) Suppose that property (i) holds, i.e., x_1 and x_2 are the projections of x on C and C^* , respectively. Then, by part (a), we have

$$x_1 \in C$$
, $(x - x_1)'x_1 = 0$, $x - x_1 \in C^*$.

Let $y = x - x_1$, so that the preceding relation can equivalently be written as

$$x - y \in C = (C^*)^*, \quad y'(x - y) = 0, \quad y \in C^*.$$

By using part (a), we conclude that y is the projection of x on C^* . Since by the Projection Theorem, the projection of a vector on a closed convex set is unique, it follows that $y = x_2$. Thus, we have $x = x_1 + x_2$ and in view of the preceding two relations, we also have $x_1 \in C$, $x_2 \in C^*$, and $x_1'x_2 = 0$. Hence, property (ii) holds.

Conversely, suppose that property (ii) holds, i.e., $x = x_1 + x_2$ with $x_1 \in C$, $x_2 \in C^*$, and $x_1'x_2 = 0$. Then, evidently the relations

$$x_1 \in C$$
, $(x - x_1)'x_1 = 0$, $x - x_1 \in C^*$,

$$x_2 \in C^*, \qquad (x - x_2)'x_2 = 0, \qquad x - x_2 \in C$$

are satisfied, so that by part (a), x_1 and x_2 are the projections of x on C and C^* , respectively. Hence, property (i) holds.

2.9

Let C be a nonempty closed convex cone in \Re^n and let a be a vector in \Re^n . Show that for any scalars $\beta > 0$ and $\gamma \geq 0$, we have

$$\max_{\|x\| \leq \beta, \, x \in C} a'x \leq \gamma \qquad \text{if and only if} \qquad a \in C^* + \big\{x \mid \|x\| \leq \gamma/\beta\big\}.$$

(This may be viewed as an "approximate" version of the Polar Cone Theorem, which is obtained for $\gamma = 0$.)

Solution: If $a \in C^* + \{x \mid ||x|| \le \gamma/\beta\}$, then

$$a = \hat{a} + \overline{a}$$
 with $\hat{a} \in C^*$ and $\|\overline{a}\| \le \gamma/\beta$.

Since C is a closed convex cone, by the Polar Cone Theorem (Prop. 2.2.1), we have $(C^*)^* = C$, implying that for all x in C with $||x|| \leq \beta$,

$$\hat{a}'x \leq 0$$
 and $\overline{a}'x \leq \|\overline{a}\| \cdot \|x\| \leq \gamma$.

Hence,

$$a'x = (\hat{a} + \overline{a})'x \le \gamma, \quad \forall x \in C \text{ with } ||x|| \le \beta,$$

thus implying that

$$\max_{\|x\| \le \beta, \ x \in C} a'x \le \gamma.$$

Conversely, assume that $a'x \leq \gamma$ for all $x \in C$ with $||x|| \leq \beta$. Let \hat{a} and \overline{a} be the projections of a on C^* and C, respectively. By the Cone Decomposition Theorem (cf. Exercise 2.8), we have $a = \hat{a} + \overline{a}$ with $\hat{a} \in C^*$, $\overline{a} \in C$, and $\hat{a}'\overline{a} = 0$. Since $a'x \leq \gamma$ for all $x \in C$ with $||x|| \leq \beta$ and $\overline{a} \in C$, we obtain

$$a'\frac{\overline{a}}{\|\overline{a}\|}\beta = (\hat{a} + \overline{a})'\frac{\overline{a}}{\|\overline{a}\|}\beta = \|\overline{a}\|\beta \le \gamma,$$

implying that $\|\overline{a}\| \le \gamma/\beta$, and showing that $a \in C^* + \{x \mid \|x\| \le \gamma/\beta\}$.

2.10 (Dimension and Lineality Space of Polar Cones)

Show that for any nonempty cone C in \Re^n , we have

$$L_{C^*} = (\operatorname{aff}(C))^{\perp},$$

$$\dim(C) + \dim(L_{C^*}) = n,$$

$$\dim(C^*) + \dim(L_{\operatorname{conv}(C)}) \le \dim(C^*) + \dim(L_{cl(\operatorname{conv}(C))}) = n,$$

where L_X denotes the lineality space of a convex set X.

Solution: Note that aff(C) is a subspace of \Re^n because C is a cone in \Re^n . We first show that

$$L_{C^*} = (\operatorname{aff}(C))^{\perp}.$$

Let $y \in L_{C^*}$. Then, by the definition of the lineality space (see Chapter 1), both vectors y and -y belong to the recession cone R_{C^*} . Since $0 \in C^*$, it follows that 0 + y and 0 - y belong to C^* . Therefore,

$$y'x \le 0, \qquad (-y)'x \le 0, \qquad \forall \ x \in C,$$

implying that

$$y'x = 0, \qquad \forall \ x \in C. \tag{2.3}$$

Let the dimension of the subspace $\operatorname{aff}(C)$ be m. By Prop. 1.3.2, there exist vectors x_0, x_1, \ldots, x_m in $\operatorname{ri}(C)$ such that $x_1 - x_0, \ldots, x_m - x_0$ span $\operatorname{aff}(C)$. Thus, for any $z \in \operatorname{aff}(C)$, there exist scalars β_1, \ldots, β_m such that

$$z = \sum_{i=1}^{m} \beta_i (x_i - x_0).$$

By using this relation and Eq. (2.3), for any $z \in aff(C)$, we obtain

$$y'z = \sum_{i=1}^{m} \beta_i y'(x_i - x_0) = 0,$$

implying that $y \in (\operatorname{aff}(C))^{\perp}$. Hence, $L_{C^*} \subset (\operatorname{aff}(C))^{\perp}$.

Conversely, let $y \in (\operatorname{aff}(C))^{\perp}$, so that in particular, we have

$$y'x = 0,$$
 $(-y)'x = 0,$ $\forall x \in C.$

Therefore, $0 + \alpha y \in C^*$ and $0 + \alpha(-y) \in C^*$ for all $\alpha \geq 0$, and since C^* is a closed convex set, by the Recession Cone Theorem [Prop. 1.4.1(b)], it follows that y and -y belong to the recession cone R_{C^*} . Hence, y belongs

to the lineality space of C^* , showing that $\left(\operatorname{aff}(C)\right)^{\perp} \subset L_{C^*}$ and completing the proof of the equality $L_{C^*} = \left(\operatorname{aff}(C)\right)^{\perp}$.

By definition, we have $\dim(C) = \dim(\operatorname{aff}(C))$ and since $L_{C^*} = (\operatorname{aff}(C))^{\perp}$, we have $\dim(L_{C^*}) = \dim((\operatorname{aff}(C))^{\perp})$. This implies that

$$\dim(C) + \dim(L_{C^*}) = n.$$

By replacing C with C^* in the preceding relation, and by using the Polar Cone Theorem (Prop. 2.2.1), we obtain

$$\dim(C^*) + \dim(L_{(C^*)^*}) = \dim(C^*) + \dim(L_{\operatorname{cl}(\operatorname{conv}(C))}) = n.$$

Furthermore, since

$$L_{\operatorname{conv}(C)} \subset L_{\operatorname{cl}(\operatorname{conv}(C))},$$

it follows that

$$\dim(C^*) + \dim(L_{\operatorname{conv}(C)}) \le \dim(C^*) + \dim(L_{\operatorname{cl}(\operatorname{conv}(C))}) = n.$$

2.11 (Polar Cone Operations)

Show the following:

(a) For any nonempty cones $C_i \subset \Re^{n_i}$, i = 1, ..., m, we have

$$(C_1 \times \cdots \times C_m)^* = C_1^* \times \cdots \times C_m^*.$$

(b) For any collection of nonempty cones $\{C_i \mid i \in I\}$, we have

$$\left(\bigcup_{i\in I} C_i\right)^* = \bigcap_{i\in I} C_i^*.$$

(c) For any two nonempty cones C_1 and C_2 , we have

$$(C_1 + C_2)^* = C_1^* \cap C_2^*.$$

(d) For any two nonempty closed convex cones C_1 and C_2 , we have

$$(C_1 \cap C_2)^* = \operatorname{cl}(C_1^* + C_2^*).$$

Furthermore, if $ri(C_1) \cap ri(C_2) \neq \emptyset$, then the cone $C_1^* + C_2^*$ is closed and the closure operation in the preceding relation can be omitted.

(e) Consider the following cones in \Re^3

$$C_1 = \{(x_1, x_2, x_3) \mid x_1^2 + x_2^2 \le x_3^2, \ x_3 \le 0\},$$

$$C_2 = \{(x_1, x_2, x_3) \mid x_2 = -x_3\}.$$

Verify that $ri(C_1) \cap ri(C_2) = \emptyset$, $(1,1,1) \in (C_1 \cap C_2)^*$, and $(1,1,1) \notin$ $C_1^* + C_2^*$, thus showing that the closure operation in the relation of part (c) may not be omitted when $ri(C_1) \cap ri(C_2) = \emptyset$.

Solution: (a) It suffices to consider the case where m=2. Let $(y_1,y_2) \in$ $(C_1 \times C_2)^*$. Then, we have $(y_1, y_2)'(x_1, x_2) \leq 0$ for all $(x_1, x_2) \in C_1 \times C_2$, or equivalently

$$y_1'x_1 + y_2'x_2 \le 0, \quad \forall x_1 \in C_1, \quad \forall x_2 \in C_2.$$

Since C_2 is a cone, 0 belongs to its closure, so by letting $x_2 \to 0$ in the preceding relation, we obtain $y_1'x_1 \leq 0$ for all $x_1 \in C_1$, showing that $y_1 \in C_1^*$. Similarly, we obtain $y_2 \in C_2^*$, and therefore $(y_1, y_2) \in C_1^* \times C_2^*$, implying that $(C_1 \times C_2)^* \subset C_1^* \times C_2^*$. Conversely, let $y_1 \in C_1^*$ and $y_2 \in C_2^*$. Then, we have

$$(y_1, y_2)'(x_1, x_2) = y_1'x_1 + y_2'x_2 \le 0, \quad \forall x_1 \in C_1, \quad \forall x_2 \in C_2,$$

implying that $(y_1, y_2) \in (C_1 \times C_2)^*$, and showing that $C_1^* \times C_2^* \subset (C_1 \times C_2)^*$.

- (b) A vector y belongs to the polar cone of $\bigcup_{i \in I} C_i$ if and only if $y'x \leq 0$ for all $x \in C_i$ and all $i \in I$, which is equivalent to having $y \in C_i^*$ for every $i \in I$. Hence, y belongs to $(\bigcup_{i \in I} C_i)^*$ if and only if y belongs to $\bigcap_{i \in I} C_i^*$.
- (c) Let $y \in (C_1 + C_2)^*$, so that

$$y'(x_1 + x_2) \le 0, \quad \forall x_1 \in C_1, \quad \forall x_2 \in C_2.$$
 (2.4)

Since the zero vector is in the closures of C_1 and C_2 , by letting $x_2 \to 0$ with $x_2 \in C_2$ in Eq. (2.4), we obtain

$$y'x_1 \leq 0, \quad \forall x_1 \in C_1,$$

and similarly, by letting $x_1 \to 0$ with $x_1 \in C_1$ in Eq. (2.4), we obtain

$$y'x_2 \le 0, \quad \forall \ x_2 \in C_2.$$

Thus, $y \in C_1^* \cap C_2^*$, showing that $(C_1 + C_2)^* \subset C_1^* \cap C_2^*$. Conversely, let $y \in C_1^* \cap C_2^*$. Then, we have

$$y'x_1 \le 0, \qquad \forall \ x_1 \in C_1,$$

$$y'x_2 < 0, \qquad \forall \ x_2 \in C_2,$$

implying that

$$y'(x_1 + x_2) \le 0, \quad \forall x_1 \in C_1, \quad \forall x_2 \in C_2.$$

Hence $y \in (C_1 + C_2)^*$, showing that $C_1^* \cap C_2^* \subset (C_1 + C_2)^*$.

(d) Since C_1 and C_2 are closed convex cones, by the Polar Cone Theorem (Prop. 2.2.1) and by part (b), it follows that

$$C_1 \cap C_2 = (C_1^*)^* \cap (C_2^*)^* = (C_1^* + C_2^*)^*.$$

By taking the polars and by using the Polar Cone Theorem, we obtain

$$(C_1 \cap C_2)^* = ((C_1^* + C_2^*)^*)^* = \operatorname{cl}(\operatorname{conv}(C_1^* + C_2^*)).$$

The cone $C_1^* + C_2^*$ is convex, so that

$$(C_1 \cap C_2)^* = \operatorname{cl}(C_1^* + C_2^*).$$

Suppose now that $\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2) \neq \emptyset$. We will show that $C_1^* + C_2^*$ is closed by using Prop. 1.4.14. According to this proposition, if for any nonempty closed convex sets \overline{C}_1 and \overline{C}_2 in \Re^n , the equality $y_1 + y_2 = 0$ with $y_1 \in R_{\overline{C}_1}$ and $y_2 \in R_{\overline{C}_2}$ implies that y_1 and y_2 belong to the lineality spaces of \overline{C}_1 and \overline{C}_2 , respectively, then the vector sum $\overline{C}_1 + \overline{C}_2$ is closed.

Let $y_1+y_2=0$ with $y_1\in R_{C_1^*}$ and $y_2\in R_{C_2^*}$. Because C_1^* and C_2^* are closed convex cones, we have $R_{C_1^*}=C_1^*$ and $R_{C_2^*}=C_2^*$, so that $y_1\in C_1^*$ and $y_2\in C_2^*$. The lineality space of a cone is the set of vectors y such that y and -y belong to the cone, so that in view of the preceding discussion, to show that $C_1^*+C_2^*$ is closed, it suffices to prove that $-y_1\in C_1^*$ and $-y_2\in C_2^*$.

Since $y_1 = -y_2$ and $y_1 \in C_1^*$, it follows that

$$y_2'x \ge 0, \qquad \forall \ x \in C_1, \tag{2.5}$$

and because $y_2 \in C_2^*$, we have

$$y_2'x \leq 0, \quad \forall x \in C_2,$$

which combined with the preceding relation yields

$$y_2'x = 0, \qquad \forall \ x \in C_1 \cap C_2. \tag{2.6}$$

In view of the fact $\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2) \neq \emptyset$, and Eqs. (2.5) and (2.6), it follows that the linear function $y_2'x$ attains its minimum over the convex set C_1 at a point in the relative interior of C_1 , implying that $y_2'x = 0$ for all $x \in C_1$ (cf. Prop. 1.3.4). Therefore, $y_2 \in C_1^*$ and since $y_2 = -y_1$, we have $-y_1 \in C_1^*$. By exchanging the roles of y_1 and y_2 in the preceding analysis, we similarly show that $-y_2 \in C_2^*$, completing the proof.

(e) By drawing the cones C_1 and C_2 , it can be seen that $ri(C_1) \cap ri(C_2) = \emptyset$ and

$$C_1 \cap C_2 = \{(x_1, x_2, x_3) \mid x_1 = 0, \ x_2 = -x_3, \ x_3 \le 0\},\$$

$$C_1^* = \{ (y_1, y_2, y_3) \mid y_1^2 + y_2^2 \le y_3^2, \ y_3 \ge 0 \},$$

$$C_2^* = \{ (z_1, z_2, z_3) \mid z_1 = 0, \ z_2 = z_3 \}.$$

Clearly, $x_1+x_2+x_3=0$ for all $x\in C_1\cap C_2$, implying that $(1,1,1)\in (C_1\cap C_2)^*$. Suppose that $(1,1,1)\in C_1^*+C_2^*$, so that $(1,1,1)=(y_1,y_2,y_3)+(z_1,z_2,z_3)$ for some $(y_1,y_2,y_3)\in C_1^*$ and $(z_1,z_2,z_3)\in C_2^*$, implying that $y_1=1,\ y_2=1-z_2,\ y_3=1-z_2$ for some $z_2\in\Re$. However, this point does not belong to C_1^* , which is a contradiction. Therefore, (1,1,1) is not in $C_1^*+C_2^*$. Hence, when $\mathrm{ri}(C_1)\cap\mathrm{ri}(C_2)=\emptyset$, the relation

$$(C_1 \cap C_2)^* = C_1^* + C_2^*$$

may fail.

2.12 (Linear Transformations and Polar Cones)

Let C be a nonempty cone in \Re^n , K be a nonempty closed convex cone in \Re^m , and A be a linear transformation from \Re^n to \Re^m . Show that

$$(AC)^* = (A')^{-1} \cdot C^*, \qquad (A^{-1} \cdot K)^* = \operatorname{cl}(A'K^*).$$

Show also that if $ri(K) \cap R(A) \neq \emptyset$, then the cone $A'K^*$ is closed and $(A')^{-1}$ and the closure operation in the above relation can be omitted.

Solution: We have $y \in (AC)^*$ if and only if $y'Ax \leq 0$ for all $x \in C$, which is equivalent to $(A'y)'x \leq 0$ for all $x \in C$. This is in turn equivalent to $A'y \in C^*$. Hence, $y \in (AC)^*$ if and only if $y \in (A')^{-1} \cdot C^*$, showing that

$$(AC)^* = (A')^{-1} \cdot C^*. \tag{2.7}$$

We next show that for a closed convex cone $K \subset \Re^m$, we have

$$(A^{-1} \cdot K)^* = \operatorname{cl}(A'K^*).$$

Let $y \in (A^{-1} \cdot K)^*$ and to arrive at a contradiction, assume that $y \notin \operatorname{cl}(A'K^*)$. By the Strict Separation Theorem (Prop. 1.5.3), the closed convex cone $\operatorname{cl}(A'K^*)$ and the vector y can be strictly separated, i.e., there exist a vector $a \in \Re^n$ and a scalar b such that

$$a'x < b < a'y$$
, $\forall x \in \operatorname{cl}(A'K^*)$.

If a'x>0 for some $x\in \operatorname{cl}(A'K^*)$, then since $\operatorname{cl}(A'K^*)$ is a cone, we would have $\lambda x\in \operatorname{cl}(A'K^*)$ for all $\lambda>0$, implying that $a'(\lambda x)\to\infty$ when $\lambda\to\infty$, which contradicts the preceding relation. Thus, we must have $a'x\le 0$ for all $x\in\operatorname{cl}(A'K^*)$, and since $0\in\operatorname{cl}(A'K^*)$, it follows that

$$\sup_{x \in cl(A'K^*)} a'x = 0 \le b < a'y.$$
 (2.8)

Therefore, $a \in (\operatorname{cl}(A'K^*))^*$, and since $(\operatorname{cl}(A'K^*))^* \subset (A'K^*)^*$, it follows that $a \in (A'K^*)^*$. In view of Eq. (2.7) and the Polar Cone Theorem (Prop. 2.2.1), we have

$$(A'K^*)^* = A^{-1}(K^*)^* = A^{-1} \cdot K,$$

implying that $a \in A^{-1} \cdot K$. Because $y \in (A^{-1} \cdot K)^*$, it follows that $y'a \leq 0$, contradicting Eq. (2.8). Hence, we must have $y \in \operatorname{cl}(A'K^*)$, showing that

$$(A^{-1} \cdot K)^* \subset \operatorname{cl}(A'K^*).$$

To show the reverse inclusion, let $y \in A'K^*$ and assume, to arrive at a contradiction, that $y \notin (A^{-1} \cdot K)^*$. By the Strict Separation Theorem (Prop. 1.5.3), the closed convex cone $(A^{-1} \cdot K)^*$ and the vector y can be strictly separated, i.e., there exist a vector $\overline{a} \in \Re^n$ and a scalar \overline{b} such that

$$\overline{a}'x < \overline{b} < \overline{a}'y, \qquad \forall \ x \in (A^{-1} \cdot K)^*.$$

Similar to the preceding analysis, since $(A^{-1} \cdot K)^*$ is a cone, it can be seen that

$$\sup_{x \in (A^{-1} \cdot K)^*} \overline{a}' x = 0 \le \overline{b} < \overline{a}' y, \tag{2.9}$$

implying that $\overline{a} \in \left((A^{-1} \cdot K)^* \right)^*$. Since K is a closed convex cone and A is a linear (and therefore continuous) transformation, the set $A^{-1} \cdot K$ is a closed convex cone. Furthermore, by the Polar Cone Theorem, we have that $\left((A^{-1} \cdot K)^* \right)^* = A^{-1} \cdot K$. Therefore, $\overline{a} \in A^{-1} \cdot K$, implying that $A\overline{a} \in K$. Since $y \in A'K^*$, we have y = A'v for some $v \in K^*$, and it follows that

$$y'\overline{a} = (A'v)'\overline{a} = v'A\overline{a} \le 0,$$

contradicting Eq. (2.9). Hence, we must have $y \in (A^{-1} \cdot K)^*$, implying that

$$A'K^* \subset (A^{-1} \cdot K)^*$$
.

Taking the closure of both sides of this relation, we obtain

$$\operatorname{cl}(A'K^*) \subset (A^{-1} \cdot K)^*,$$

completing the proof.

Suppose that $ri(K^*) \cap R(A) \neq \emptyset$. We will show that the cone $A'K^*$ is closed by using Prop. 1.4.13. According to this proposition, if $R_{K^*} \cap N(A')$ is a subspace of the lineality space L_{K^*} of K^* , then

$$cl(A'K^*) = A'K^*.$$

Thus, it suffices to verify that $R_{K^*} \cap N(A')$ is a subspace of L_{K^*} . Indeed, we will show that $R_{K^*} \cap N(A') = L_{K^*} \cap N(A')$.

Let $y \in K^* \cap N(A')$. Because $y \in K^*$, we obtain

$$(-y)'x \ge 0, \qquad \forall \ x \in K. \tag{2.10}$$

For $y \in N(A')$, we have $-y \in N(A')$ and since $N(A') = R(A)^{\perp}$, it follows that

$$(-y)'z = 0, \qquad \forall \ z \in R(A). \tag{2.11}$$

In view of the relation $\operatorname{ri}(K) \cap R(A) \neq \emptyset$, and Eqs. (2.10) and (2.11), the linear function (-y)'x attains its minimum over the convex set K at a point in the relative interior of K, implying that (-y)'x = 0 for all $x \in K$ (cf. Prop. 1.3.4). Hence $(-y) \in K^*$, so that $y \in L_{K^*}$ and because $y \in N(A')$, we see that $y \in L_{K^*} \cap N(A')$. The reverse inclusion follows directly from the relation $L_{K^*} \subset R_{K^*}$, thus completing the proof.

2.13 (Pointed Cones and Bases)

Let C be a closed convex cone in \Re^n . We say that C is a *pointed cone* if $C \cap (-C) = \{0\}$. A convex set $D \subset \Re^n$ is said to be a *base* for C if C = cone(D) and $0 \notin \text{cl}(D)$. Show that the following properties are equivalent:

- (a) C is a pointed cone.
- (b) $cl(C^* C^*) = \Re^n$.
- (c) $C^* C^* = \Re^n$.
- (d) C^* has nonempty interior.
- (e) There exist a nonzero vector $\hat{x} \in \Re^n$ and a positive scalar δ such that $\hat{x}'x \geq \delta ||x||$ for all $x \in C$.
- (f) C has a bounded base.

Hint: Use Exercise 2.11 to show the implications (a) \Rightarrow (b) \Rightarrow (c) \Rightarrow (d) \Rightarrow (e) \Rightarrow (f) \Rightarrow (a).

Solution: (a) \Rightarrow (b) Since C is a pointed cone, $C \cap (-C) = \{0\}$, so that

$$(C \cap (-C))^* = \Re^n.$$

On the other hand, by Exercise 2.11, it follows that

$$(C \cap (-C))^* = \operatorname{cl}(C^* - C^*),$$

which when combined with the preceding relation yields $\operatorname{cl}(C^*-C^*)=\Re^n$.

(b) \Rightarrow (c) Since C is a closed convex cone, by the polar cone operations of Exercise 2.11, it follows that

$$(C \cap (-C))^* = \operatorname{cl}(C^* - C^*) = \Re^n.$$

By taking the polars and using the Polar Cone Theorem (Prop. 2.2.1), we obtain

$$\left(\left(C \cap (-C) \right)^* \right)^* = C \cap (-C) = \{0\}. \tag{2.12}$$

Now, to arrive at a contradiction assume that there is a vector $\hat{x} \in \mathbb{R}^n$ such that $\hat{x} \notin C^* - C^*$. Then, by the Separating Hyperplane Theorem (Prop. 1.5.2), there exists a nonzero vector $a \in \mathbb{R}^n$ such that

$$a'\hat{x} \ge a'x, \quad \forall \ x \in C^* - C^*.$$

If a'x > 0 for some $x \in C^* - C^*$, then since $C^* - C^*$ is a cone, the right hand-side of the preceding relation can be arbitrarily large, a contradiction. Thus, we have $a'x \leq 0$ for all $x \in C^* - C^*$, implying that

$$a \in (C^* - C^*)^*$$
.

By the polar cone operations of Exercise 2.11(b) and the Polar Cone Theorem, it follows that

$$(C^* - C^*)^* = (C^*)^* \cap (-C^*)^* = C \cap (-C).$$

Thus, $a \in C \cap (-C)$ with $a \neq 0$, contradicting Eq. (2.12). Hence, we must have $C^* - C^* = \Re^n$.

- (c) \Rightarrow (d) Because $C^* \subset \operatorname{aff}(C^*)$ and $-C^* \subset \operatorname{aff}(C^*)$, we have $C^* C^* \subset \operatorname{aff}(C^*)$ and since $C^* C^* = \Re^n$, it follows that $\operatorname{aff}(C^*) = \Re^n$, showing that C^* has nonempty interior.
- (d) \Rightarrow (e) Let v be a vector in the interior of C^* . Then, there exists a positive scalar δ such that the vector $v + \delta \frac{y}{\|y\|}$ is in C^* for all $y \in \Re^n$ with $y \neq 0$, i.e.,

$$\left(v + \delta \frac{y}{\|y\|}\right)' x \le 0, \quad \forall \ x \in C, \quad \forall \ y \in \Re^n, \ y \ne 0.$$

By taking y = x, it follows that

$$\left(v + \delta \frac{x}{\|x\|}\right)' x \le 0, \quad \forall x \in C, \ x \ne 0,$$

implying that

$$v'x + \delta ||x|| \le 0, \qquad \forall \ x \in C, \ x \ne 0.$$

Clearly, this relation holds for x = 0, so that

$$v'x \le -\delta ||x||, \quad \forall \ x \in C.$$

Multiplying the preceding relation with -1 and letting $\hat{x} = -v$, we obtain

$$\hat{x}'x \ge \delta ||x||, \quad \forall \ x \in C.$$

$$(e) \Rightarrow (f)$$
 Let

$$D = \{ y \in C \mid \hat{x}'y = 1 \}.$$

Then, D is a closed convex set since it is the intersection of the closed convex cone C and the closed convex set $\{y \mid \hat{x}'y = 1\}$. Obviously, $0 \notin D$. Thus, to show that D is a base for C, it remains to prove that $C = \operatorname{cone}(D)$. Take any $x \in C$. If x = 0, then $x \in \operatorname{cone}(D)$ and we are done, so assume that $x \neq 0$. We have by hypothesis

$$\hat{x}'x \ge \delta \|x\| > 0, \qquad \forall \ x \in C, \ x \ne 0,$$

so we may define

$$\hat{y} = \frac{x}{\hat{x}'x}.$$

Clearly, $\hat{y} \in D$ and $x = (\hat{x}'x)\hat{y}$ with $\hat{x}'x > 0$, showing that $x \in \text{cone}(D)$ and that $C \subset \text{cone}(D)$. Since $D \subset C$, the inclusion $\text{cone}(D) \subset C$ is obvious. Thus, C = cone(D) and D is a base for C. Furthermore, for every y in D, since y is also in C, we have

$$1 = \hat{x}'y \ge \delta ||y||,$$

showing that D is bounded and completing the proof.

(f) \Rightarrow (a) Since C has a bounded base, C = cone(D) for some bounded convex set D with $0 \notin \text{cl}(D)$. To arrive at a contradiction, we assume that the cone C is not pointed, so that there exists a nonzero vector

$$d \in C \cap (-C)$$
,

implying that d and -d are in C. Let $\{\lambda_k\}$ be a sequence of positive scalars. Since $\lambda_k d \in C$ for all k and D is a base for C, there exist a sequence $\{\mu_k\}$ of positive scalars and a sequence $\{y_k\}$ of vectors in D such that

$$\lambda_k d = \mu_k y_k, \quad \forall k.$$

Therefore,

$$y_k = \frac{\lambda_k}{\mu_k} d \in D$$

for all k and because D is bounded, the sequence $\{y_k\}$ has a subsequence converging to some $y \in \operatorname{cl}(D)$. Without loss of generality, we may assume that $y_k \to y$, which in view of $y_k = \frac{\lambda_k}{\mu_k} d$ for all k, implies that $y = \alpha d$ and $\alpha d \in \operatorname{cl}(D)$ for some $\alpha \geq 0$. Furthermore, by the definition of base, we have

 $0 \notin \operatorname{cl}(D)$, so that $\alpha > 0$. Similar to the preceding, by replacing d with -d, we can show that $\tilde{\alpha}(-d) \in \operatorname{cl}(D)$ for some positive scalar $\tilde{\alpha}$. Therefore, $\alpha d \in \operatorname{cl}(D)$ and $\tilde{\alpha}(-d) \in \operatorname{cl}(D)$ with $\alpha > 0$ and $\tilde{\alpha} > 0$. Since D is convex, its closure $\operatorname{cl}(D)$ is also convex, implying that $0 \in \operatorname{cl}(D)$, contradicting the definition of a base. Hence, the cone C must be pointed.

SECTION 2.3: Polyhedral Sets and Functions

2.14

Show that a closed convex cone is polyhedral if and only if its polar cone is polyhedral.

Solution: Let the closed convex cone C be polyhedral, and of the form

$$C = \{ x \mid a_j' x \le 0, \ j = 1, \dots, r \},\$$

for some vectors a_i in \Re^n . By Farkas' Lemma, we have

$$C^* = \operatorname{cone}(\{a_1, \dots, a_r\}),$$

so the polar cone of a polyhedral cone is finitely generated. Conversely, using the Polar Cone Theorem, we have

cone
$$(\{a_1, \dots, a_r\})^* = \{x \mid a'_j x \le 0, \ j = 1, \dots, r\},\$$

so the polar of a finitely generated cone is polyhedral. Thus, a closed convex cone is polyhedral if and only if its polar cone is finitely generated. By the Minkowski-Weyl Theorem (Prop. 2.3.2), a cone is finitely generated if and only if it is polyhedral. Therefore, a closed convex cone is polyhedral if and only if its polar cone is polyhedral.

2.15 (Closedness of Finitely Generated Cones)

This exercise proves that a finitely generated cone is closed without invoking Prop. 1.4.13. Let a_1, \ldots, a_r be vectors in \Re^n and let A be the $n \times r$ matrix that has as columns these vectors. Consider the cone generated by a_1, \ldots, a_r :

cone
$$(\{a_1, \dots, a_r\}) = \{A\mu \mid \mu \ge 0\}.$$

(a) Show that if a_1, \ldots, a_r are linearly independent, then cone $(\{a_1, \ldots, a_r\})$ is closed. *Hint*: Show that if $y_k = \{A\mu_k\}$ and $y_k \to y$, then $y = A\mu$ with

$$\mu = \lim_{k \to \infty} \mu_k = \lim_{k \to \infty} (A'A)^{-1} A' y_k = (A'A)^{-1} A' y.$$

(b) Show that $cone(\{a_1, \ldots, a_r\})$ is closed without the linear independence assumption of part (a). *Hint*: Use Caratheodory's Theorem to show that $cone(\{a_1, \ldots, a_r\})$ is equal to the union of a finite number of cones generated by linearly independent vectors.

Solution: (a) Consider a sequence $\{y_k\} \subset \text{cone}(\{a_1, \ldots, a_r\})$ with $y_k \to y$. We will show that $y \in \text{cone}(\{a_1, \ldots, a_r\})$. For each k, we have $y_k = A\mu_k$ for some $\mu_k \geq 0$, from which we obtain,

$$A'y_k = A'A\mu_k.$$

Since a_1, \ldots, a_r are assumed linearly independent, the matrix A'A is invertible, and we have

$$\mu_k = (A'A)^{-1}A'y_k.$$

It follows that

$$\mu_k \to \mu$$
,

where

$$\mu = (A'A)^{-1}A'y.$$

Furthermore, since $\mu_k \geq 0$, we have $\mu \geq 0$. Taking the limit in the relation $y_k = A\mu_k$, we obtain $y = A\mu$ with $\mu \geq 0$, so $y \in \text{cone}(\{a_1, \dots, a_r\})$.

(b) By Caratheodory's Theorem, every vector in cone $(\{a_1, \ldots, a_r\})$ is a positive combination of linearly independent vectors. Thus, cone $(\{a_1, \ldots, a_r\})$ is the union of cone $(\{a_j \mid j \in J\})$ as J ranges over all subsets of $\{1, \ldots, r\}$ such that the set $\{a_j \mid j \in J\}$ is linearly independent. Each of these cones is closed by part (a), so their union is also closed.

2.16

Let P be a polyhedral set in \Re^n , with a Minkowski-Weyl Representation

$$P = \left\{ x \mid x = \sum_{j=1}^{m} \mu_j v_j + y, \ \sum_{j=1}^{m} \mu_j = 1, \ \mu_j \ge 0, \ j = 1, \dots, m, \ y \in C \right\},$$

where v_1, \ldots, v_m are some vectors in \Re^n and C is a finitely generated cone in \Re^n (cf. Prop. 2.3.3). Show that:

- (a) The recession cone of P is equal to C.
- (b) Each extreme point of P is equal to some vector v_i that cannot be represented as a convex combination of the vectors v_j with $v_j \neq v_i$.

Solution: (a) We first show that C is a subset of R_P , the recession cone of P. Let $\overline{y} \in C$, and choose any $\alpha \geq 0$ and $x \in P$ of the form $x = \sum_{j=1}^{m} \mu_j v_j$.

Since C is a cone, $\alpha \overline{y} \in C$, so that $x + \alpha \overline{y} \in P$ for all $\alpha \geq 0$. It follows that $\overline{y} \in R_P$. Hence $C \subset R_P$. Conversely, to show that $R_P \subset C$, let $\overline{y} \in R_P$ and take any $x \in P$. Then $x + k \overline{y} \in P$ for all $k \geq 1$. Since P = V + C, where $V = \text{conv}(\{v_1, \ldots, v_m\})$, it follows that

$$x + k\overline{y} = v^k + y^k, \qquad \forall \ k \ge 1,$$

with $v^k \in V$ and $y^k \in C$ for all $k \ge 1$. Because V is compact, the sequence $\{v^k\}$ has a limit point $v \in V$, and without loss of generality, we may assume that $v^k \to v$. Then

$$\lim_{k \to \infty} ||k\overline{y} - y^k|| = \lim_{k \to \infty} ||v^k - x|| = ||v - x||,$$

implying that

$$\lim_{k \to \infty} \left\| \overline{y} - (1/k)y^k \right\| = 0.$$

Therefore, the sequence $\{(1/k)y^k\}$ converges to \overline{y} . Since $y^k \in C$ for all $k \geq 1$, the sequence $\{(1/k)y^k\}$ is in C, and by the closedness of C, it follows that $\overline{y} \in C$. Hence, $R_P \subset C$.

(b) Any point in P has the form v + y with $v \in \text{conv}(\{v_1, \ldots, v_m\})$ and $y \in C$, or equivalently

$$v + y = \frac{1}{2}v + \frac{1}{2}(v + 2y),$$

with v and v + 2y being two distinct points in P if $y \neq 0$. Therefore, none of the points v + y, with $v \in \text{conv}(\{v_1, \ldots, v_m\})$ and $y \in C$, is an extreme point of P if $y \neq 0$. Hence, an extreme point of P must be in the set $\{v_1, \ldots, v_m\}$. Since by definition, an extreme point of P is not a convex combination of points in P, an extreme point of P must be equal to some v_i that cannot be expressed as a convex combination of the remaining vectors v_i , $j \neq i$.

2.17 (Compact Polyhedral Sets)

Show that a nonempty compact convex set is polyhedral if and only if it has a finite number of extreme points. Show by example that the compactness assumption is essential.

Solution: By the Minkowski-Weyl Representation Theorem (Prop. 2.3.3), a polyhedral set has a finite number of extreme points. Conversely, let P be a compact convex set having a finite number of extreme points $\{v_1, \ldots, v_m\}$. By the Krein-Milman Theorem (Exercise 2.2), a compact convex set is equal to the convex hull of its extreme points, so that $P = \text{conv}(\{v_1, \ldots, v_m\})$, which is a polyhedral set by Minkowski-Weyl Representation Theorem.

As an example showing that the assertion fails if compactness of the set is replaced by a weaker assumption that the set is closed and contains no lines, consider the set $D \subset \Re^3$ given by

$$D = \{(x_1, x_2, x_3) \mid x_1^2 + x_2^2 \le 1, \ x_3 = 1\}.$$

Let $C = \operatorname{cone}(D)$. It can seen that C is not a polyhedral set. On the other hand, C is closed, convex, does not contain a line, and has a unique extreme point at the origin.

[For a more formal argument, note that if C were polyhedral, then the set

$$D = C \cap \{(x_1, x_2, x_3) \mid x_3 = 1\}$$

would also be polyhedral by Prop. 2.3.4, since both C and $\{(x_1, x_2, x_3) \mid x_3 = 1\}$ are polyhedral sets. Thus, by Prop. 2.3.3, it would follow that D has a finite number of extreme points. But this is a contradiction because the set of extreme points of D coincides with $\{(x_1, x_2, x_3) \mid x_1^2 + x_2^2 = 1, x_3 = 1\}$, which contains an infinite number of points. Thus, C is not a polyhedral cone, and therefore not a polyhedral set, while C is closed, convex, does not contain a line, and has a unique extreme point at the origin.]

2.18 (Polyhedral Set Decomposition)

Show that a polyhedral set can be written as the vector sum of a subspace and a polyhedral set that contains at least one extreme point. *Hint*: Use the decomposition result of Prop. 1.4.4.

Solution: By the remarks following Prop. 1.4.4, a convex set C can be written as

$$C = L_C + (C \cap L_C^{\perp}),$$

where L_C is the lineality space of C. Furthermore, the set $C \cap L_C^{\perp}$ contains no lines. If C is polyhedral, then $C \cap L_C^{\perp}$ is also polyhedral and since it contains no lines, by Prop. 2.1.2, it must contain an extreme point.

2.19 (Cones Generated by Polyhedral Sets)

Show that if P is a polyhedral set in \Re^n containing the origin, then $\operatorname{cone}(P)$ is a polyhedral cone. Give an example showing that if P does not contain the origin, then $\operatorname{cone}(P)$ may not be a polyhedral cone.

Solution: We give two proofs. The first is based on the Minkowski-Weyl Representation of a polyhedral set P (cf. Prop. 2.3.3), while the second is based on a representation of P by a system of linear inequalities.

Let P be a polyhedral set with Minkowski-Weyl representation

$$P = \left\{ x \mid x = \sum_{j=1}^{m} \mu_j v_j + y, \ \sum_{j=1}^{m} \mu_j = 1, \ \mu_j \ge 0, \ j = 1, \dots, m, \ y \in C \right\},$$

where v_1, \ldots, v_m are some vectors in \Re^n and C is a finitely generated cone in \Re^n . Let C be given by

$$C = \left\{ y \mid y = \sum_{i=1}^{r} \lambda_i a_i, \ \lambda_i \ge 0, \ i = 1, \dots, r \right\},$$

where a_1, \ldots, a_r are some vectors in \Re^n , so that

$$P = \left\{ x \mid x = \sum_{j=1}^{m} \mu_j v_j + \sum_{i=1}^{r} \lambda_i a_i, \ \sum_{j=1}^{m} \mu_j = 1, \ \mu_j \ge 0, \ \forall \ j, \ \lambda_i \ge 0, \ \forall \ i \right\}.$$

We claim that

$$cone(P) = cone(\{v_1, \dots, v_m, a_1, \dots, a_r\}).$$

Since $P \subset \text{cone}(\{v_1, \ldots, v_m, a_1, \ldots, a_r\})$, it follows that

$$cone(P) \subset cone(\{v_1, \ldots, v_m, a_1, \ldots, a_r\}).$$

Conversely, let $y \in \text{cone}(\{v_1, \ldots, v_m, a_1, \ldots, a_r\})$. Then, we have

$$y = \sum_{j=1}^{m} \overline{\mu}_{j} v_{j} + \sum_{i=1}^{r} \overline{\lambda}_{i} a_{i},$$

with $\overline{\mu}_j \geq 0$ and $\overline{\lambda}_i \geq 0$ for all i and j. If $\overline{\mu}_j = 0$ for all j, then $y = \sum_{i=1}^r \overline{\lambda}_i a_i \in C$, and since $C = R_P$ (cf. Exercise 2.16), it follows that $y \in R_P$. Because the origin belongs to P and $y \in R_P$, we have $0 + y \in P$, implying that $y \in P$, and consequently $y \in \text{cone}(P)$. If $\overline{\mu}_j > 0$ for some j, then by setting $\overline{\mu} = \sum_{j=1}^m \overline{\mu}_j$, $\mu_j = \overline{\mu}_j/\overline{\mu}$ for all j, and $\lambda_i = \overline{\lambda}_i/\overline{\mu}$ for all i, we obtain

$$y = \overline{\mu} \left(\sum_{j=1}^{m} \mu_j v_j + \sum_{i=1}^{r} \lambda_i a_i \right),$$

where $\overline{\mu} > 0$, $\mu_j \ge 0$ with $\sum_{j=1}^m \mu_j = 1$, and $\lambda_i \ge 0$. Therefore $y = \overline{\mu} \overline{x}$ with $\overline{x} \in P$ and $\overline{\mu} > 0$, implying that $y \in \text{cone}(P)$ and showing that

$$\operatorname{cone}(\{v_1,\ldots,v_m,a_1,\ldots,a_r\})\subset\operatorname{cone}(P).$$

We now give an alternative proof using the representation of P by a system of linear inequalities. Let P be given by

$$P = \{x \mid a_j' x \le b_j, \ j = 1, \dots, r\},\$$

where a_1, \ldots, a_r are vectors in \Re^n and b_1, \ldots, b_r are scalars. Since P contains the origin, it follows that $b_j \geq 0$ for all j. Define the index set J as follows

$$J = \{j \mid b_j = 0\}.$$

We consider separately the two cases where $J \neq \emptyset$ and $J = \emptyset$. If $J \neq \emptyset$, then we will show that

$$cone(P) = \{x \mid a_j' x \le 0, \ j \in J\}.$$

To see this, note that since $P \subset \{x \mid a_i'x \leq 0, j \in J\}$, we have

$$cone(P) \subset \{x \mid a'_j x \le 0, \ j \in J\}.$$

Conversely, let $\overline{x} \in \{x \mid a'_j x \leq 0, \ j \in J\}$. We will show that $\overline{x} \in \text{cone}(P)$. If $\overline{x} \in P$, then $\overline{x} \in \text{cone}(P)$ and we are done, so assume that $\overline{x} \notin P$, implying that the set

$$\overline{J} = \{ j \notin J \mid a_i' \overline{x} > b_j \} \tag{2.13}$$

is nonempty. By the definition of J, we have $b_j > 0$ for all $j \notin J$, so let

$$\mu = \min_{j \in \overline{J}} \frac{b_j}{a_j' \overline{x}},$$

and note that $0 < \mu < 1$. We have

$$a_j'(\mu \overline{x}) \le 0, \quad \forall j \in J,$$

$$a'_j(\mu \overline{x}) \le b_j, \quad \forall j \in \overline{J}.$$

For $j \notin \overline{J} \cup J$ and $a'_j \overline{x} \leq 0 < b_j$, since $\mu > 0$, we still have $a'_j (\mu \overline{x}) \leq 0 < b_j$. For $j \notin \overline{J} \cup J$ and $0 < a'_j \overline{x} \leq b_j$, since $\mu < 1$, we have $0 < a'_j (\mu \overline{x}) < b_j$. Therefore, $\mu \overline{x} \in P$, implying that $\overline{x} = \frac{1}{\mu} (\mu \overline{x}) \in \text{cone}(P)$. It follows that

$$\{x \mid a_j'x \le 0, \ j \in J\} \subset \operatorname{cone}(P),$$

and hence, cone $(P) = \{x \mid a'_j x \le 0, j \in J\}.$

If $J = \emptyset$, then we will show that $\operatorname{cone}(P) = \Re^n$. To see this, take any $\overline{x} \in \Re^n$. If $\overline{x} \in P$, then clearly $\overline{x} \in \operatorname{cone}(P)$, so assume that $\overline{x} \notin P$, implying that the set \overline{J} as defined in Eq. (2.13) is nonempty. Note that

 $b_j > 0$ for all j, since J is empty. The rest of the proof is similar to the preceding case.

As an example, where cone(P) is not polyhedral when P does not contain the origin, consider the polyhedral set $P \subset \Re^2$ given by

$$P = \{(x_1, x_2) \mid x_1 \ge 0, \ x_2 = 1\}.$$

Then, we have

cone(P) =
$$\{(x_1, x_2) \mid x_1 > 0, x_2 > 0\} \cup \{(x_1, x_2) \mid x_1 = 0, x_2 \ge 0\},\$$

which is not closed and therefore not polyhedral.

2.20

Show that if P is a polyhedral set in \Re^n containing the origin, then $\operatorname{cone}(P)$ is a polyhedral cone. Give an example showing that if P does not contain the origin, then $\operatorname{cone}(P)$ may not be a polyhedral cone.

Solution: We give two proofs. The first is based on the Minkowski-Weyl Representation of a polyhedral set P (cf. Prop. 2.3.3), while the second is based on a representation of P by a system of linear inequalities.

Let P be a polyhedral set with Minkowski-Weyl representation

$$P = \left\{ x \mid x = \sum_{j=1}^{m} \mu_j v_j + y, \ \sum_{j=1}^{m} \mu_j = 1, \ \mu_j \ge 0, \ j = 1, \dots, m, \ y \in C \right\},$$

where v_1, \ldots, v_m are some vectors in \Re^n and C is a finitely generated cone in \Re^n . Let C be given by

$$C = \left\{ y \mid y = \sum_{i=1}^{r} \lambda_i a_i, \ \lambda_i \ge 0, \ i = 1, \dots, r \right\},$$

where a_1, \ldots, a_r are some vectors in \Re^n , so that

$$P = \left\{ x \mid x = \sum_{j=1}^{m} \mu_j v_j + \sum_{i=1}^{r} \lambda_i a_i, \ \sum_{j=1}^{m} \mu_j = 1, \ \mu_j \ge 0, \ \forall \ j, \ \lambda_i \ge 0, \ \forall \ i \right\}.$$

We claim that

$$cone(P) = cone(\{v_1, \dots, v_m, a_1, \dots, a_r\}).$$

Since $P \subset \text{cone}(\{v_1,\ldots,v_m,a_1,\ldots,a_r\})$, it follows that

$$cone(P) \subset cone(\{v_1, \dots, v_m, a_1, \dots, a_r\}).$$

Conversely, let $y \in \text{cone}(\{v_1, \ldots, v_m, a_1, \ldots, a_r\})$. Then, we have

$$y = \sum_{j=1}^{m} \overline{\mu}_{j} v_{j} + \sum_{i=1}^{r} \overline{\lambda}_{i} a_{i},$$

with $\overline{\mu}_j \geq 0$ and $\overline{\lambda}_i \geq 0$ for all i and j. If $\overline{\mu}_j = 0$ for all j, then $y = \sum_{i=1}^r \overline{\lambda}_i a_i \in C$, and since $C = R_P$ (cf. Exercise 2.16), it follows that $y \in R_P$. Because the origin belongs to P and $y \in R_P$, we have $0 + y \in P$, implying that $y \in P$, and consequently $y \in \text{cone}(P)$. If $\overline{\mu}_j > 0$ for some j, then by setting $\overline{\mu} = \sum_{j=1}^m \overline{\mu}_j$, $\mu_j = \overline{\mu}_j/\overline{\mu}$ for all j, and $\lambda_i = \overline{\lambda}_i/\overline{\mu}$ for all i, we obtain

$$y = \overline{\mu} \left(\sum_{j=1}^{m} \mu_j v_j + \sum_{i=1}^{r} \lambda_i a_i \right),$$

where $\overline{\mu} > 0$, $\mu_j \ge 0$ with $\sum_{j=1}^m \mu_j = 1$, and $\lambda_i \ge 0$. Therefore $y = \overline{\mu} \overline{x}$ with $\overline{x} \in P$ and $\overline{\mu} > 0$, implying that $y \in \text{cone}(P)$ and showing that

$$\operatorname{cone}(\{v_1,\ldots,v_m,a_1,\ldots,a_r\})\subset\operatorname{cone}(P).$$

We now give an alternative proof using the representation of P by a system of linear inequalities. Let P be given by

$$P = \{x \mid a_j' x \le b_j, \ j = 1, \dots, r\},\$$

where a_1, \ldots, a_r are vectors in \Re^n and b_1, \ldots, b_r are scalars. Since P contains the origin, it follows that $b_j \geq 0$ for all j. Define the index set J as follows

$$J = \{j \mid b_j = 0\}.$$

We consider separately the two cases where $J \neq \emptyset$ and $J = \emptyset$. If $J \neq \emptyset$, then we will show that

$$cone(P) = \{x \mid a_j' x \le 0, \ j \in J\}.$$

To see this, note that since $P \subset \{x \mid a_j'x \leq 0, j \in J\}$, we have

$$cone(P) \subset \{x \mid a'_j x \le 0, \ j \in J\}.$$

Conversely, let $\overline{x} \in \{x \mid a'_j x \leq 0, \ j \in J\}$. We will show that $\overline{x} \in \text{cone}(P)$. If $\overline{x} \in P$, then $\overline{x} \in \text{cone}(P)$ and we are done, so assume that $\overline{x} \notin P$, implying that the set

$$\overline{J} = \{ j \notin J \mid a_i' \overline{x} > b_i \} \tag{2.14}$$

is nonempty. By the definition of J, we have $b_j > 0$ for all $j \notin J$, so let

$$\mu = \min_{j \in \overline{J}} \frac{b_j}{a_j' \overline{x}},$$

and note that $0 < \mu < 1$. We have

$$a'_{j}(\mu \overline{x}) \leq 0, \quad \forall j \in J,$$

 $a'_{j}(\mu \overline{x}) \leq b_{j}, \quad \forall j \in \overline{J}.$

For $j \notin \overline{J} \cup J$ and $a'_j \overline{x} \leq 0 < b_j$, since $\mu > 0$, we still have $a'_j (\mu \overline{x}) \leq 0 < b_j$. For $j \notin \overline{J} \cup J$ and $0 < a'_j \overline{x} \leq b_j$, since $\mu < 1$, we have $0 < a'_j (\mu \overline{x}) < b_j$. Therefore, $\mu \overline{x} \in P$, implying that $\overline{x} = \frac{1}{\mu} (\mu \overline{x}) \in \text{cone}(P)$. It follows that

$$\{x \mid a_j'x \le 0, \ j \in J\} \subset \operatorname{cone}(P),$$

and hence, $cone(P) = \{x \mid a_i'x \leq 0, \ j \in J\}.$

If $J=\emptyset$, then we will show that $\operatorname{cone}(P)=\Re^n$. To see this, take any $\overline{x}\in\Re^n$. If $\overline{x}\in P$, then clearly $\overline{x}\in\operatorname{cone}(P)$, so assume that $\overline{x}\not\in P$, implying that the set \overline{J} as defined in Eq. (2.14) is nonempty. Note that $b_j>0$ for all j, since J is empty. The rest of the proof is similar to the preceding case.

As an example, where cone(P) is not polyhedral when P does not contain the origin, consider the polyhedral set $P \subset \Re^2$ given by

$$P = \{(x_1, x_2) \mid x_1 \ge 0, \ x_2 = 1\}.$$

Then, we have

cone(P) =
$$\{(x_1, x_2) \mid x_1 > 0, x_2 > 0\} \cup \{(x_1, x_2) \mid x_1 = 0, x_2 \ge 0\}$$
, which is not closed and therefore not polyhedral.

2.21 (Support Function of a Polyhedral Set)

Show that the support function of a polyhedral set is a polyhedral function.

Solution: Let X be a polyhedral set with Minkowski-Weyl representation

$$X = \operatorname{conv}(\{v_1, \dots, v_m\}) + \operatorname{cone}(\{d_1, \dots, d_r\})$$

for some vectors $v_1, \ldots, v_m, d_1, \ldots, d_r$ (cf. Prop. 2.3.3). The support function of X takes the form

$$\begin{split} \sigma_X(y) &= \sup_{x \in X} y'x \\ &= \sup_{\substack{\alpha_1, \dots, \alpha_m, \beta_1, \dots, \beta_r \geq 0 \\ \sum_{i=1}^m \alpha_i = 1}} \left\{ \sum_{i=1}^m \alpha_i v_i' y + \sum_{j=1}^r \beta_j d_j' y \right\} \\ &= \left\{ \max_{i=1, \dots, m} v_i' y & \text{if } d_j' y \leq 0, \ j = 1, \dots, r, \\ \infty & \text{otherwise.} \end{split} \right.$$

Thus the support function is polyhedral.

2.22 (Conjugate of a Polyhedral Function)

(a) Show that the conjugate of a function can be specified in terms of the support function of its epigraph with the formula

$$f^{\star}(y) = \sigma_{\operatorname{epi}(f)}(y, -1).$$

(b) Use part (a) to show that the conjugate of a polyhedral function is polyhedral.

Solution: (a) We have

$$f^{\star}(y) = \sup_{x \in \Re^n} \{ x'y - f(x) \},$$

which can equivalently be written as

$$f^{\star}(y) = \sup_{(x,w) \in \operatorname{epi}(f)} \{x'y - w\}.$$

Since the expression in braces in the right-hand side is the inner product of the vectors (x, w) and (y, -1), the supremum above is the value of the support function of epi(f) at (y, -1):

$$f^{\star}(y) = \sigma_{\text{epi}(f)}(y, -1), \quad \forall y \in \Re^n.$$

(See Fig. 2.2.)

(b) Let us apply the result of part (a) to the case where f is a polyhedral function, so that $\operatorname{epi}(f)$ is a polyhedral set. From Exercise 2.21, the support function $\sigma_{\operatorname{epi}(f)}$ is a polyhedral function, and it can be seen that $\sigma_{\operatorname{epi}(f)}(y,-1)$, viewed as a function of y, is polyhedral.

2.23 (Polar Sets)

This exercise introduces a notion of polar set that generalizes the notion of polar cone. Polar sets originated in Euclidean geometry, where they can be used to provide elegant proofs to many classical theorems. Given a nonempty set $C \subset \Re^n$, the *polar set of* C is defined as

$$C^{\circ} = \{y \mid y'x \leq 1, \forall \ x \in C\}.$$

Thus the polar set C° is the level set $\{y \mid \sigma_C(y) \leq 1\}$ of the support function σ_C of C. Since a single level set is sufficient to characterize all level

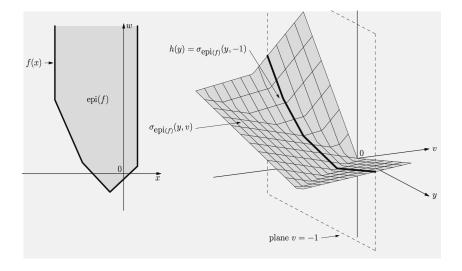


Figure 2.2. Construction of the conjugate of a function f from the support function $\sigma_{\text{epi}(f)}(y,v)$ of epi(f) (cf. Exercise 2.22). The conjugate is obtained by setting v=-1:

$$h(y) = \sigma_{\text{epi}(f)}(y, -1), \quad \forall y \in \Re^n.$$

If f is polyhedral as in the figure, then $\operatorname{epi}(f)$ and $\sigma_{\operatorname{epi}(f)}(y,v)$ are polyhedral, so the conjugate is also polyhedral.

sets of a support function (in view of positive homogeneity), it follows from the Conjugacy Theorem (Prop. 1.6.1), that any set is fully characterized by its polar up to convex closure, i.e., two sets with the same polar set have the same convex closure.

- (a) Show that C° is a closed convex set. Furthermore, C° is bounded if and only if the origin in an interior point of $\operatorname{conv}(C)$.
- (b) Show that the polar set of a cone is equal to its polar cone.
- (c) Consider the subset \hat{C} of \Re^{n+1} obtained from C via the lifting procedure,

$$\hat{C} = \left\{ (x, 1) \mid x \in C \right\}.$$

Show that C° is obtained from the polar of the cone generated by \hat{C} , by "slicing" at the level -1:

$$C^{\circ} = \big\{ y \mid (y, -1) \in \big(\operatorname{cone}(\hat{C}) \big)^* \big\}.$$

- (d) Show that if C is a finite set, then C° is a polyhedral set.
- (e) Show that

$$(C^{\circ})^{\circ} = \operatorname{cl}(\operatorname{conv}(\{0\} \cup C)),$$

so if C is a closed convex set containing the origin, we have $(C^{\circ})^{\circ} = C$.

(f) Consider a bounded polyhedral set P. For each extreme point v of P, consider the halfspace

$$H_v = \{ y \mid y'v \le 1 \}.$$

Show that the polar set P° is the intersection of the halfspaces H_v , where v ranges over the extreme points of P.

(g) Consider a circle in the plane that is centered at the origin, and a convex polygon that is inscribed in the circle and contains the origin in its interior. Show that the polar set is a polygon that can be circumscribed around some circle centered at the origin.

Solution: (a) Clearly, we have

$$0 \in \operatorname{int}(\operatorname{conv}(C)) \iff \sigma_C(y) > 0, \quad \forall y \neq 0.$$

Since σ_C is positively homogeneous, it is equal to its recession function, so

$$0 \in \operatorname{int}(\operatorname{conv}(C)) \iff R_{\sigma_C} = \{0\}.$$

Since $R_{\sigma_C} = \{0\}$ if and only if the nonempty level sets of σ_C are compact, and C° is a level set, we have

$$0 \in \operatorname{int}(\operatorname{conv}(C)) \iff \{y \mid \sigma_C(y) \leq 1\} = C^{\circ} \text{ is compact.}$$

- (b) If C is a cone, by Example 5.2.2, σ_C is the indicator function of the polar set C^* . Since C° is a nonempty level set of σ_C , it follows that $C^{\circ} = C^*$.
- (c) Using the definition of cone:

$$\mathrm{cone}(\hat{C}) = \big\{ (\lambda x, \lambda) \mid x \in C, \ \lambda > 0 \big\}.$$

Using the definition of polar cone:

$$(\operatorname{cone}(\hat{C}))^{\star} = \left\{ (y,w) \mid y'\lambda x + w'\lambda \leq 0, x \in C, \lambda > 0 \right\}$$

Therefore

$$\{y \mid (y, -1) \in (\operatorname{cone}(\hat{C}))^*\} = \{y \mid y'\lambda x - \lambda \le 0, x \in C, \lambda > 0\}$$

$$= \{y \mid y'x \le 1\}$$

$$= C^{\circ}$$

(d) If C is a finite set, C° is the intersection of a finite number of halfspaces. Furthermore, C° is nonempty since it contains the origin, so it is polyhedral.

(e) We first show that

$$C^{\circ} = (\operatorname{cl}(C)^{\circ}) = (\operatorname{conv}(C)^{\circ}) = (\{0\} \cap C)^{\circ}.$$

The first two equations hold because C, cl(C), and conv(C) have the same support function. The third equation is true by the definition of polar set.

Assume that C is closed, convex, and contains the origin. To show that $(C^{\circ})^{\circ}=C,$ note that

$$\begin{split} (C^{\circ})^{\circ} &= \{x \mid x'y \leq 1, \ \forall \ y \in C^{\circ} \} \\ &= \big\{ x \mid (x,1)'(y,-1) \leq 0, \ \forall \ (y,-1) \in D^{*} \big\}, \end{split}$$

where $D = \operatorname{cone}(\hat{C})$ and the second equation follows from part (c). Since D is a cone, its polar set is a cone by part (b). We write the above equation as

$$(C^{\circ})^{\circ} = \{x \mid (x,1)'(\lambda y, -\lambda) \leq 0, \ \forall \ (y,-1) \in D^*, \ \forall \ \lambda > 0)\},\$$

or equivalently,

$$(C^{\circ})^{\circ} = \{x \mid (x,1)'(\bar{y}) \le 0, \ \forall \ \bar{y} \in D^{\star}\},\$$

and note that

$$\{(\lambda x, \lambda) \mid (x, 1)' \bar{y} \le 0, \ \forall \ \bar{y} \in D^{\star}\} = (D^{*})^{*}$$

and $(D^*)^* = D$ because \hat{C} is closed and convex. Now it follows that

$$(C^{\circ})^{\circ} = \{x \mid (x,1) \in (D^{\star})^{\star} = D\}$$

where D is the cone of "lifted" C. Therefore

$$(C^{\circ})^{\circ} = C.$$

For an arbitrary set C without any assumption,

$$\operatorname{cl}(\operatorname{conv}(\{0\}\cap C)) = \left(\operatorname{cl}(\operatorname{conv}(\{0\}\cap C)^\circ)\right)^\circ = (C^\circ)^\circ.$$

(f) By the Minkowski-Weyl representation, a bounded polyhedral set P is the convex hull of its extreme points. Thus

$$P^{\circ} = \{ y \mid y'x < 1, \ \forall \ x \in \text{conv}(\{v_1, v_2, ..., v_r\}) \}.$$

We have

$$y'x \le 1$$
, $\forall x \in \text{conv}(\{v_1, ..., v_r\})$ \iff $y'v_i \le 1, \forall i = 1, ..., r$.

Therefore

$$P^{\circ} = \{y \mid y'x \leq 1, \ \forall \ x \in \mathrm{conv}(\{v_1, v_2, ..., v_r\})\} = H_{v_1} \cap ... \cap H_{v_r}.$$

(g) Using part (f), the polar set of the convex polygon P is the intersection of

$$H_v = \{ y \mid y'v \le 1 \},\$$

where v ranges over the extreme points of P. Furthermore, P° is bounded because 0 is an interior point of P, and it is polyhedral because it is the intersection of a finite number of halfspaces.

If P is inscribed in the circle

$${x \mid ||x|| = r},$$

all extreme points v satisfy ||v|| = r, and H_v corresponds to a tangent hyperplane of the circle centered at origin with radius 1/r. Thus, the intersection of H_v can be circumscribed around the circle $\{x \mid ||x|| = 1/r\}$.

SECTION 2.4: Polyhedral Aspects of Optimization

2.24 (Gordan's Theorem of the Alternative [Gor1873])

Let a_1, \ldots, a_r be vectors in \Re^n .

- (a) Show that exactly one of the following two conditions holds:
 - (i) There exists a vector $x \in \Re^n$ such that

$$a_1'x < 0, \dots, a_r'x < 0.$$

(ii) There exists a vector $\mu \in \Re^r$ such that $\mu \neq 0$, $\mu \geq 0$, and

$$\mu_1 a_1 + \dots + \mu_r a_r = 0.$$

(b) Show that an equivalent statement of part (a) is the following: a polyhedral cone has nonempty interior if and only if its polar cone does not contain a line, i.e., a set of the form $\{x + \alpha z \mid \alpha \in \Re\}$, where x lies in the polar cone and z is a nonzero vector.

Note: This result is also given with an alternative proof in Section 5.6.

Solution: (a) Assume that there exist $\hat{x} \in \mathbb{R}^n$ and $\mu \in \mathbb{R}^r$ such that both conditions (i) and (ii) hold, i.e.,

$$a_j'\hat{x} < 0, \qquad \forall \ j = 1, \dots, r, \tag{2.15}$$

$$\mu \neq 0, \qquad \mu \geq 0, \qquad \sum_{j=1}^{r} \mu_j a_j = 0.$$
 (2.16)

By premultiplying Eq. (2.15) with $\mu_j \geq 0$ and summing the obtained inequalities over j, we have

$$\sum_{j=1}^{r} \mu_j a_j' \hat{x} < 0.$$

On the other hand, from Eq. (2.16), we obtain

$$\sum_{j=1}^{r} \mu_j a_j' \hat{x} = 0,$$

which is a contradiction. Hence, both conditions (i) and (ii) cannot hold simultaneously.

The proof will be complete if we show that the conditions (i) and (ii) cannot *fail* to hold simultaneously. Assume that condition (i) fails to hold, and consider the sets given by

$$C_1 = \left\{ w \in \Re^r \mid a_j' x \le w_j, \ j = 1, \dots, r, \ x \in \Re^n \right\},$$
$$C_2 = \left\{ \xi \in \Re^r \mid \xi_j < 0, \ j = 1, \dots, r \right\}.$$

It can be seen that both C_1 and C_2 are convex. Furthermore, because the condition (i) does not hold, C_1 and C_2 are disjoint sets. Therefore, by the Separating Hyperplane Theorem (Prop. 1.5.2), C_1 and C_2 can be separated, i.e., there exists a nonzero vector $\mu \in \Re^r$ such that

$$\mu'w \ge \mu'\xi, \quad \forall \ w \in C_1, \quad \forall \ \xi \in C_2,$$

implying that

$$\inf_{w \in C_1} \mu' w \ge \mu' \xi, \qquad \forall \ \xi \in C_2.$$

Since each component ξ_j of $\xi \in C_2$ can be any negative scalar, for the preceding relation to hold, μ_j must be nonnegative for all j. Furthermore, by letting $\xi \to 0$, in the preceding relation, it follows that

$$\inf_{w \in C_1} \mu' w \ge 0,$$

implying that

$$\mu_1 w_1 + \dots + \mu_r w_r \ge 0, \quad \forall \ w \in C_1.$$

By setting $w_j = a'_i x$ for all j, we obtain

$$(\mu_1 a_1 + \dots + \mu_r a_r)' x \ge 0, \qquad \forall \ x \in \Re^n,$$

and because this relation holds for all $x \in \Re^n$, we must have

$$\mu_1 a_1 + \cdots + \mu_r a_r = 0.$$

Hence, the condition (ii) holds, showing that the conditions (i) and (ii) cannot fail to hold simultaneously.

Alternative proof: We will show the equivalent statement of part (b), i.e., that a polyhedral cone contains an interior point if and only if the polar C^* does not contain a line.

Let

$$C = \{x \mid a_i' x \le 0, j = 1, \dots, r\},\$$

where $a_j \neq 0$ for all j. Assume that C contains an interior point, and to arrive at a contradiction, assume that C^* contains a line. Then there exists a $d \neq 0$ such that d and -d belong to C^* , i.e., $d'x \leq 0$ and $-d'x \leq 0$ for all $x \in C$, so that d'x = 0 for all $x \in C$. Thus for the interior point $\overline{x} \in C$, we have $d'\overline{x} = 0$, and since $d \in C^*$ and $d = \sum_{j=1}^r \mu_j a_j$ for some $\mu_j \geq 0$, we have

$$\sum_{j=1}^{r} \mu_j a_j' \overline{x} = 0.$$

This is a contradiction, since \overline{x} is an interior point of C, and we have $a'_{i}\overline{x} < 0$ for all j.

Conversely, assume that C^* does not contain a line. Then C^* has an extreme point, and since the origin is the only possible extreme point of a cone, it follows that the origin is an extreme point of C^* , which is the cone generated by $\{a_1, \ldots, a_r\}$. Therefore $0 \notin \text{conv}(\{a_1, \ldots, a_r\})$, and there exists a hyperplane that strictly separates the origin from $\text{conv}(\{a_1, \ldots, a_r\})$. Thus, there exists a vector x such that y'x < 0 for all $y \in \text{conv}(\{a_1, \ldots, a_r\})$, so in particular,

$$a_i'x < 0, \quad \forall j = 1, \dots, r,$$

and x is an interior point of C.

(b) Let C be a polyhedral cone given by

$$C = \{x \mid a_j' x \le 0, j = 1, \dots, r\},\$$

where $a_j \neq 0$ for all j. The interior of C is given by

$$\operatorname{int}(C) = \big\{ x \mid a_j' x < 0, \ j = 1, \dots, r \big\},\,$$

so that C has nonempty interior if and only if the condition (i) of part (a) holds.

By Farkas' Lemma, the polar cone of C is given by

$$C^* = \left\{ x \mid x = \sum_{j=1}^r \mu_j a_j, \mu_j \ge 0, \ j = 1, \dots, r \right\}.$$

We now show that C^* contains a line if and only if there is a $\mu \in \Re^r$ such that $\mu \neq 0$, $\mu \geq 0$, and $\sum_{j=1}^r \mu_j a_j = 0$ [condition (ii) of part (a) holds]. Suppose that C^* contains a line, i.e., a set of the form $\{x + \alpha z \mid \alpha \in \Re\}$, where $x \in C^*$ and z is a nonzero vector. Since C^* is a closed convex cone, by the Recession Cone Theorem (Prop. 1.4.1), it follows that z and -z belong to R_{C^*} . This, implies that $0 + z = z \in C^*$ and $0 - z = -z \in C^*$, and therefore z and -z can be represented as

$$z = \sum_{j=1}^{r} \mu_j a_j, \quad \forall j, \ \mu_j \ge 0, \ \mu_j \ne 0 \text{ for some } j,$$

$$-z = \sum_{j=1}^r \overline{\mu}_j a_j, \qquad \forall \ j, \ \overline{\mu}_j \geq 0, \ \overline{\mu}_j \neq 0 \ \text{for some} \ j.$$

Thus, $\sum_{j=1}^{r} (\mu_j + \overline{\mu}_j) a_j = 0$, where $(\mu_j + \overline{\mu}_j) \ge 0$ for all j and $(\mu_j + \overline{\mu}_j) \ne 0$ for at least one j, showing that the condition (ii) of part (a) holds.

Conversely, suppose that $\sum_{j=1}^{r} \mu_j a_j = 0$ with $\mu_j \geq 0$ for all j and $\mu_j \neq 0$ for some j. Assume without loss of generality that $\mu_1 > 0$, so that

$$-a_1 = \sum_{j \neq 1} \frac{\mu_j}{\mu_1} a_j,$$

with $\mu_j/\mu_1 \geq 0$ for all j, which implies that $-a_1 \in C^*$. Since $a_1 \in C^*$, $-a_1 \in C^*$, and $a_1 \neq 0$, it follows that C^* contains a line, completing the proof.

2.25 (Linear System Alternatives)

Let a_1, \ldots, a_r be vectors in \Re^n and let $b_1, \ldots b_r$ be scalars. Show that exactly one of the following two conditions holds:

(i) There exists a vector $x \in \Re^n$ such that

$$a_1'x < b_1, \ldots, a_r'x < b_r$$
.

(ii) There exists a vector $\mu \in \Re^r$ such that $\mu \geq 0$ and

$$\mu_1 a_1 + \dots + \mu_r a_r = 0, \qquad \mu_1 b_1 + \dots + \mu_r b_r < 0.$$

Note: This result is a special case of Motzkin's Transposition Theorem, given with an alternative proof in Section 5.6.

Solution: Assume that there exist $\hat{x} \in \Re^n$ and $\mu \in \Re^r$ such that both conditions (i) and (ii) hold, i.e.,

$$a_i'\hat{x} \le b_j, \qquad \forall \ j = 1, \dots, r,$$
 (2.17)

$$\mu \ge 0, \qquad \sum_{j=1}^{r} \mu_j a_j = 0, \qquad \sum_{j=1}^{r} \mu_j b_j < 0.$$
 (2.18)

By premultiplying Eq. (2.17) with $\mu_j \geq 0$ and summing the obtained inequalities over j, we have

$$\sum_{j=1}^r \mu_j a_j' \hat{x} \le \sum_{j=1}^r \mu_j b_j.$$

On the other hand, by using Eq. (2.18), we obtain

$$\sum_{j=1}^{r} \mu_j a_j' \hat{x} = 0 > \sum_{j=1}^{r} \mu_j b_j,$$

which is a contradiction. Hence, both conditions (i) and (ii) cannot hold simultaneously.

The proof will be complete if we show that conditions (i) and (ii) cannot *fail* to hold simultaneously. Assume that condition (i) fails to hold, and consider the sets given by

$$P_1 = \{ \xi \in \Re^r \mid \xi_i \le 0, \ j = 1, \dots, r \},\$$

$$P_2 = \{ w \in \Re^r \mid a_i' x - b_j = w_j, \ j = 1, \dots, r, \ x \in \Re^n \}.$$

Clearly, P_1 is a polyhedral set. For the set P_2 , we have

$$P_2 = \{ w \in \Re^r \mid Ax - b = w, \ x \in \Re^n \} = R(A) - b,$$

where A is the matrix with rows a'_j and b is the vector with components b_j . Thus, P_2 is an affine set and is therefore polyhedral. Furthermore, because the condition (i) does not hold, P_1 and P_2 are disjoint polyhedral sets, and they can be strictly separated [Prop. 1.5.3 under condition (3)]. Hence, there exists a vector $\mu \in \Re^r$ such that

$$\sup_{\xi \in P_1} \mu' \xi < \inf_{w \in P_2} \mu' w.$$

Since each component ξ_j of $\xi \in P_1$ can be any negative scalar, for the preceding relation to hold, μ_j must be nonnegative for all j. Furthermore, since $0 \in P_1$, it follows that

$$0 < \inf_{w \in P_2} \mu' w,$$

implying that

$$0 < \mu_1 w_1 + \dots + \mu_r w_r, \qquad \forall \ w \in P_2.$$

By setting $w_j = a'_j x - b_j$ for all j, we obtain

$$\mu_1 b_1 + \dots + \mu_r b_r < (\mu_1 a_1 + \dots + \mu_r a_r)' x, \quad \forall x \in \Re^n.$$

Since this relation holds for all $x \in \mathbb{R}^n$, we must have

$$\mu_1 a_1 + \dots + \mu_r a_r = 0,$$

implying that

$$\mu_1 b_1 + \dots + \mu_r b_r < 0.$$

Hence, the condition (ii) holds, showing that the conditions (i) and (ii) cannot fail to hold simultaneously.

2.26 (Integer Programming and Unimodular Matrices)

Integer programming problems are optimization problems, which as part of their constraints include the requirement that the optimization variables take integer values, such as 0 or 1. An important method for solving such problems relies on the solution of a continuous optimization problem, called the relaxed problem, which is derived from the original by neglecting the integer constraints while maintaining all the other constraints. If the relaxed problem happens to have integer components, it will then solve optimally not just the relaxed problem, but also the original integer programming problem. Thus, polyhedral sets whose extreme points have integer components are of special significance. We will characterize an important class of such sets.

Let us say that a square matrix with integer components is unimodular if its determinant is 0, 1, or -1, and let us say that a rectangular matrix with integer components is totally unimodular if each of its square submatrices is unimodular. If A is an invertible matrix, by Cramer's rule, its inverse A^{-1} has components of the form

$$\left[A^{-1}\right]_{ij} = \frac{\text{polynomial in the components of } A}{\text{determinant of } A}.$$

It follows that if A is an invertible matrix with integer components that is unimodular, its inverse has integer components. Furthermore, for any vector b with integer components, the unique solution $A^{-1}b$ of the system

$$Ax = b$$

has integer components.

Let P be a polyhedral set of the form

$$P = \{x \mid Ax = b, c \le x \le d\},\$$

where A is an $m \times n$ matrix, b is a vector in \Re^m , and c and d are vectors in \Re^n . Assume that all the components of A, b, c, and d are integer, and that A is totally unimodular. Show that all the extreme points of P have integer components.

Solution: Let v be an extreme point of P. Consider the subset of indices

$$I = \{i \mid c_i < v_i < d_i\},\$$

and without loss of generality, assume that

$$I = \{1, \dots, \overline{m}\}$$

for some integer \overline{m} . Let \overline{A} be the matrix consisting of the first \overline{m} columns of A and let \overline{v} be the vector consisting of the first \overline{m} components of v. Note that each of the last $n-\overline{m}$ components of v is equal to either the corresponding component of c or to the corresponding component of d, which are integer. Thus the extreme point v has integer components if and only if the subvector \overline{v} has integer components.

By Prop. 2.1.4, \overline{A} has linearly independent columns, so \overline{v} is the unique solution of the system of equations

$$\overline{A}y = \overline{b},$$

where \overline{b} is equal to b minus the last $n-\overline{m}$ columns of A multiplied with the corresponding components of v (each of which is equal to either the corresponding component of c or the corresponding component of d, so that \overline{b} has integer components). Equivalently, there exists an invertible $\overline{m} \times \overline{m}$ submatrix \tilde{A} of \overline{A} and a subvector \tilde{b} of \overline{b} with \overline{m} components such that

$$\overline{v} = (\tilde{A})^{-1}\tilde{b}.$$

Since by hypothesis, A is totally unimodular, the invertible submatrix \tilde{A} is unimodular, and it follows that \overline{v} (and hence also the extreme point v) has integer components.

2.27 (Unimodularity I)

Let A be an $n \times n$ invertible matrix with integer entries. Show that A is unimodular if and only if the solution of the system Ax = b has integer components for every vector $b \in \Re^n$ with integer components. Hint: To prove that A is unimodular when the given property holds, use the system $Ax = u_i$, where u_i is the *i*th unit vector, to show that A^{-1} has integer components, and then use the equality $\det(A) \cdot \det(A^{-1}) = 1$. To prove the converse, use Cramer's rule.

Solution: Suppose that the system Ax = b has integer components for every vector $b \in \mathbb{R}^n$ with integer components. Since A is invertible, it follows that the vector $A^{-1}b$ has integer components for every $b \in \mathbb{R}^n$ with integer components. For $i = 1, \ldots, n$, let e_i be the vector with ith component equal to 1 and all other components equal to 0. Then, for $b = e_i$, the vectors $A^{-1}e_i$, $i = 1, \ldots, n$, have integer components, implying that the columns of A^{-1} are vectors with integer components, so that A^{-1} has integer entries. Therefore, $\det(A^{-1})$ is integer, and since $\det(A)$ is also integer and $\det(A) \cdot \det(A^{-1}) = 1$, it follows that either $\det(A) = 1$ or $\det(A) = -1$, showing that A is unimodular.

Suppose now that A is unimodular. Take any vector $b \in \mathbb{R}^n$ with integer components, and for each $i \in \{1, ..., n\}$, let A_i be the matrix obtained from A by replacing the ith column of A with b. Then, according to Cramer's rule, the components of the solution \hat{x} of the system Ax = b are given by

$$\hat{x}_i = \frac{\det(A_i)}{\det(A)}, \quad i = 1, \dots, n.$$

Since each matrix A_i has integer entries, it follows that $\det(A_i)$ is integer for all i = 1, ..., n. Furthermore, because A is invertible and unimodular, we have either $\det(A) = 1$ or $\det(A) = -1$, implying that the vector \hat{x} has integer components.

2.28 (Unimodularity II)

Let A be an $m \times n$ matrix.

- (a) Show that A is totally unimodular if and only if its transpose A' is totally unimodular.
- (b) Show that A is totally unimodular if and only if every subset J of $\{1, \ldots, n\}$ can be partitioned into two subsets J_1 and J_2 such that

$$\left| \sum_{j \in J_1} a_{ij} - \sum_{j \in J_2} a_{ij} \right| \le 1, \quad \forall i = 1, \dots, m.$$

Solution: (a) The proof is straightforward from the definition of the totally unimodular matrix and the fact that B is a submatrix of A if and only if B' is a submatrix of A'.

(b) Suppose that A is totally unimodular. Let J be a subset of $\{1, \ldots, n\}$. Define z by $z_j = 1$ if $j \in J$, and $z_j = 0$ otherwise. Also let w = Az, $c_i = d_i = \frac{1}{2}w_i$ if w_i is even, and $c_i = \frac{1}{2}(w_i - 1)$ and $d_i = \frac{1}{2}(w_i + 1)$ if w_i is odd. Consider the polyhedral set

$$P = \{x \mid c \le Ax \le d, \ 0 \le x \le z\},\$$

and note that $P \neq \emptyset$ because $\frac{1}{2}z \in P$. Since A is totally unimodular, the polyhedron P has integer extreme points. Let $\hat{x} \in P$ be one of them. Because $0 \leq \hat{x} \leq z$ and \hat{x} has integer components, it follows that $\hat{x}_j = 0$ for $j \notin J$ and $\hat{x}_j \in \{0,1\}$ for $j \in J$. Therefore, $z_j - 2\hat{x}_j = \pm 1$ for $j \in J$. Define $J_1 = \{j \in J \mid z_j - 2\hat{x}_j = 1\}$ and $J_2 = \{j \in J \mid z_j - 2\hat{x}_j = -1\}$. We have

$$\sum_{j \in J_1} a_{ij} - \sum_{j \in J_2} a_{ij} = \sum_{j \in J} a_{ij} (z_j - 2\hat{x}_j)$$

$$= \sum_{j=1}^n a_{ij} (z_j - 2\hat{x}_j)$$

$$= [Az]_i - 2[A\hat{x}]_i$$

$$= w_i - 2[A\hat{x}]_i,$$

where $[Ax]_i$ denotes the *i*th component of the vector Ax. If w_i is even, then since $c_i \leq [A\hat{x}]_i \leq d_i$ and $c_i = d_i = \frac{1}{2}w_i$, it follows that $[A\hat{x}]_i = w_i$, so that

$$w_i - 2[A\hat{x}]_i = 0,$$
 when w_i is even.

If w_i is odd, then since $c_i \leq [A\hat{x}]_i \leq d_i$, $c_i = \frac{1}{2}(w_i - 1)$, and $d_i = \frac{1}{2}(w_i + 1)$, it follows that

$$\frac{1}{2}(w_i - 1) \le [A\hat{x}]_i \le \frac{1}{2}(w_i + 1),$$

implying that

$$-1 \le w_i - 2[A\hat{x}]_i \le 1.$$

Because $w_i - 2[A\hat{x}]_i$ is integer, we conclude that

$$w_i - 2[A\hat{x}]_i \in \{-1, 0, 1\},$$
 when w_i is odd.

Therefore,

$$\left| \sum_{j \in J_1} a_{ij} - \sum_{j \in J_2} a_{ij} \right| \le 1, \quad \forall \ i = 1, \dots, m.$$
 (2.19)

Suppose now that the matrix A is such that any $J \subset \{1, \ldots, n\}$ can be partitioned into two subsets so that Eq. (2.19) holds. We prove that A is totally unimodular, by showing that each of its square submatrices is unimodular, i.e., the determinant of every square submatrix of A is -1, 0, or 1. We use induction on the size of the square submatrices of A.

To start the induction, note that for $J \subset \{1, \ldots, n\}$ with J consisting of a single element, from Eq. (2.19) we obtain $a_{ij} \in \{-1, 0, 1\}$ for all i and j. Assume now that the determinant of every $(k-1) \times (k-1)$ submatrix of A is -1, 0, or 1. Let B be a $k \times k$ submatrix of A. If $\det(B) = 0$, then we are done, so assume that B is invertible. Our objective is to prove that $|\det B| = 1$. By Cramer's rule and the induction hypothesis, we have

 $B^{-1}=\frac{B^*}{\det(B)},$ where $b^*_{ij}\in\{-1,0,1\}.$ By the definition of $B^*,$ we have $Bb^*_1=\det(B)e_1,$ where b^*_1 is the first column of B^* and $e_1=(1,0,\ldots 0)'.$ Let $J=\{j\mid b^*_{j1}\neq 0\}$ and note that $J\neq \varnothing$ since B is invertible. Let $\overline{J}_1=\{j\in J\mid b^*_{j1}=1\}$ and $\overline{J}_2=\{j\in J\mid j\not\in \overline{J}_1\}.$ Then, since $[Bb^*_1]_i=0$ for $i=2,\ldots,k,$ we have

$$[Bb_1^*]_i = \sum_{j=1}^k b_{ij}b_{j1}^* = \sum_{j \in \overline{J}_1} b_{ij} - \sum_{j \in \overline{J}_2} b_{ij} = 0, \quad \forall i = 2, \dots, k.$$

Thus, the cardinality of the set J is even, so that for any partition $(\tilde{J}_1, \tilde{J}_2)$ of J, it follows that $\sum_{j \in \tilde{J}_1} b_{ij} - \sum_{j \in \tilde{J}_2} b_{ij}$ is even for all $i = 2, \ldots, k$. By assumption, there is a partition (J_1, J_2) of J such that

$$\left| \sum_{j \in J_1} b_{ij} - \sum_{j \in J_2} b_{ij} \right| \le 1 \qquad \forall \ i = 1, \dots, k, \tag{2.20}$$

implying that

$$\sum_{j \in J_1} b_{ij} - \sum_{j \in J_2} b_{ij} = 0, \qquad \forall \ i = 2, \dots, k.$$
 (2.21)

Consider now the value

$$\alpha = \left| \sum_{j \in J_1} b_{1j} - \sum_{j \in J_2} b_{1j} \right|,$$

for which in view of Eq. (2.20), we have either $\alpha=0$ or $\alpha=1$. Define $y\in\Re^k$ by $y_i=1$ for $i\in J_1,\ y_i=-1$ for $i\in J_2$, and $y_i=0$ otherwise. Then, we have $|[By]_1|=\alpha$ and by Eq. (2.21), $[By]_i=0$ for all $i=2,\ldots,k$. If $\alpha=0$, then By=0 and since B is invertible, it follows that y=0, implying that $J=\emptyset$, which is a contradiction. Hence, we must have $\alpha=1$ so that $By=\pm e_1$. Without loss of generality assume that $By=e_1$ (if $By=-e_1$, we can replace y by -y). Then, since $Bb_1^*=\det(B)e_1$, we see that $B(b_1^*-\det(B)y)=0$ and since B is invertible, we must have $b_1^*=\det(B)y$. Because y and b_1^* are vectors with components -1, 0, or 1, it follows that $b_1^*=\pm y$ and $|\det(B)|=1$, completing the induction and showing that A is totally unimodular.

2.29 (Unimodularity III)

Show that a matrix A is totally unimodular if one of the following holds:

(a) The entries of A are -1, 0, or 1, and there are exactly one 1 and exactly one -1 in each of its columns.

(b) The entries of A are 0 or 1, and in each of its columns, the entries that are equal to 1 appear consecutively.

Solution: (a) We show that the determinant of any square submatrix of A is -1, 0, or 1. We prove this by induction on the size of the square submatrices of A. In particular, the 1×1 submatrices of A are the entries of A, which are -1, 0, or 1. Suppose that the determinant of each $(k-1) \times (k-1)$ submatrix of A is -1, 0, or 1, and consider a $k \times k$ submatrix B of A. If B has a zero column, then $\det(B) = 0$ and we are done. If B has a column with a single nonzero component (1 or -1), then by expanding its determinant along that column and by using the induction hypothesis, we see that $\det(B) = 1$ or $\det(B) = -1$. Finally, if each column of B has exactly two nonzero components (one 1 and one -1), the sum of its rows is zero, so that B is singular and $\det(B) = 0$, completing the proof and showing that A is totally unimodular.

(b) The proof is based on induction as in part (a). The 1×1 submatrices of A are the entries of A, which are 0 or 1. Suppose now that the determinant of each $(k-1) \times (k-1)$ submatrix of A is -1, 0, or 1, and consider a $k \times k$ submatrix B of A. Since in each column of A, the entries that are equal to 1 appear consecutively, the same is true for the matrix B. Take the first column b_1 of B. If $b_1 = 0$, then B is singular and det(B) = 0. If b_1 has a single nonzero component, then by expanding the determinant of Balong b_1 and by using the induction hypothesis, we see that det(B) = 1 or det(B) = -1. Finally, let b_1 have more than one nonzero component (its nonzero entries are 1 and appear consecutively). Let l and p be rows of Bsuch that $b_{i1} = 0$ for all i < l and i > p, and $b_{i1} = 1$ for all $l \le i \le p$. By multiplying the lth row of B with (-1) and by adding it to the l+1st, l+2nd, ..., kth row of B, we obtain a matrix \overline{B} such that $\det(B) = \det(\overline{B})$ and the first column \overline{b}_1 of \overline{B} has a single nonzero component. Furthermore, the determinant of every square submatrix of B is -1, 0, or 1 (this follows from the fact that the determinant of a square matrix is unaffected by adding a scalar multiple of a row of the matrix to some of its other rows, and from the induction hypothesis). Since \overline{b}_1 has a single nonzero component, by expanding the determinant of \overline{B} along \overline{b}_1 , it follows that $\det(\overline{B}) = 1$ or $\det(\overline{B}) = -1$, implying that $\det(B) = 1$ or $\det(B) = -1$, completing the induction and showing that A is totally unimodular.

2.30 (Unimodularity IV)

Let A be a matrix with entries -1, 0, or 1, and exactly two nonzero entries in each of its columns. Show that A is totally unimodular if and only if the rows of A can be divided into two subsets such that for each column the following hold: if the two nonzero entries in the column have the same

sign, their rows are in different subsets, and if they have the opposite sign, their rows are in the same subset.

Solution: If A is totally unimodular, then by Exercise 2.28(a), its transpose A' is also totally unimodular, and by Exercise 2.28(b), the set $I = \{1, \ldots, m\}$ can be partitioned into two subsets I_1 and I_2 such that

$$\left| \sum_{i \in I_1} a_{ij} - \sum_{i \in I_2} a_{ij} \right| \le 1, \quad \forall j = 1, \dots, n.$$

Since $a_{ij} \in \{-1, 0, 1\}$ and exactly two of a_{1j}, \ldots, a_{mj} are nonzero for each j, it follows that

$$\sum_{i \in I_1} a_{ij} - \sum_{i \in I_2} a_{ij} = 0, \quad \forall \ j = 1, \dots, n.$$

Take any $j \in \{1, ..., n\}$, and let l and p be such that $a_{ij} = 0$ for all $i \neq l$ and $i \neq p$, so that in view of the preceding relation and the fact $a_{ij} \in \{-1, 0, 1\}$, we see that: if $a_{lj} = -a_{pj}$, then both l and p are in the same subset (I_1 or I_2); if $a_{lj} = a_{pj}$, then l and p are not in the same subset.

Suppose now that the rows of A can be divided into two subsets such that for each column the following property holds: if the two nonzero entries in the column have the same sign, they are in different subsets, and if they have the opposite sign, they are in the same subset. By multiplying all the rows in one of the subsets by -1, we obtain the matrix \overline{A} with entries $\overline{a}_{ij} \in \{-1,0,1\}$, and exactly one 1 and exactly one -1 in each of its columns. Therefore, by Exercise 2.29(a), \overline{A} is totally unimodular, so that every square submatrix of \overline{A} has determinant -1, 0, or 1. Since the determinant of a square submatrix of \overline{A} and the determinant of the corresponding submatrix of A differ only in sign, it follows that every square submatrix of A has determinant -1, 0, or 1, showing that A is totally unimodular.

2.31 (Elementary Vectors [Roc69])

Given a vector $z = (z_1, ..., z_n)$ in \Re^n , the *support* of z is the set of indices $\{j \mid z_j \neq 0\}$. We say that a nonzero vector z of a subspace S of \Re^n is *elementary* if there is no vector $\overline{z} \neq 0$ in S that has smaller support than z, i.e., for all nonzero $\overline{z} \in S$, $\{j \mid \overline{z}_j \neq 0\}$ is not a strict subset of $\{j \mid z_j \neq 0\}$. Show that:

- (a) Two elementary vectors with the same support are scalar multiples of each other.
- (b) For every nonzero vector y, there exists an elementary vector with support contained in the support of y.

(c) (Conformal Realization Theorem) We say that a vector x is in harmony with a vector z if

$$x_j z_j \ge 0, \quad \forall j = 1, \dots, n.$$

Show that every nonzero vector x of a subspace S can be written in the form

$$x = z^1 + \ldots + z^m.$$

where z^1, \ldots, z^m are elementary vectors of S, and each of them is in harmony with x and has support contained in the support of x. Note: Among other subjects, this result finds significant application in network optimization algorithms (see Rockafellar [Roc69] and Bertsekas [Ber98]).

Solution: (a) If two elementary vectors z and \overline{z} had the same support, the vector $z - \gamma \overline{z}$ would be nonzero and have smaller support than z and \overline{z} for a suitable scalar γ . If z and \overline{z} are not scalar multiples of each other, then $z - \gamma \overline{z} \neq 0$, which contradicts the definition of an elementary vector.

- (b) We note that either y is elementary or else there exists a nonzero vector \overline{z} with support strictly contained in the support of y. Repeating this argument for at most n-1 times, we must obtain an elementary vector.
- (c) We first show that every nonzero vector $y \in S$ has the property that there exists an elementary vector of S that is in harmony with y and has support that is contained in the support of y.

We show this by induction on the number of nonzero components of y. Let V_k be the subset of nonzero vectors in S that have k or less nonzero components, and let \overline{k} be the smallest k for which V_k is nonempty. Then, by part (b), every vector $y \in V_{\overline{k}}$ must be elementary, so it has the desired property. Assume that all vectors in V_k have the desired property for some $k \geq \overline{k}$. We let y be a vector in V_{k+1} and we show that it also has the desired property. Let z be an elementary vector whose support is contained in the support of y. By using the negative of z if necessary, we can assume that $y_j z_j > 0$ for at least one index j. Then there exists a largest value of γ , call it $\overline{\gamma}$, such that

$$y_j - \gamma z_j \ge 0, \quad \forall j \text{ with } y_j > 0,$$

$$y_j - \gamma z_j \le 0, \quad \forall j \text{ with } y_j < 0.$$

The vector $y - \overline{\gamma}z$ is in harmony with y and has support that is strictly contained in the support of y. Thus either $y - \overline{\gamma}z = 0$, in which case the elementary vector z is in harmony with y and has support equal to the support of y, or else $y - \overline{\gamma}z$ is nonzero. In the latter case, we have

 $y - \overline{\gamma}z \in V_k$, and by the induction hypothesis, there exists an elementary vector \overline{z} that is in harmony with $y - \overline{\gamma}z$ and has support that is contained in the support of $y - \overline{\gamma}z$. The vector \overline{z} is also in harmony with y and has support that is contained in the support of y. The induction is complete.

Consider now the given nonzero vector $x \in S$, and choose any elementary vector \overline{z}^1 of S that is in harmony with x and has support that is contained in the support of x (such a vector exists by the property just shown). By using the negative of \overline{z}^1 if necessary, we can assume that $x_j\overline{z}_j^1 > 0$ for at least one index j. Let $\overline{\gamma}$ be the largest value of γ such that

$$x_j - \gamma \overline{z}_j^1 \ge 0, \quad \forall j \text{ with } x_j > 0,$$

$$x_j - \gamma \overline{z}_j^1 \le 0, \quad \forall j \text{ with } x_j < 0.$$

The vector $x - z^1$, where

$$z^1 = \overline{\gamma} \, \overline{z}^1$$
,

is in harmony with x and has support that is strictly contained in the support of x. There are two cases: (1) $x=z^1$, in which case we are done, or (2) $x \neq z^1$, in which case we replace x by $x-z^1$ and we repeat the process. Eventually, after m steps where $m \leq n$ (since each step reduces the number of nonzero components by at least one), we will end up with the desired decomposition $x=z^1+\cdots+z^m$.

2.32 (Combinatorial Separation Theorem [Cam68], [Roc69])

Let S be a subspace of \Re^n . Consider a set B that is a Cartesian product of n nonempty intervals, and is such that $B\cap S^\perp=\emptyset$ (by an interval, we mean a convex set of scalars, which may be open, closed, or neither open nor closed.) Show that there exists an elementary vector z of S (cf. Exercise 2.31) such that

$$t'z < 0, \quad \forall \ t \in B,$$

i.e., a hyperplane that separates B and S^{\perp} , and does not contain any point of B. Note: There are two points here: (1) The set B need not be closed, as required for application of the Strict Separation Theorem (cf. Prop. 1.5.3), and (2) the hyperplane normal can be one of the elementary vectors of S (not just any vector of S). For application of this result in duality theory for network optimization and monotropic programming, see Rockafellar [Roc84] and Bertsekas [Ber98].

Solution: For simplicity, assume that B is the Cartesian product of bounded open intervals, so that B has the form

$$B = \{t \mid \underline{b}_j < t_j < \overline{b}_j, j = 1, \dots, n\},\$$

where \underline{b}_j and \overline{b}_j are some scalars. The proof is easily modified for the case where B has a different form.

Since $B \cap S^{\perp} = \emptyset$, there exists a hyperplane that separates B and S^{\perp} . The normal of this hyperplane is a nonero vector $d \in S$ such that

$$t'd < 0$$
, $\forall t \in B$.

Since B is open, this inequality implies that actually

$$t'd < 0$$
, $\forall t \in B$.

Equivalently, we have

$$\sum_{\{j|d_j>0\}} (\overline{b}_j - \epsilon)d_j + \sum_{\{j|d_j<0\}} (\underline{b}_j + \epsilon)d_j < 0, \tag{2.22}$$

for all $\epsilon > 0$ such that $\underline{b}_j + \epsilon < \overline{b}_j - \epsilon$. Let

$$d = z^1 + \dots + z^m,$$

be a decomposition of d, where z^1, \ldots, z^m are elementary vectors of S that are in harmony with x, and have supports that are contained in the support of d [cf. part (c) of the Exercise 2.31]. Then the condition (2.22) is equivalently written as

$$\begin{split} 0 &> \sum_{\{j|d_j>0\}} (\overline{b}_j - \epsilon) d_j + \sum_{\{j|d_j<0\}} (\underline{b}_j + \epsilon) d_j \\ &= \sum_{\{j|d_j>0\}} (\overline{b}_j - \epsilon) \left(\sum_{i=1}^m z_j^i\right) + \sum_{\{j|d_j<0\}} (\underline{b}_j + \epsilon) \left(\sum_{i=1}^m z_j^i\right) \\ &= \sum_{i=1}^m \left(\sum_{\{j|z_j^i>0\}} (\overline{b}_j - \epsilon) z_j^i + \sum_{\{j|z_j^i<0\}} (\underline{b}_j + \epsilon) z_j^i\right), \end{split}$$

where the last equality holds because the vectors z^i are in harmony with d and their supports are contained in the support of d. From the preceding relation, we see that for at least one elementary vector z^i , we must have

$$0 > \sum_{\{j \mid z_j^i > 0\}} (\overline{b}_j - \epsilon) z_j^i + \sum_{\{j \mid z_j^i < 0\}} (\underline{b}_j + \epsilon) z_j^i,$$

for all $\epsilon > 0$ that are sufficiently small and are such that $\underline{b}_j + \epsilon < \overline{b}_j - \epsilon$, or equivalently

$$0 > t'z^i, \quad \forall \ t \in B.$$

2.33 (Tucker's Complementarity Theorem)

(a) Let S be a subspace of \Re^n . Show that there exist disjoint index sets I and \overline{I} with $I \cup \overline{I} = \{1, \dots, n\}$, and vectors $x \in S$ and $y \in S^{\perp}$ such that

$$x_i > 0, \quad \forall i \in I, \qquad x_i = 0, \quad \forall i \in \overline{I},$$

 $y_i = 0, \quad \forall i \in I, \qquad y_i > 0, \quad \forall i \in \overline{I}.$

 $y_i = 0, \quad \forall \ i \in I, \qquad y_i > 0, \quad \forall \ i \in I.$

Furthermore, the index sets I and \overline{I} with this property are unique. In addition, we have

$$x_i = 0, \quad \forall i \in \overline{I}, \quad \forall x \in S \text{ with } x \ge 0,$$

$$y_i = 0, \quad \forall i \in I, \quad \forall y \in S^{\perp} \text{ with } y \geq 0.$$

Hint: Use a hyperplane separation argument based on Exercise 2.32.

(b) Let A be an $m \times n$ matrix and let b be a vector in \Re^n . Assume that the set $F = \{x \mid Ax = b, \ x \ge 0\}$ is nonempty. Apply part (a) to the subspace

$$S = \{(x, w) \mid Ax - bw = 0, \ x \in \Re^n, \ w \in \Re\},\$$

and show that there exist disjoint index sets I and \overline{I} with $I \cup \overline{I} = \{1, \ldots, n\}$, and vectors $x \in F$ and $z \in \Re^m$ such that b'z = 0 and

$$x_i > 0, \quad \forall \ i \in I, \qquad x_i = 0, \quad \forall \ i \in \overline{I},$$

$$y_i = 0, \quad \forall \ i \in I, \qquad y_i > 0, \quad \forall \ i \in \overline{I},$$

where y = A'z. Note: A special choice of A and b yields an important result, which relates optimal primal and dual solutions in linear programming: the Goldman-Tucker Complementarity Theorem [GoT56] (see the exercises of Chapter 5).

Solution: (a) Fix an index k and consider the following two assertions:

- (1) There exists a vector $x \in S$ with $x_i \ge 0$ for all i, and $x_k > 0$.
- (2) There exists a vector $y \in S^{\perp}$ with $y_i \geq 0$ for all i, and $y_k > 0$.

We claim that one and only one of the two assertions holds. Clearly, assertions (1) and (2) cannot hold simultaneously, since then we would have x'y>0, while $x\in S$ and $y\in S^{\perp}$. We will show that they cannot fail simultaneously. Indeed, if (1) does not hold, the Cartesian product $B=\prod_{i=1}^n B_i$ of the intervals

$$B_i = \begin{cases} (0, \infty) & \text{if } i = k, \\ [0, \infty) & \text{if } i \neq k, \end{cases}$$

does not intersect the subspace S, so by the result of Exercise 2.32, there exists a vector z of S^{\perp} such that x'z < 0 for all $x \in B$. For this to hold, we must have $z \in B^*$ or equivalently $z \leq 0$, while by choosing $x = (0, \ldots, 0, 1, 0, \ldots, 0) \in B$, with the 1 in the kth position, the inequality x'z < 0 yields $z_k < 0$. Thus assertion (2) holds with y = -z. Similarly, we show that if (2) does not hold, then (1) must hold.

Let now I be the set of indices k such that (1) holds, and for each $k \in I$, let x(k) be a vector in S such that $x(k) \geq 0$ and $x_k(k) > 0$ (note that we do not exclude the possibility that one of the sets I and \overline{I} is empty). Let \overline{I} be the set of indices such that (2) holds, and for each $k \in \overline{I}$, let y(k) be a vector in S^{\perp} such that $y(k) \geq 0$ and $y_k(k) > 0$. From what has already been shown, I and \overline{I} are disjoint, $I \cup \overline{I} = \{1, \ldots, n\}$, and the vectors

$$x = \sum_{k \in I} x(k), \qquad y = \sum_{k \in \overline{I}} y(k),$$

satisfy

$$x_i > 0, \quad \forall i \in I, \qquad x_i = 0, \quad \forall i \in \overline{I},$$

 $y_i = 0, \quad \forall i \in I, \qquad y_i > 0, \quad \forall i \in \overline{I}.$

The uniqueness of I and \overline{I} follows from their construction and the preceding arguments. In particular, if for some $k \in \overline{I}$, there existed a vector $x \in S$ with $x \geq 0$ and $x_k > 0$, then since for the vector y(k) of S^{\perp} we have $y(k) \geq 0$ and $y_k(k) > 0$, assertions (a) and (b) must hold simultaneously, which is a contradiction.

The last assertion follows from the fact that for each k, exactly one of the assertions (1) and (2) holds.

(b) Consider the subspace

$$S = \{(x, w) \mid Ax - bw = 0, x \in \Re^n, w \in \Re\}.$$

Its orthogonal complement is the range of the transpose of the matrix [A - b], so it has the form

$$S^{\perp} = \big\{ (A'z, -b'z) \mid z \in \Re^m \big\}.$$

By applying the result of part (a) to the subspace S, we obtain a partition of the index set $\{1, \ldots, n+1\}$ into two subsets. There are two possible cases:

- (1) The index n+1 belongs to the first subset.
- (2) The index n+1 belongs to the second subset.

In case (2), the two subsets are of the form I and $\overline{I} \cup \{n+1\}$ with $I \cup \overline{I} = \{1, \ldots, n\}$, and by the last assertion of part (a), we have w = 0 for all (x, w) such that $x \geq 0$, $w \geq 0$ and Ax - bw = 0. This, however, contradicts the

fact that the set $F = \{x \mid Ax = b, x \geq 0\}$ is nonempty. Therefore, case (1) holds, i.e., the index n+1 belongs to the first index subset. In particular, we have that there exist disjoint index sets I and \overline{I} with $I \cup \overline{I} = \{1, \dots, n\}$, and vectors (x, w) with Ax - bw = 0, and $z \in \mathbb{R}^m$ such that

$$w > 0,$$
 $b'z = 0,$
$$x_i > 0, \quad \forall i \in I, \qquad x_i = 0, \quad \forall i \in \overline{I},$$

$$y_i = 0, \quad \forall i \in I, \qquad y_i > 0, \quad \forall i \in \overline{I},$$

where y = A'z. By dividing (x, w) with w if needed, we may assume that w = 1 so that Ax - b = 0, and the result follows.

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Convex Optimization Theory

Chapter 3 Exercises and Solutions

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CHAPTER 3: EXERCISES AND SOLUTIONS

SECTION 3.1: Constrained Optimization

3.1 (Local Minima Along Lines)

- (a) Consider a vector x^* such that a given function $f: \mathbb{R}^n \to \mathbb{R}$ is convex over a sphere centered at x^* . Show that x^* is a local minimum of f if and only if it is a local minimum of f along every line passing through x^* [i.e., for all $d \in \mathbb{R}^n$, the function $g: \mathbb{R} \to \mathbb{R}$, defined by $g(\alpha) = f(x^* + \alpha d)$, has $\alpha^* = 0$ as its local minimum].
- (b) Consider the nonconvex function $f: \Re^2 \mapsto \Re$ given by

$$f(x_1, x_2) = (x_2 - px_1^2)(x_2 - qx_1^2),$$

where p and q are scalars with $0 , and <math>x^* = (0,0)$. Show that $f(y, my^2) < 0$ for $y \neq 0$ and m satisfying p < m < q, so x^* is not a local minimum of f even though it is a local minimum along every line passing through x^* .

Solution: (a) If x^* is a local minimum of f, evidently it is also a local minimum of f along any line passing through x^* . Conversely, let x^* be a local minimum of f along any line passing through x^* . Assume, to arrive at a contradiction, that x^* is not a local minimum of f and that we have $f(\overline{x}) < f(x^*)$ for some \overline{x} in the sphere centered at x^* within which f is assumed convex. Then, by convexity of f, for all $\alpha \in (0,1)$, we have

$$f(\alpha x^* + (1 - \alpha)\overline{x}) \le \alpha f(x^*) + (1 - \alpha)f(\overline{x}) < f(x^*),$$

so f decreases monotonically along the line segment connecting x^* and \overline{x} . This contradicts the hypothesis that x^* is a local minimum of f along any line passing through x^* .

(b) We first show that the function $g: \Re \mapsto \Re$ defined by $g(\alpha) = f(x^* + \alpha d)$ has a local minimum at $\alpha = 0$ for all $d \in \Re^2$. We have

$$g(\alpha) = f(x^* + \alpha d) = (\alpha d_2 - p\alpha^2 d_1^2)(\alpha d_2 - q\alpha^2 d_1^2) = \alpha^2 (d_2 - p\alpha d_1^2)(d_2 - q\alpha d_1^2).$$
 Also,

$$g'(\alpha) = 2\alpha(d_2 - p\alpha d_1^2)(d_2 - q\alpha d_1^2) + \alpha^2(-pd_1^2)(d_2 - q\alpha d_1^2) + \alpha^2(d_2 - p\alpha d_1^2)(-qd_1^2).$$

Thus g'(0) = 0. Furthermore,

$$g''(\alpha) = 2(d_2 - p\alpha d_1^2)(d_2 - q\alpha d_1^2) + 2\alpha(-pd_1^2)(d_2 - q\alpha d_1^2)$$

$$+ 2\alpha(d_2 - p\alpha d_1^2)(-qd_1^2) + 2\alpha(-pd_1^2)(d_2 - q\alpha d_1^2) + \alpha^2(-pd_1^2)(-qd_1^2)$$

$$+ 2\alpha(d_2 - p\alpha d_1^2)(-qd_1^2) + \alpha^2(-pd_1^2)(-qd_1^2).$$

Thus $g''(0) = 2d_2^2$, which is positive if $d_2 \neq 0$. If $d_2 = 0$, $g(\alpha) = pq\alpha^4d_1^4$, which is clearly minimized at $\alpha = 0$. Therefore, (0,0) is a local minimum of f along every line that passes through (0,0). We now show that if p < m < q, $f(y,my^2) < 0$ if $y \neq 0$ and that $f(y,my^2) = 0$ otherwise. Consider a point of the form (y,my^2) . We have $f(y,my^2) = y^4(m-p)(m-q)$. Clearly, $f(y,my^2) < 0$ if and only if p < m < q and $y \neq 0$. In any ϵ -neighborhood of (0,0), there exists a $y \neq 0$ such that for some $m \in (p,q)$, (y,my^2) also belongs to the neighborhood. Since f(0,0) = 0, we see that (0,0) is not a local minimum.

3.2 (Equality-Constrained Quadratic Programming)

(a) Consider the quadratic program

minimize
$$\frac{1}{2}||x||^2 + c'x$$

subject to $Ax = 0$, (3.1)

where $c \in \Re^n$ and A is an $m \times n$ matrix of rank m. Use the Projection Theorem to show that

$$x^* = -(I - A'(AA')^{-1}A)c (3.2)$$

is the unique solution.

(b) Consider the more general quadratic program

minimize
$$\frac{1}{2}(x-\overline{x})'Q(x-\overline{x}) + c'(x-\overline{x})$$

subject to $Ax = b$, (3.3)

where c and A are as before, Q is a symmetric positive definite matrix, $b \in \mathbb{R}^m$, and \overline{x} is a vector in \mathbb{R}^n , which is feasible, i.e., satisfies $A\overline{x} = b$. Use the transformation $y = Q^{1/2}(x - \overline{x})$ to write this problem in the form of part (a) and show that the optimal solution is

$$x^* = \overline{x} - Q^{-1}(c - A'\lambda),$$

where λ is given by

$$\lambda = (AQ^{-1}A')^{-1} AQ^{-1}c.$$

(c) Apply the result of part (b) to the program

minimize
$$\frac{1}{2}x'Qx + c'x$$

subject to $Ax = b$,

and show that the optimal solution is

$$x^* = -Q^{-1} (c - A'\lambda - A'(AQ^{-1}A')^{-1}b).$$

Solution: (a) By adding the constant term $\frac{1}{2}||c||^2$ to the cost function, we can equivalently write this problem as

minimize
$$\frac{1}{2} ||c + x||^2$$

subject to $Ax = 0$,

which is the problem of projecting the vector -c on the subspace $X = \{x \mid Ax = 0\}$. By the optimality condition for projection, a vector x^* such that $Ax^* = 0$ is the unique projection if and only if

$$(c+x^*)'x = 0, \quad \forall x \text{ with } Ax = 0.$$

It can be seen that the vector

$$x^* = -(I - A'(AA')^{-1}A)c$$
(3.4)

satisfies this condition and is thus the unique solution of the quadratic programming problem (3.1). (The matrix AA' is invertible because A has rank m.)

(b) By introducing the transformation $y = Q^{1/2}(x - \overline{x})$, we can write the problem as

minimize
$$\frac{1}{2}||y||^2 + \left(Q^{-1/2}c\right)'y$$

subject to $AQ^{-1/2}y = 0$.

Using Eq. (3.4) we see that the solution of this problem is

$$y^* = -\left(I - Q^{-1/2}A'\left(AQ^{-1}A'\right)^{-1}AQ^{-1/2}\right)Q^{-1/2}c$$

and by passing to the x-coordinate system through the inverse transformation $x^* - \overline{x} = Q^{-1/2}y^*$, we obtain the optimal solution

$$x^* = \overline{x} - Q^{-1}(c - A'\lambda), \tag{3.5}$$

where λ is given by

$$\lambda = (AQ^{-1}A')^{-1} AQ^{-1}c. \tag{3.6}$$

The quadratic program (3.3) contains as a special case the program

minimize
$$\frac{1}{2}x'Qx + c'x$$

subject to $Ax = b$. (3.7)

This special case is obtained when \overline{x} is given by

$$\overline{x} = Q^{-1}A'(AQ^{-1}A')^{-1}b.$$
 (3.8)

Indeed \overline{x} as given above satisfies $A\overline{x} = b$ as required, and for all x with Ax = b, we have

$$x'Q\overline{x} = x'A'(AQ^{-1}A')^{-1}b = b'(AQ^{-1}A')^{-1}b,$$

which implies that for all x with Ax = b,

$$\frac{1}{2}(x-\overline{x})'Q(x-\overline{x})+c'(x-\overline{x}) = \frac{1}{2}x'Qx+c'x+\left(\frac{1}{2}\overline{x}'Q\overline{x}-c'\overline{x}-b'(AQ^{-1}A')^{-1}b\right).$$

The last term in parentheses on the right-hand side above is constant, thus establishing that the programs (3.3) and (3.7) have the same optimal solution when \overline{x} is given by Eq. (3.8). By combining Eqs. (3.5) and (3.8), we obtain the optimal solution of program (3.7):

$$x^* = -Q^{-1}(c - A'\lambda - A'(AQ^{-1}A')^{-1}b),$$

where λ is given by Eq. (3.6).

3.3 (Approximate Minima of Convex Functions)

Let X be a closed convex subset of \Re^n , and let $f: \Re^n \mapsto (-\infty, \infty]$ be a closed convex function such that $X \cap \operatorname{dom}(f) \neq \emptyset$. Assume that f and X have no common nonzero direction of recession. Let X^* be the set of minima of f over X (which is nonempty and compact by Prop. 3.2.2), and let $f^* = \inf_{x \in X} f(x)$. Show that:

- (a) For every $\epsilon > 0$ there exists a $\delta > 0$ such that every vector $x \in X$ with $f(x) \leq f^* + \delta$ satisfies $\min_{x^* \in X^*} ||x x^*|| \leq \epsilon$.
- (b) If f is real-valued, for every $\delta > 0$ there exists an $\epsilon > 0$ such that every vector $x \in X$ with $\min_{x^* \in X^*} ||x x^*|| \le \epsilon$ satisfies $f(x) \le f^* + \delta$.
- (c) Every sequence $\{x_k\} \subset X$ satisfying $f(x_k) \to f^*$ is bounded and all its limit points belong to X^* .

Solution: (a) Let $\epsilon > 0$ be given. Assume, to arrive at a contradiction, that for any sequence $\{\delta_k\}$ with $\delta_k \downarrow 0$, there exists a sequence $\{x_k\} \in X$ such that for all k

$$f^* \le f(x_k) \le f^* + \delta_k, \qquad \min_{x^* \in X^*} ||x_k - x^*|| \ge \epsilon.$$

It follows that, for all k, x_k belongs to the set $\{x \in X \mid f(x) \leq f^* + \delta_0\}$, which is compact since f and X are closed and have no common nonzero

direction of recession. Therefore, the sequence $\{x_k\}$ has a limit point $\overline{x} \in X$, which using also the lower semicontinuity of f, satisfies

$$f(\overline{x}) \le \liminf_{k \to \infty} f(x_k) = f^*, \qquad ||\overline{x} - x^*|| \ge \epsilon, \quad \forall \ x^* \in X^*,$$

a contradiction.

(b) Let $\delta > 0$ be given. Assume, to arrive at a contradiction, that there exist sequences $\{x_k\} \subset X$, $\{x_k^*\} \subset X^*$, and $\{\epsilon_k\}$ with $\epsilon_k \downarrow 0$ such that

$$f(x_k) > f^* + \delta, \qquad ||x_k - x_k^*|| \le \epsilon_k, \qquad \forall \ k = 0, 1, \dots$$

(here x_k^* is the projection of x_k on X^*). Since X^* is compact, there is a subsequence $\{x_k^*\}_{\mathcal{K}}$ that converges to some $x^* \in X^*$. It follows that $\{x_k\}_{\mathcal{K}}$ also converges to x^* . Since f is real-valued, it is continuous, so we must have $f(x_k) \to f(x^*)$, a contradiction.

(c) Let \overline{x} be a limit point of the sequence $\{x_k\} \subset X$ satisfying $f(x_k) \to f^*$. By lower semicontinuity of f, we have that

$$f(\overline{x}) \le \liminf_{k \to \infty} f(x_k) = f^*.$$

Because $\{x_k\} \subset X$ and X is closed, we have $\overline{x} \in X$, which in view of the preceding relation implies that $f(\overline{x}) = f^*$, i.e., $\overline{x} \in X^*$.

SECTION 3.2: Existence of Optimal Solutions

3.4 (Minimization of Quasiconvex Functions)

We say that a function $f: \Re^n \mapsto (-\infty, \infty]$ is quasiconvex if all its level sets

$$V_{\gamma} = \left\{ x \mid f(x) \le \gamma \right\}$$

are convex. Let X be a convex subset of \Re^n , let f be a quasiconvex function such that $X \cap \text{dom}(f) \neq \emptyset$, and denote $f^* = \inf_{x \in X} f(x)$.

- (a) Assume that f is not constant on any line segment of X, i.e., we do not have f(x) = c for some scalar c and all x in the line segment connecting any two distinct points of X. Show that every local minimum of f over X is also global.
- (b) Assume that X is closed, and f is closed and proper. Let Γ be the set of all $\gamma > f^*$, and denote

$$R_f = \bigcap_{\gamma \in \Gamma} R_{\gamma}, \qquad L_f = \bigcap_{\gamma \in \Gamma} L_{\gamma},$$

where R_{γ} and L_{γ} are the recession cone and the lineality space of V_{γ} , respectively. Use the line of proof of Prop. 3.2.4 to show that

f attains a minimum over X if any one of the following conditions holds:

- (1) $R_X \cap R_f = L_X \cap L_f$.
- (2) $R_X \cap R_f \subset L_f$, and X is a polyhedral set.

Solution: (a) Let x^* be a local minimum of f over X and assume, to arrive at a contradiction, that there exists a vector $\overline{x} \in X$ such that $f(\overline{x}) < f(x^*)$. Then, \overline{x} and x^* belong to the set $X \cap V_{\gamma^*}$, where $\gamma^* = f(x^*)$. Since this set is convex, the line segment connecting x^* and \overline{x} belongs to the set, implying that

$$f(\alpha \overline{x} + (1 - \alpha)x^*) \le \gamma^* = f(x^*), \qquad \forall \ \alpha \in [0, 1]. \tag{3.9}$$

For each integer $k \geq 1$, there must exist an $\alpha_k \in (0, 1/k]$ such that

$$f(\alpha_k \overline{x} + (1 - \alpha_k)x^*) < f(x^*), \quad \text{for some } \alpha_k \in (0, 1/k];$$
 (3.10)

otherwise, in view of Eq. (3.9), we would have that f(x) is constant for x on the line segment connecting x^* and $(1/k)\overline{x} + (1 - (1/k))x^*$. Equation (3.10) contradicts the local optimality of x^* .

(b) We consider the level sets

$$V_{\gamma} = \left\{ x \mid f(x) \le \gamma \right\}$$

for $\gamma > f^*$. Let $\{\gamma_k\}$ be a scalar sequence such that $\gamma_k \downarrow f^*$. Using the fact that for two nonempty closed convex sets C and D such that $C \subset D$, we have $R_C \subset R_D$, it can be seen that

$$R_f = \bigcap_{\gamma \in \Gamma} R_{\gamma} = \bigcap_{k=1}^{\infty} R_{\gamma_k}.$$

Similarly, L_f can be written as

$$L_f = \cap_{\gamma \in \Gamma} L_\gamma = \cap_{k=1}^\infty L_{\gamma_k}.$$

Under each of the conditions (1) and (2), we will show that the set of minima of f over X, which is given by

$$X^* = \bigcap_{k=1}^{\infty} (X \cap V_{\gamma_k})$$

is nonempty. Let condition (1) hold. The sets $X \cap V_{\gamma_k}$ are nonempty, closed, convex, and nested. Furthermore, for each k, their recession cone is given by $R_X \cap R_{\gamma_k}$ and their lineality space is given by $L_X \cap L_{\gamma_k}$. We have that

$$\cap_{k=1}^{\infty} (R_X \cap R_{\gamma_k}) = R_X \cap R_f,$$

and

$$\cap_{k=1}^{\infty} (L_X \cap L_{\gamma_k}) = L_X \cap L_f,$$

while by assumption $R_X \cap R_f = L_X \cap L_f$. Then it follows by Prop. 3.2.4 that X^* is nonempty. Let condition (2) hold. The sets V_{γ_k} are nested and the intersection $X \cap V_{\gamma_k}$ is nonempty for all k. We also have by assumption that $R_X \cap R_f \subset L_f$ and X is a polyhedral set. By Prop. 3.2.4, it follows that X^* is nonempty.

3.5 (Properties of Quasiconvex Functions)

Show the following properties of quasiconvex functions (cf. Exercise 3.4):

(a) A function $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ is quasiconvex if and only if

$$f(\alpha x + (1 - \alpha)y) \le \max\{f(x), f(y)\}, \qquad \forall x, y \in \Re^n, \alpha \in [0, 1].$$
(3.11)

(b) A differentiable function $f: \Re^n \mapsto (-\infty, \infty)$ is quasiconvex if and only if

$$f(y) \le f(x) \quad \Rightarrow \quad \nabla f(x)'(y-x) \le 0, \qquad \forall \ x, y \in \Re^n.$$

(c) If $f: \Re^n \mapsto [-\infty, \infty]$ is a quasiconvex function, A is a matrix and b is a vector, the function

$$g(x) = f(Ax + b), \qquad x \in \Re^n,$$

is quasiconvex.

Solution: (a) Let f be quasiconvex. If either $f(x) = \infty$ or $f(y) = \infty$, then Eq. (3.11) holds, so assume that

$$\max\{f(x), f(y)\} < \infty.$$

Consider the level set

$$L = \left\{ z \mid f(z) \le \max\{f(x), f(y)\} \right\}.$$

Since x and y belong to L, and by quasiconvexity, L is convex, we have $\alpha x + (1 - \alpha)y \in L$ for all $\alpha \in [0, 1]$, so Eq. (3.11) holds. Conversely, let Eq. (3.11) hold, and consider two points x and y in a level set $\{x \mid f(x) \leq \gamma\}$. For any $\alpha \in [0, 1]$, we have

$$f(\alpha x + (1 - \alpha)y) \le \max\{f(x), f(y)\} \le \gamma,$$

so the level set $\{x \mid f(x) \leq \gamma\}$ is convex, implying that f is quasiconvex.

(b) Let f be quasiconvex, and let x and y be such that $f(y) \leq f(x)$. Assume, to arrive at a contradiction, that $\nabla f(x)'(y-x) > 0$. Then by Taylor's Theorem, we have $f(x+\epsilon(y-x)) > f(x)$ for all sufficiently small $\epsilon > 0$. This contradicts Eq. (3.11), which holds by quasiconvexity of f, as shown in part (a). Conversely, assume that

$$f(y) \le f(x) \quad \Rightarrow \quad \nabla f(x)'(y-x) \le 0, \qquad \forall \ x, y \in \Re^n.$$
 (3.12)

Note that in one dimension, this relation implies that if we have

$$f(z) > \max\{f(x), f(y)\}\$$

for some point z in the line segment connecting x and y, then the slope of f at z is zero. This shows that we cannot have $f(z) > \max\{f(x), f(y)\}$ for any z in the line segment connecting x and y, and by part (a) proves that f is quasiconvex. This is the idea underlying the following argument. Assume, to arrive at a contradiction, that f is not quasiconvex, so that, by part (a), there exist x and y, and $\overline{\alpha} \in (0,1)$ such that $f(y) \leq f(x)$ and

$$f(z) > \max\{f(x), f(y)\} = f(x),$$

where

$$z = \overline{\alpha}x + (1 - \overline{\alpha})y.$$

By the continuity of f, there exists a scalar $\beta \in (\overline{\alpha}, 1]$ such that

$$f(\alpha x + (1 - \alpha)y) > f(x), \quad \forall \alpha \in (\overline{\alpha}, \beta),$$

and

$$f(w) = f(x),$$

where

$$w = \beta x + (1 - \beta)y.$$

Using the Mean Value Theorem, it follows that there exists $\gamma \in (\overline{\alpha}, \beta)$ such that

$$\nabla f(\overline{w})'(z-w) = f(z) - f(w) > 0,$$

where

$$\overline{w} = \gamma x + (1 - \gamma)y$$
.

Since the vector z-w is colinear with the vector $y-\overline{w}$, follows that

$$\nabla f(\overline{w})'(y - \overline{w}) > 0,$$

which contradicts Eq. (3.12), since by construction we have

$$f(y) \le f(x) = f(w) \le f(\overline{w}).$$

(c) To show that g is quasiconvex, we must show that for all $\gamma \in \Re$, the set

$$V_{\gamma} = \{ x \mid f(Ax + b) \le \gamma \}$$

is convex. We have

$$V_{\gamma} = \{ x \mid Ax + b = y, y \in L_{\gamma} \},$$

where

$$L_{\gamma} = \{ y \mid f(y) \le \gamma \}.$$

Since f is convex, it follows that L_{γ} is convex, which implies that V_{γ} is convex.

3.6 (Directions Along Which a Function is Flat)

The purpose of the exercise is to provide refinements of results relating to set intersections and existence of optimal solutions. Let $f: \Re^n \mapsto (-\infty, \infty]$ be a closed proper convex function, and let F_f be the set of all directions g such that for every $g \in \text{dom}(f)$, the limit $\lim_{n\to\infty} f(x+ng)$ exists. We refer to F_f as the set of directions along which f is flat. Note that

$$L_f \subset F_f \subset R_f$$
,

where L_f and R_f are the constancy space and recession cone of f, respectively. Let X be a subset of \Re^n specified by linear inequality constraints, i.e.,

$$X = \{x \mid a_i' x \le b_j, j = 1, \dots, r\},\$$

where a_i are vectors in \Re^n and b_i are scalars. Assume that

$$R_X \cap F_f \subset L_f$$
,

where R_X is the recession cone of X.

(a) Let

$$C_k = \{x \mid f(x) \le w_k\},\$$

where $\{w_k\}$ is a monotonically decreasing and convergent scalar sequence, and assume that $X \cap C_k \neq \emptyset$ for all k. Show that

$$X \cap (\cap_{k=0}^{\infty} C_k) \neq \emptyset.$$

- (b) Show that if $\inf_{x \in X} f(x)$ is finite, the function f attains a minimum over the set X.
- (c) Show by example that f need not attain a minimum over X if we just assume that $X \cap \text{dom}(f) \neq \emptyset$.

Solution: (a) We use induction on the dimension of the set X. Suppose that the dimension of X is 0. Then X consists of a single point. This point belongs to $X \cap C_k$ for all k, and hence belongs to the intersection $X \cap \left(\bigcap_{k=0}^{\infty} C_k\right)$. Assume that, for some l < n, the intersection $\overline{X} \cap \left(\bigcap_{k=0}^{\infty} C_k\right)$ is nonempty for every set \overline{X} of dimension less than or equal to l that is specified by linear inequality constraints, and is such that $\overline{X} \cap C_k$ is nonempty for all k and $R_{\overline{X}} \cap F_f \subset L_f$. Let X be of the form

$$X = \{x \mid a'_j x \le b_j, j = 1, \dots, r\},\$$

and be such that $X \cap C_k$ is nonempty for all k, satisfy $R_X \cap F_f \subset L_f$, and have dimension l+1. We will show that the intersection $X \cap (\bigcap_{k=0}^{\infty} C_k)$

is nonempty. If $L_X \cap L_f = R_X \cap R_f$, then by Prop. 3.2.4 applied to the sets $X \cap C_k$, we have that $X \cap \left(\bigcap_{k=0}^{\infty} C_k \right)$ is nonempty, and we are done. We may thus assume that $L_X \cap L_f \neq R_X \cap R_f$. Let $\overline{y} \in R_X \cap R_f$ with $-\overline{y} \notin R_X \cap R_f$. If $\overline{y} \notin F_f$, then, since $\overline{y} \in R_X \cap R_f$, for all $x \in X \cap \text{dom}(f)$ we have $\lim_{\alpha \to \infty} f(x + \alpha y) = -\infty$ and $x + \alpha y \in X$ for all $\alpha \geq 0$. Therefore, $x + \alpha y \in X \cap \left(\bigcap_{k=0}^{\infty} C_k \right)$ for sufficiently large α , and we are done. We may thus assume that $\overline{y} \in F_f$, so that $\overline{y} \in R_X \cap F_f$ and therefore also $\overline{y} \subset L_f$, in view of the hypothesis $R_X \cap F_f \subset L_f$. Since $-\overline{y} \notin R_X \cap R_f$, it follows that $-\overline{y} \notin R_X$. Thus, we have

$$\overline{y} \in R_X, \quad -\overline{y} \notin R_X, \quad \overline{y} \in L_f.$$

Using Prop. 1.4.2(c), it is seen that the recession cone of X is

$$R_X = \{ y \mid a_i' y \le 0, \ j = 1, \dots, r \},\$$

so the fact $\overline{y} \in R_X$ implies that

$$a_i' \overline{y} \le 0, \quad \forall j = 1, \dots, r,$$

while the fact $-\overline{y} \notin R_X$ implies that the index set

$$J = \{ j \mid a_i' \overline{y} < 0 \}$$

is nonempty. Consider a sequence $\{x_k\}$ such that

$$x_k \in X \cap C_k, \quad \forall \ k.$$

We then have

$$a'_j x_k \le b_j, \quad \forall j = 1, \dots, r, \quad \forall k.$$

We may assume that

$$a_i'x_k < b_j, \quad \forall j \in J, \quad \forall k;$$

otherwise we can replace x_k with $x_k + \overline{y}$, which belongs to $X \cap C_k$ (since $\overline{y} \in R_X$ and $\overline{y} \in L_f$). Suppose that for each k, we start at x_k and move along $-\overline{y}$ as far as possible without leaving the set X, up to the point where we encounter the vector

$$\overline{x}_k = x_k - \beta_k \overline{y},$$

where β_k is the positive scalar given by

$$\beta_k = \min_{j \in J} \frac{a_j' x_k - b_j}{a_j' \overline{y}}.$$

Since $a'_j \overline{y} = 0$ for all $j \notin J$, we have $a'_j \overline{x}_k = a'_j x_k$ for all $j \notin J$, so the number of linear inequalities of X that are satisfied by \overline{x}_k as equalities is strictly

larger than the number of those satisfied by x_k . Thus, there exists $j_0 \in J$ such that $a'_{j_0}\overline{x}_k = b_{j_0}$ for all k in an infinite index set $\mathcal{K} \subset \{0, 1, \ldots\}$. By reordering the linear inequalities if necessary, we can assume that $j_0 = 1$, i.e.,

$$a_1'\overline{x}_k = b_1, \quad a_1'x_k < b_1, \quad \forall \ k \in \mathcal{K}.$$

To apply the induction hypothesis, consider the set

$$\overline{X} = \{x \mid a_1' x = b_1, \ a_j' x \le b_j, \ j = 2, \dots, r\},\$$

and note that $\{\overline{x}_k\}_{\mathcal{K}} \subset \overline{X}$. Since $\overline{x}_k = x_k - \beta_k \overline{y}$ with $x_k \in C_k$ and $\overline{y} \in L_f$, we have $\overline{x}_k \in C_k$ for all k, implying that $\overline{x}_k \in \overline{X} \cap C_k$ for all $k \in \mathcal{K}$. Thus, $\overline{X} \cap C_k \neq \emptyset$ for all k. Because the sets C_k are nested, so are the sets $\overline{X} \cap C_k$. Furthermore, the recession cone of \overline{X} is

$$R_{\overline{X}} = \{ y \mid a_1' y = 0, \ a_j' y \le 0, \ j = 2, \dots, r \},$$

which is contained in R_X , so that

$$R_{\overline{X}} \cap F_f \subset R_X \cap F_f \subset L_f$$
.

Finally, to show that the dimension of \overline{X} is smaller than the dimension of X, note that the set $\{x \mid a_1'x = b_1\}$ contains \overline{X} , so that a_1 is orthogonal to the subspace $S_{\overline{X}}$ that is parallel to $\operatorname{aff}(\overline{X})$. Since $a_1'\overline{y} < 0$, it follows that $\overline{y} \notin S_{\overline{X}}$. On the other hand, \overline{y} belongs to S_X , the subspace that is parallel to $\operatorname{aff}(X)$, since for all k, we have $x_k \in X$ and $x_k - \beta_k \overline{y} \in X$. Based on the preceding, we can use the induction hypothesis to assert that the intersection $\overline{X} \cap \left(\cap_{k=0}^{\infty} C_k \right)$ is nonempty. Since $\overline{X} \subset X$, it follows that $X \cap \left(\cap_{k=0}^{\infty} C_k \right)$ is nonempty.

(b) Denote

$$f^* = \inf_{x \in X} f(x),$$

and assume without loss of generality that $f^* = 0$ [otherwise, we replace f(x) by $f(x) - f^*$]. We choose a scalar sequence $\{w_k\}$ such that $w_k \downarrow f^*$, and we consider the (nonempty) level sets

$$C_k = \left\{ x \in \Re^n \mid f(x) \le w_k \right\}.$$

The set $X \cap C_k$ is nonempty for all k. Furthermore, by assumption, $R_X \cap F_f \subset L_f$ and X is specified by linear inequality constraints. By part (a), it follows that $X \cap (\bigcap_{k=0}^{\infty} C_k)$, the set of minimizers of f over X, is nonempty.

(c) Let $X = \Re$ and f(x) = x. Then

$$F_f = L_f = \{ y \mid y = 0 \},$$

so the condition $R_X \cap F_f \subset L_f$ is satisfied. However, we have $\inf_{x \in X} f(x) = -\infty$ and f does not attain a minimum over X.

3.7 (Bidirectionally Flat Functions - Intersections of Sets Defined by Quadratic Functions)

The purpose of the exercise is to provide refinements of various results relating to set intersections, closedness under linear transformations, existence of optimal solutions, and closedness under partial minimization. Important special cases arise when the sets involved are defined by convex quadratic functions.

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a closed proper convex function, and let F_f be the set of directions along which f is flat (cf. Exercise 3.6). We say that f is bidirectionally flat if $L_f = F_f$ (i.e., if it is flat in some direction it must be flat, and hence constant, in the opposite direction). Note that every convex quadratic function is bidirectionally flat. More generally, a function of the form

$$f(x) = h(Ax) + c'x,$$

where A is an $m \times n$ matrix and $h : \Re^m \mapsto (-\infty, \infty]$ is a coercive closed proper convex function, is bidirectionally flat. In this case, we have

$$L_f = F_f = \{ y \mid Ay = 0, c'y = 0 \}.$$

Let $g_j: \Re^n \mapsto (-\infty, \infty], \ j=0,1,\ldots,r$, be closed proper convex functions that are bidirectionally flat.

(a) Assume that each vector x such that $g_0(x) \leq 0$ belongs to $\bigcap_{j=1}^r \text{dom}(g_j)$, and that for some scalar sequence $\{w_k\}$ with $w_k \downarrow 0$, the set

$$C_k = \{x \mid g_0(x) \le w_k, g_j(x) \le 0, j = 1, \dots, r\}$$

is nonempty for each k. Show that the intersection $\bigcap_{k=0}^{\infty} C_k$ is nonempty.

(b) Assume that each g_j , j = 1, ..., r, is real-valued and the set

$$C = \{x \mid q_i(x) < 0, j = 1, \dots, r\}$$

is nonempty. Show that for any $m \times n$ matrix A, the set AC is closed.

(c) Show that a closed proper convex function $f: \Re^n \mapsto (-\infty, \infty]$ that is bidirectionally flat attains a minimum over the set C of part (b), provided that $\inf_{x \in C} f(x)$ is finite.

Solution: (a) As a first step, we will show that either $\bigcap_{k=1}^{\infty} C_k \neq \emptyset$ or else

there exists
$$\overline{j} \in \{1,\dots,r\}$$
 and $y \in \cap_{j=0}^r R_{g_j}$ with $y \notin F_{g_{\overline{j}}}$.

Let \overline{x} be a vector in C_0 , and for each $k \geq 1$, let x_k be the projection of \overline{x} on C_k . If $\{x_k\}$ is bounded, then since the g_j are closed, any limit point \tilde{x} of $\{x_k\}$ satisfies

$$g_j(\tilde{x}) \le \liminf_{k \to \infty} g_j(x_k) \le 0,$$

so $\tilde{x} \in \bigcap_{k=1}^{\infty} C_k$, and $\bigcap_{k=1}^{\infty} C_k \neq \emptyset$. If $\{x_k\}$ is unbounded, let y be a limit point of the sequence $\{(x_k - \overline{x})/\|x_k - \overline{x}\| \mid x_k \neq \overline{x}\}$, and without loss of generality, assume that

 $\frac{x_k - \overline{x}}{\|x_k - \overline{x}\|} \to y.$

We claim that

$$y\in \cap_{j=0}^r R_{g_j}.$$

Indeed, if for some j, we have $y \notin R_{g_j}$, then there exists $\alpha > 0$ such that $g_j(\overline{x} + \alpha y) > w_0$. Let

 $z_k = \overline{x} + \alpha \frac{x_k - \overline{x}}{\|x_k - \overline{x}\|},$

and note that for sufficiently large k, z_k lies in the line segment connecting \overline{x} and x_k , so that $g_1(z_k) \leq w_0$. On the other hand, we have $z_k \to \overline{x} + \alpha y$, so using the closedness of g_j , we must have

$$g_j(\overline{x} + \alpha y) \le \liminf_{k \to \infty} g_1(z_k) \le w_0,$$

which contradicts the choice of α to satisfy $g_j(\overline{x}+\alpha y)>w_0$. If $y\in \cap_{j=0}^r F_{g_j}$, since all the g_j are bidirectionally flat, we have $y\in \cap_{j=0}^r L_{g_j}$. If the vectors \overline{x} and $x_k, k\geq 1$, all lie in the same line [which must be the line $\{\overline{x}+\alpha y\mid\alpha\in\Re\}$], we would have $g_j(\overline{x})=g_j(x_k)$ for all k and j. Then it follows that \overline{x} and x_k all belong to $\cap_{k=1}^\infty C_k$. Otherwise, there must be some x_k , with k large enough, and such that, by the Projection Theorem, the vector $x_k-\alpha y$ makes an angle greater than $\pi/2$ with $x_k-\overline{x}$. Since the g_j are constant on the line $\{x_k-\alpha y\mid\alpha\in\Re\}$, all vectors on the line belong to C_k , which contradicts the fact that x_k is the projection of \overline{x} on C_k . Finally, if $y\in R_{g_0}$ but $y\notin F_{g_0}$, we have $g_0(x+\alpha y)\to -\infty$ as $\alpha\to\infty$, so that $\cap_{k=1}^\infty C_k\neq\emptyset$. This completes the proof that

$$\bigcap_{k=1}^{\infty} C_k = \emptyset \implies \text{there exists } \overline{j} \in \{1, \dots, r\} \text{ and } y \in \bigcap_{j=0}^r R_j \text{ with } y \notin F_{g_{\overline{j}}}.$$
(3.13)

We now use induction on r. For r=0, the preceding proof shows that $\bigcap_{k=1}^{\infty} C_k \neq \emptyset$. Assume that $\bigcap_{k=1}^{\infty} C_k \neq \emptyset$ for all cases where $r < \overline{r}$. We will show that $\bigcap_{k=1}^{\infty} C_k \neq \emptyset$ for $r=\overline{r}$. Assume the contrary. Then, by Eq. (3.13), there exists $\overline{j} \in \{1,\ldots,r\}$ and $y \in \bigcap_{j=0}^r R_j$ with $y \notin F_{g_{\overline{j}}}$. Let us consider the sets

$$\overline{C}_k = \{x \mid g_0(x) \le w_k, g_j(x) \le 0, j = 1, \dots, r, j \ne \overline{j}\}.$$

Since these sets are nonempty, by the induction hypothesis, $\bigcap_{k=1}^{\infty} \overline{C}_k \neq \emptyset$. For any $\tilde{x} \in \bigcap_{k=1}^{\infty} \overline{C}_k$, the vector $\tilde{x} + \alpha y$ belongs to $\bigcap_{k=1}^{\infty} \overline{C}_k$ for all $\alpha > 0$, since $y \in \bigcap_{j=0}^r R_j$. Since $g_0(\tilde{x}) \leq 0$, we have $\tilde{x} \in \text{dom}(g_{\overline{j}})$, by the hypothesis regarding the domains of the g_j . Since $y \in \bigcap_{j=0}^r R_j$ with $y \notin F_{g_{\overline{j}}}$, it follows

that $g_{\overline{j}}(\tilde{x} + \alpha y) \to -\infty$ as $\alpha \to \infty$. Hence, for sufficiently large α , we have $g_{\overline{i}}(\tilde{x} + \alpha y) \leq 0$, so $\tilde{x} + \alpha y$ belongs to $\bigcap_{k=1}^{\infty} C_k$.

Note: To see that the assumption

$$\{x \mid g_0(x) \le 0\} \subset \cap_{i=1}^r \operatorname{dom}(g_i)$$

is essential for the result to hold, consider an example in \Re^2 . Let

$$g_0(x_1, x_2) = x_1,$$
 $g_1(x_1, x_2) = \phi(x_1) - x_2,$

where the function $\phi : \Re \mapsto (-\infty, \infty]$ is convex, closed, and coercive with $dom(\phi) = (0, 1)$ [for example,

$$\phi(t) = -\ln t - \ln(1-t)$$

for 0 < t < 1]. Then it can be verified that $C_k \neq \emptyset$ for every k and sequence $\{w_k\} \subset (0,1)$ with $w_k \downarrow 0$ [take $x_1 \downarrow 0$ and $x_2 \geq \phi(x_1)$]. On the other hand, we have

$$\cap_{k=0}^{\infty} C_k = \emptyset.$$

The difficulty here is that the set $\{x \mid g_0(x) \leq 0\}$, which is equal to

$$\{x \mid x_1 \le 0, x_2 \in \Re\},\$$

is not contained in $dom(g_1)$, which is equal to

$$\{x \mid 0 < x_1 < 1, x_2 \in \Re\}$$

(in fact the two sets are disjoint).

(b) Let $\{y_k\}$ be a sequence in AC converging to some $\overline{y} \in \mathbb{R}^n$. We will show that $\overline{y} \in AC$. We let

$$g_0(x) = ||Ax - \overline{y}||^2, \qquad w_k = ||y_k - \overline{y}||^2,$$

and

$$C_k = \{x \mid g_0(x) \le w_k, g_j(x) \le 0, j = 1, \dots, r\}.$$

The functions involved in the definition of C_k are bidirectionally flat, and each C_k is nonempty by construction. By applying part (a), we see that the intersection $\bigcap_{k=0}^{\infty} C_k$ is nonempty. For any x in this intersection, we have $Ax = \overline{y}$ (since $y_k \to \overline{y}$), showing that $\overline{y} \in AC$.

(c) Denote

$$f^* = \inf_{x \in C} f(x),$$

and assume without loss of generality that $f^* = 0$ [otherwise, we replace f(x) by $f(x) - f^*$]. We choose a scalar sequence $\{w_k\}$ such that $w_k \downarrow f^*$, and we consider the (nonempty) sets

$$C_k = \{x \in \Re^n \mid f(x) \le w_k, g_j(x) \le 0, j = 1, \dots, r\}.$$

By part (a), it follows that $\bigcap_{k=0}^{\infty} C_k$, the set of minimizers of f over C, is nonempty.

SECTION 3.3: Partial Minimization of Convex Functions

3.8 (Asymptotic Slopes of Functions of Two Vectors)

Let $F: \Re^{n+m} \mapsto (-\infty, \infty]$ be a closed proper convex function of two vectors $x \in \Re^n$ and $z \in \Re^m$, and let

$$X = \left\{ x \mid \inf_{z \in \Re^m} F(x, z) < \infty \right\}.$$

Assume that the function $F(x,\cdot)$ is closed for each $x \in X$. Show that:

- (a) If for some $\overline{x} \in X$, the minimum of $F(\overline{x}, \cdot)$ over \Re^m is attained at a nonempty and compact set, the same is true for all $x \in X$.
- (b) If the functions $F(x,\cdot)$ are differentiable for all $x \in X$, they have the same asymptotic slopes along all directions, i.e., for each $d \in \mathbb{R}^m$, the value of $\lim_{\alpha \to \infty} \nabla_z F(x,z+\alpha d)'d$ is the same for all $x \in X$ and $z \in \mathbb{R}^m$.

Solution: By Prop. 1.4.5(a), the recession cone of F has the form

$$R_F = \{ (d_x, d_z) \mid (d_x, d_z, 0) \in R_{\text{epi}(F)} \}.$$

The (common) recession cone of the nonempty level sets of $F(x,\cdot)$, $x\in X$, has the form

$$\{d_z \mid (0, d_z) \in R_F\},\$$

for all $x \in X$, where R_F is the recession cone of F. Furthermore, the recession function of $F(x,\cdot)$ is the same for all $x \in X$.

- (a) By the compactness hypothesis, the recession cone of $F(\overline{x}, \cdot)$ consists of just the origin, so the same is true for the recession cones of all $F(x, \cdot)$, $x \in X$. Thus the nonempty level sets of $F(x, \cdot)$, $x \in X$, are all compact.
- (b) This is a consequence of the fact that the recession function of $F(x,\cdot)$ is the same for all $x \in X$, and the comments following Prop. 1.4.7.

3.9 (Partial Minimization)

(a) Let $f: \Re^n \mapsto [-\infty, \infty]$ be a function and consider the subset of \Re^{n+1} given by

$$E_f = \{(x, w) \mid f(x) < w\}.$$

Show that E_f is related to the epigraph of f as follows:

$$E_f \subset \operatorname{epi}(f) \subset \operatorname{cl}(E_f).$$

Show also that f is convex if and only if E_f is convex.

(b) Let $F: \Re^{m+n} \mapsto [-\infty, \infty]$ be a function and let

$$f(x) = \inf_{z \in \mathbb{R}^m} F(x, z), \qquad x \in \mathbb{R}^n.$$

Show that E_f is the projection of the set $\{(x, z, w) \mid F(x, z) < w\}$ on the space of (x, w).

(c) Use parts (a) and (b) to show that convexity of F implies convexity of f [this shows Prop. 3.3.1(a) by a somewhat different argument, which does not rely on the assumption $F(x,z) > -\infty$ for all (x,z)].

Solution: (a) We clearly have $E_f \subset \operatorname{epi}(f)$. Let $(x, w) \in \operatorname{epi}(f)$, so that $f(x) \leq w$. Let $\{w_k\}$ be a decreasing sequence such that $w_k \to w$, so that $f(x) < w_k$ and $(x, w_k) \in E_f$ for all k. Since $(x, w_k) \to (x, w)$, it follows that $(x, w) \in \operatorname{cl}(E_f)$. Hence, $\operatorname{epi}(f) \subset \operatorname{cl}(E_f)$. Let f be convex, and let $(x, w), (y, v) \in E_f$. Then for any $\alpha \in [0, 1]$, we have

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y) < \alpha w + (1 - \alpha)v.$$

It follows that $(\alpha x + (1 - \alpha)y, \alpha w + (1 - \alpha)v) \in E_f$, so that E_f is convex. Let E_f be convex, and for any $x, y \in \text{dom}(f)$, let $\{(x, w_k)\}$, $\{(y, v_k)\}$ be sequences in E_f such that $w_k \downarrow f(x)$ and $v_k \downarrow f(y)$, respectively [we allow the possibility that f(x) and/or f(y) are equal to $-\infty$]. Since E_f is convex, we have

$$f(\alpha x + (1 - \alpha)y) < \alpha w_k + (1 - \alpha)v_k, \qquad \forall \ \alpha \in [0, 1], \ k = 0, 1, \dots$$

Taking the limit as $k \to \infty$, we obtain

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y), \quad \forall \alpha \in [0, 1],$$

so f is convex.

(b) We have f(x) < w if and only if there exists $z \in \Re^m$ such that F(x,z) < w:

$$E_f = \{(x, w) \mid F(x, z) < w \text{ for some } z \in \Re^m \}.$$

Thus E_f is the projection of the set $\{(x, z, w) \mid F(x, z) < w\}$ on the space of (x, w).

(c) If F is convex, by part (a) the set $\{(x, z, w) \mid F(x, z) < w\}$ is convex, and by part (b), E_f is the image of this set under a linear transformation. Therefore, E_f is convex, so by part (a), f is convex.

3.10 (Closures of Partially Minimized Functions)

Consider a function $F: \mathbb{R}^{n+m} \mapsto (-\infty, \infty]$ and the function $f: \mathbb{R}^n \mapsto [-\infty, \infty]$ defined by

$$f(x) = \inf_{z \in \Re^m} F(x, z).$$

(a) Show that

$$P(\operatorname{cl}(\operatorname{epi}(F))) \subset \operatorname{cl}(\operatorname{epi}(f)),$$

where $P(\cdot)$ denotes projection on the space of (x, w).

(b) Let \overline{f} be defined by

$$\overline{f}(x) = \inf_{z \in \Re^m} (\operatorname{cl} F)(x, z),$$

where $\operatorname{cl} F$ is the closure of F. Show that the closures of f and \overline{f} coincide.

Solution: (a) We first note that from Prop. 3.3.1, we have

$$P(\operatorname{epi}(F)) \subset \operatorname{epi}(f) \subset \operatorname{cl}(P(\operatorname{epi}(F))),$$
 (3.14)

where $P(\cdot)$ denotes projection on the space of (x, w). To show the equation

$$P(\operatorname{cl}(\operatorname{epi}(F))) \subset \operatorname{cl}(\operatorname{epi}(f)),$$
 (3.15)

let $(\overline{x}, \overline{w})$ belong to $P(\operatorname{cl}(\operatorname{epi}(F)))$. Then there exists \overline{z} such that $(\overline{x}, \overline{z}, \overline{w}) \in \operatorname{cl}(\operatorname{epi}(F))$, and hence there is a sequence $(x_k, z_k, w_k) \in \operatorname{epi}(F)$ such that $x_k \to \overline{x}, z_k \to \overline{z}$, and $w_k \to \overline{w}$. Thus we have $f(x_k) \leq F(x_k, z_k) \leq w_k$, implying that $(x_k, w_k) \in \operatorname{epi}(f)$ for all k. It follows that $(\overline{x}, \overline{w}) \in \operatorname{cl}(\operatorname{epi}(f))$.

(b) By taking closure in Eq. (3.14), we see that

$$\operatorname{cl}(\operatorname{epi}(f)) = \operatorname{cl}(P(\operatorname{epi}(F))).$$
 (3.16)

Denoting $\overline{F} = \operatorname{cl} F$ and replacing F with \overline{F} , we also have

$$\operatorname{cl}(\operatorname{epi}(\overline{f})) = \operatorname{cl}(P(\operatorname{epi}(\overline{F}))).$$
 (3.17)

On the other hand, by taking closure in Eq. (3.15), we have

$$\operatorname{cl}\left(P\left(\operatorname{epi}(\overline{F})\right)\right) \subset \operatorname{cl}\left(P\left(\operatorname{epi}(F)\right)\right),$$

which, in view of $\operatorname{epi}(\overline{F}) \supset \operatorname{epi}(F)$, implies that

$$\operatorname{cl}\left(P\left(\operatorname{epi}(\overline{F})\right)\right) = \operatorname{cl}\left(P\left(\operatorname{epi}(F)\right)\right).$$
 (3.18)

By combining Eqs. (3.16)-(3.18), we see that

$$\operatorname{cl}(\operatorname{epi}(f)) = \operatorname{cl}(\operatorname{epi}(\overline{f})).$$

3.11 (Counterexample for Partially Minimized Functions)

Consider the function of $x \in \Re$ and $z \in \Re$ given by

$$F(x,z) = \begin{cases} e^{-\sqrt{xz}} & \text{if } x \ge 0, \ z \ge 0, \\ \infty & \text{otherwise.} \end{cases}$$

Verify that F convex and closed, but that the function $f(x) = \inf_{z \in \Re} F(x, z)$ is convex but not closed.

Solution: To prove that F is convex, we note that it is the composition of the monotonically increasing exponential function, and the function $-\sqrt{xz}$ which can be shown to convex over $\{(x,z)\mid x\geq 0,\ z\geq 0\}$ (one way to do this is to use Prop. 1.1.7). It is straightforward to verify that f is the nonclosed function

$$f(x) = \inf_{z \in \Re} F(x, z) = \begin{cases} 0 & \text{if } x > 0, \\ 1 & \text{if } x = 0, \\ \infty & \text{if } x < 0. \end{cases}$$

In particular, if x > 0 and $\epsilon > 0$, we have $0 < e^{-\sqrt{xz}} < \epsilon$, provided z is sufficiently large.

3.12 (Image Operation)

Let $h: \Re^m \mapsto (-\infty, \infty]$ be a function and A be an $m \times n$ matrix. Consider the function $A \circ h: \Re^n \mapsto [-\infty, \infty]$ defined by

$$(A \circ h)(x) = \inf_{Az=x} h(z), \qquad x \in \Re^n.$$
 (3.19)

We refer to $A \circ h$ as the *image of h under A*. (The terminology comes from the special case where h is the indicator function of a set; then $A \circ h$ is the indicator function of the image of the set under A.)

- (a) Show that if h is convex, then $A \circ h$ is convex.
- (b) Assume that h is closed proper convex and that $R(A) \cap \text{dom}(h) \neq \emptyset$, where R(A) is the range of A. Assume further that every direction of recession of h that belongs to the nullspace of A is a direction along which h is constant. Then $A \circ h$ is closed proper convex, and for every $x \in \text{dom}(f)$, the infimum in the definition of $(A \circ h)(x)$ is attained. Hint: Use Prop. 3.3.4.

Solution: Consider the function $F: \Re^{n+m} \mapsto (-\infty, \infty]$ given by

$$F(x,z) = \begin{cases} h(z) & \text{if } Az = x, \\ \infty & \text{otherwise.} \end{cases}$$

We have

$$(A\circ h)(x)=\inf_{z\in\Re^m}F(x,z),\qquad x\in\Re^n,$$

so $A \circ h$ is obtained from by partial minimization, and its properties follow by applying the theory of Section 3.3.

- (a) By Prop. 3.3.1(a), $A \circ h$ is convex.
- (b) We will use Prop. 3.3.4. Let \overline{x} be a point in $R(A) \cap \text{dom}(h)$. Choose $\overline{\gamma}$ so that the set $\{z \mid h(z) \leq \overline{\gamma}, Az = \overline{x}\}$ is nonempty, and note that

$$\{z \mid h(z) \le \overline{\gamma}, Az = \overline{x}\} = \{z \mid F(\overline{x}, z) \le \overline{\gamma}\}.$$

From this we see that our assumption that every direction of recession of h that belongs to the nullspace of A is a direction along which h is constant is equivalent to the recession cone of $\{z \mid F(\overline{x},z) \leq \overline{\gamma}\}$ being equal to its lineality space. Thus Prop. 3.3.4 applies, and by the conclusion of the proposition, $A \circ h$ is closed proper convex, and for every $x \in \text{dom}(A \circ h)$, the infimum in the definition of $(A \circ h)(x)$ is attained.

3.13 (Infimal Convolution Operation)

Consider the function $f: \Re^n \mapsto [-\infty, \infty]$ defined by

$$f(x) = \inf_{x_1 + \dots + x_m = x} \{ f_1(x_1) + \dots + f_m(x_m) \}, \quad x \in \Re^n,$$

where for $i=1,\ldots,m,\ f_i:\Re^n\mapsto (-\infty,\infty]$ is a closed proper convex function. The function f is called the *infimal convolution of* f_1,\ldots,f_m . The terminology comes from the resemblance with the classical formula of integral convolution in the case of two functions when f(x) takes the form

$$\inf_{z\in\Re^n}\big\{f_1(x-z)+f_2(z)\big\}.$$

Assume that $\bigcap_{i=1}^m \text{dom}(f_i) \neq \emptyset$, and furthermore that every $d = (d_1, \dots, d_m)$ that is a direction of recession of $f_1 + \dots + f_m$ and satisfies

$$d_1 + \dots + d_m = 0,$$

is a direction along which $f_1 + \cdots + f_m$ is constant. Show that f is closed proper convex, and that the infimum in the definition of f(x) is attained for every $x \in \text{dom}(f)$. *Hint*: Show that infimal convolution is a special case of the image operation of Exercise 3.12, and apply the result of that exercise.

Solution: We can write f as the result of an image operation, $(A \circ h)$ [cf. Eq. (3.19)], where $h: \Re^{mn} \mapsto (-\infty, \infty]$ is the function given by

$$h(x_1, \ldots, x_m) = f_1(x_1) + \cdots + f_m(x_m),$$

and A is the $n \times (mn)$ matrix defined by

$$A(x_1,\ldots,x_m)=x_1+\cdots+x_m.$$

Our assumption on directions of recession is the same as the assumption of Exercise 3.12(b), specialized to the infimal convolution context.

3.14 (Epigraph Relations for Image and Infimal Convolution)

Consider the image operation $A \circ h$ of Exercise 3.12, and the relation

$$(A \circ h)(x) = \inf_{z \in \Re^m} F(x, z), \qquad x \in \Re^n,$$

where $F: \Re^{n+m} \mapsto (-\infty, \infty]$ is given by

$$F(x,z) = \begin{cases} h(z) & \text{if } Az = x, \\ \infty & \text{otherwise.} \end{cases}$$

(a) Show that the epigraph relation

$$P(\operatorname{epi}(F)) \subset \operatorname{epi}(f) \subset \operatorname{cl}(P(\operatorname{epi}(F)))$$

where $P(\cdot)$ denotes projection on the space of (x,w) [cf. Prop. 3.3.1(b)] takes the form

$$\overline{A}\operatorname{epi}(h) \subset \operatorname{epi}(f) \subset \operatorname{cl}(\overline{A}\operatorname{epi}(h)),$$

where \overline{A} is the linear transformation that maps a vector $(x, w) \in \mathbb{R}^{n+1}$ to (Ax, w).

(b) Show that in the special case of the infimal convolution operation (cf. Exercise 3.13) the epigraph relation takes the form

$$\operatorname{epi}(f_1) + \cdots + \operatorname{epi}(f_m) \subset \operatorname{epi}(f) \subset \operatorname{cl}(\operatorname{epi}(f_1) + \cdots + \operatorname{epi}(f_m)).$$

Solution: In the case of the image operation, we have

$$epi(F) = \{(x, z, w) \mid Az = x, h(z) \le w\},\$$

so that

$$P(\operatorname{epi}(F)) = \{(x, w) \mid \text{for some } z \text{ with } Az = x, h(z) \le w\}$$
$$= \{(Az, w) \mid h(z) \le w\}$$
$$= \overline{A} \operatorname{epi}(h).$$

We thus obtain

$$\overline{A}\operatorname{epi}(h) \subset \operatorname{epi}(f) \subset \operatorname{cl}(\overline{A}\operatorname{epi}(h)).$$

(b) Similar to part (a).

3.15 (Conjugates of Linear Composition and Image)

Consider the composition of a closed proper convex function $f: \Re^m \mapsto (-\infty, \infty]$ with a linear transformation A, an $m \times n$ matrix, i.e., the function $f \circ A: \Re^n \mapsto [-\infty, \infty]$ given by

$$(f \circ A)(x) = f(Ax).$$

Consider also the function $A' \circ f^*$ given by

$$(A' \circ f^*)(x) = \inf_{A'z = x} f^*(z), \qquad x \in \Re^n,$$

which was called the image function of f^* under A' in Example 3.13. Show that the conjugate of $A' \circ f^*$ is $f \circ A$, and that the conjugate of $f \circ A$ is the closure of $A' \circ f^*$, provided $f \circ A$ is proper [which is true if and only if the range of A contains a point in dom(f)]. Give an example where the conjugate of $f \circ A$ is different from $A' \circ f^*$.

Solution: We have for all $z \in \Re^n$,

$$\begin{split} f(Az) &= \sup_{y} \left\{ z'A'y - f^{\star}(y) \right\} \\ &= \sup_{x} \sup_{A'y = x} \left\{ z'x - f^{\star}(y) \right\} \\ &= \sup_{x} \left\{ z'x - \inf_{A'y = x} f^{\star}(y) \right\} \\ &= \sup_{x} \left\{ z'x - (A' \circ f^{\star})(x) \right\}. \end{split}$$

Thus the conjugate of $A' \circ f^*$ is $f \circ A$. By the Conjugacy Theorem [Prop. 1.6.1(c)], the conjugate of $f \circ A$ is the closure of $A' \circ f^*$, provided $f \circ A$ is proper. Let $f^* : \Re^2 \mapsto (-\infty, \infty]$ be the closed proper convex function

$$f^{\star}(y_1, y_2) = \begin{cases} e^{-\sqrt{y_1 y_2}} & \text{if } y_1 \ge 0, y_2 \ge 0, \\ \infty & \text{otherwise,} \end{cases}$$

and let A' be projection on the space of y_1 , i.e., $A'(y_1, y_2) = y_1$ (cf. Exercise 3.11). Then for $x \ge 0$,

$$(A' \circ f^*)(x) = \inf_{A'y=x} f^*(y) = \inf_{y_1=x} e^{-\sqrt{y_1 y_2}} = \begin{cases} 0 & \text{if } x > 0, \\ 1 & \text{if } x = 0, \end{cases}$$

while for x < 0,

$$(A' \circ f^*)(x) = \infty.$$

Thus $A' \circ f^*$ is convex proper but not closed.

3.16 (Conjugates of Sum and Infimal Convolution)

Consider the function

$$(f_1 + \dots + f_m)(x) = f_1(x) + \dots + f_m(x),$$

where $f_i: \mathbb{R}^n \mapsto (-\infty, \infty], i = 1, \dots, m$, are closed proper convex functions, with conjugates denoted by f_i^* . Consider also the function $f_1^* \oplus \cdots \oplus f_m^*$ given by

$$(f_1^{\star} \oplus \cdots \oplus f_m^{\star})(x) = \inf_{x_1 + \cdots + x_m = x} \{ f_1^{\star}(x_1) + \cdots + f_m^{\star}(x_m) \}, \quad x \in \Re^n$$

which was called the infimal convolution of f_1^*, \ldots, f_m^* in Exercise 3.13. Show that the conjugate of $f_1^* \oplus \cdots \oplus f_m^*$ is $f_1 + \cdots + f_m$, and the conjugate of $f_1 + \cdots + f_m$ is the closure of $f_1^* \oplus \cdots \oplus f_m^*$, provided $f_1 + \cdots + f_m$ is proper [which is true if and only if $\bigcap_{i=1}^m \text{dom}(f_i) \neq \emptyset$].

Solution: We have for all $y \in \Re^n$,

$$f_{1}(y) + \dots + f_{m}(y) = \sup_{x_{1}} \{y'x_{1} - f_{1}^{\star}(x_{1})\} + \dots + \sup_{x_{m}} \{y'(x_{m} - f_{m}^{\star}(x_{m}))\}$$

$$= \sup_{x_{1},\dots,x_{m}} \{y'(x_{1} + \dots + x_{m}) - f_{1}^{\star}(x_{1}) - \dots - f_{m}^{\star}(x_{m})\}$$

$$= \sup_{x} \sup_{x_{1} + \dots + x_{m} = x} \{y'x - f_{1}^{\star}(x_{1}) - \dots - f_{m}^{\star}(x_{m})\}$$

$$= \sup_{x} \{y'x - \inf_{x_{1} + \dots + x_{m} = x} \{f_{1}^{\star}(x_{1}) + \dots + f_{m}^{\star}(x_{m})\}\}$$

$$= \sup_{x} \{y'x - (f_{1}^{\star} \oplus \dots \oplus f_{m}^{\star})(x)\}.$$

Thus the conjugate of $f_1^{\star} \oplus \cdots \oplus f_m^{\star}$ is $f_1 + \cdots + f_m$. By the Conjugacy Theorem [Prop. 1.6.1(c)] and given the properness of $f_1 + \cdots + f_m$, it follows that the conjugate of $f_1 + \cdots + f_m$ is the closure of $f_1^{\star} \oplus \cdots \oplus f_m^{\star}$, provided $f_1 + \cdots + f_m$ is proper.

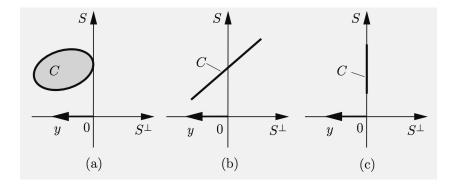


Figure 3.1. Visualization and proof of the Proper Separation Lemma of Exercise 3.17. As figure (a) illustrates, we have $S \cap \operatorname{ri}(C) = \emptyset$ if and only if S and C can be properly separated (necessarily by a hyperplane whose normal belongs to S^{\perp}). Thus, $S \cap \operatorname{ri}(C) = \emptyset$ if and only if there exists $y \in S^{\perp}$ such that $\sigma_C(y) > 0$ and $\sigma_C(-y) \leq 0$, as in figure (a). Equivalently, $S \cap \operatorname{ri}(C) \neq \emptyset$ if and only if for all $y \in S^{\perp}$ with $y \neq 0$, we have either

$$\sigma_C(y) > 0, \qquad \sigma_C(-y) > 0$$

as in figure (b), or

$$\sigma_C(y) = \sigma_C(-y) = 0$$

as in figure (c). This is equivalent to the statement of Lemma.

3.17 (Proper Separation Lemma)

This exercise and the next one provide machinery that will be used to develop sharper characterizations of the conjugates of the image and infimal convolution operations (see Fig. 3.1). Let S be a subspace and let C be a convex set. Then

$$S \cap ri(C) \neq \emptyset$$

if and only if the support function σ_C satisfies

$$\sigma_C(y) = \sigma_C(-y) = 0, \quad \forall y \in S^{\perp} \text{ with } \sigma_C(y) \leq 0.$$

Solution: Since $\operatorname{ri}(S) = S$, by the Proper Separation Theorem (Prop. 1.5.6), we have $S \cap \operatorname{ri}(C) = \emptyset$ if and only if there exists a hyperplane that properly separates S and C, i.e., a $y \in S^{\perp}$ such that

$$\sup_{x \in C} x'y > 0, \qquad \inf_{x \in C} x'y \ge 0,$$

or equivalently,

$$\sigma_C(y) > 0, \qquad \sigma_C(-y) \le 0,$$

since the sup and inf of x'y over C are the support function values $\sigma_C(y)$ and $-\sigma_C(-y)$. This shows the result, as illustrated also in Fig. 3.1.

3.18 (Relative Interior of Domain and Directions of Recession)

Let S be a subspace of \Re^n , let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let f^* be its conjugate.

(a) We have

$$S \cap \operatorname{ri}(\operatorname{dom}(f)) \neq \emptyset$$

if and only if every direction of recession of f^* that belongs to S^{\perp} is a direction along which f^* is constant.

(b) If f is closed, we have

$$S \cap \operatorname{ri}(\operatorname{dom}(f^*)) \neq \emptyset$$

if and only if every direction of recession of f that belongs to S^{\perp} is a direction along which f is constant.

Solution: (a) Use Exercise 3.17 with C = dom(f), and with σ_C equal to the recession function r_{f^*} . Note that the support function of dom(f) is the recession function of f^* , according to one of the exercises for Chapter 1.

(b) If f is closed, then by the Conjugacy Theorem [Prop. 1.6.1(c)], it is the conjugate of f^* , and the result follows by applying part (a) with the roles of f and f^* interchanged.

3.19 (Closedness of Image)

Let $f: \mathbb{R}^n \mapsto (-\infty, \infty]$ be a closed proper convex function, and let f^* be its conjugate. Consider the image function

$$(A \circ f)(x) = \inf_{Az=x} f(z), \qquad x \in \Re^n,$$

where A is an $m \times n$ matrix. Note that the convexity of $A \circ f$ was established in Exercise 3.12, and the calculation of its conjugate was given as $(f^* \circ A')(y) = f^*(A'y)$ in Exercise 3.15. This exercise provides criteria for closedness of $A \circ f$.

(a) If

$$R(A') \cap ri(dom(f^*)) \neq \emptyset$$
,

then $A \circ f$ is closed and proper, and the infimum in its definition is attained for all $x \in \text{dom}(A \circ f)$. Furthermore, $A \circ f$ is the conjugate of $f^* \circ A'$.

(b) If f is polyhedral and

$$R(A') \cap \operatorname{dom}(f^*) \neq \emptyset$$
,

then $A \circ f$ is polyhedral, and the infimum in its definition is attained for all $x \in \text{dom}(A \circ f)$. Furthermore, $A \circ f$ is the conjugate of $f^* \circ A'$.

Solution: To simplify notation, we denote $g = A \circ f$.

- (a) By Exercise 3.18, every direction of recession of f that belongs to $R(A')^{\perp}$ [which is N(A)] is a direction along which f is constant. From Exercise 3.12(b), it follows that g is closed and proper, and the infimum in its definition is attained for all $x \in \text{dom}(g)$. Furthermore, g is the conjugate of $f^* \circ A'$ in view of the closedness of g and Exercise 3.15.
- (b) We view g as the result of partial minimization of a polyhedral function. Therefore, using Prop. 3.3.1(b) and the fact that projection of a polyhedral set is polyhedral, the epigraph of g is polyhedral. Furthermore, the assumption $R(A') \cap \text{dom}(h) \neq \emptyset$ implies that there exists $\overline{x} \in \Re^m$ such that $f^*(A'\overline{x}) < \infty$ or

$$\sup_{y \in \Re^n} \left\{ y' A' \overline{x} - f(y) \right\} < \infty.$$

It follows that

$$y'A'\overline{x} \le f(y), \quad \forall y \in \Re^n,$$

so using the definition of g, we have $x'\overline{x} \leq g(x)$ for all $x \in \Re^m$. Therefore, $-\infty < g(x)$ for all $x \in \Re^m$, so that g is proper, and hence polyhedral by Prop. 3.3.5. Finally, since minimizing f(z) over $\{z \mid Az = x\}$ is a linear program whose optimal value is finite for all $x \in \text{dom}(g)$, by Prop. 1.4.12, it follows that the infimum in the definition of g is attained for all $x \in \text{dom}(g)$.

3.20 (Closedness of Infimal Convolution)

Let $f_i: \Re^n \mapsto (-\infty, \infty]$, i = 1, ..., m, be closed proper convex functions, and let f_i^* be their corresponding conjugates. Consider the infimal convolution

$$(f_1 \oplus \cdots \oplus f_m)(x) = \inf_{x_1 + \cdots + x_m = x} \{ f_1(x_1) + \cdots + f_m(x_m) \}, x \in \Re^n.$$

Note that the convexity of $f_1 \oplus \cdots \oplus f_m$ was established in Exercise 3.13, and the calculation of its conjugate was given as $f_1^* + \cdots + f_m^*$ in Exercise 3.16. This exercise provides criteria for closedness of $f_1 \oplus \cdots \oplus f_m$. Show that if for some k, the functions f_1, \ldots, f_k are polyhedral, the functions f_{k+1}, \ldots, f_m are closed, and

$$\left(\bigcap_{i=1}^{k} \operatorname{dom}(f_{i}^{\star})\right) \cap \left(\bigcap_{i=k+1}^{m} \operatorname{ri}(\operatorname{dom}(f_{i}^{\star}))\right) \neq \emptyset, \tag{3.20}$$

then $f_1 \oplus \cdots \oplus f_m$ is closed and proper, and the infimum in its definition is attained for all $x \in \text{dom}(f_1 \oplus \cdots \oplus f_m)$. In this case, $f_1 \oplus \cdots \oplus f_m$ is the conjugate of $f_1^* + \cdots + f_m^*$.

Solution: We first show the result in the two special cases where all the functions f_i are nonpolyhedral and polyhedral, respectively, and we then combine the two cases to show the result in its full generality. To simplify notation, we denote

$$g = f_1 \oplus \cdots \oplus f_m$$
.

We first assume that $\bigcap_{i=1}^{m} \operatorname{ri}(\operatorname{dom}(f_i^*)) \neq \emptyset$, and we show that the result follows from the corresponding result for the image operation (Exercise 3.19). Indeed, the infimal convolution operation is the special case of the image operation $A \circ f$, where f and A are given by

$$f(x) = f_1(x_1) + \dots + f_m(x_m), \qquad A(x_1, \dots, x_m) = x_1 + \dots + x_m.$$

It is straightforward to verify that the condition $\bigcap_{i=1}^m \text{ri}(\text{dom}(f_i^*)) \neq \emptyset$ is equivalent to

$$R(A') \cap ri(dom(f^*)) \neq \emptyset,$$

where f^* is the conjugate of f. Hence, by Exercise 3.19(a), g is closed and proper, and the infimum in its definition is attained for all $x \in \text{dom}(g)$. A similar proof, using Exercise 3.19(b), shows the result under the assumption that $\bigcap_{i=1}^m \text{dom}(f_i^*) \neq \emptyset$ and that all functions f_i are polyhedral. Finally, to show the result in the general case, we introduce the polyhedral function

$$p(y) = f_1^{\star}(y) + \dots + f_k^{\star}(y)$$

and the closed proper convex function

$$q(y) = f_{k+1}^{\star}(y) + \dots + f_m^{\star}(y).$$

The assumption (3.20) then amounts to $\operatorname{dom}(p) \cap \operatorname{ri}(\operatorname{dom}(q)) \neq \emptyset$, since $\operatorname{dom}(p) = \bigcap_{i=1}^k \operatorname{dom}(f_i^{\star})$, and the condition $\bigcap_{i=1}^m \operatorname{ri}(\operatorname{dom}(f_i^{\star})) \neq \emptyset$ implies that

$$\operatorname{ri}(\operatorname{dom}(q)) = \bigcap_{i=k+1}^{m} \operatorname{ri}(\operatorname{dom}(f_i^{\star}))$$

by Prop. 1.3.8. Let M be the affine hull of dom(q) and assume for simplicity, and without loss of generality, that M is a subspace. It can be seen that the condition $dom(p) \cap ri(dom(q)) \neq \emptyset$ is equivalent to

$$\operatorname{ri}(\operatorname{dom}(p) \cap M) \cap \operatorname{ri}(\operatorname{dom}(q)) \neq \emptyset$$

[see the exercises for Chapter 1; the proof is to choose $\overline{y} \in \text{ri}(\text{dom}(p) \cap M)$ and $y \in \text{dom}(p) \cap \text{ri}(\text{dom}(q))$ [which belongs to $\text{dom}(p) \cap M$], consider the line segment connecting y and \overline{y} , and use the Line Segment Principle to

conclude that points close to y belong to $dom(p) \cap ri(dom(q))$]. Therefore, if we replace p with the function \hat{p} given by

$$\hat{p}(y) = \begin{cases} p(y) & \text{if } y \in M, \\ \infty & \text{otherwise,} \end{cases}$$

whose domain is $dom(p) \cap M$, the special case of the result shown earlier for nonpolyhedral functions applies. We thus obtain that the function

$$\hat{g}(x) = \inf_{y \in \Re^n} \{ r(y) + s(x - y) \}$$

is closed and the infimum over y is attained for all $x \in \text{dom}(\hat{g})$, where r and s are the conjugates of \hat{p} and q, respectively. It can be seen by a straightforward calculation that $\hat{g} = g$ and the result follows.

3.21 (Preservation of Closedness Under Linear Transformation)

Let C, C_i , i = 1, ..., m, be nonempty closed convex subsets of \Re^n , and let D, D_i , i = 1, ..., m, denote the domains of their support functions:

$$D = dom(\sigma_C), \qquad D_i = dom(\sigma_{C_i}), \quad i = 1, \dots, m.$$

- (a) Let A be an $m \times n$ matrix. Then AC is closed if $R(A') \cap ri(D) \neq \emptyset$.
- (b) The vector sum $C_1 + \cdots + C_m$ is closed if $\bigcap_{i=1}^m \operatorname{ri}(D_i) \neq \emptyset$.
- (c) The vector sum $C_1 + \cdots + C_m$ is closed if for some $k \geq 1, C_1, \ldots, C_k$ is polyhedral, and $\left(\bigcap_{i=1}^k D_i\right) \cap \left(\bigcap_{i=k+1}^m \operatorname{ri}(D_i)\right) \neq \emptyset$.

Solution: (a) We apply Exercise 3.19(a) with f and f^* being the support and indicator functions of C, respectively.

(b), (c) It is sufficient to consider the case of two sets (m=2). We apply the result of Exercise 3.20 with f_1 , f_2 equal to the support functions of C_1 , C_2 , and f_1^* , f_2^* equal to the indicator functions of C_1 , C_2 . Then the indicator function of $C_1 + C_2$ is obtained as the infimal convolution of f_1^* , f_2^* . Under the given relative interior assumptions, Exercise 3.20 shows that the infimal convolution function is closed, so $C_1 + C_2$ is closed.

3.22 (Partial Minimization of Quasiconvex Functions)

Consider a function $F: \mathbb{R}^{n+m} \mapsto [-\infty, \infty]$ of vectors $x \in \mathbb{R}^n$ and $z \in \mathbb{R}^m$, which is quasiconvex (cf. Exercise 3.4). Show that the function $f: \mathbb{R}^n \mapsto [-\infty, \infty]$ given by

$$f(x) = \inf_{z \in \Re^m} F(x, z)$$

is quasiconvex. *Hint*: The level sets of f are obtained by projection of level sets of F on \mathbb{R}^n as follows: for any $\gamma \in \mathbb{R}$ and monotonically decreasing scalar sequence $\{\gamma_k\}$ with $\gamma_k \to \gamma$,

$$\{x \mid f(x) \le \gamma\} = \bigcap_{k=1}^{\infty} \{x \mid \text{there exists } (x, z) \text{ with } F(x, z) \le \gamma_k\}.$$

Solution: We follow the hint. Since intersection and projection preserve convexity, $\{x \mid f(x) \leq \gamma\}$ is convex for all $\gamma \in \Re$, so f is quasiconvex.

SECTION 3.4: Saddle Point and Minimax Theory

3.23 (Saddle Points in Two Dimensions)

Consider a function ϕ of two real variables x and z taking values in compact intervals X and Z, respectively. Assume that for each $z \in Z$, the function $\phi(\cdot,z)$ is minimized over X at a unique point denoted $\hat{x}(z)$. Similarly, assume that for each $x \in X$, the function $\phi(x,\cdot)$ is maximized over Z at a unique point denoted $\hat{z}(x)$. Assume further that the functions $\hat{x}(z)$ and $\hat{z}(x)$ are continuous over Z and X, respectively. Show that ϕ has a saddle point (x^*, z^*) . Use this to investigate the existence of saddle points of $\phi(x, z) = x^2 + z^2$ over X = [0, 1] and Z = [0, 1].

Solution: We consider a function ϕ of two real variables x and z taking values in compact intervals X and Z, respectively. We assume that for each $z \in Z$, the function $\phi(\cdot, z)$ is minimized over X at a unique point denoted $\hat{x}(z)$, and for each $x \in X$, the function $\phi(x, \cdot)$ is maximized over Z at a unique point denoted $\hat{z}(x)$,

$$\hat{x}(z) = \arg\min_{x \in X} \phi(x, z), \qquad \hat{z}(x) = \arg\max_{z \in Z} \phi(x, z).$$

Consider the composite function $f: X \mapsto X$ given by

$$f(x) = \hat{x}(\hat{z}(x)),$$

which is a continuous function in view of the assumption that the functions $\hat{x}(z)$ and $\hat{z}(x)$ are continuous over Z and X, respectively. Assume that the compact interval X is given by [a, b]. We now show that the function f has a fixed point, i.e., there exists some $x^* \in [a, b]$ such that

$$f(x^*) = x^*.$$

Define the function $g: X \mapsto X$ by

$$g(x) = f(x) - x.$$

Assume that f(a) > a and f(b) < b, since otherwise [in view of the fact that f(a) and f(b) lie in the range [a, b], we must have f(a) = a and f(b) = b, and we are done. We have

$$g(a) = f(a) - a > 0,$$

$$g(b) = f(b) - b < 0.$$

Since g is a continuous function, the preceding relations imply that there exists some $x^* \in (a,b)$ such that $g(x^*) = 0$, i.e., $f(x^*) = x^*$. Hence, we have

$$\hat{x}(\hat{z}(x^*)) = x^*.$$

Denoting $\hat{z}(x^*)$ by z^* , we obtain

$$x^* = \hat{x}(z^*), \qquad z^* = \hat{z}(x^*).$$
 (3.21)

By definition, a pair $(\overline{x}, \overline{z})$ is a saddle point if and only if

$$\max_{z \in Z} \phi(\overline{x}, z) = \phi(\overline{x}, \overline{z}) = \min_{x \in X} \phi(x, \overline{z}),$$

or equivalently, if $\overline{x} = \hat{x}(\overline{z})$ and $\overline{z} = \hat{z}(\overline{x})$. Therefore, from Eq. (3.21), we see that (x^*, z^*) is a saddle point of ϕ . We now consider the function $\phi(x, z) = x^2 + z^2$ over X = [0, 1] and Z = [0, 1]. For each $z \in [0, 1]$, the function $\phi(\cdot, z)$ is minimized over [0, 1] at a unique point $\hat{x}(z) = 0$, and for each $x \in [0, 1]$, the function $\phi(x, \cdot)$ is maximized over [0, 1] at a unique point $\hat{z}(x) = 1$. These two curves intersect at $(x^*, z^*) = (0, 1)$, which is the unique saddle point of ϕ .

3.24 (Saddle Points for Quadratic Functions)

Consider a quadratic function $\phi: X \times Z \mapsto \Re$ of the form

$$\phi(x,z) = x'Qx + c'x + z'Mx - z'Rz - d'z,$$

where Q and R are symmetric positive semidefinite $n \times n$ and $m \times m$ matrices, respectively, M is an $n \times m$ matrix, $c \in \Re^n$, $d \in \Re^m$, and X and Z are polyhedral subsets of \Re^n and \Re^m , respectively. Show that if either $\sup_{z \in Z} \inf_{x \in X} \phi(x,z)$ is finite or $\inf_{x \in X} \sup_{z \in Z} \phi(x,z)$ is finite, there exists a saddle point. Use the case where $\phi(x,z) = x + z$, $X = Z = \Re$, to show that the finiteness assumption is essential.

Solution: Here, the domain of the function

$$\sup_{z \in Z} \Phi(x, z)$$

is polyhedral and the function is convex quadratic within its domain. The result follows from the existence of solutions result for quadratic programs (Prop. 1.4.12).

Convex Optimization Theory

Chapter 4 Exercises and Solutions

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CHAPTER 4: EXERCISES AND SOLUTIONS

4.1 (Augmented Lagrangian Duality for Equality Constraints)

Construct an augmented Lagrangian framework and derive the dual function similar to the one of Section 4.2.4 for the case of the equality constraints, i.e. when we have the constraint Ex = d in place of $g(x) \leq 0$, where E is an $m \times n$ matrix and $d \in \Re^m$.

Solution: The dual function has the form

$$q_c(\mu) = \inf_{x \in X} \left\{ f(x) + \mu'(Ex - d) + \frac{c}{2} ||Ex - d||^2 \right\}.$$

4.2

In the context of Section 4.2.2, let $F(x,u) = f_1(x) + f_2(Ax + u)$, where A is an $m \times n$ matrix, and $f_1 : \Re^n \mapsto (-\infty, \infty]$ and $f_2 : \Re^m \mapsto (-\infty, \infty]$ are closed convex functions. Show that the dual function is

$$q(\mu) = -f_1^{\star}(A'\mu) - f_2^{\star}(-\mu),$$

where f_1^{\star} and f_2^{\star} are the conjugate functions of f_1 and f_2 , respectively. *Note:* This is the Fenchel duality framework discussed in Section 5.3.5.

Solution: From Section 4.2.1, the dual function is

$$q(\mu) = -p^{\star}(-\mu),$$

where p^* is the conjugate of the function

$$p(u) = \inf_{x \in \Re^n} F(x, u).$$

By using the change of variables z = Ax + u in the following calculation, we have

$$p^*(-\mu) = -\sup_{u} \left\{ -\mu' u - \inf_{x} \left\{ f_1(x) + f_2(Ax + u) \right\} \right\}$$
$$= \sup_{z,x} \left\{ -\mu' (z - Ax) - f_1(x) - f_2(z) \right\}$$
$$= f_1^*(A'\mu) + f_2^*(-\mu),$$

where f_1^{\star} and f_2^{\star} are the conjugate functions of f_1 and f_2 , respectively. Thus,

$$q(\mu) = -f_1^{\star}(A'\mu) - f_2^{\star}(-\mu).$$

4.3 (An Example of Lagrangian Duality)

Consider the problem

minimize
$$f(x)$$

subject to $x \in X$, $e'_i x = d_i$, $i = 1, ..., m$, (4.1)

where $f: \mathbb{R}^n \to \mathbb{R}$ is a convex function, X is a nonempty convex set, and e_i and d_i are given vectors and scalars, respectively. Consider the min common/max crossing framework where M is the subset of \mathbb{R}^{m+1} given by

$$M = \{ (e'_1 x - d_1, \dots, e'_m x - d_m, f(x)) \mid x \in X \},\$$

and assume that $w^* < \infty$.

(a) Show that w^* is equal to the optimal value of problem (4.1), and that the max crossing problem is to maximize $q(\mu)$ given by

$$q(\mu) = \inf_{x \in X} \left\{ f(x) + \sum_{i=1}^{m} \mu_i (e_i' x - d_i) \right\}.$$

- (b) Show that the corresponding set \overline{M} is convex.
- (c) Show that if X is compact, then $q^* = w^*$.
- (d) Show that if there exists a vector $\overline{x} \in \text{ri}(X)$ such that $e'_i \overline{x} = d_i$ for all i = 1, ..., m, then $q^* = w^*$ and the max crossing problem has an optimal solution.

Solution: (a) It is easily seen that w^* is the minimal value of f(x) subject to $x \in \text{conv}(X)$ and $e'_i x = d_i$, i = 1, ..., m. The corresponding max crossing problem is given by

$$q^* = \sup_{\mu \in \Re^m} q(\mu),$$

where $q(\mu)$ is given by

$$q(\mu) = \inf_{(u,w) \in M} \{ w + \mu' u \} = \inf_{x \in X} \left\{ f(x) + \sum_{i=1}^{m} \mu_i (e_i' x - d_i) \right\}.$$

(b) Consider the set

$$\overline{M} = \Big\{ \big(u_1, \dots, u_m, w\big) \mid \text{there exists } x \in X \text{ such that } e_i'x - d_i = u_i, \ \forall i, \ f(x) \le w \Big\}.$$

We show that \overline{M} is convex. To this end, we consider vectors $(u, w) \in \overline{M}$ and $(\tilde{u}, \tilde{w}) \in \overline{M}$, and we show that their convex combinations lie in \overline{M} . The definition of \overline{M} implies that for some $x \in X$ and $\tilde{x} \in X$, we have

$$f(x) \le w$$
, $e'_i x - d_i = u_i$, $i = 1, \dots, m$,
 $f(\tilde{x}) \le \tilde{w}$, $e'_i \tilde{x} - d_i = \tilde{u}_i$, $i = 1, \dots, m$.

For any $\alpha \in [0,1]$, we multiply these relations with α and 1- α , respectively, and add. By using the convexity of f, we obtain

$$f(\alpha x + (1 - \alpha)\tilde{x}) \le \alpha f(x) + (1 - \alpha)f(\tilde{x}) \le \alpha w + (1 - \alpha)\tilde{w},$$

$$e'_{i}(\alpha x + (1 - \alpha)\tilde{x}) - d_{i} = \alpha u_{i} + (1 - \alpha)\tilde{u}_{i}, \quad i = 1, \dots, m.$$

In view of the convexity of X, we have $\alpha x + (1-\alpha)\tilde{x} \in X$, so these equations imply that the convex combination of (u,w) and (\tilde{u},\tilde{w}) belongs to \overline{M} , thus proving that \overline{M} is convex.

(c) We prove this result by showing that all the assumptions of Min Common/Max Crossing Theorem I are satisfied. By assumption, w^* is finite. It follows from part (b) that the set \overline{M} is convex. Therefore, we only need to show that for every sequence $\{(u_k,w_k)\}\subset M$ with $u_k\to 0$, there holds $w^*\leq \liminf_{k\to\infty}w_k$. Consider a sequence $\{(u_k,w_k)\}\subset M$ with $u_k\to 0$. Since X is compact and f is convex by assumption (which implies that f is continuous by Prop. 1.4.6), it follows from Prop. 1.1.9(c) that set M is compact. Hence, the sequence $\{(u_k,w_k)\}$ has a subsequence that converges to some $(0,\overline{w})\in M$. Assume without loss of generality that $\{(u_k,w_k)\}$ converges to $(0,\overline{w})$. Since $(0,\overline{w})\in M$, we get

$$w^* = \inf_{(0,w) \in M} w \le \overline{w} = \liminf_{k \to \infty} w_k,$$

proving the desired result, and thus showing that $q^* = w^*$.

(d) We prove this result by showing that all the assumptions of Min Common/Max Crossing Theorem II are satisfied. By assumption, w^* is finite. It follows from part (b) that the set \overline{M} is convex. Therefore, we only need to show that the set

$$D = \{(e'_1 x - d_1, \dots, e'_m x - d_m) \mid x \in X\}$$

contains the origin in its relative interior. The set D can equivalently be written as

$$D = E \cdot X - d,$$

where E is a matrix, whose rows are the vectors e'_i , i = 1, ..., m, and d is a vector with entries equal to d_i , i = 1, ..., m. By Prop. 1.3.6, it follows that

$$ri(D) = E \cdot ri(X) - d.$$

Hence the assumption that there exists a vector $\overline{x} \in \text{ri}(X)$ such that $E\overline{x} - d = 0$ implies that $0 \in \text{ri}(D)$, thus showing that $q^* = w^*$ and that the max crossing problem has an optimal solution.

4.4 (Lagrangian Duality and Compactness of the Constraint Set)

Consider the problem of Exercise 4.3, but assume that f is linear and X is compact (instead of f and X being convex). Show that q^* is equal to the minimal value of f(x) subject to $x \in \text{conv}(X)$ and $e'_i x = d_i$, $i = 1, \ldots, m$. Hint: Show that

$$\operatorname{conv}(M) = \left\{ \left(e_1' x - d_1, \dots, e_m' x - d_m, f(x) \right) \mid x \in \operatorname{conv}(X) \right\},\,$$

and use Exercise 4.3(c).

Solution: Clearly, the max crossing values corresponding to M and $\operatorname{conv}(M)$ are equal [this is true generically, since closed halfspaces containing M also contain $\operatorname{conv}(M)$]. The expression for $\operatorname{conv}(M)$ in the hint follows from the linearity of f. Thus, the min $\operatorname{common/max}$ crossing framework for $\operatorname{conv}(M)$ corresponds to the problem of minimizing f(x) subject to $x \in \operatorname{conv}(X)$ and $e_i'x = d_i, i = 1, \ldots, m$. Since M is compact, $\operatorname{conv}(M)$ is also compact, and the result follows from Exercise 4.3(c).

Convex Optimization Theory

Chapter 5 Exercises and Solutions

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CHAPTER 5: EXERCISES AND SOLUTIONS

5.1 (Extended Representation)

Consider the convex programming problem

minimize
$$f(x)$$

subject to $x \in X$, $l(x) = 0$, $g(x) \le 0$, (5.1)

of Section 5.3, which involves a mixture of m linear equality and r convex inequality constraints. Let $l = (l_1, \ldots, l_m)$ and $g = (g_1, \ldots, g_r)$. Assume that the set X is described by equality and inequality constraints as

$$X = \{x \mid l_i(x) = 0, i = m + 1, \dots, \overline{m}, g_j(x) \le 0, j = r + 1, \dots, \overline{r}\}.$$

Then the problem can alternatively be described without an abstract set constraint, in terms of all of the constraint functions

$$l_i(x) = 0, \quad i = 1, \dots, \overline{m}, \qquad g_j(x) \le 0, \quad j = 1, \dots, \overline{r}.$$

We call this the *extended representation* of (P). Show if there is no duality gap and there exists a dual optimal solution for the extended representation, the same is true for the original problem (5.1).

Solution: Assume that there exists a dual optimal solution in the extended representation. Thus there exist nonnegative scalars $\lambda_1^*, \ldots, \lambda_m^*, \lambda_{m+1}^*, \ldots, \lambda_{\overline{m}}^*$ and $\mu_1^*, \ldots, \mu_r^*, \mu_{r+1}^*, \ldots, \mu_{\overline{r}}^*$ such that

$$f^* = \inf_{x \in \Re^n} \left\{ f(x) + \sum_{i=1}^{\overline{m}} \lambda_i^* h_i(x) + \sum_{j=1}^{\overline{r}} \mu_j^* g_j(x) \right\},$$

from which we have

$$f^* \leq f(x) + \sum_{i=1}^{\overline{m}} \lambda_i^* h_i(x) + \sum_{j=1}^{\overline{r}} \mu_j^* g_j(x), \qquad \forall \ x \in \Re^n.$$

For any $x \in X$, we have $h_i(x) = 0$ for all $i = m+1, \ldots, \overline{m}$, and $g_j(x) \leq 0$ for all $j = r+1, \ldots, \overline{r}$, so that $\mu_j^* g_j(x) \leq 0$ for all $j = r+1, \ldots, \overline{r}$. Therefore, it follows from the preceding relation that

$$f^* \le f(x) + \sum_{i=1}^m \lambda_i^* h_i(x) + \sum_{j=1}^r \mu_j^* g_j(x), \quad \forall \ x \in X.$$

Taking the infimum over all $x \in X$, it follows that

$$f^* \leq \inf_{x \in X} \left\{ f(x) + \sum_{i=1}^m \lambda_i^* h_i(x) + \sum_{j=1}^r \mu_j^* g_j(x) \right\}$$

$$\leq \inf_{\substack{x \in X, \ h_i(x) = 0, \ i = 1, \dots, r \\ g_j(x) \leq 0, \ j = 1, \dots, r}} \left\{ f(x) + \sum_{i=1}^m \lambda_i^* h_i(x) + \sum_{j=1}^r \mu_j^* g_j(x) \right\}$$

$$\leq \inf_{\substack{x \in X, \ h_i(x) = 0, \ i = 1, \dots, m \\ g_j(x) \leq 0, \ j = 1, \dots, r}} f(x)$$

$$= f^*.$$

Hence, equality holds throughout above, showing that the scalars $\lambda_1^*, \ldots, \lambda_m^*$, μ_1^*, \ldots, μ_r^* constitute a dual optimal solution for the original representation.

5.2 (Dual Function for Linear Constraints)

Consider the problem

minimize
$$f(x)$$

subject to $Ax \le b$, $Ex = d$,

and the dual function

$$q(\lambda,\mu) = \inf_{x \in \Re n} \{ f(x) + \lambda'(Ex - d) + \mu'(Ax - b) \}.$$

Show that

$$q(\lambda, \mu) = -d'\lambda - b'\mu - f^*(-E'\lambda - A'\mu),$$

where f^* is the conjugate function of f.

Solution: Evident from the definition of a conjugate convex function

$$f^{\star}(y) = \sup_{x \in \Re^n} \{ y'x - f(x) \}.$$

5.3 (Quadratic Programming Duality)

Consider the quadratic program

$$\begin{array}{ll} \text{minimize} & c'x+\frac{1}{2}x'Qx\\ \text{subject to} & x\in X, \qquad a'_jx\leq b_j, \quad j=1,\ldots,r, \end{array}$$

where X is a polyhedral set, Q is a symmetric positive semidefinite $n \times n$ matrix, c, a_1, \ldots, a_r are vectors in \Re^n , and b_1, \ldots, b_r are scalars, and assume that its optimal value is finite. Then there exist at least one optimal solution and at least one dual optimal solution. *Hint*: Use the extended representation of Exercise 5.1.

Solution: Consider the extended representation of the problem in which the linear inequalities that represent the polyhedral part are lumped with the remaining linear inequality constraints (cf. Exercise 5.1). From Prop. 1.4.12, finiteness of the optimal value implies that there exists a primal optimal solution and from the analysis of Section 5.3.1, there exists a dual optimal solution. From Exercise 5.1, it follows that there exists a dual optimal solution for the original representation of the problem.

5.4 (Sensitivity)

Consider the class of problems

minimize
$$f(x)$$

subject to $x \in X$, $g_j(x) \le u_j$, $j = 1, ..., r$,

where $u=(u_1,\ldots,u_r)$ is a vector parameterizing the right-hand side of the constraints. Given two distinct values \bar{u} and \tilde{u} of u, let \bar{f} and \tilde{f} be the corresponding optimal values, and assume that \bar{f} and \hat{f} are finite. Assume further that $\bar{\mu}$ and $\tilde{\mu}$ are corresponding dual optimal solutions and that there is no duality gap. Show that

$$\tilde{\mu}'(\tilde{u} - \bar{u}) \le \overline{f} - \tilde{f} \le \overline{\mu}'(\tilde{u} - \overline{u}).$$

Solution: We have

$$\overline{f} = \inf_{x \in X} \{ f(x) + \overline{\mu}' (g(x) - \overline{u}) \},\$$

$$\tilde{f} = \inf_{x \in X} \{ f(x) + \tilde{\mu}' (g(x) - \tilde{u}) \}.$$

Let $\overline{q}(\mu)$ denote the dual function of the problem corresponding to \overline{u} :

$$\overline{q}(\mu) = \inf_{x \in X} \{ f(x) + \mu' (g(x) - \overline{u}) \}.$$

We have

$$\begin{split} \overline{f} - \widetilde{f} &= \inf_{x \in X} \left\{ f(x) + \overline{\mu}' \big(g(x) - \overline{u} \big) \right\} - \inf_{x \in X} \left\{ f(x) + \widetilde{\mu}' \big(g(x) - \widetilde{u} \big) \right\} \\ &= \inf_{x \in X} \left\{ f(x) + \overline{\mu}' \big(g(x) - \overline{u} \big) \right\} - \inf_{x \in X} \left\{ f(x) + \widetilde{\mu}' \big(g(x) - \overline{u} \big) \right\} + \widetilde{\mu}' (\widetilde{u} - \overline{u}) \\ &= \overline{q}(\overline{\mu}) - \overline{q}(\widetilde{\mu}) + \widetilde{\mu}' (\widetilde{u} - \overline{u}) \\ &\geq \widetilde{\mu}' (\widetilde{u} - \overline{u}), \end{split}$$

where the last inequality holds because $\overline{\mu}$ maximizes \overline{q} . This proves the left-hand side of the desired inequality. Interchanging the roles of \overline{f} , \overline{u} , $\overline{\mu}$, and \tilde{f} , \tilde{u} , $\tilde{\mu}$, shows the desired right-hand side.

5.5 (Duality and Zero Sum Games)

Let A be an $n \times m$ matrix, and let X and Z be the unit simplices in \Re^n and \Re^m , respectively:

$$X = \left\{ x \mid \sum_{i=1}^{n} x_i = 1, x_i \ge 0, i = 1, \dots, n \right\},$$
$$Z = \left\{ z \mid \sum_{j=1}^{m} z_j = 1, z_j \ge 0, j = 1, \dots, m \right\}.$$

Show that the minimax equality

$$\max_{z \in Z} \min_{x \in X} x'Az = \min_{x \in X} \max_{z \in Z} x'Az$$

is a special case of linear programming duality. *Hint*: For a fixed z, $\min_{x \in X} x'Az$ is equal to the minimum component of the vector Az, so

$$\max_{z \in Z} \min_{x \in X} x' A z = \max_{z \in Z} \min \{ (Az)_1, \dots, (Az)_n \} = \max_{\xi e \le Az, z \in Z} \xi,$$
 (5.2)

where e is the unit vector in \Re^n (all components are equal to 1). Similarly,

$$\min_{x \in X} \max_{z \in Z} x' A z = \min_{\zeta e \ge A' x, \, x \in X} \zeta. \tag{5.3}$$

Show that the linear programs in the right-hand sides of Eqs. (5.2) and (5.3) are dual to each other.

Solution: Consider the linear program

$$\min_{\substack{\zeta e \geq A'x \\ \sum_{i=1}^{n} x_i = 1, x_i \geq 0}} \zeta,$$

whose optimal value is equal to $\min_{x \in X} \max_{z \in Z} x'Az$. Introduce dual variables $z \in \Re^m$ and $\xi \in \Re$, corresponding to the constraints $A'x - \zeta e \leq 0$ and $\sum_{i=1}^n x_i = 1$, respectively. The dual function is

$$\begin{split} q(z,\xi) &= \inf_{x_i \geq 0, \ i=1,\dots,n} \left\{ \zeta + z'(A'x - \zeta e) + \xi \left(1 - \sum_{i=1}^n x_i \right) \right\} \\ &= \inf_{x_i \geq 0, \ i=1,\dots,n} \left\{ \zeta \left(1 - \sum_{j=1}^m z_j \right) + x'(Az - \xi e) + \xi \right\} \\ &= \left\{ \xi \quad \text{if } \sum_{j=1}^m z_j = 1, \ \xi e - Az \leq 0, \\ -\infty \quad \text{otherwise.} \end{split}$$

Thus the dual problem, which is to maximize $q(z,\xi)$ subject to $z \geq 0$ and $\xi \in \Re$, is equivalent to the linear program

$$\max_{\xi e < Az, z \in Z} \xi,$$

whose optimal value is equal to $\max_{z \in Z} \min_{x \in X} x'Az$.

5.6 (Inconsistent Convex Systems of Inequalities)

Let $g_j: \Re^n \mapsto \Re$, $j=1,\ldots,r$, be convex functions over the nonempty convex subset of \Re^n . Show that the system

$$g_j(x) < 0, \qquad j = 1, \dots, r,$$

has no solution within X if and only if there exists a vector $\mu \in \Re^r$ such that

$$\sum_{j=1}^{r} \mu_j = 1, \qquad \mu \ge 0,$$

$$\mu'g(x) \ge 0, \quad \forall \ x \in X.$$

Hint: Consider the convex program

minimize ysubject to $x \in X$, $y \in \Re$, $g_i(x) \le y$, j = 1, ..., r.

Solution: The dual function for the problem in the hint is

$$q(\mu) = \inf_{y \in \Re, x \in X} \left\{ y + \sum_{j=1}^{r} \mu_j (g_j(x) - y) \right\}$$
$$= \left\{ \inf_{x \in X} \sum_{j=1}^{r} \mu_j g_j(x) & \text{if } \sum_{j=1}^{r} \mu_j = 1, \\ -\infty & \text{if } \sum_{j=1}^{r} \mu_j \neq 1. \right\}$$

The problem in the hint satisfies the Slater condition, so the dual problem has an optimal solution μ^* and there is no duality gap. Clearly the problem in the hint has an optimal value that is greater or equal to 0 if and only if the system of inequalities

$$g_j(x) < 0, j = 1, \dots, r,$$

has no solution within X. Since there is no duality gap, we have

$$\max_{\mu \ge 0, \sum_{j=1}^r \mu_j = 1} q(\mu) \ge 0$$

if and only if the system of inequalities $g_j(x) < 0, j = 1, ..., r$, has no solution within X. This is equivalent to the statement we want to prove.

5.7 (Finiteness of the Optimal Dual Value)

Consider the problem

minimize
$$f(x)$$

subject to $x \in X$, $g_j(x) \le 0$, $j = 1, ..., r$,

where X is a convex set, and f and g_j are convex over X. Assume that the problem has at least one feasible solution. Show that the following are equivalent.

- (i) The dual optimal value $q^* = \sup_{\mu \in \Re^r} q(\mu)$ is finite.
- (ii) The primal function p is proper.
- (iii) The set

$$M = \{(u, w) \in \Re^{r+1} \mid \text{there is an } x \in X \text{ such that } g(x) \le u, f(x) \le w\}$$

does not contain a vertical line.

Solution: We note that -q is closed and convex, and that

$$q(\mu) = \inf_{u \in \Re^r} \bigl\{ p(u) + \mu' u \bigr\}, \qquad \forall \; \mu \in \Re^r.$$

Since $q(\mu) \leq p(0)$ for all $\mu \in \Re^r$, given the feasibility of the problem [i.e., $p(0) < \infty$], we see that q^* is finite if and only if q is proper. Since q is the conjugate of p(-u) and p is convex, by the Conjugacy Theorem [Prop. 1.6.1(b)], q is proper if and only if p is proper. Hence (i) is equivalent to (ii). We note that the epigraph of p is the closure of p. Hence, given the feasibility of the problem, (ii) is equivalent to the closure of p not containing a vertical line. Since p is convex, its closure does not contain a line if and only if p does not contain a line (since the closure and the relative interior of p have the same recession cone). Hence (ii) is equivalent to (iii).

5.8

Show that for the function f(x) = ||x||, we have

$$\partial f(x) = \begin{cases} \left\{ x/\|x\| \right\} & \text{if } x \neq 0, \\ \left\{ g \mid \|g\| \leq 1 \right\} & \text{if } x = 0. \end{cases}$$

Solution: For $x \neq 0$, the function f(x) = ||x|| is differentiable with $\nabla f(x) = x/||x||$, so that $\partial f(x) = {\nabla f(x)} = {x/||x||}$. Consider now

the case x = 0. If a vector d is a subgradient of f at x = 0, then $f(z) \ge f(0) + d'z$ for all z, implying that

$$||z|| \ge d'z, \qquad \forall \ z \in \Re^n.$$

By letting z = d in this relation, we obtain $||d|| \le 1$, showing that $\partial f(0) \subset \{d \mid ||d|| \le 1\}$. On the other hand, for any $d \in \Re^n$ with $||d|| \le 1$, we have

$$d'z \le ||d|| \cdot ||z|| \le ||z||, \qquad \forall \ z \in \Re^n,$$

which is equivalent to $f(0) + d'z \le f(z)$ for all z, so that $d \in \partial f(0)$, and therefore $\{d \mid ||d|| \le 1\} \subset \partial f(0)$.

5.9

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let x and y be vectors in \Re^n . Show that if $g_x \in \partial f(x)$ and $g_y \in \partial f(y)$, then

$$(g_x - g_y)'(x - y) \ge 0.$$

Hint: Write the subgradient inequalities for x, y, g_x and for x, y, g_y , and add.

Solution: Following the hint, we write

$$f(y) \ge f(x) + g_x'(y - x),$$

$$f(x) \ge f(y) + g_y'(x - y).$$

By adding, we obtain

$$(g_x - g_y)'(x - y) \ge 0.$$

5.10

Let $f: \Re^n \to \Re$ be a convex function, and let x and y be given vectors in \Re^n . Consider the scalar function $\varphi: \Re \to \Re$ defined by $\varphi(t) = f(tx + (1-t)y)$ for all $t \in \Re$, and show that

$$\partial \varphi(t) = \{(x-y)'g \mid g \in \partial f(tx + (1-t)y)\}, \quad \forall t \in \Re.$$

Hint: Apply the chain rule of Prop. 5.4.5.

Solution: We can view the function

$$\varphi(t) = f(tx + (1-t)y), \qquad t \in \Re$$

as the composition of the form

$$\varphi(t) = f(g(t)), \qquad t \in \Re,$$

where $g(t): \Re \mapsto \Re^n$ is an affine function given by

$$g(t) = y + t(x - y), \qquad t \in \Re.$$

By using the Chain Rule (Prop. 5.4.5), where A = (x - y), we obtain

$$\partial \varphi(t) = A' \, \partial f(g(t)), \quad \forall t \in \Re,$$

or equivalently

$$\partial \varphi(t) = \{ (x - y)'d \mid d \in \partial f(tx + (1 - t)y) \}, \quad \forall t \in \Re.$$

5.11 (Partial Differentiation)

Consider a proper convex function F of two vectors $x \in \Re^n$ and $y \in \Re^m$. For a fixed $(\overline{x}, \overline{y}) \in \text{dom}(F)$, let $\partial_x F(\overline{x}, \overline{y})$ and $\partial_y F(\overline{x}, \overline{y})$ be the subdifferentials of the functions $F(\cdot, \overline{y})$ and $F(\overline{x}, \cdot)$ at \overline{x} and \overline{y} , respectively.

(a) Show that

$$\partial F(\overline{x}, \overline{y}) \subset \partial_x F(\overline{x}, \overline{y}) \times \partial_y F(\overline{x}, \overline{y}),$$

and give an example showing that the inclusion may be strict in general

(b) Assume that F has the form

$$F(x,y) = h_1(x) + h_2(y) + h(x,y),$$

where h_1 and h_2 are proper convex functions, and h is convex, real-valued, and differentiable. Show that the formula of part (a) holds with equality.

Solution: (a) We have $(g_x, g_y) \in \partial F(\overline{x}, \overline{y})$ if and only if

$$F(x,y) \ge F(\overline{x},\overline{y}) + g_x'(x-\overline{x}) + g_y'(y-\overline{y}), \qquad \forall \ x \in \Re^n, \ y \in \Re^m.$$

By setting $y = \overline{y}$, we obtain that $g_x \in \partial_x F(\overline{x}, \overline{y})$, and by setting $x = \overline{x}$, we obtain that $g_y \in \partial_y F(\overline{x}, \overline{y})$, so that $(g_x, g_y) \in \partial_x F(\overline{x}, \overline{y}) \times \partial_y F(\overline{x}, \overline{y})$. For an example where the inclusion is strict, consider any function whose subdifferential is not a Cartesian product at some point, such as F(x, y) = |x + y| at points $(\overline{x}, \overline{y})$ with $\overline{x} + \overline{y} = 0$.

(b) Since F is the sum of functions of the given form, Prop. 5.4.6 applies and shows that

$$\partial F(\overline{x}, \overline{y}) = \{(g_x, 0) \mid g_x \in \partial h_1(\overline{x})\} + \{(0, g_y) \mid g_y \in \partial h_2(\overline{y})\} + \{\nabla h(\overline{x}, \overline{y})\}$$

[the relative interior condition of the proposition is clearly satisfied]. Since

$$\nabla h(\overline{x}, \overline{y}) = (\nabla_x h(\overline{x}, \overline{y}), \nabla_y h(\overline{x}, \overline{y})),$$

$$\partial_x F(\overline{x}, \overline{y}) = \partial h_1(\overline{x}) + \nabla_x h(\overline{x}, \overline{y}),$$

$$\partial_y F(\overline{x}, \overline{y}) = \partial h_2(\overline{y}) + \nabla_y h(\overline{x}, \overline{y}),$$

the result follows.

5.12 (Normal Cones of Level Sets)

Let $f: \Re^n \mapsto (-\infty, \infty]$ be a proper convex function, and let x be a vector that does not achieve a minimum of f, and is such that $\partial f(x) \neq \emptyset$.

(a) Show that the normal cone $N_V(x)$, where

$$V = \{ z \mid f(z) \le f(x) \},$$

is the closure of the convex cone generated by $\partial f(x)$.

(b) Show that if $x \in \text{int}(\text{dom}(f))$, then $N_V(x)$ is the convex cone generated by $\partial f(x)$.

Solution: It is easy to show that $\operatorname{cl}(\operatorname{cone}(\partial f(x))) \subset N_v(x)$, so to prove that $N_v(x) = \operatorname{cl}(\operatorname{cone}(\partial f(x)))$, we show $N_v(x) \subset \operatorname{cl}(\operatorname{cone}(\partial f(x)))$. We will prove this by contradiction. Suppose the set $N_v(x)$ is strictly larger than $\operatorname{cl}(\operatorname{cone}(\partial f(x)))$. Since both sets are closed, the polar cone $(\operatorname{cl}(\operatorname{cone}(\partial f(x))))^*$ is strictly larger than the polar cone $(N_v(x))^*$. Thus there exists a pair of directions (y, d_1) such that

$$y'd_1 > 0$$
, $y \in N_v(x) \setminus \text{clcone}(\partial f(x))$, $d_1 \in (\text{cl}(\text{cone}(\partial f(x)))^*$.

Furthermore, since x is not the minimum of f, we have $0 \notin \partial f(x)$, therefore the set

$$(\operatorname{cl}(\operatorname{cone}(\partial f(x))))^* \cap \operatorname{aff}(\operatorname{cone}(\partial f(x)))$$

contains points other than the origin. Hence we can choose a direction d_2 such that

$$d_2 \in (\operatorname{cl}(\operatorname{cone}(\partial f(x))))^* \cap \operatorname{aff}(\operatorname{cone}(\partial f(x))),$$

and

$$y'd > 0$$
, $d \in (\operatorname{cl}(\operatorname{cone}(\partial f(x))))^*$, $d \notin (\operatorname{aff}(\operatorname{cone}(\partial f(x))))^{\perp}$.

where $d = d_1 + d_2$. Next, we show that $\sup_{z \in \partial f(x)} d'z < -\delta$ for some positive δ . Since $d \in (\operatorname{cl}(\operatorname{cone}(\partial f(x))))^*$, it can be seen that $\sup_{z \in \partial f(x)} d'z \leq 0$. To claim that strict inequality holds, since $0 \notin \partial f(x)$ and $\partial f(x)$ is a closed set, and furthermore, $d \notin (\operatorname{aff}(\operatorname{cone}(\partial f(x))))^{\perp}$, it follows that there exists some $\delta > 0$ such that

$$\sup_{z \in \partial f(x)} d'z < -\delta.$$

Finally, since the closure of the directional derivative function $f'(x;\cdot)$ is the support function of $\partial f(x)$, we have

$$(clf')(x;d) = \sup_{z \in \partial f(x)} d'z < -\delta.$$

Therefore, by the definition of the closure of a function, we can choose a direction \bar{d} sufficiently close to d such that

$$f'(x; \bar{d}) < 0, \quad y'\bar{d} > 0,$$

where the second inequality is possible because y'd > 0. The fact $f'(x; \bar{d}) < 0$ implies that along the direction \bar{d} we can find a point z such that f(z) < f(x). Since $y \in N_v$, we have $y'(z-x) \le 0$, and this contradicts $y'\bar{d} > 0$.

5.13 (Functions of Legendre Type)

A convex function f on \Re^n is said to be of *Legendre type* if it is real-valued, everywhere differentiable, and is such that for all $g \in \Re^n$, the minimum of f(x) - g'x over $x \in \Re^n$ is attained at a unique point.

- (a) Show that if f is real-valued, everywhere differentiable, and strictly convex, and f(x) g'x is coercive as a function of x for each g, it is of Legendre type. Furthermore, this is true if f is strongly convex [for example it is twice differentiable, and is such that for some $\alpha > 0$, $\nabla^2 f(x) \alpha I$ is positive definite for all x; see the exercises of Chapter 1, Exercise 1.11].
- (b) Show that for 1 , the p-th power of the p-norm

$$f(x) = \sum_{i=1}^{n} |x_i|^p,$$

and all positive definite quadratic functions are functions of Legendre type.

- (c) Give an example of a twice differentiable function, with everywhere positive definite Hessian, which is not of Legendre type.
- (d) A convex function is of Legendre type if and only if its conjugate is of Legendre type.

Solution: (a) Consider the function $f_g(x) = f(x) - g'x$, which is real-valued and convex for any value of g. Since $f_g(x)$ is convex and coercive, it has an nonempty and compact set of minima. Also $f_g(x)$ is strictly convex by strict convexity of f, which implies that it has an unique global minimal point. To sum up, f(x) is real-valued, everywhere differentiable, and is such that for all $g \in \mathbb{R}^n$, the minimum of $f(x) - g'x = f_g(x)$ over $x \in \mathbb{R}^n$ is attained at a unique point. It follows that f is of Legendre type. If f is strongly convex, the same is true for f_g , so by the result of Exercise 1.11, the minimum of f_g is attained at a unique point. In the special case where f is twice differentiable, and is such that for some a > 0, $\nabla^2 f(x) - aI$ is positive definite for all $x \in \mathbb{R}^n$. A more direct proof is possible. Then the Hessian matrix of $f_g(x)$ also satisfies $\nabla^2 f_g(x) - aI > 0$ for all x. It follows that $\lim_{\|x\| \to \infty} f_g(x) \ge \lim_{\|x\| \to \infty} (f_g(0) + x'(aI)x) = \infty$, where $f_g(0) \ne -\infty$ since f is real-valued. Thus f_g is coercive and it follows that f is of Legendre type.

(b) Consider the p-th power of the p-norm,

$$f(x) = \sum_{i=1}^{n} ||x_i||^p,$$

where $x = (x_1, x_2, ..., x_n)'$ and 1 . Note that <math>f is real-valued, everywhere differentiable, and we have

$$\partial f / \partial x_i = \begin{cases} 0, & \text{if } x_i = 0 \\ p x_i^{p-1}, & \text{if } x_i > 0 \\ -p(-x_i)^{p-1}, & \text{if } x_i < 0 \end{cases}.$$

The function f(x) - g'x is coercive, because

$$\lim_{x \leftarrow \infty} \sum_{i=1}^{n} ||x_i||^n - g'x \ge \sum_{i=1}^{n} ||x_i||^p - ||g'x|| = \infty.$$

Using part (a) it follows that f is of Legendre type. Consider an arbitrary positive definite quadratic function

$$f(x) = x'Hx + Ax + b.$$

The Hessian of f is the positive definite matrix H, so f is strongly convex. Using part (a) it follows that f is of Legendre type.

- (c) Consider the exponential function $f(x) = e^{-x}$. Then f is twice differentiable, with everywhere positive definite Hessian. However, f is not of Legendre type, because the function f(x) g'x with g > 0 does not have a minimum
- (d) Let $f: \Re^n \to \Re$ be of Legendre type, and let h be its conjugate. Then, for every $g \in \Re^n$, the minimum of f(x) g'x over $x \in \Re^n$ is attained at a unique point, denoted x(g). It follows that h(g) is the real number

g'x(g) - f(x(g)), and by the equivalence of (i) and (iii) in Prop. 5.4.3, the unique subgradient of h at g is x(g), so h is everywhere differentiable. By the equivalence of (i) and (ii) in Prop. 5.4.3, for every x, $\nabla f(x)$ is the unique point at which h(g) - x'g attains its minimum over $g \in \mathbb{R}^n$. Hence h is of Legendre type.

By interchanging the roles of f and h in the preceding argument, and by using the Conjugacy Theorem [Prop. 1.6.1(c)], it follows that if h is of Legendre type, then f is also of Legendre type.

5.14 (Duality Gap and Nondifferentiabilities)

This exercise shows how a duality gap results in nondifferentiability of the dual function. Consider the problem

minimize
$$f(x)$$

subject to $x \in X$, $g_j(x) \le 0$, $j = 1, ..., r$,

and assume that for all $\mu \geq 0$, the infimum of the Lagrangian $L(x,\mu)$ over X is attained by at least one $x_{\mu} \in X$. Show that if there is a duality gap, then the dual function $q(\mu) = \inf_{x \in X} L(x,\mu)$ is nondifferentiable at every dual optimal solution. Hint: If q is differentiable at a dual optimal solution μ^* , by the theory of Section 6.2.3, we must have $\partial q(\mu^*)/\partial \mu_j \leq 0$ and $\mu_j^* \partial q(\mu^*)/\partial \mu_j = 0$ for all j. Use optimality conditions for μ^* , together with any vector x_{μ^*} that minimizes $L(x,\mu^*)$ over X, to show that there is no duality gap.

Solution: To obtain a contradiction, assume that q is differentiable at some dual optimal solution $\mu^* \in M$, where $M = \{ \mu \in \Re^r \mid \mu \geq 0 \}$. Then

$$\nabla q(\mu^*)(\mu^* - \mu) \ge 0, \qquad \forall \ \mu \ge 0.$$

If $\mu_j^* = 0$, then by letting $\mu = \mu^* + \gamma e_j$ for a scalar $\gamma \geq 0$, and the vector e_j whose jth component is 1 and the other components are 0, from the preceding relation we obtain $\partial q(\mu^*)/\partial \mu_j \leq 0$. Similarly, if $\mu_j^* > 0$, then by letting $\mu = \mu^* + \gamma e_j$ for a sufficiently small scalar γ (small enough so that $\mu^* + \gamma e_j \in M$), from the preceding relation we obtain $\partial q(\mu^*)/\partial \mu_j = 0$. Hence

$$\partial q(\mu^*)/\partial \mu_j \le 0, \qquad \forall \ j = 1, \dots, r,$$

 $\mu_j^* \partial q(\mu^*)/\partial \mu_j = 0, \qquad \forall \ j = 1, \dots, r.$

Since q is differentiable at μ^* , we have that

$$\nabla q(\mu^*) = g(x^*),$$

for some vector $x^* \in X$ such that $q(\mu^*) = L(x^*, \mu^*)$. This and the preceding two relations imply that x^* and μ^* satisfy the necessary and sufficient optimality conditions for an optimal primal and dual optimal solution pair. It follows that there is no duality gap, a contradiction.

5.15 (Saddle Points for Quadratic Functions)

Consider a quadratic function $\phi: X \times Z \mapsto \Re$ of the form

$$\phi(x,z) = x'Qx + c'x + z'Mx - z'Rz - d'z,$$

where Q and R are symmetric positive semidefinite $n \times n$ and $m \times m$ matrices, respectively, M is an $n \times m$ matrix, $c \in \mathbb{R}^n$, $d \in \mathbb{R}^m$, and X and X are subsets of \mathbb{R}^n and \mathbb{R}^m , respectively. Derive conditions based on Prop. 5.5.6 for ϕ to have at least one saddle point.

Solution: We assume that X and Z are closed and convex sets. We will calculate the recession cone and the lineality space of the functions

$$t(x) = \begin{cases} \sup_{z \in Z} \phi(x, z) & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$

and

$$r(z) = \begin{cases} -\inf_{x \in X} \phi(x, z) & \text{if } z \in Z, \\ \infty & \text{if } z \notin Z, \end{cases}$$

and we will then use Prop. 5.5.6. For each $z \in Z$, the function $t_z : \Re^n \mapsto (-\infty, \infty]$ defined by

$$t_z(x) = \begin{cases} \phi(x, z) & \text{if } x \in X, \\ \infty & \text{if } x \notin X, \end{cases}$$

is closed and convex in view of the assumption that Q is positive semidefinite symmetric. Similarly, for each $x \in X$, the function $r_x : \Re^m \mapsto (-\infty, \infty]$ defined by

$$r_x(z) = \begin{cases} -\phi(x, z) & \text{if } z \in Z, \\ \infty & \text{if } z \notin Z, \end{cases}$$

is closed and convex in view of the assumption that R is a positive semidefinite symmetric. Hence, Assumption 3.3.1 is satisfied. By the positive semidefiniteness of Q and the calculations of Example 2.1.1, it can be seen that for each $z \in Z$, the recession cone of the function t_z is given by

$$R_{t_z} = R_X \cap N(Q) \cap \{y \mid y'(Mz + c) \le 0\},\$$

where R_X is the recession cone of the convex set X and N(Q) is the null space of the matrix Q. Similarly, for each $z \in Z$, the constancy space of the function t_z is given by

$$L_{tz} = L_X \cap N(Q) \cap \{y \mid y'(Mz + c) = 0\},\$$

where L_X is the lineality space of the set X. By the positive semidefiniteness of R, for each $x \in X$, it can be seen that the recession cone of the function r_x is given by

$$R_{r_x} = R_Z \cap N(R) \cap \{y \mid y'(Mx - d) \ge 0\},\$$

where R_Z is the recession cone of the convex set Z and N(R) is the null space of the matrix R. Similarly, for each $x \in X$, the constancy space of the function r_x is given by

$$L_{r_x} = L_Z \cap N(R) \cap \{y \mid y'(Mx - d) = 0\},\$$

where L_Z is the lineality space of the set Z. Since $t(x) = \sup_{z \in Z} t_z(x)$, the recession cone of t is

$$R_t = \bigcap_{z \in Z} R_{t_z} = R_X \cap N(Q) \cap (\bigcap_{z \in Z} \{ y \mid y'(Mz + c) \le 0 \}),$$

or

$$R_t = R_X \cap N(Q) \cap (MZ + c)^*,$$

where

$$(MZ + c)^* = \{ y \mid y'w \le 0, \forall w \in MZ + c \}.$$

[Note that $(MZ + c)^*$ is the polar cone of the set MZ + c, as defined in Chapter 4.] Similarly, the lineality space of t is

$$L_t = L_X \cap N(Q) \cap (MZ)^{\perp},$$

where

$$(MZ + c)^{\perp} = \{ y \mid y'w = 0, \forall w \in MZ + c \}.$$

By the same calculations, we also have

$$R_r = R_Z \cap N(R) \cap (-MX + d)^*, \qquad L_r = L_Z \cap N(R) \cap (-MX + d)^{\perp},$$

where

$$\begin{split} & (-MX+d)^* = \{y \mid y'w \le 0, \, \forall \; w \in -MX+d\}, \\ & (-MX+d)^\perp = \{y \mid y'w = 0, \, \forall \; w \in -MX+d\}. \end{split}$$

If

$$R_t = R_r = \{0\},\,$$

then it follows from Prop. 5.5.6(a) that the set of saddle points of ϕ is nonempty and compact. (In particular, this condition holds if Q and R are positive definite matrices, or if X and Z are compact.) Similarly, if

$$R_t = L_t, \qquad R_r = L_r,$$

then it follows from Prop. 5.5.6(b) that the set of saddle points of ϕ is nonempty.

5.16 (Goldman-Tucker Complementarity Theorem [GoT56])

Consider the linear programming problem

minimize
$$c'x$$

subject to $Ax = b$, $x \ge 0$, (LP)

where A is an $m \times n$ matrix, c is a vector in \Re^n , and b is a vector in \Re^m . Consider also the dual problem

Assume that the sets of optimal solutions of LP and DLP, denoted X^* and Λ^* , respectively, are nonempty. Show that the index set $\{1, \ldots, n\}$ can be partitioned into two disjoint subsets I and \overline{I} with the following two properties:

(1) For all $x^* \in X^*$ and $\lambda^* \in \Lambda^*$, we have

$$x_i^* = 0, \quad \forall i \in \overline{I}, \quad (A'\lambda^*)_i = c_i, \quad \forall i \in I,$$

where x_i^* and $(A'\lambda^*)_i$ are the *i*th components of x^* and $A'\lambda^*$, respectively.

(2) There exist vectors $x^* \in X^*$ and $\lambda^* \in \Lambda^*$ such that

$$\begin{split} x_i^* > 0, \quad \forall \ i \in I, \qquad x_i^* = 0, \quad \forall \ i \in \overline{I}, \\ (A'\lambda^*)_i = c_i, \quad \forall \ i \in I, \qquad (A'\lambda^*)_i < c_i, \quad \forall \ i \in \overline{I}. \end{split}$$

Hint: Apply the Tucker Complementarity Theorem (Exercises of Chapter 2).

Solution: Consider the subspace

$$S = \{(x, w) \mid bw - Ax = 0, c'x = wv, x \in \mathbb{R}^n, w \in \mathbb{R}\},$$

where v is the optimal value of (LP). Its orthogonal complement is the range of the matrix

$$\begin{pmatrix} -A' & c \\ b & -v \end{pmatrix},$$

so it has the form

$$S^{\perp} = \{ (c\zeta - A'\lambda, b'\lambda - v\zeta) \mid \lambda \in \Re^m, \zeta \in \Re \}.$$

Applying the Tucker Complementarity Theorem (Exercises of Chapter 2) for this choice of S, we obtain a partition of the index set $\{1,\ldots,n+1\}$ in two subsets. There are two possible cases: (1) the index n+1 belongs to the first subset, or (2) the index n+1 belongs to the second subset. Since the vectors (x,1) such that $x \in X^*$ satisfy Ax - bw = 0 and c'x = wv, we see that case (1) holds, i.e., the index n+1 belongs to the first index subset. In particular, we have that there exist disjoint index sets I and \overline{I} such that $I \cup \overline{I} = \{1,\ldots,n\}$ and the following properties hold:

- (a) There exist vectors $(x, w) \in S$ and $(\lambda, \zeta) \in \mathbb{R}^{m+1}$ with the property $x_i > 0$, $\forall i \in I$, $x_i = 0$, $\forall i \in \overline{I}$, w > 0, (5.4) $c_i \zeta (A'\lambda)_i = 0$, $\forall i \in I$, $c_i \zeta (A'\lambda)_i > 0$, $\forall i \in \overline{I}$, $b'\lambda = v\zeta$. (5.5)
- (b) For all $(x, w) \in S$ with $x \ge 0$, and $(\lambda, \zeta) \in \mathbb{R}^{m+1}$ with $c\zeta A'\lambda \ge 0$, $v\zeta b'\lambda \ge 0$, we have

$$x_i = 0, \quad \forall i \in \overline{I},$$
 $c_i \zeta - (A'\lambda)_i = 0, \quad \forall i \in I, \quad b'\lambda = v\zeta.$

By dividing (x, w) by w, we obtain [cf. Eq. (5.4)] an optimal primal solution $x^* = x/w$ such that

$$x_i^* > 0, \quad \forall \ i \in I, \qquad x_i^* = 0, \quad \forall \ i \in \overline{I}.$$

Similarly, if the scalar ζ in Eq. (5.5) is positive, by dividing with ζ in Eq. (5.5), we obtain an optimal dual solution $\lambda^* = \lambda/\zeta$, which satisfies the desired property

$$c_i - (A'\lambda^*)_i = 0, \quad \forall \ i \in I, \qquad c_i - (A'\lambda^*)_i > 0, \quad \forall \ i \in \overline{I}.$$

If the scalar ζ in Eq. (5.5) is nonpositive, we choose any optimal dual solution λ^* , and we note, using also property (b), that we have

$$c_i - (A'\lambda^*)_i = 0, \quad \forall i \in I, \quad c_i - (A'\lambda^*)_i \ge 0, \quad \forall i \in \overline{I}, \quad b'\lambda^* = v.$$

$$(5.6)$$

Consider the vector

$$\tilde{\lambda} = (1 - \zeta)\lambda^* + \lambda.$$

By multiplying Eq. (5.6) with the positive number $1-\zeta$, and by combining it with Eq. (5.5), we see that

$$c_i - (A'\tilde{\lambda})_i = 0, \quad \forall i \in I, \quad c_i - (A'\tilde{\lambda})_i > 0, \quad \forall i \in \overline{I}, \quad b'\tilde{\lambda} = v.$$

Thus, $\tilde{\lambda}$ is an optimal dual solution that satisfies the desired property.

REFERENCES

[GoT56] Goldman, A. J., and Tucker, A. W., 1956. "Theory of Linear Programming," in Linear Inequalities and Related Systems, H. W. Kuhn and A. W. Tucker, eds., Princeton University Press, Princeton, N.J., pp. 53-97.